

# Windows Heap Exploitation

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# Windows Heap Exploitation

- Windows Memory Allocator
  - NT heap
  - BackEnd
    - Exploitation
  - FrontEnd
    - LowFragmentationHeap
    - Exploitation

# Windows memory allocator

- Win 10 的 memory allocator 非常複雜，基本上主要有下列兩種
- Nt Heap
  - Default memory allocator
- SegmentHeap
  - Win10 中全新的 memory allocator 機制
  - 部分系統程式及 UWP 會使用

# Windows memory allocator

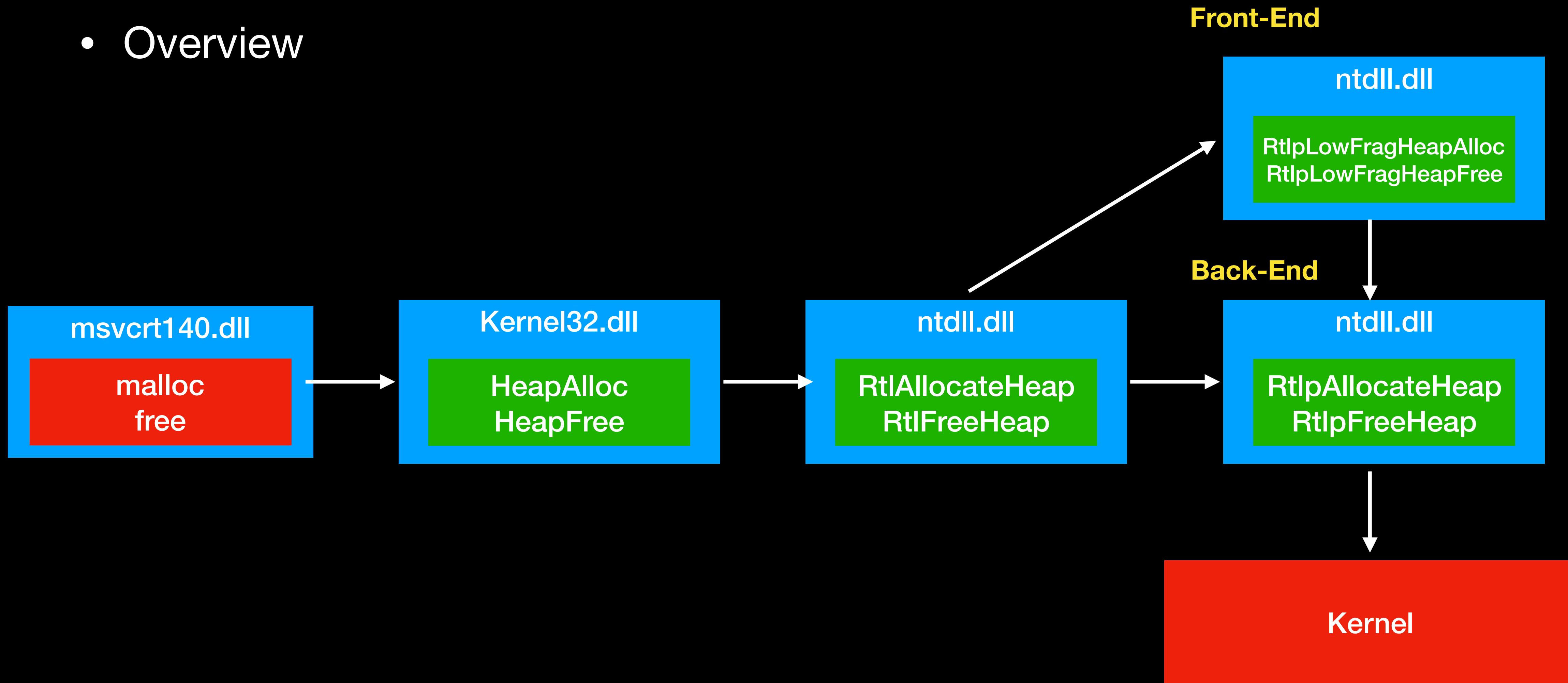
- Nt Heap 中又可分
  - 後端管理器 (Back-End)
  - 前端管理器 (Front-End)
    - LowFragmentationHeap
    - 基本上是為了不讓記憶體過於破碎，相同大小的記憶體區塊分配到一定數量後就會使用 (連續分配相同記憶體區塊 18 次)
    - Size <= 0x4000

# Windows memory allocator

- 目前會以 win 10 (1809) 為主
  - OS build - 17763.379
  - 之後版本可能不一定完全正確
  - 結構常常改動很大，有個概念後，到新的發行版本，需要自己去追一下結構

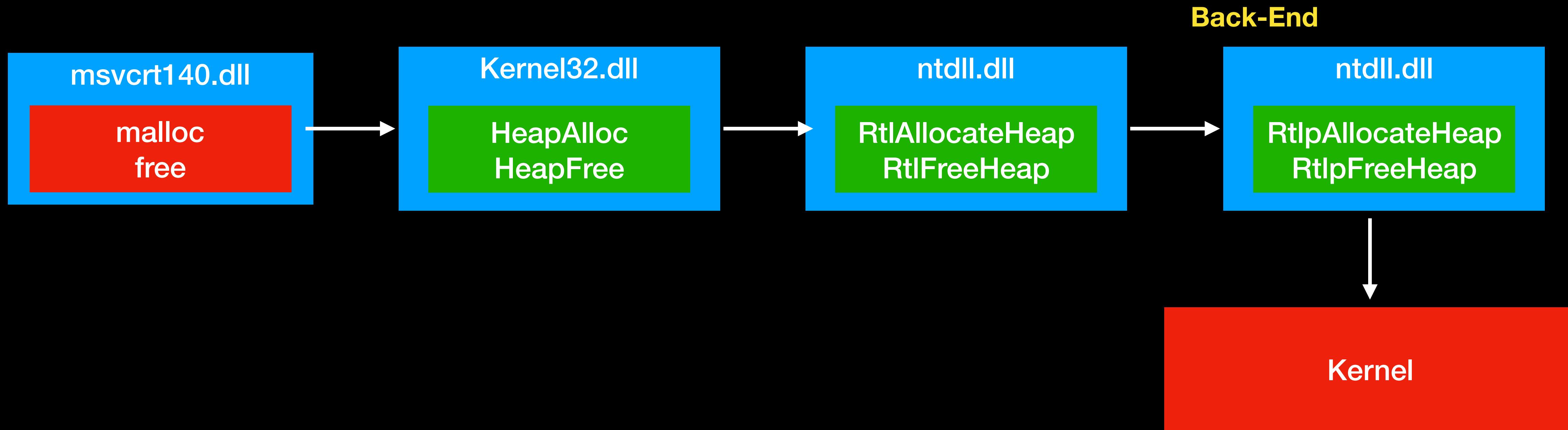
# Windows memory allocator

- Overview



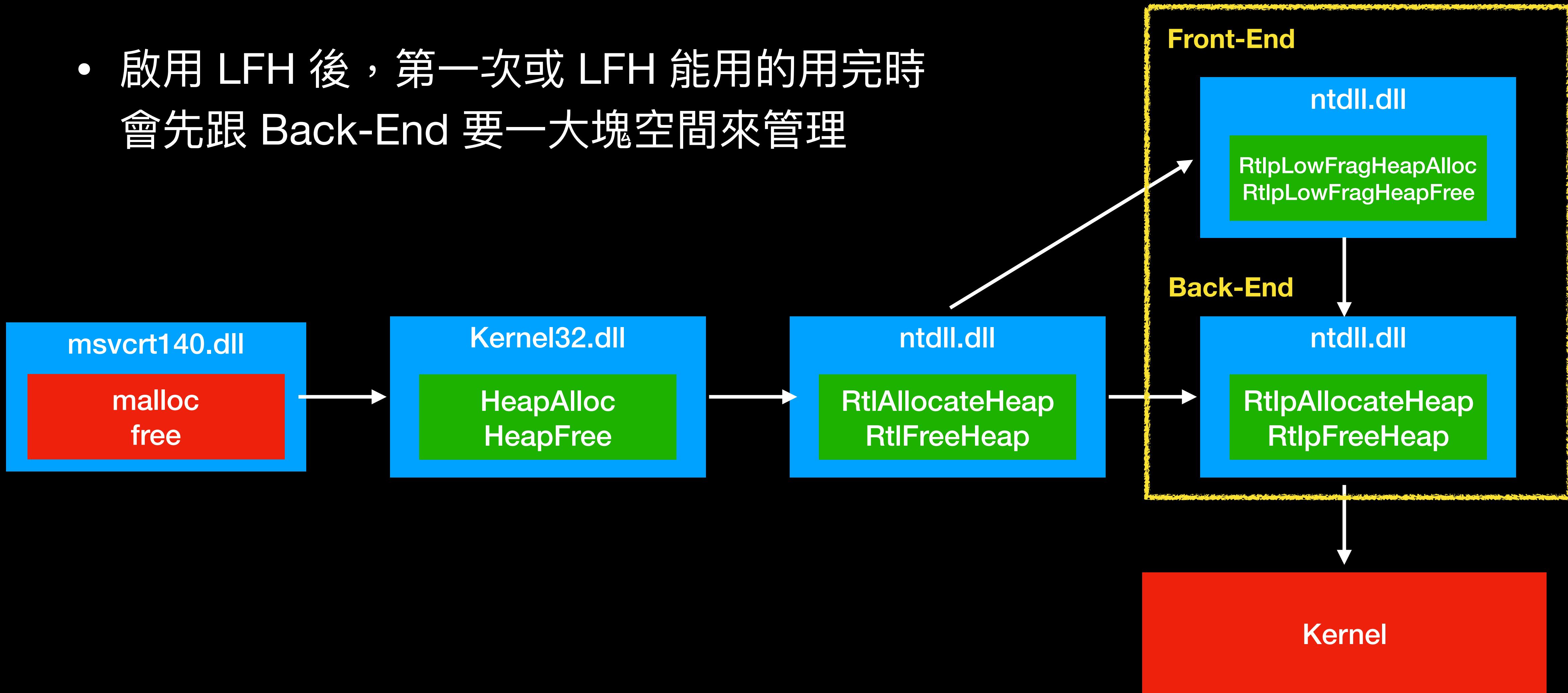
# Windows memory allocator

- 在 LFH 沒啟用下，我們呼叫 malloc



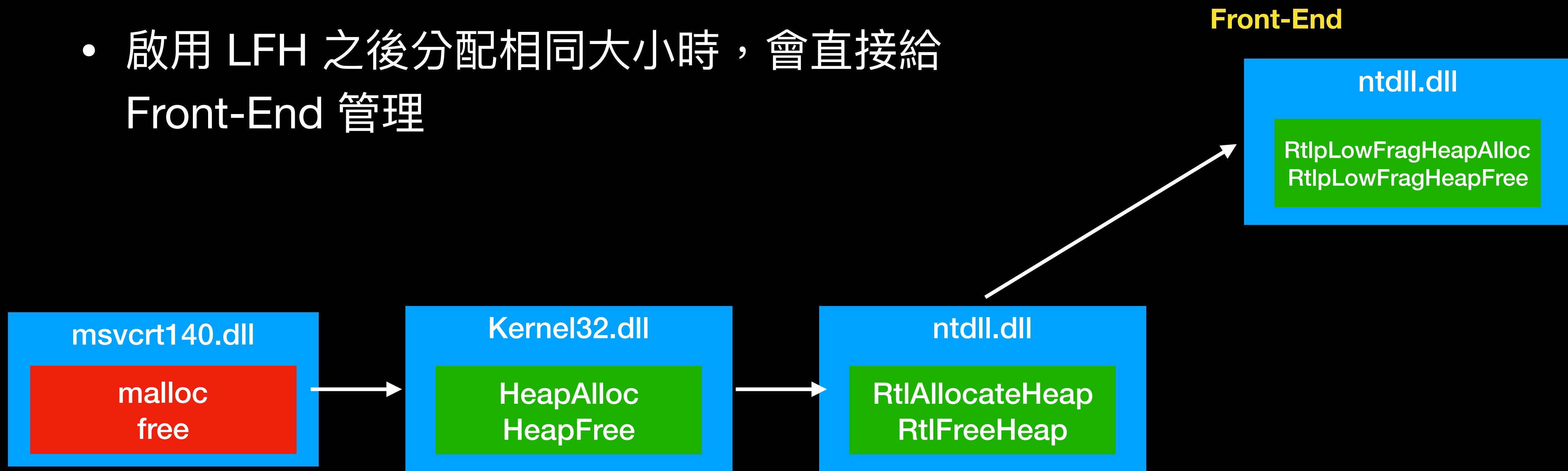
# Windows memory allocator

- 啟用 LFH 後，第一次或 LFH 能用的用完時會先跟 Back-End 要一大塊空間來管理



# Windows memory allocator

- 啟用 LFH 之後分配相同大小時，會直接給 Front-End 管理



# Windows memory allocator

- HEAP 可分為
  - Process Heap
  - Default heap
  - 整個 process 共享，呼叫 api 時會用到
    - 會存在 \_PEB 結構中
  - CRT 中的函式時也會用到
    - 存在 crt\_heap

# Windows memory allocator

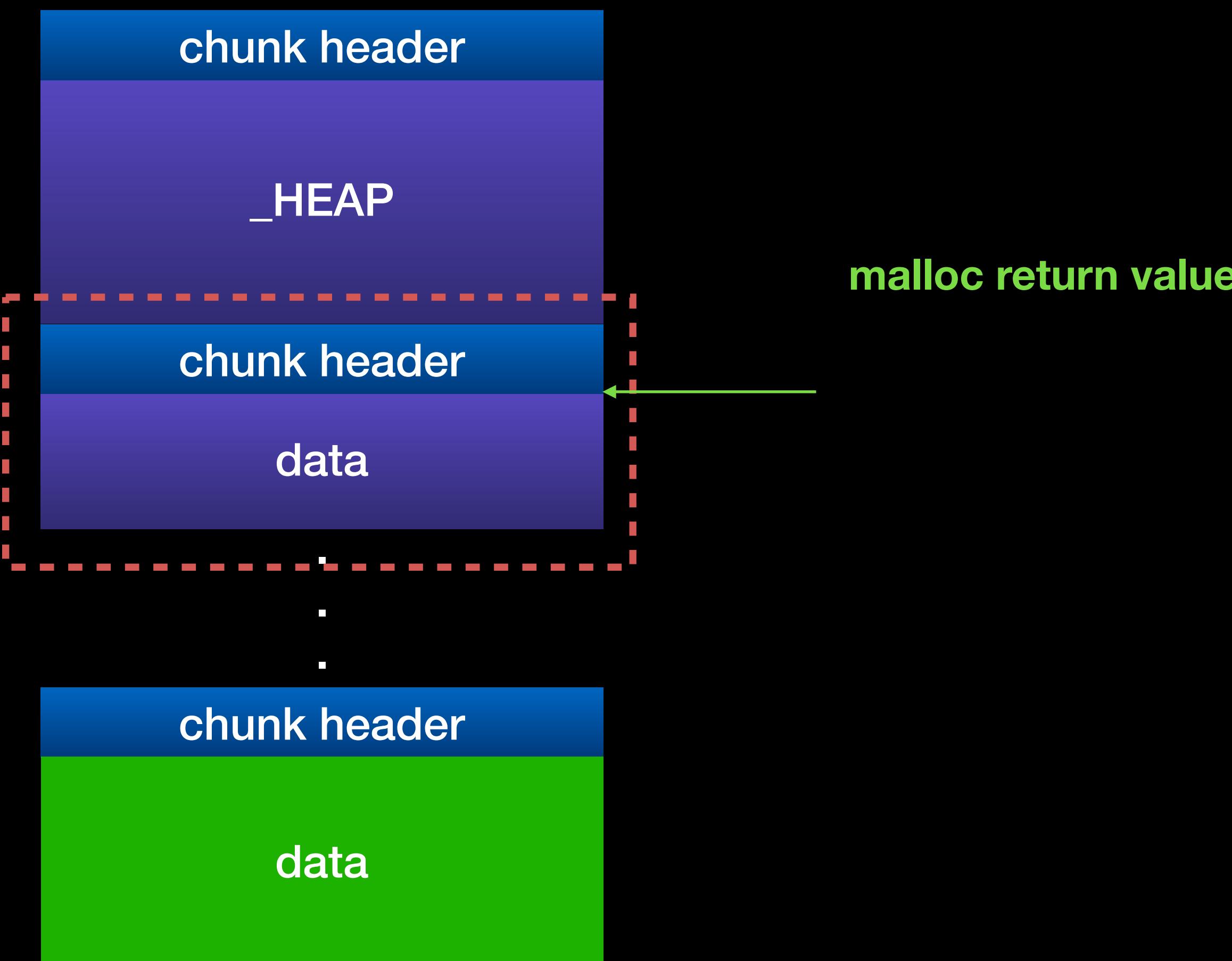
- HEAP 可分為
  - Private heap
  - 另外創建的 heap
    - HeapCreate

# Windows memory allocator

- Core data structure
  - `_HEAP_ENTRY` (chunk)
    - Malloc 出來的基本結構
    - 在前後端分配器中，長相會不太一樣，但都是同一個名字

# Windows memory allocator

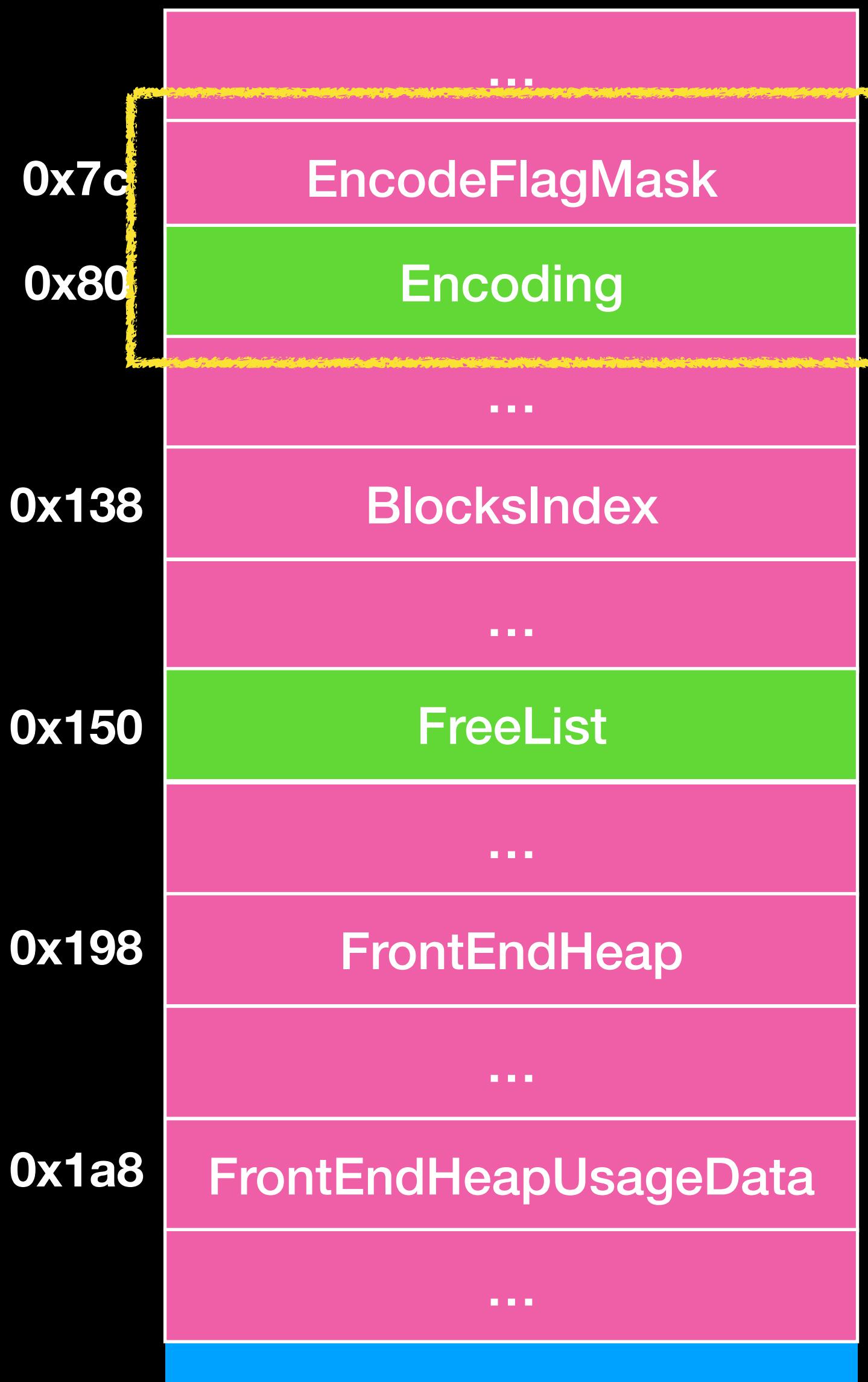
- `_HEAP_ENTRY` (chunk)



# Windows memory allocator

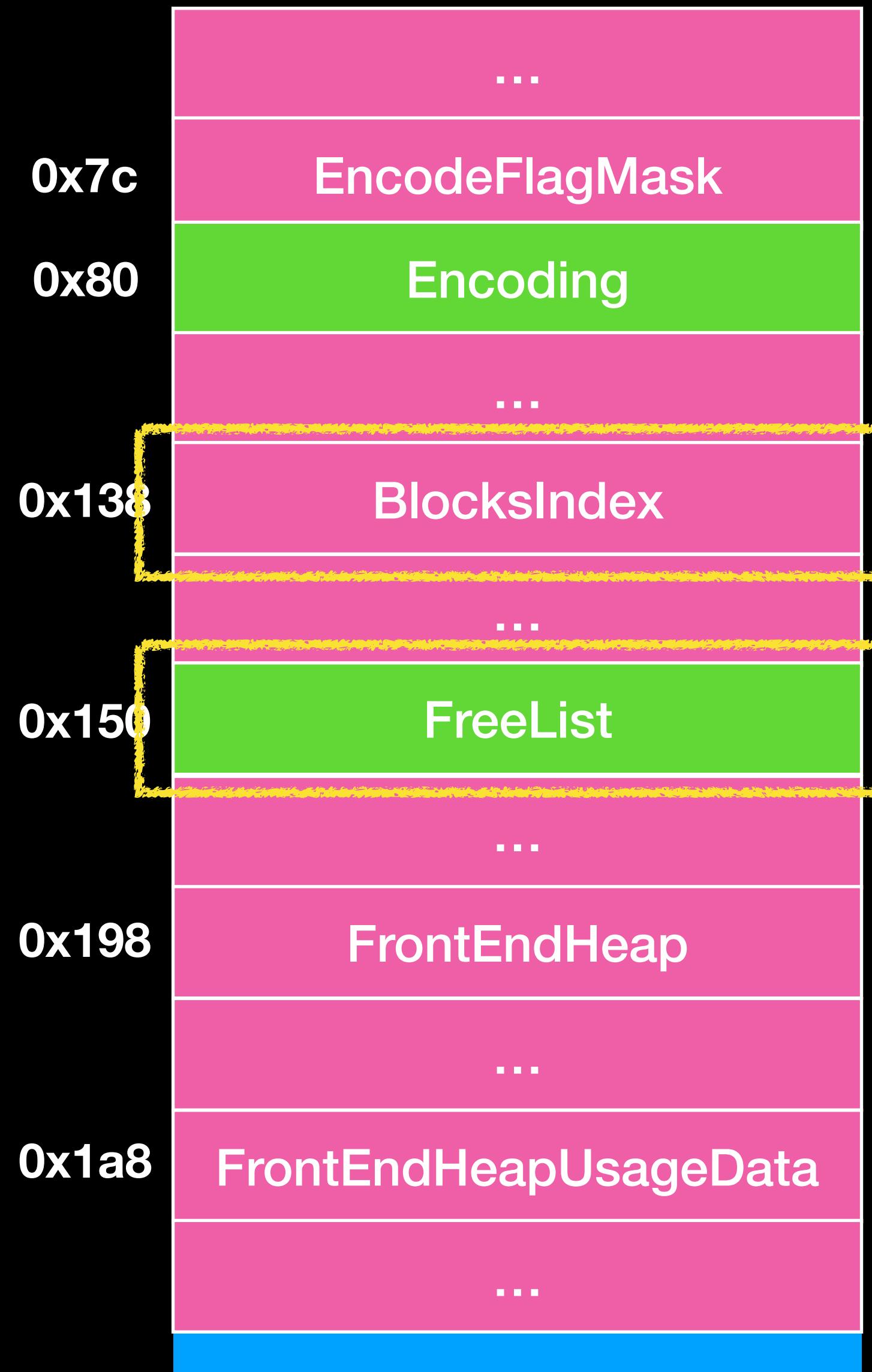
- Core data structure
  - \_HEAP
    - 整個 memory allocator 中最核心的結構，用來管理該 heap
    - 每個 Heap 都會有一個 \_HEAP 通常在 heap 開頭

# Windows memory allocator



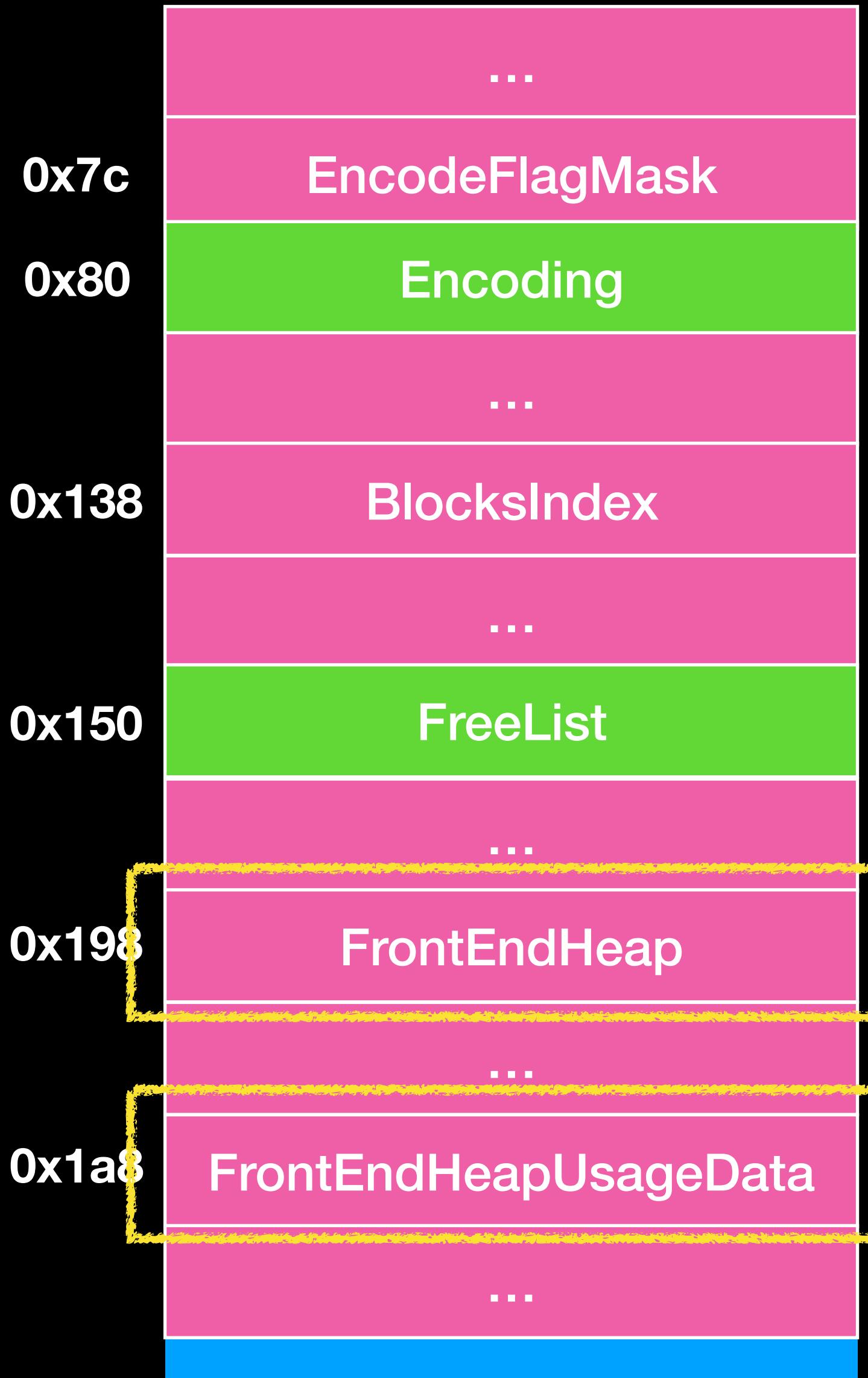
- `_HEAP`
- `EncodeFlagMask`
  - Heap 初始化後會設置為 0x10000 在判斷是否要 encode 該 heap 中 chunk 的 header
- `Encoding (_HEAP_ENTRY)`
  - 用來與 chunk header 做 xor 的 cookies

# Windows memory allocator



- `_HEAP`
- `BlocksIndex` (`_HEAP_LIST_LOOKUP_`)
  - Back-End 中，重要的結構之一，用來管理 Back-End 中的 chunk，後面會詳細敘述
- `FreeList` (`_HEAP_ENTRY`)
  - 串接 Back-End 中的所有 free chunk，類似 unsorted bin
- Sorted list

# Windows memory allocator



- **\_HEAP**
- **FrontEndHeap**
- 指向管理 FrontEnd 的 Heap 的結構
- **FrontEndHeapUsageData**
- 指向一個對應各大小 chunk 的陣列
- 紀錄各種大小 chunk 使用次數，到達某個程度時會去 enable 該對應大小 chunk 的 Front-End allocator

# Windows memory allocator

- Nt Heap
  - 後端管理器 (Back-End)
  - 前端管理器 (Front-End)

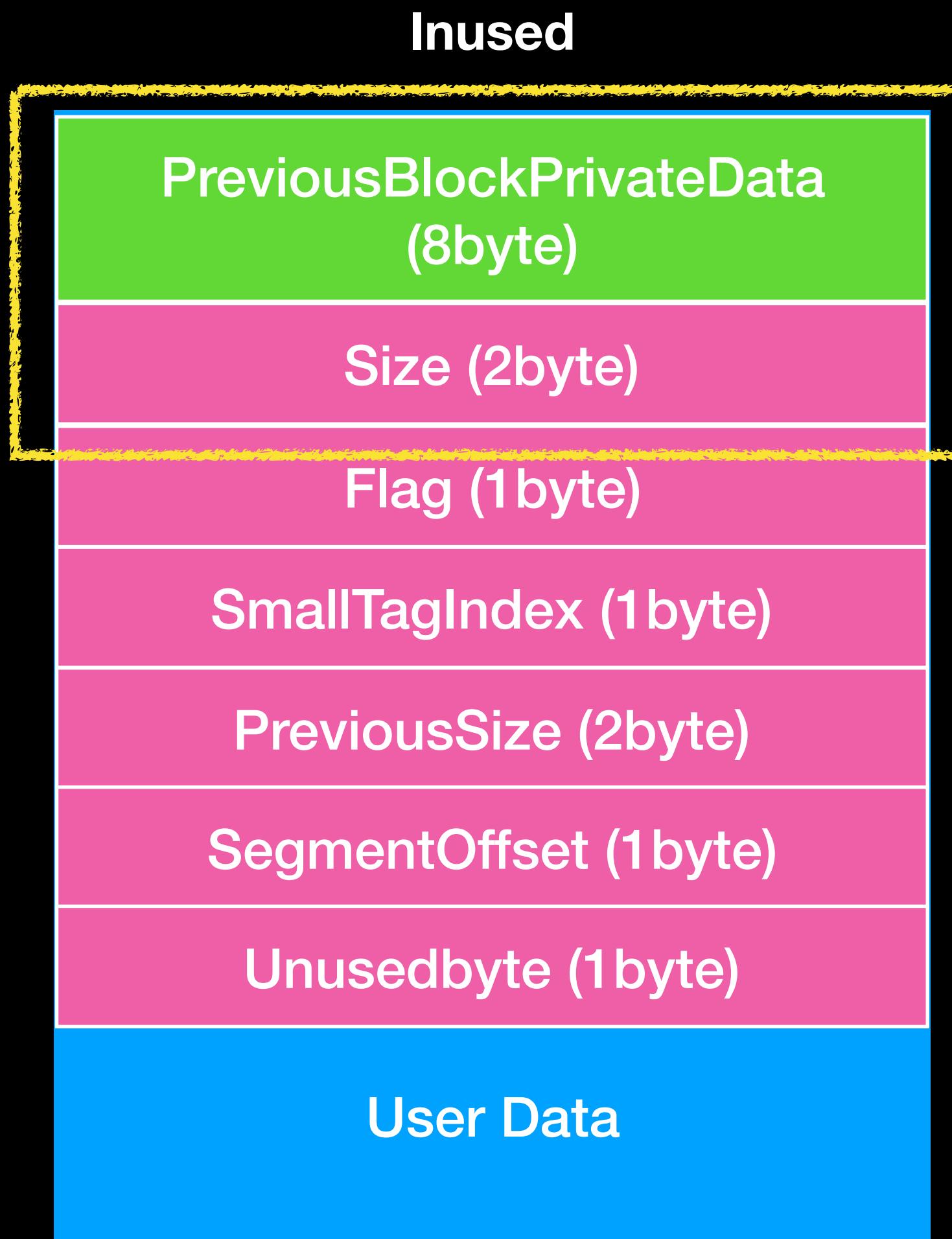
# Nt heap

- Nt Heap
  - 後端管理器 (Back-End)
    - Data structure
    - 分配機制

# Windows memory allocator

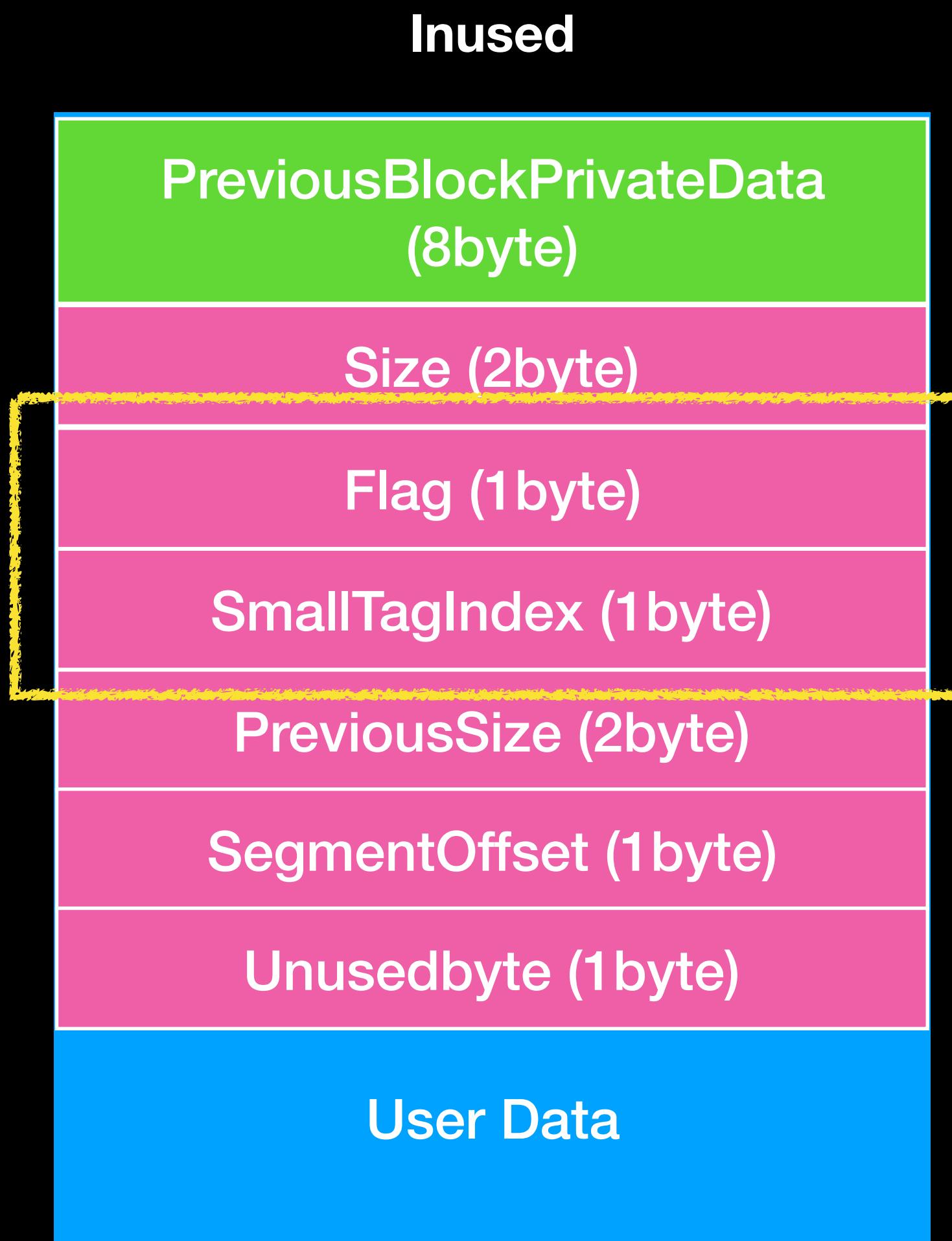
- \_HEAP\_ENTRY (chunk)
  - 分為三種狀況
    - Allocated chunk
    - Freed chunk
    - VirtualAlloc chunk

# Nt heap



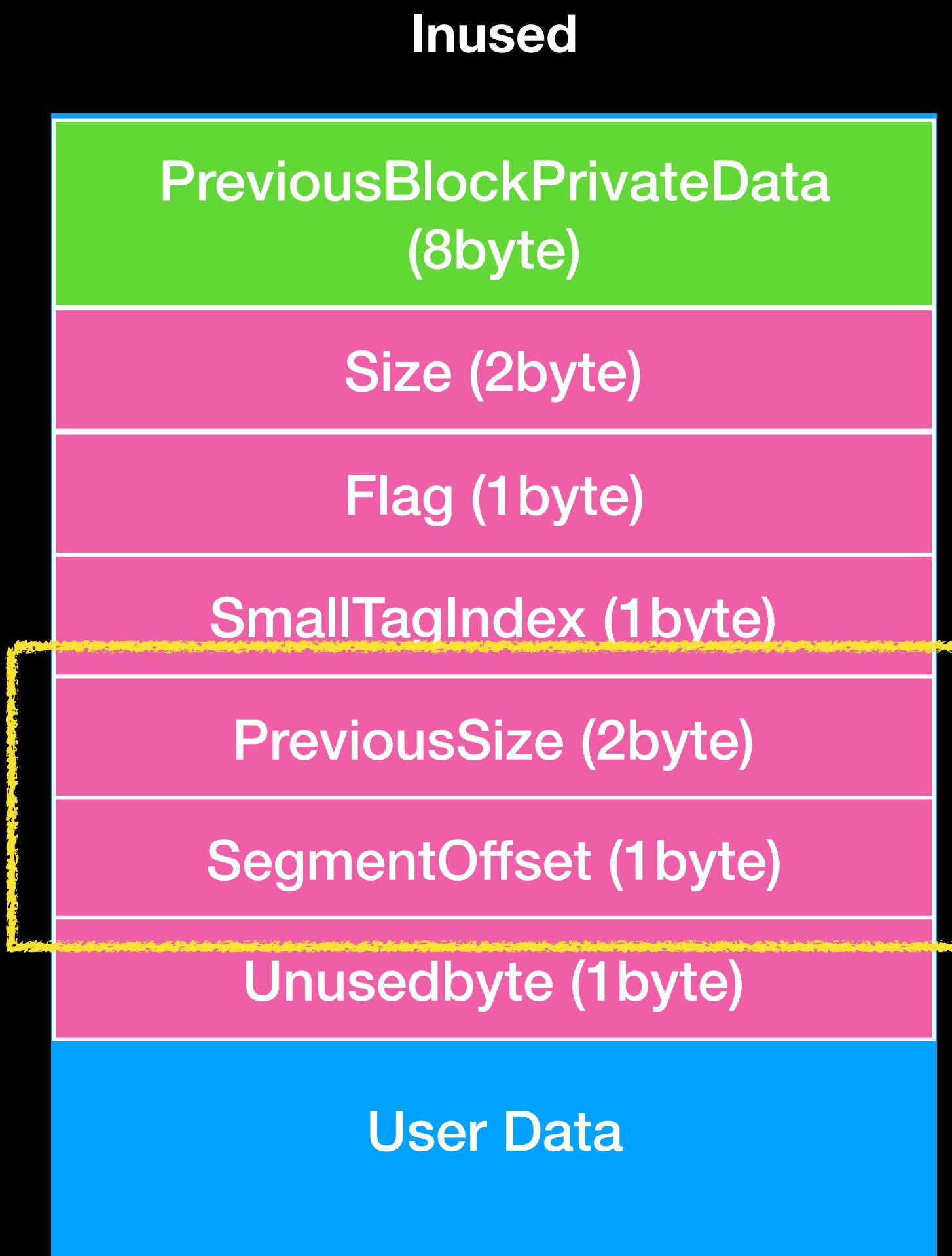
- **\_HEAP\_ENTRY (chunk)**
  - **PreviousBlockPrivateData**
    - 基本上可為前一塊 chunk 的 data，因為 chunk 必須對齊 0x10
  - **Size**
    - Chunk 的 size 但存入方式並不是直接存 size 而是存 (size >>4) 後的值

# Nt heap



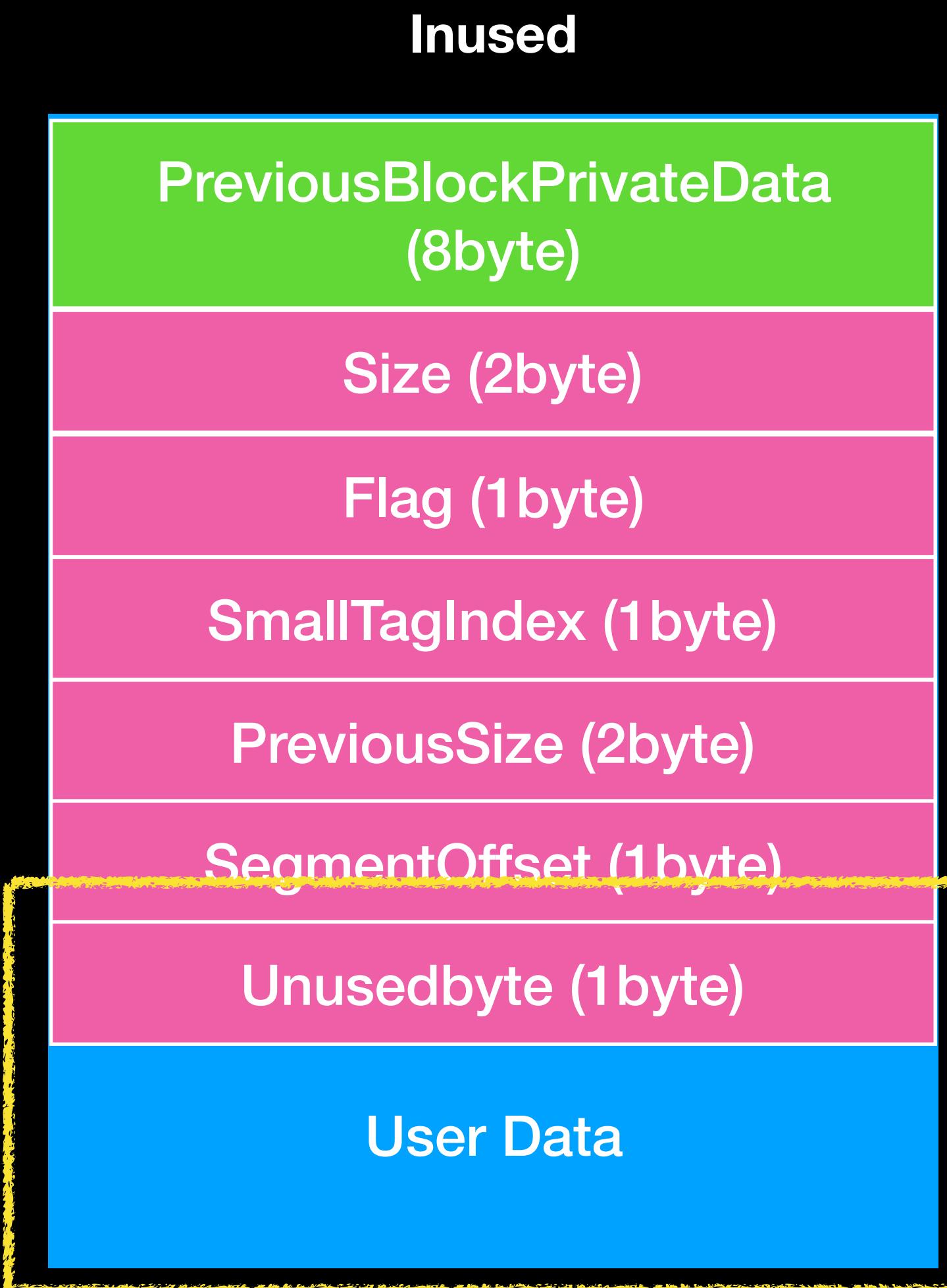
- **\_HEAP\_ENTRY (chunk)**
  - Flag
    - 表示該 chunk 是否 inused
  - SmallTagIndex
    - 前面三個 byte 的 Checksum(三個 byte 做 xor)，之後拿 header 出來時會驗證

# Nt heap



- \_HEAP\_ENTRY (chunk)
  - PreviousSize
    - 相鄰的前一塊 chunk 的 Size ，這邊一樣是 ( $\text{Size} \gg 4$ ) 過後的數值
  - SegmentOffset
    - 某些情況下用來找 segment 的

# Nt heap



- `_HEAP_ENTRY` (chunk)
  - Unusedbyte
    - 紀錄 user malloc 後所剩 chunk 的空間
    - 可以用來判斷 chunk 的狀態判斷 chunk 是 FrontEnd or BackEnd
  - User Data
    - User 所使用的區塊

# Nt heap



- `_HEAP_ENTRY` (chunk)
  - Flink
    - 指向 linked list 中下一塊 chunk
  - Blink
    - 指向 linked list 中上一塊 chunk
  - Unusedbyte
  - 恒為 0

# Nt heap



- `_HEAP_VIRTUAL_ALLOC_ENTRY` (mmap chunk)
  - **Flink**
    - 指向下一塊 VirtualAlloc 出來的 chunk
  - **Blink**
    - 指向上一塊 VirtualAlloc 出來的 chunk

