


# Lecture #8: More Recursion

## Announcements:

- Project #1 due next Thursday (13 Feb).
- Test #1 Tuesday, 18 Feb at 8PM.
- AWE 61A Party this Sunday (9 Feb) in the Woz, 1-3PM.
- Guerilla Sections this weekend (see Piazza).
- Self-assessment quiz will be released tonight, due Monday. Watch the website and Piazza.

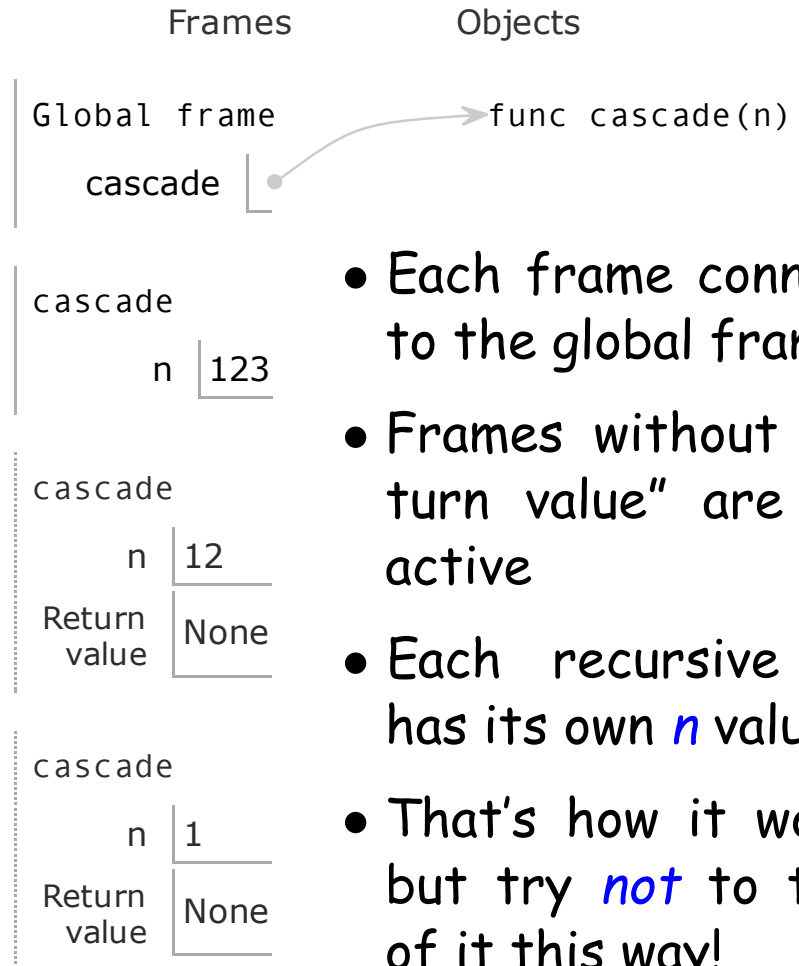
# A Simple Recursion

```
1 def cascade(n):
2     if n < 10:
3         print(n)
4     else:
5         print(n)
6         cascade(n//10)
7         print(n)
8
9 cascade(123)
```



