

# Itic notes

Vincent Siebler

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# 1 Motivation

Fix  $0 \leq m \leq n \leq \infty$ . An  $(n, m)$  category is a "category-like" structure consisting of a class of objects, notions of 1- morphism, 2-morphism, ... ,  $n$ -morphism (i.e.  $k$ -morphsim  $0 < k \leq n$ ) with a "suitable composition law" (satisfying "suitable axioms") and such that  $\forall m < k \leq n$  the  $k$ -morphisms are "invertible".

- $(0, 0)$ -cat. = set
- $(1, 0)$ -cat. = groupoid, i.e. 1-groupoid
- $(1, 1)$ -cat. = category, i.e. 1-category
- $(2, 0)$ -cat. = 2-groupoids
- $(2, 1)$ -cat.
- $(2, 2)$ -cat. = 2-categories
- $\vdots$
- $(n, 0)$ -cat. =  $n$ -groupoid
- $(n, n)$ -cat. =  $n$ -cat.
- $\vdots$
- $(\infty, 0)$  =  $\infty$ -groupoids
- $(\infty, 1)$  =  $\infty$ -categories (see. Boardman-Vogt)

**Reminder 1.1.** A map  $f: X \rightarrow Y$  between topological spaces is a weak homotopy equivalence if  $\forall x \in X, \forall n \in \mathbb{N}$  the map

$$\pi_n(f): \pi_n(X, x) \rightarrow \pi_n(Y, f(x))$$

is a bijection.

**Theorem 1.2** ( Grothendieck's Homotopy Hypothesis ). *There is an  $(\infty, 1)$ -category of topological spaces up to weak homotopy equivalence and there is an  $(\infty, 1)$ -category of  $\infty$ -groupoids up to equivalence. There is furthermore an  $\infty$ -functor assigning to each topological space  $X$  its Poincare  $\infty$ -groupoid  $\pi_\infty(X)$ , this is an equivalence.*

**Remark 1.3.** Let  $\mathcal{C}$  be an  $(\infty, 1)$ -category, then for all  $X, Y \in \text{Ob}(\mathcal{C})$  we get that  $\text{Hom}_{\mathcal{C}}(X, Y)$  is an  $\infty$ -groupoid/ "space". We have the homotopy category of  $\mathcal{C}$  denoted by  $\text{Ho}(\mathcal{C})$  whose objects are those of  $\mathcal{C}$  and for all objects  $X, Y \in \text{Ho}(\mathcal{C})$  we have that  $\text{Hom}_{\text{Ho}(\mathcal{C})}(X, Y) = \pi_0(\text{Hom}_{\mathcal{C}}(X, Y))$ .

**Warning! 1.4.** The passage from  $\mathcal{C}$  often results in a tremendous loss of information, that is essential for various purposes.

- Computing co-/limits within  $\mathcal{C}$ .
- Computing co-/limits with  $\mathcal{C}$ .

- Define invariants associated to  $\mathcal{C}$  ( f.e. Hochschild cohomology).

**Reminder 1.5.** *Many important 1-categories arise as homotopy categories of genuine  $(\infty, 1)$ -categories, that is derived categories Recall for  $R$ : ring,  $\text{mod}_R$  its category of right  $R$ -modules and  $\text{Ch}(\text{mod}_R)$  the category of chain complexes in  $\text{mod}_R$ , that a morphism of chain complexes  $f^\bullet: X^\bullet \rightarrow Y^\bullet$  in  $\text{Ch}(\text{mod}_R)$  is a quasi-isomorphism if  $\forall n \in \mathbb{Z}$  we have that  $H_n(f^\bullet) = H_n(X^\bullet) \xrightarrow{\sim} H_n(Y^\bullet)$  is an isomorphism. The derived category is defined as follows  $D(\text{mod}_R) := \text{Ch}(\text{mod}_R)[qiso^{-1}]$  the localisation at the quasi-isomorphisms. Furthermore we have that  $\text{Ho}(\mathcal{D}(\text{mod}_R)) = D(\text{mod}_R)$ . In the first case we obtain it by building it from the ground up so to say and in the second case we obtain by forgetting information from a higher structure.*

**Warning! 1.6.** The homotopy theory of  $(\infty, 1)$ -categories has many equivalent (Quillen) implementations:

- Topological categories (Ilias)
- Simplicial categories (Bergner)
- Complete Segal spaces (Rezk)
- Relative categories (Barwick-Kan)
- Pre-derivations
- $\infty$ -categories ( Joyal, Lurie )

In the  $k$ -linear setting, for  $k$  a field we have:

1. Differentially graded  $k$ -categories
2.  $A_\infty$ -categories (Lefèvre-Hasegawa)

The Plan for the lecture is to start of with investigating 2-categories, then give definition and examples of  $\infty$ -categories, then do enriched category theory and end on the homotopy theory of  $\infty$ -categories.