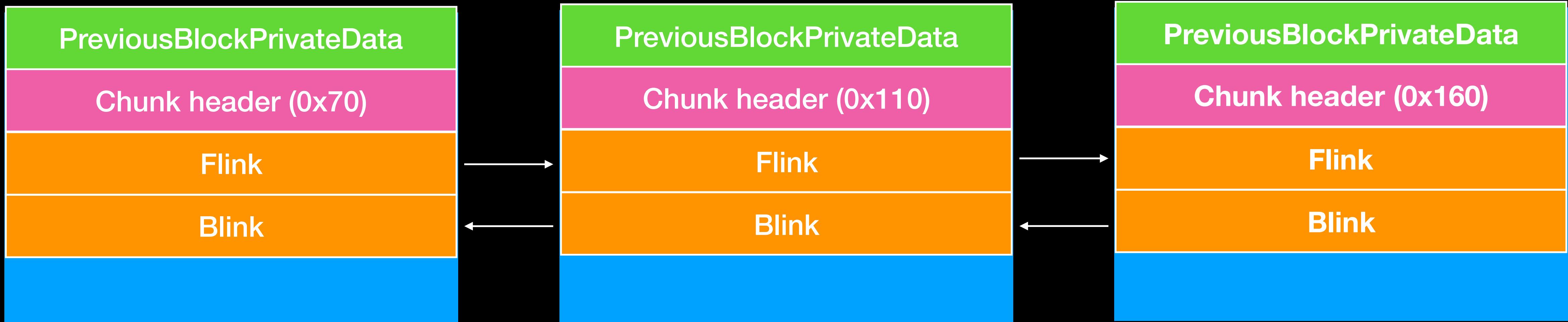
# Nt heap



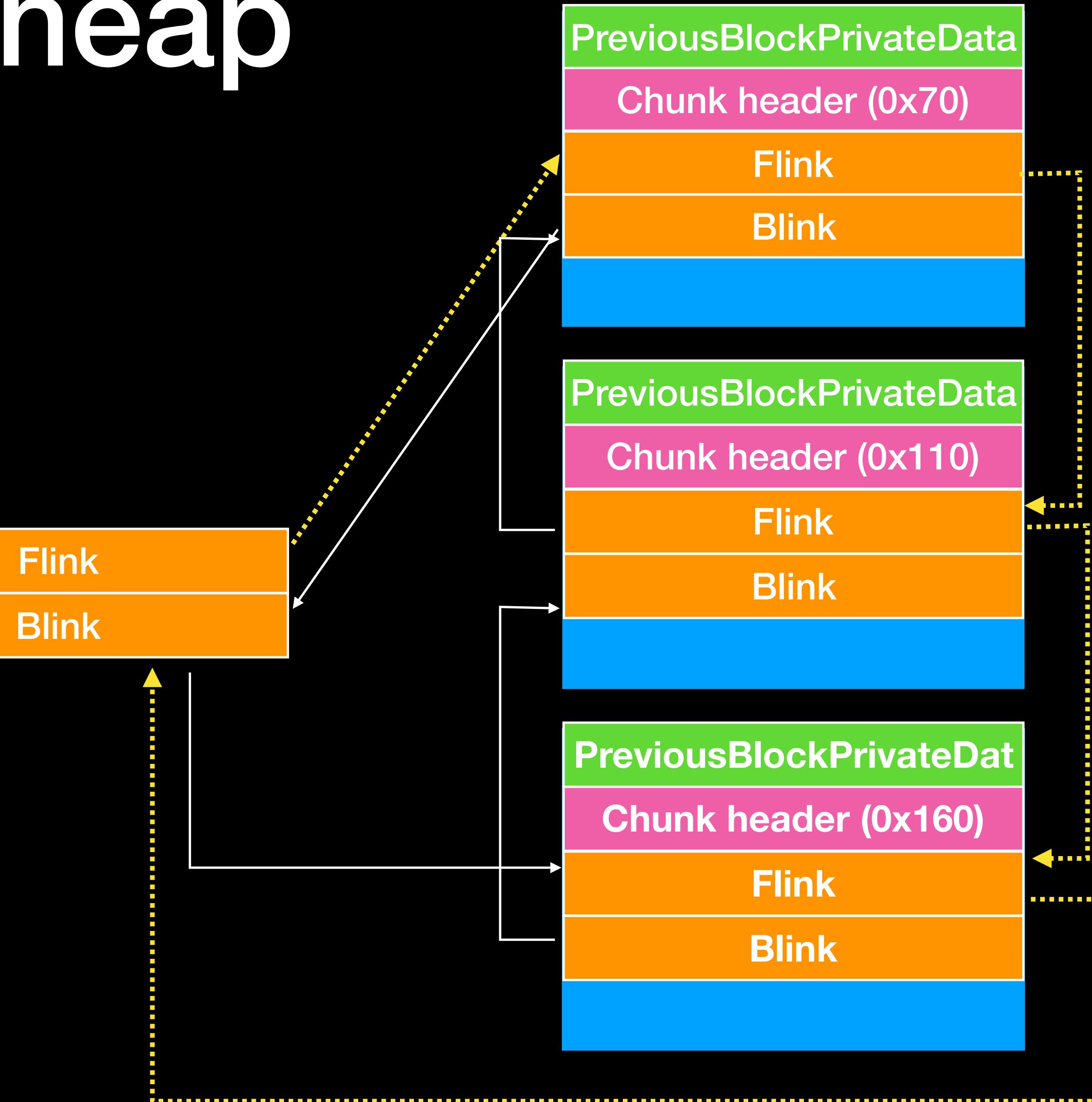
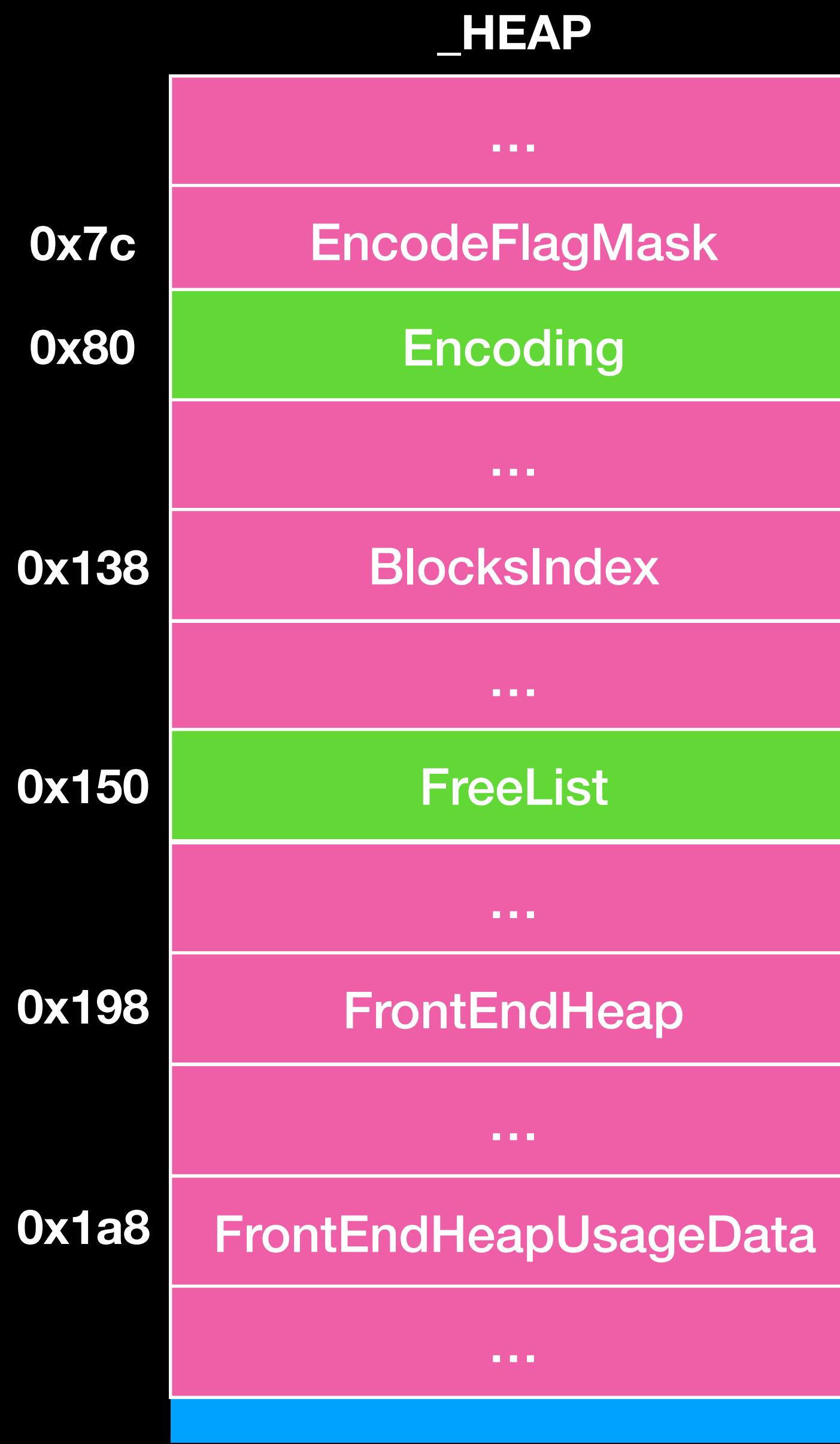
- `_HEAP_VIRTUAL_ALLOC_ENTRY` (`mmap` chunk)
- Size
  - 這邊的 size 指的是 unused size，且沒有 shift 過
- Unusedbyte
  - 恒為 4

# Nt heap

- FreeLists (\_HEAP\_ENTRY)
  - 在 Free 完一塊 chunk 後，會將該 chunk 放到 FreeLists 中，並會依照大小決定插在 FreeLists 中的位置



# Nt heap



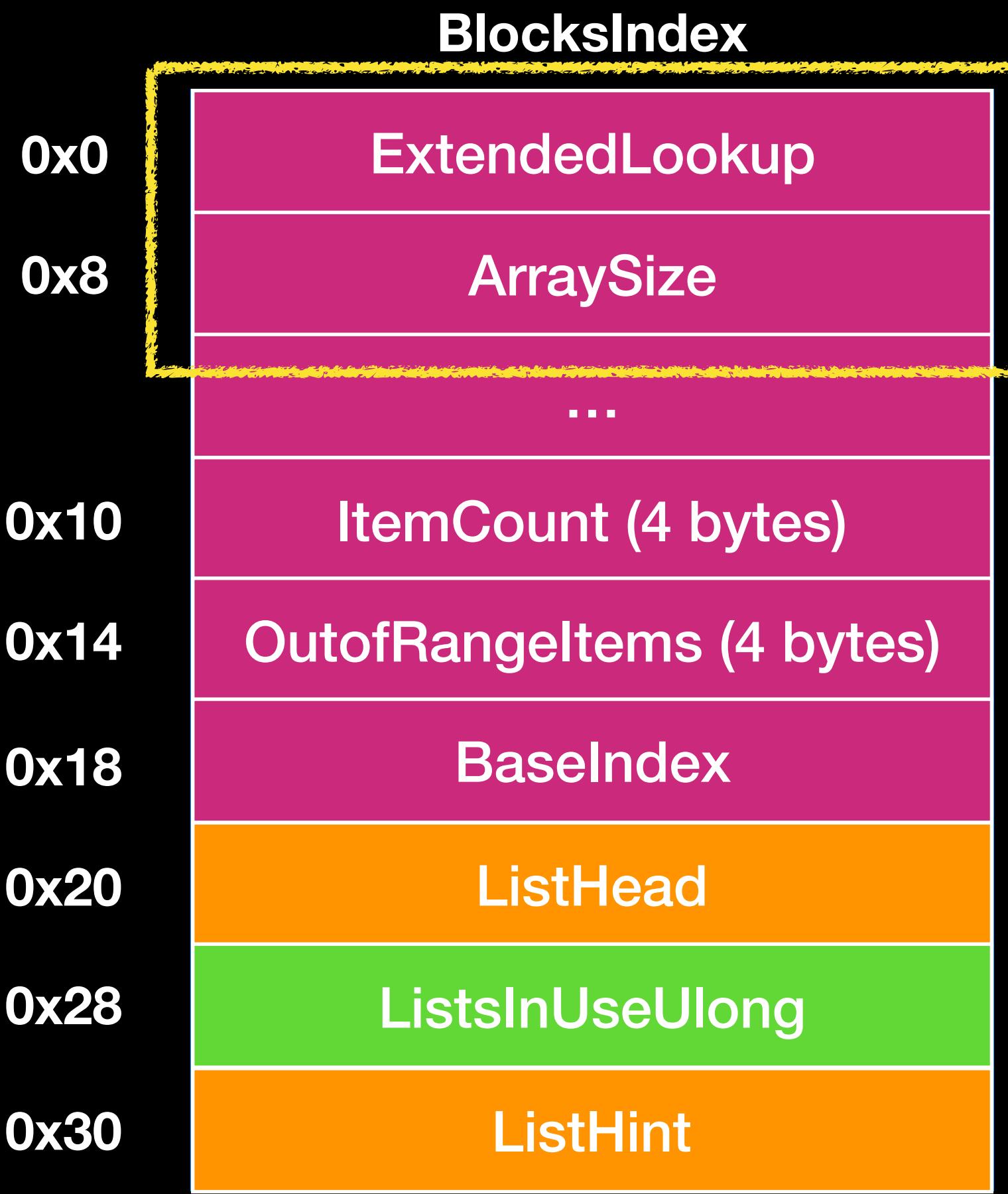
# Nt heap

- Remark
  - 關於 header encoding
    - 所有的 chunk 都會經過 xor 過，在存 chunk header 時，會將整個 header $\wedge$ (\_HEAP->Encoding) 在存入
    - decode 時會驗證 check sum 確保沒被改掉
    - 驗證方式為 encode\_header $\wedge$ (\_HEAP->Encoding) 後的值，前 3 byte xor 後和第 4 byte 比對

# Nt heap

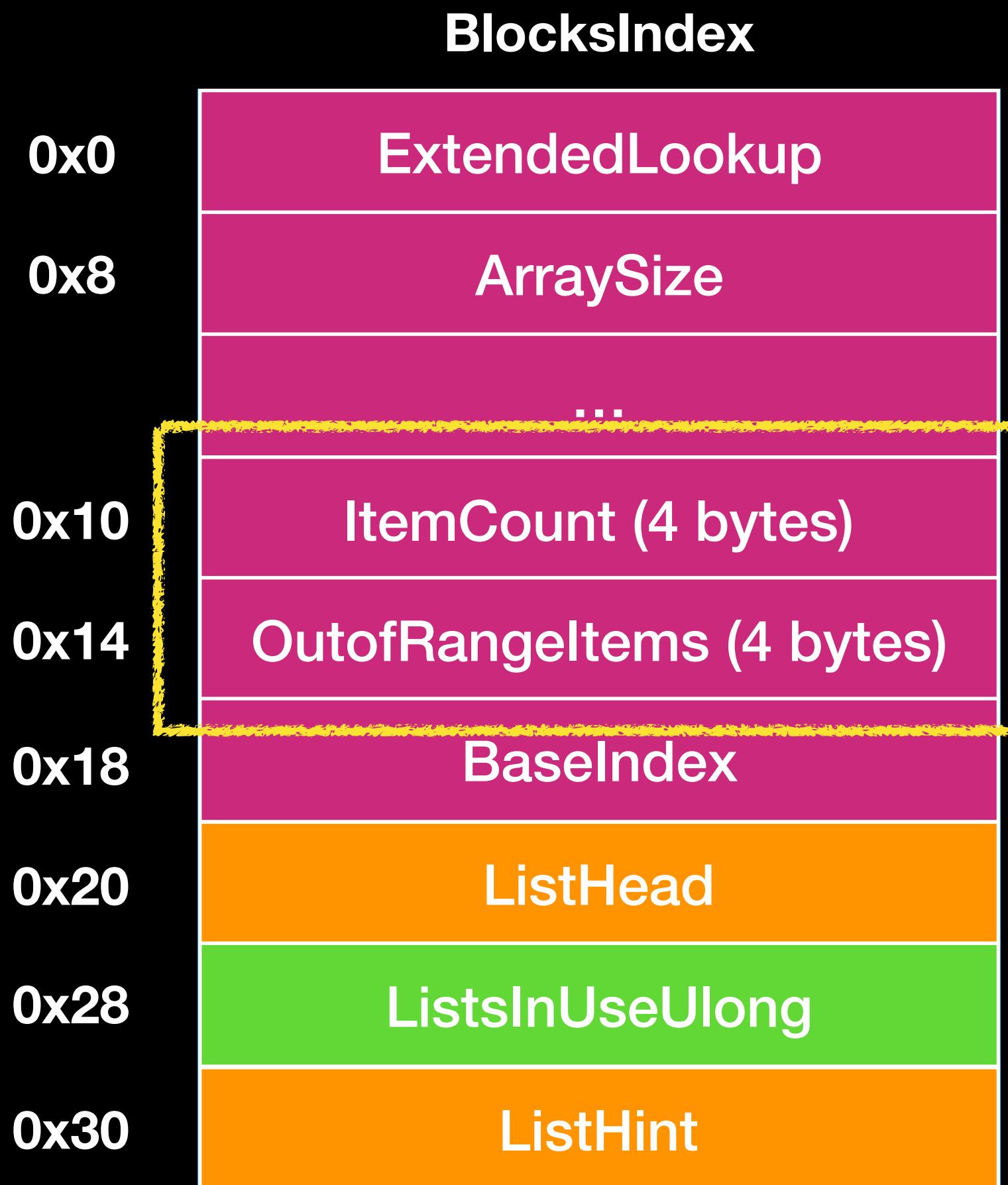
- `BlocksIndex (_HEAP_LIST_LOOKUP)`
  - 主要用來管理各種不同大小的 freed chunk 的結構，方便在配置記憶體時，能快速的找到適合的 chunk

# Nt heap



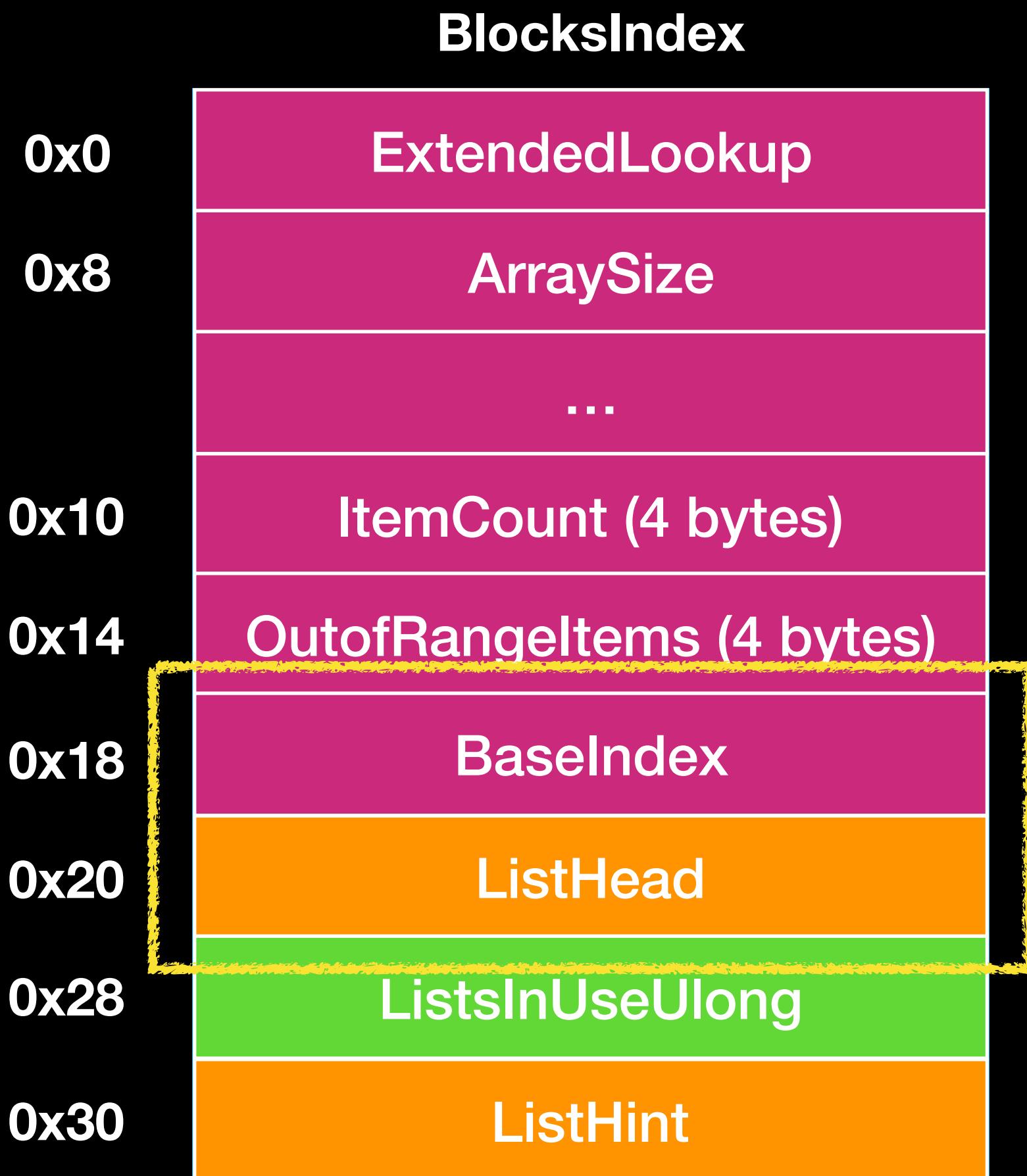
- BlocksIndex (\_HEAP\_LIST\_LOOKUP)
  - ExtendedLookup (\_HEAP\_LIST\_LOOKUP)
    - 指向下一个 ExtendedLookup
    - 通常下一个会管理更大块 chunk 所用的结构
  - ArraySize
    - 该结构会管理的，最大 chunk 的大小，通常为 0x80 (实际上为 0x800)

# Nt heap



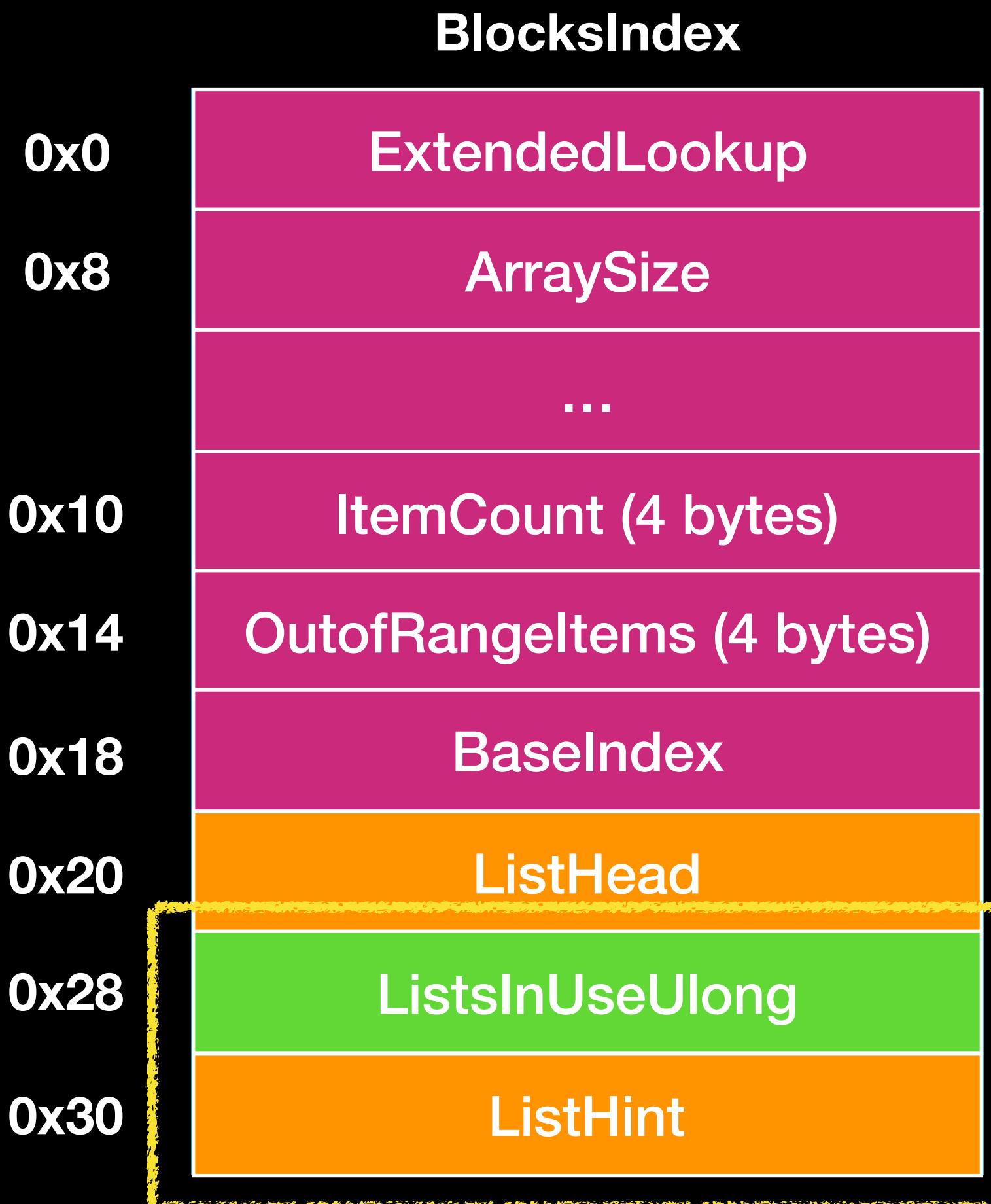
- `BlocksIndex` (`_HEAP_LIST_LOOKUP`)
  - `ItemCount`
    - 目前該結構所管理的 chunk 數
  - `OutofRangeItems`
    - 超出該結構所管理大小的 chunk 的數量

# Nt heap

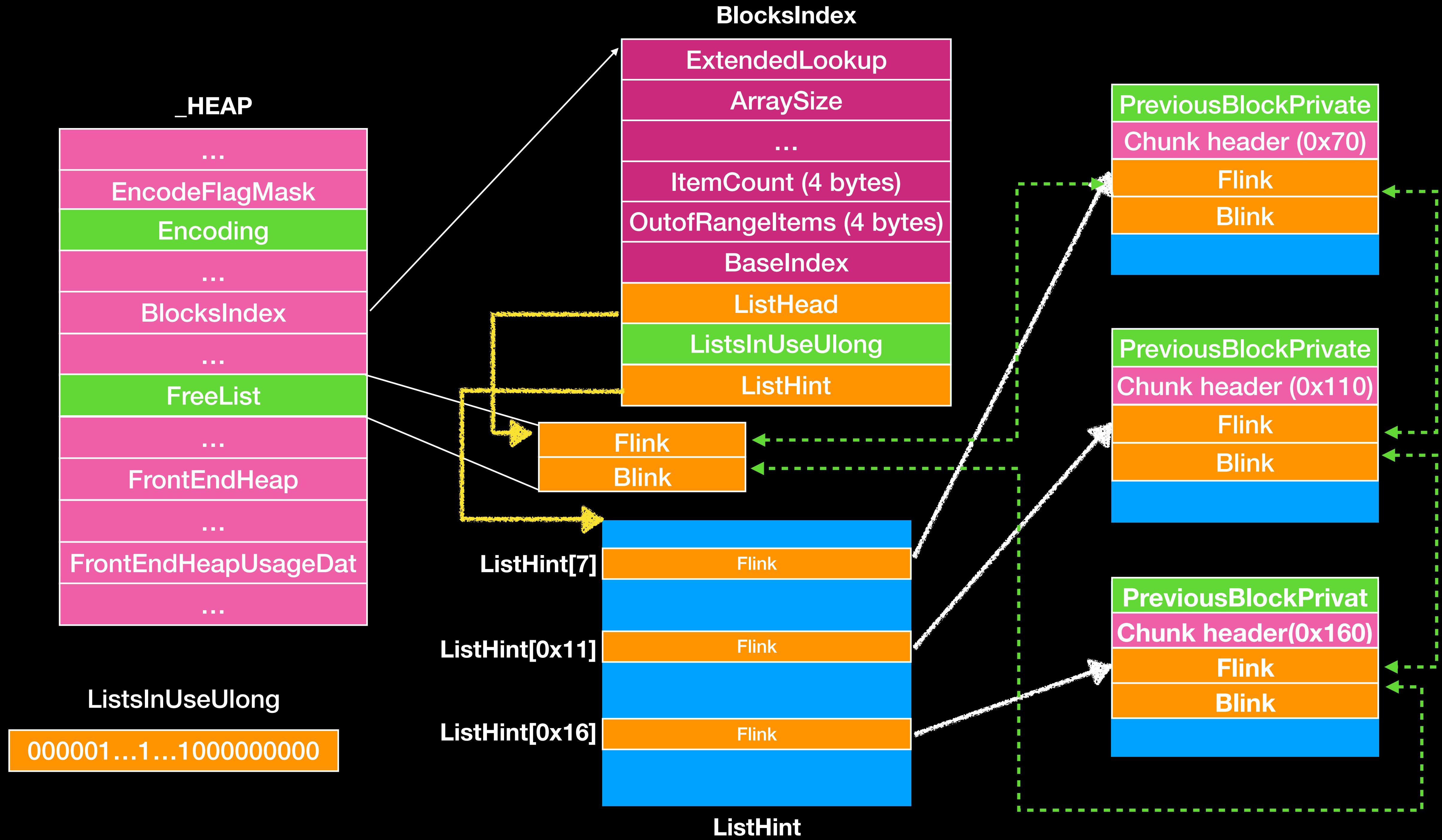


- **BlocksIndex (\_HEAP\_LIST\_LOOKUP)**
  - **BaselIndex**
    - 該結構所管理的 chunk 的起始 index，用來從 ListHint 找適合大小的 free chunk 用的，下一個 BlocksIndex 將從這個 index 的最大值作為 BaselIndex
- **ListHead (\_HEAP\_ENTRY)**
  - FreeList 的 Head

# Nt heap



- **BlocksIndex (\_HEAP\_LIST\_LOOKUP)**
  - **ListsInUseUlong**
    - 用在判斷 ListHint 中是否有適合大小的 chunk，是一個 bitmap
  - **ListHint**
    - 重要結構，用來指向相對應大小的 chunk array
    - 其目的就在於更快速找到適合大小的 chunk 大小為 0x10 為一個間隔



# Nt heap

- Nt Heap
  - 後端管理器 (Back-End)
    - Data structure
    - 分配機制

# Nt heap

- Allocate (RtlpAllocateHeap)
  - 基本上分為三種
    - Size <= 0x4000
    - 0x4000 < size <= 0xff000
    - Size > 0xff000

# Nt heap

- Allocate (RtlpAllocateHeap)
  - Size <= 0x4000
  - 基本上主要分配都會在 RtlpAllocateHeap
  - 接著會去看該 size 對應到的 FrontEndHeapStatusBitmap 使否有啟用 LFH
  - 沒有的話會對對應到的 FrontEndHeapUsageData 加上 0x21
  - 並且檢查值是否超過 0xff00 或者 & 0x1f 後超過 0x10  
通過條件就會啟用 LFH

# Nt heap

- Allocate (RtlpAllocateHeap)
  - Size <= 0x4000
  - 接下來會先看對應到的 ListHint 是否有值，會以 ListHint 中的 chunk 為優先
  - 如果有適合剛好的 chunk 在 ListHint 上則移除 ListHint，並且會看該 chunk 的 Flink 大小是否剛好也是同樣 size
  - 是的話就將 ListHint 填上 Flink
  - 不是則清空
  - 最後則 unlink 該 chunk 把這塊 chunk 從 linked list 中移除返還給使用者，並將 header xor 回去

# Nt heap

- Allocate (RtlpAllocateHeap)
  - Size <= 0x4000
  - 如果沒有剛好適合的
    - 從比較大的 ListHint 中找，有找到比較大的後如果該 chunk 將會從 ListHint 則移除之，並看下一塊 chunk size 有沒有剛好大小，有則補上，並且 unlink 將該 chunk 移除 freelist
    - 最後將該 chunk 做切割，剩下的大小重新加入 Freelist ，如果可以放進 ListHint 就會放進去
    - 然後回傳切好的 chunk 紿使用者，這邊也會對切割後跟回傳使用者的 chunk header xor

# Nt heap

- Allocate (RtlpAllocateHeap)
  - Size <= 0x4000
  - 如果 FreeList 中都沒有
    - 嘗試 ExtendHeap 加大 heap 空間
    - 再從 extend 出來的 chunk 拿
    - 接著後面一樣切割，放回 ListHint，xor header

# Nt heap

- Allocate (RtlpAllocateHeap)
  - $0x4000 < \text{size} \leq 0xff000$ 
    - 除了沒有對 LFH 相關的操作外，其餘都跟  $0x4000$  一樣

# Nt heap

- Allocate (RtlpAllocateHeap)
  - Size > 0xff000 (VirtualMemoryThreshold << 4)
  - 直接使用 ZwAllocateVirtualMemory
  - 類似 mmap 直接給一大塊，並且會插入 \_HEAP->VirtualAllocdBlocks 這個 linked list 中
  - 這個 linked list 是串接該 Heap VirtualAllocate 出來的區段用的

# Nt heap

- Free (RtlpFreeHeap)
  - 可分為
    - Size <= 0xff000
    - Size > 0xff000

# Nt heap

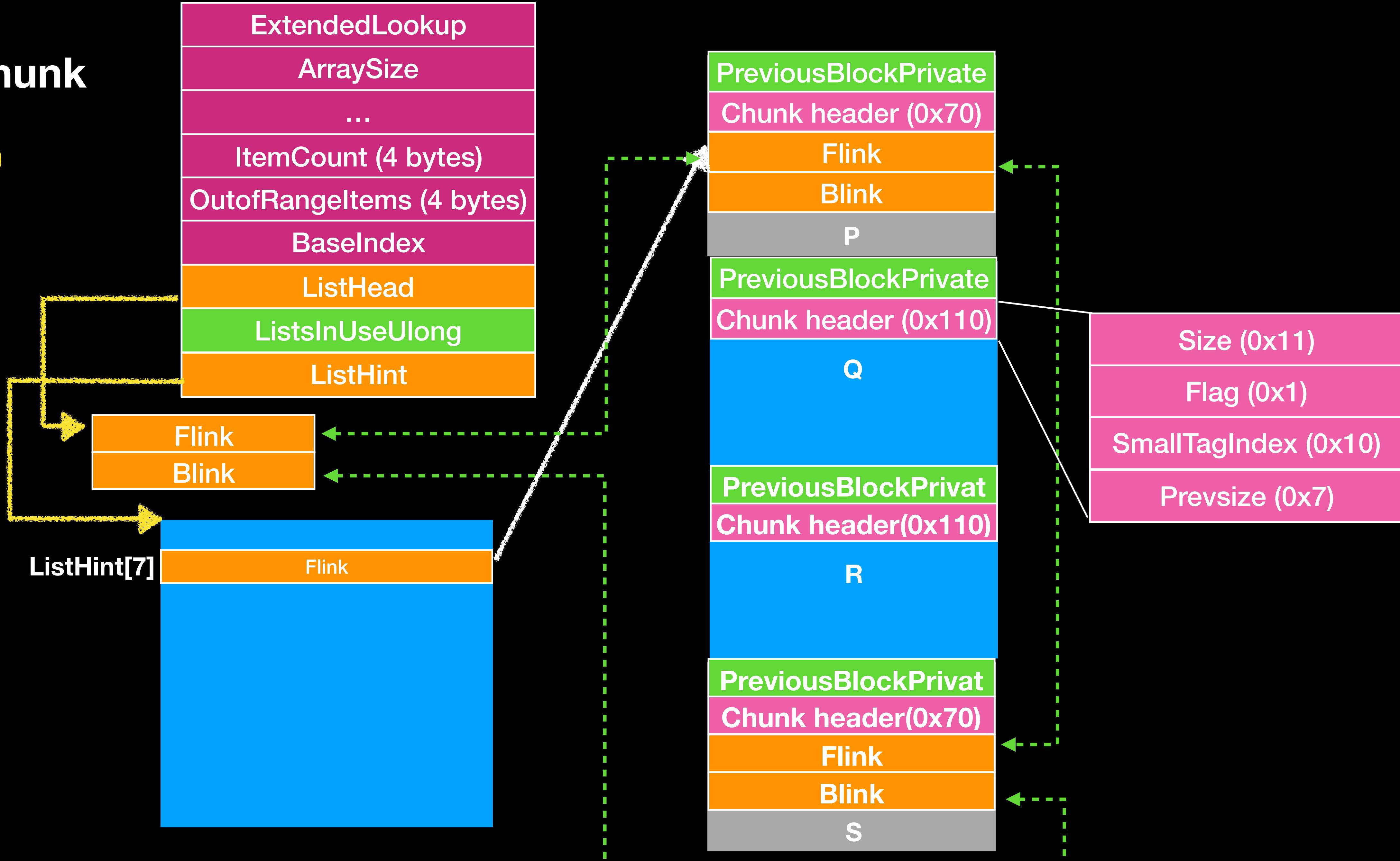
- Free (RtlpFreeHeap)
  - Size <= 0xff000
  - 會先檢查 alignment、利用 unused byte 判斷該 chunk 狀態
  - 如果是非 LFH 下，會對對應到的 FrontEndHeapUsageData 減 1
  - 接著會判斷前後的 chunk 是否為 freed ，是的話就合併
    - 此時會把可以合併的 chunk 做 unlink ，並從 ListHint 移除
    - 移除方式與前面相同，看看下一塊是不是同樣大小，是的話補上 ListHint

# Nt heap

- Free (RtlpFreeHeap)
  - Size <= 0xff000
  - 合併完後， update size & preysize ，然後會看看是不是最前跟最後，是就插入，不行插入就從 ListHint 中插入，並且 update ListHint ，插入時也會對 linked list 做檢查
  - 但這檢查不會 abort
    - 其原因主要因為不做 unlink 寫入

