

Discussion 6

Roadmap

- FAQ on hw3

FAQ on hw3

Note1: Pseudocode just provide you main idea of my solution. If you want to follow it, you need to do some adjustment and consider the edge cases on your own

Note2: Don't be limited by the pseudocode. We are glad to see various solutions.

Double Pointer

```
// You can think this is an array of pointers which point to TreeNode
TreeNode **children

// Demo
// Goal: Create an array whose length is 10 and each element is an int pointer
int length = 10;
int **array = new int*[length];
// Code below just to do some operations on one element in the array
int a = 1;
array[0] = &a;
std::cout << *array[0] << std::endl; // Result is 1
```

Tree Insertion

What is `max_width` in the tree?

It's the maximum number of children a `TreeNode` can hold

Example for `max_width`:

<code>max_width: 3</code>	<code>max_width: 4</code>
1	1
2 3 4	2 3 4 5

Example for Tree insertion:

`max_width` is 2

Insert 1 -> 2 -> 3 -> 4 -> 5

1		1		1		1		1
		/		/ \		/ \		/ \
->	2		->	2 3	->	2 3	->	2 3
						/		/ \
					4			4 5

```
# One way for tree insertion:
# Use Levelorder traversal -> You can find more info from last discussion
def insert_node(root, queue, inode):
    """Pseudocode to insert a node in a tree

    Args:
        root: root of the tree
        queue: a queue to store the node
        inode: the node you want to insert
    """
    queue.enqueue(root)
    while queue.head != None:
        node = queue.dequeue()
        # If the node is not full, insert the node
        # Full means you can not insert child on that node
        if not check_full(node):
            Insert the inode as a child of node
            return
        for child in node.children:
            queue.enqueue(child)
```

When you want to insert 6 in previous tree:
 Queue: 1 -> 2 3 -> 3 4 5 -> Find 3 is not full -> Insert 6 (left child of 3)

Right & left view of the tree

Another variant of levelorder:

```

      1          -> 1
     / \
    2   3       -> 2 3
   / \ / \
  4 5 6 7      -> 4 5 6 7
 / \
8  9          -> 8 9
```

The right view is 1 3 7 9, and the left view is 1 2 4 8

Step1: Get elements in each level (Using queue)

Step2: Get the leftmost/rightmost element in each level

Why left and right view of the tree are the same in input/1.txt?

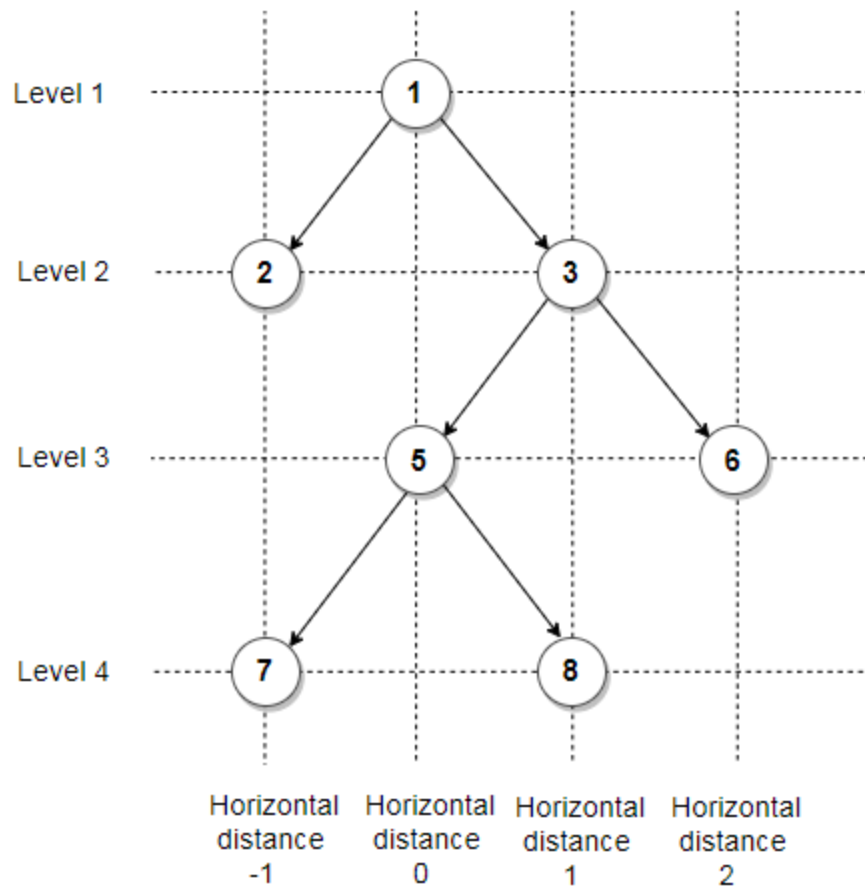
The tree degenerates to a linked list:

```

      1
      2
     3   <- 1 2 3 4 5
     4
     5
```

Bottom & Top view of the tree

- A binary tree as an example:



The bottom view above $\rightarrow 7, 5, 8, 6$

The top view above $\rightarrow 2, 1, 3, 6$

Horizontal distance:

- Horizontal distance of the root = 0
- Horizontal distance of a left child = horizontal distance of its parent - 1
- Horizontal distance of a right child = horizontal distance of its parent + 1

Algorithms: (One method to solve this kind of problem)

```
def printBottom(root, dist, level, dict):
    """Pseudocode of bottom view

    Print what you can see from the bottom of the tree.
    BTW, you need to print all of the overlapped nodes.

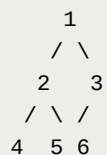
    Args:
        root: root of the tree/subtree
        dist: horizontal distance of the root
        level: level of the root
        dict:
            key: horizontal distance
            value: (nodes, level)
    """
    # For the base case
    if root is None:
        return

    # Main operations
    # 1. dist is not in dict
    if dist not in dict:
        Add (root, level) in dict[dist]
    # 2. dist is in dict and present node is at the higher level
    elif level > dict[dist][1]:
        Replace dict[dist] with (root, level)
    # 3. dist is in dict and present node is at the same level
    elif level == dict[dist][1]:
        Add node in dict[dist]

    # Recursion
    printBottom(root.left, dist - 1, level + 1, dict)
    printBottom(root.right, dist + 1, level + 1, dict)

# Actually it's a pre-order traversal for a tree
# For the top view, it's quite similar
```

Simple example:



Steps	node	dist	dict (Ignore level here for clarity)
0	1	0	{0:1}
1	2	-1	{0:1, -1:2}
2	4	-2	{0:1, -1:2, -2:4}
3	5	0	{0:5, -1:2, -2:4}
4	3	1	{0:5, -1:2, -2:4, 1:3}
5	6	0	{0:(5,6), -1:2, -2:4, 1:3}

The bottom view is 4 2 5 6 3

Time complexity: $O(n)$
Space complexity: $O(n)$