Lecture 37: Declarative Programming

Announcements:

- Autograder should start running this weekend.
- Remember: you still have to provide your own tests!
- We have updated files other than scheme.py a few times. I generally announce if the modification is important. You can use Unix diff to check for differences between what you have and the version in ~cs61a/lib/projects/scheme.
- Submit your Project 4 contest entries as "proj4-contest." by next Wednesday. Assuming we get entries, we'll ask the class to judge these entries

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Imperative vs. Declarative

- So far, our programs are explicit directions for solving a problem; the problem itself is *implicit* in the program.
- Declarative programming turns this around:
 - A "program" is a description of the desired characteristics of a solution.
 - It is up to the system to figure out how to achieve these characteristics.
- Taken to the extreme, this is a very difficult problem in AI.
- However, people have come up with interesting compromises for small problems.
- For example, constraint solvers allow you to specify relationships between objects (like minimum or maximum distances) and then try to find configurations of those objects that meet the constraints.

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Prolog and Predecessors

- Way back in 1959, researchers at CMU created GPS (General Problem Solver [A. Newell, J. C. Shaw, H. A. Simon])
 - Input defined objects and allowable operations on them, plus a description of the desired outcome.
 - Output consisted of a sequence of operations to bring the outcome about.
 - Only worked for small problems, unsurprisingly.
- Planner at MIT [C. Hewitt, 1969] was another programming language for theorem proving: one specified desired goal assertion, and system would find rules to apply to demonstrate the assertion. Again, this didn't scale all that well.
- Planner was one inspiration for the development of the *logic-programming* language *Prolog*.

Prolog (Lisp Style)

- Let's interpret Scheme expressions as logical assertions.
- For example, (likes brian potstickers) might be such an assertion: likes is a *predicate* that relates brian and potstickers.
- We don't interpret the arguments of the predicate: they are just uninterpreted data structures.
- We also allow one other type of expression: a symbol that starts with an underscore will indicate a *logical variable*.
- An assertion such as (likes brian _X) asserts that there is some replacement for _X that makes the assertion true.

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Facts and Rules

- We will make queries in the form of assertions, possibly with logical variables.
- The system will look to see if the queries are true based on a database of facts (axioms or postulates) about the predicates.
- It will inform us of what replacements for logical variables make the assertion true.
- Each fact will have the form

```
(fact Conclusion Hypothesis1 Hypothesis2...)
```

Meaning "For any substitution of logical variables in the Conclusion and Hypotheses, we may derive the conclusion if we can derive each of the hypotheses."

Example: Family Relations

First, some facts with no hypotheses:

```
(fact (parent george paul))
(fact (parent martin george))
(fact (parent martin martin_jr))
(fact (parent martin donald))
(fact (parent george ann))
```

Now some general rules about relations:

```
(fact (ancestor _X _Y) (parent _X _Y))
(fact (ancestor _X _Y) (parent _X _Z) (ancestor _Z _Y))
```

From these, we ought to be able to conclude that Martin is an ancestor of Ann, for example.

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Relations, Not Functions In this style of programming, we don't define functions, but rather relations. Instead of saying abs(-3) == 3, we say abs(-3, 3) (that is, "3 stands in the abs relation to -3." Instead of add(x, y) == z, we say add(x, y, z). This will allow us to run programs "both ways": from inputs to outputs, or from outputs to inputs.	
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