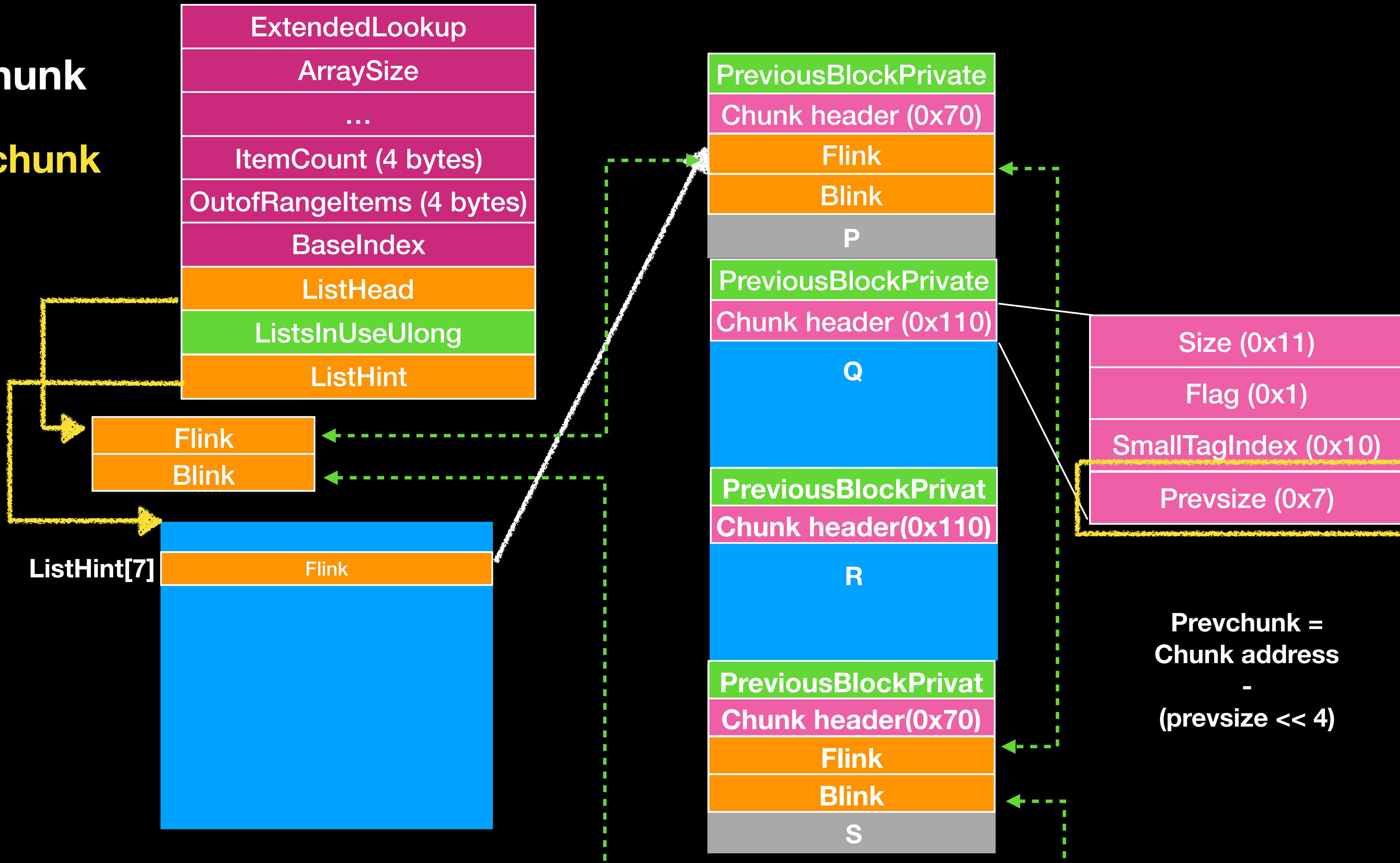
**merge chunk**

**Free(Q)**



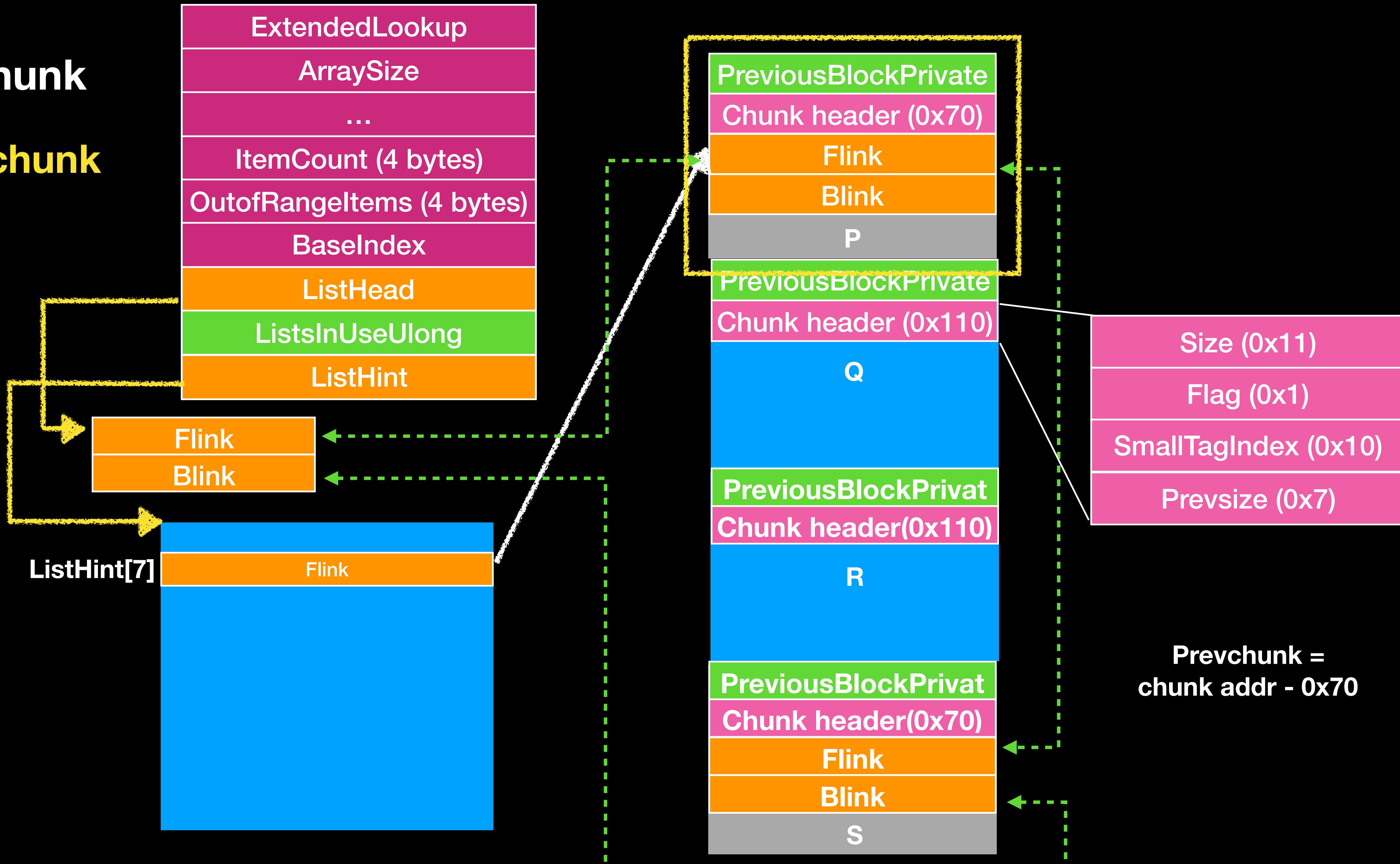
**merge chunk**

**Find Prevchunk**



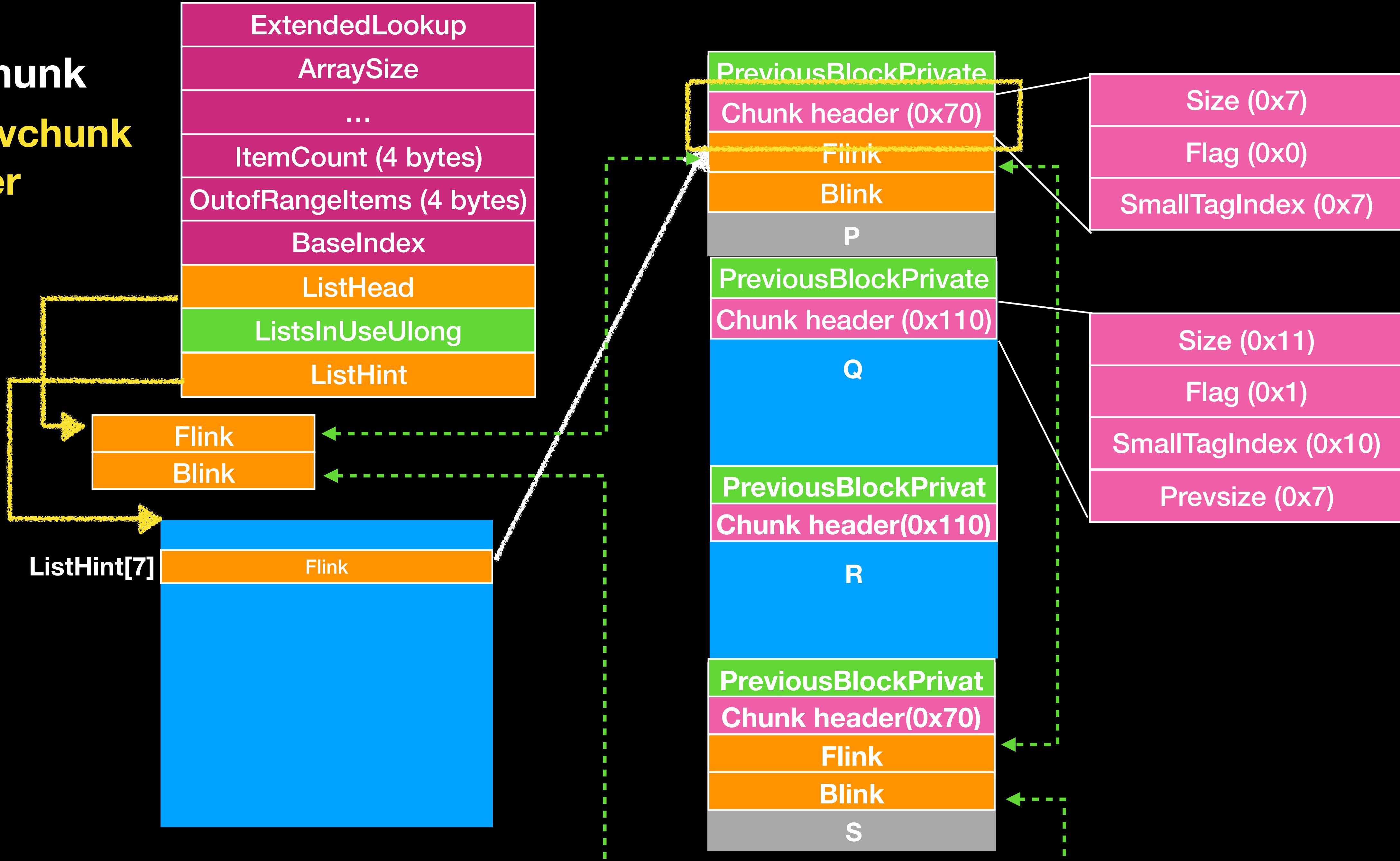
**merge chunk**

**Find Prevchunk**



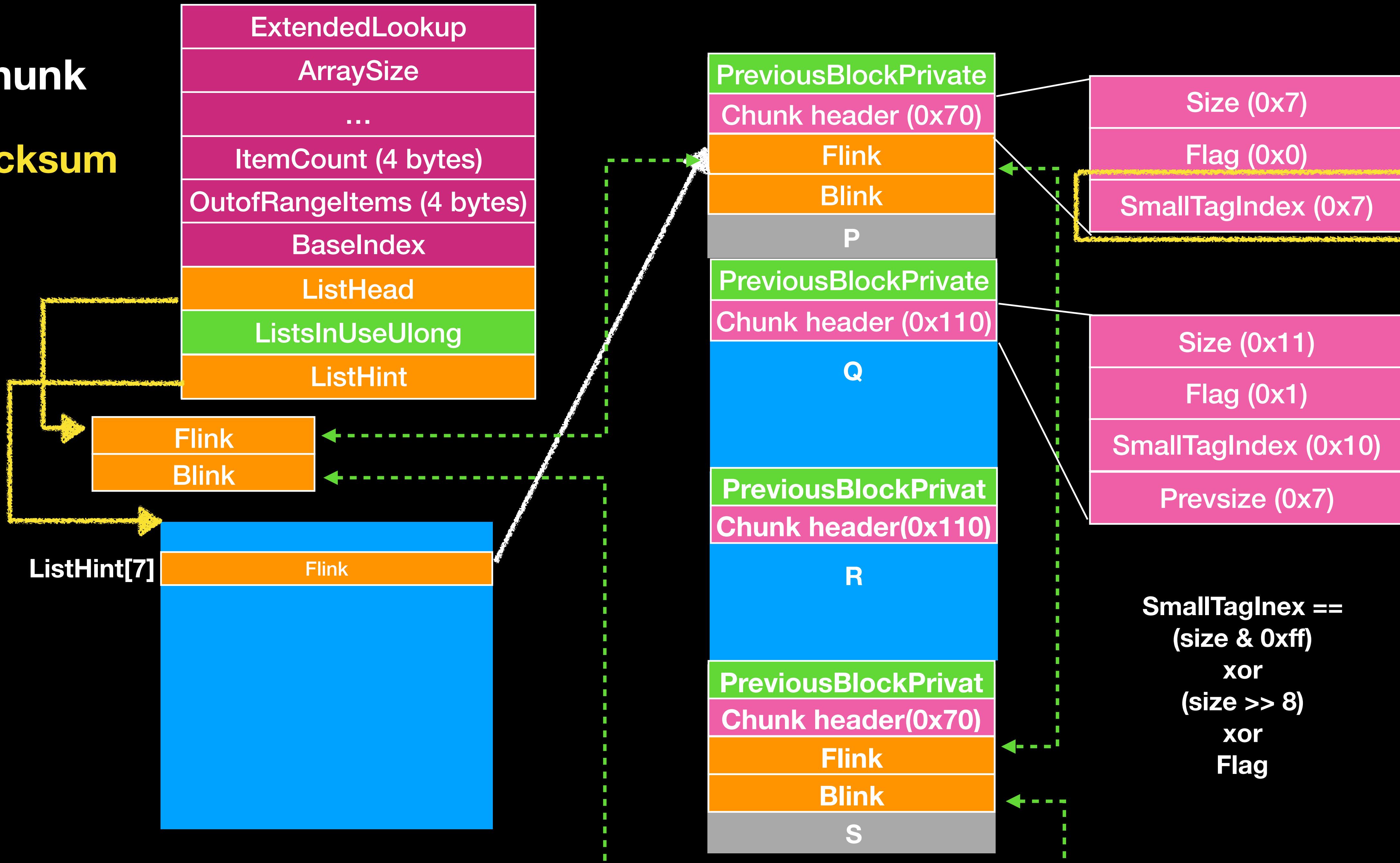
**merge chunk**

**Decode prevchunk header**



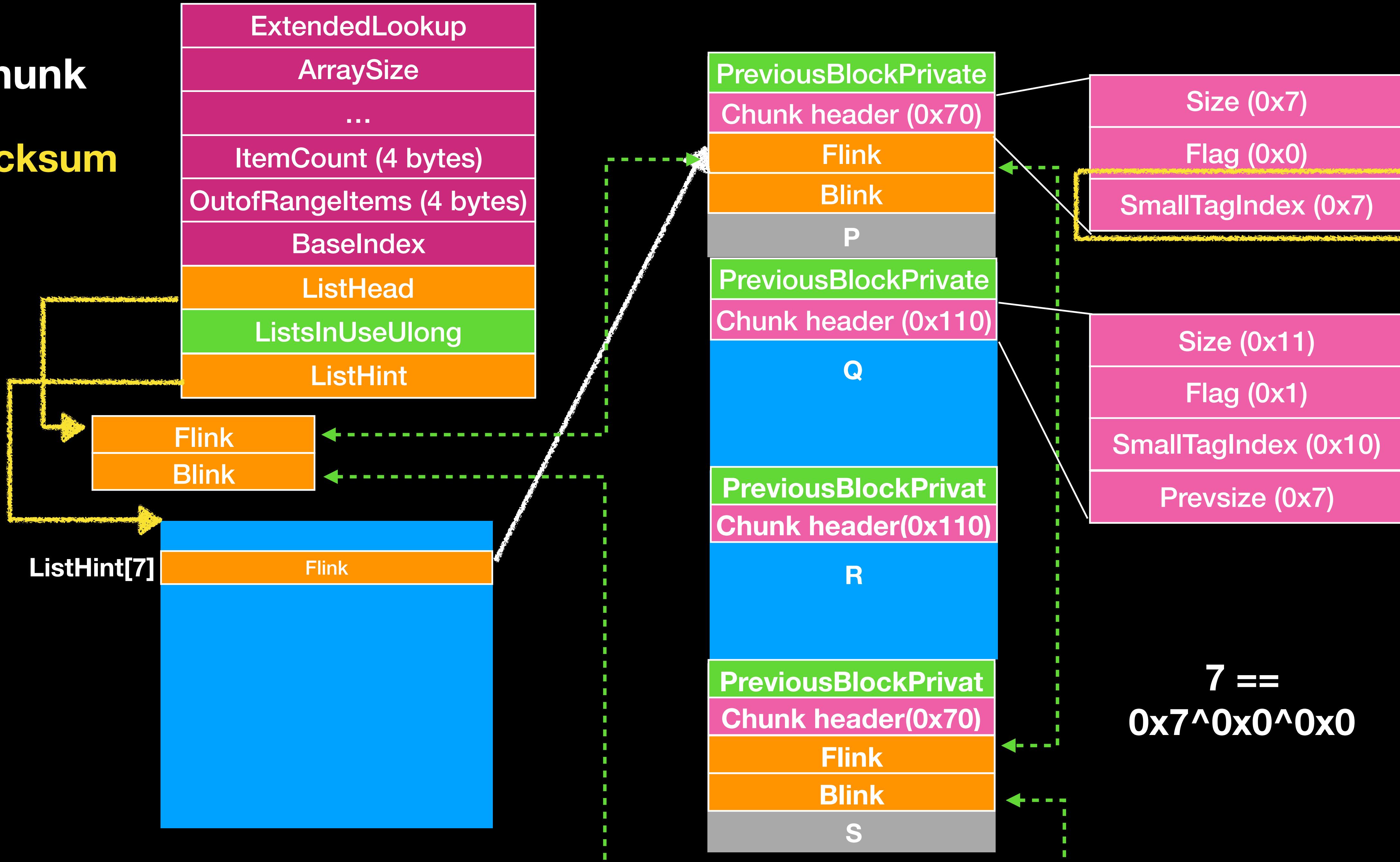
**merge chunk**

**Check checksum**

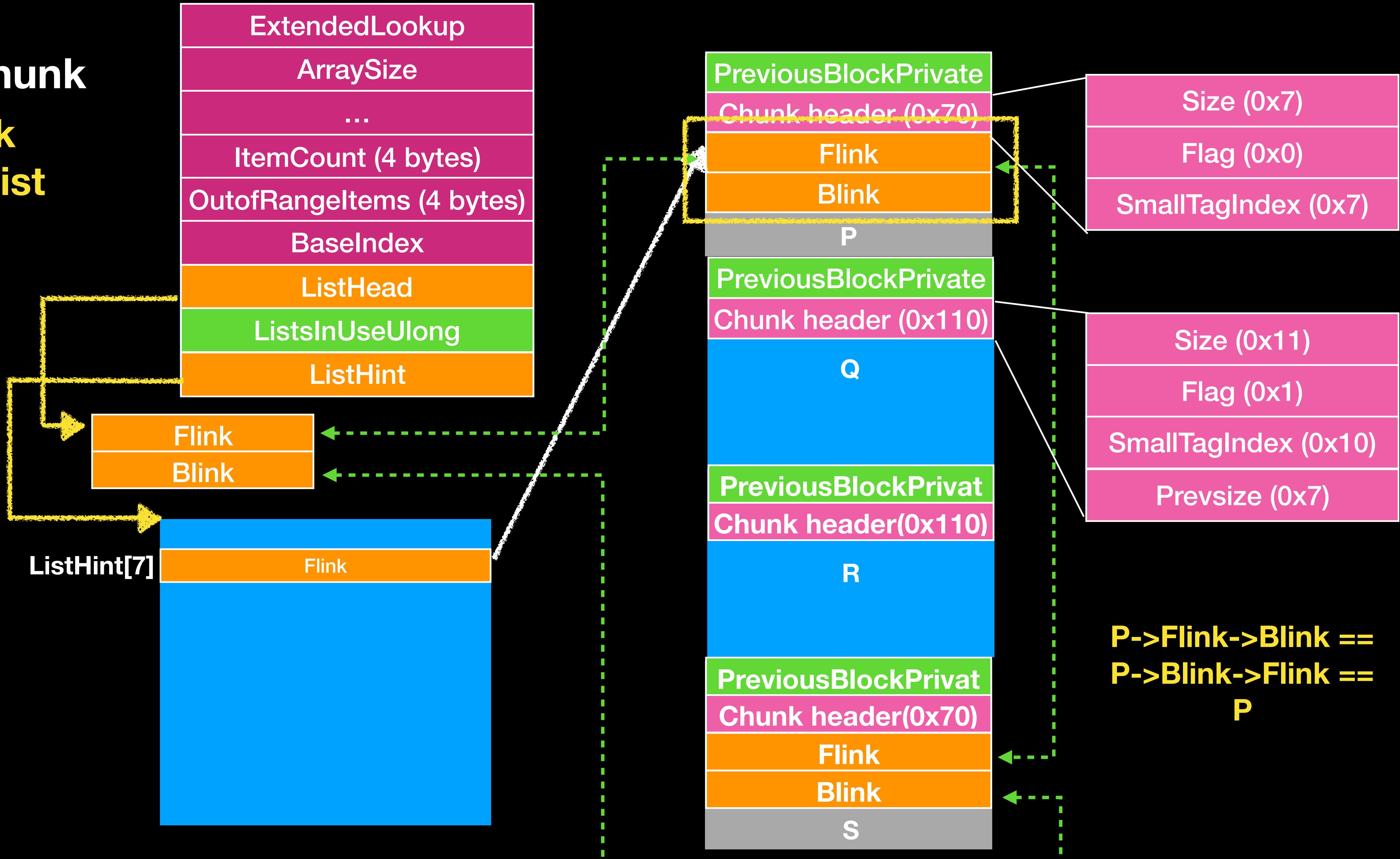


# merge chunk

# Check checksum

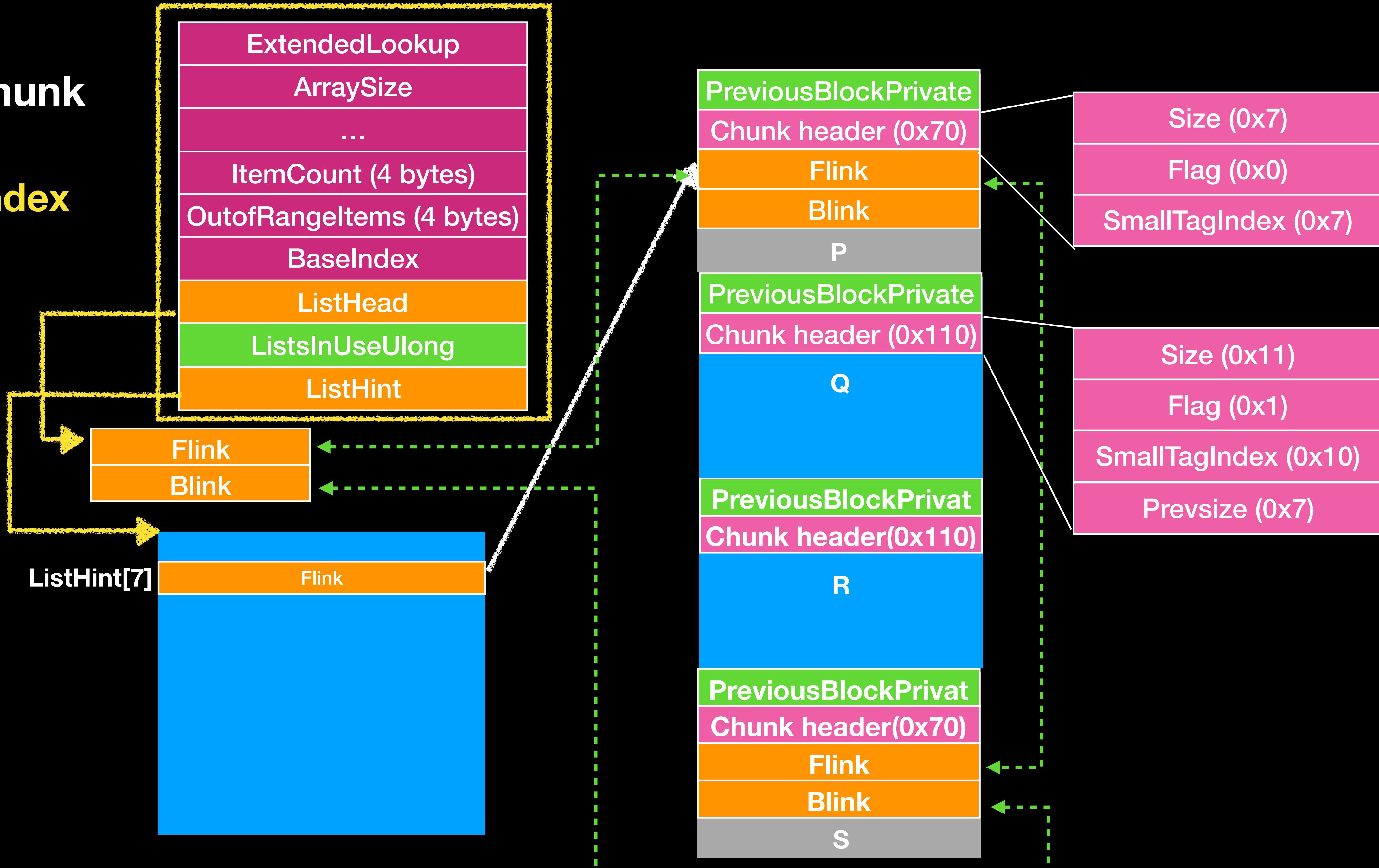


**merge chunk**  
**Check  
linked list**

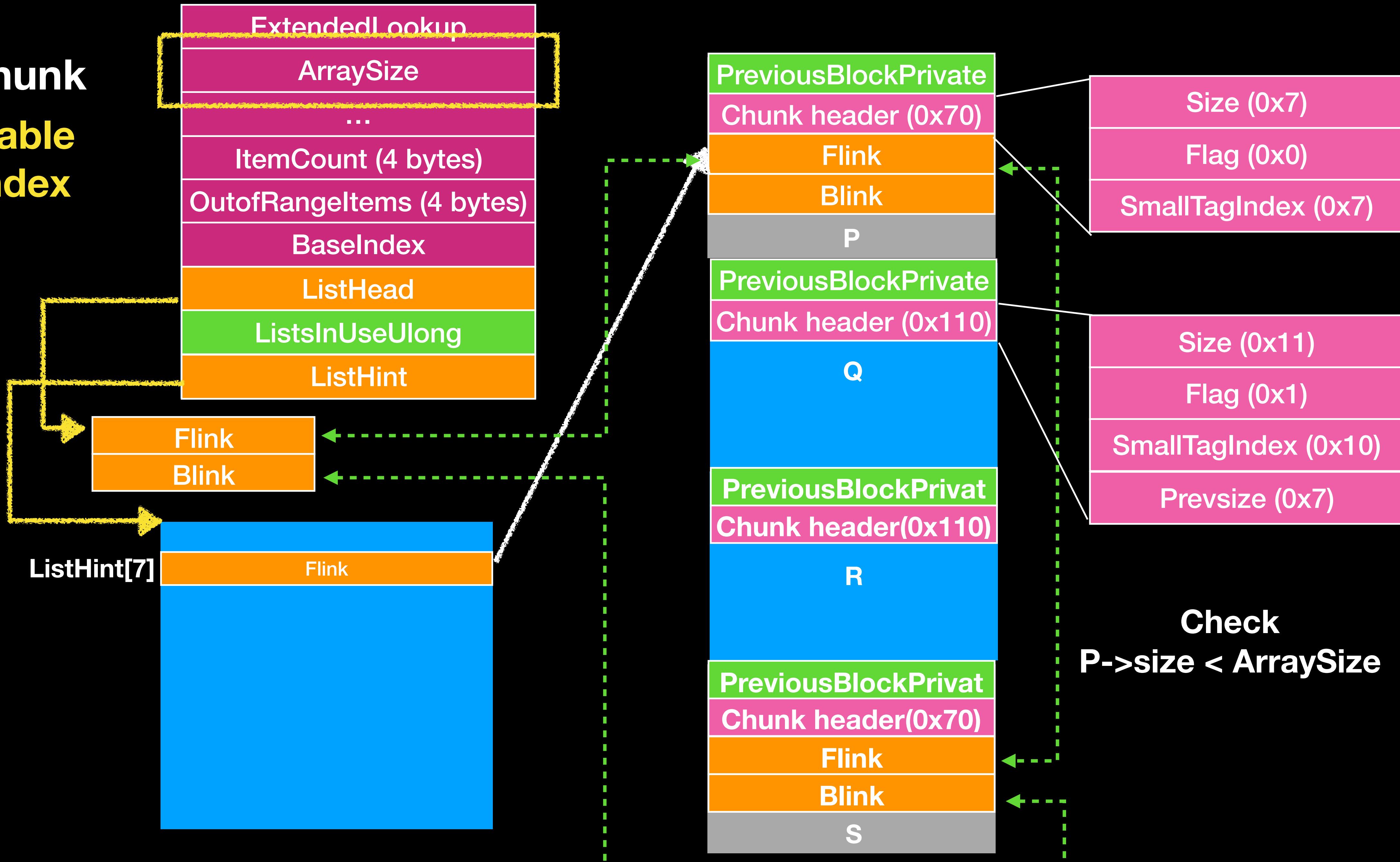


# merge chunk

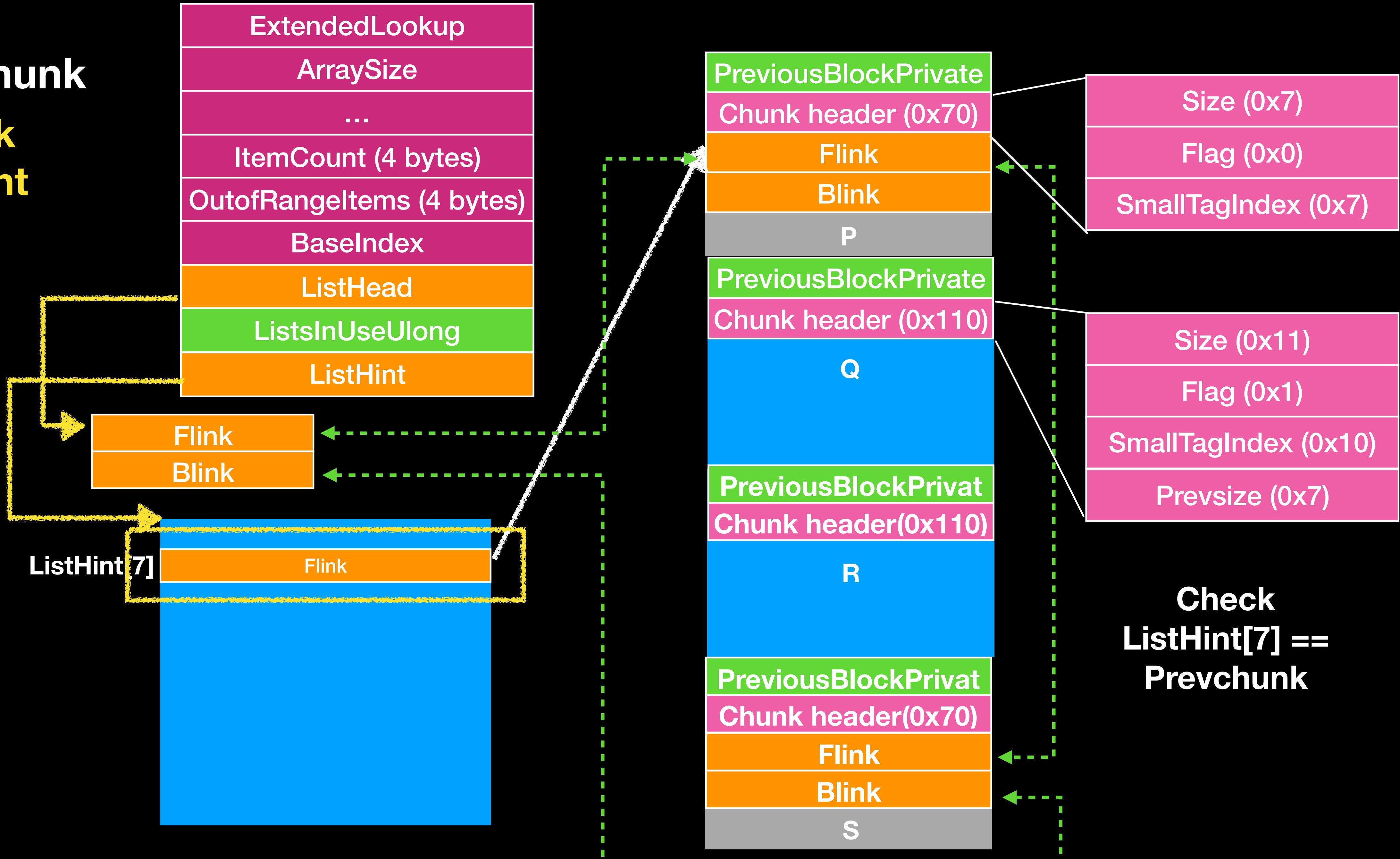
## Find BlocksIndex



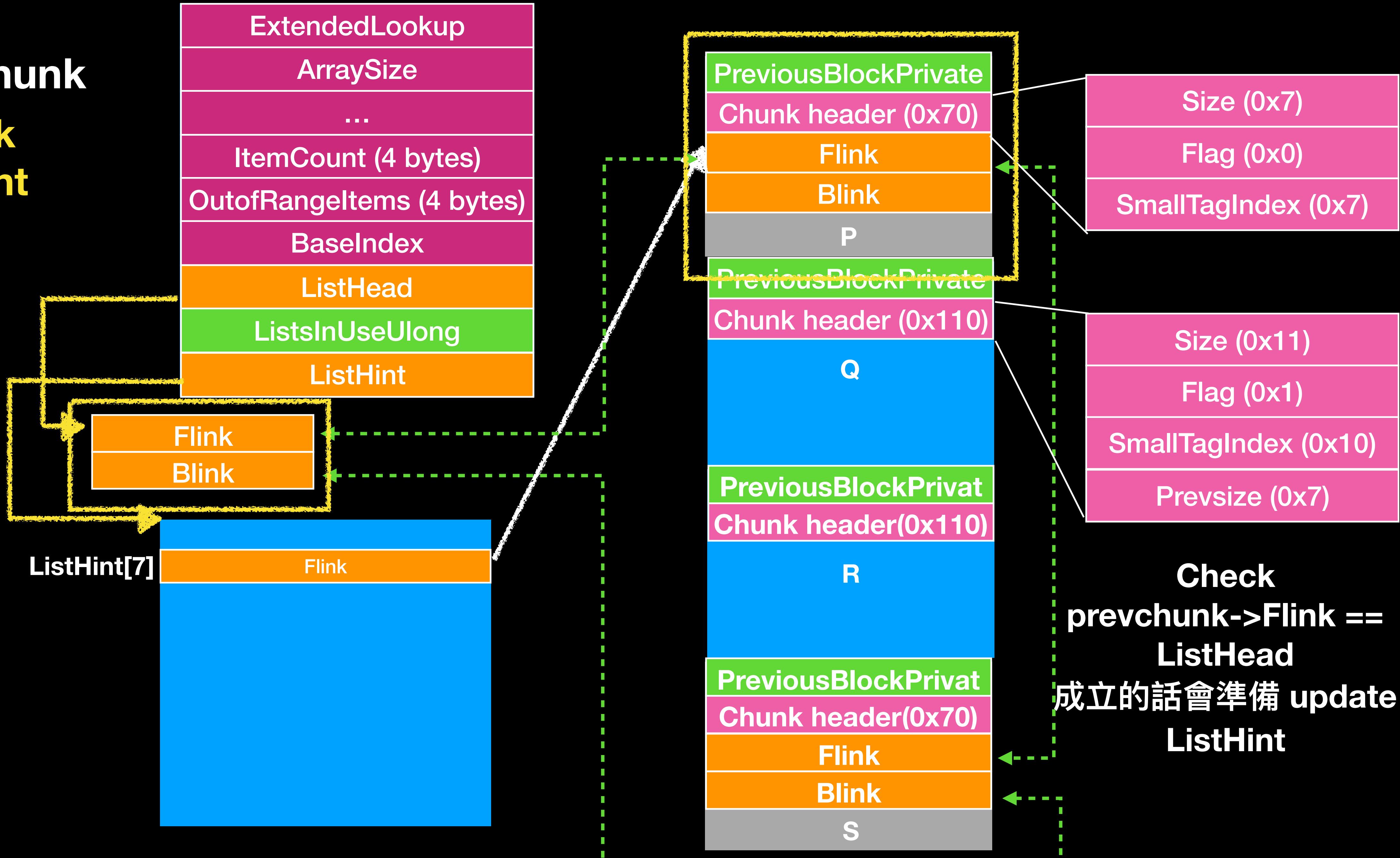
**merge chunk**  
**Find suitable  
BlocksIndex**



**merge chunk**  
**Check**  
**ListHint**

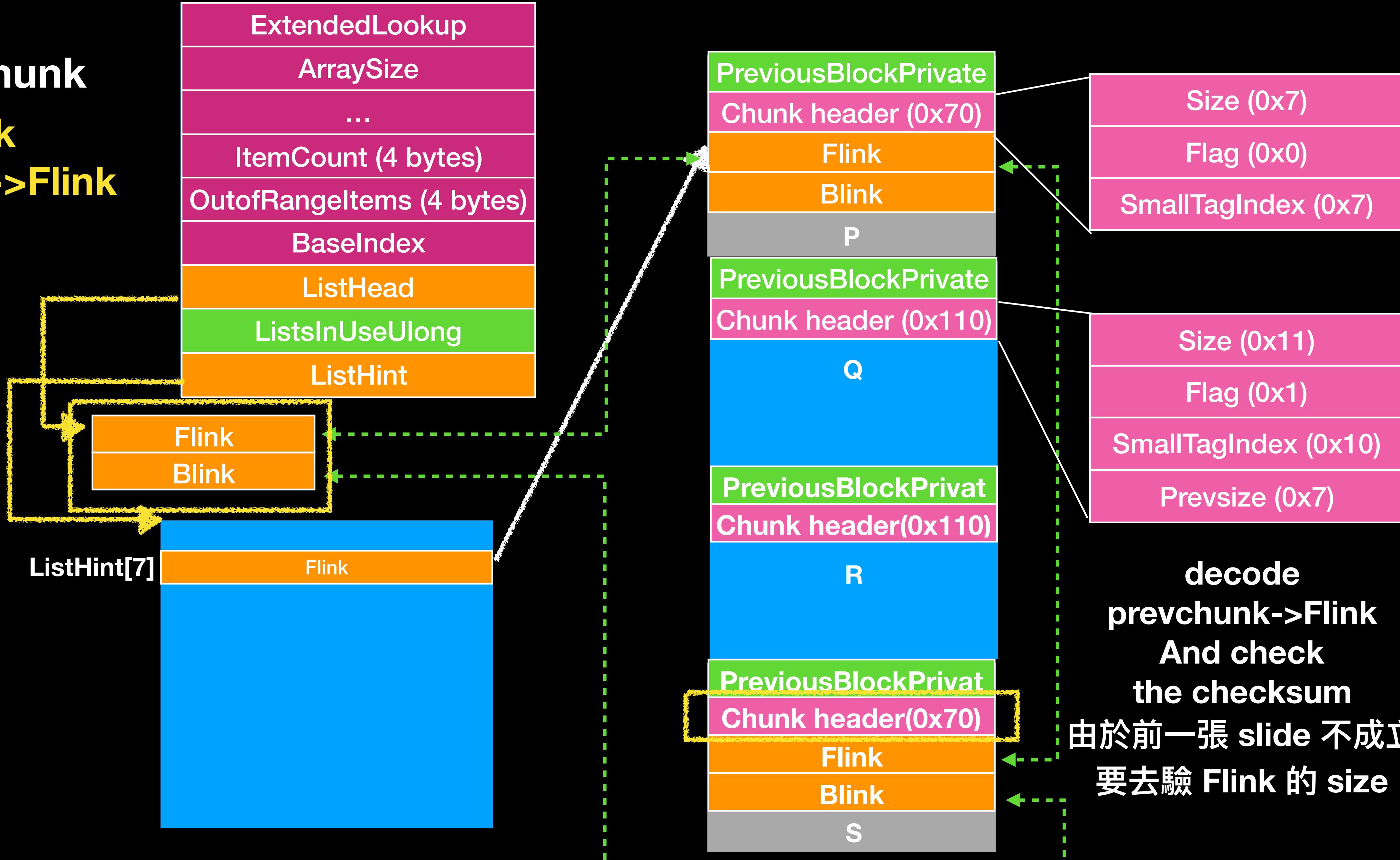


**merge chunk**  
**Check**  
**ListHint**

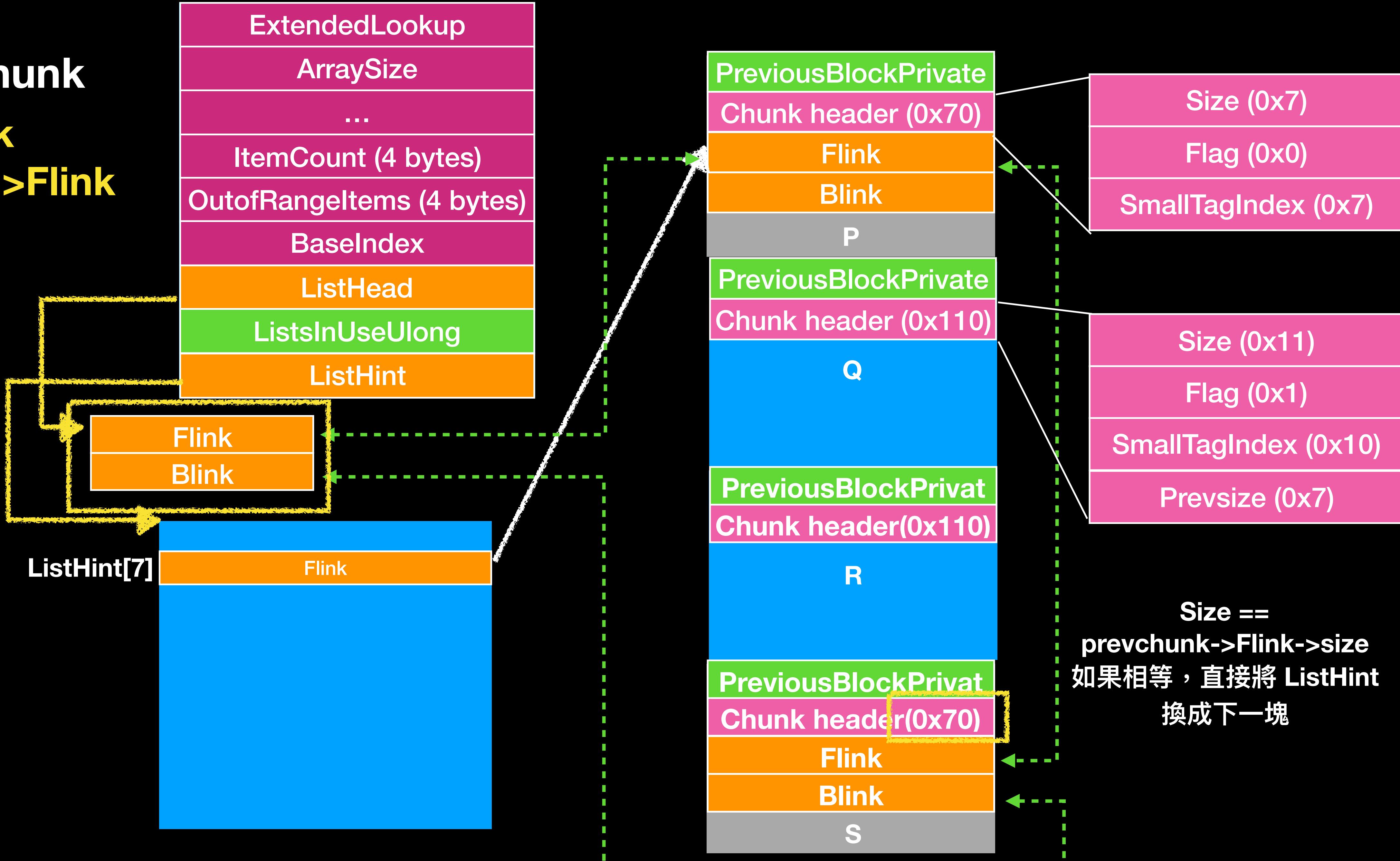


**Check**  
**prevchunk->Flink ==**  
**ListHead**  
成立的話會準備 update  
**ListHint**

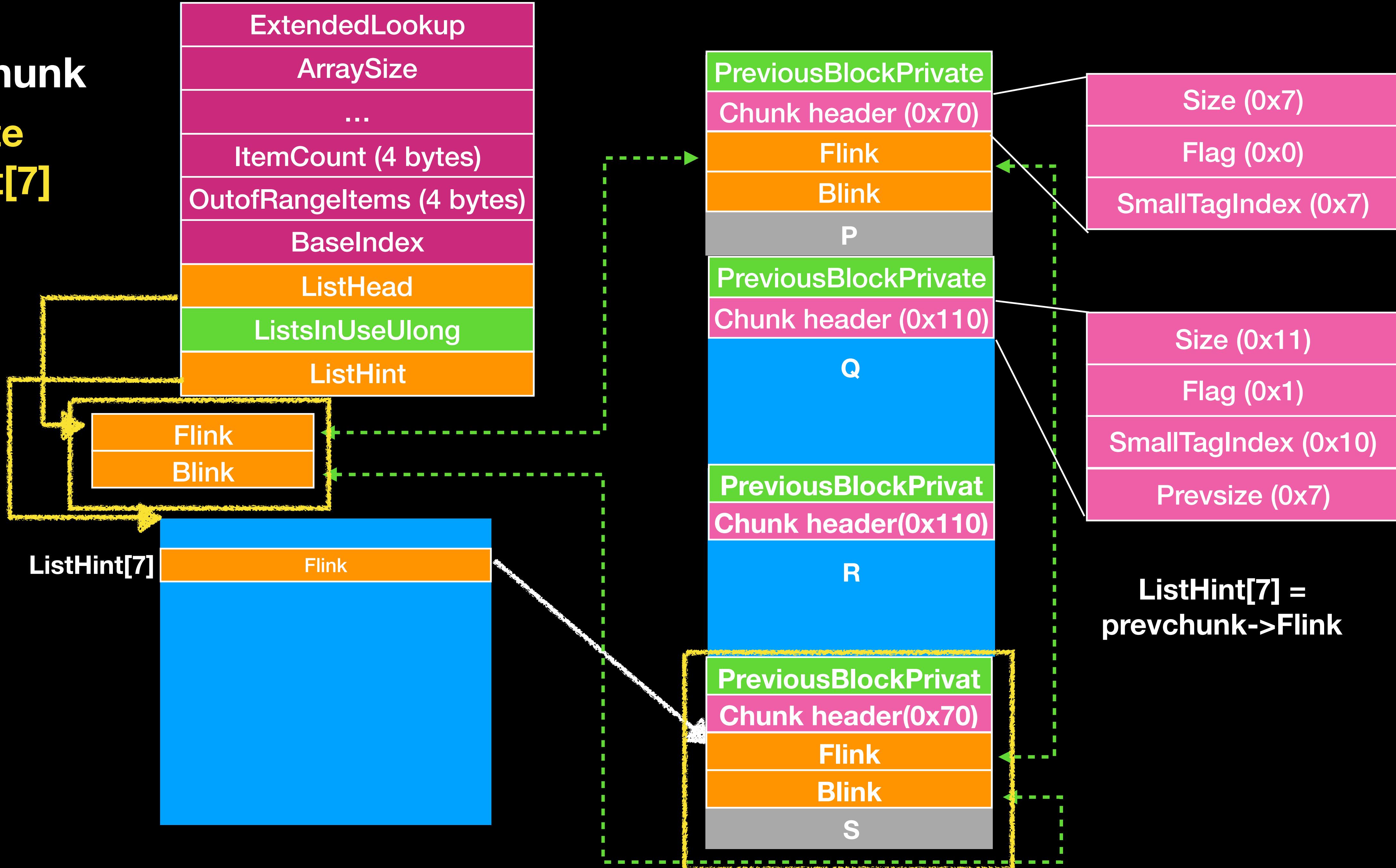
merge chunk  
Check  
ListHint[7]->Flink



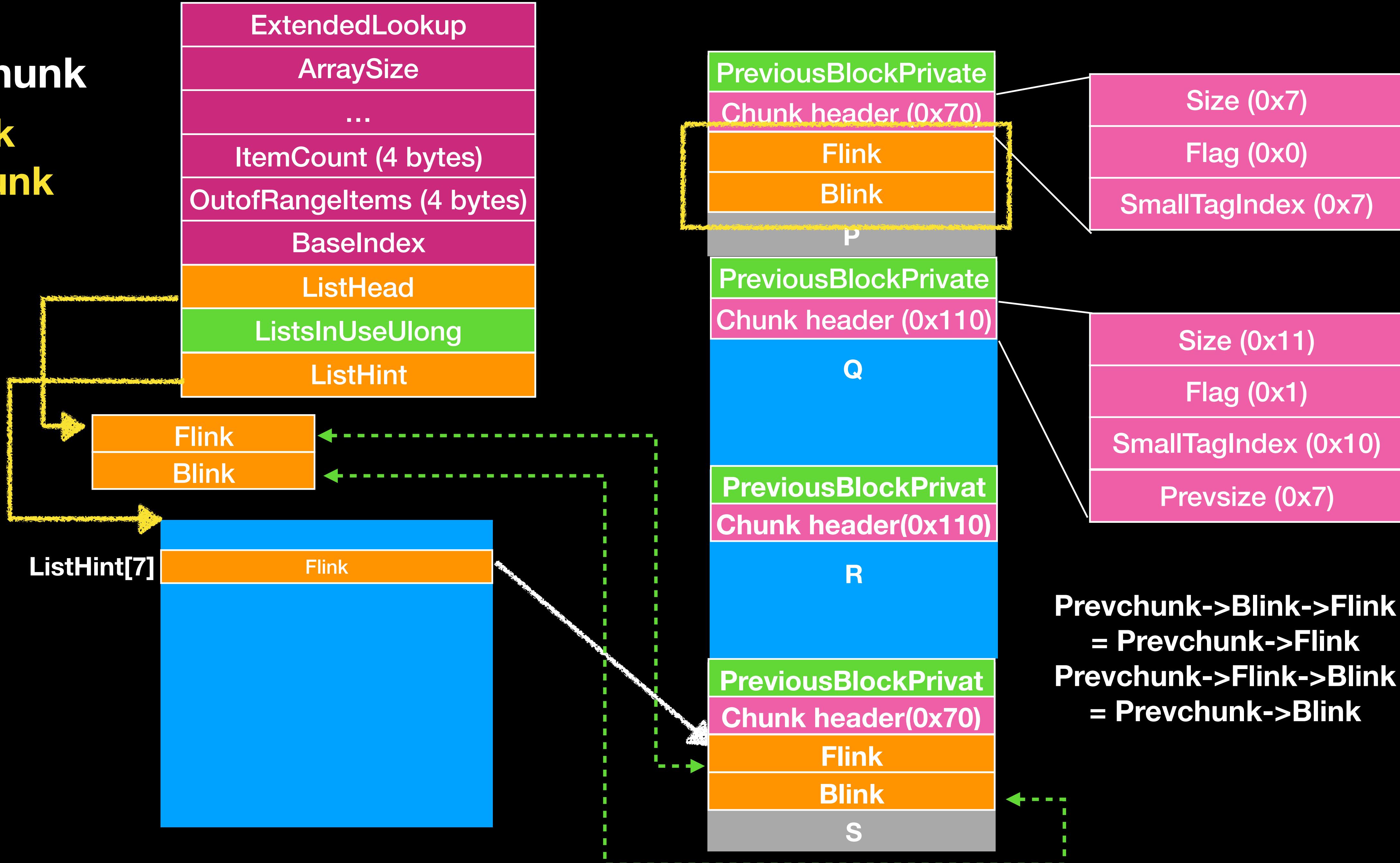
merge chunk  
Check  
listHint[7]->Flink



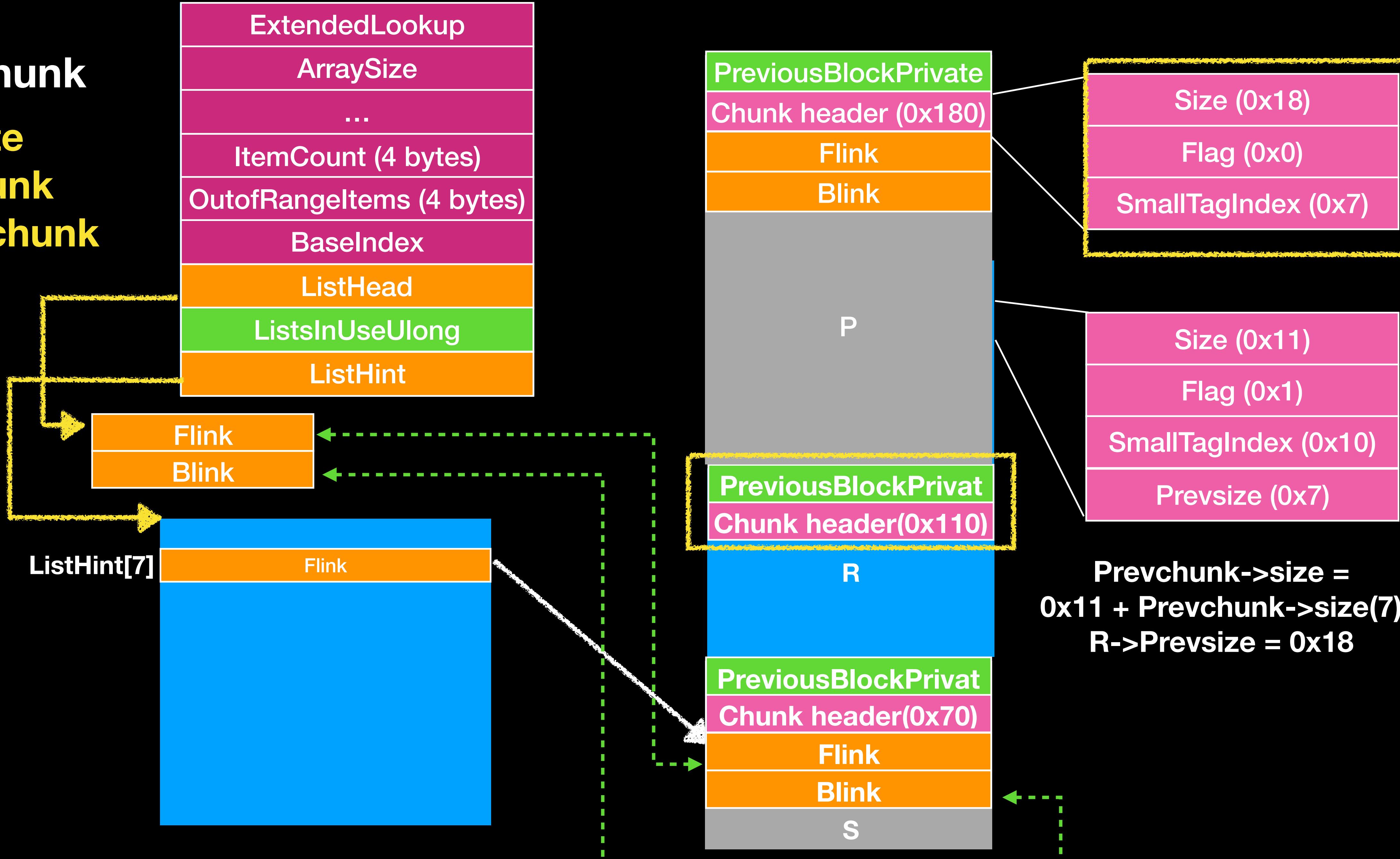
**merge chunk**  
**Update**  
**ListHint[7]**



**merge chunk**  
**Unlink**  
**prevchunk**

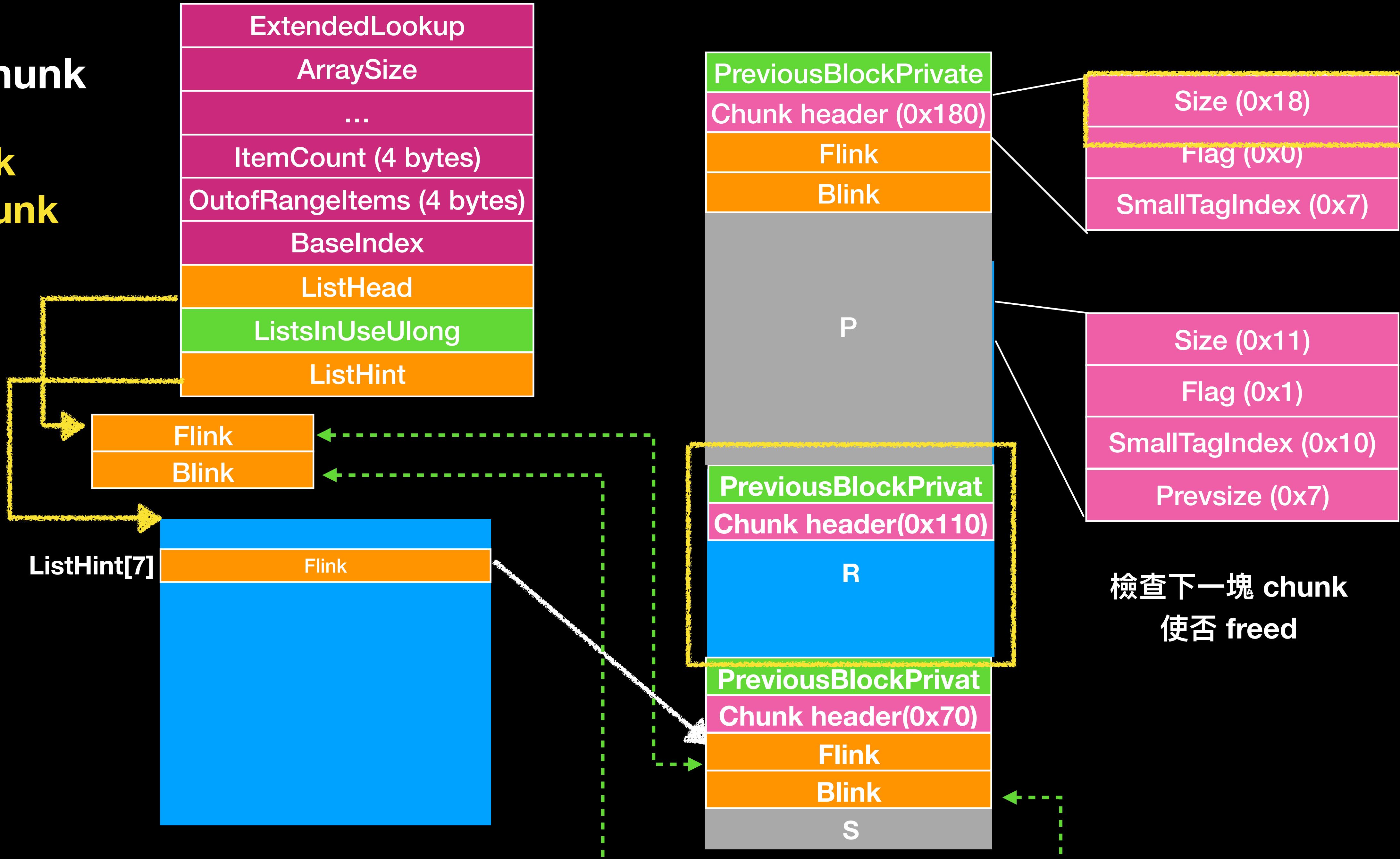


**merge chunk**  
**Update  
prevchunk  
and next chunk**

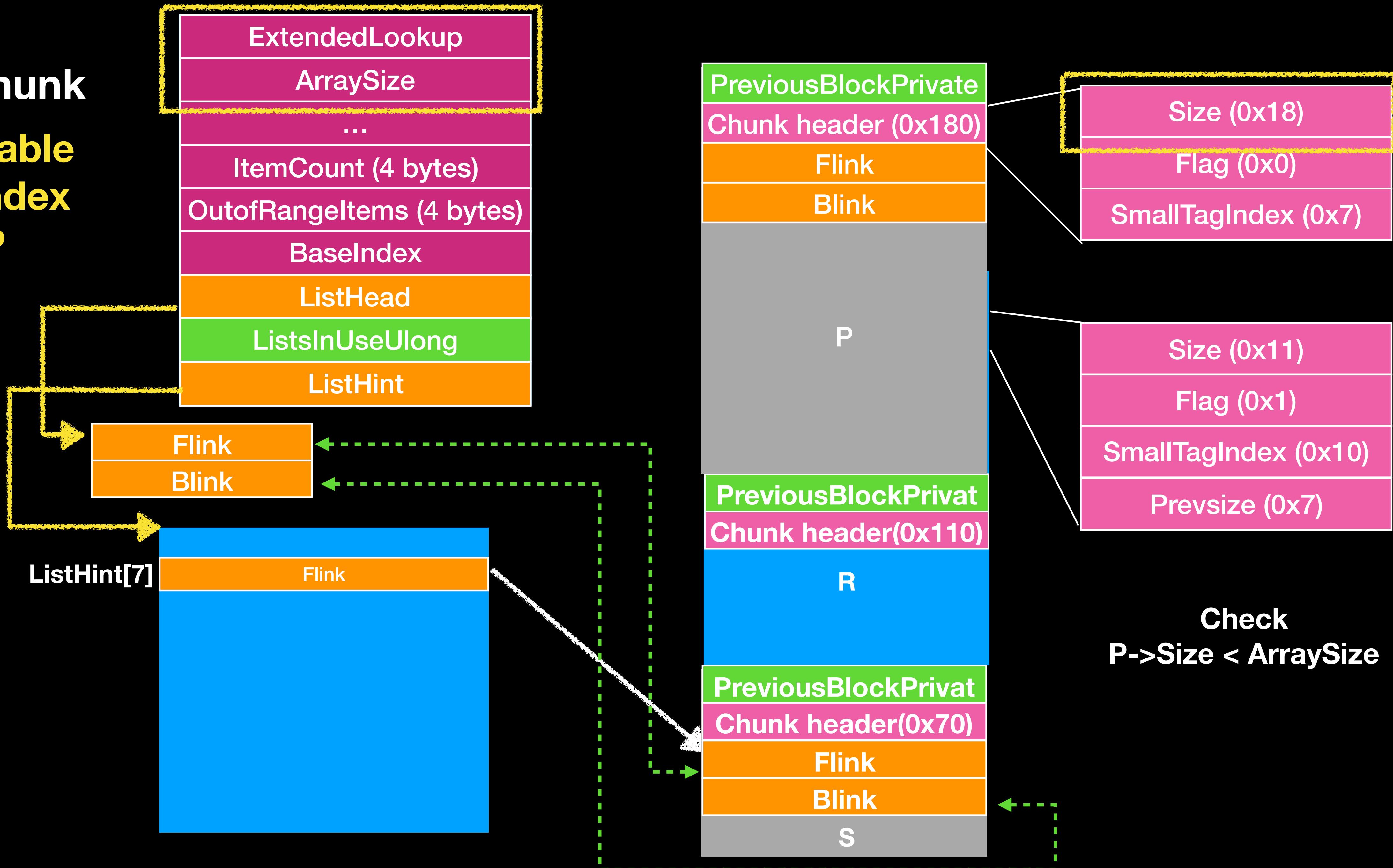


merge chunk

Check  
next chunk

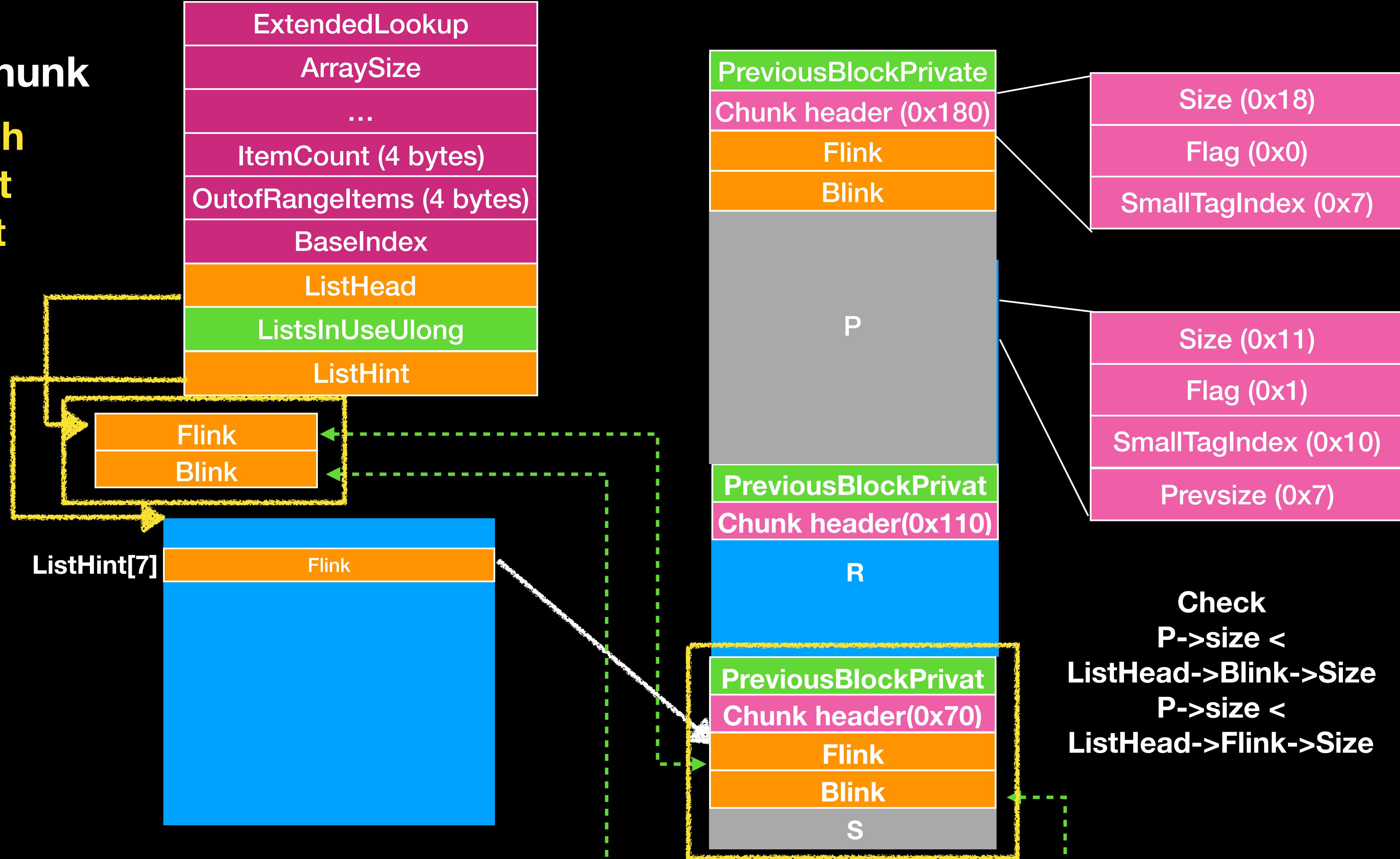


merge chunk  
Find suitable  
**BlocksIndex**  
For P



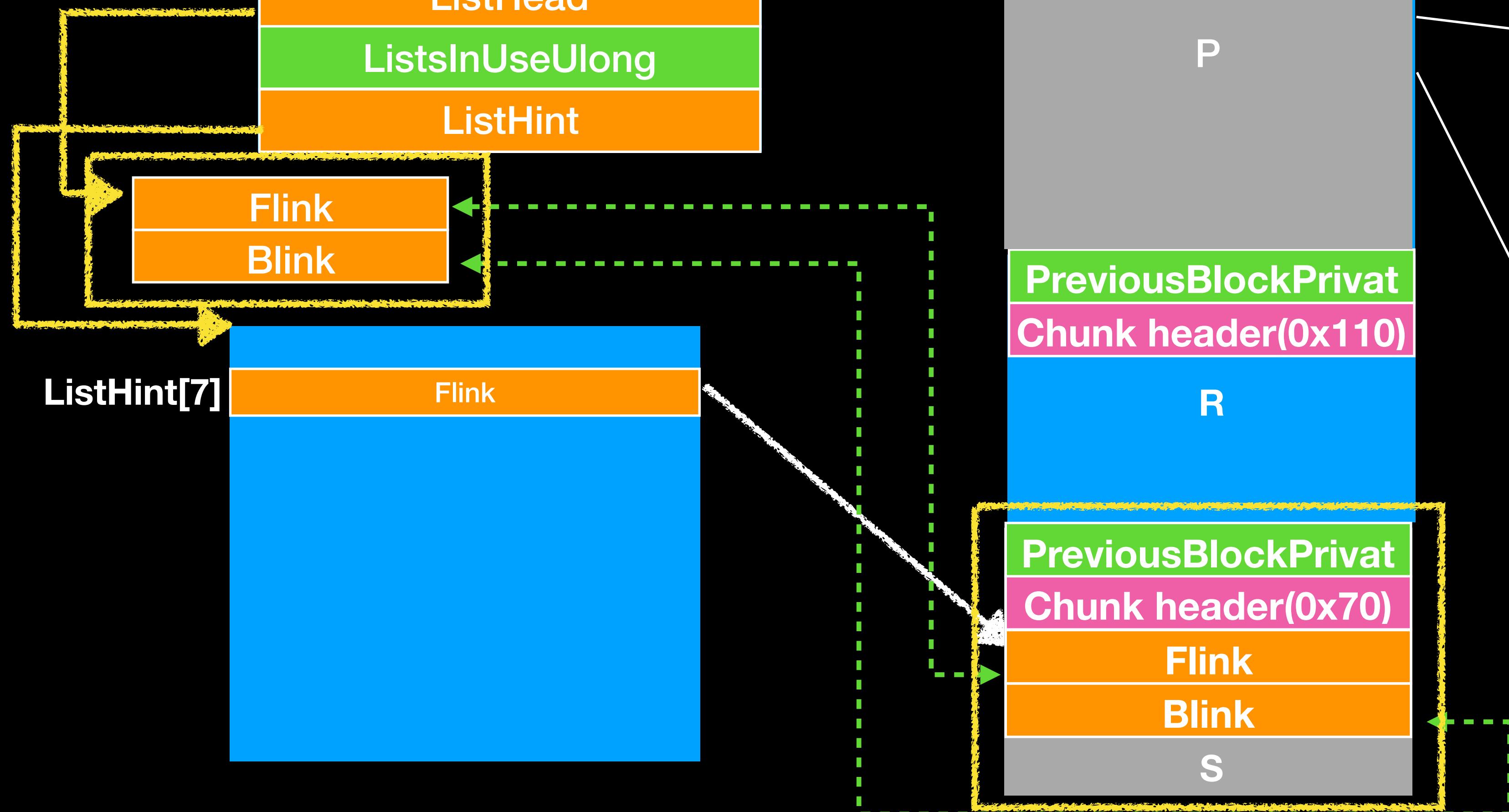
**merge chunk**

**Search  
insert  
point**



**merge chunk**

**Insert  
linked list**



Size (0x18)

