Epidemic Algorithms for Replicated Database Maintenance

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EPIDEMIC ALGORITHMS FOR REPLICATED DATABASE MAINTENANCE Alan Demers, Dan Greene, Carl Hauser, Wes Irish, John Larson, Scott Shenker, Howard Sturgis, Dan Swinehart, and Doug Terry

Epidemic Algorithms For Replicated Database Maintenance

• Alan Demers Retired Professor at Cornell University

• **Dan Greene** At Xerox PARC – Vehicle networks

• Carl Hauser Associate Professor, Washington State University

• Wes Irish Coyote Hill Consulting

• **Scott Shenker** Professor at UCBerkeley

• **Doug Terry** Microsoft Research

John Larson, Howard Sturgis, Dan Swinehart

Summary of the Research

- Database management for distributed systems
 - Consistent data records
- 3 methods
 - Direct Mail
 - Anti-Entropy
 - Rumor Mongering
- CAP Theorem
- Real world applications
 - Vegvisir blockchain
 - Amazon

Research Motivation

- Clearinghouse servers on Xerox Corporate Internet (CIN)
 - Hundreds of ethernets connected by gateways and phone lines
 - Ex Message: Japan -> Europe goes through 14 gateways and 7 phone lines
 - Organized by Hierarchical name (domains)
 - Remailing Inefficiency during disagreement among participants

Points of Differentiation

- Eventual delivery of repeated messages and do not require data structures at one server to describe information held at other servers
- Algorithm are randomized

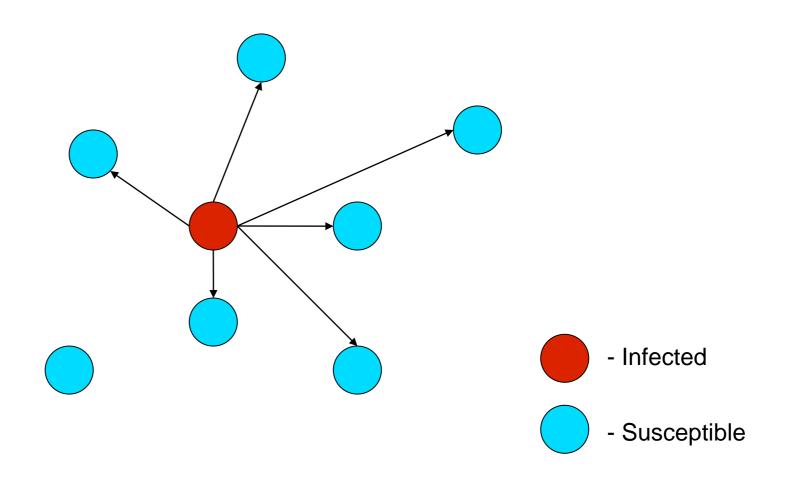
Vocabulary

- Infected Knows the update and *spreads* it
- Susceptible Does not know the update
- **Removed** Knows the update but *not* able to spread it anymore /
- Push Tells an updates to another node
- Pull Asks for an update from another node