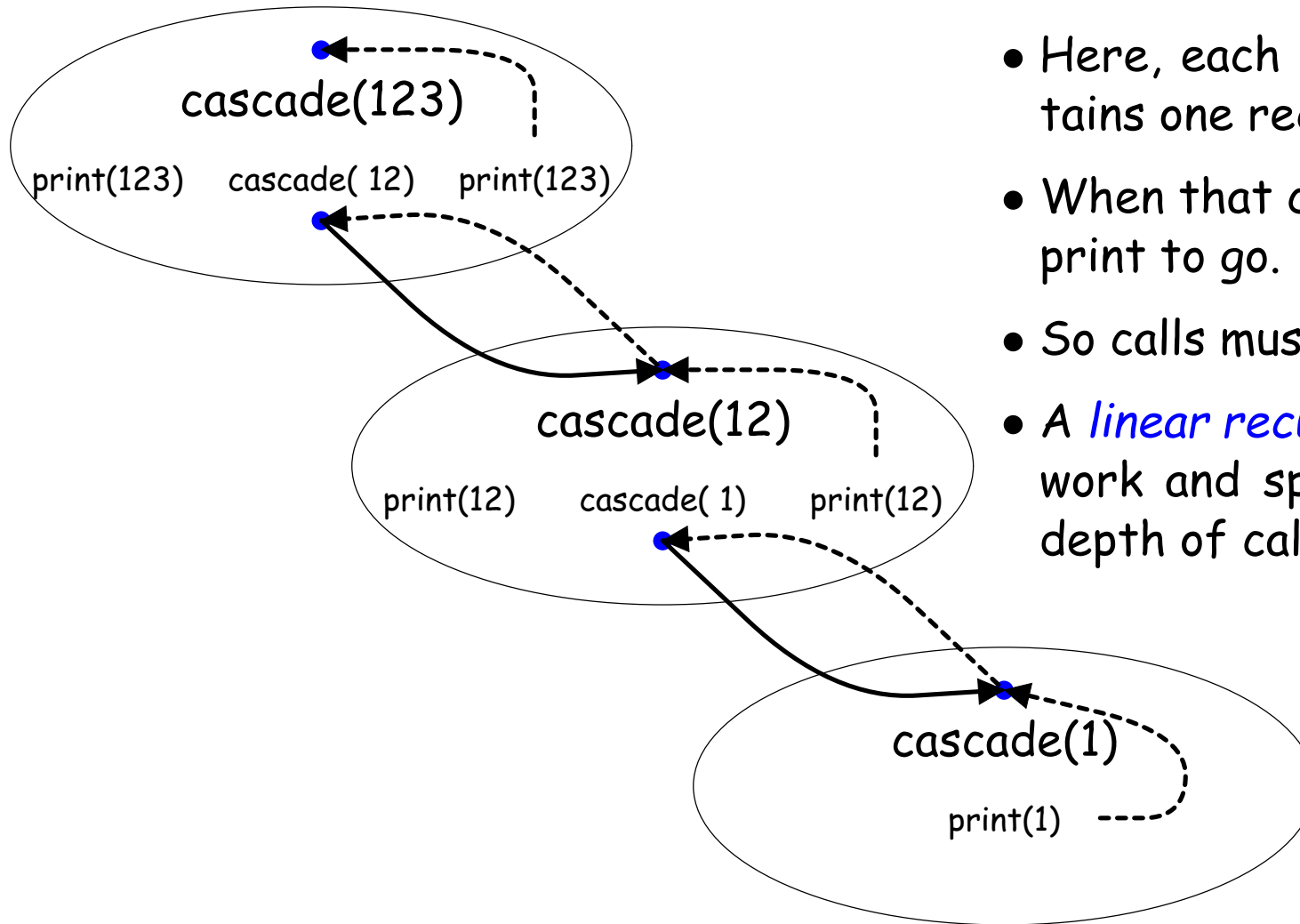
Program output:

```
123
12
1
12
```



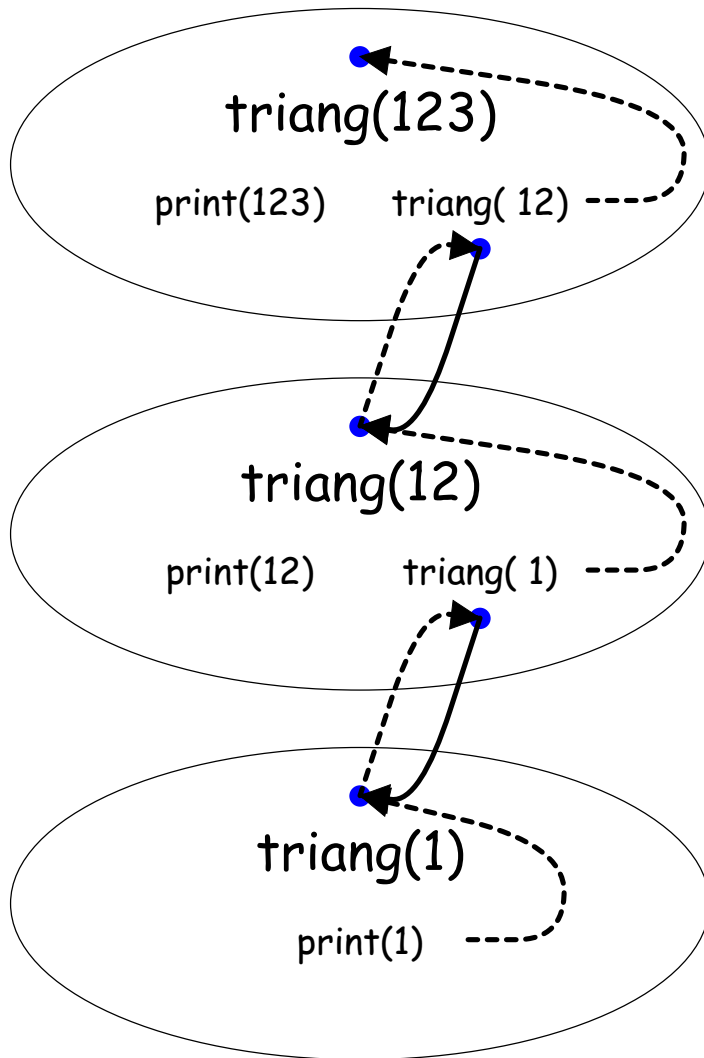
- Each frame connects to the global frame.
- Frames without "Return value" are still active
- Each recursive call has its own *n* value.
- That's how it works, but try *not* to think of it this way!
- Think recursively instead.

# Classifying Recursions: Linear Recursions



- Here, each call of `cascade` contains one recursive call.
- When that call completes, still a print to go.
- So calls must remain pending.
- A *linear recursive process*: total work and space proportional to depth of calls.