Flag (0x0)

SmallTagIndex (0x7)

Size (0x11)

Flag (0x1)

SmallTagIndex (0x10)

Preysize (0x7)

Check

**S->Blink->Flink**

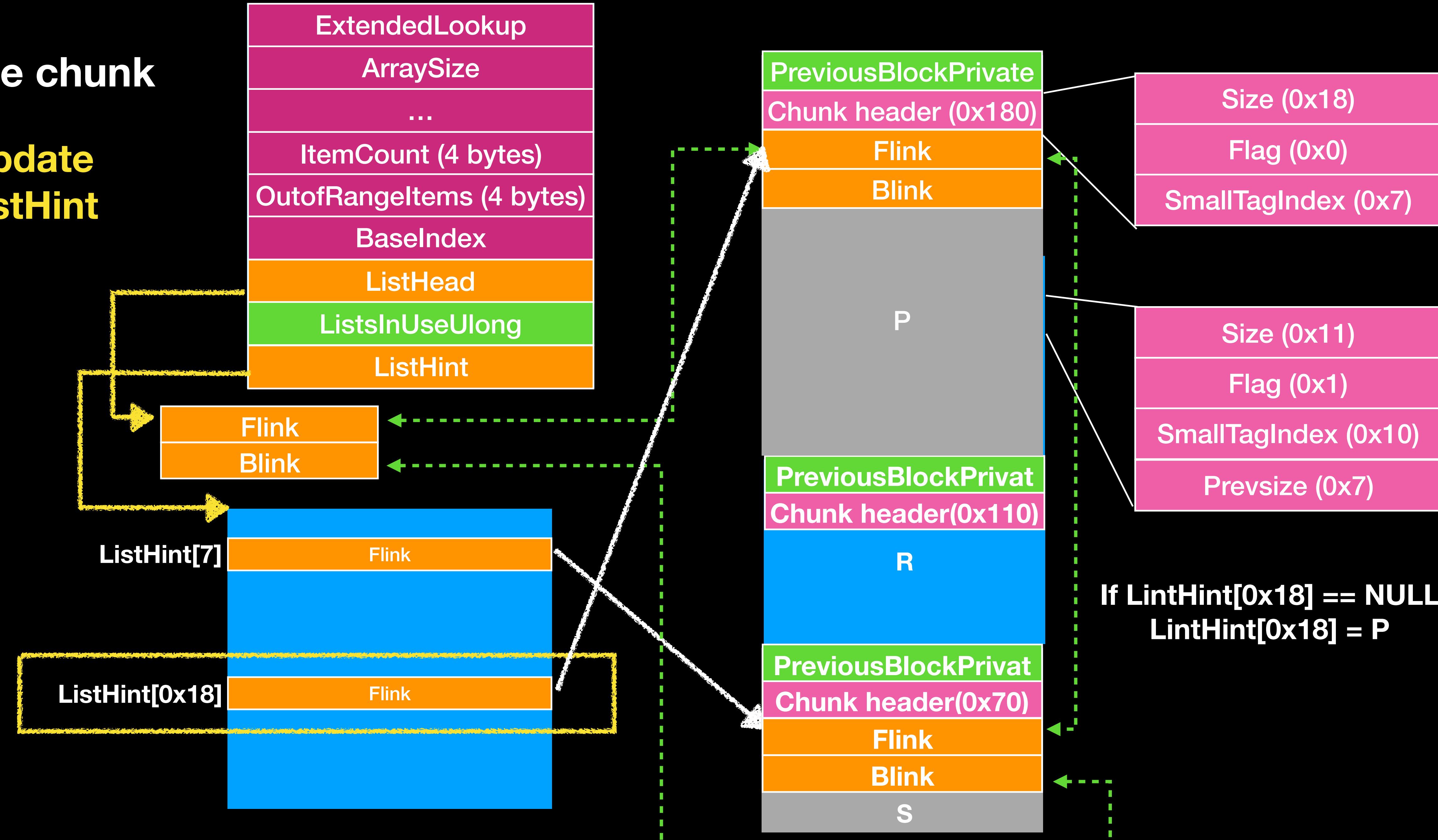
**== S**

**If not pass it will not abort**

**But it will not unlink**

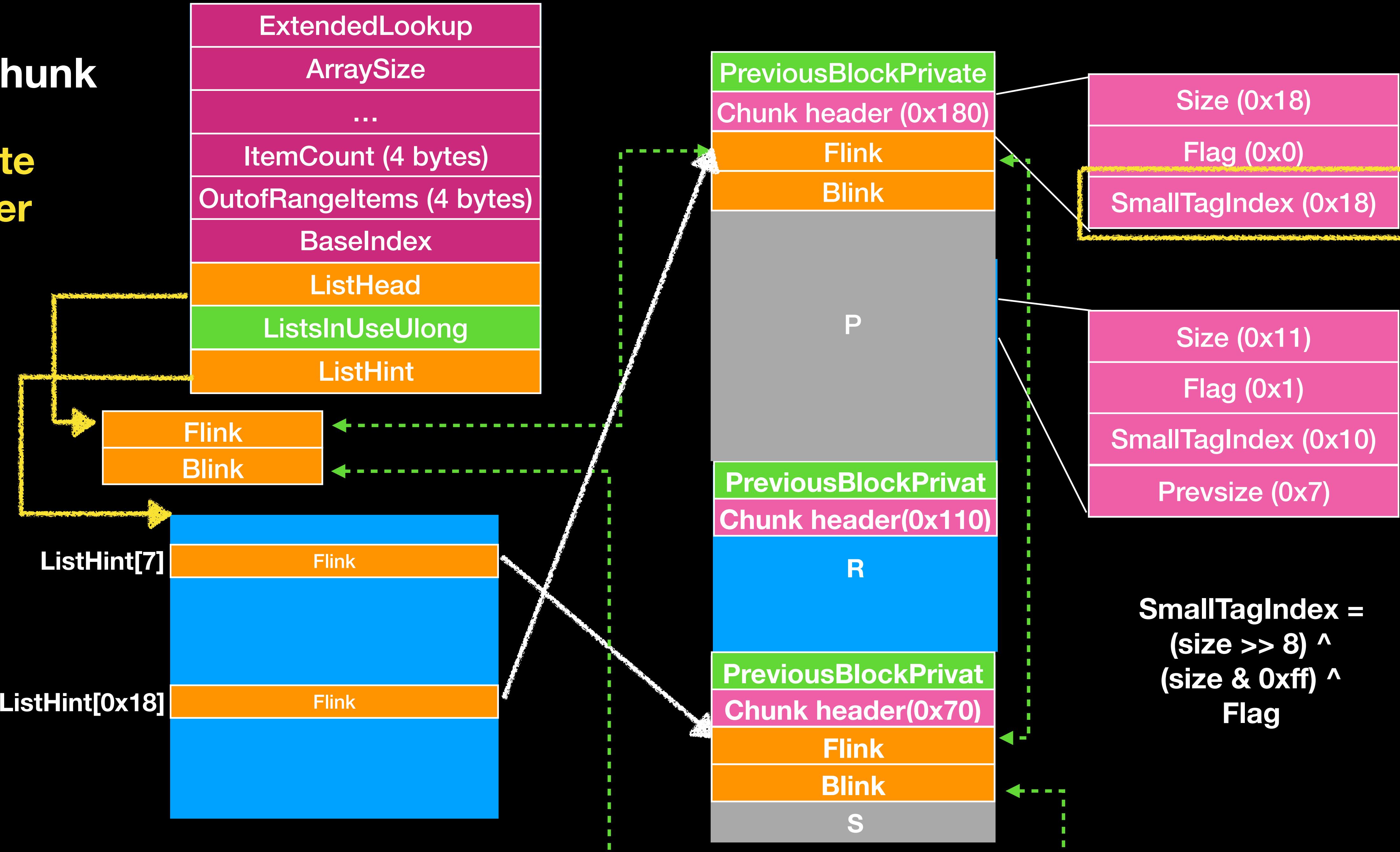
**merge chunk**

**Update  
ListHint**



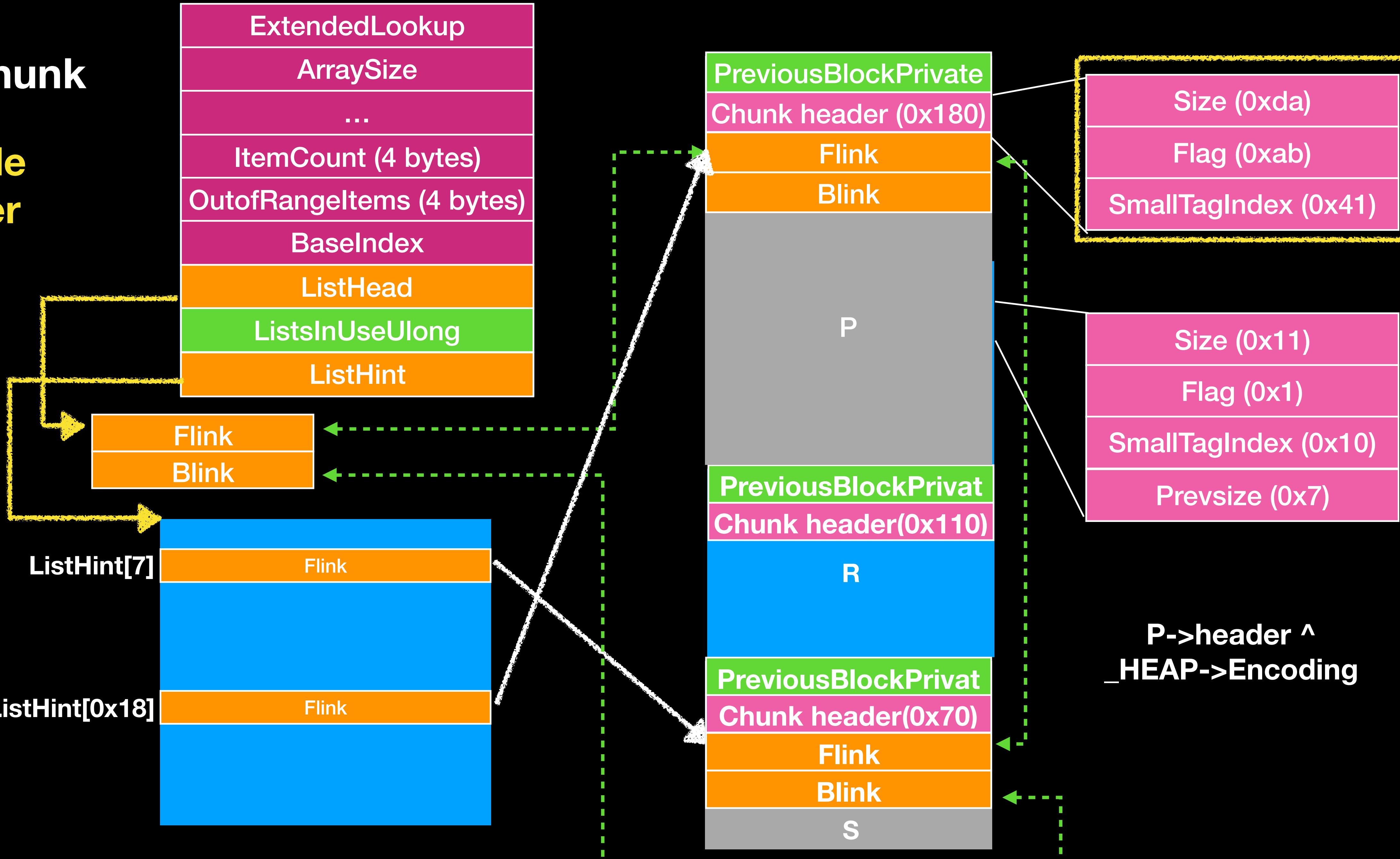
merge chunk

Update header



# merge chunk

# Encode header



# Nt heap

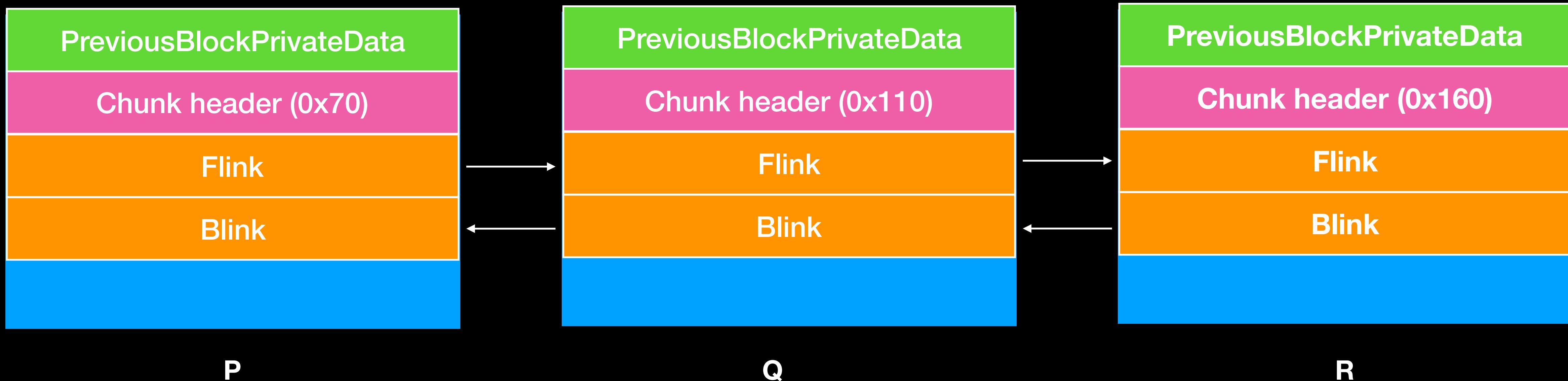
- Free (RtlpFreeHeap)
  - Size > 0xff000
  - 檢查該 chunk 的 linked list 並從 \_HEAP->VirtualAllocdBlocks 移除
  - 接著使用 RtlpSecMemFreeVirtualMemory 將 chunk 整個 munmap 掉

# Back-End Exploitation

- Unlink
  - 基本上與 Linux 中的 unlink 很類似，繞過方法也差不多
  - 簡而言之就是利用從 linked listed 移除 node 的行為來做有限制的寫入
  - 要注意的是在會 decode 的地方，都要讓他正常 decode
    - 也就是 check sum 要過
  - 另外一點是 Flink 及 Blink 並不是指向 chunk 開頭，而是直接指向 User data 部分，也就是不太需要做偏移偽造 chunk
    - 所以找到一個指向該 userdata 的 pointer 讓他繞過 double linked list 的驗證就好了

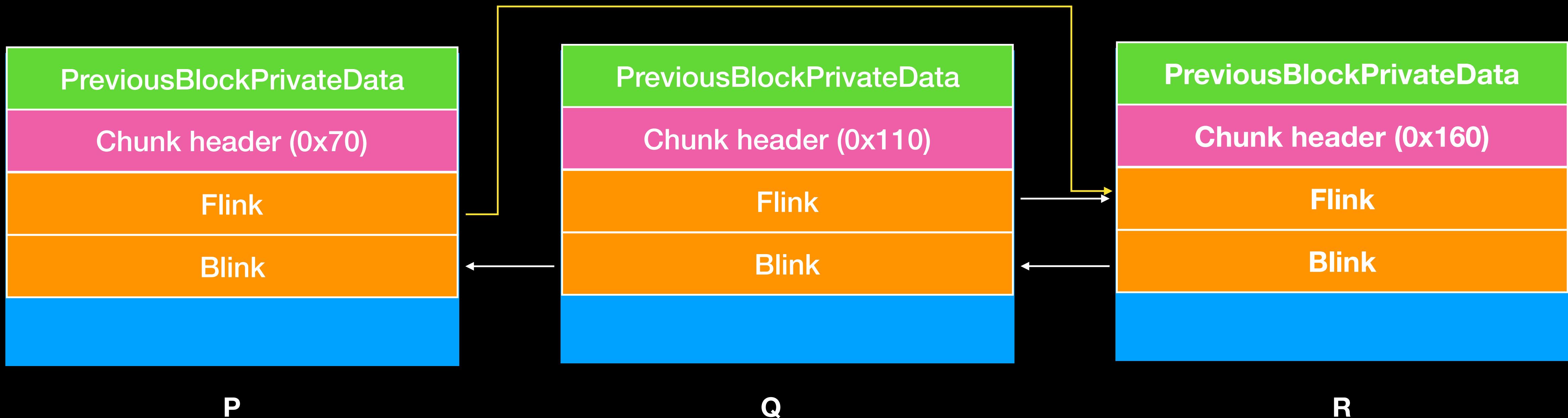
# Back-End Exploitation

- Unlink
  - $Q->\text{Blink}->\text{Flink} = Q->\text{Flink}$
  - $Q->\text{Flink}->\text{Blink} = Q->\text{Blink}$



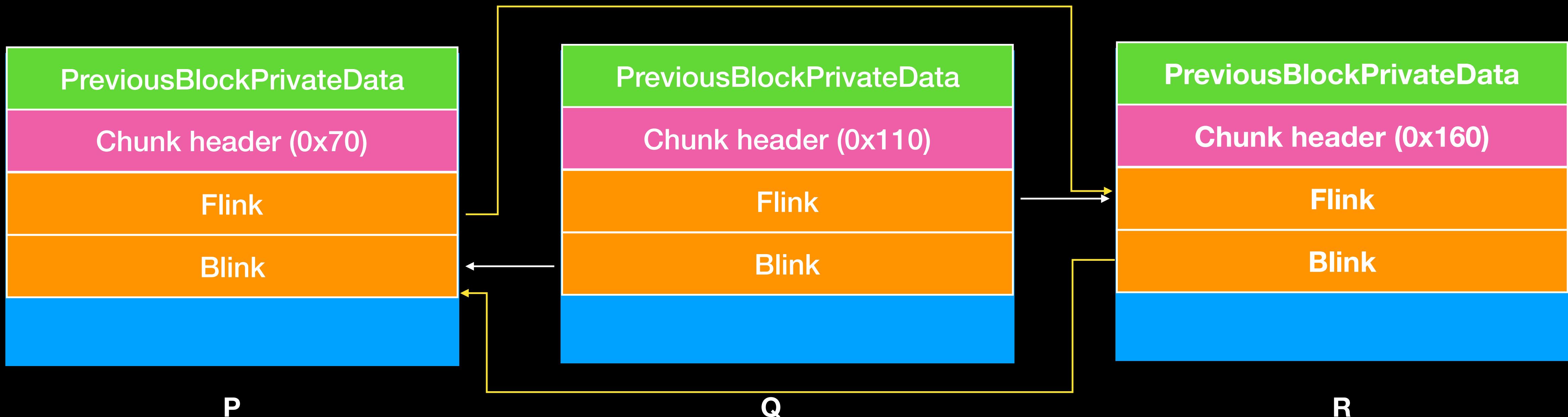
# Back-End Exploitation

- Unlink
  - $Q \rightarrow \text{Blink} \rightarrow \text{Flink} = Q \rightarrow \text{Flink}$
  - $Q \rightarrow \text{Flink} \rightarrow \text{Blink} = Q \rightarrow \text{Blink}$



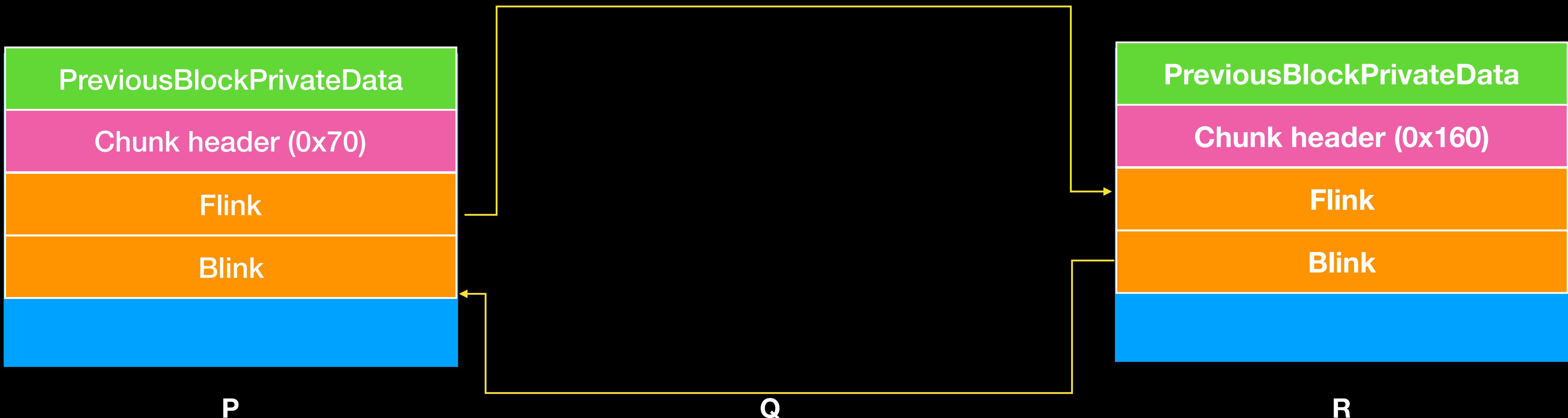
# Back-End Exploitation

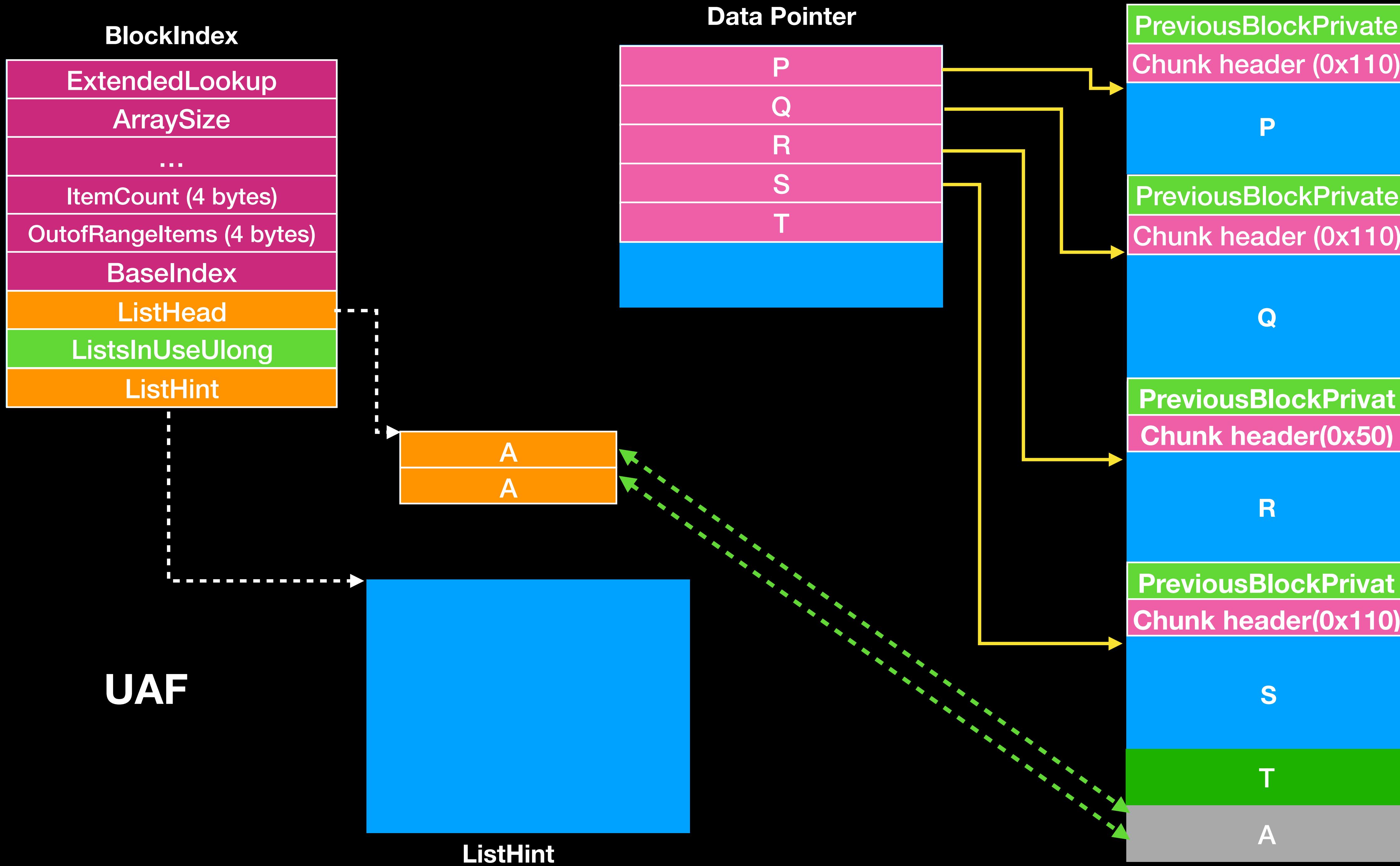
- Unlink
  - $Q \rightarrow \text{Blink} \rightarrow \text{Flink} = Q \rightarrow \text{Flink}$
  - $Q \rightarrow \text{Flink} \rightarrow \text{Blink} = Q \rightarrow \text{Blink}$

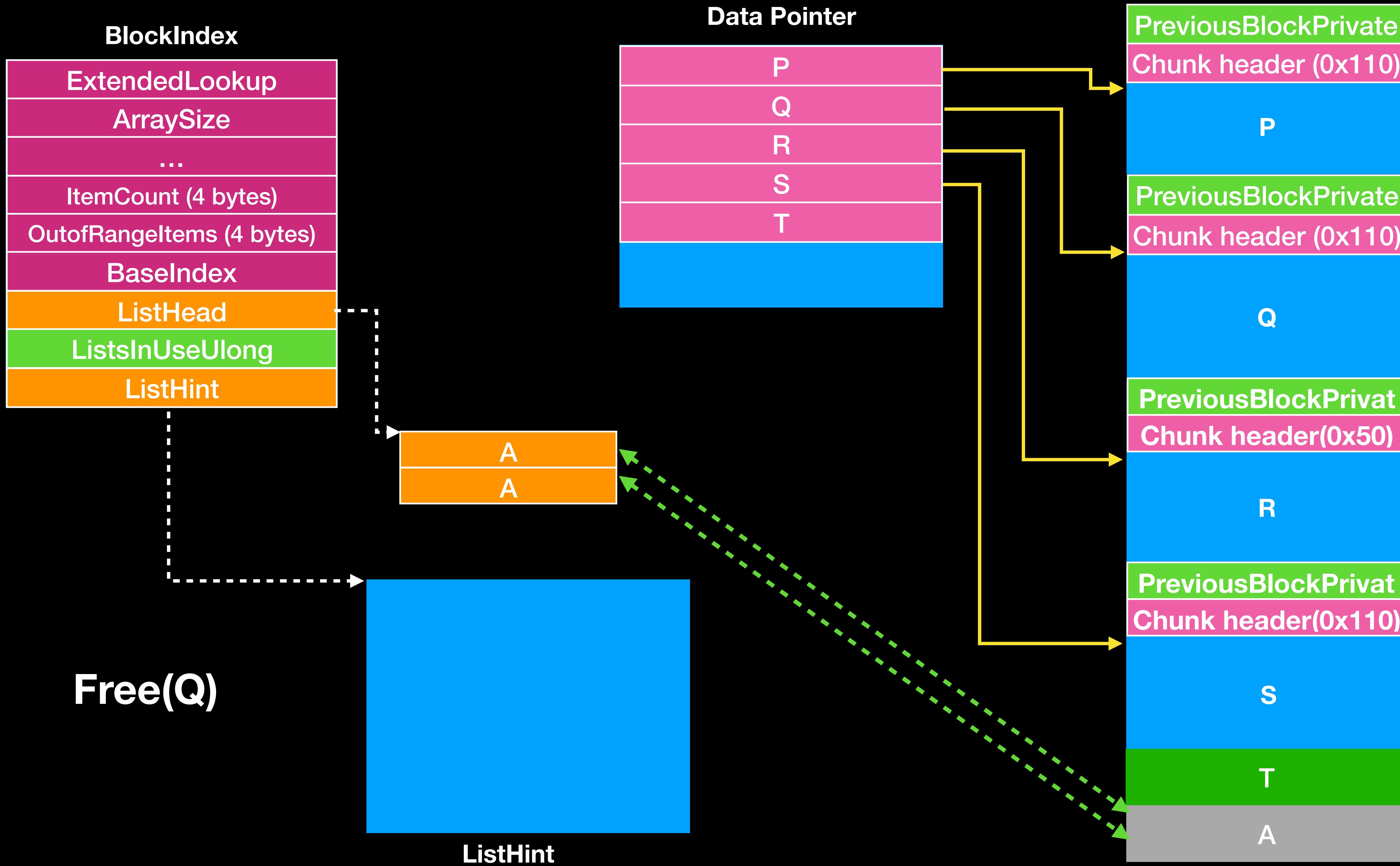


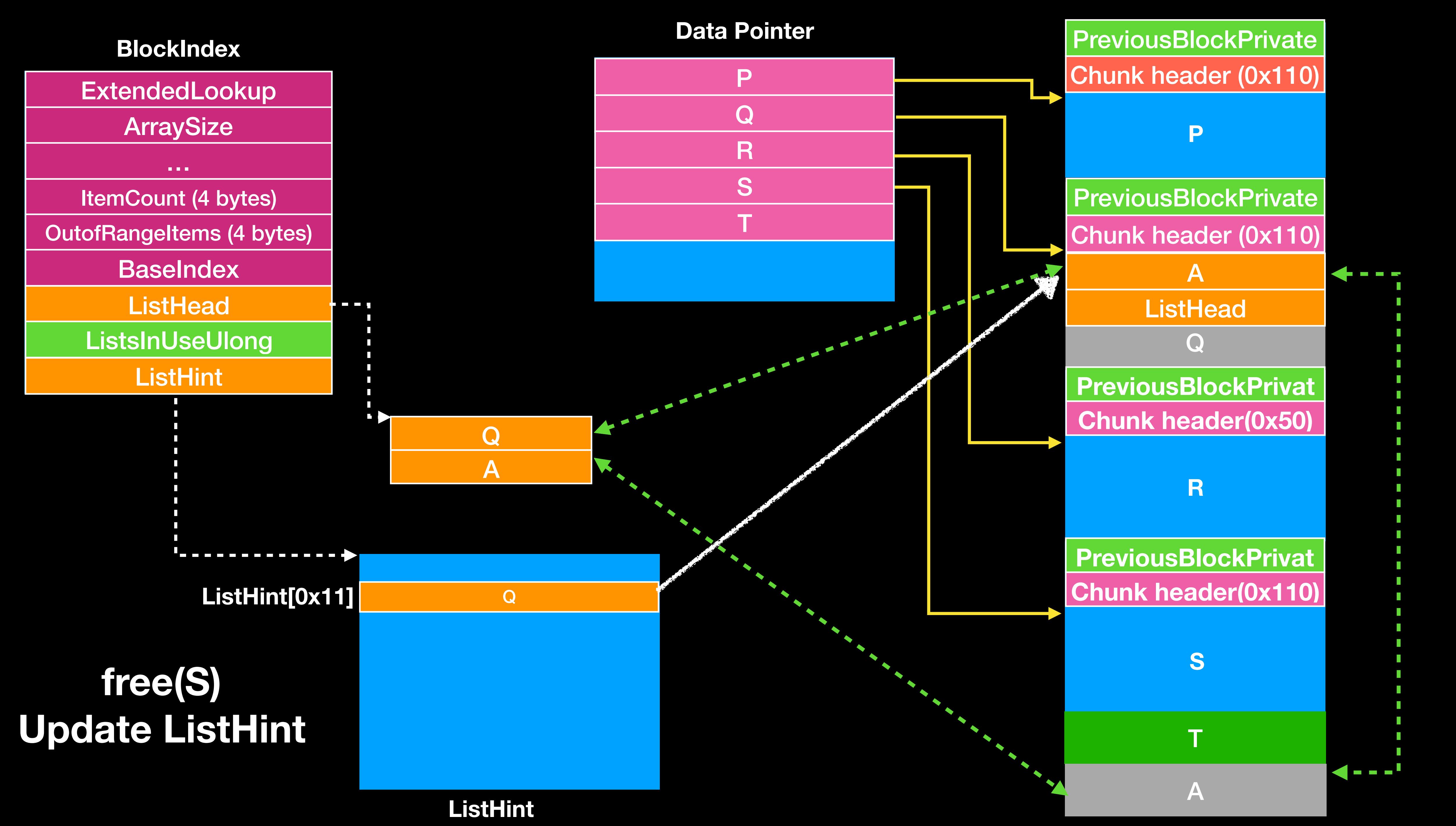
# Back-End Exploitation

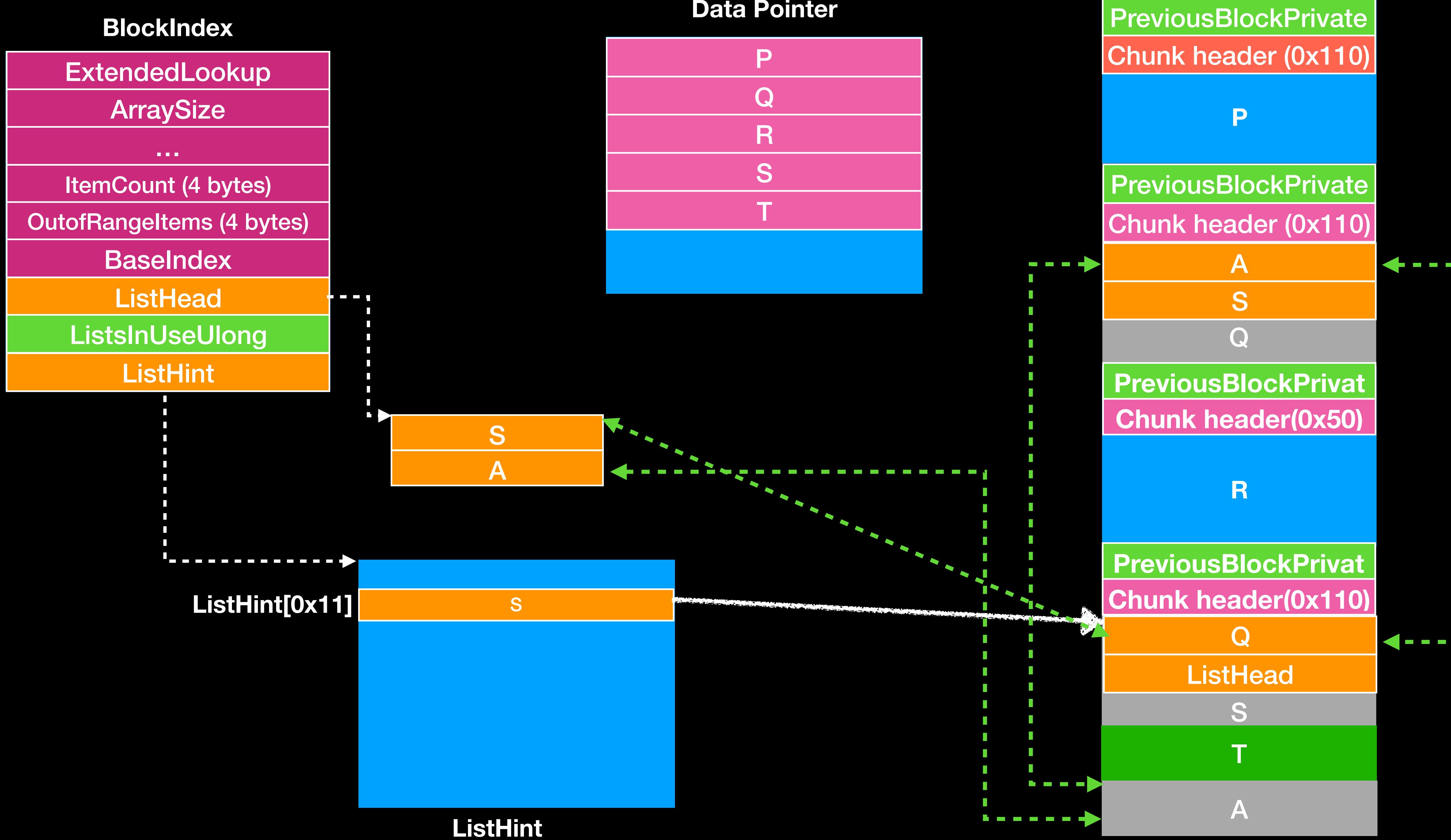
- Unlink
  - $Q \rightarrow \text{Blink} \rightarrow \text{Flink} = Q \rightarrow \text{Flink}$
  - $Q \rightarrow \text{Flink} \rightarrow \text{Blink} = Q \rightarrow \text{Blink}$