Direct Mail

Direct Mail – Sends update to all nodes in the network

• Traffic proportional to the number of sites * average distance between sites



Direct Mail

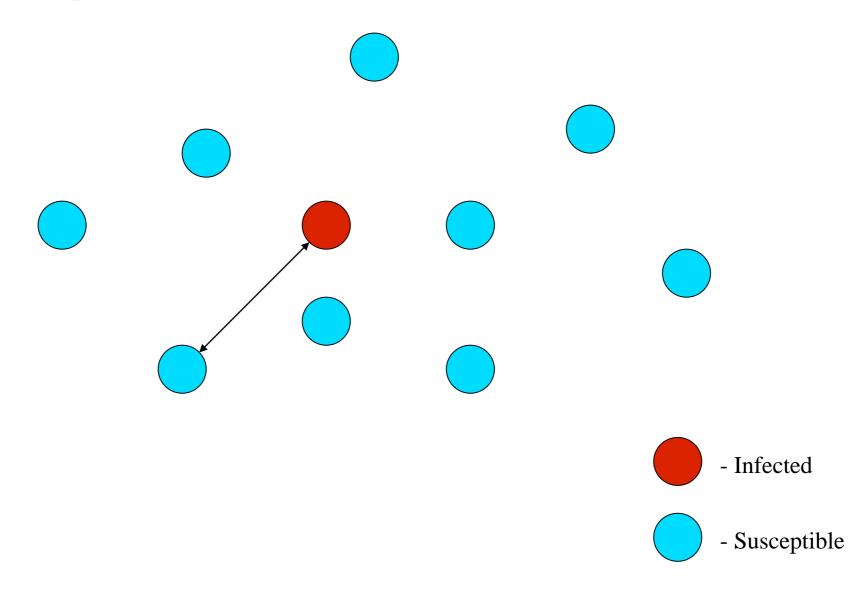
Failure Modes

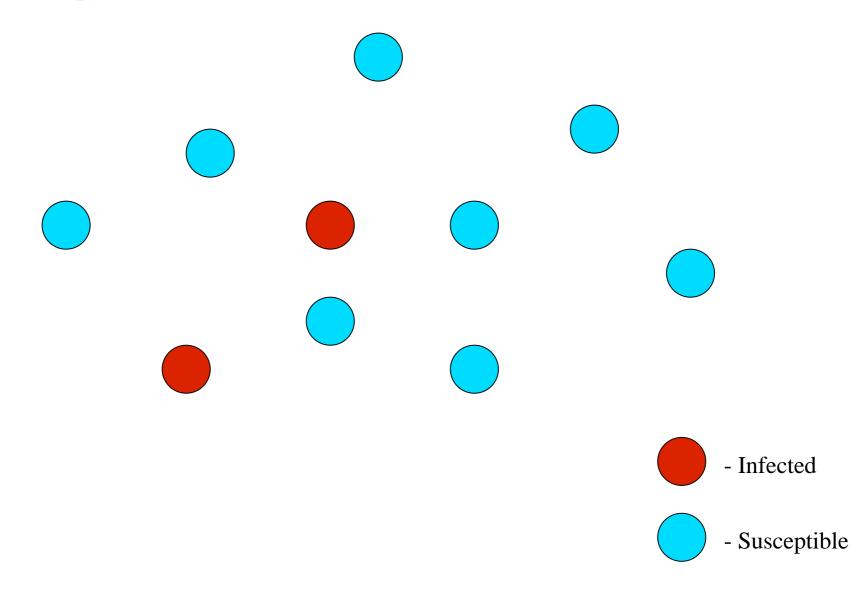
- Message discarded for nodes
 - Que overflows
 - Extended period of inaccessibility

