# Using the code generator language $\mathcal{CSTL}$

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## 1 Introduction

 $\mathcal{CSTL}$  is a simple text-based language for defining code generators for the UMLRSDS subset of UML.

Rules in  $\mathcal{CSTL}$  have the form:

```
Source syntax |--> Target syntax <when> Condition
```

The left side of the rule is some text syntax of a UML class, expression or statement, where KM3 text syntax is used for class declarations, and OCL syntax for expressions. The statement syntax of UML-RSDS is used for statements. The right side of the rule is written in the syntax of the target language. The *Condition* is expressed in terms of the source language syntax categories and stereotypes.

Rules are grouped based on the source syntax category: types, expressions, etc.

As an example, some rules from UML to Java could be written as:

The  $\_i$  for positive integer i denote variables which hold concrete syntax fragments. They are termed 'metavariables'. On the LHS of a rule,  $\_i$  represents some UML fragment, such as a name, or the contents of a class. On the RHS  $\_i$  denotes the corresponding code fragment derived by applying the mapping rules recursively. Variables should be named successively  $\_1$ ,  $\_2$ , etc, up to  $\_9$ .

Specialised rules are listed before more general rules. Given a source text element elem, the first rule whose LHS matches elem is applied to elem, with metavariables  $\_i$  of the LHS being bound to fragments within elem. These fragments are then themselves mapped by the rule set and the result of transformation substituted for  $\_i$  on the RHS of the rule. If no rule applies, an element is mapped to itself.

As an example, the following KM3 text of a UML package:

```
package App
{ class CDO
  { reference sectors : Set(Sector); }
  abstract class Sector
  { attribute probDefault : double;
    attribute lossAmount : int;
  class IndustrySector extends Sector
  { reference companys : Set(Company); }
  class BankingSector extends Sector
  { reference banks : Set(Bank); }
  class Bank
  { attribute name : String; }
  class Company
  { attribute name : String; }
Would be converted by the above rules into Java code beginning:
package App;
class CDO
  Set<Sector> sectors = new HashSet<Sector>();
```

```
abstract class Sector
{
   double probDefault;
   int lossAmount;
}
```

## 2 Writing CSTL specifications

The following categories of rules are available:

• Type rules, defining the target language equivalents of usages of UML types int, double, long, String, boolean, enumerations, class types and collection types. These rules are listed together, following the header Type ::.

For example:

```
Type::
Set(_1) |-->HashSet<_1>
Sequence(_1) |-->ArrayList<_1>
```

map UML collection types to Java 7+ templated types.

- Enumeration rules, defining the interpretation of enumerated type definitions and their literal values.
- Expression rules, defining the interpretation of OCL expressions. These are divided into categories of basic expressions, binary expressions, unary expressions, conditional expressions and set expressions. For example:

```
BasicExpression::
_1._2(_3) |-->_1._2(_3)
_1(_2) |-->_1(_2)
_1._2 |-->_1.get_2()
```

define rules for operation calls with and without navigation in the function, and navigation of data features. These would map a call obj.op(x.att) into obj.op(x.getatt()).

- Package and class rules, which define how the top-level structure of a class diagram maps to a program.
- Attribute and operation rules, defining how class data and operation features are mapped. Attribute rules either have the format

```
attribute _1 : _2;
```

for data features of string, numeric, boolean or enumerated type, or the format

```
reference _1 : _2;
```

for data features of collection or class types. The keywords identity or static can be used for the attribute case.

- Statement rules, defining how the high-level activity statements of UML-RSDS are interpreted as program code.
- Use case rules defining how use cases are interpreted as program code.
- Text rules, defining pattern matching and replacement based on unstructured text, eg.:

```
Text::
createByPK_1(_2) |-->_1.createByPK_1(_2)
```

These apply to basic expressions and statements.

Rules are applied repeatedly in order to transform arbitrarily complex expressions, statements and classes, each application replaces a metavariable  $_{-i}$  on the LHS with the transformed text of whatever source text was in the  $_{-i}$  place. Thus, given some binary and unary expression rules:

## UnaryExpression::

```
-_1 |-->-_1
+_1 |-->+_1
_1->log() |-->Math.log(_1)
_1->exp() |-->Math.exp(_1)
_1->sin() |-->Math.sin(_1)
_1->cos() |-->Math.cos(_1)
_1->tan() |-->Math.tan(_1)
_1->sqr() |-->(_1)*(_1)
_1->sqrt() |-->Math.sqrt(_1)
_1->cbrt() |-->Math.cbrt(_1)
_1->floor() |-->((int) Math.floor(_1))
_1->ceil() |-->((int) Math.ceil(_1))
_1->round() |-->((int) Math.round(_1))
_1->size() |-->_1.size()
_1->first() |-->0cl.first(_1)
_1->last() |-->0cl.last(_1)
_1->tail() |-->Ocl.tail(_1)
_1->front() |-->Ocl.front(_1)
```

```
_1->reverse() |-->Ocl.reverse(_1)
_1->max() |-->Ocl.max(_1)
_1->min() |-->Ocl.min(_1)
_1->sum() |-->Ocl.sum(_1)
_1->sort() |-->Ocl.sort(_1)
not(_1) |-->!(_1)
_1->isEmpty() |-->(_1.size() == 0)
_1->notEmpty() |-->(_1.size() > 0)
_1->display() |-->
                     System.out.println(_1 + "");
_1->asSet() |-->Ocl.asSet(_1)
_1->flatten() |-->Ocl.flatten(_1)
BinaryExpression::
_1 & _2 |-->_1 && _2
_1 or _2 |-->_1 || _2
_1 xor _2 |-->((_1 || _2) && !(_1 && _2))
_1 + _2 |-->_1 + _2
_1 - _2 |-->_1 - _2<when> _1 numeric, _2 numeric
_1 - _2 |-->Ocl.stringSubtract(_1,_2)<when> _1 String, _2 String
_1 - _2 |-->Ocl.setSubtract(_1,_2)<when> _1 Set, _2 collection
_1 - _2 |-->Ocl.sequenceSubtract(_1,_2)<when> _1 Sequence, _2 collection
_1 mod _2 |-->_1 % _2
_1 * _2 |-->_1 * _2
_1 / _2 |-->_1 / _2
```

