

What is Sorting?

By definition sorting refers to arranging data in a particular format : either ascending or descending.

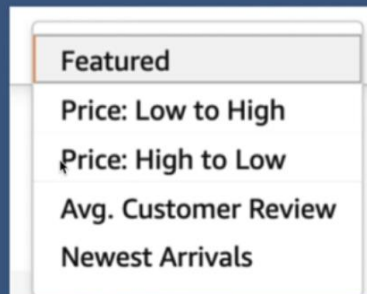
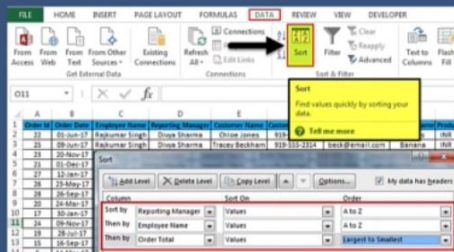


What is Sorting?

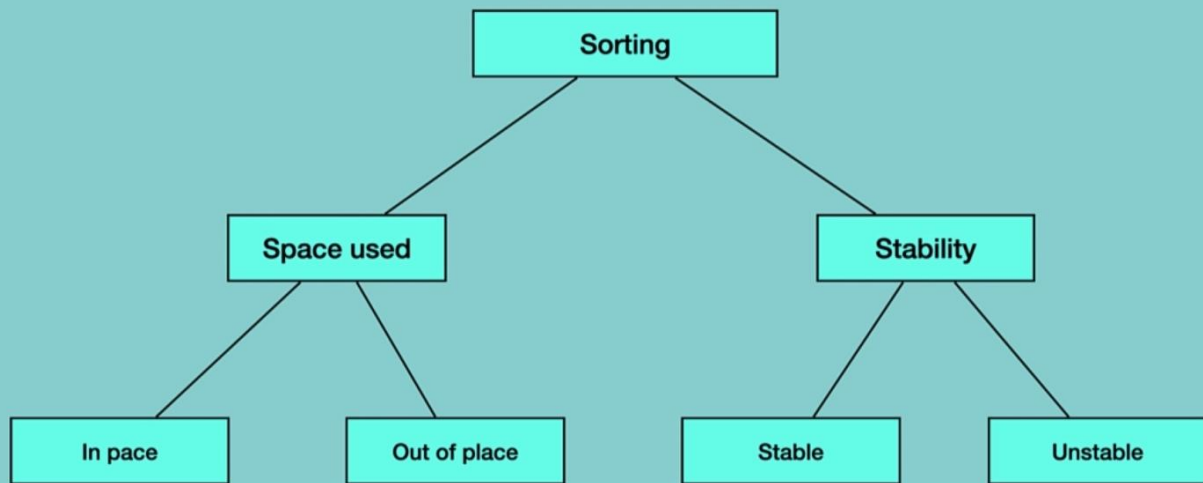
Practical Use of Sorting

Microsoft Excel : Built in functionality to sort data

Online Stores: Online shopping websites generally have option for sorting by price, review, ratings..



Types of Sorting?



Space used

In place sorting : Sorting algorithms which does not require any extra space for sorting

Example : Bubble Sort

70	10	80	30	20	40	60	50	90
----	----	----	----	----	----	----	----	----

10	20	30	40	50	60	70	80	90
----	----	----	----	----	----	----	----	----

Out place sorting : Sorting algorithms which requires an extra space for sorting

Example : Merge Sort

70	10	80	30	20	40	60	50	90
----	----	----	----	----	----	----	----	----

10	20	30	40	50	60	70	80	90
----	----	----	----	----	----	----	----	----

Stability

Stable sorting : If a sorting algorithm after sorting the contents does not change the sequence of similar content in which they appear, then this sorting is called stable sorting.

Example : Insertion Sort

70	10	80	40	20	40	60	50	90
----	----	----	----	----	----	----	----	----

10	20	40	40	50	60	70	80	90
----	----	----	----	----	----	----	----	----

UnStable sorting : If a sorting algorithm after sorting the content changes the sequence of similar content in which they appear, then it is called unstable sort.

Example : Quick Sort

70	10	80	40	20	40	60	50	90
----	----	----	----	----	----	----	----	----

10	20	40	40	50	60	70	80	90
----	----	----	----	----	----	----	----	----

Stability

UnStable sorting example

Unsorted data		Sorted by name		Sorted by age (stable)		Sorted by age (unstable)	
Name	Age	Name	Age	Name	Age	Name	Age
Renad	7	Nick	6	Nick	6	Nick	6
Nick	6	Parker	7	Richard	6	Richard	6
Richard	6	Renad	7	Parker	7	Renad	7
Parker	7	Richard	6	Renad	7	Parker	7
Sofia	7	Sofia	7	Sofia	7	Sofia	7

GROUP BY command in DBMS is based on this concept.

Sorting Terminology

Increasing Order

- If successive element is greater than the previous one
- Example : 1, 3, 5, 7, 9, 11

Decreasing Order

- If successive element is less than the previous one
- Example : 11, 9, 7, 5, 3, 1

Non Increasing Order

- If successive element is less than or equal to its previous element in the sequence.
- Example : 11, 9, 7, 5, 5, 3, 1

Non Decreasing Order

- If successive element is greater than or equal to its previous element in the sequence
- Example : 1, 3, 5, 7, 7, 9, 11

When we see “Non” then we should assume that there may be duplicate elements

Sorting Algorithms

Bubble sort

Selection sort

Insertion sort

Bucket sort

Merge sort

Quick sort

Heap sort

Which one to select?

- Stability
- Space efficient
- Time efficient

Different sorting algorithms are chosen at different instances based on our requirement. We should choose the best algorithm for our need whether it should be space efficient or it should be time efficient, etc.

