

Transactions

Members:

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1. Four Non-Conflicting:

1. Display All Electronics Products from Products Table:

START TRANSACTION;

SELECT * FROM Products WHERE category = 'Electronics';

COMMIT;

product_id	stock	product_name	review_id	category	subcategory	description
1000	1	7 Laptop	4	Electronics	Computers	High-performance laptop
700	4	30 LED TV	7	Electronics	Television	Smart LED TV with HD display
50	7	5 Headphones	10	Electronics	Audio	Wireless headphones with noise cancellation
80	10	12 Rice Cooker	13	Electronics	Kitchen Appliances	Multi-functional rice cooker

4 rows in set (0.00 sec)

Explanation:-

This transaction only retrieves data from the Products table based on specific condition (category = 'Electronics') without modifying it, making it non-conflicting as it doesn't impact other transactions.

2. Update Stocks:

START TRANSACTION;

UPDATE Products SET stock = stock - 5 WHERE product_id = 1;

COMMIT;

```
mysql> UPDATE Products SET stock = stock - 5 WHERE product_id = 1;
Query OK, 1 row affected (0.00 sec)
Rows matched: 1 Changed: 1 Warnings: 0
```

Explanation:-

This transaction updates the stock of a specific product, which is independent of other product updates, ensuring concurrency without conflicts.

3. Update Price:

START TRANSACTION;

UPDATE Products SET price = 1200 WHERE product_id = 4;

COMMIT;

```
mysql> UPDATE Products SET price = 1200 WHERE product_id = 4;  
Query OK, 1 row affected (0.00 sec)  
Rows matched: 1 Changed: 1 Warnings: 0
```

Explanation:-

This transaction updates the price of a different product, operating on separate data from the previous transaction, thus non-conflicting.

4. Display Customers from New York:

START TRANSACTION;

SELECT * FROM customer WHERE city = 'New York';

COMMIT;

```
mysql> SELECT * FROM customer WHERE city = 'New York';  
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+  
| user_id | phone_number | password | first_name | last_name | name | flat_no | city | state | zip_code | address |  
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+  
| 1 | 1234567890 | password1 | John | Doe | John Doe | Apt 101 | New York | NY | 10001 | Apt 101, New York, NY 10001 |  
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+  
1 row in set (0.00 sec)
```

Explanation:-

This transaction retrieves customer data based on a specific condition (city = 'New York'), which is separate from product-related transactions, allowing it to execute concurrently without conflict.

2. Two Conflicting:

1. Two employees are updating the price of the same product at the same time without knowing of each other's actions.

Transaction A (Promotional Discount) :

START TRANSACTION;

```
SELECT price INTO @current_price FROM Products WHERE product_id = 1;

SET @new_price = @current_price * 0.90;

UPDATE Products SET price = @new_price WHERE product_id = 1;

COMMIT;
```

Transaction B (Cost Adjustment) :

```
START TRANSACTION;

SELECT price INTO @current_price FROM Products WHERE product_id = 1;

SET @new_price = @current_price + 5;

UPDATE Products SET price = @new_price WHERE product_id = 1;

COMMIT;
```

Explanation:

If these transactions are executed simultaneously and without proper isolation, one of the following might happen:

- **Employee A's** transaction executes first, setting the new price to \$900. Then, **Employee B's** transaction, which reads the old price of \$1000, executes and sets the new price to \$1005, effectively overwriting the promotional discount.
- Or vice versa, where **Employee B** goes first, followed by **Employee A**, and the price set by **Employee B** is inadvertently reduced by the promotional discount, resulting in a price not reflecting the updated supplier cost.

This is a classic **lost update problem**. One transaction's changes are lost because another transaction overwrites the changes without knowing the first transaction's actions.

2. Read-Write Conflict (Dirty Read)

Transaction-C: Read product information of product-1 and then update its stock.

```
START TRANSACTION;

SELECT * FROM Products WHERE product_id = 1;

-- Simulated delay to allow another transaction to intervene
```

DO SLEEP(5);

UPDATE Products SET stock = stock - 3 WHERE product_id = 1;

COMMIT;

Transaction-D: Simultaneously update the same product's stock

START TRANSACTION;

UPDATE Products SET stock = stock + 8 WHERE product_id = 1;

COMMIT;

Explanation:-

In this scenario, **Transaction-C** initially reads product information for product-1 and then updates its stock after a delay. However, before **Transaction-C** completes its update, **Transaction-D** simultaneously updates the stock of the same product.

The conflict arises due to the possibility of **Transaction-D** updating the stock of product-1 while **Transaction-C** is still processing. If **Transaction-D**'s update occurs before **Transaction-C** completes its update, **Transaction-C** may act upon outdated data (i.e Previous Stock) , resulting in a "dirty read" where **Transaction-C** makes decisions based on stale information. This can lead to inconsistencies in the stock quantity, as **Transaction-C** may unintentionally overwrite or conflict with the changes made by **Transaction-D**, potentially resulting in incorrect stock levels or lost updates.

When the admin updates the products table a conflict may arise on deciding the price for that we use lock in the products table to avoid such a conflict
Similar lock exist for payment table