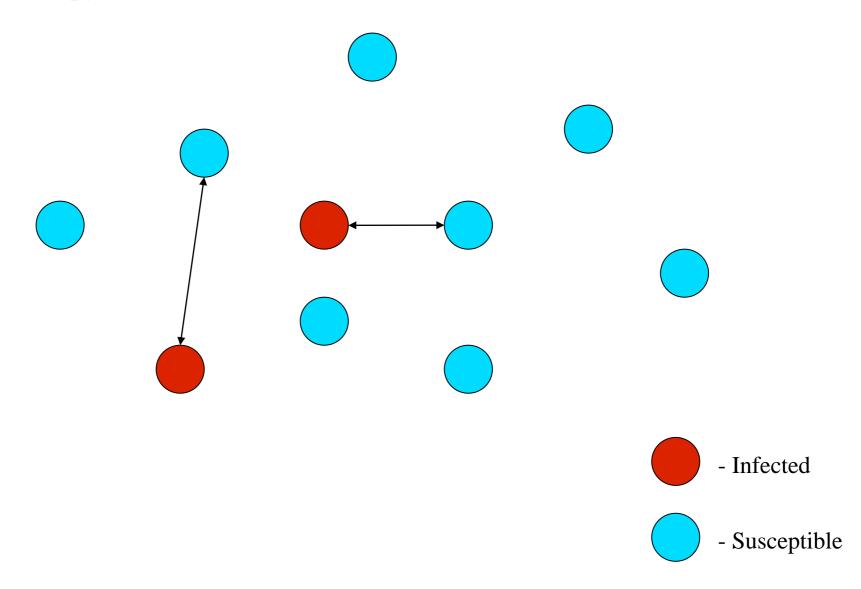
Anti-Entropy

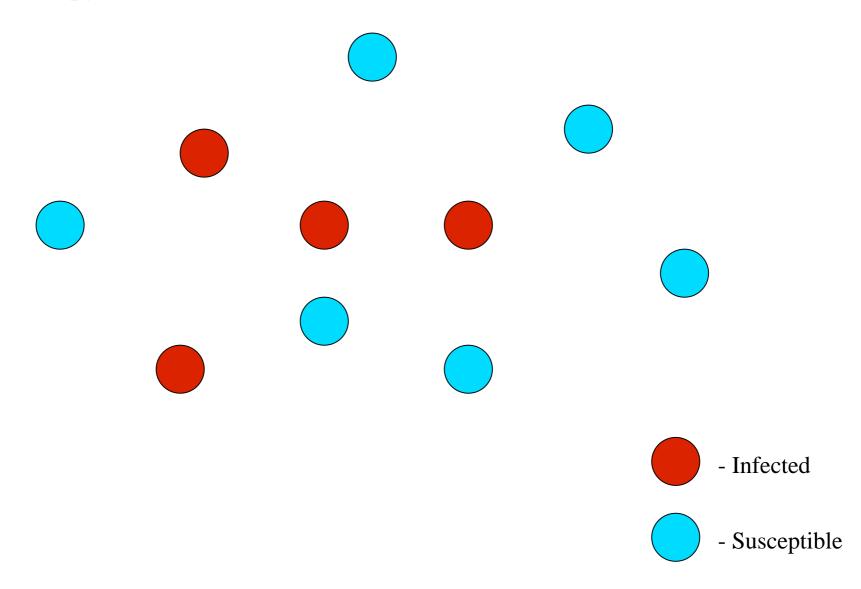
Anti-Entropy – Nodes exchange messages with a random node through the methods below:

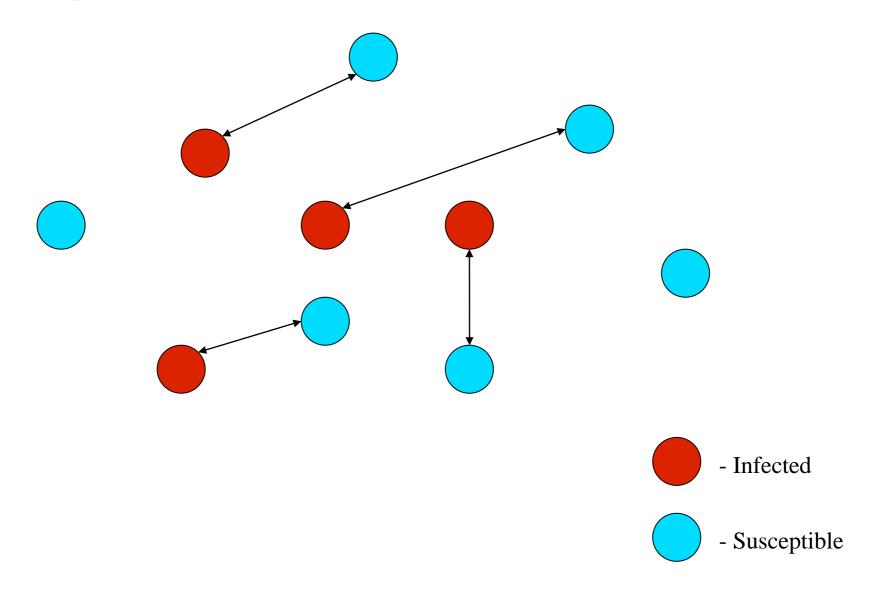
- Pull Grows fast but slows down overtime
- Push Grows slowly but speeds up overtime
- Push-pull Most efficient and every node receives the message

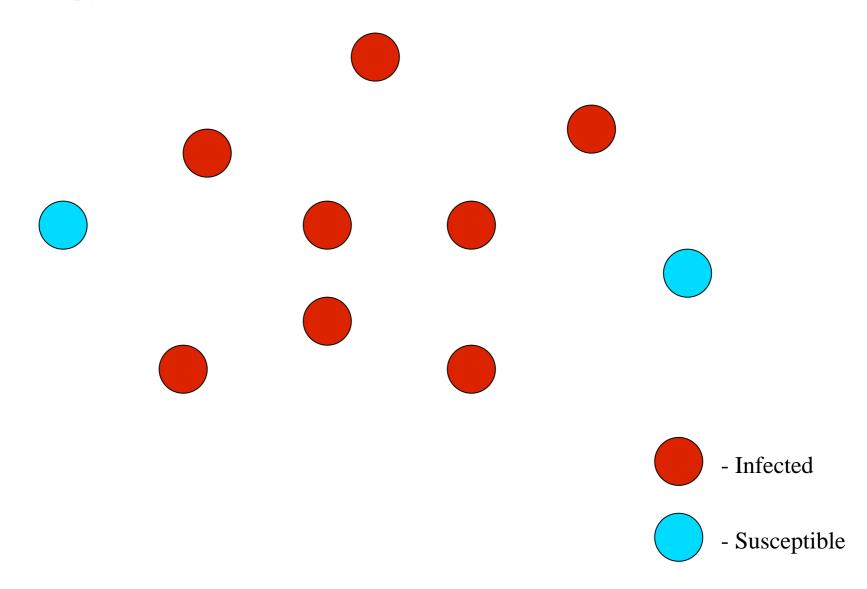












Anti-Entropy

Pro:

• Eventually everyone receives the message

Con:

• Large overhead due to external requests for updates

Rumor Mongering

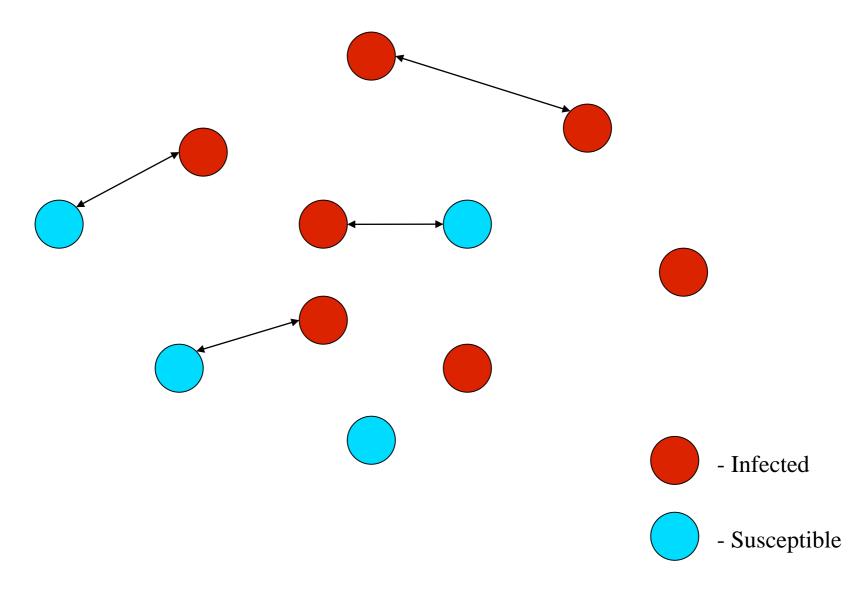
Rumor mongering – Optimized algorithm for spreading messages

- When a node receives a new update (rumor)
- Periodically choose another site at random to infect other nodes
- When enough nodes have seen the rumor it is removed