Bubble Sort

- Bubble sort is also referred as Sinking sort
- We repeatedly compare each pair of adjacent items and swap them if they are in the wrong order



Initial array

Bubble Sort

- Bubble sort is also referred as Sinking sort
- We repeatedly compare each pair of adjacent items and swap them if they are in the wrong order



Intermediate State of Bubble Sort

Here the numbers are not sorted. In Bubble sort, we check for the biggest element between two numbers and swap the biggest number to the right side keeping the smallest number to the left. After one iteration the largest element will be at the last position indicating that it is sorted. Similarly proceeding in this way, the numbers will get sorted.

Animation link:

<https://upload.wikimedia.org/wikipedia/commons/c/c8/Bubble-sort-example-300px.gif>

Visualization of the code:

[https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+BubbleSort%7B%0A++++public+static+void+main\(String%5B%5D+args\)+%7B%0A++++++int+arr%5B%5D+%3D+%7B5,1,3,9,0%7D%3B%0A++++++int+n%3Darr.length%3B%0A++++++for\(int+i%3D0%3Bi%3Cn-1%3Bi%2B%2B\)%7B%0A++++++for\(int+j%3D0%3Bj%3Cn-1%3Bj%2B%2B\)%7B%0A++++++if\(arr%5Bj%5D%3Earr%5Bj%2B1%5D\)%7B%0A++++++int+temp+%3D+arr%5Bj%5D%3B%0A++++++arr%5Bj%5D%3Darr%5Bj%2B1%5D%3B%0A++++++arr%5Bj%2B1%5D%3Dtemp%3B%0A++++++%7D%0A++++++%7D%0A++++++%7D%0A%0A++++%7D%0A%7D&mode=display&curInstr= 88](https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+BubbleSort%7B%0A++++public+static+void+main(String%5B%5D+args)+%7B%0A++++++int+arr%5B%5D+%3D+%7B5,1,3,9,0%7D%3B%0A++++++int+n%3Darr.length%3B%0A++++++for(int+i%3D0%3Bi%3Cn-1%3Bi%2B%2B)%7B%0A++++++for(int+j%3D0%3Bj%3Cn-1%3Bj%2B%2B)%7B%0A++++++if(arr%5Bj%5D%3Earr%5Bj%2B1%5D)%7B%0A++++++int+temp+%3D+arr%5Bj%5D%3B%0A++++++arr%5Bj%5D%3Darr%5Bj%2B1%5D%3B%0A++++++arr%5Bj%2B1%5D%3Dtemp%3B%0A++++++%7D%0A++++++%7D%0A++++++%7D%0A%0A++++%7D%0A%7D&mode=display&curInstr= 88)

Time complexity: $O(N^2)$

Space complexity: $O(1)$

Bubble Sort

When to use Bubble Sort?

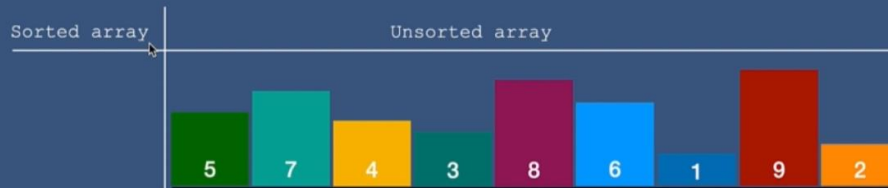
- When the input is almost sorted
- Space is a concern
- Easy to implement

When to avoid Bubble Sort?

- Average time complexity is poor

Selection Sort

- In case of selection sort we repeatedly find the minimum element and move it to the sorted part of array to make unsorted part sorted.



Initial array

Selection Sort

- In case of selection sort we repeatedly find the minimum element and move it to the sorted part of array to make unsorted part sorted.



Intermediate array in the Selection sort

In Selection Sort, there are two areas in an array one is the sorted area and the other is the unsorted area. We check for the smallest element in the unsorted array and swap the left most element of the unsorted array which will next be the part of the sorted array.

Animation of Selection Sort:

<https://upload.wikimedia.org/wikipedia/commons/9/94/Selection-Sort-Animation.gif>

Visualization of the code:

[https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+SelectionSort+%7B%0A++++public+static+void+main\(String%5B%5D+args\)+%7B%0A++++++++int+arr%5B%5D+%3D+%7B10,2,5,8,1,9%7D%3B%0A++++++++int+n%3Darr.length%3B%0A%0A++++++++for\(int+i%3D0%3Bi%3Cn%3Bi%2B%2B\)%7B](https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+SelectionSort+%7B%0A++++public+static+void+main(String%5B%5D+args)+%7B%0A++++++++int+arr%5B%5D+%3D+%7B10,2,5,8,1,9%7D%3B%0A++++++++int+n%3Darr.length%3B%0A%0A++++++++for(int+i%3D0%3Bi%3Cn%3Bi%2B%2B)%7B)


```
%0A++++++int+minIndex%3Di%3B%0A++++++for(int+j%3Di%2B1%3Bj%3Cn%3Bj%2B%2B)%7B%0A++++++if(arr%5Bi%5D%3Carr%5BminIndex%5D)%7B%0A++++++minIndex%3Di%3B%0A++++++%7D%0A++++++%7D%0A++++++if(minIndex!%3Di)%7B+//optimization%0A++++++int+temp%3Darr%5Bi%5D%3B%0A++++++arr%5Bi%5D%3Darr%5BminIndex%5D%3B%0A++++++arr%5BminIndex%5D%3Dtemp%3B%0A++++++%7D%0A++++++%7D%0A%0A++++++for(int+k%3D0%3Bk%3Cn%3Bk%2B%2B)%7B%0A++++++System.out.print(arr%5Bk%5D%2B%22+%22)%3B%0A++++++%7D%0A++++%7D%0A%7D%0A&mode=display&curlInstr=116
```

Time complexity: $O(N^2)$

Space complexity: $O(1)$

Selection Sort

When to use Selection Sort?