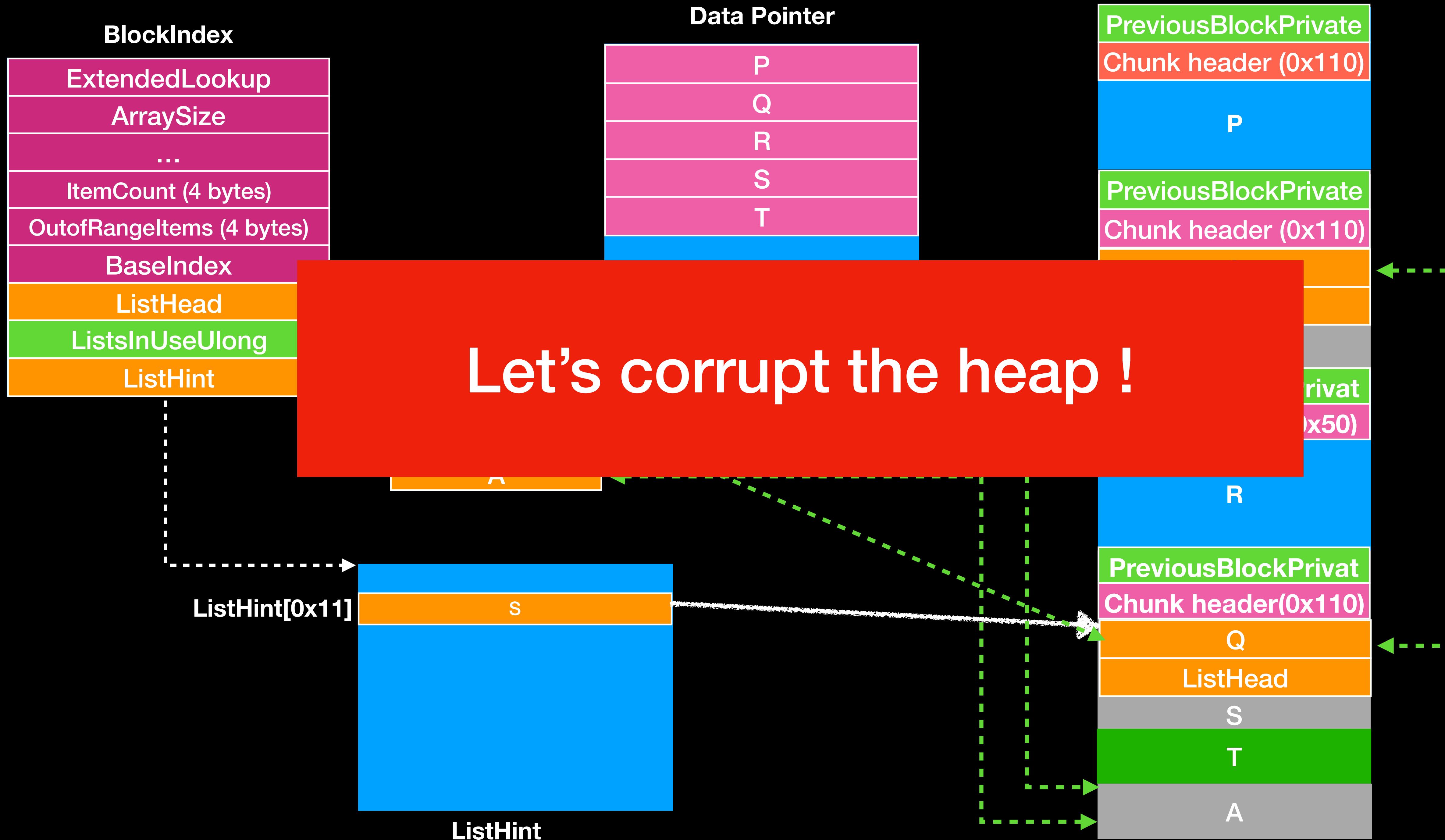


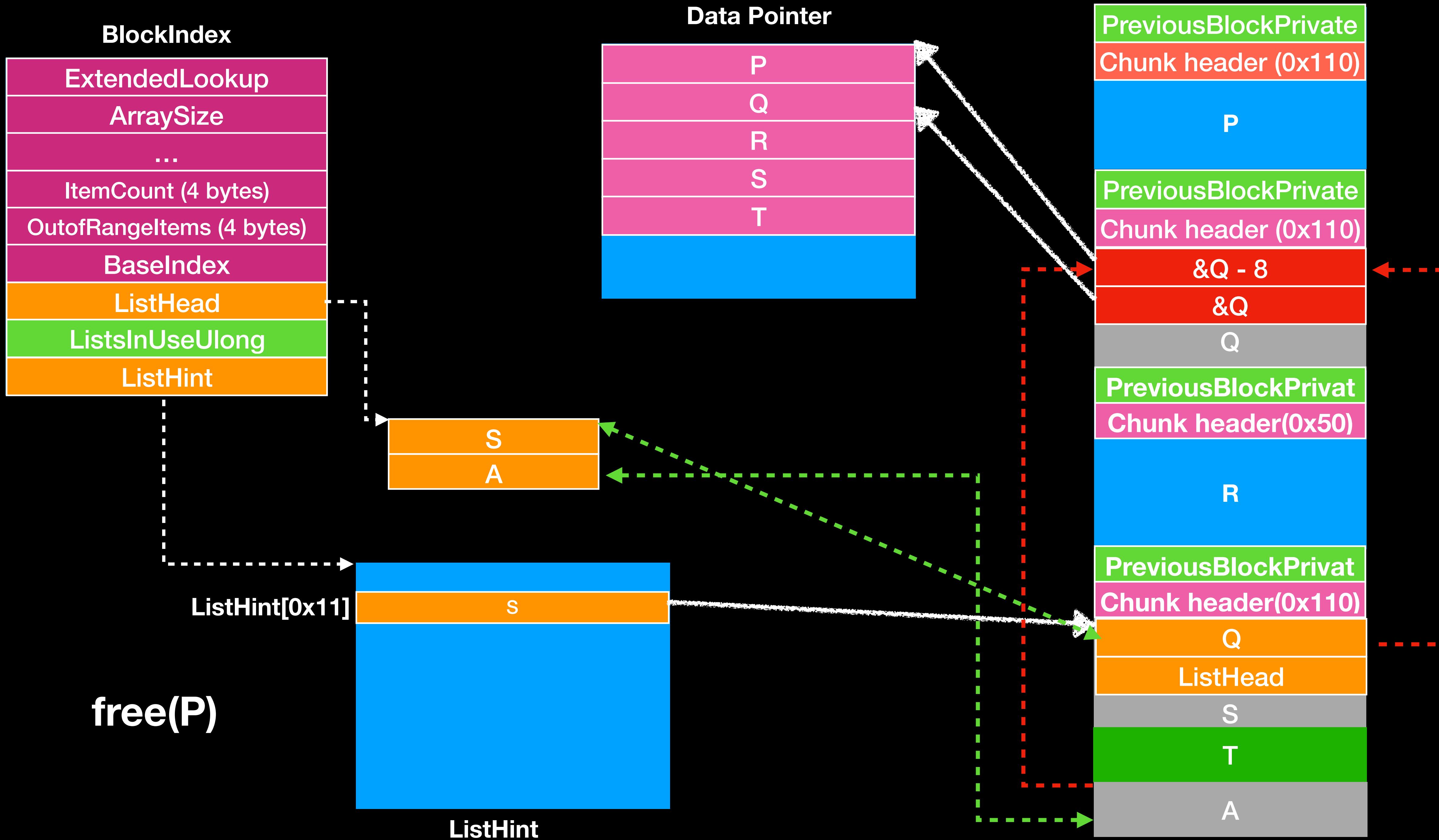


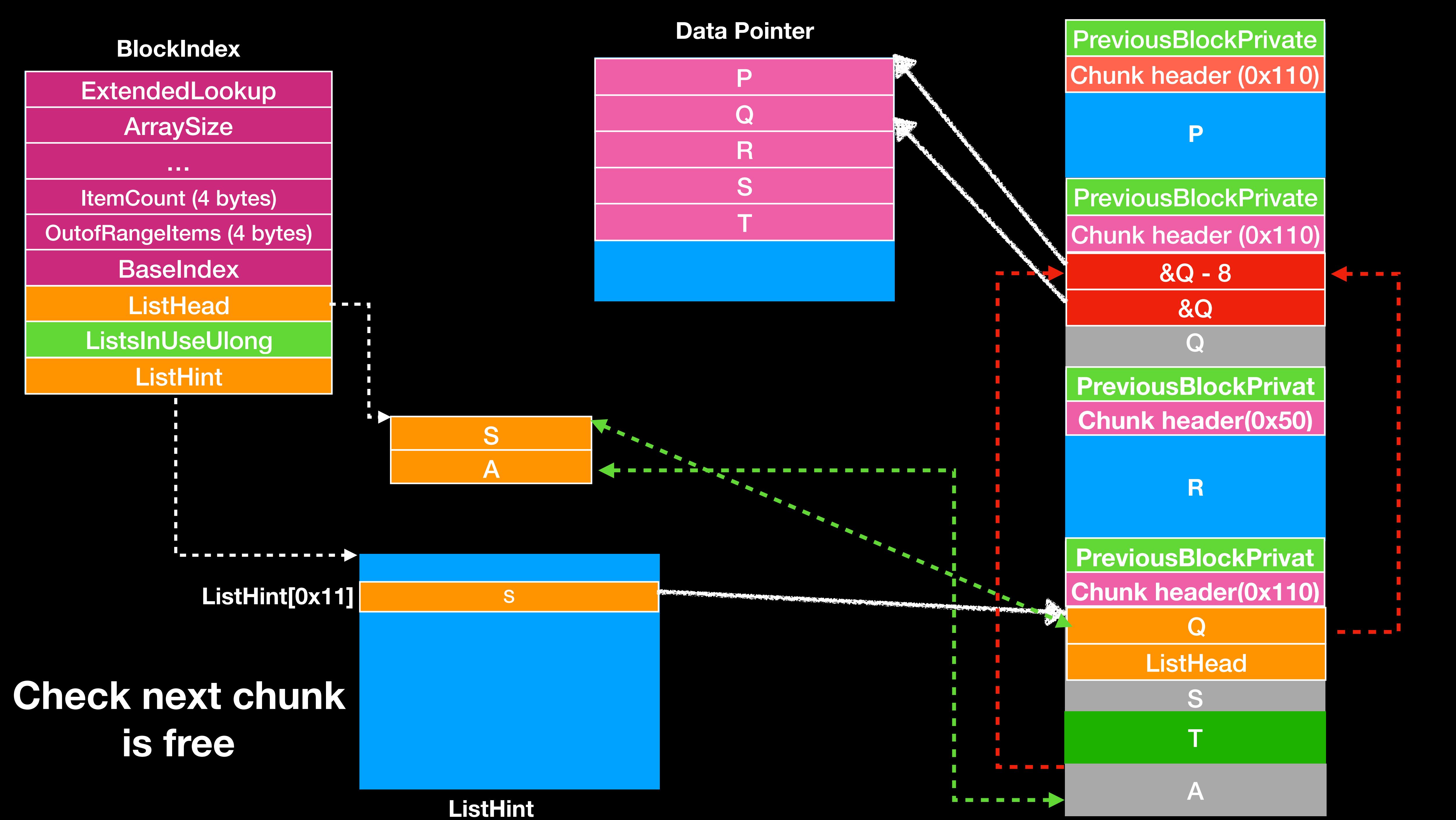


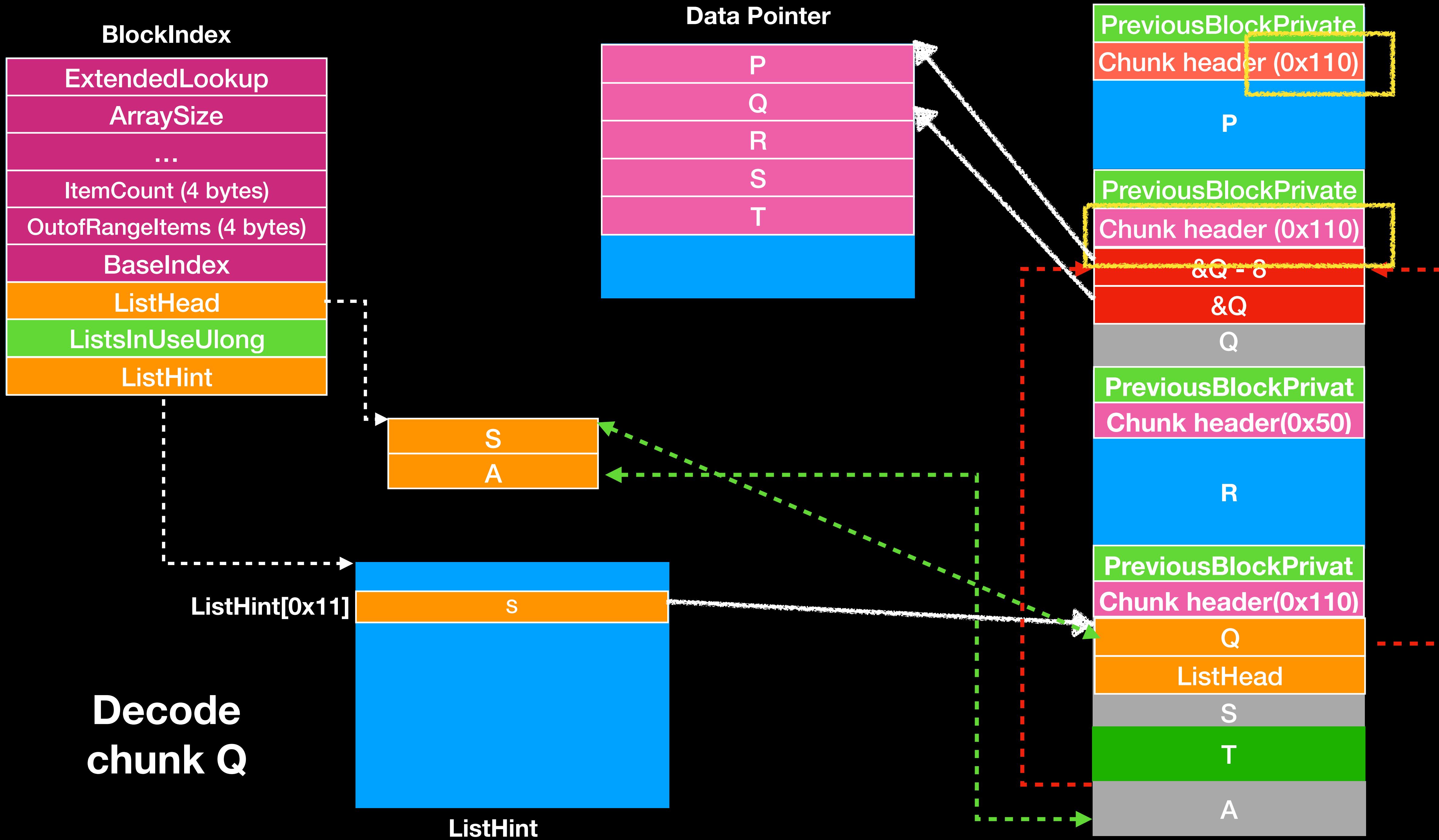


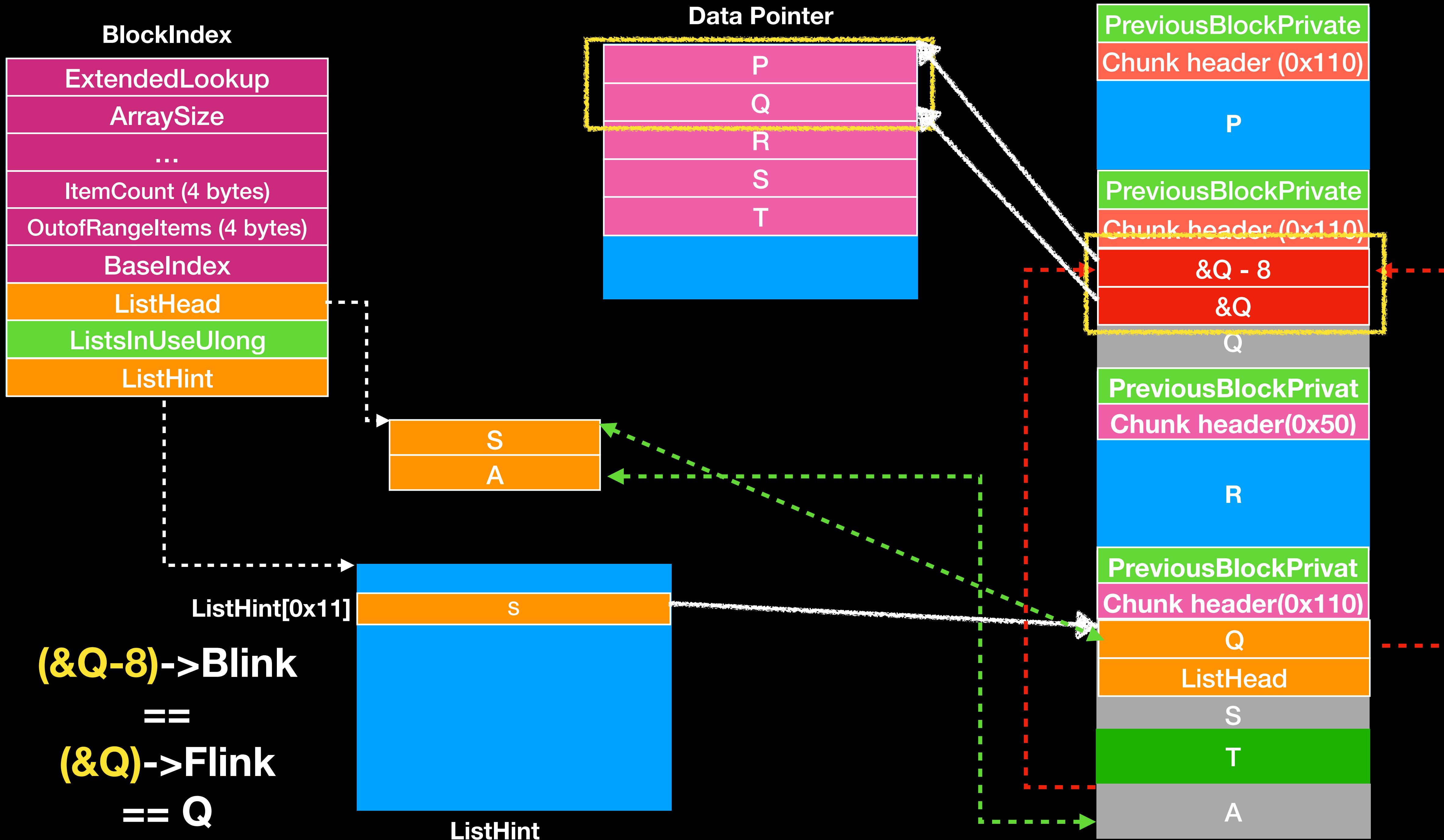


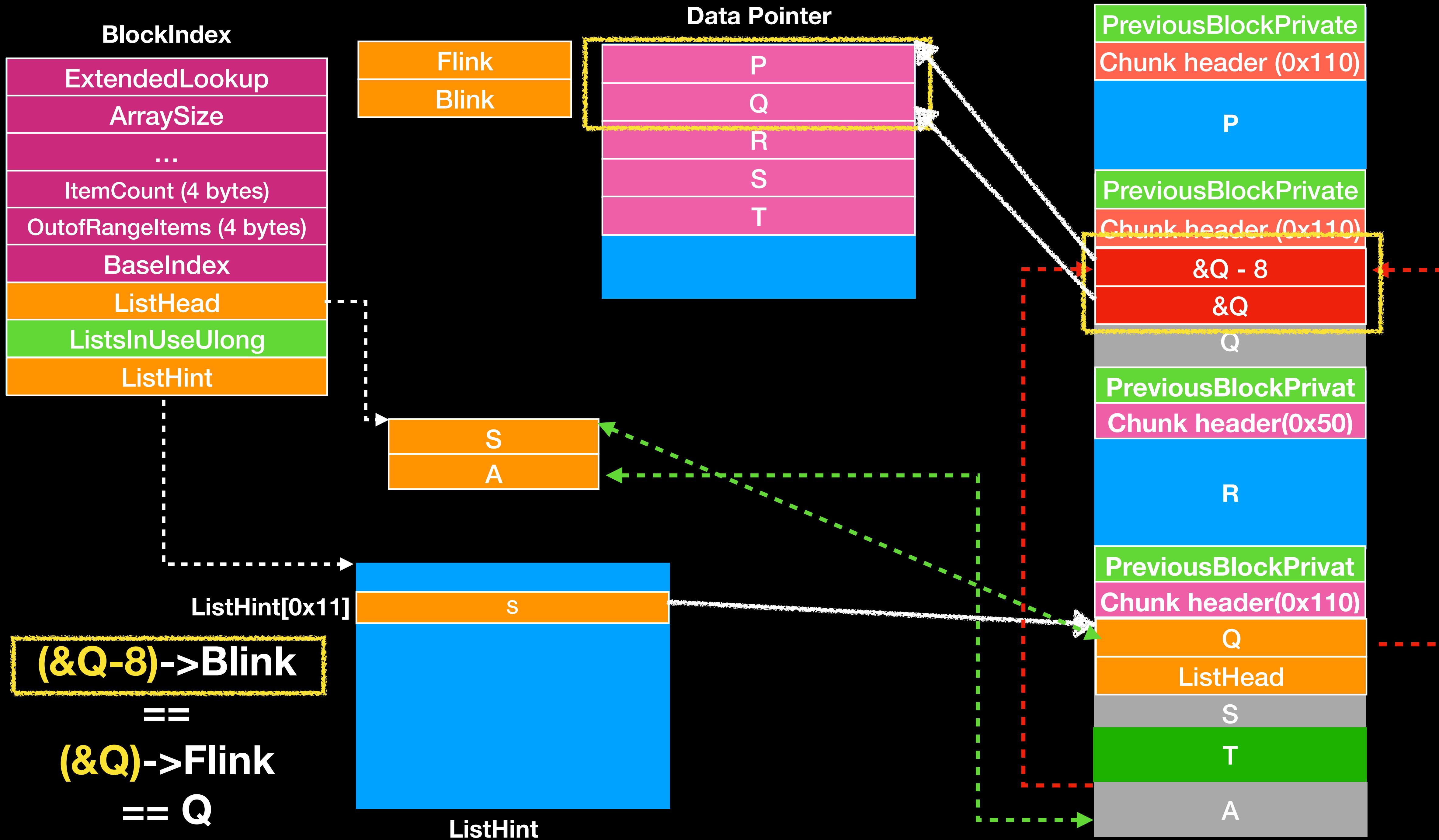


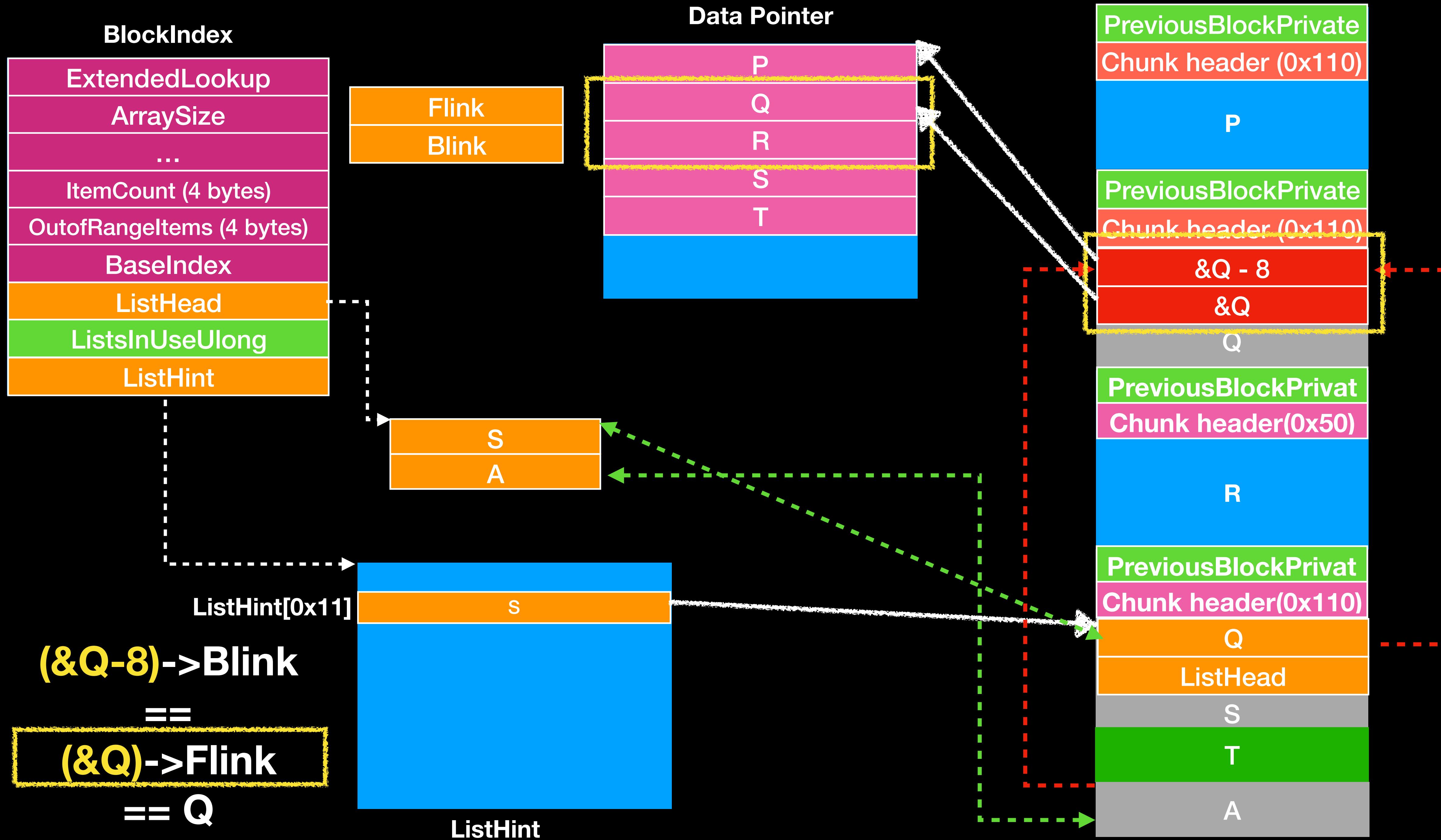


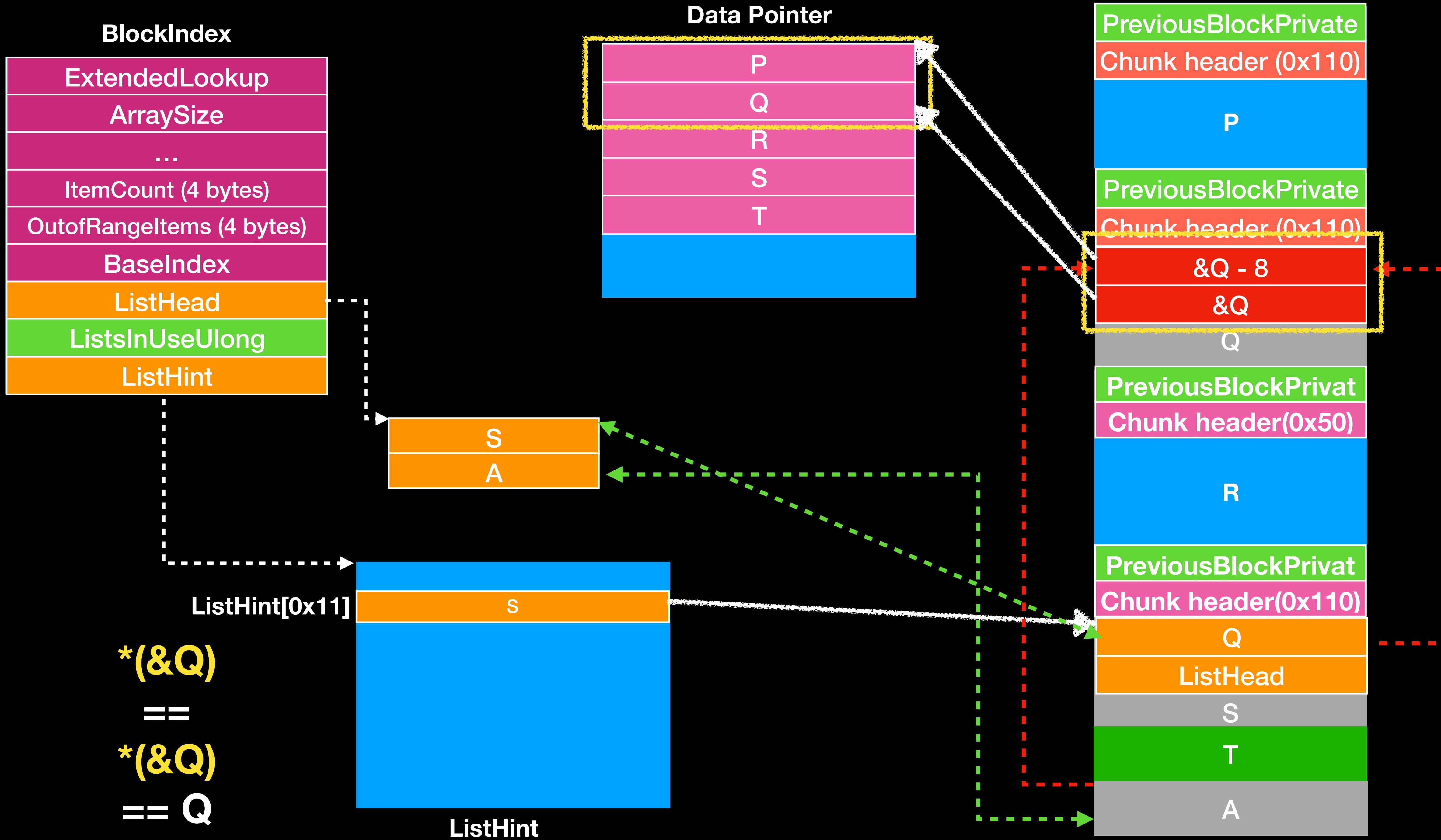


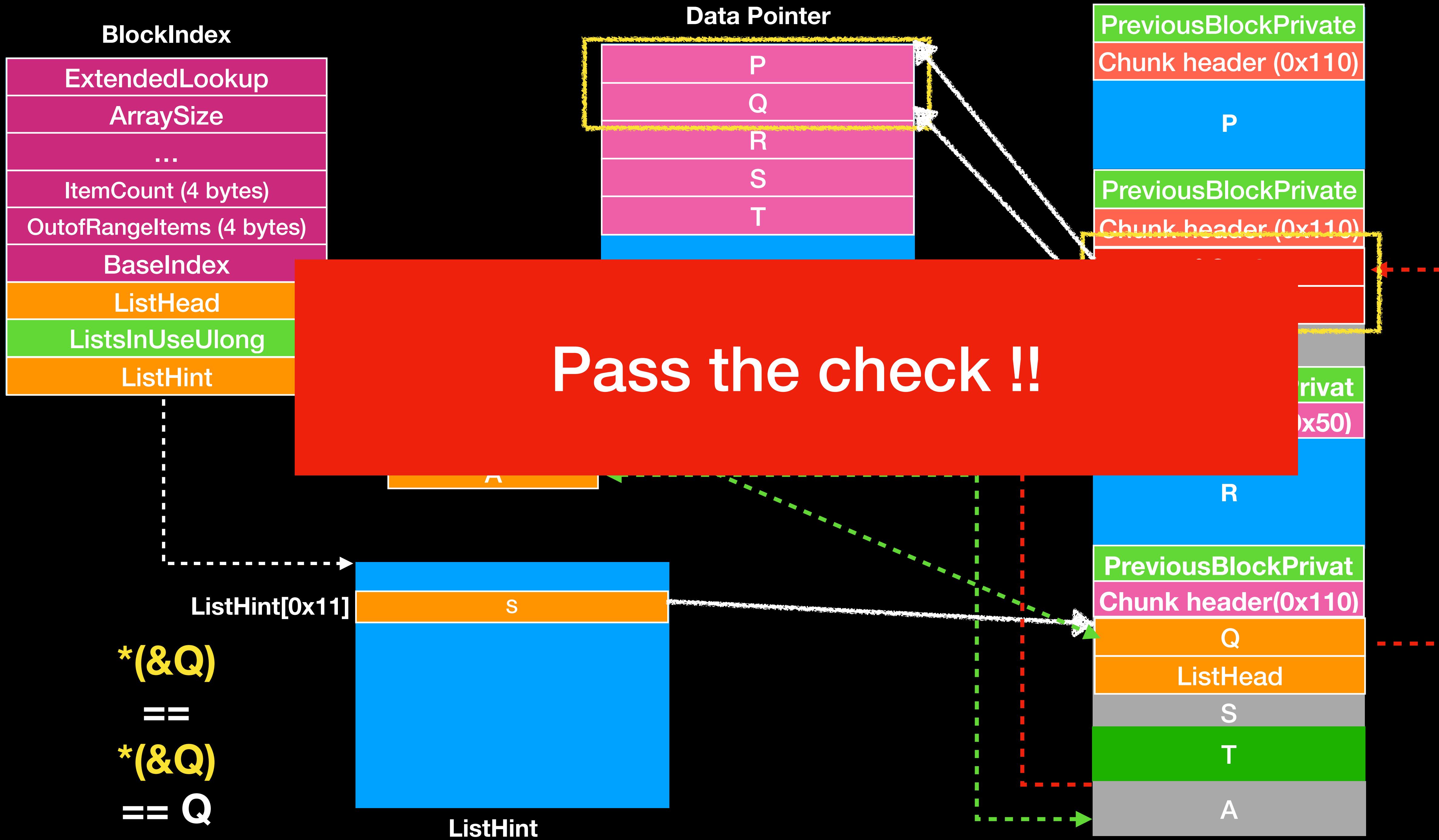


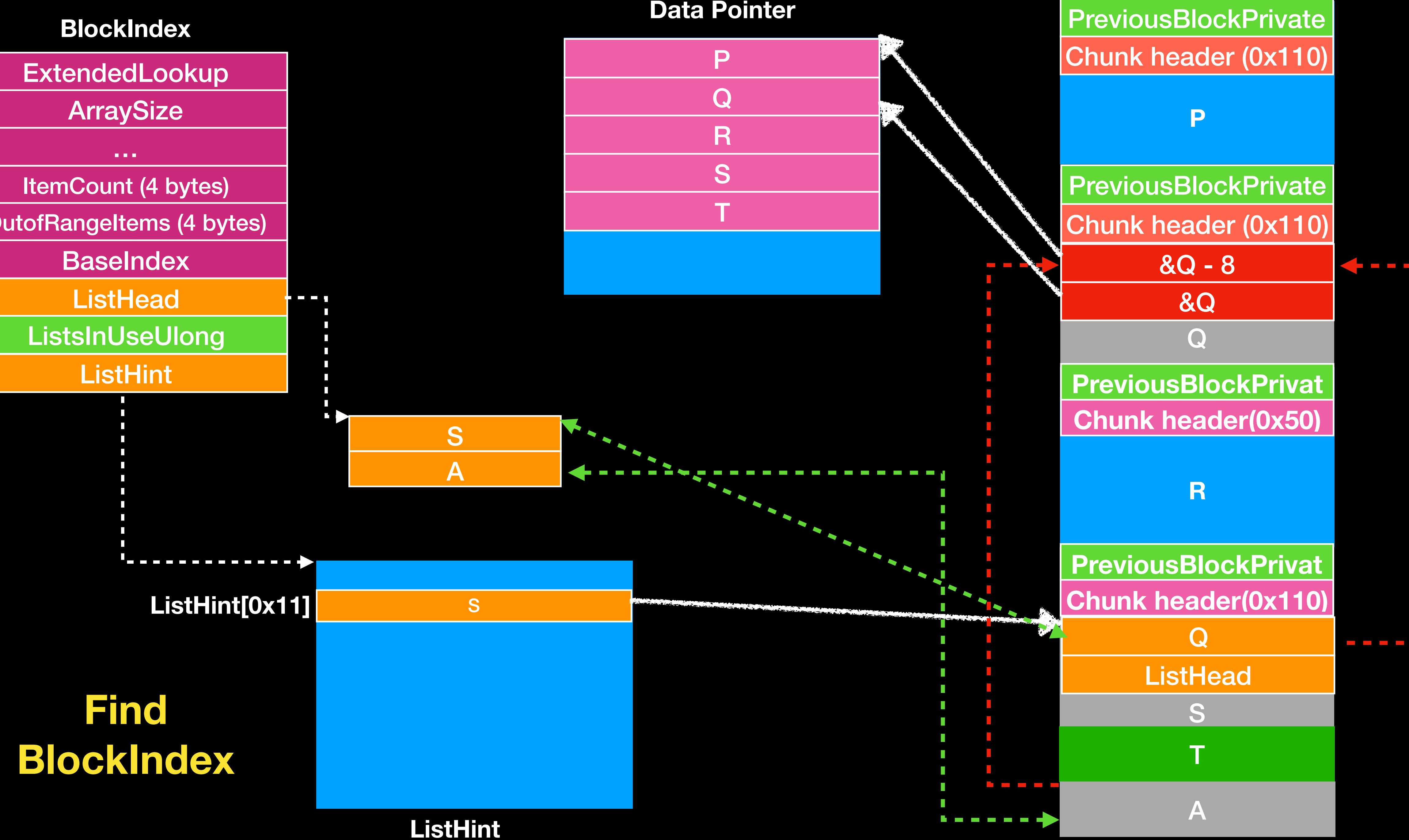


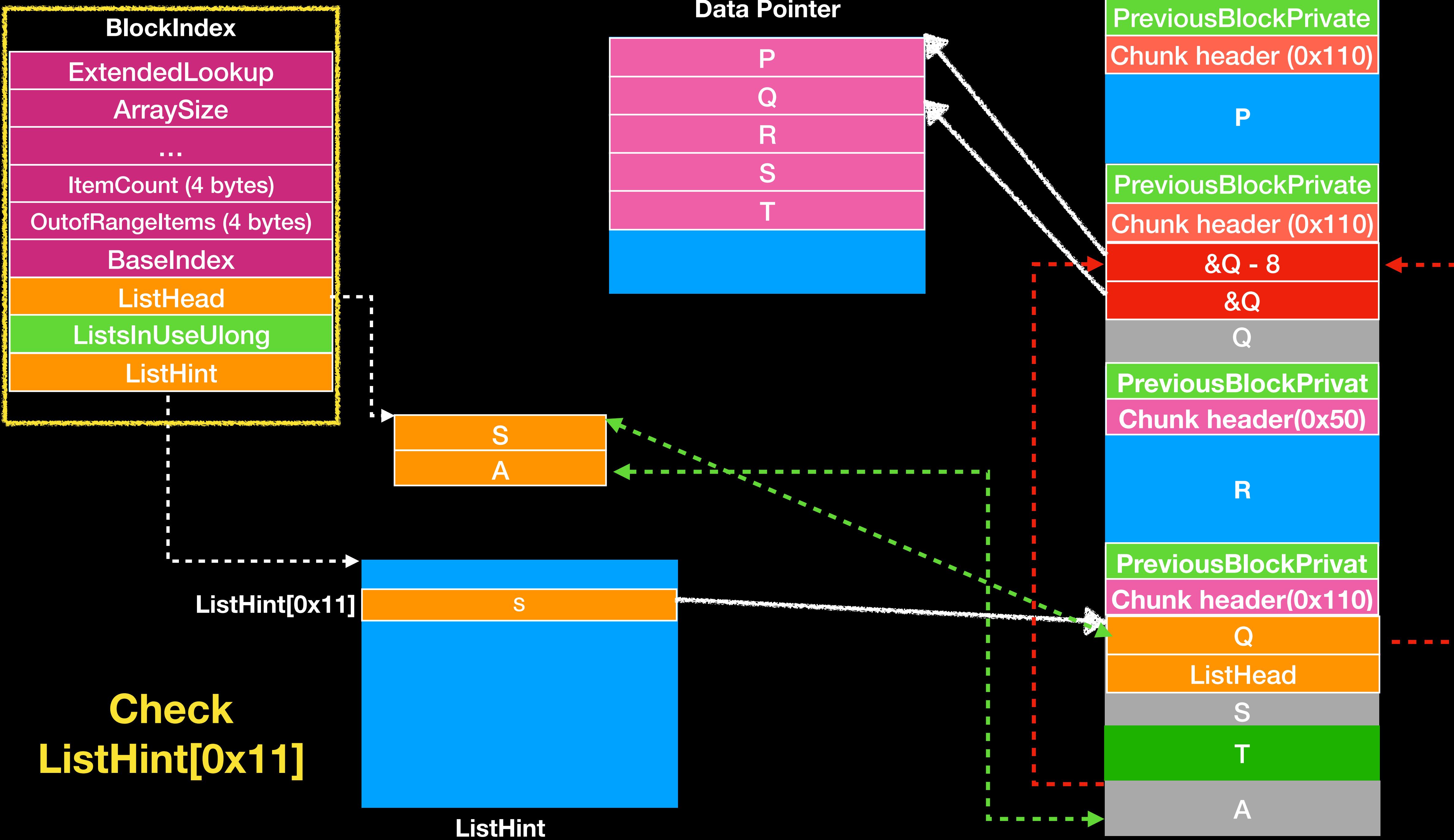


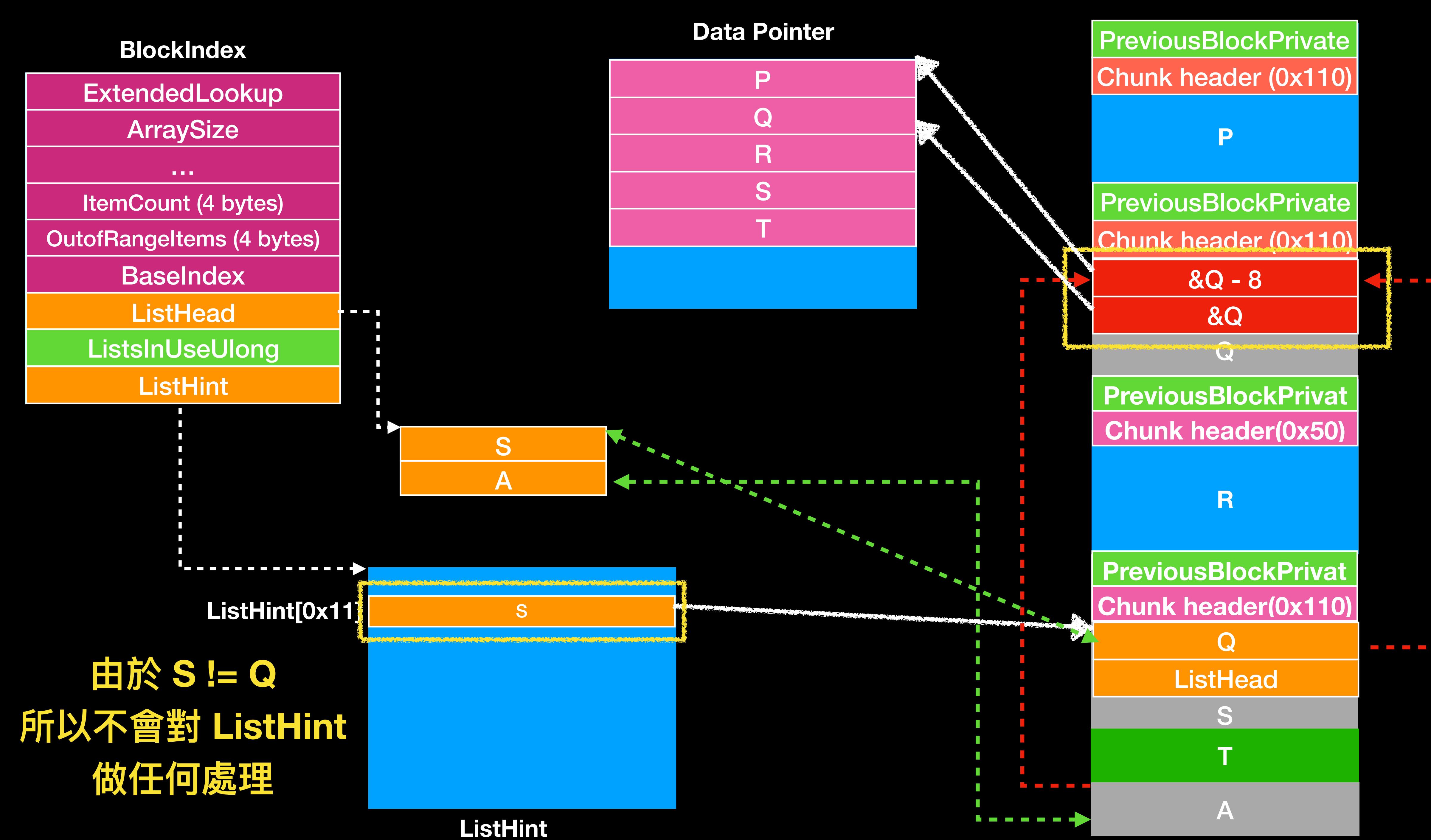


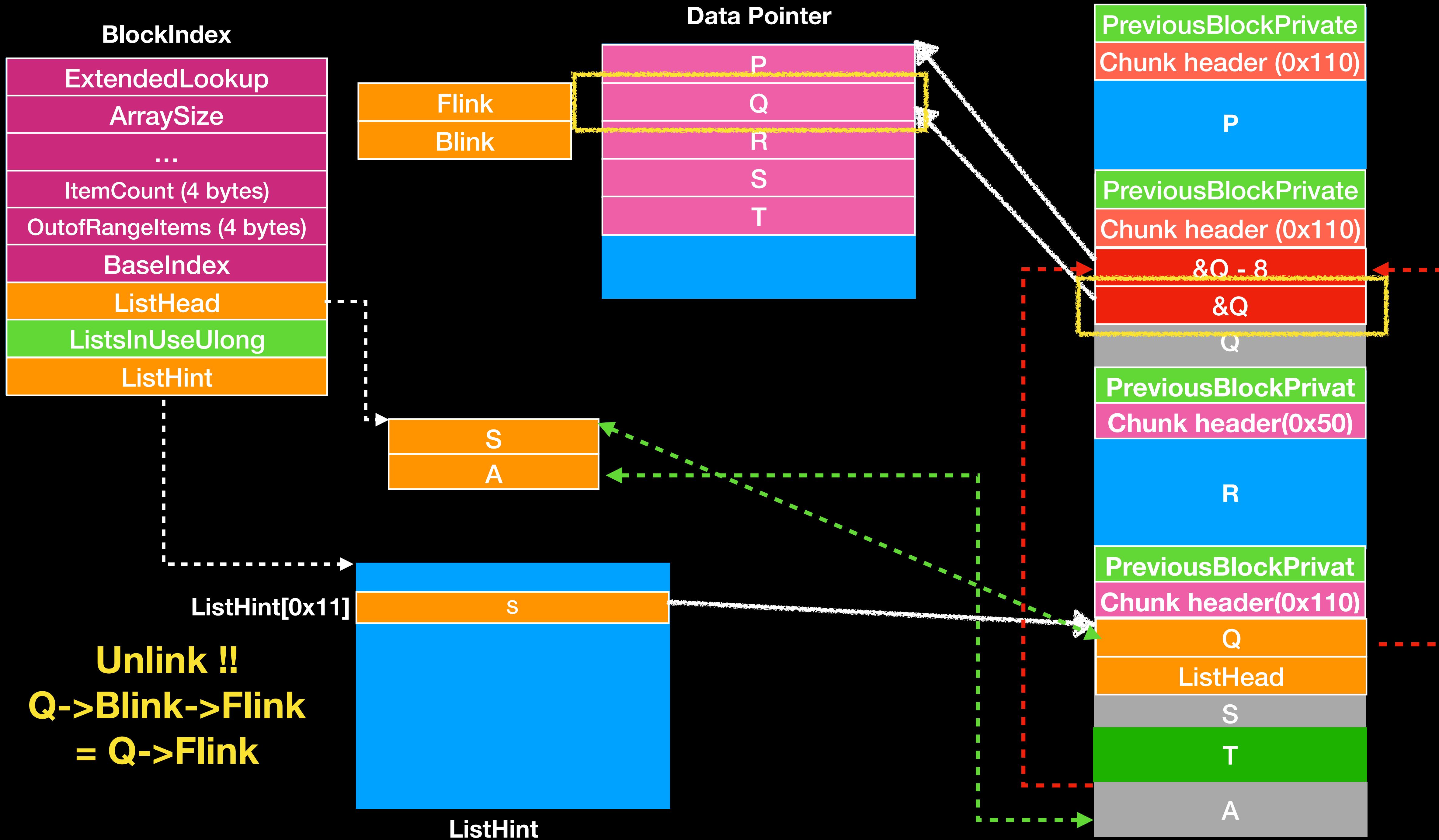


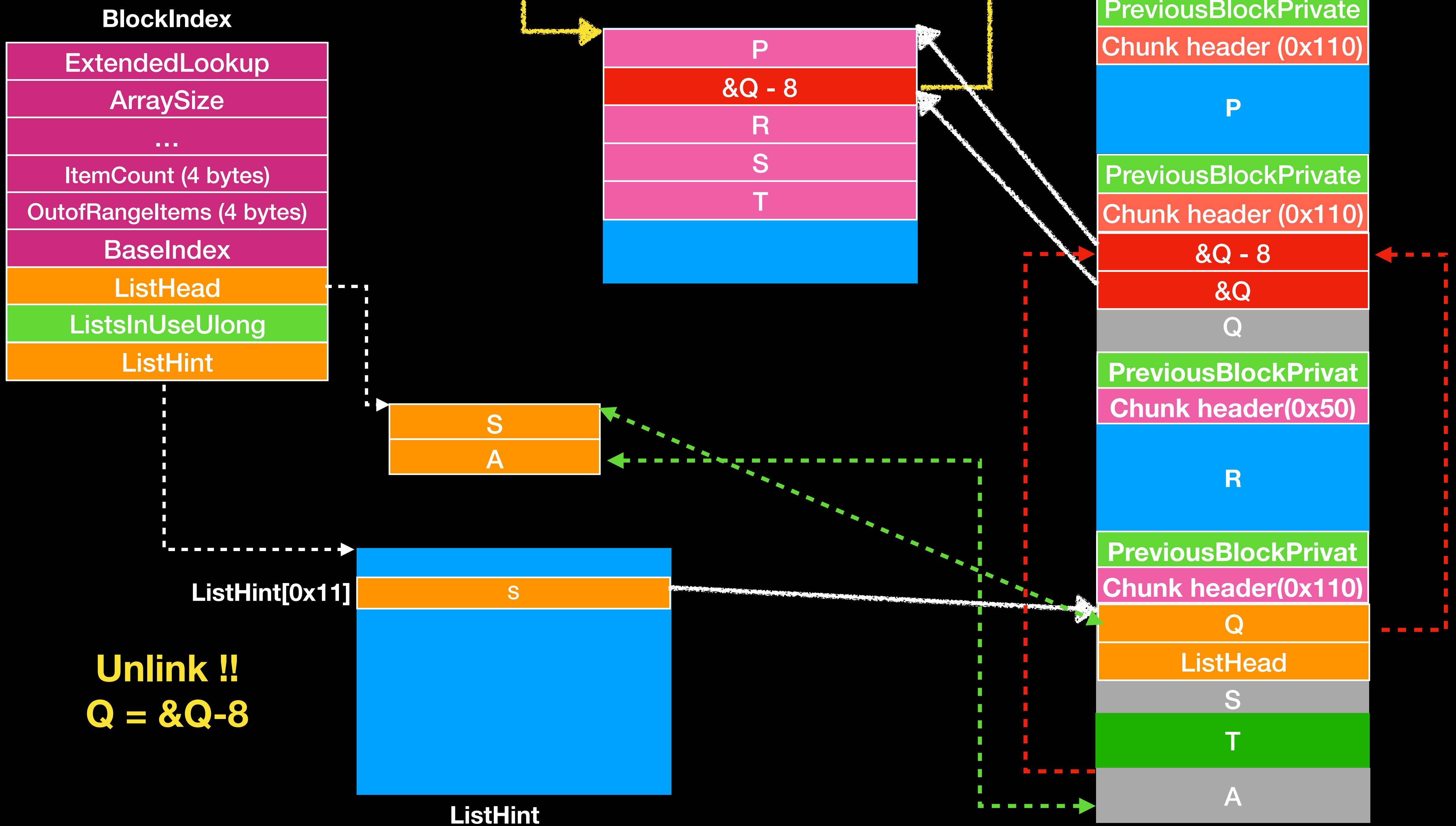


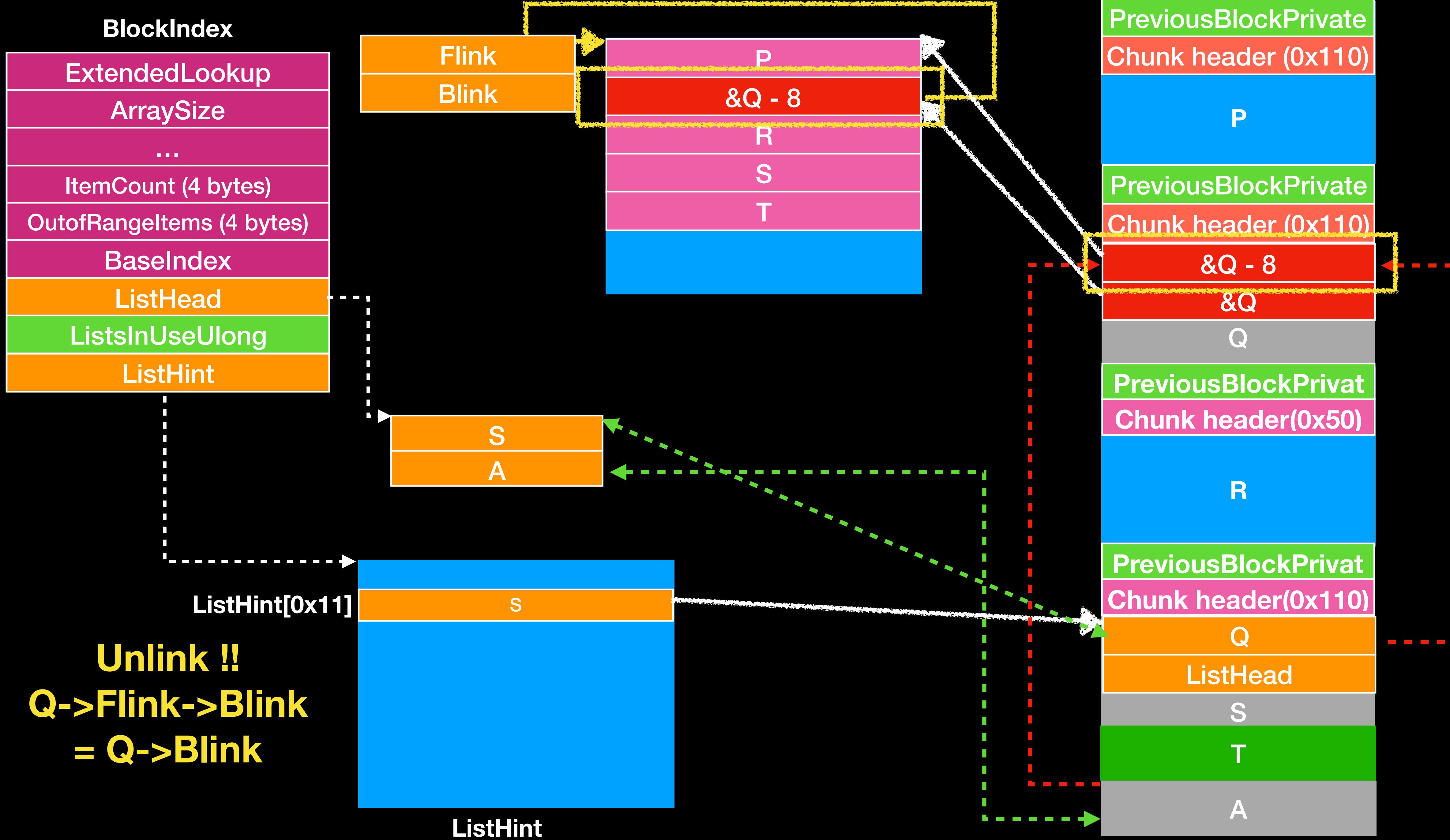


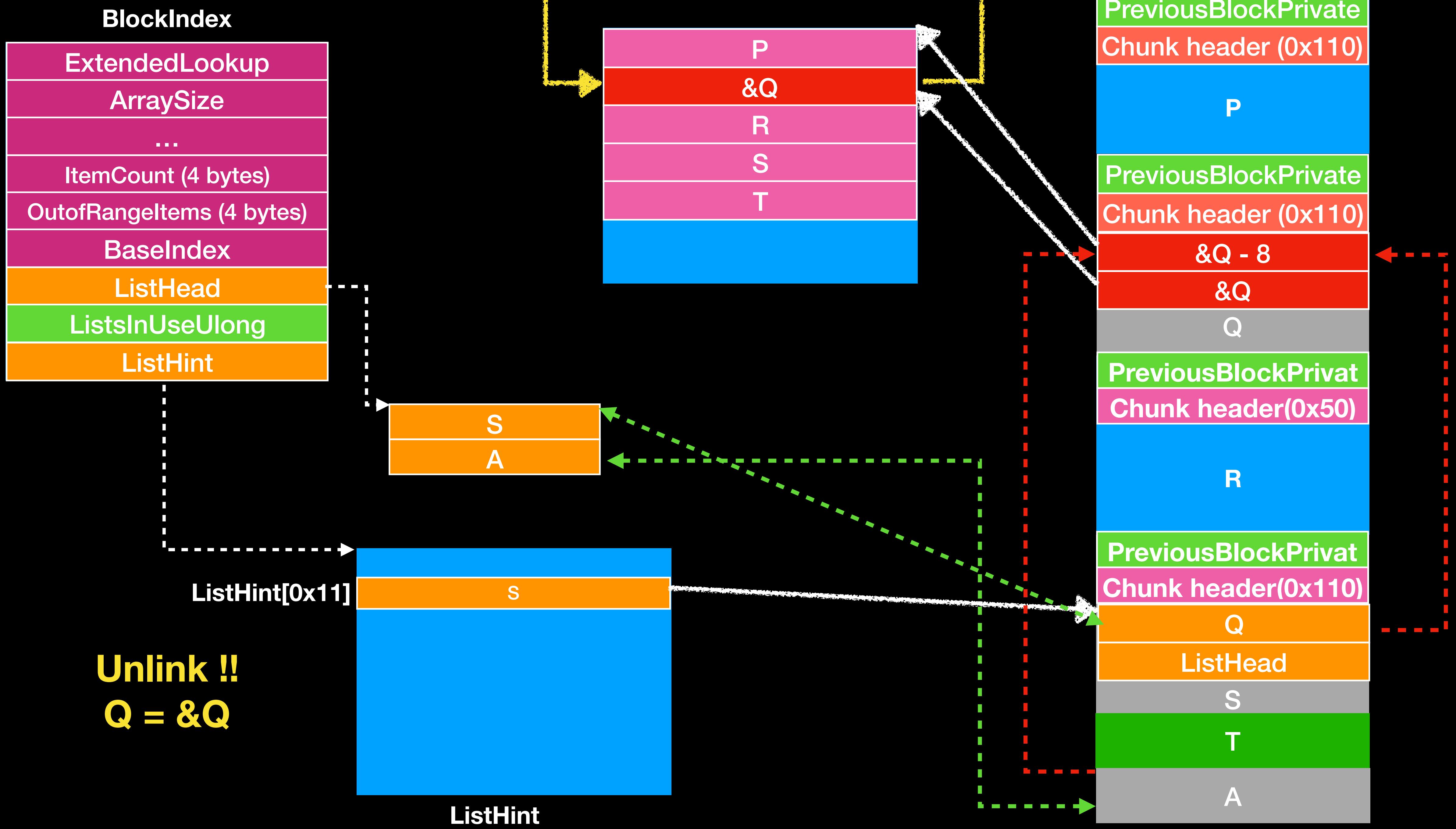


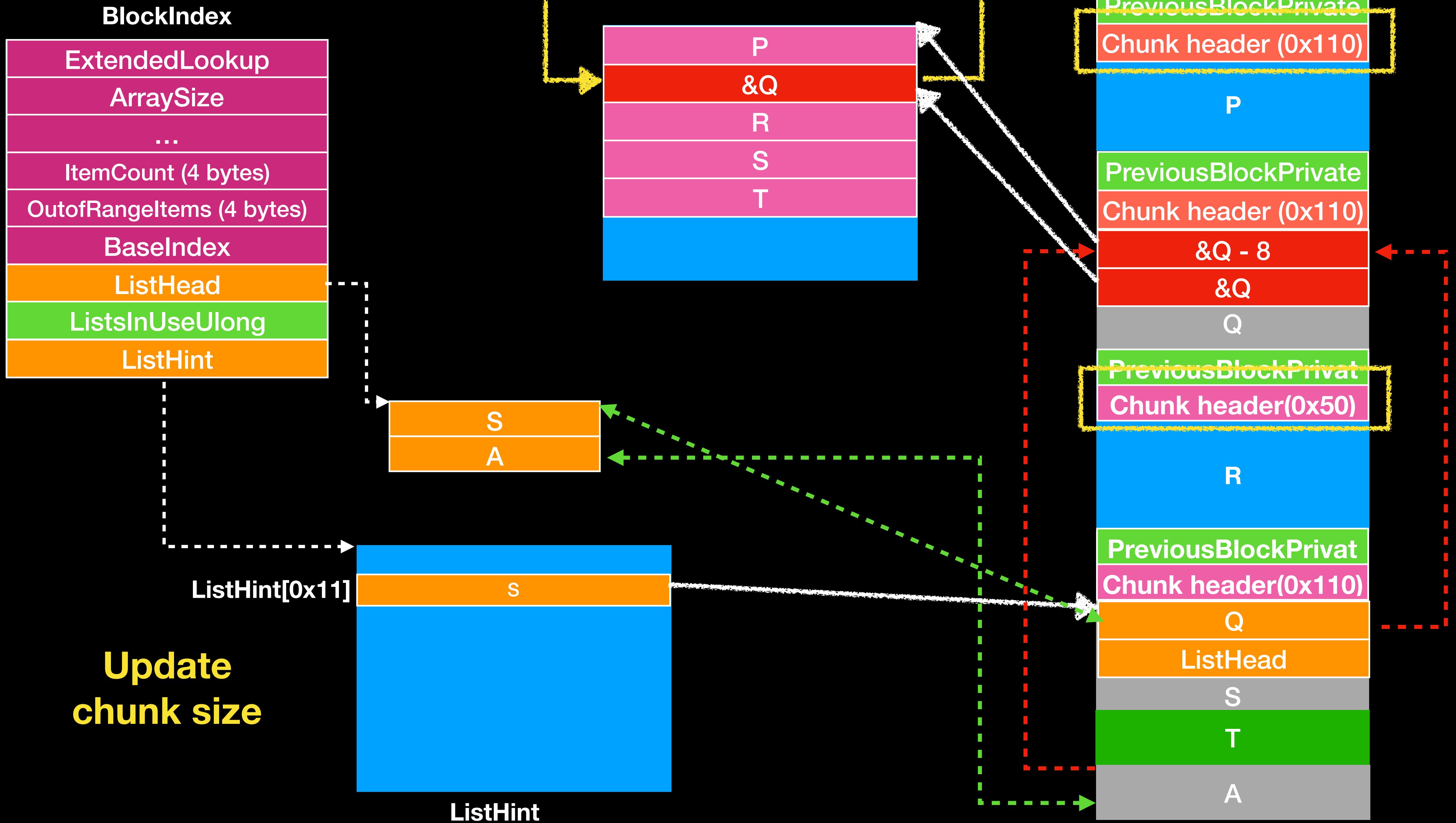




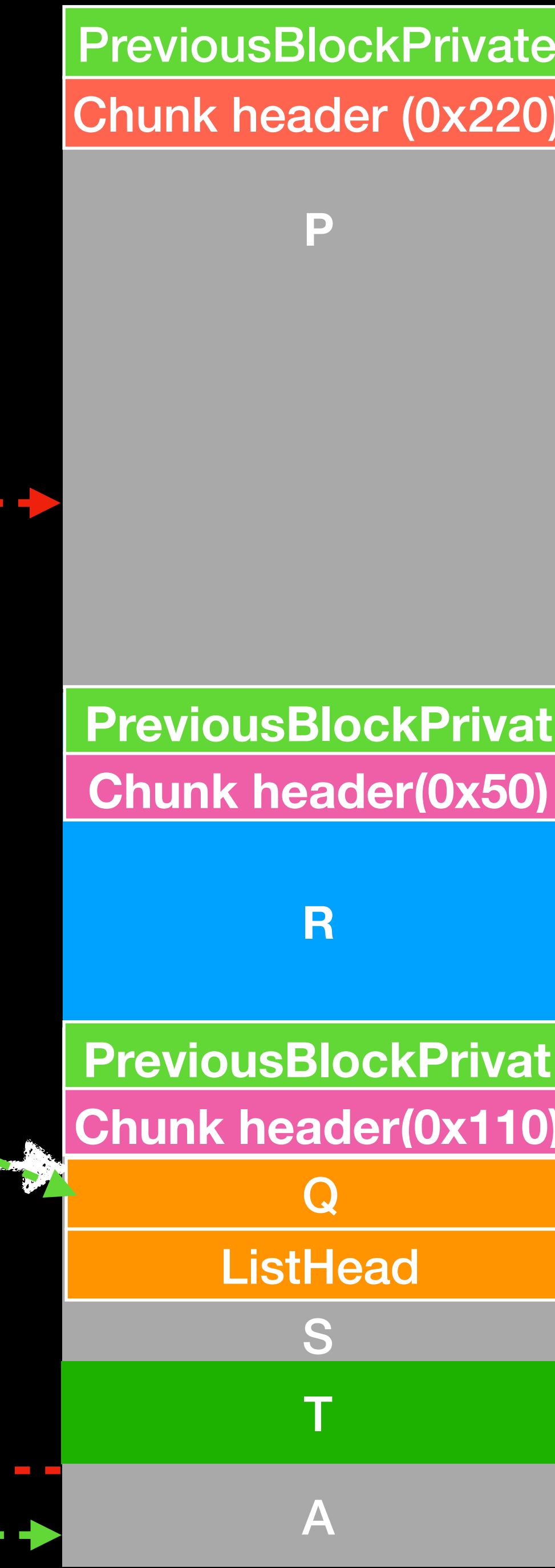






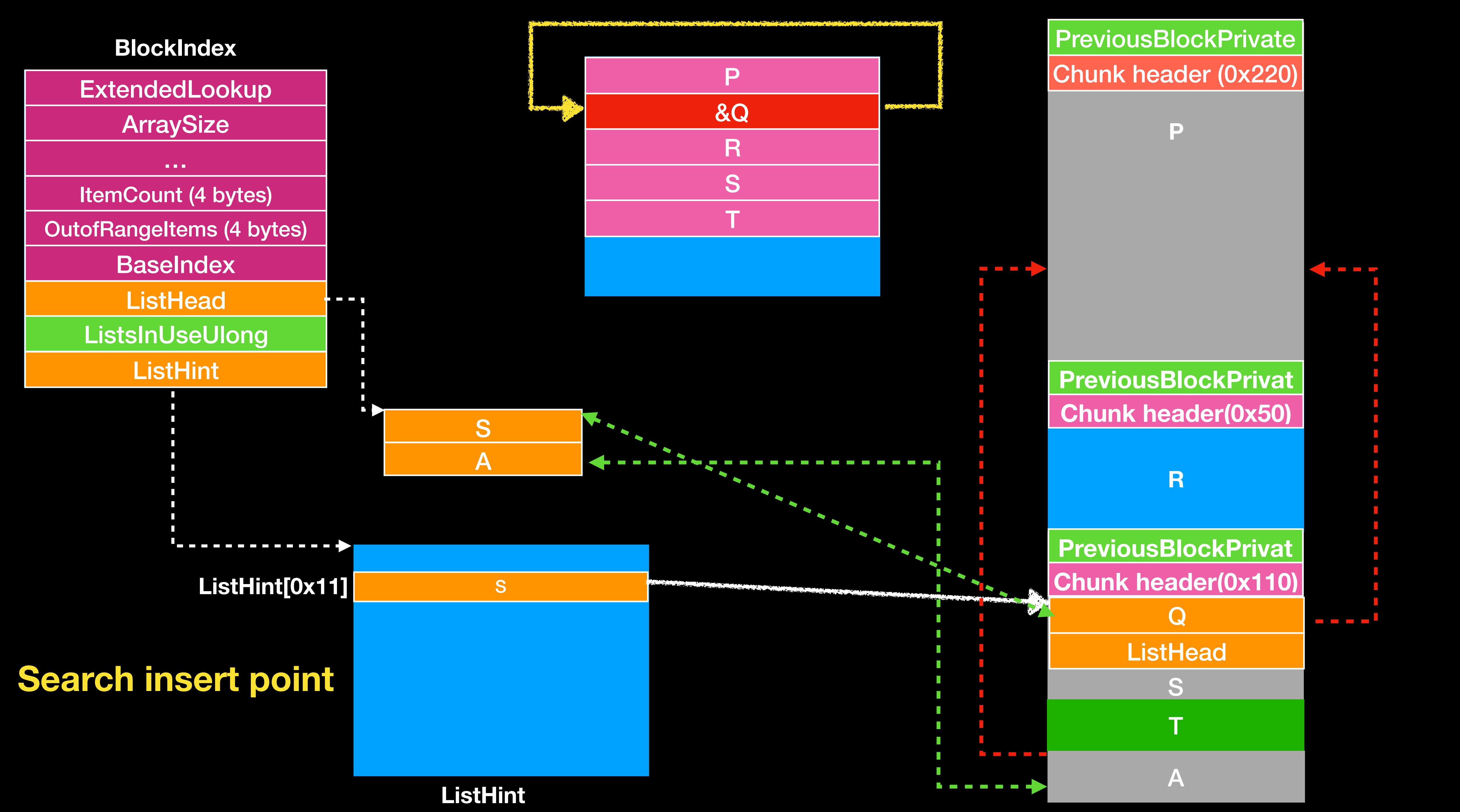


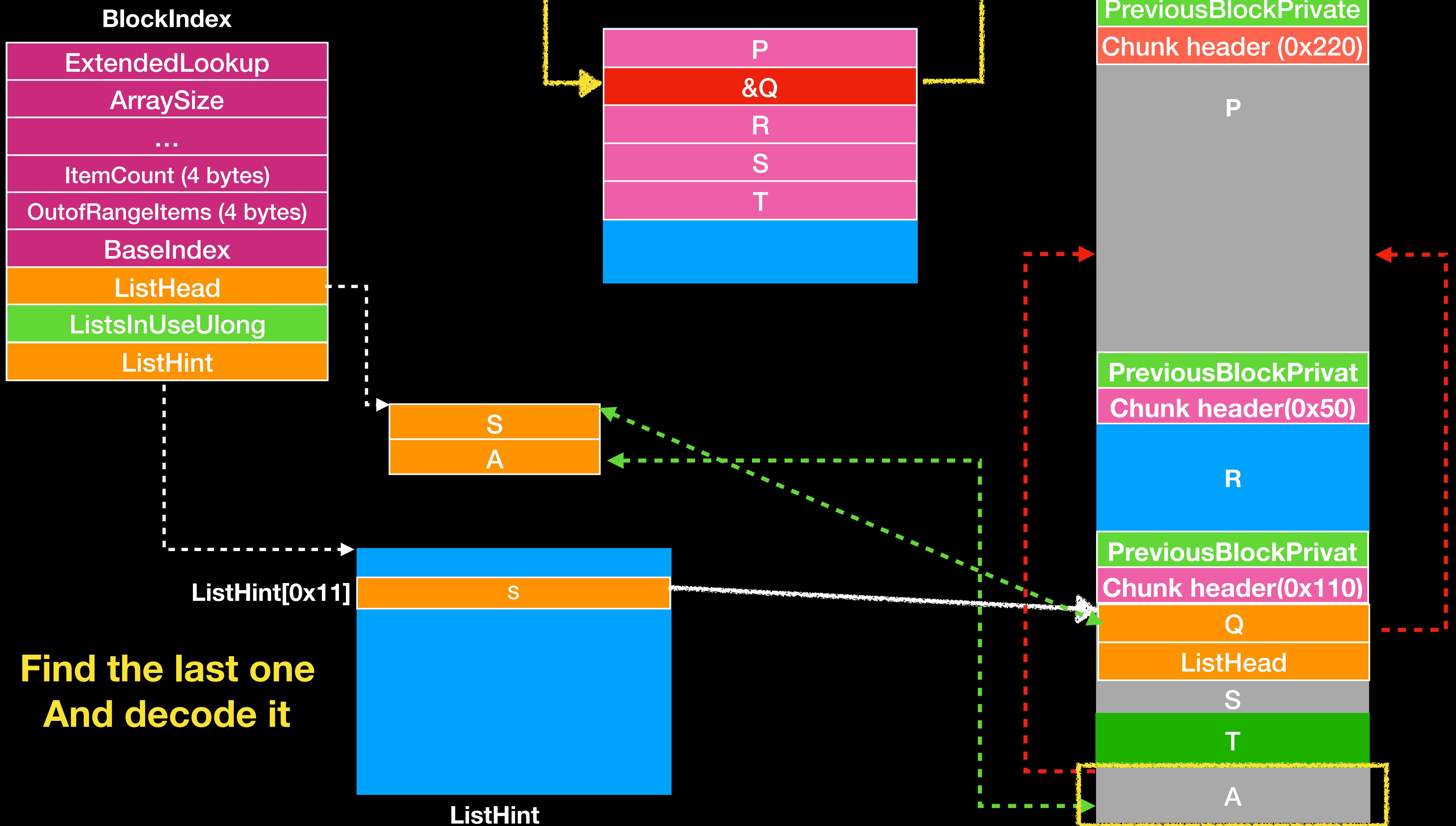