# Classifying Recursions: Iterative Processes



```
def triang(n):  
    print(n)  
    if n < 10: triang(n-1)
```

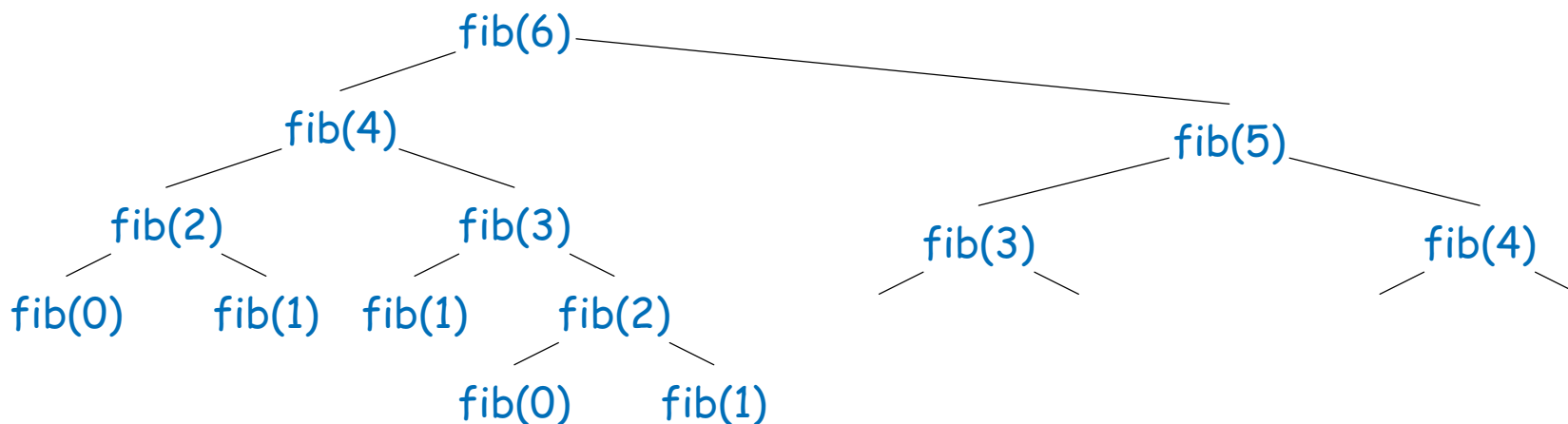
- Again, each call of `triang` contains one recursive call.
- So this is a type of linear recursive process.
- But there's no more to do when that call completes (*tail recursive*)
- So in principle, calls need not remain pending.
- An *iterative process*: total work still proportional to depth of calls, but total space need not be.
- This kind is suitable for a loop.

# Classifying Recursion: Tree Recursions

- Previously, we looked at a program for computing values in the Fibonacci sequence:

```
def fib(n):  
    """The Nth Fibonacci number, N>=0."""  
    assert n >= 0  
    if n <= 1:  
        return n  
    else:  
        return fib(n-2) + fib(n-1)
```

Here, each invocation of `fib` makes *two* calls: work is exponential in depth of calls: A *tree-recursive process*.



# A Tree Recursion: Partitions

- *partitions*( $n, k$ ): The number of non-decreasing sequences of two or more positive integers between 1 and  $k$  that add up to  $n$ .
- For example, *partitions*(6, 4) is 9:

$$2 + 4 = 6$$

$$1 + 1 + 4 = 6$$

$$3 + 3 = 6$$

$$1 + 2 + 3 = 6$$

$$1 + 1 + 1 + 3 = 6$$

$$2 + 2 + 2 = 6$$

$$1 + 1 + 2 + 2 = 6$$

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# Computing Partitions

- Observation: can choose sizes  $1-k$  for the last partition.
- If we choose size  $k$  for the last partition, then how many ways are there to partition the rest?  
.
- Suppose we choose not to use size  $k$  for the last partition, then how many choices are there?  
.
- Finally, there is only one way to partition 0 items or to partition a negative number of items or a positive number of items with maximum partition size of 0.

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  - The number of ways of partitioning  $n - k$  items of maximum size  $k$ .
  - Suppose we choose not to use size  $k$  for the last partition, then how many choices are there?
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- Finally, there is only one way to partition 0 items or to partition a negative number of items or a positive number of items with maximum partition size of 0.



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- The number of ways of partitioning  $n - k$  items of maximum size  $k$ .
- Suppose we choose not to use size  $k$  for the last partition, then how many choices are there?
- The number of ways of partitioning  $n$  items of maximum size  $k - 1$ .
- Finally, there is only one way to partition  $0$  items or to partition a negative number of items or a positive number of items with maximum partition size of  $0$ .

# Partitions, concluded

This leads to the following program:

```
def partitions(n, k):  
    """The number of ways of partitioning N items into partitions of size  
    <=K."""  
    if n == 0:  
        return 1  
    elif n < 0 or k <= 0:  
        return 0  
    else:  
        with_k =  
        without_k =  
        return with_k + without_k
```

# Partitions, concluded

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def partitions(n, k):  
    """The number of ways of partitioning N items into partitions of size  
    <=K."""  
    if n == 0:  
        return 1  
    elif n < 0 or k <= 0:  
        return 0  
    else:  
        with_k = partitions(n-k, k)  
        without_k = partitions(n, k-1)  
        return with_k + without_k
```