- When we have insufficient memory
- Easy to implement

When to avoid Selection Sort?

- When time is a concern

Insertion Sort

- Divide the given array into two part
- Take first element from unsorted array and find its correct position in sorted array
- Repeat until unsorted array is empty



Initial array

Insertion Sort

- Divide the given array into two part
- Take first element from unsorted array and find its correct position in sorted array
- Repeat until unsorted array is empty



Intermediate array in Insertion Sort

Insertion sort is quite similar to Selection Sort. In Insertion Sort, we have sorted area and unsorted area in an array. We select the minimum element from the unsorted area of the array and find the right position for the element in the sorted area of the array. By proceeding in this similar way, the elements get sorted in the array.

Animation of Insertion Sort:

https://en.wikipedia.org/wiki/Insertion_sort#/media/File:Insertion-sort-example-300px.gif

Visualization of the code:

```
https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+InsertionSort+%7B%0A++++public+static+void+main(String%5B%5D+args)+%7B%0A++++++int+arr%5B%5D+%3D%7B10,3,2,5,8,4,3,1%7D%3B%0A++++++int+n%3Darr.length%3B%0A++++++for(int+i%3D1%3Bi%3Cn%3Bi%2B%2B)%7B%0A++++++int+temp%3Darr%5B%5D,j%3Di%3B+//temp+important%0A++++++while(j%3E0+%26%26+arr%5Bj-1%5D%3Etemp)%7B%0A++++++arr%5Bj%5D%3Darr%5Bj-1%5D%3B%0A++++++j--%3B%0A++++++%7D%0A++++++arr%5Bj%5D%3Dtemp%3B%0A++++++%7D%0A%0A++++++for(int+i%3D0%3Bi%3Cn%3Bi%2B%2B)%7B%0A++++++System.out.print(arr%5B%5D%2B%22+%22)%3B%0A++++++%7D%0A++++++%7D%0A%7D%0A&mode=display&curlInstr=0
```

Time complexity: $O(N^2)$

Space complexity: $O(1)$

Insertion Sort

When to use Insertion Sort?

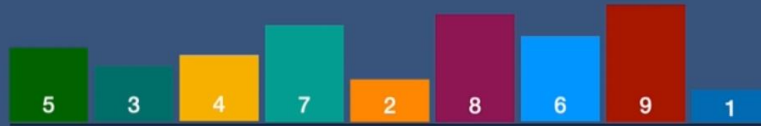
- When we have insufficient memory
- Easy to implement
- When we have continuous inflow of numbers and we want to keep them sorted

When to avoid Insertion Sort?

- When time is a concern

Bucket Sort

- Create buckets and distribute elements of array into buckets
- Sort buckets individually
- Merge buckets after sorting



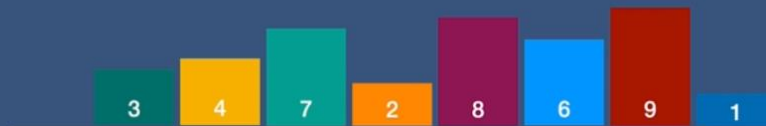
- Number of buckets = $\text{round}(\sqrt{\text{number of elements}})$
 $\text{round}(\sqrt{9}) = 3$
- Appropriate bucket = $\text{ceil}(\text{Value} * \text{number of buckets} / \text{maxValue})$



Initial state of the array

Bucket Sort

- Create buckets and distribute elements of array into buckets
- Sort buckets individually
- Merge buckets after sorting



- Number of buckets = $\text{round}(\sqrt{\text{number of elements}})$
 $\text{round}(\sqrt{9}) = 3$
- Appropriate bucket = $\text{ceil}(\text{Value} * \text{number of buckets} / \text{maxValue})$
 $\text{ceil}(5 * 3 / 9) = \text{ceil}(1.6) = 2$



For the first element in the array which is 5

In Bucket Sort, we create buckets for elements of the array. The number of buckets is calculated by the formula in the above image. Appropriate bucket for the selected element will also be calculated by the above formula in the above images. Then, sort the bucket with anyone of the sorting algorithm. Then merging the bucket to get the sorted array as the result.

Bucket Sort

- Create buckets and distribute elements of array into buckets
- Sort buckets individually
- Merge buckets after sorting

Number of buckets = $\text{round}(\sqrt{\text{number of elements}})$

$\text{round}(\sqrt{9}) = 3$

Appropriate bucket = $\text{ceil}(\text{Value} * \text{number of buckets} / \text{maxValue})$

Sort all buckets (using any sorting algorithm)



Initial state of the bucket before sorting

Bucket Sort

- Create buckets and distribute elements of array into buckets
- Sort buckets individually
- Merge buckets after sorting

Number of buckets = $\text{round}(\sqrt{\text{number of elements}})$

$\text{round}(\sqrt{9}) = 3$

Appropriate bucket = $\text{ceil}(\text{Value} * \text{number of buckets} / \text{maxValue})$

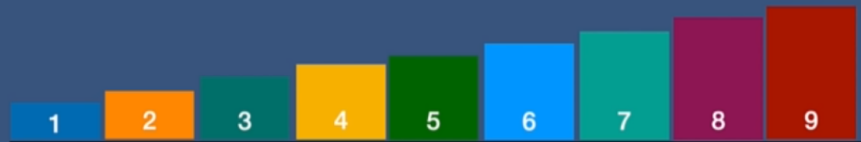
Sort all buckets (using any sorting algorithm)



After sorting the individual buckets with any sorting algorithms.

Bucket Sort

- Create buckets and distribute elements of array into buckets
- Sort buckets individually
- Merge buckets after sorting



Number of buckets = $\text{round}(\sqrt{\text{number of elements}})$

$\text{round}(\sqrt{9}) = 3$

Appropriate bucket = $\text{ceil}(\text{Value} * \text{number of buckets} / \text{maxValue})$

Sort all buckets (using any sorting algorithm)



After final step, the array elements are as follows.