## 2 Heuristics: Isomorphism vs. Equivalence

Makkai's Principle of Isomorphism (1998) says:

"All grammatically correct properties about objects in a fixed category are to be invariant under isomorphism."

Fix a category  $\mathcal{C}$  and a small category  $A$  and take the functor category  $\text{Fun}(A, \mathcal{C})$ , that is  $A$ -shaped diagrams in  $\mathcal{C}$ . For  $X \in \mathcal{C}$  define:

$$\lim_{a \in A} \text{Hom}_{\mathcal{C}}(X, D(a)) := \{(p_a : X \rightarrow D(a))_{a \in A} \mid \begin{array}{ccc} X & & \\ f_a \downarrow & \searrow f_b & \\ D(a) & \longrightarrow & D(b) \end{array} \quad \forall f : a \rightarrow b \text{ in } A\}$$

For  $\phi : Y \rightarrow X$  define the function

$$\begin{aligned} \phi^* : \lim_{a \in A} \text{Hom}_{\mathcal{C}}(X, D(a)) &\rightarrow \lim_{a \in A} \text{Hom}_{\mathcal{C}}(Y, D(a)) \\ p = (p_a : X \rightarrow D(a))_{a \in A} &\mapsto \phi^*(p) = (p_a \circ \phi : Y \rightarrow D(a))_{a \in A} \end{aligned}$$

Thus  $\lim_{a \in A} \text{Hom}_{\mathcal{C}}(-, D(a)) : \mathcal{C}^{\text{op}} \rightarrow \text{Set}$  is a presheaf of sets on  $\mathcal{C}$ .

**Definition 2.1.** A limit of a diagram  $D : A \rightarrow \mathcal{C}$  is a cone  $p \in \lim_{a \in A} \text{Hom}_{\mathcal{C}}(X, D(a))$  that is universal in the sense that  $\forall Y \in \mathcal{C}, \forall q \in \lim_{a \in A} \text{Hom}_{\mathcal{C}}(Y, D(a)) \exists ! \varphi : Y \rightarrow X$  such that  $\varphi^*(p) = q$ . We write  $\lim_{a \in A} D(a)$  for any limit of  $D$  (which may or may not exist).

**Reminder 2.2** (Yoneda Lemma). *Let  $X \in \mathcal{C}$  and let*

$$\begin{aligned} \nu : \text{Nat}(\text{Hom}_{\mathcal{C}}(-, X), \lim_{a \in A} \text{Hom}_{\mathcal{C}}(-, D(a))) &\xrightarrow{\sim} \lim_{a \in A} \text{Hom}_{\mathcal{C}}(X, D(a)) \\ \eta = (\eta_Y : \text{Hom}_{\mathcal{C}}(Y, X) \rightarrow \lim_{a \in \mathcal{C}} (Y, D(a)))_{Y \in \mathcal{C}} &\mapsto \eta_X(\text{id}_X) \\ (\eta_Y^p(p) := \varphi^*(p))_{Y \in \mathcal{C}} &\leftarrow p \end{aligned}$$

If  $p \in \lim_{a \in A} \text{Hom}_{\mathcal{C}}(X, D(a))$  is a limit of  $D : A \rightarrow \mathcal{C}$  then

$\eta^p : \text{Hom}_{\mathcal{C}}(-, X) \xrightarrow{\sim} \lim_{a \in A} \text{Hom}_{\mathcal{C}}(-, D(a))$  is a natural isomorphism. We furthermore obtain, that for an isomorphism  $\psi : Y \rightarrow X$  in  $\mathcal{C}$  we have

$$\begin{array}{ccc} & \text{Hom}_{\mathcal{C}}(-, X) & \\ \psi^* \nearrow & & \searrow \eta^p \\ \text{Hom}_{\mathcal{C}}(-, Y) & \xrightarrow{\sim} & \lim_{a \in A} \text{Hom}_{\mathcal{C}}(-, D(a)) \end{array}$$