 $x \rightarrow exp() * x \rightarrow cbrt()$  is rewritten to Math.exp(x) \* Math.cbrt(x) by applying the binary expression rule for \*, then applying the unary expression rules for exp and cbrt to the two arguments of the \* expression. Rules for operators of lower precedence (such as or) should be listed before those of higher precedence.

A built-in rule is that a list of expressions \_1, \_2 is mapped to \_1, \_2, where \_1 is an individual expression and \_2 a list of expressions.

Rules may have *conditions*, for example the rules for — above distinguish between numeric subtraction and other meanings of the — symbol by inspecting the type of elements in the metavariables. If \_1 and \_2 are both numeric (have *int*, *long* or *double* type) then the rule

```
_1 - _2 |-->_1 - _2<when> _1 numeric, _2 numeric
```

applies, however if the metavariable elements have String type, the next rule applies:

```
_1 - _2 |-->Ocl.stringSubtract(_1,_2)<when> _1 String, _2 String
```

Possible conditions for expression metavariables are:

numeric String

```
Set
Sequence
collection
object
classid
value
variable
enumerationLiteral
```

The first 6 concern the type of the expression, the last 4 concern the syntactic category of a basic expression: classId is true for an identifier which is a class name, value for a literal value, etc.

Conditions can be negated, eg:

```
_1 = _2 |-->_1 == _2<when> _1 not String, _1 not object, _1 not collection
```

Possible conditions for type rules are:

#### enumerated

class

Class means a class type. The condition *collection* can be used for reference rules, to distinguish many-valued reference declarations from single-valued references.

A condition for an attribute is *primary*, meaning that it is the first identity attribute listed in its owner class.

Stereotypes can also be used as conditions, for classes and attributes. For example:

### Attribute::

```
_1 : _2 |--> let _1 : _2<when>_1 readOnly
```

for a mapping from UML to Swift.

The RHS of rules can also use metafeatures, which compute some function of the source element held in a metavariable. The notation is  $\_i$ 'f for metafeature f. Supported metafeatures are type, typename and elementType for expressions, attributes and operations, owner and ownername for attributes and operations, and name for classes, types, operations, use cases, attributes and references. An example rule is:

```
for _1 : _2 do _3 |-->for (_2'elementType _1 : _2) { _3}
```

In general, this mechanism provides access to any feature of the abstract syntax of UML-RSDS class diagrams.

The var'name notation can also be used when name.cstl is an alternative/additional  $\mathcal{CSTL}$  file which can be used to map var. For example, the Swift code generation rule:

```
Class::
```

```
class _1 { _2 } |-->protocol _1 { _2'cgprotocol }\n<when> _1 interface
```

maps the content of an interface using the rules of *cgprotocol.cstl*, instead of the *cg.cstl* file. This is necessary, because Swift interface definitions are restricted in their formats.

Another point to notice in the above example rules is the use of a library, called *Ocl*, which contains definitions of specialised operations such as *String stringSubtract(String s1, String s2)*. The developer of the code generator needs to provide definitions of these operations in the target language – or reuse them from a similar language. For example, the *stringSubtract* operation is already defined in *Ocl* libraries for Java 6. Java 7 and other languages.

## 3 Using CSTL specifications

The LHS of rules have a fixed format and these cannot be varied. The files output/cgJava8.cstl and output/cgSwift.cstl give examples of all the available rules. The RHS and conditions of rules can be varied to generate text in many different languages.

Write your main translation rules in a file output/cg.cstl. The syntax of this can be checked by loading it into the Agile UML tools via the File menu option  $Load\ transformation \Rightarrow Load\ CSTL$ . Additional CSTL files in the output directory are also loaded by this step and can be invoked within cg.cstl by their names.

If you have defined a class diagram model, this can be translated to code via cg.cstl by selecting the Build menu option  $Use\ CSTL\ specification$ . The output text is stored in output/cgout.txt.

For some languages, multiple output files need to be produced from the same source model, for example, C header and code files. Separate code generator  $\mathcal{CSTL}$  files should be written to produce each output.

The latest version of the tools can be obtained from: https://nms.kcl.ac.uk/kevin.lano/uml2web/.