Bucket Sort

When to use Bucket Sort?

- When input uniformly distributed over range

1,2,4,5,3,8,7,9

When to avoid Bucket Sort?

- When space is a concern

Bucket Sort

When to use Bucket Sort?

- When input uniformly distributed over range

1,2,4,5,3,8,7,9

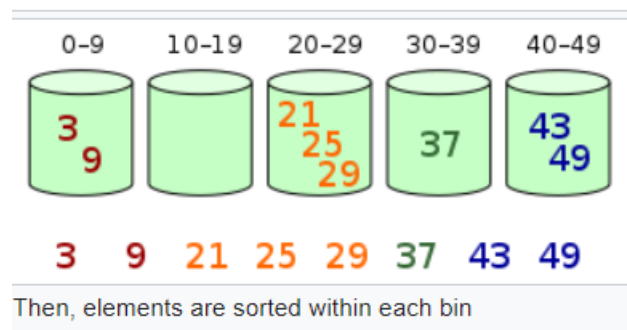
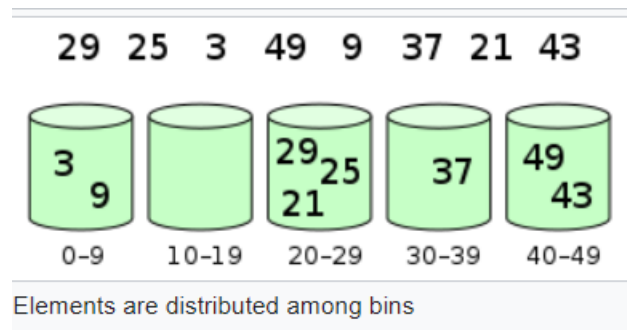
1,2,4,91,43,95

Not Uniformly distributed which means that the difference between the elements should not be of greater value. Eg: $93 - 1 = 92$ (Greater Value).

When to avoid Bucket Sort?

- When space is a concern

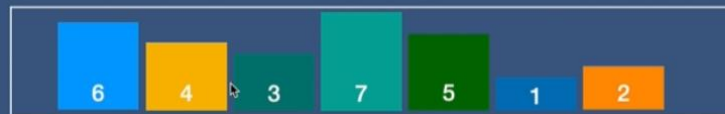
Example Scenario from Wikipedia – Quite Different: so follow the above procedure.



Merge Sort

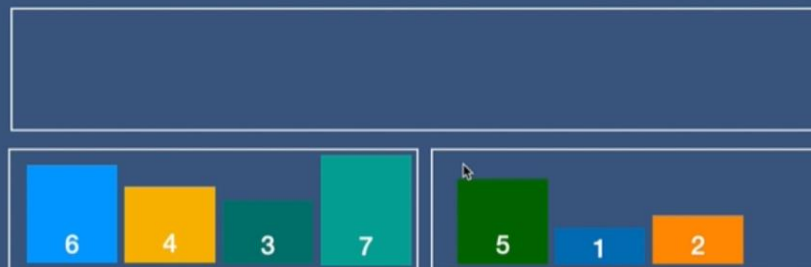
- Merge sort is a divide and conquer algorithm
- Divide the input array in two halves and we keep halving recursively until they become too small that cannot be broken further
- Merge halves by sorting them