BlockIndex
ExtendedLookup
ArraySize
...
ItemCount (4 bytes)
OutofRangeItems (4 bytes)
BaseIndex
ListHead
ListsInUseUlong
ListHint

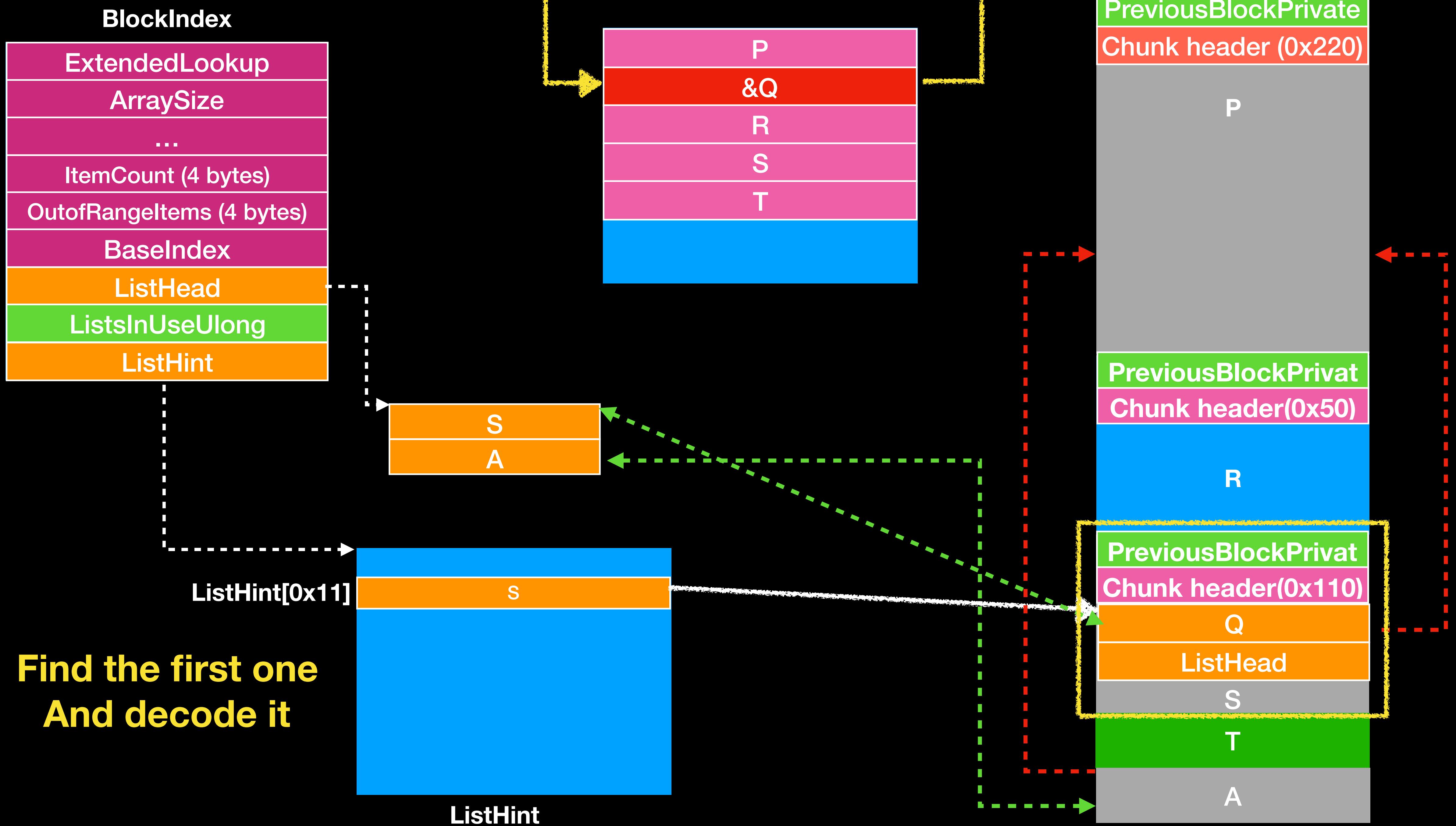


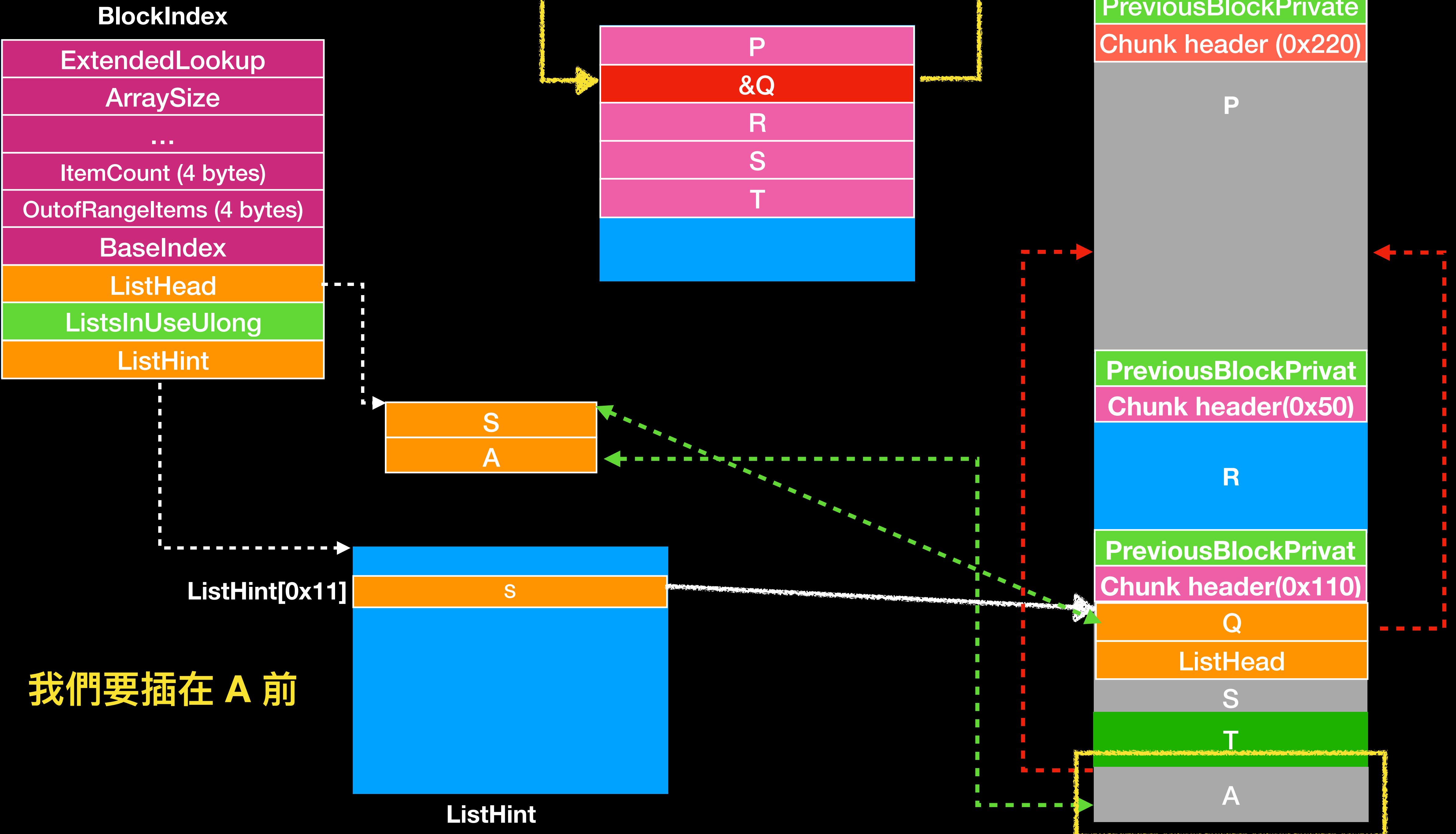
**Update  
chunk size**

`ListHint`

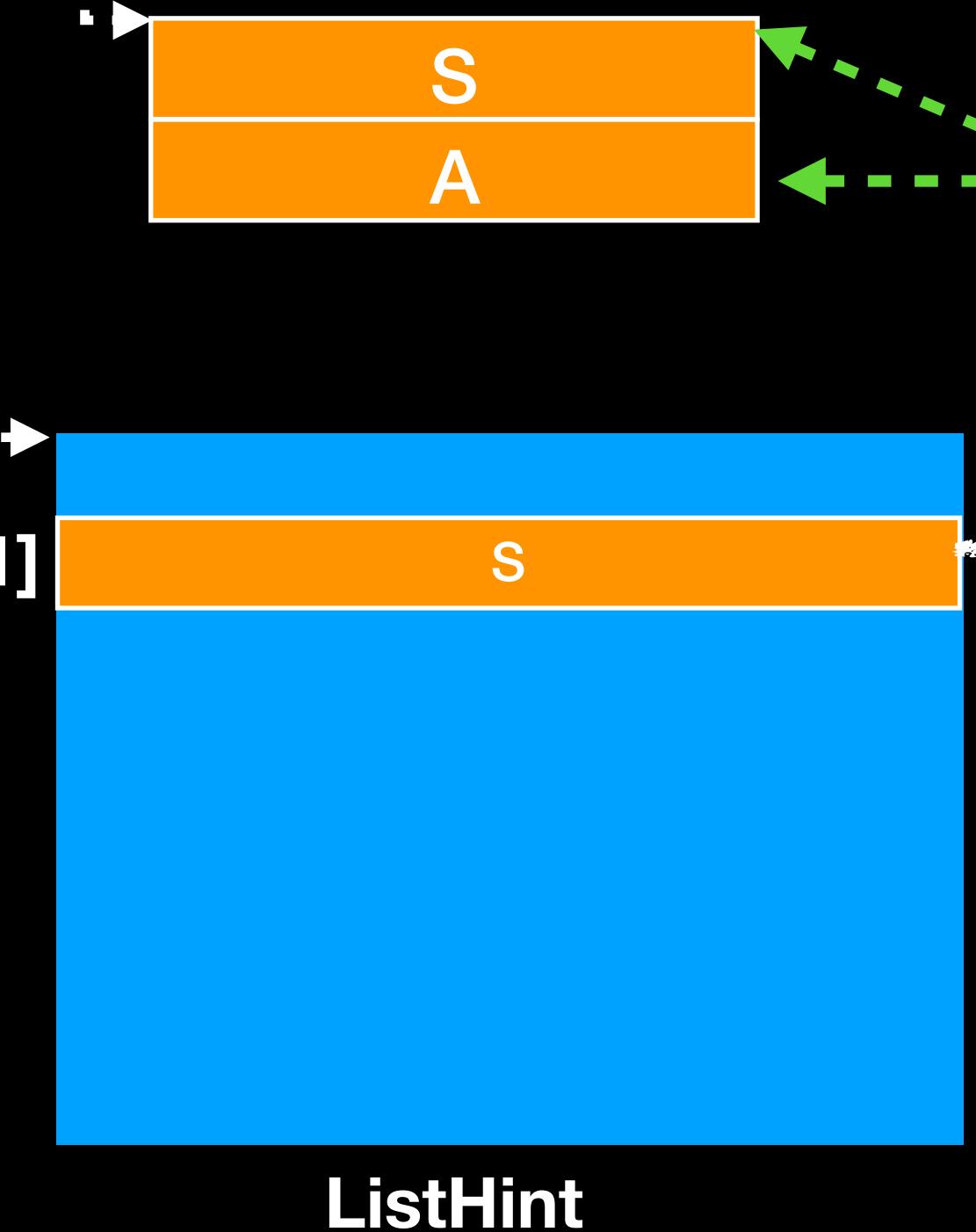
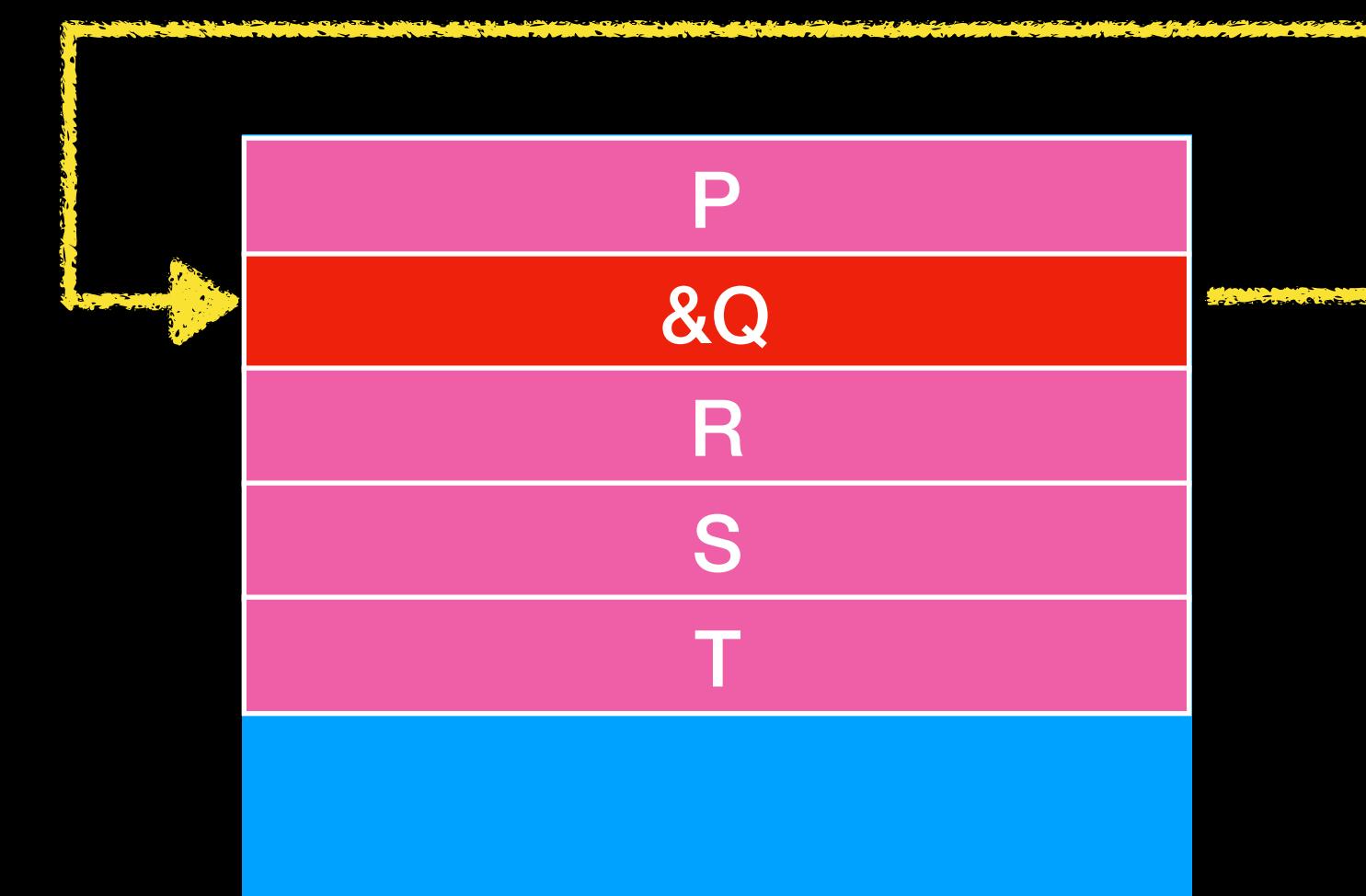




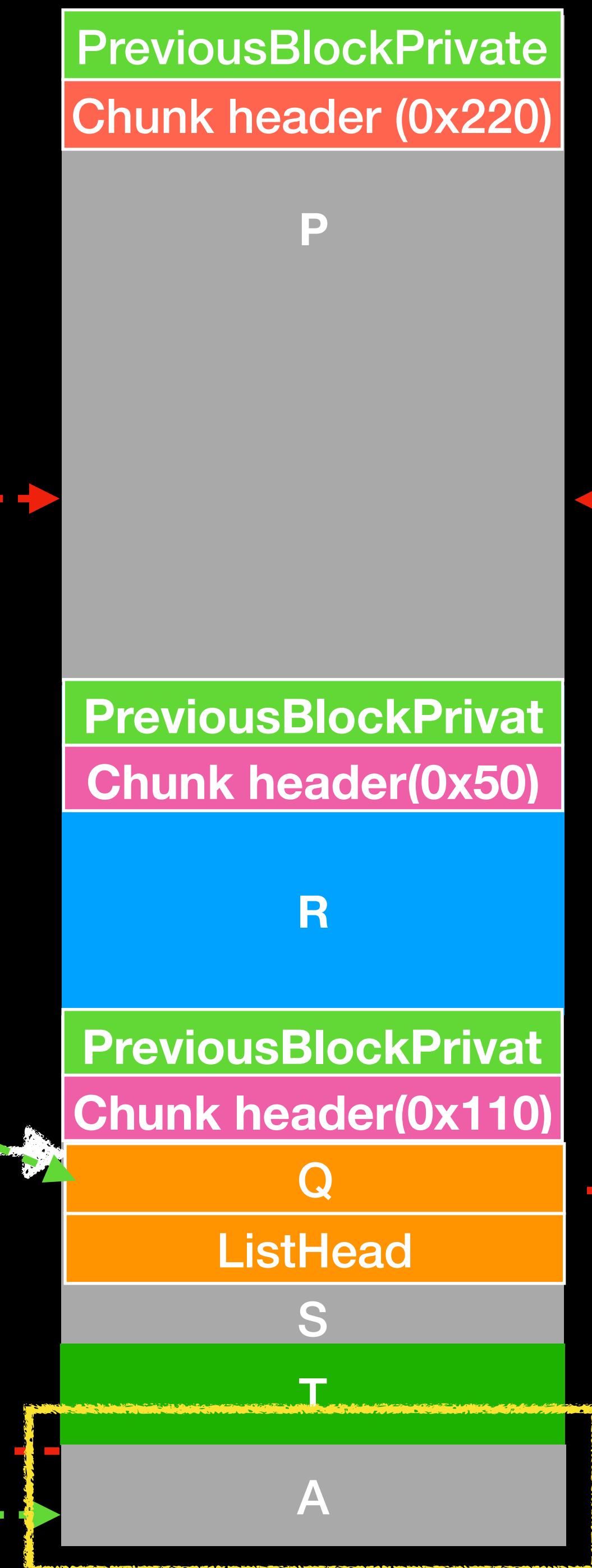




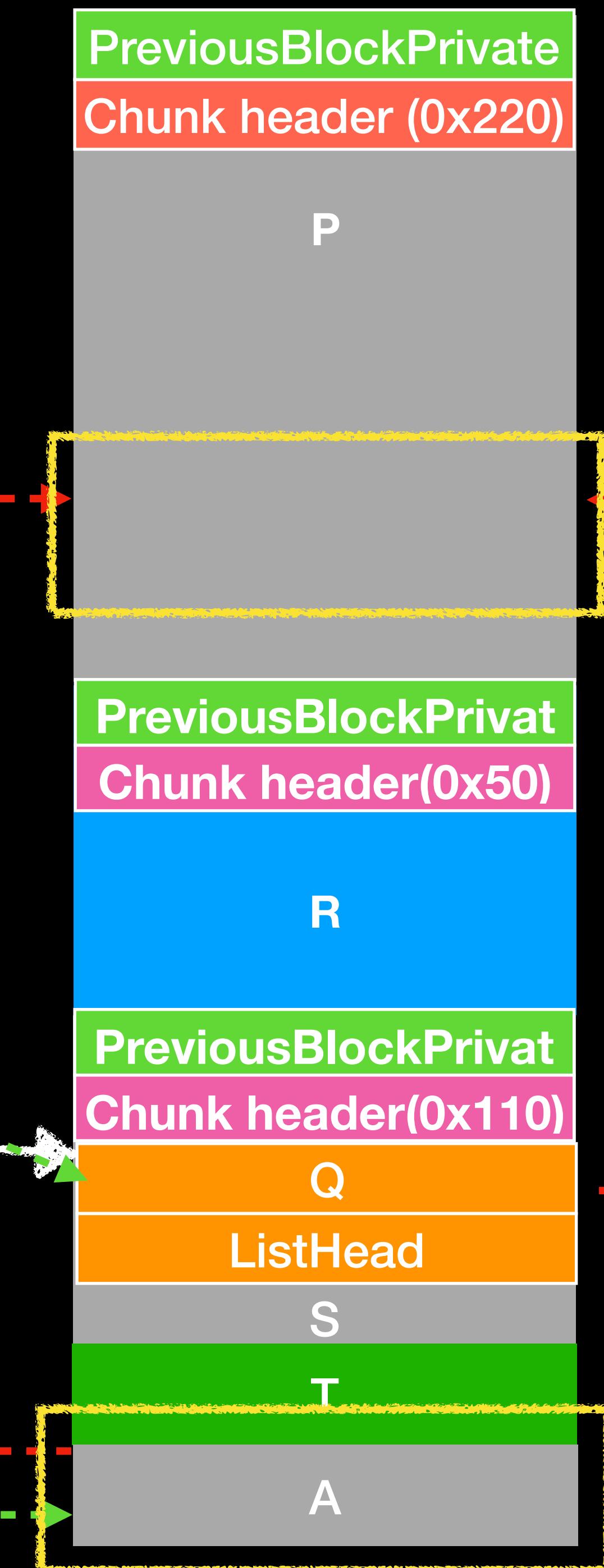
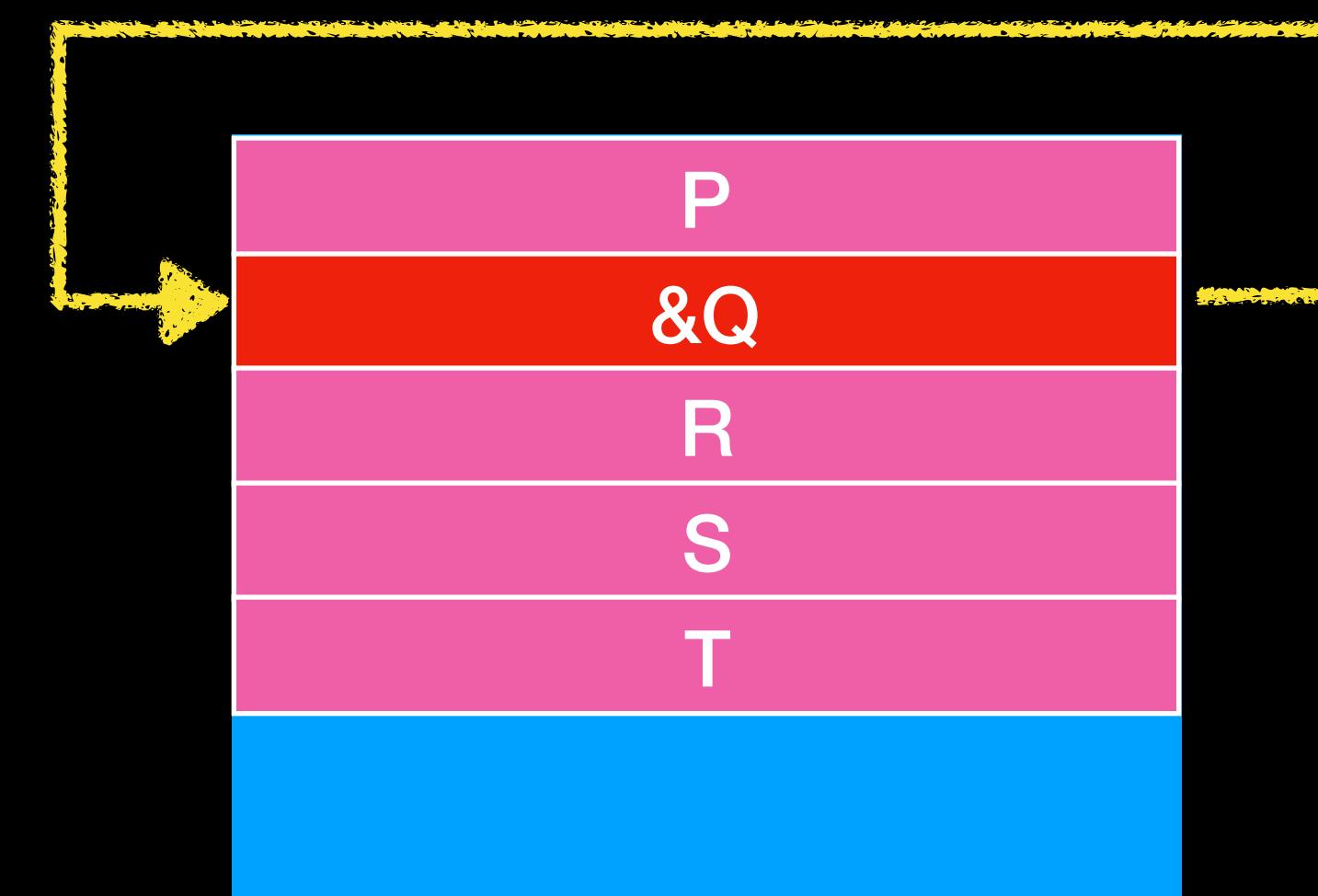
BlockIndex
ExtendedLookup
ArraySize
...
ItemCount (4 bytes)
OutofRangeItems (4 bytes)
BaseIndex
ListHead
ListsInUseUlong
ListHint



**Check**  
**A->Blink->Flink**  
 $\equiv A$

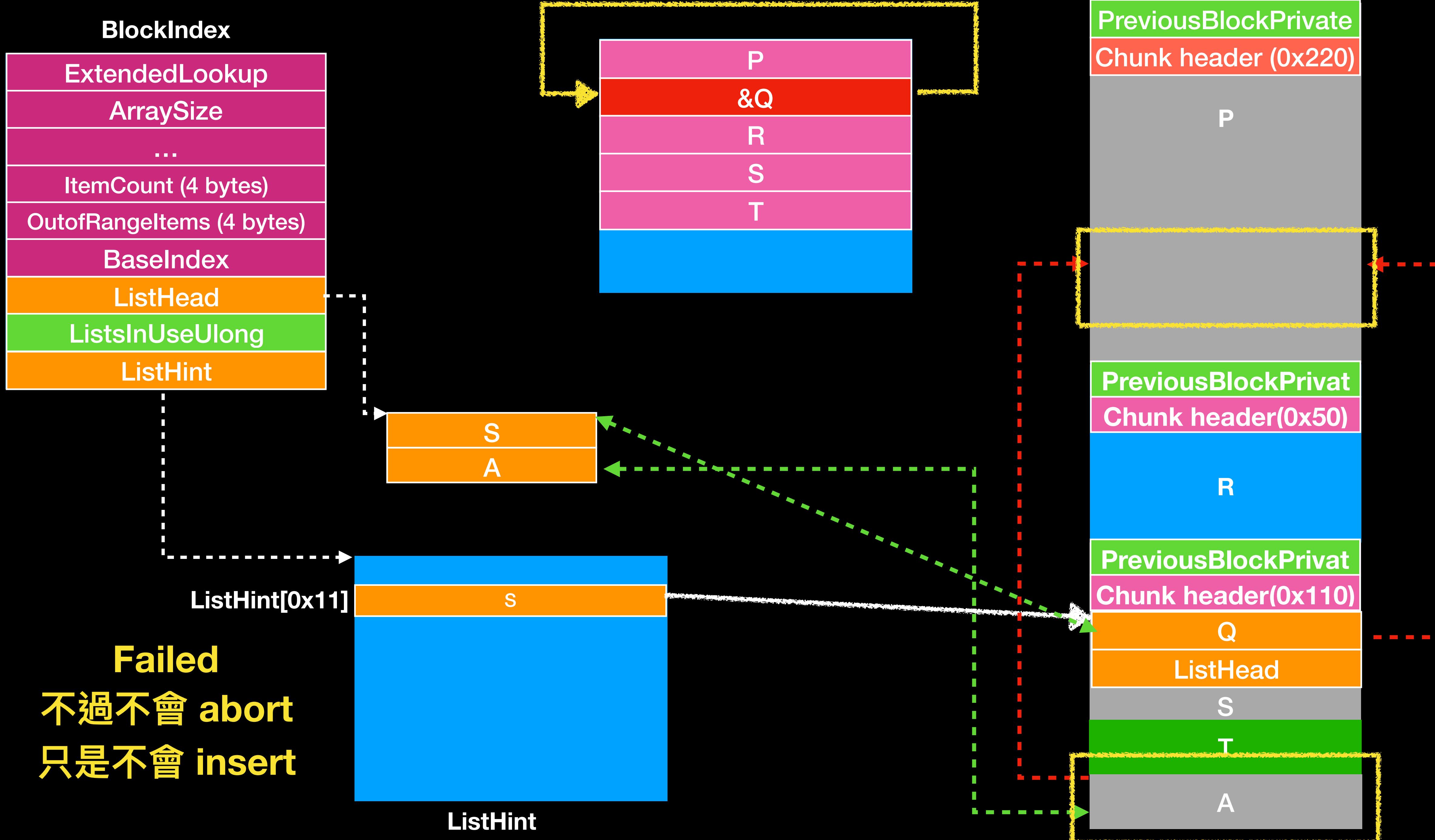


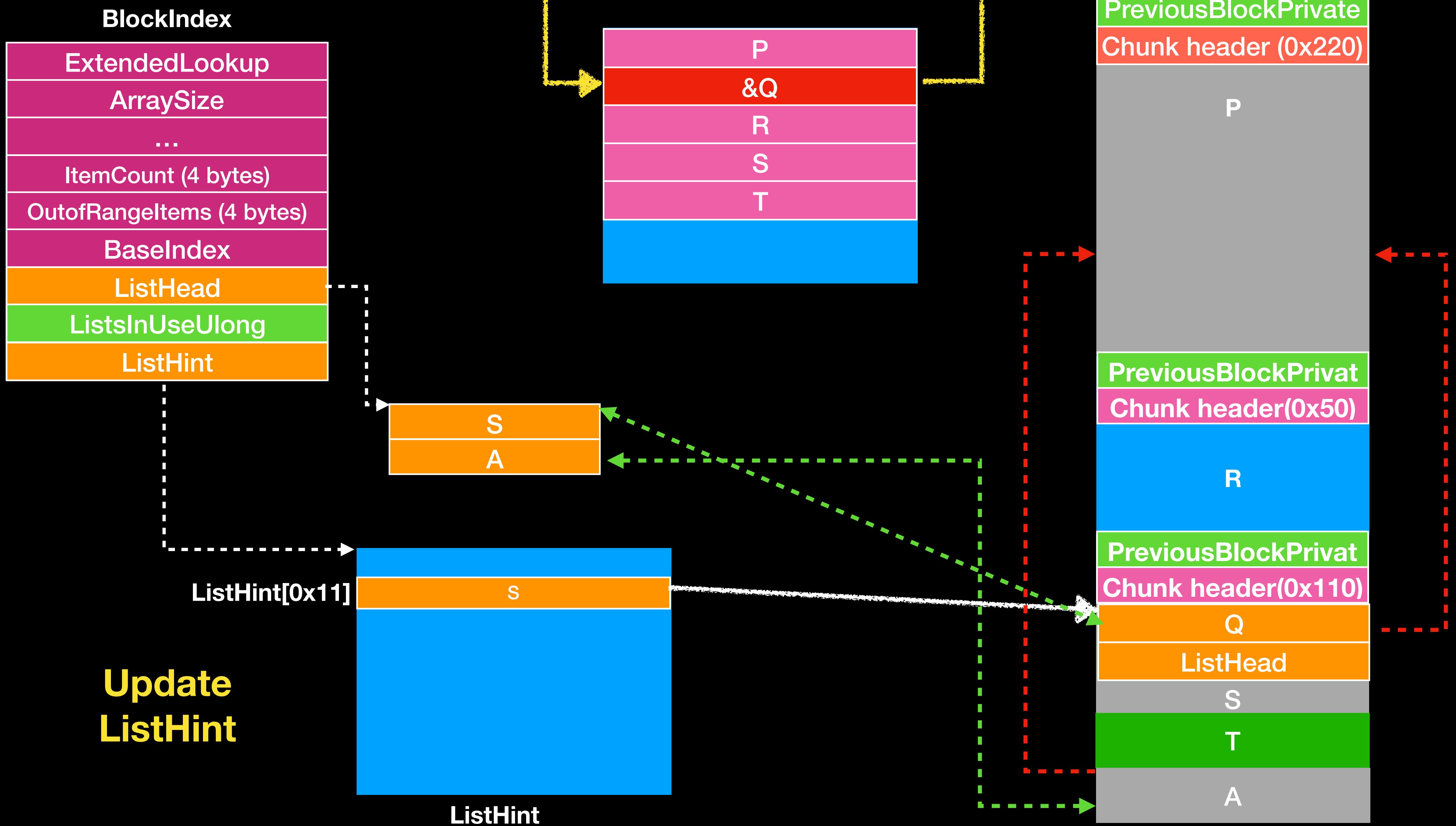
BlockIndex
ExtendedLookup
ArraySize
...
ItemCount (4 bytes)
OutofRangeItems (4 bytes)
BaseIndex
ListHead
ListsInUseUlong
ListHint