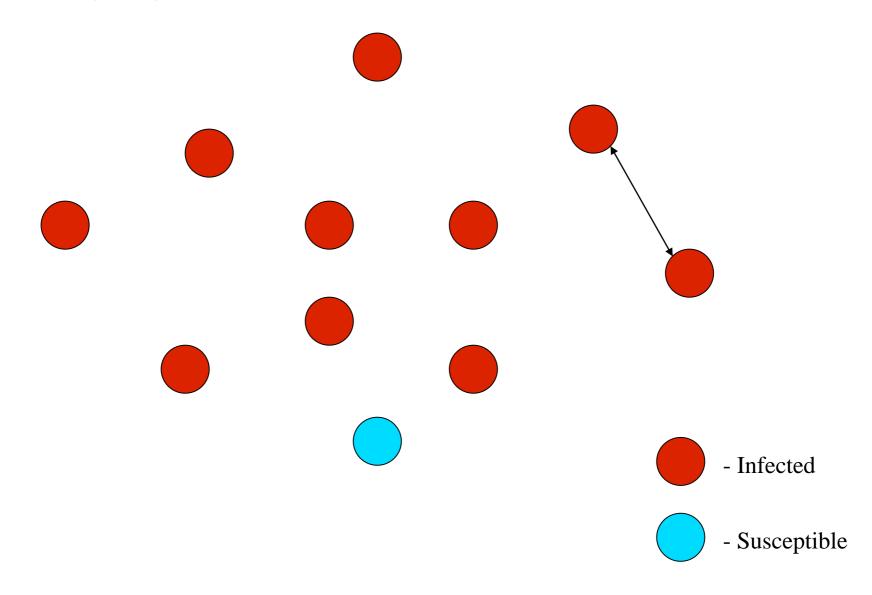
Problem of convergence

• Fix with anti-entropy combination

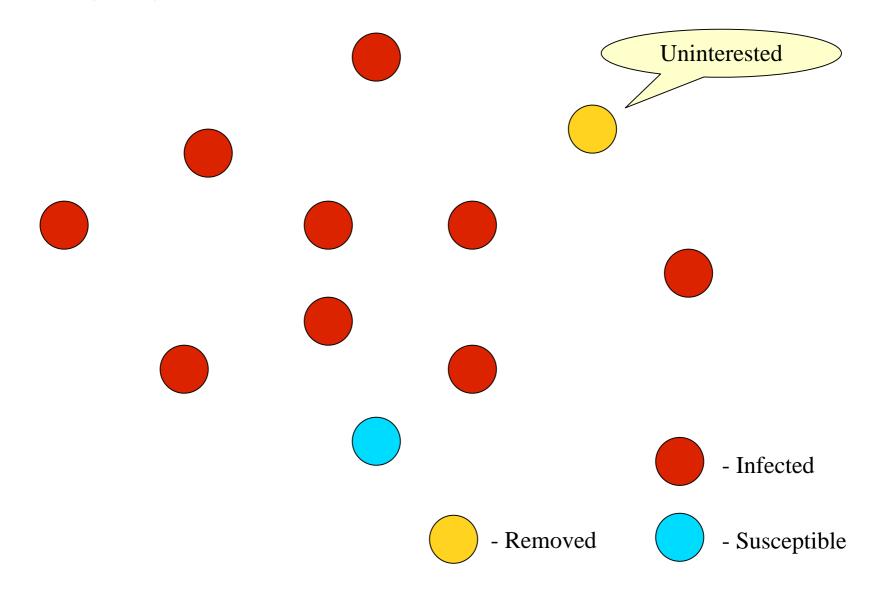
Rumor mongering



Rumor mongering



Rumor mongering



Points of differentiation

Death certificates

- Shows when a node is decommissioned
- Verified with a timestamp

Spatial distribution

• Favors sending updates to closest nodes first

Anti-Entropy Results

Table 4. Simulation results for anti-entropy, no connection limit.

Spatial	t_{last}	tare	Compare Traffic		Update Traffic	
Distribution		}	Average	Bushey	Average	Bushey
uniform	7.81	5.27	5.87	75.74	5.85	74.43
a = 1.2	10.04	6.29	2.00	11.19	2.61	17.52
a = 1.4	10.31	6.39	1.93	8.77	2.49	14.10
a = 1.6	10.94	6.70	1.71	5.72	2.27	10.88
a = 1.8	11.97	7.21	1.52	3.74	2.07	7.68
a = 2.0	13.32	7.76	1.36	2.38	1.89	5.87

Table 5. Simulation results for anti-entropy, connection limit 1.

Spatial	tinst	tave	Compare Traffic		Update Traffic	
Distribution			Average	Bushey	Average	Bushey
uniform	11.00	6.97	3.71	47.54	5.83	75.17
a = 1.2	16.89	9.92	1.14	6.39	2.69	18.03
a = 1.4	17.34	10.15	1.08	4.68	2.55	13.68
a = 1.6	19.06	11.06	0.94	2.90	2.32	10.20
a = 1.8	21.46	12.37	0.82	1.68	2.12	7.03
a = 2.0	24.64	14.14	0.72	0.94	1.94	4.85

Rumor Mongering Results

Table 6. Simulation results for push-pull rumor mongering.

Spatial	k	tiast	tave	Compare Traffic		Update Traffic	
Dist				Avg	Bushey	Avg	Bushey
uniform	4	7.83	5.32	8.87	114.0	5.84	75.87
a = 1.2	6	10.14	6.33	3.20	18.0	2.60	17.25
a = 1.4	5	10.27	6.31	2.86	13.0	2.49	14.05
a = 1.6	8	11.24	6.90	2.94	9.80	2.27	10.54
a = 1.8	7	12.04	7.24	2.40	5.91	2.08	7.69
a = 2.0	6	13.09	7.74	1.99	3.44	1.90	5.94

CAP Theorem

Only 2/3 are achievable

- Consistency Every node has the most recent message
- Accessibility Every node receives a message but no guarantee that it is the most recent

 Partition Tolerance – System continues to operate even if messages are lost

Applications of the Research

- 1. Vegvisir Agriculture specific blockchain that reconciles with random nodes within a specific range
- 2. Amazon S3 storage system uses gossip to disseminate information

Questions?