Merge Sort



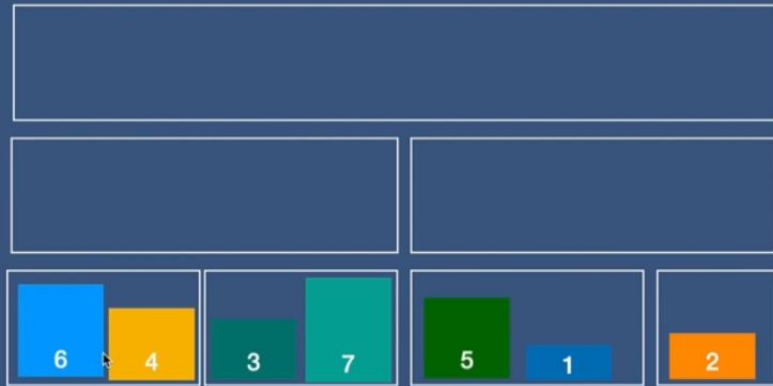
Initial state of the array

Merge Sort

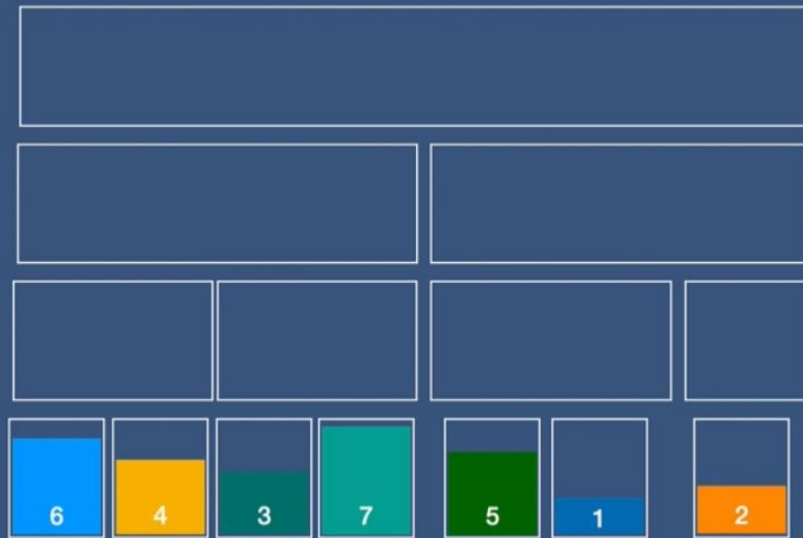


Dividing

Merge Sort

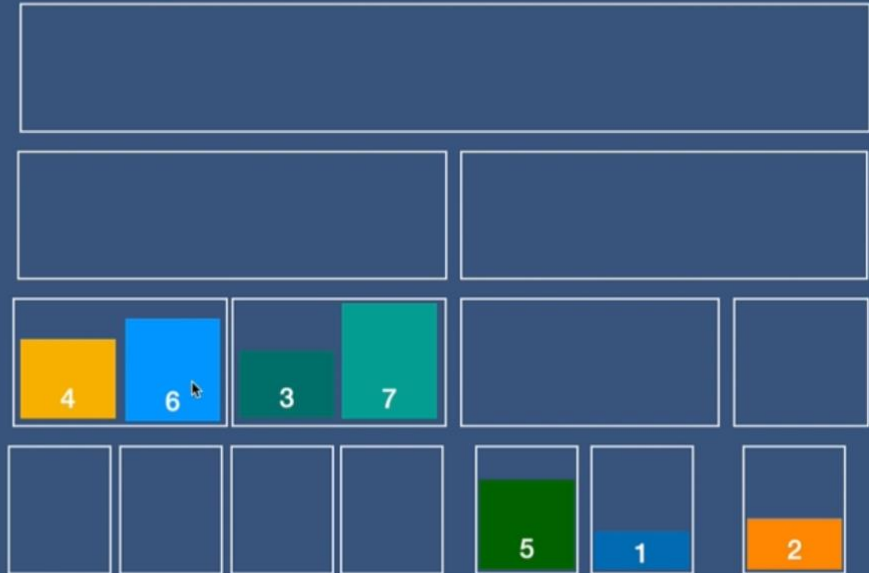


Merge Sort

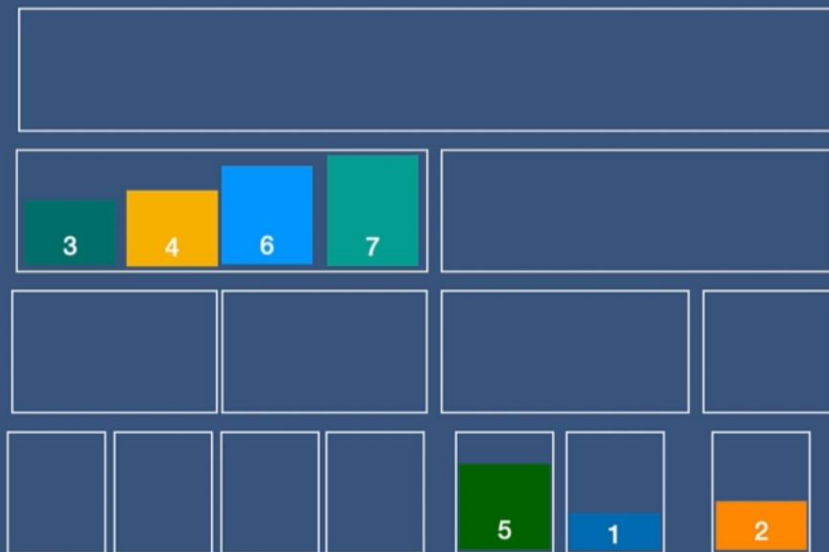


Merging:

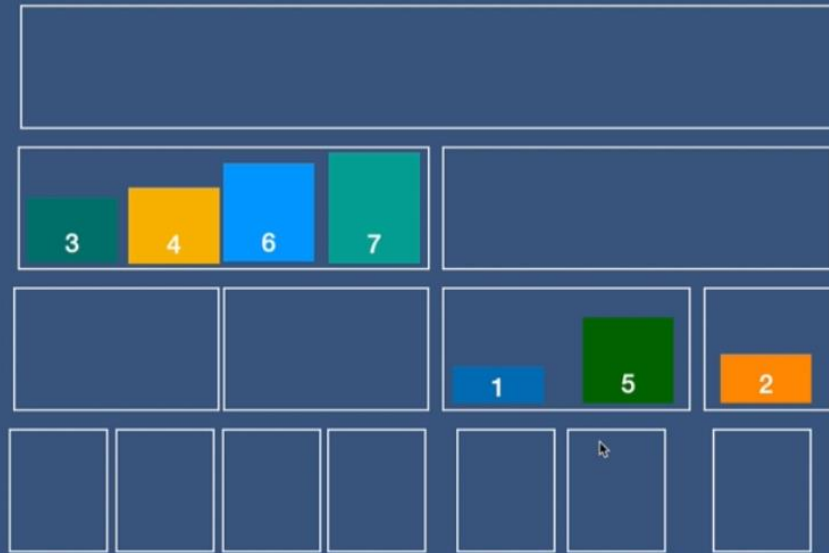
Merge Sort



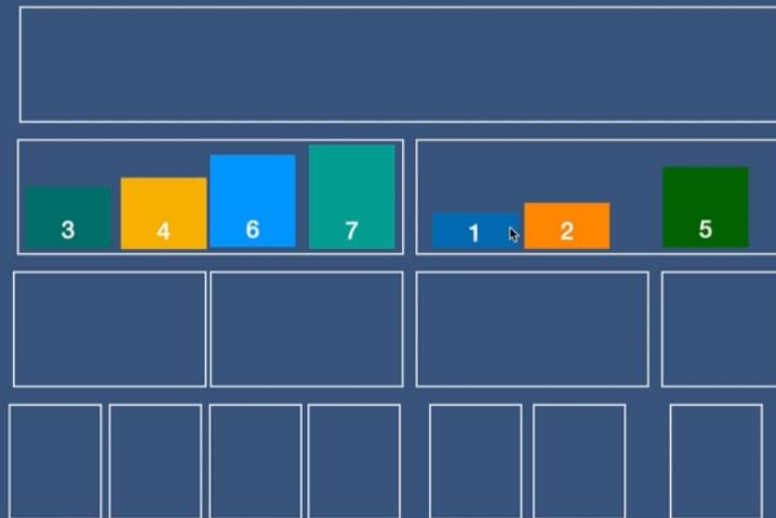
Merge Sort



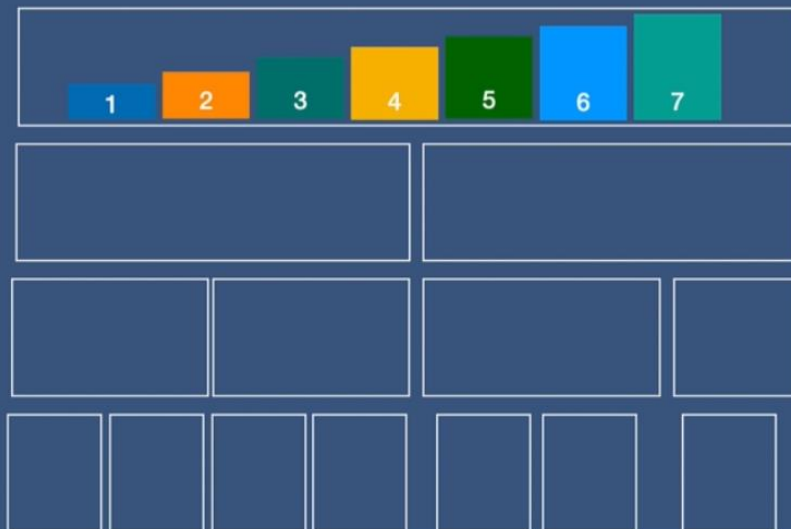
Merge Sort



Merge Sort



Merge Sort



Merge Sort

When to use Merge Sort?

- When you need stable sort
- When average expected time is $O(N \log N)$

When to avoid Merge Sort?

- When space is a concern

Time complexity: $O(N \log N)$

Space complexity: $O(N)$

Animation of Merge Sort:

https://en.wikipedia.org/wiki/Merge_sort#/media/File:Merge-sort-example-300px.gif

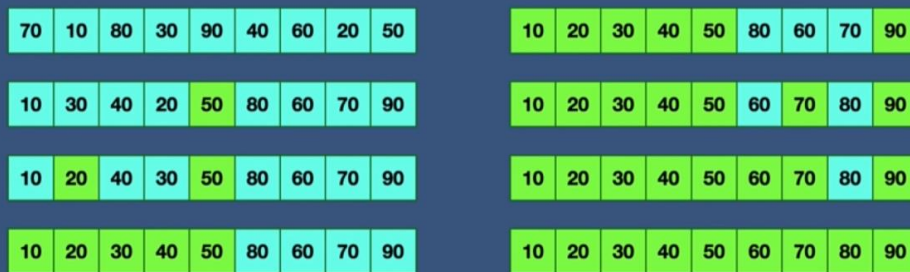
Visualization link of the code:

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Taking the right most element as the pivot element. The single green colored elements are the pivot elements and the continuous green coloured elements indicates that the array is sorted (partially or completely).

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Initial state of the array

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
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Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Pointers stop up here. The left pointer moves until it find the element greater than the pivot element and the right pointer moves until it find the element less than the pivot element. Then, the right and the left pointer get swapped.

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



We stop when the right and the left pointer meets up (9 → pointers meet here). Then, we swap the meeting element (9) of the left and the right pointer to the pivot element (6). So, the meeting point of the left and the right pointer is considered to be fully sorted.

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



And now we can observe that the left of the sorted element (6) is smaller than the sorted element (6) and to the right are the larger element than the sorted element (6).

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



We then continue the same process with the left part of the sorted element (6). Refer the above image.

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Swapping the element 1 and 3.

Quick Sort

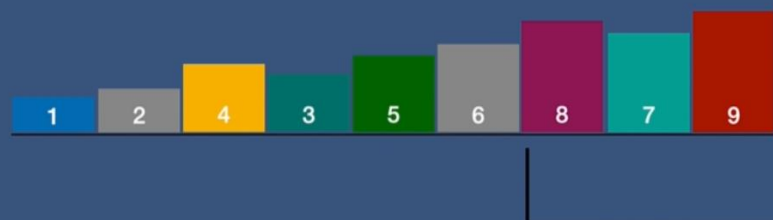
- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Now swap the pivot number(2) with the meeting element of the left pointer and the right pointer.

Quick Sort

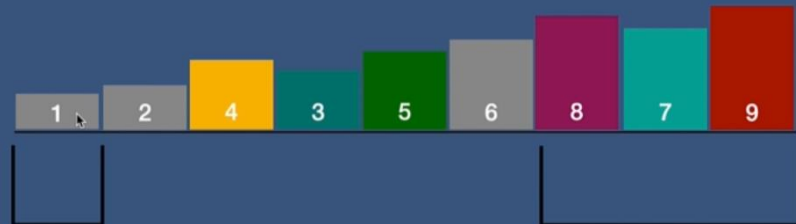
- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Now the element 2 becomes fully sorted.

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



To the left of the element 2 we have only one element so now it becomes fully sorted.

Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
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Quick Sort

- Quick sort is a divide and conquer algorithm
- Find pivot number and make sure smaller numbers located at the left of pivot and bigger numbers are located at the right of the pivot.
- Unlike merge sort extra space is not required



Now the array is fully sorted

Quick sort is the most effective sorting algorithm and it is the best even when the size of the array is large.

Quick Sort

When to use Quick Sort?

- When average expected time is $O(N \log N)$

When to avoid Quick Sort?