**Example 2.3.** Let the following be a diagram in  $\mathcal{C}$

$$\begin{array}{ccc} & Y & \\ & \downarrow g & \\ X & \xrightarrow{f} & Z \end{array}$$

then the limit of the diagram (if it exists ) is called a pullback.

$$\begin{array}{ccccc}
 W & & & & \\
 & \searrow \exists! p & & \nearrow \forall f'' & \\
 & X \times_Z Y & \xrightarrow{f'} & Y & \\
 & \downarrow g' & & \downarrow g & \\
 & X & \xrightarrow{f} & Z & \\
 & \nearrow \forall g'' & & \nwarrow &
 \end{array}$$

For example if  $\mathcal{C} = \text{Set}$  then  $X \times_Z Y = \{(X, Y) \in X \times Y \mid f(x) = g(y)\}$ .

The commutativity condition takes place in  $\text{Hom}_{\mathcal{C}}(X \times_Z Y, Z) \ni g \circ f' = f \circ g'$   
Makkai's Principle of Equivalence All grammatically correct properties of objects in a fixed 2-category are to be invariant under equivalence.

**Remark 2.4.** We want to  $\text{Cat}$  be the strict 2-category of (small) categories with functors as 1-morphisms and natural transformations as 2-morphisms. Now natural transformation allow for a notion of equivalence of morphisms, that is in a 1-category we only knew what it meant for two morphisms to be equal, but now we can talk about two functors being naturally isomorphic given us a notion of equivalence of 1-morphisms, via the 2-morphisms.

**Definition/Proposition 2.5** (Godement Product). Consider natural transformations

$$\begin{array}{ccccc}
 & F_1 & & G_1 & \\
 \mathcal{C} & \xrightarrow{\quad} & \mathcal{D} & \xrightarrow{\quad} & \mathcal{E} \\
 & F_2 & & G_2 & \\
 & \Downarrow \alpha & & \Downarrow \beta &
 \end{array}$$

Their Godement product is the natural transformation.

$$\begin{array}{ccc}
 & G_1 \circ F_1 & \\
 \mathcal{C} & \xrightarrow{\quad} & \mathcal{E} \\
 & G_2 \circ F_2 & \\
 & \Downarrow \beta * \alpha &
 \end{array}$$

Let  $X \in \mathcal{C}$ , we obtain the following diagram

$$\begin{array}{ccccc}
 F_1(X) & & G_1(F_1(X)) & \xrightarrow{\beta_{F_1(X)}} & G_2(F_1(X)) \\
 \downarrow \alpha_x & & \downarrow G_1(\alpha_X) & \searrow (\beta * \alpha)_X & \downarrow G_2(\alpha_X) \\
 F_2(X) & & G_1(F_2(X)) & \xrightarrow{\beta_{F_2(X)}} & G_2(F_2(X))
 \end{array}$$

in  $\mathcal{D}$ .

*Proof.* We show that  $\beta * \alpha: G_1 \circ F_1 \Rightarrow g_2 \circ F_2$  is indeed a natural transformation. For that we take the following diagram

$$\begin{array}{ccccccc}
X & & G_1(F_1(X)) & \xrightarrow{G_1(\alpha_X)} & G_1(F_2(X)) & \xrightarrow{\beta_{F_2(X)}} & G_2(F_2(X)) \\
\downarrow f & & \downarrow G_1(F_1(f)) & & \downarrow G_1(F_2(f)) & & \downarrow G_2(F_2(f)) \\
Y & & G_1(F_1(Y)) & \xrightarrow{G_1(\alpha_1)} & G_1(F_2(Y)) & \xrightarrow{\beta_{F_2(Y)}} & G_2(F_2(Y))
\end{array}$$

□

**Proposition 2.6.** *Consider natural transformations*

$$\begin{array}{ccccc}
& F_1 & & G_1 & \\
& \downarrow \alpha & & \downarrow \beta & \\
\mathcal{C} & \xrightarrow{F_2} & \mathcal{D} & \xrightarrow{G_2} & \mathcal{E} \\
& \downarrow \gamma & & \downarrow \delta & \\
& F_3 & & G_3 & 
\end{array}$$

Then  $(\delta\beta) * (\gamma\alpha) = (\delta * \gamma) \circ (\beta * \alpha)$ .

