

### Problem 5

Tollbooths are 75 km apart, and the cars propagate at 100km/hr. A tollbooth services a car at a rate of one car every 12 seconds.

- a) There are ten cars. It takes 120 seconds, or 2 minutes, for the first tollbooth to service the 10 cars. Each of these cars has a propagation delay of 45 minutes (travel 75 km) before arriving at the second tollbooth. Thus, all the cars are lined up before the second tollbooth after 47 minutes. The whole process repeats itself for traveling between the second and third tollbooths. It also takes 2 minutes for the third tollbooth to service the 10 cars. Thus the total delay is 96 minutes.
- b) Delay between tollbooths is  $8 \times 12$  seconds plus 45 minutes, i.e., 46 minutes and 36 seconds. The total delay is twice this amount plus  $8 \times 12$  seconds, i.e., 94 minutes and 48 seconds.

### Problem 6

- a)  $d_{prop} = m / s$  seconds.
- b)  $d_{trans} = L / R$  seconds.
- c)  $d_{end-to-end} = (m / s + L / R)$  seconds.
- d) The bit is just leaving Host A.
- e) The first bit is in the link and has not reached Host B.
- f) The first bit has reached Host B.
- g) Want

$$m = \frac{L}{R} s = \frac{120}{56 \times 10^3} (2.5 \times 10^8) = 536 \text{ km.}$$

### Problem 8

- a) 20 users can be supported.
- b)  $p = 0.1$ .
- c)  $\binom{120}{n} p^n (1-p)^{120-n}$ .
- d)  $1 - \sum_{n=0}^{20} \binom{120}{n} p^n (1-p)^{120-n}$ .

We use the central limit theorem to approximate this probability. Let  $X_j$  be independent random variables such that  $P(X_j = 1) = p$ .

$$P(\text{"21 or more users"}) = 1 - P\left(\sum_{j=1}^{120} X_j \leq 21\right)$$

### Problem 13

- a) The queuing delay is 0 for the first transmitted packet,  $L/R$  for the second transmitted packet, and generally,  $(n-1)L/R$  for the  $n^{th}$  transmitted packet. Thus, the average delay for the  $N$  packets is:

$$\begin{aligned} & (L/R + 2L/R + \dots + (N-1)L/R)/N \\ &= L/(RN) * (1 + 2 + \dots + (N-1)) \\ &= L/(RN) * N(N-1)/2 \\ &= LN(N-1)/(2RN) \\ &= (N-1)L/(2R) \end{aligned}$$

Note that here we used the well-known fact:

$$1 + 2 + \dots + N = N(N+1)/2$$

- b) It takes  $LN/R$  seconds to transmit the  $N$  packets. Thus, the buffer is empty when a each batch of  $N$  packets arrive. Thus, the average delay of a packet across all batches is the average delay within one batch, i.e.,  $(N-1)L/2R$ .

$$\begin{aligned} P\left(\sum_{j=1}^{120} X_j \leq 21\right) &= P\left(\frac{\sum_{j=1}^{120} X_j - 12}{\sqrt{120 \cdot 0.1 \cdot 0.9}} \leq \frac{9}{\sqrt{120 \cdot 0.1 \cdot 0.9}}\right) \\ &\approx P\left(Z \leq \frac{9}{3.286}\right) = P(Z \leq 2.74) \\ &= 0.997 \end{aligned}$$

when  $Z$  is a standard normal r.v. Thus  $P(\text{"21 or more users"}) \approx 0.003$ .

### Problem 25

- a) 160,000 bits
- b) 160,000 bits
- c) The bandwidth-delay product of a link is the maximum number of bits that can be in the link.
- d) the width of a bit = length of link / bandwidth-delay product, so 1 bit is 125 meters long, which is longer than a football field
- e)  $s/R$