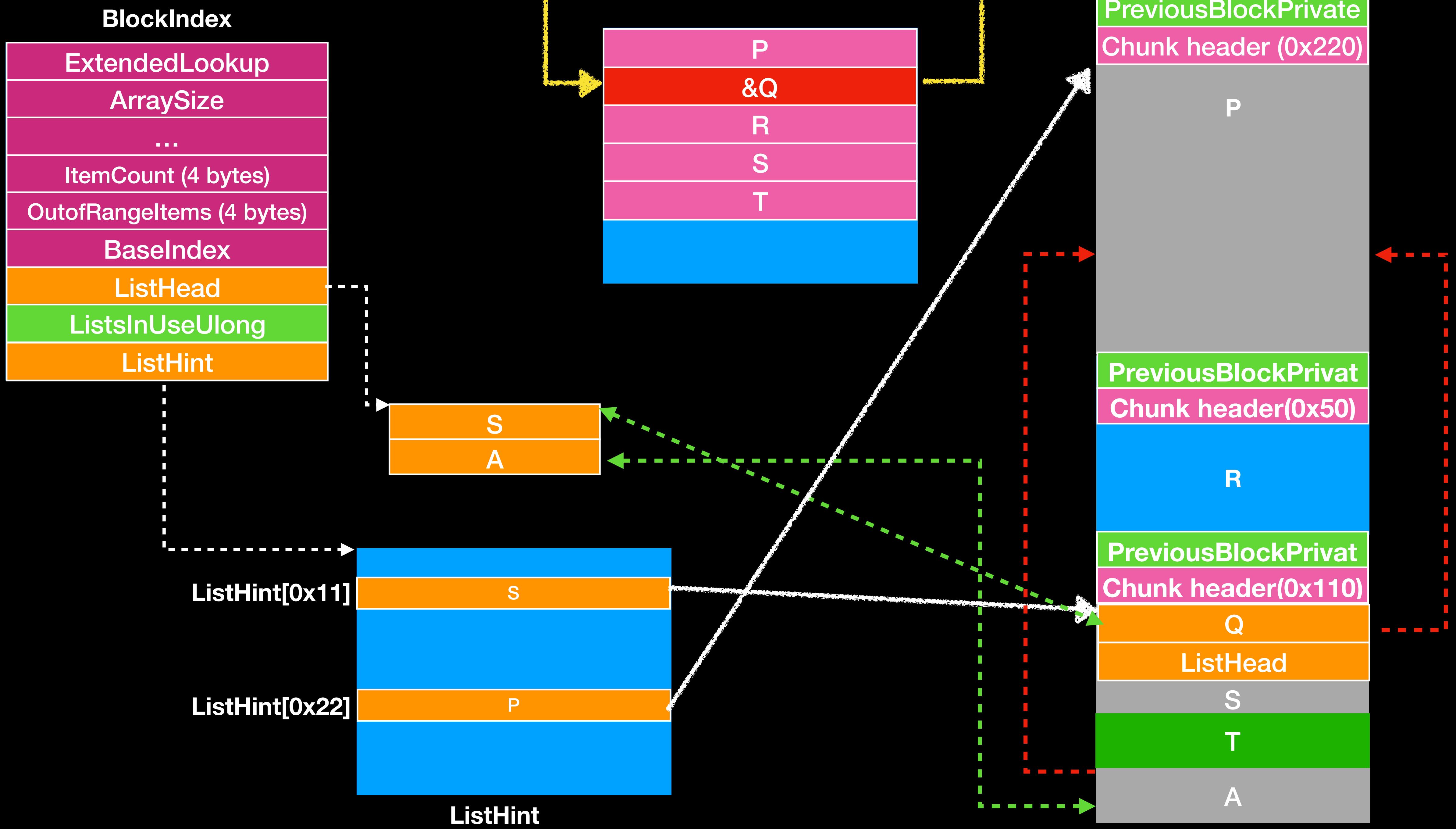


**Check  
Q->Flink  
== A**

ListHint







# Back-End Exploitation

- Edit Q
- 我們可以把 Q 周圍的 pointer 寫掉
- 最後可達成任意記憶體讀寫



# Back-End Exploitation

- Edit Q
- 我們可以把 Q 周圍的 pointer 寫掉



- 最後可達成任意記憶體讀寫
- 可對 Target 1 - 4 做寫入

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
    - ntdll.dll
    - PebLdr
    - binary address
  - Kernel32.dll

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
    - kernelbase
    - KERNELBASE!BasepFilterInfo
  - Stack address

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
    - ntdll.dll
    - \_HEAP\_LOCK
      - \_HEAP->LockVariable.Lock
    - CriticalSection->DebugInfo
    - 指向 ntdll

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
  - ntdll!PebLdr
    - \_PEB\_LDR\_DATA
    - 可找出所有 DLL 的位置
    - 不過缺點是最後一兩個 byte 可能是 0
    - 可先找出 binary base 從 IAT 找 kernel32

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
  - BinaryBase
  - 從 IAT 找出 Kernel32 位置

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
  - Kernel32
    - 重要 dll 很多常用 function 都在這 (CreateFile, ReadFile, WriteFile)
  - IAT
    - 獲得 kernelbase.dll 位置

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
    - kernelbase
    - KERNELBASE!BasepFilterInfo
      - 指向 Heap 上的結構
      - 裡面有高機率會有 stack pointer
    - <https://j00ru.vexillium.org/2016/07/disclosed-stack-data-from-the-default-heap-on-windows/>

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Leak
  - 如果 BasepFilterInfo 中沒有 stack address
    - 因 BasepFilterInfo 內容會有 stack address 是因為未初始化因素，所以有機會沒有 stack address
    - 可以從 PEB 往前或往後算 1 個 page ，通常會是 TEB 位置，上面繪有 stack 尾段的位置，可以藉由該位置找出 retn address

# Back-End Exploitation

- After arbitrary memory reading and writing
  - Write
  - Return address
  - Control RIP
  - ROP to VirtualProtect/VirtualAlloc
  - Jmp to shellcode

# Windows memory allocator

- Nt Heap
  - 後端管理器 (Back-End)
  - 前端管理器 (Front-End)

# Windows memory allocator

- Nt Heap
  - 前端管理器 (Front-End)
    - 目前 win10 主要使用
    - LowFragmentationHeap
    - 在非 Debug 下才會 enable
    - Size < 0x4000

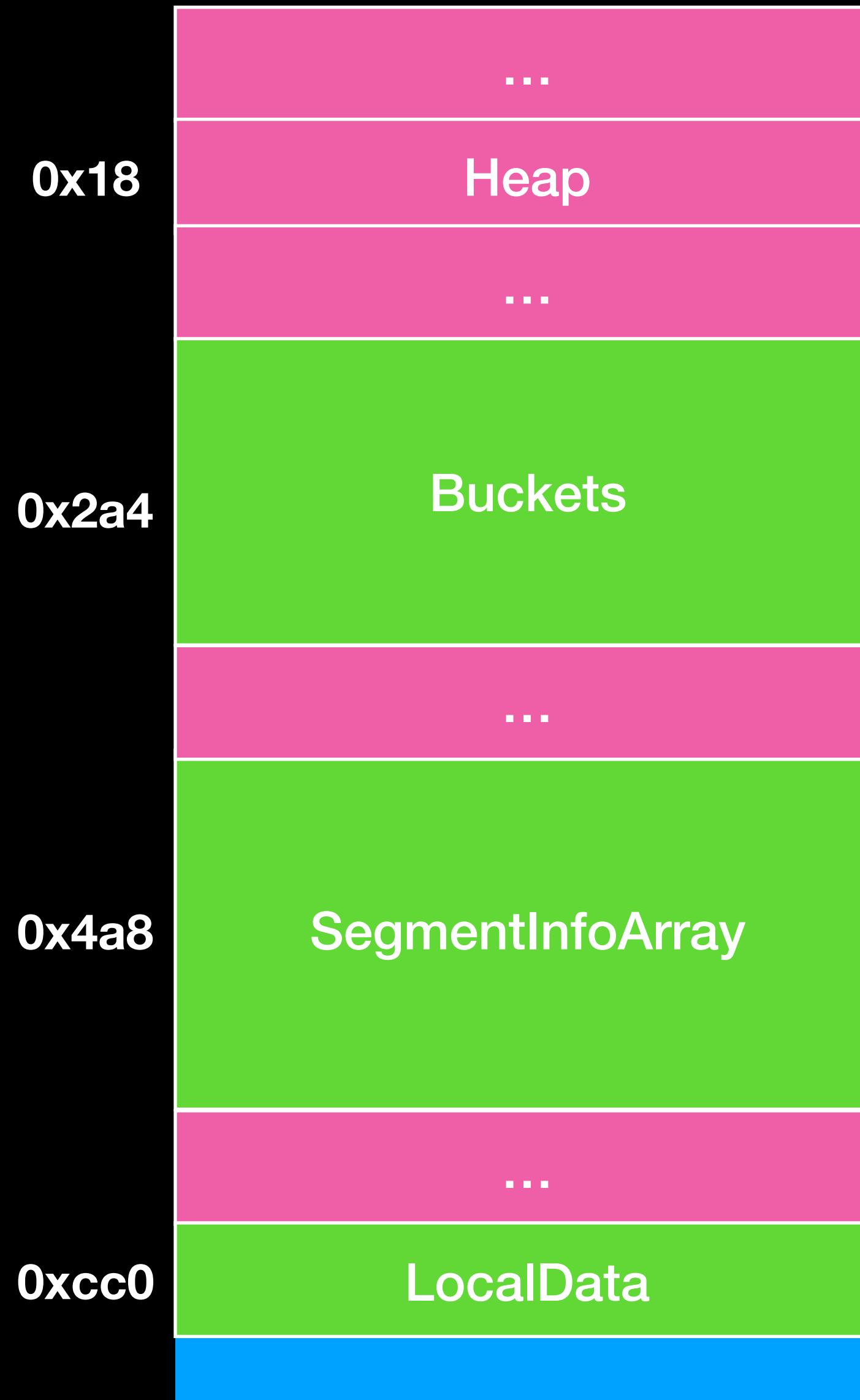
# Windows memory allocator

- Nt Heap
  - 前端管理器 (Front-End)
  - Data Structure
  - 分配機制

# Windows memory allocator

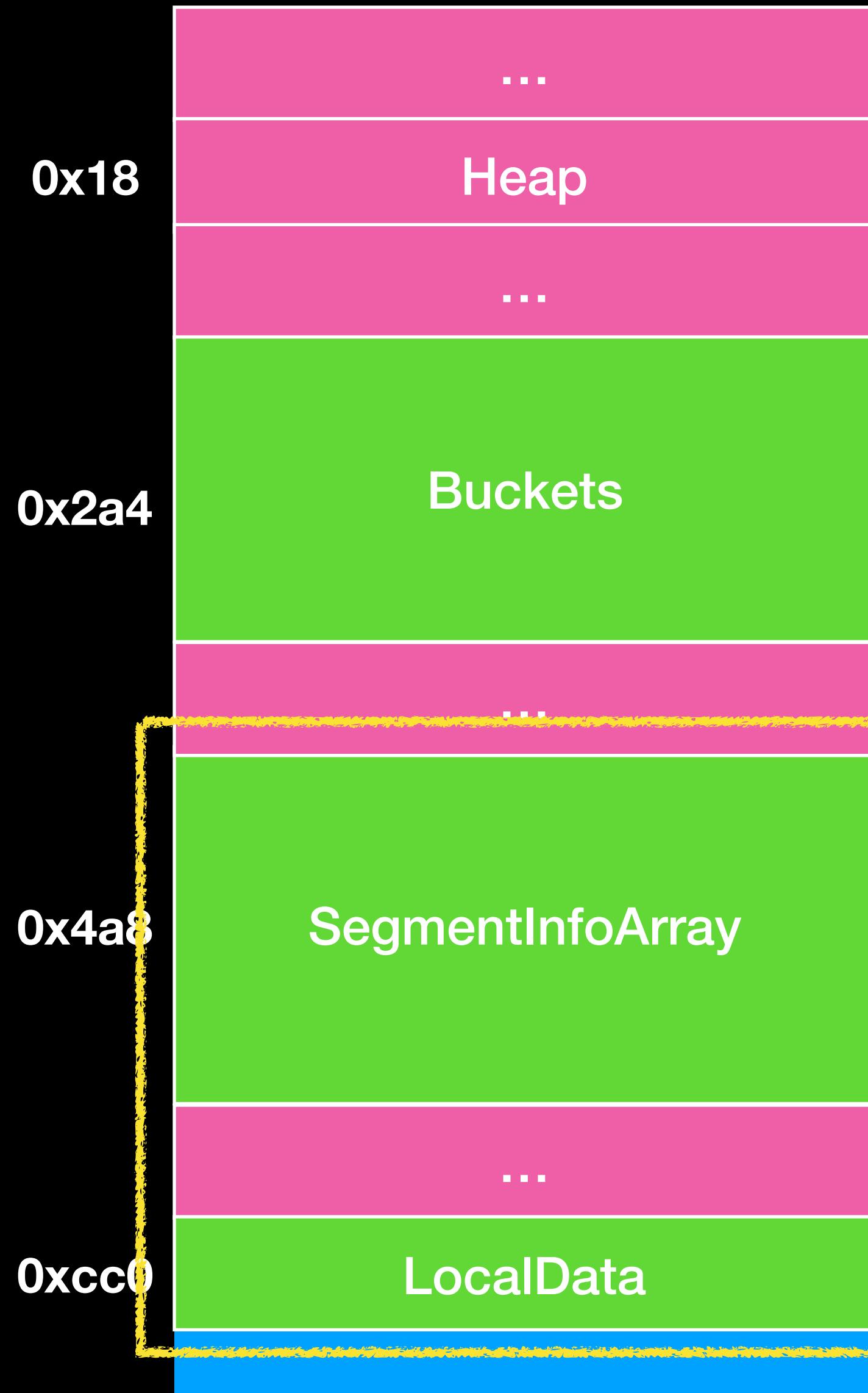
- Nt Heap
  - 前端管理器 (Front-End)
  - Data Structure
  - 分配機制

# LFH



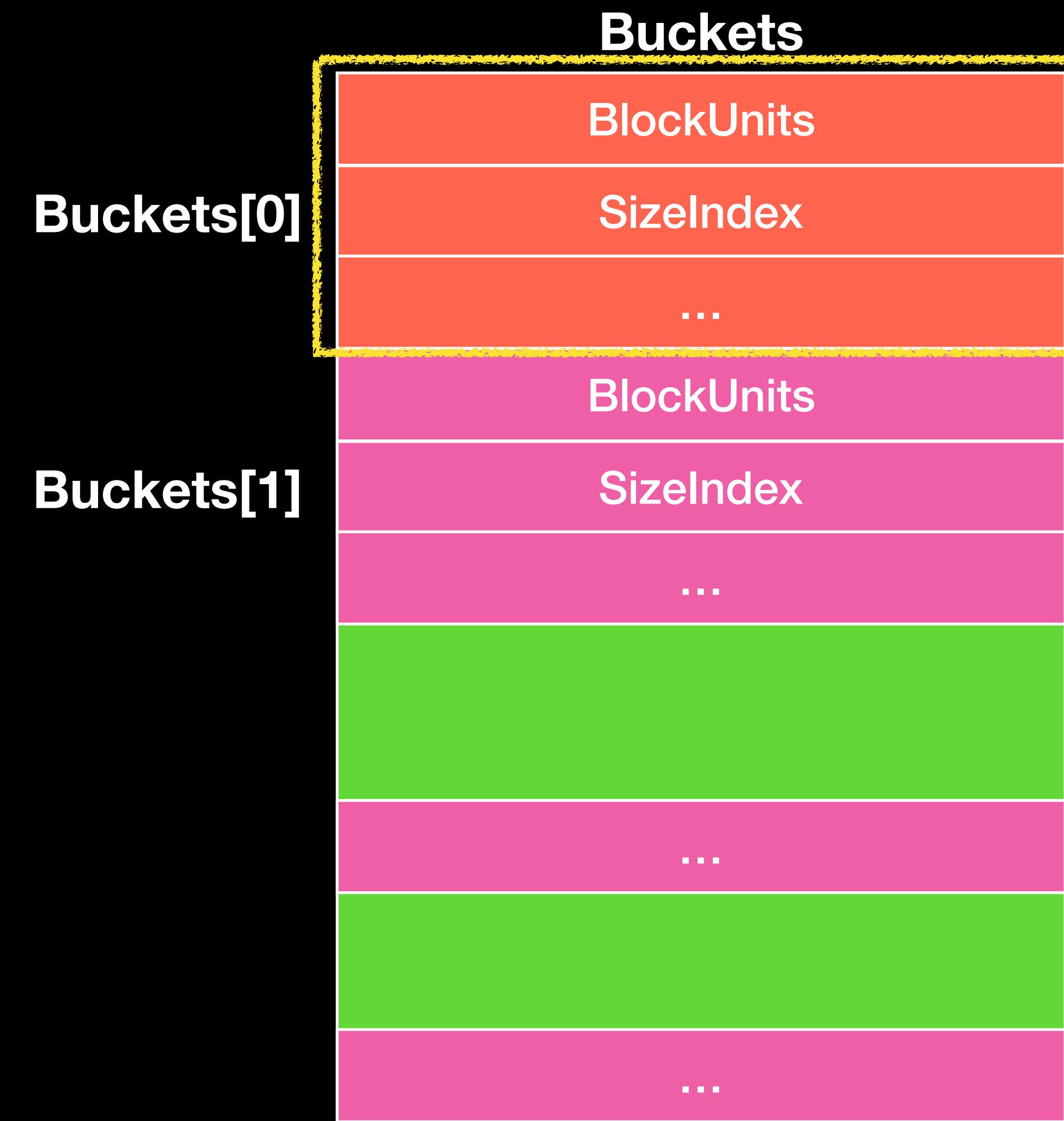
- FrontEndHeap (`_LFH_HEAP`)
  - HEAP (`_HEAP`)
  - 指向對應到的 `_HEAP`
- Buckets (`_HEAP_BUCKET`)
  - 用來尋找配置大小對應到 Block 大小的陣列結構

# LFH



- FrontEndHeap (`_LFH_HEAP`)
- SegmentInfoArray (`_HEAP_LOCAL_SEGMENT_INFO`)
  - `_HEAP_LOCAL_SEGMENT_INFO` array，不同大小對應到不同的 Segment\_info 結構，主要管理對應到的 Subsegment 的資訊
- LocalData (`_HEAP_LOCAL_DATA`)
  - 其中有個欄位指向 LFH 本身，通常用來找回 LFH

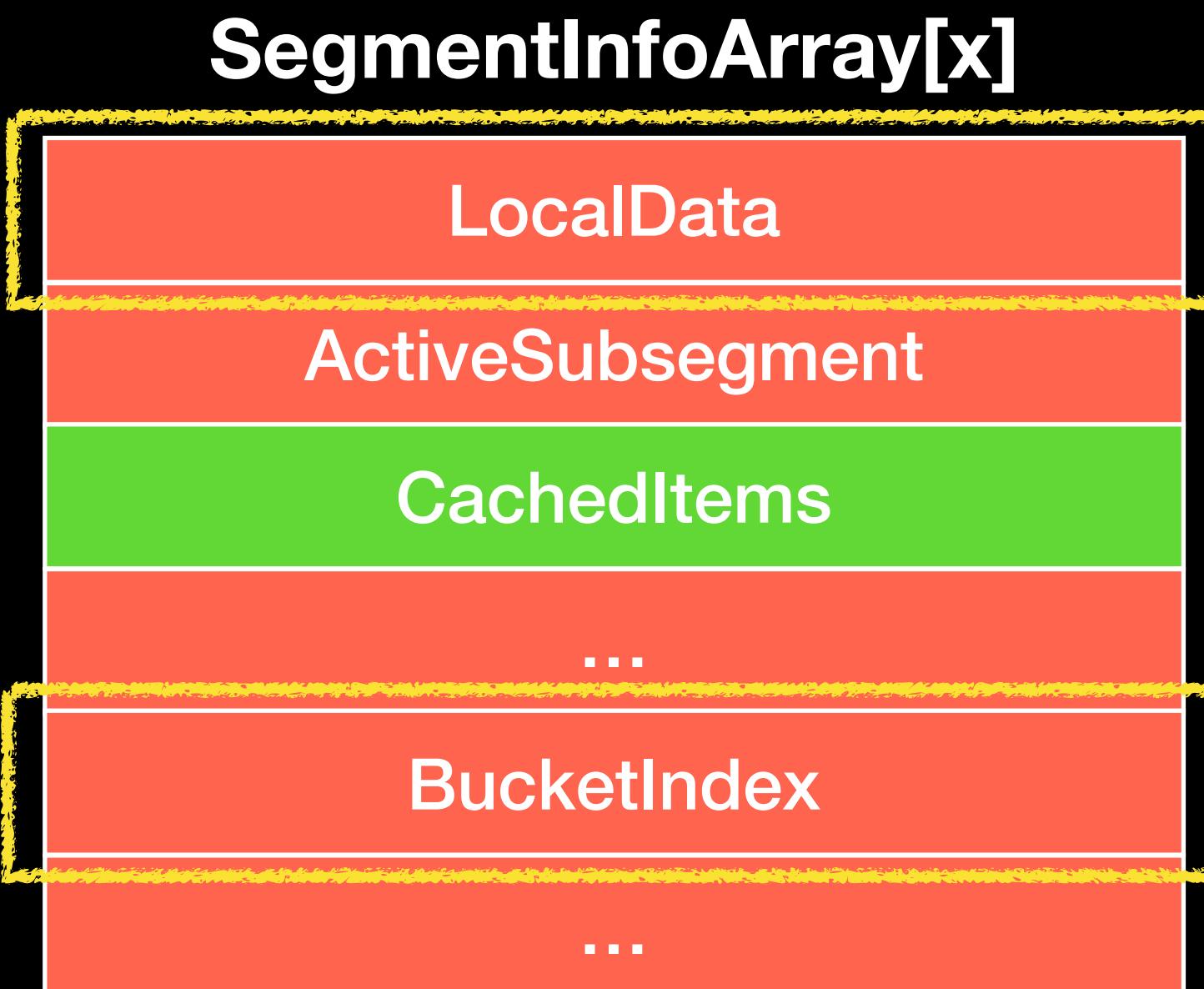
# LFH



- Buckets (`_HEAP_BUCKET`)
  - `BlockUnits`
    - 要分配出去的一個 block 大小  $>> 4$
  - `SizeIndex`
    - 使用者需要的大小  $>> 4$

# LFH

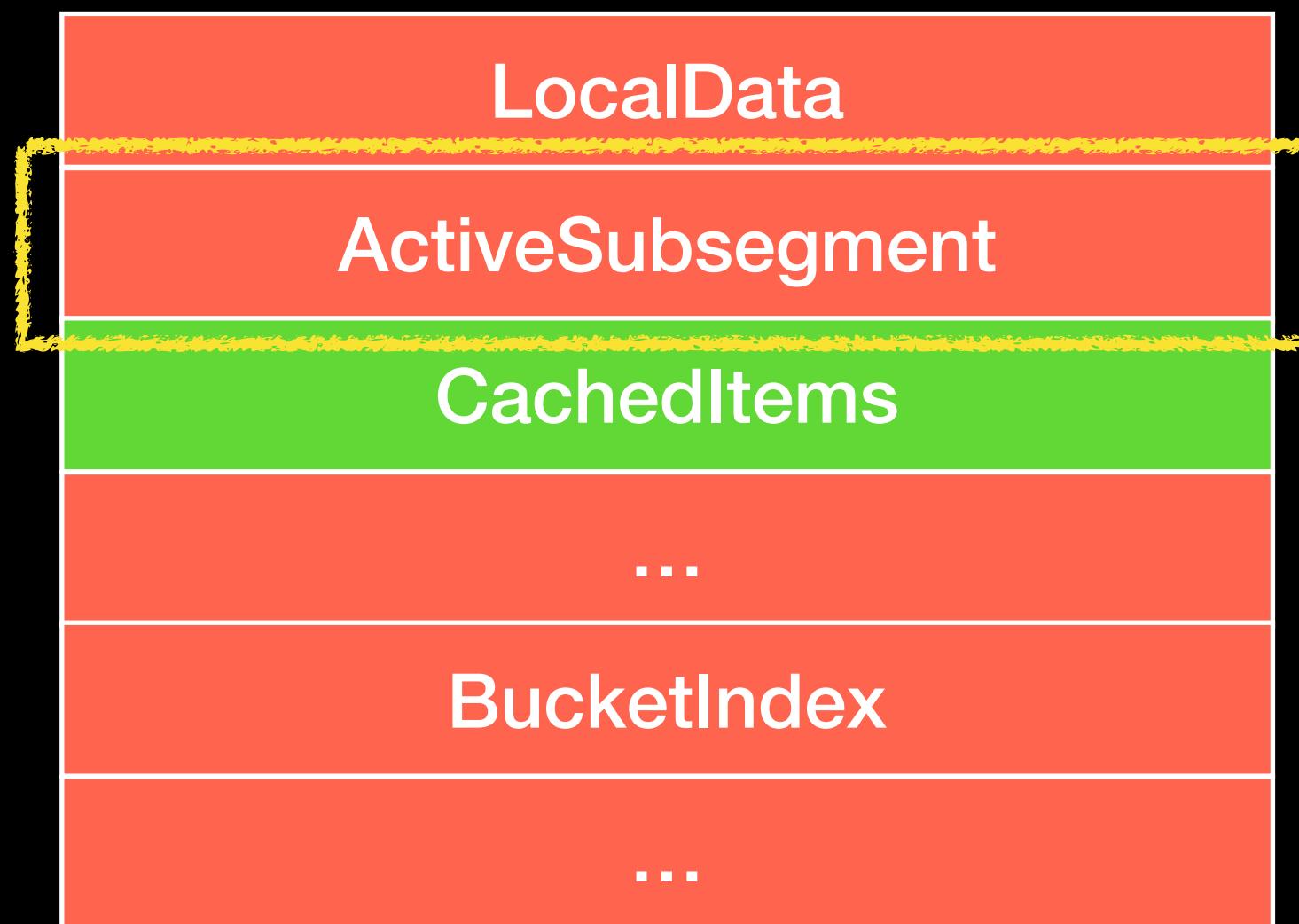
- SegmentInfoArray  
(\_HEAP\_LOCAL\_SEGMENT\_INFO)
  - LocalData (\_HEAP\_LOCAL\_DATA)
    - 對應到 \_LFH\_HEAP->LocalData 方便從 SegmentInfo 找回 \_LFH\_HEAP
  - BucketIndex
    - 對應到的 BucketIndex



# LFH

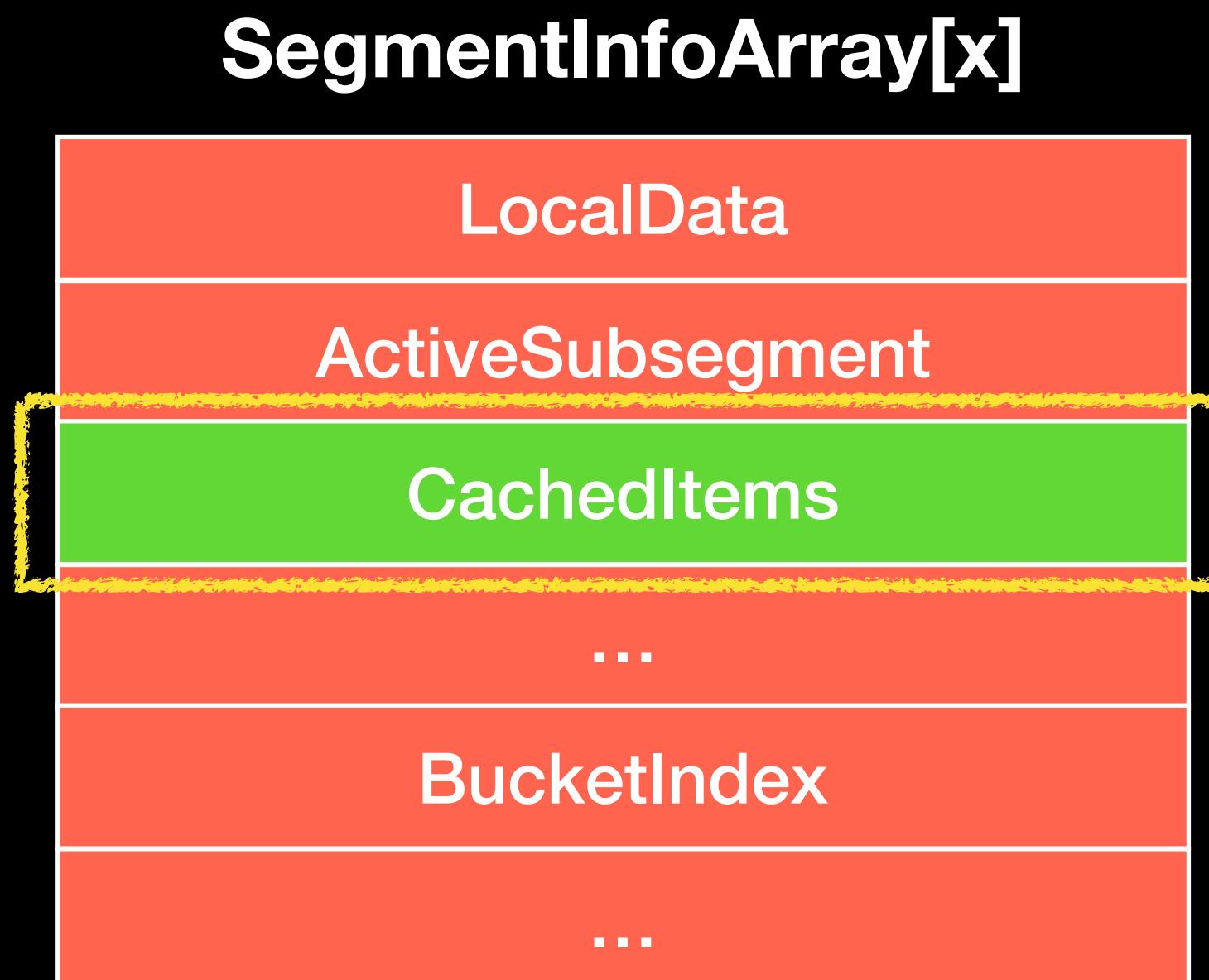
- SegmentInfoArray  
(\_HEAP\_LOCAL\_SEGMENT\_INFO)
  - ActiveSubsegment (\_HEAP\_SUBSEGMENT)
    - 極為重要結構
    - 對應到分配出去的 Subsegment
    - 其目的在於管理 Userblock
      - 記錄了剩餘多少 chunk
      - 該 Userblock 最大分配數等等

## SegmentInfoArray[x]

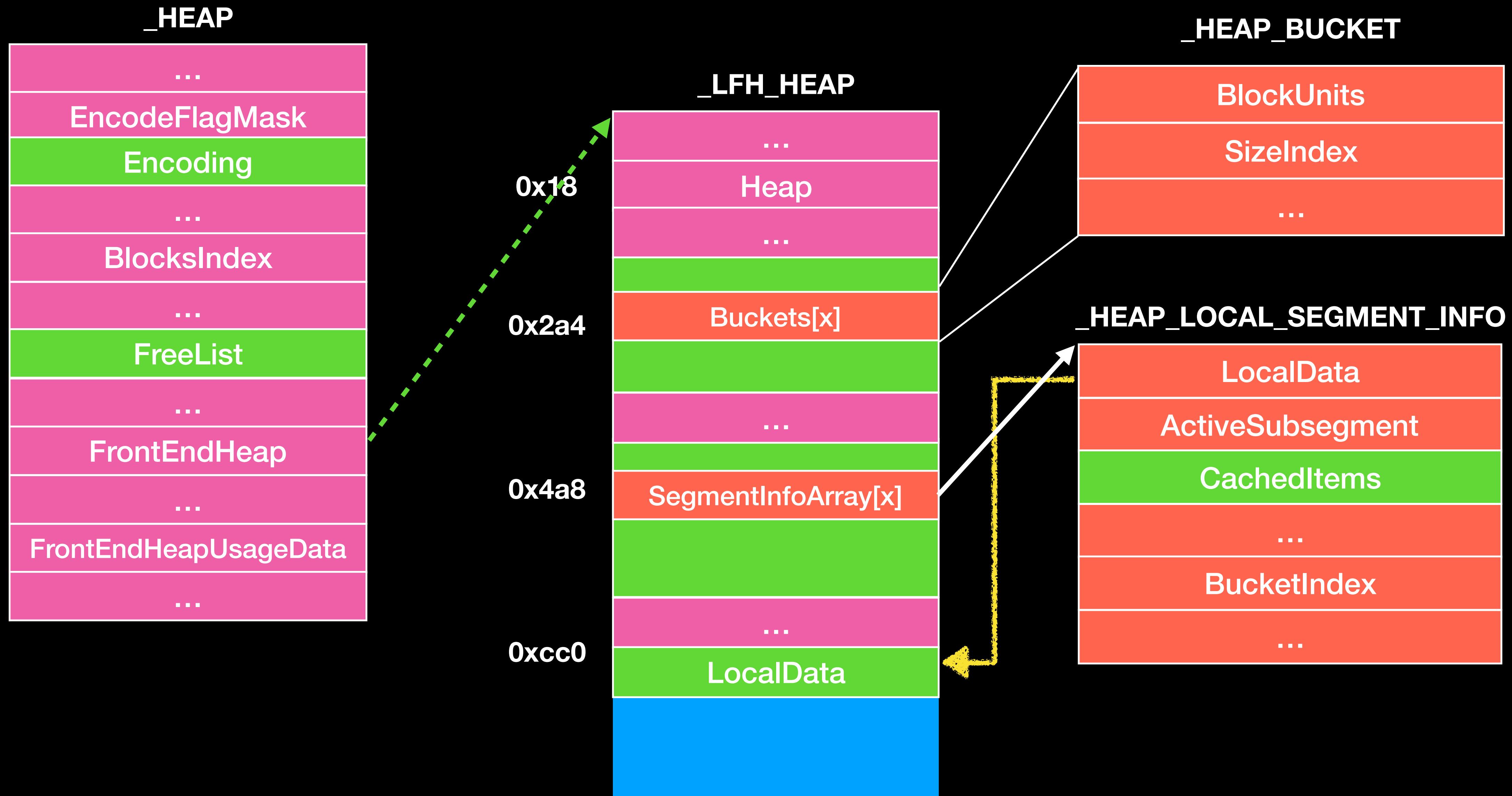


# LFH

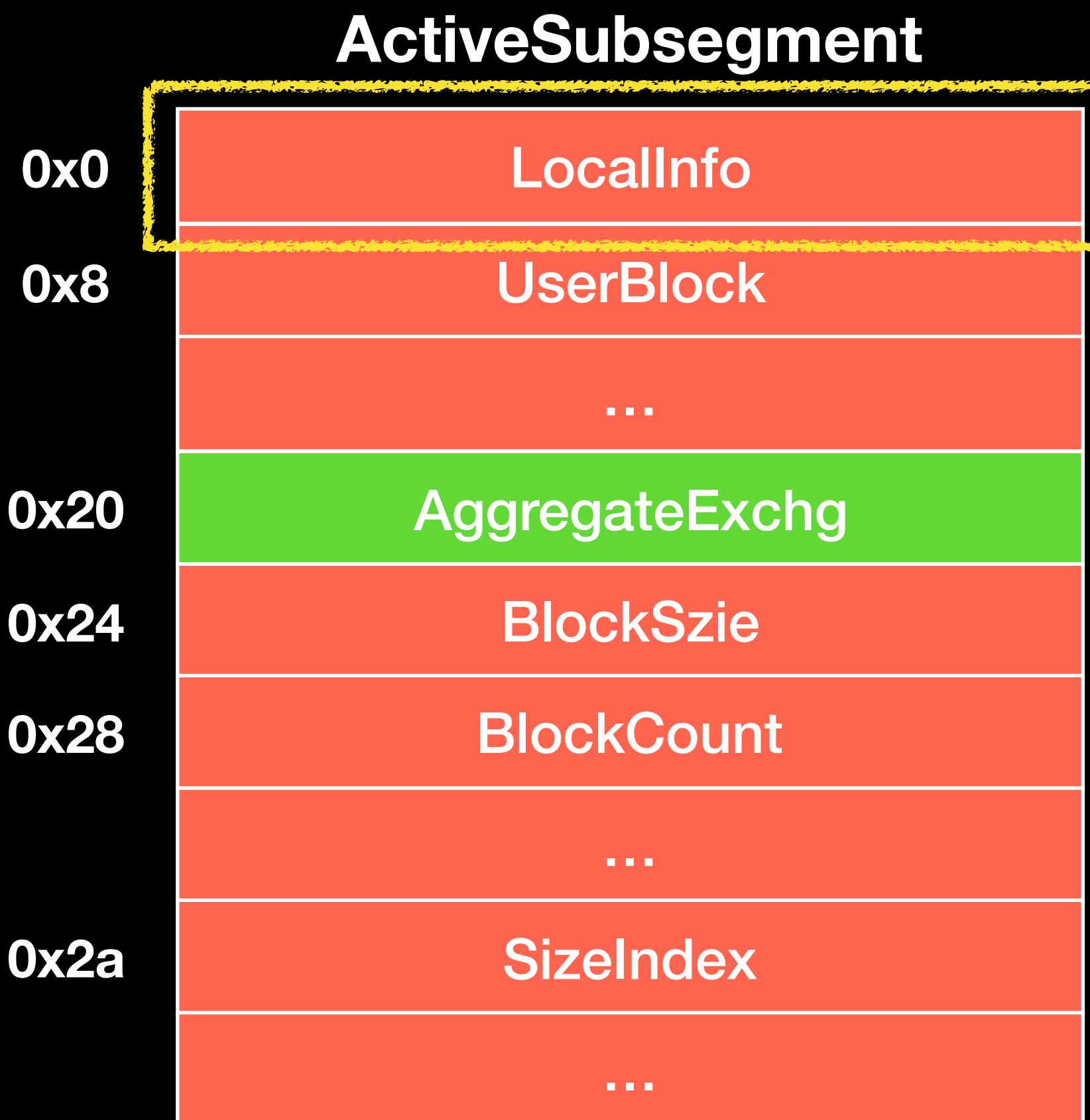
- SegmentInfoArray  
(`_HEAP_LOCAL_SEGMENT_INFO`)
- CachedItems (`_HEAP_SUBSEGMENT`)
- `_HEAP_SUBSEGMENT` array
  - 存放對映到該 SegmentInfo 且還有可以分配 chunk 給 user 的 Subsegment
  - 當 ActiveSubsegment 用完時，會從這邊填充，置換掉 ActiveSubsegment



# LFH

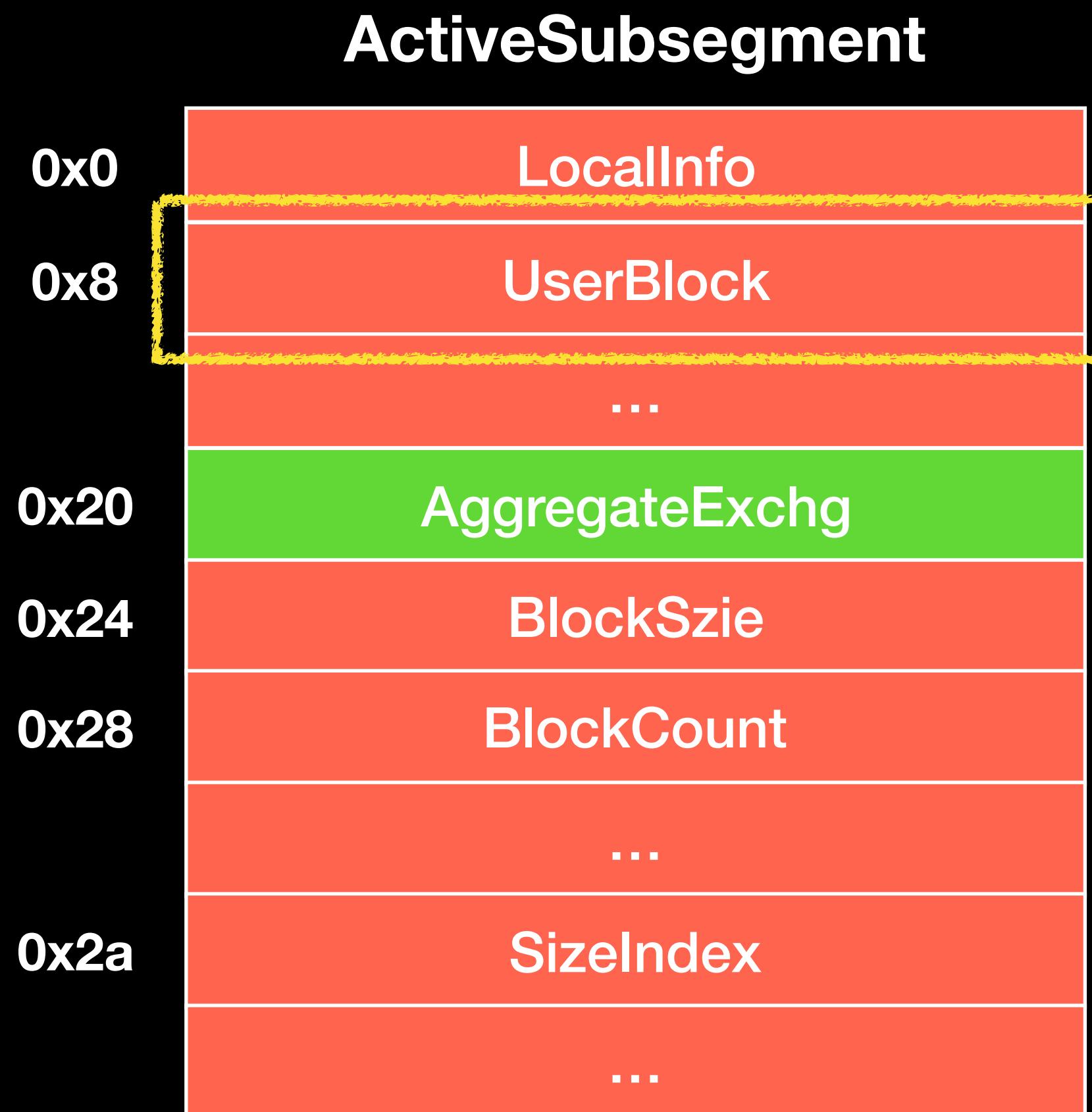


# LFH



- ActiveSubsegment (**\_HEAP\_SUBSEGMENT**)
  - LocalInfo  
(\_HEAP\_LOCAL\_SEGMENT\_INFO)
  - 指回對應到的  
**\_HEAP\_LOCAL\_SEGMENT\_INFO**

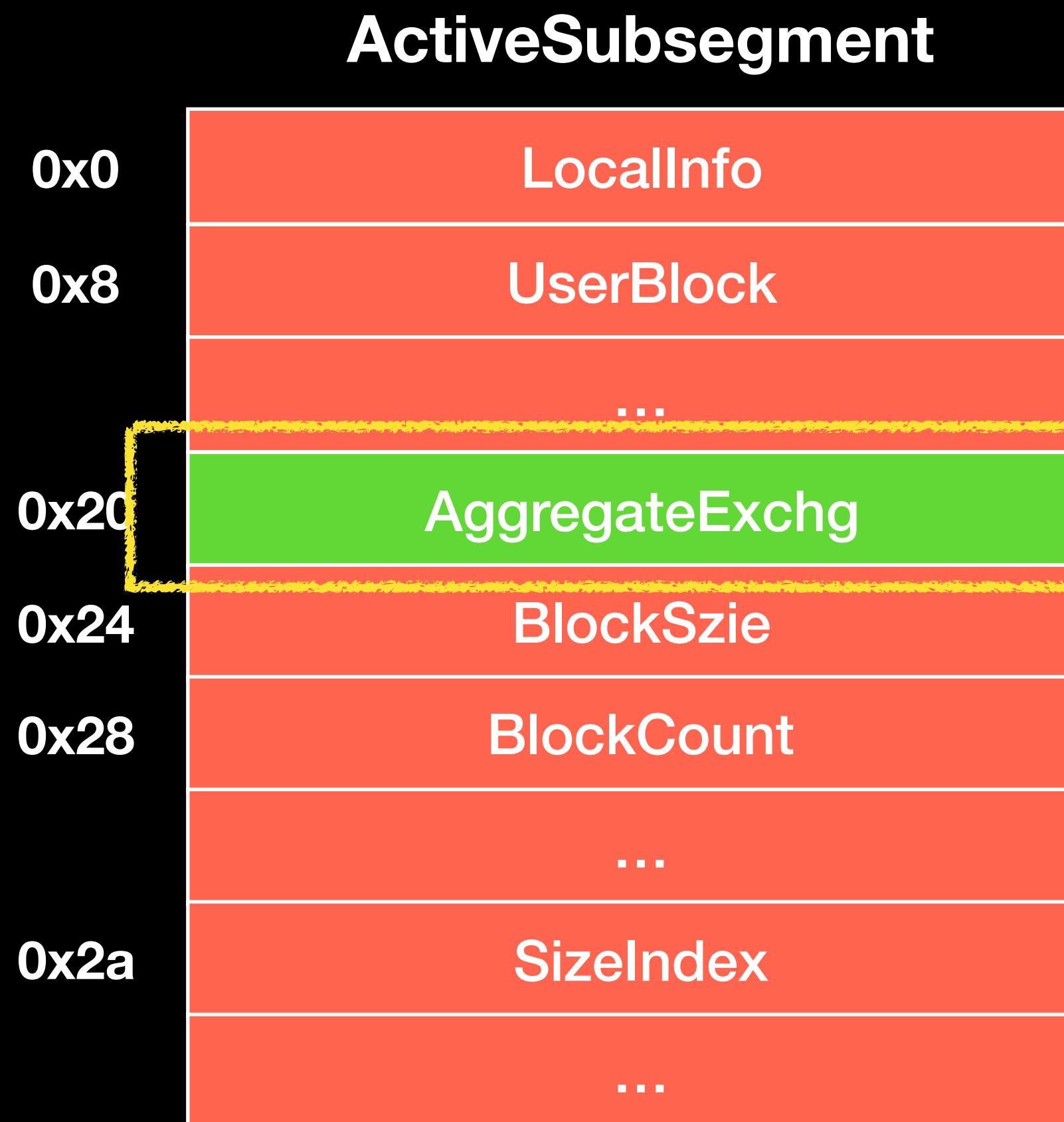
# LFH



- ActiveSubsegment ([\\_HEAP\\_SUBSEGMENT](#))
  - UserBlock ([\\_HEAP\\_USERDATA\\_HEADER](#))
    - LFH 的記憶體分配池
    - 也就是要分配出去的 Chunk (Block) 所在位置
    - 開頭會有些 metadata 管理這些 chunk
    - 重要結構

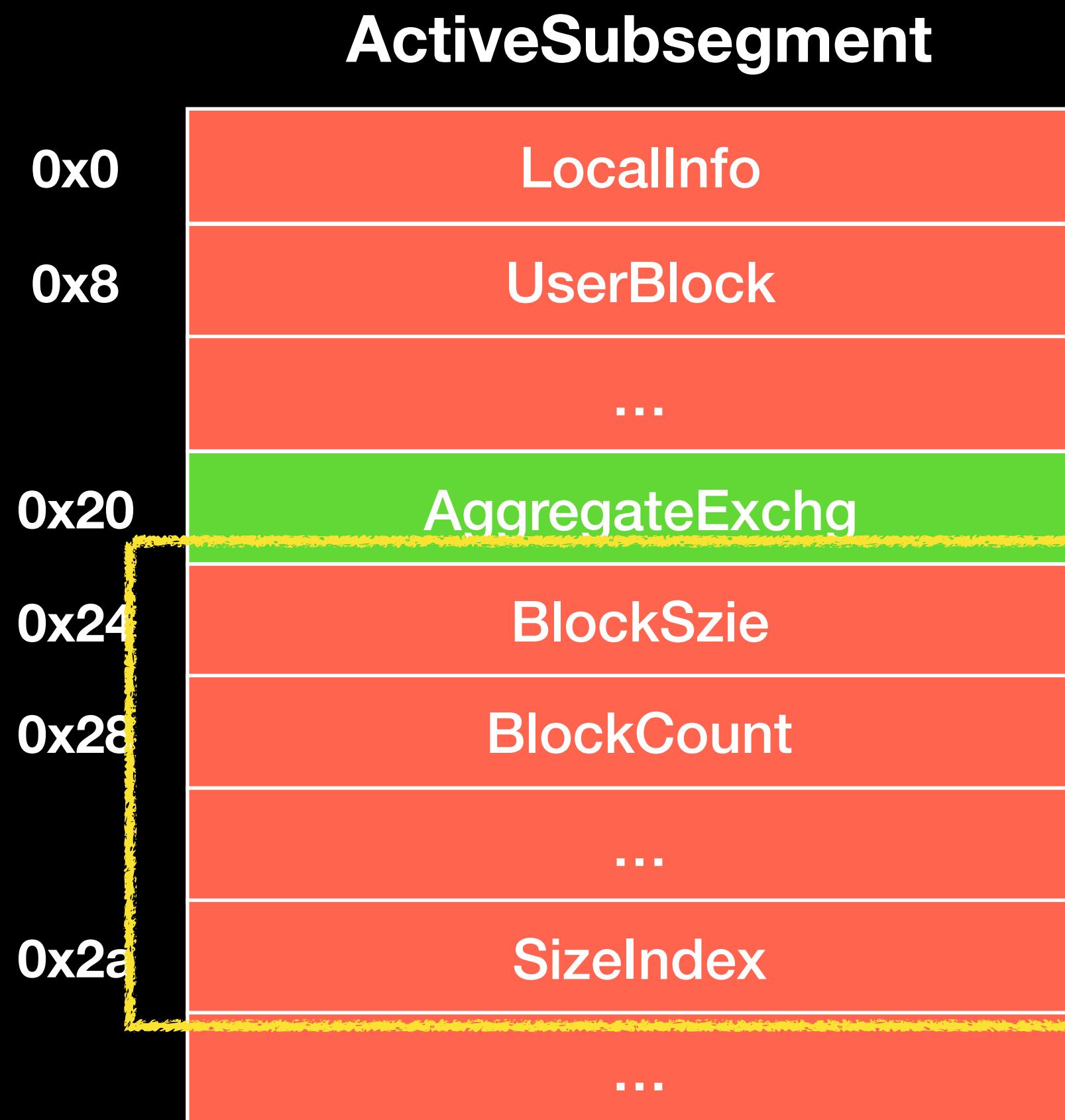
# LFH

- ActiveSubsegment (`_HEAP_SUBSEGMENT`)



- AggregateExchg (`_INTERLOCK_SEQ`)
  - 用來管理對應到的 UserBlock 中還有多少 freed chunk
  - LFH 用這個判斷是否還從該 UserBlock 分配
  - 同時也有 Lock 的作用

# LFH



- ActiveSubsegment (`_HEAP_SUBSEGMENT`)
  - BlockSize
    - 在該 UserBlock 中每個 Block (chunk) 的大小
  - BlockCount
    - 在該 UserBlock 中 Block 的總數
  - SizelIndex
    - 該 UserBlock 對應到的 SizelIndex

# LFH

- AggregateExchg (\_INTERLOCK\_SEQ)