- When space is a concern
- When you need stable sort

Animation for Quick Sort:

<https://en.wikipedia.org/wiki/Quicksort#/media/File:Quicksort-example.gif>

Visualization of the code:

[https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+QuickSort+%7B%0A%0A++static+int+partition\(int%5B%5D+array,+int+start,+int+end\)+%7B%0A++++int+pivot+%3D+end%3B%0A++++int+i+%3D+start+-+1%3B%0A++++for+\(int+j+%3D+start%3B+j%3C%3Dend%3B+j%2B%2B\)+%7B%0A++++++if+\(array%5Bj%5D+%3C%3Darray%5Bpivot%5D\)+%7B%0A++++++i%2B%2B%3B%0A++++++int+temp+%3D+array%5Bi%5D%3B%0A++++++array%5Bi%5D+%3D+array%5Bj%5D%3B%0A++++++array%5Bj%5D+%3D+temp%3B%0A++++++%7D%0A++++%7D%0A++++return+i%3B%0A++%7D%0A%0A++public+static+void+quickSort\(int%5B%5D+array,+int+start,+int+end\)+%7B%0A++++if+\(start+%3C+end\)+%7B%0A++++++int+pivot+%3D+partition\(array,+start,+end\)%3B%0A++++++quickSort\(array,+start,+pivot+-1\)%3B%0A++++++quickSort\(array,+pivot+%2B+1,+end\)%3B%0A++++%7D%0A++%7D%0A%0A%09public+static+void+printArray\(int%5B%5D+array\)+%7B%0A%09%09for+\(int+i+%3D+0%3B+i+%3C+array.length%3B+i%2B%2B\)+%7B%0A%09%09%09System.out.print\(array%5Bi%5D%2B%22++%22\)%3B%0A%09%09%7D%0A%09%7D%0A++public+static+void+main\(String%5B%5D+args\)+%7B%0A++++int+array%5B%5D+%3D+%7B10,3,2,7,8%7D%3B%0A++++quickSort\(array,+0,+array.length-1\)%3B%0A++++printArray\(array\)%3B%0A++%7D%0A%7D%0A&mode=display&curlInsr=0](https://cscircles.cemc.uwaterloo.ca/java_visualize/#code=public+class+QuickSort+%7B%0A%0A++static+int+partition(int%5B%5D+array,+int+start,+int+end)+%7B%0A++++int+pivot+%3D+end%3B%0A++++int+i+%3D+start+-+1%3B%0A++++for+(int+j+%3D+start%3B+j%3C%3Dend%3B+j%2B%2B)+%7B%0A++++++if+(array%5Bj%5D+%3C%3Darray%5Bpivot%5D)+%7B%0A++++++i%2B%2B%3B%0A++++++int+temp+%3D+array%5Bi%5D%3B%0A++++++array%5Bi%5D+%3D+array%5Bj%5D%3B%0A++++++array%5Bj%5D+%3D+temp%3B%0A++++++%7D%0A++++%7D%0A++++return+i%3B%0A++%7D%0A%0A++public+static+void+quickSort(int%5B%5D+array,+int+start,+int+end)+%7B%0A++++if+(start+%3C+end)+%7B%0A++++++int+pivot+%3D+partition(array,+start,+end)%3B%0A++++++quickSort(array,+start,+pivot+-1)%3B%0A++++++quickSort(array,+pivot+%2B+1,+end)%3B%0A++++%7D%0A++%7D%0A%0A%09public+static+void+printArray(int%5B%5D+array)+%7B%0A%09%09for+(int+i+%3D+0%3B+i+%3C+array.length%3B+i%2B%2B)+%7B%0A%09%09%09System.out.print(array%5Bi%5D%2B%22++%22)%3B%0A%09%09%7D%0A%09%7D%0A++public+static+void+main(String%5B%5D+args)+%7B%0A++++int+array%5B%5D+%3D+%7B10,3,2,7,8%7D%3B%0A++++quickSort(array,+0,+array.length-1)%3B%0A++++printArray(array)%3B%0A++%7D%0A%7D%0A&mode=display&curlInsr=0)

Time complexity: $O(N \log N)$

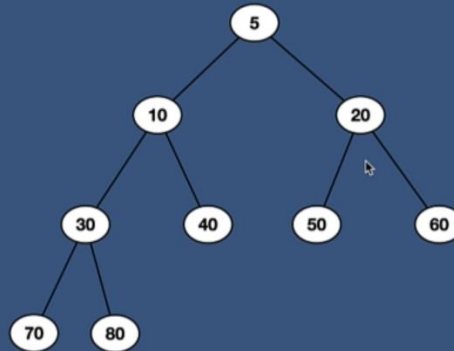
Space complexity: $O(N)$

Heap Sort

- Step 1 : Insert data to Binary Heap Tree
- Step 2 : Extract data from Binary Heap
- It is best suited with array, it does not work with Linked List

Binary Heap is a binary tree with special properties

- The value of any given node must be less or equal of its children (min heap)
- The value of any given node must be greater or equal of its children (max heap)

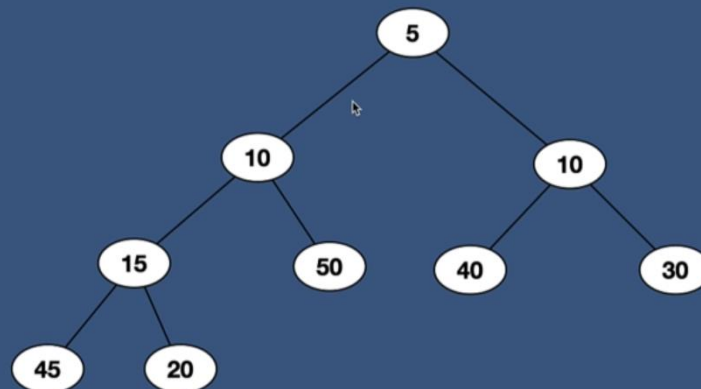


The Binary Heap Tree is an example of Min Heap.