*Proof.* Let  $X \in \mathcal{C}$

$$\begin{array}{c}
\begin{array}{ccccccc}
G_1(F_1(X)) & \xrightarrow{\beta_{F_1(X)}} & G_2(F_1(X)) & \xrightarrow{\delta_{F_1(X)}} & G_3(F_1(X)) & & \\
\downarrow G_1(\alpha_X) & \searrow (\beta * \alpha)_X & \downarrow G_2(\alpha_X) & & \downarrow G_3(\alpha_X) & & \\
G_1(F_2(X)) & \xrightarrow{\beta_{F_2(X)}} & G_2(F_2(X)) & \xrightarrow{\delta_{F_2(X)}} & G_3(F_2(X)) & & \\
\downarrow G_1(\gamma_X) & \searrow (\delta * \gamma)_X & \downarrow G_2(\gamma_X) & & \downarrow G_3(\gamma_X) & & \\
G_1(F_3(X)) & \xrightarrow{\beta_{F_3(X)}} & G_2(F_3(X)) & \xrightarrow{\delta_{F_3(X)}} & G_3(F_3(X)) & & 
\end{array} \\
\begin{array}{c}
\text{Left side: } G_1(\gamma_X \alpha_X) \text{ (curved arrow from } G_1(F_1(X)) \text{ to } G_1(F_3(X))) \\
\text{Right side: } G_3(\gamma_X \alpha_X) \text{ (curved arrow from } G_3(F_1(X)) \text{ to } G_3(F_3(X)))
\end{array}
\end{array}$$

Now the long diagonal of the diagram corresponds to  $(\delta * \gamma) \circ (\beta * \alpha)$  and the outer large square to  $(\delta \circ \beta) * (\gamma \circ \alpha)$ . □

**Definition 2.7.** The Godement products bellow are called whickerings

$$\begin{array}{ccc}
& & G_1 \circ F & \\
& G_1 & \searrow & \\
\mathcal{C} & \xrightarrow{F} & \mathcal{D} & \xrightarrow{\beta} & \mathcal{E} \\
& G_2 & \nearrow & \\
& & G_2 \circ F & 
\end{array}$$

$$\begin{array}{ccccc}
& & & & G \circ F_1 \\
& & & & \curvearrowright \\
& & & & \parallel \text{ id}_G * \alpha \\
& & & & \curvearrowright \\
& & & & G \circ F_2 \\
\mathcal{C} & \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \alpha \\ \xrightarrow{G_2} \end{array} & \mathcal{D} & \xrightarrow{G} & \mathcal{E} \mapsto \mathcal{C}
\end{array}$$

**Construction 2.8.** Given a cospan of groupoids

$$\begin{array}{ccc}
& \mathcal{B} & \\
& \downarrow G & \\
\mathcal{A} & \xrightarrow{F} & \mathcal{C}
\end{array}$$

its 2-pullback is the diagonal of groupoids.

$$\begin{array}{ccc}
\mathcal{A} \times_{\mathcal{C}} \mathcal{B} & \xrightarrow{\pi_{\mathcal{B}}} & \mathcal{B} \\
\downarrow \pi_{\mathcal{A}} & & \downarrow G \\
\mathcal{A} & \xrightarrow{F} & \mathcal{C}
\end{array}$$

The objects are given as  $\text{Ob}(\mathcal{A} \times_{\mathcal{C}} \mathcal{B}) = (a \in \mathcal{A}, b \in \mathcal{B}, \varphi: F(a) \xrightarrow{\sim} G(b) \text{ in } \mathcal{C})$  and morphisms are given by tuples of morphisms  $(u, v): (a, b, \varphi) \rightarrow (a', b', \varphi')$ , where  $u: a \rightarrow a'$  and  $v: b \rightarrow b'$  are morphisms in the respective groupoids, such that the following square commutes

$$\begin{array}{ccc}
F(a) & \xrightarrow[\sim]{\varphi} & G(b) \\
F(a) \downarrow & & \downarrow G(v) \\
F(a') & \xrightarrow[\sim]{\varphi} & G(b')
\end{array}$$