Heap Sort

- Step 1 : Insert data to Binary Heap Tree
- Step 2 : Extract data from Binary Heap

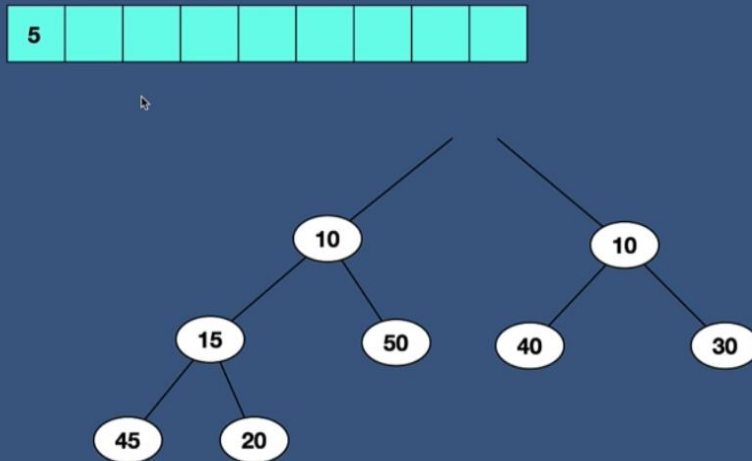
15	10	40	20	50	10	30	45	5
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Step 1: Insertion of the data to the Binary Heap Tree.

Heap Sort

- Step 1 : Insert data to Binary Heap Tree
- Step 2 : Extract data from Binary Heap

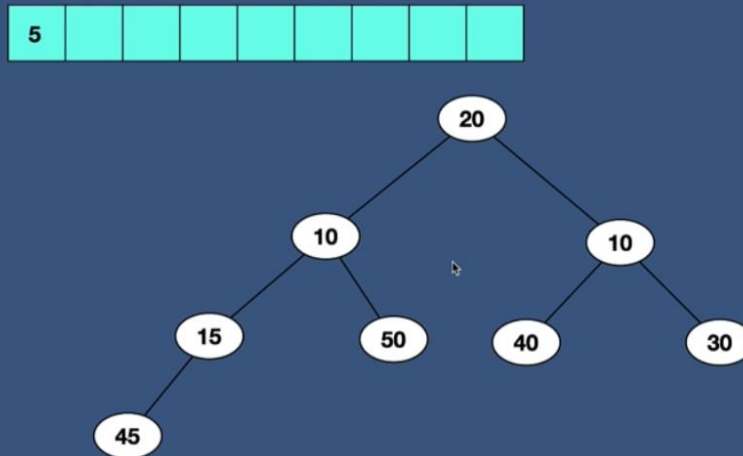


Step 2: Extracting data from the Binary Heap Tree

First, we insert the root element of the Binary Heap Tree to the array and then try to insert the left most element of the Binary Heap tree to the array

Heap Sort

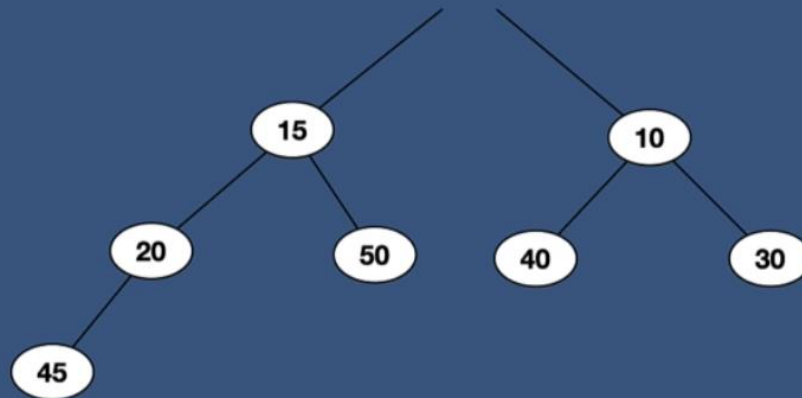
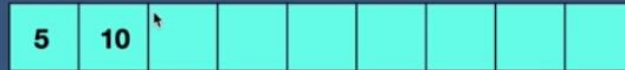
- Step 1 : Insert data to Binary Heap Tree
- Step 2 : Extract data from Binary Heap



The left most element of the Binary Heap gets to the root node and then we check for the Binary Heap property. If the Binary Heap property is not satisfied then we call Heapify method.

Heap Sort

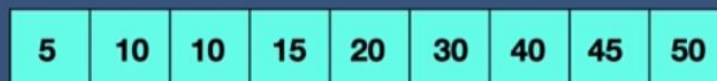
- Step 1 : Insert data to Binary Heap Tree
- Step 2 : Extract data from Binary Heap



After Heapifying the Binary Heap Tree, we then extract the root node element and then put that element to the array. This process is continued until the Binary Heap becomes null.

Heap Sort

- Step 1 : Insert data to Binary Heap Tree
- Step 2 : Extract data from Binary Heap



After the complete extraction process from the Binary Heap Tree.

Time complexity: $O(N \log N)$

Space complexity: $O(1)$

Sorting Algorithms

Name	Time Complexity	Space Complexity	Stable
Bubble Sort	$O(n^2)$	$O(1)$	Yes
Selection Sort	$O(n^2)$	$O(1)$	No
Insertion Sort	$O(n^2)$	$O(1)$	Yes
Bucket Sort	$O(n \log n)$	$O(n)$	Yes
Merge Sort	$O(n \log n)$	$O(n)$	Yes
Quick Sort	$O(n \log n)$	$O(n)$	No
Heap Sort	$O(n \log n)$	$O(1)$	No