- Depth

## AggregateExchg

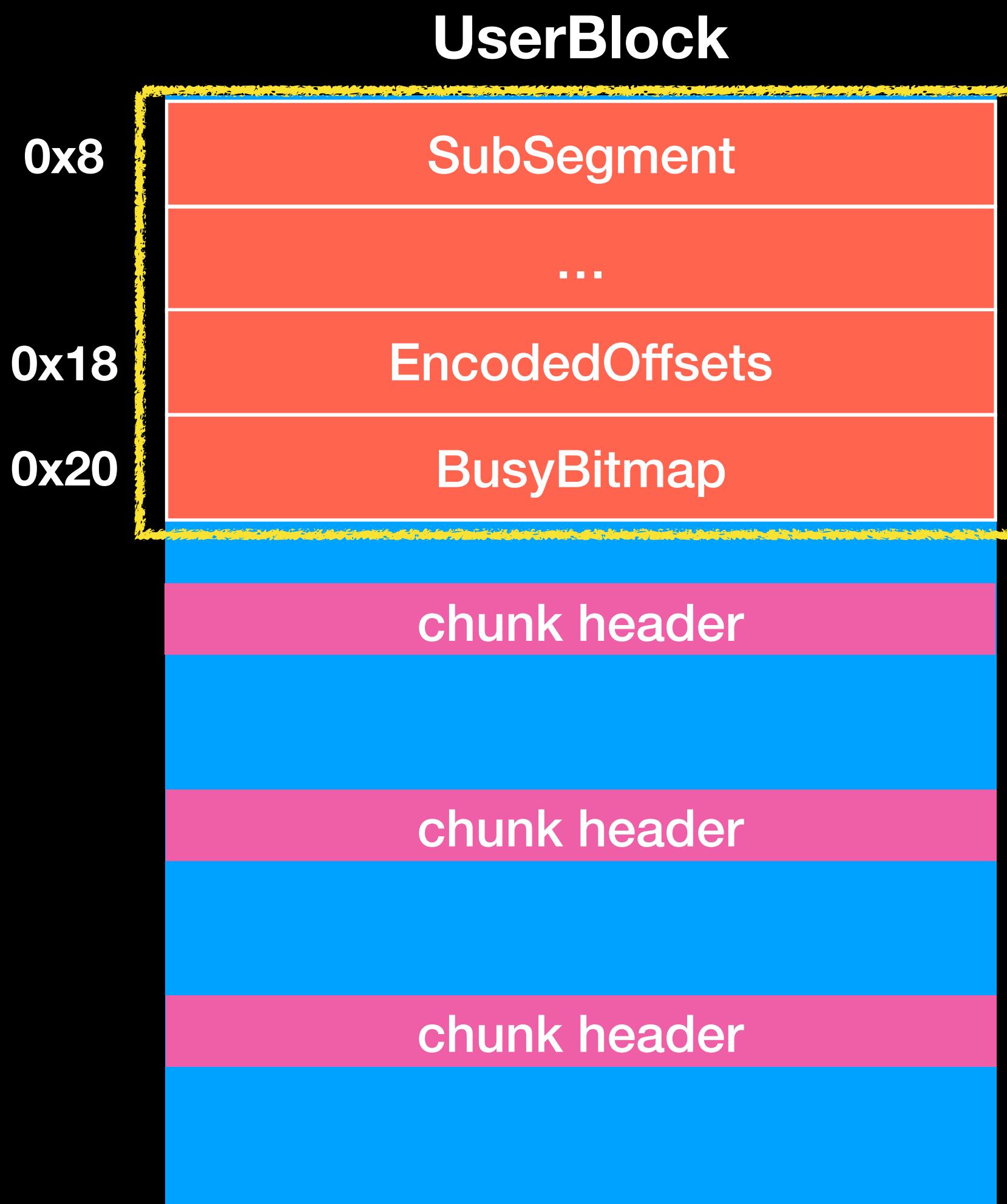
0x0	Depth
0x8	Hint (15bit)
	Lock (1bit)

- 該 UserBlock 所剩下的 Freed chunk 數量

- Lock

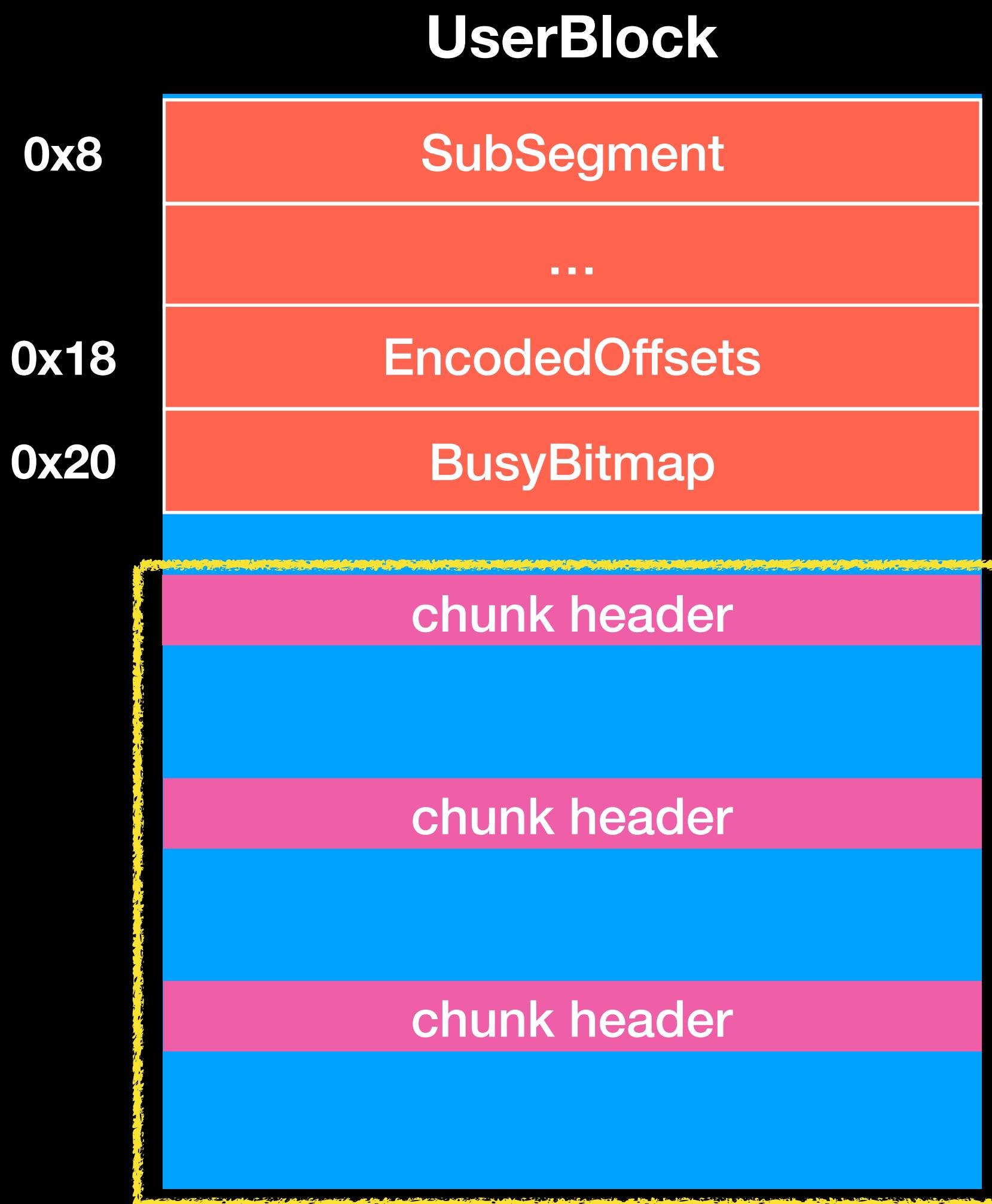
- 就是 Lock

# LFH



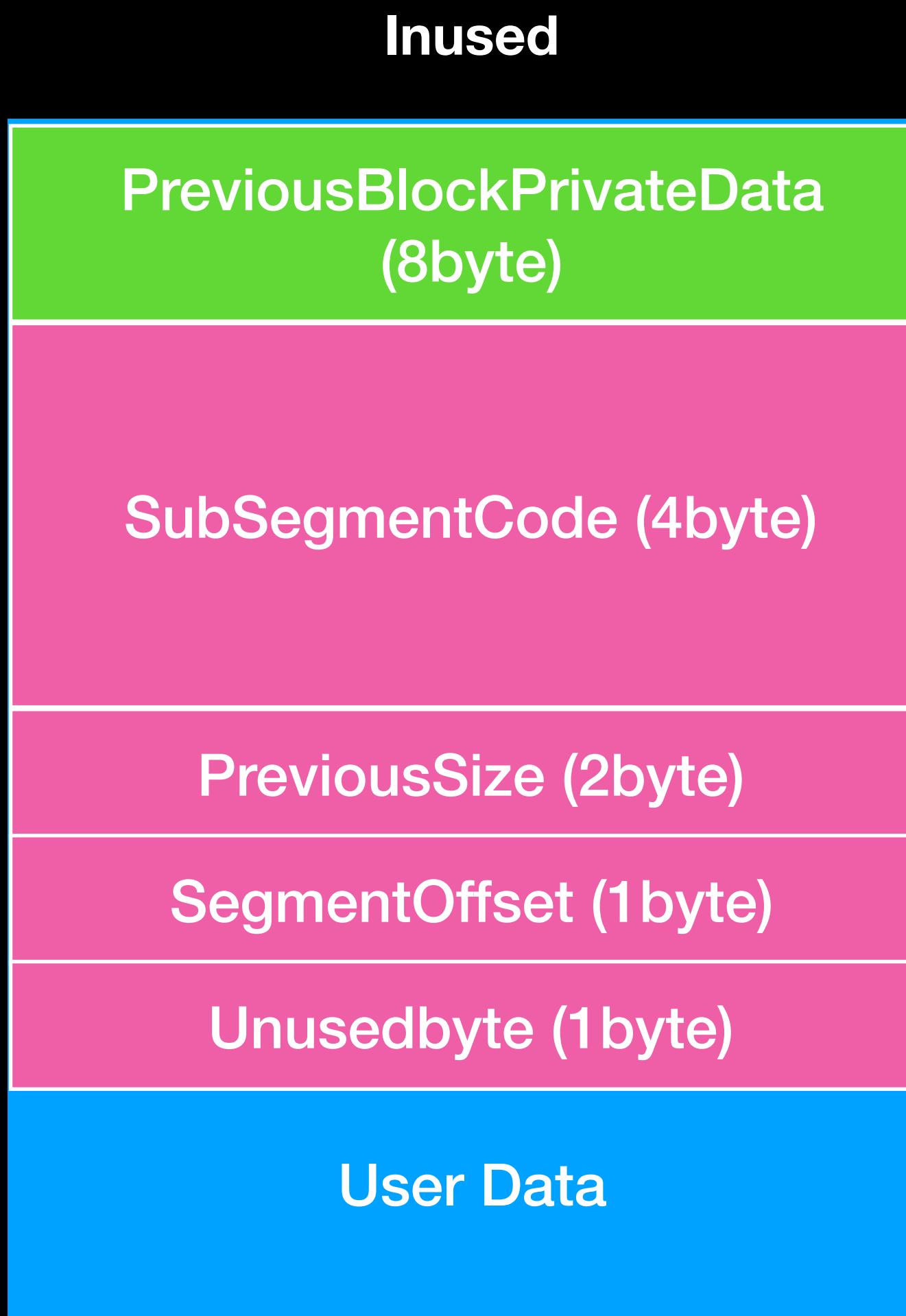
- UserBlock (\_HEAP\_USERDATA\_HEADER)
  - SubSegment
    - 指回對應的 Subsegment
  - EncodedOffsets
    - 用來驗證 chunk header 是否有被改過
  - BusyBitmap
    - 記錄該 UserBlock 哪些 chunk 有在用的 bitmap

# LFH



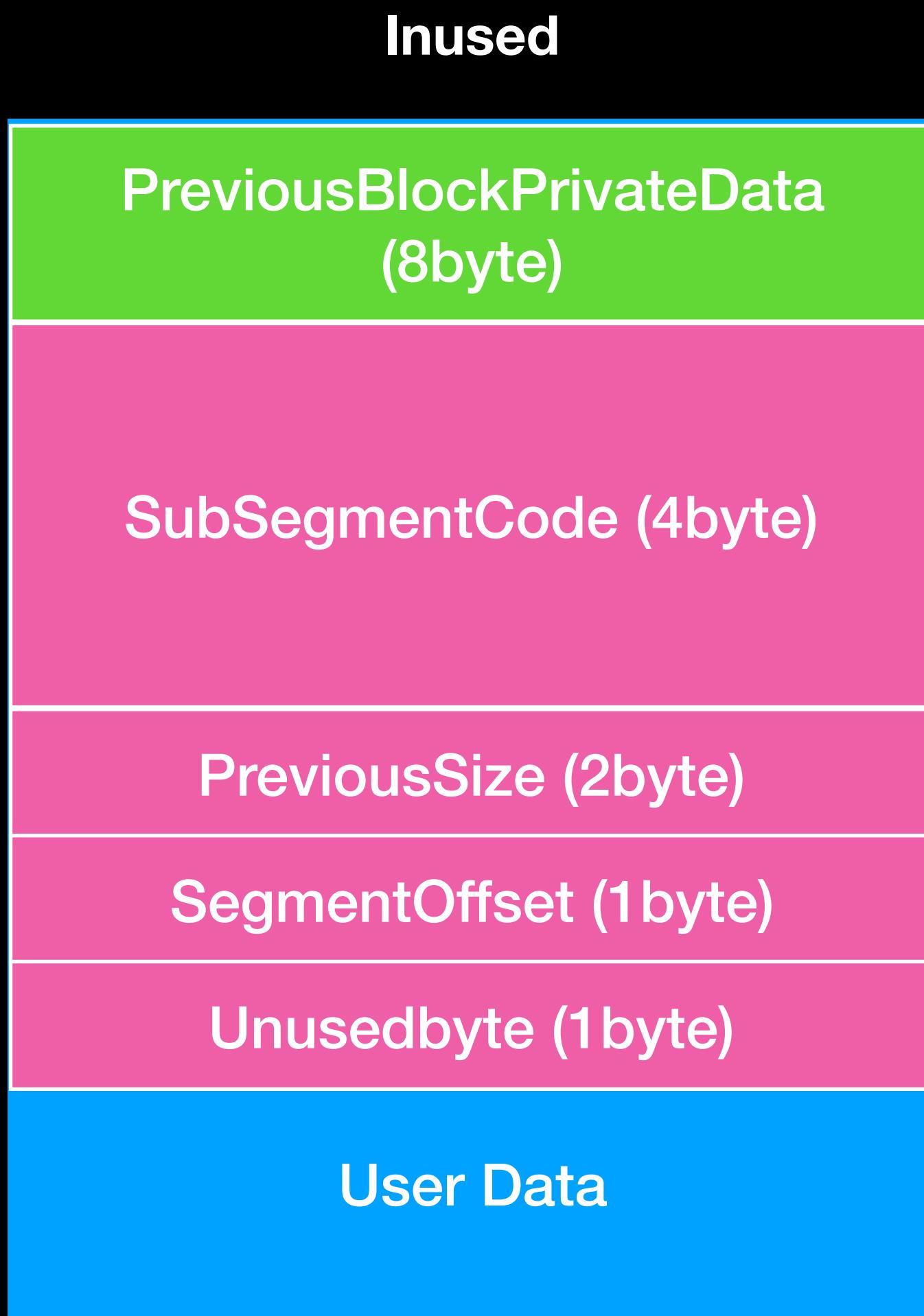
- UserBlock (\_HEAP\_USERDATA\_HEADER)
  - Block (chunk)
  - LFH 回傳給使用者的 chunk

# LFH



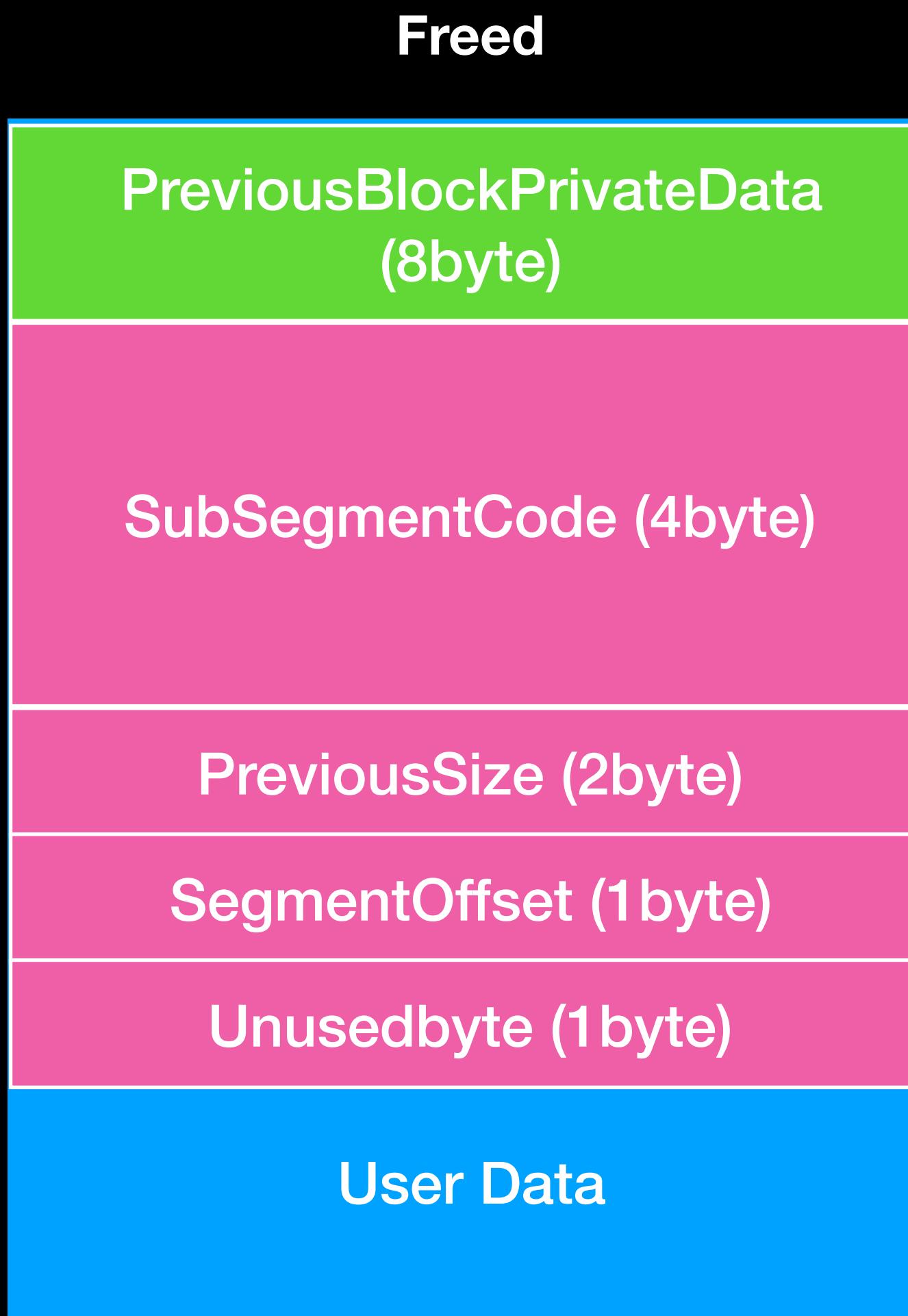
- \_HEAP\_ENTRY (chunk)
  - SubSegmentCode
    - Encode 過的 metadata 用來推回 userblock 的位置
  - PreviousSize
    - 該 chunk 在 UserBlock 中的 index

# LFH



- `_HEAP_ENTRY` (chunk)
  - Unusedbyte
  - Unusedbyte & 0x80 is true

# LFH



- `_HEAP_ENTRY` (chunk)
  - Unusedbytes
    - 恒為 0x80
    - 用來判斷是否為 LFH 的 Free chunk

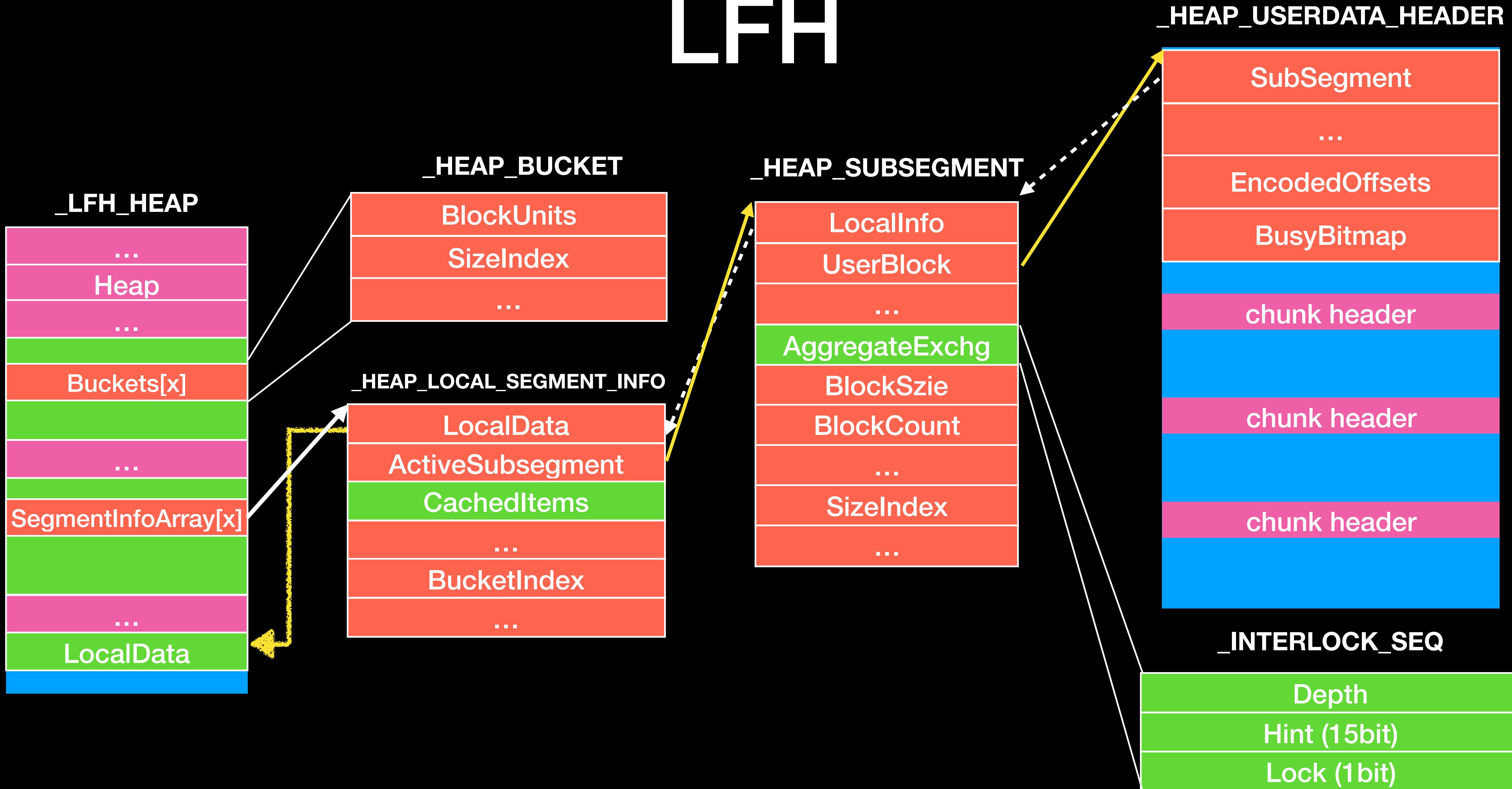
# LFH

- Remark
  - 關於 EncodedOffsets
  - 在 UserBlock 初始化時設置，其算法為下面四個值的 xor
    - $(\text{sizeof(userblock header}) \mid (\text{BlockUnit} * 0x10) << 16))$
    - LFHkey
    - Userblock address
    - \_LFH\_HEAP address

# LFH

- Remark
  - 關於 LFH header encoding
  - 所有的 chunk header 在初始化時都會經過 xor 過，算法為下面個值的 xor
    - \_HEAP address
    - LFHkey
    - Chunk address  $\gg 4$
    - $((\text{chunk address}) - (\text{UserBlock address})) \ll 12$

# LFH



# Windows memory allocator

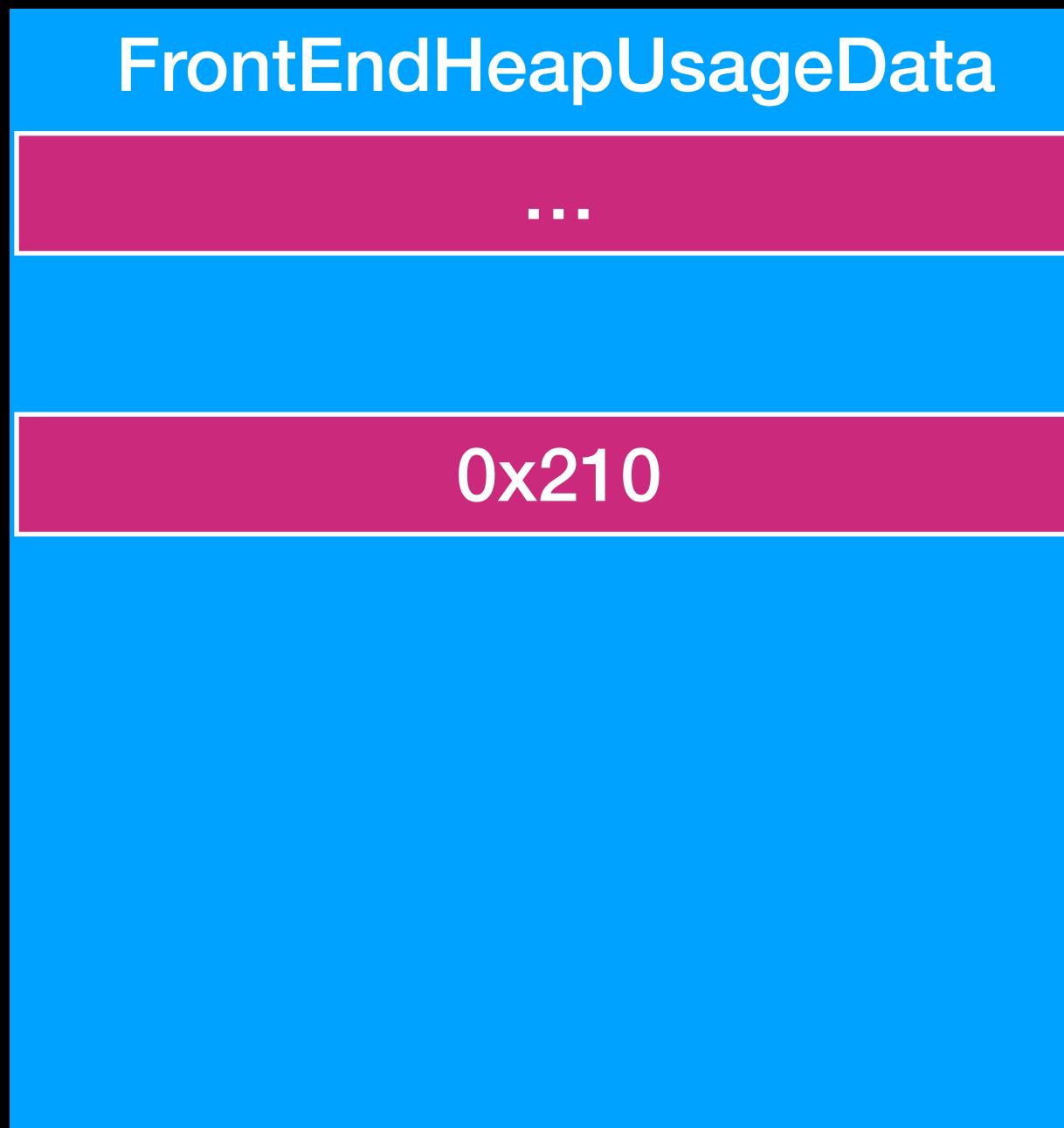
- Nt Heap
  - 前端管理器 (Front-End)
    - Data Structure
    - 分配機制

# Windows memory allocator

- Nt Heap
  - 前端管理器 (Front-End)
    - 初始话
    - 在 FrontEndHeapUsageData[x] & 0x1f > 0x10 时，下一次的 allocate 会对 LFH 做出初始化
    - 会先 ExtendFrontEndUsageData 及增加更大的 BlocksIndex (0x80-0x400)
    - 建立 FrontEndHeap
    - 初始化 SegmentInfoArrays[idx]
    - 接下来在 allocate 相同大小的 chunk 就会开始使用 LFH

# Windows memory allocator

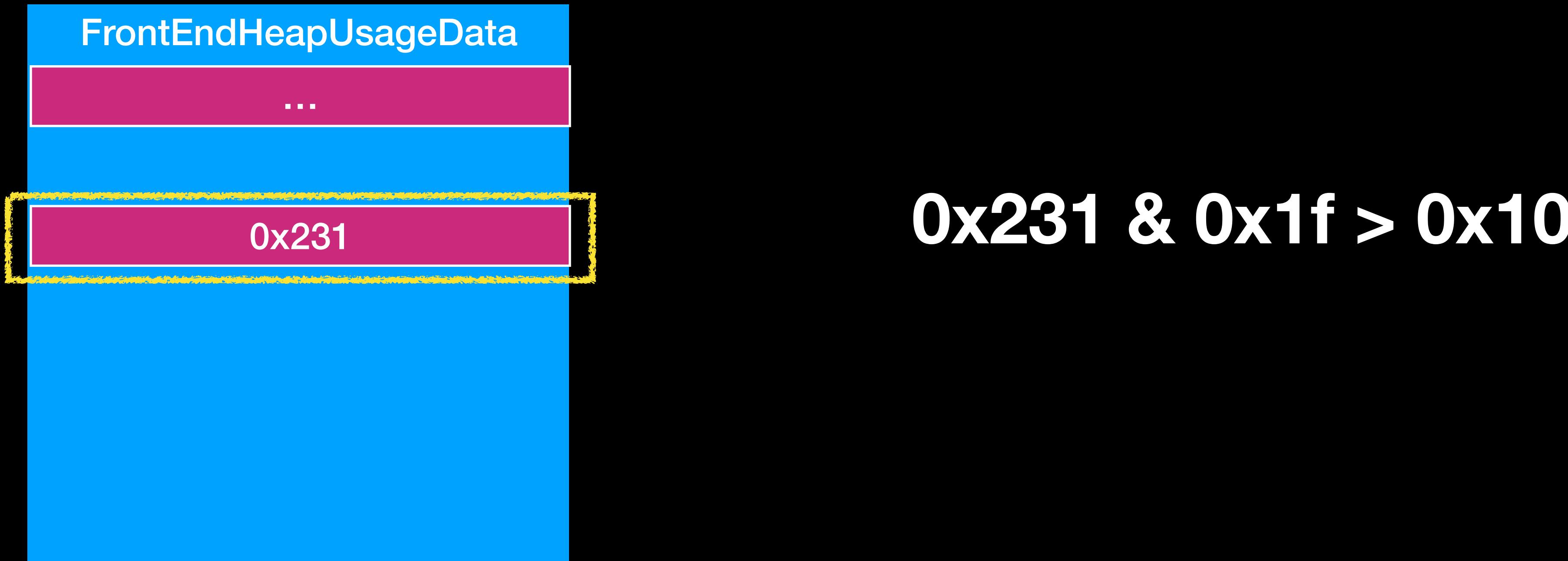
- LFH(初始化)
  - `malloc(0x40) * 16`



此時可 enable LFH 範圍為  
index 0x0 - 0x80

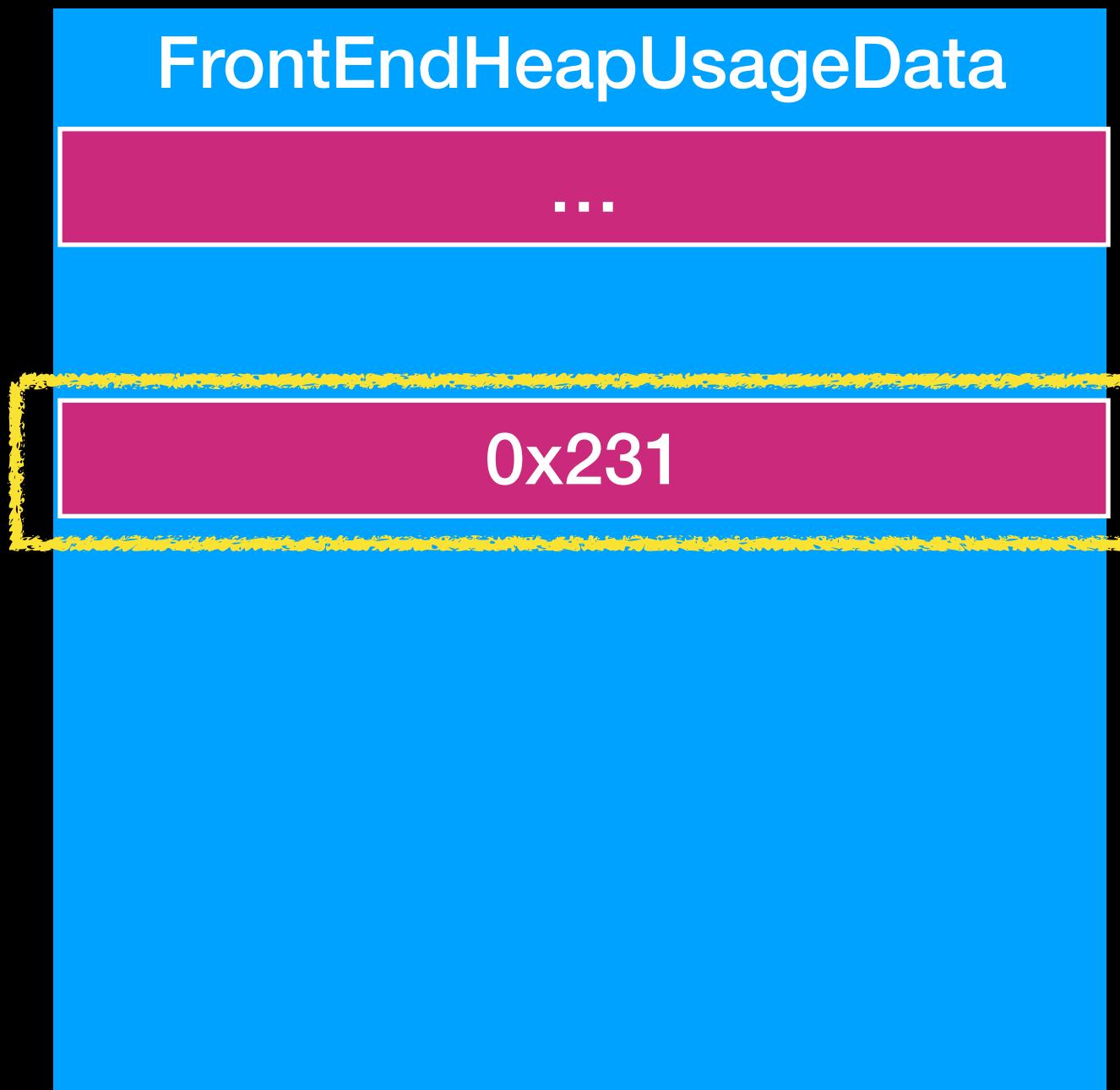
# Windows memory allocator

- LFH(初始化)
  - malloc(0x40) (第 17 次)



# Windows memory allocator

- LFH(初始化)

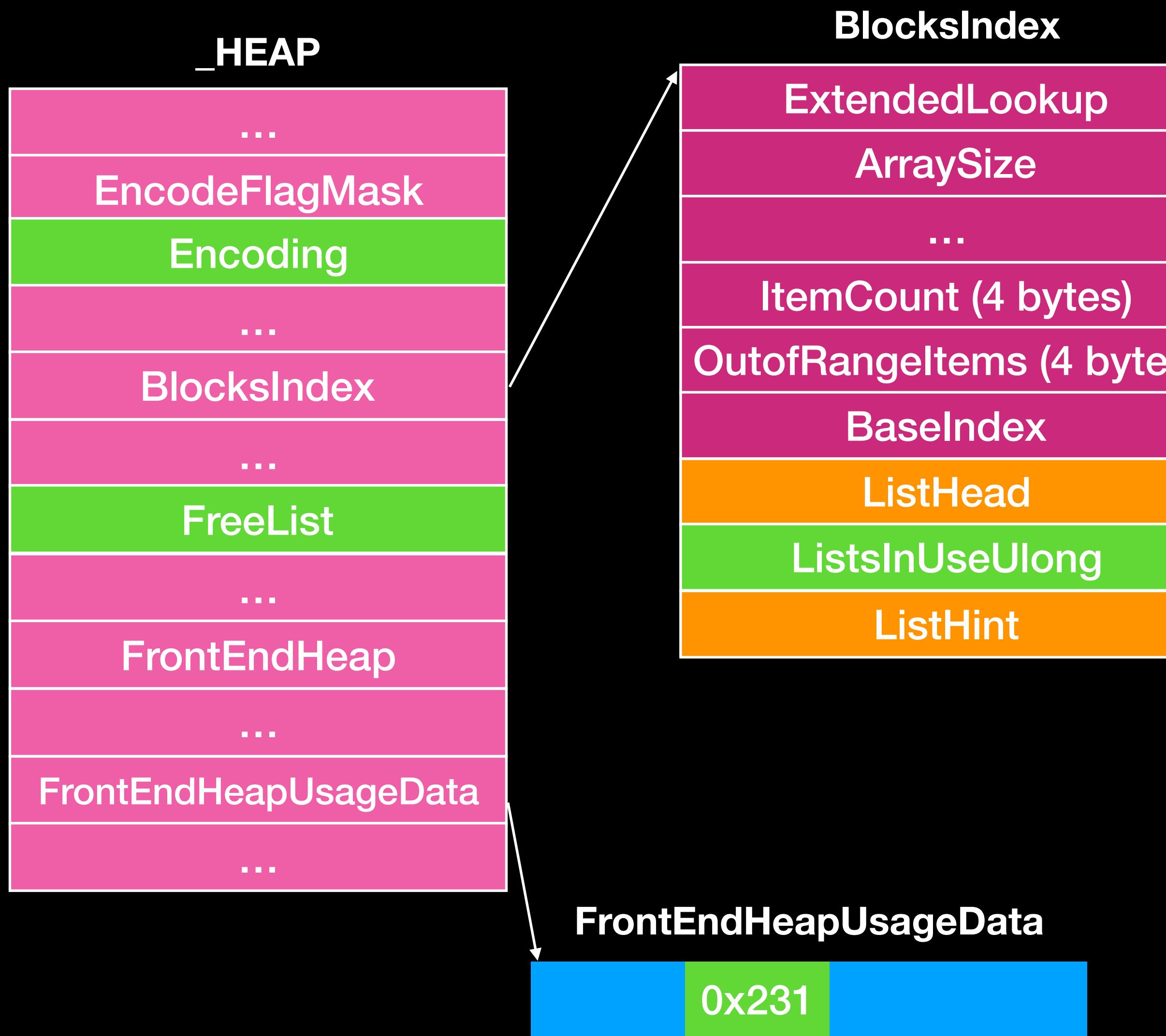


**heap->CompatibilityFlag |= 0x20000000**

設置上這個 flag 後

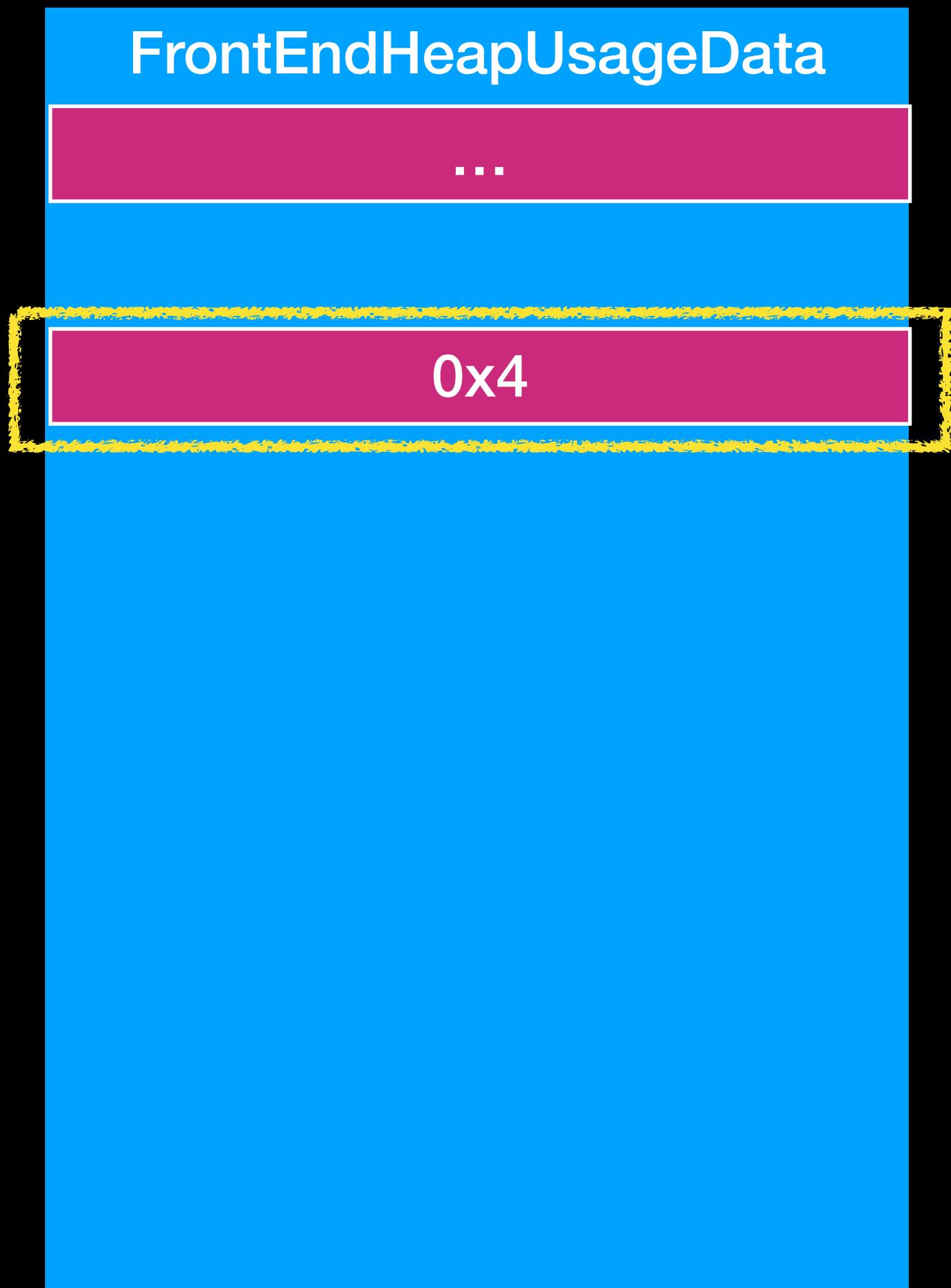
下次 allocate 時會去初始化LFH

# Windows memory allocator



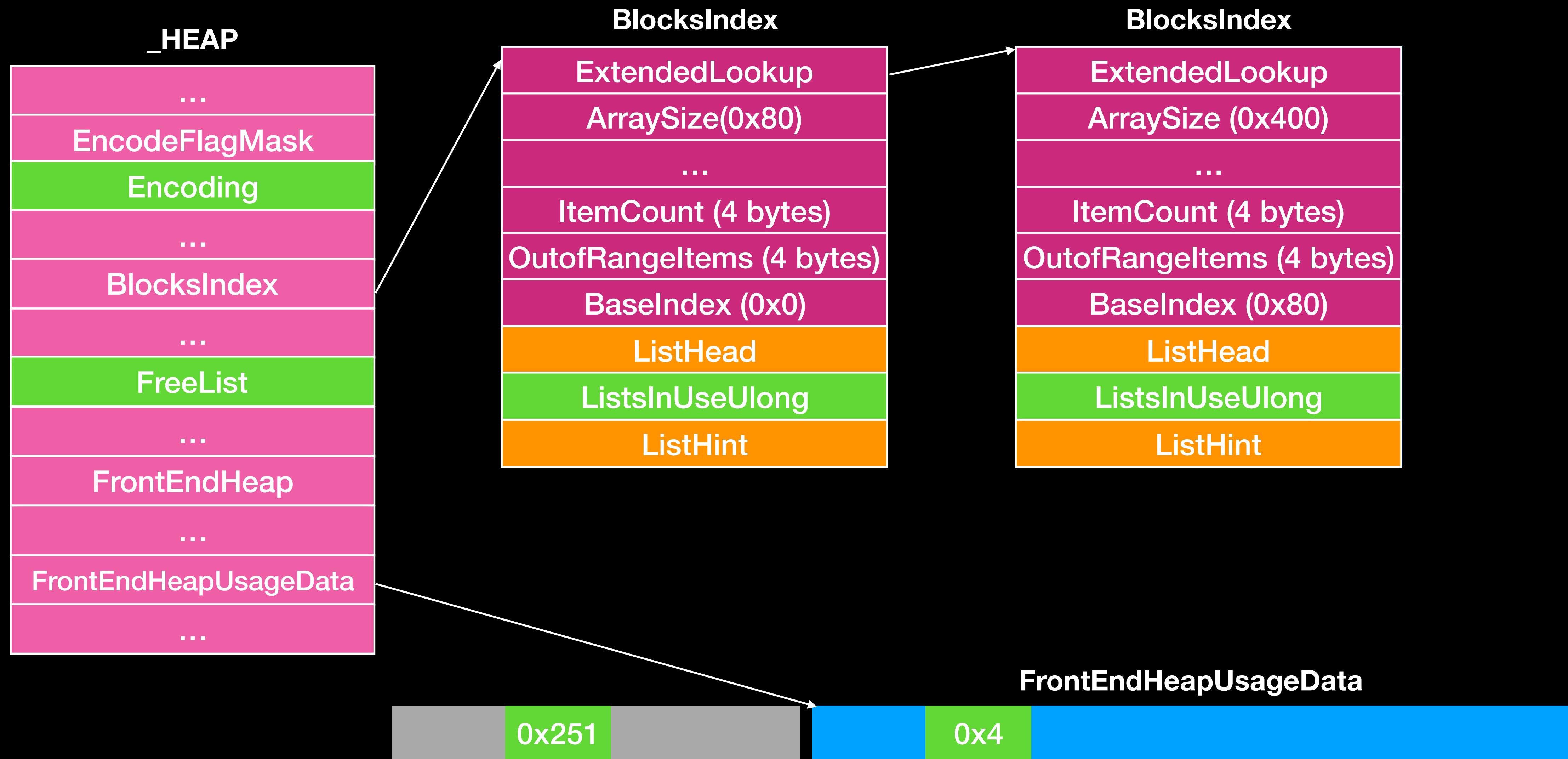
# Windows memory allocator

- LFH(初始化)

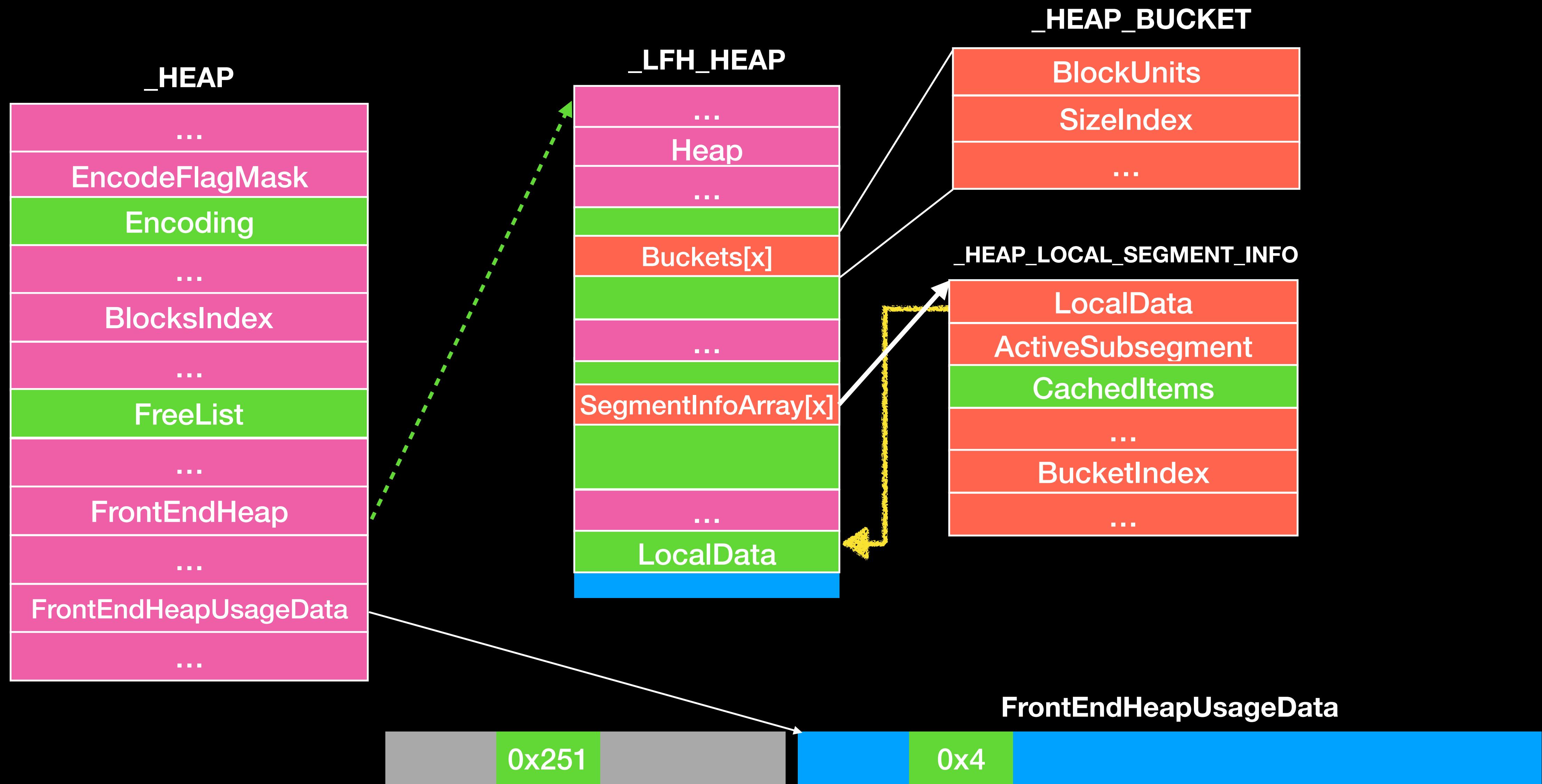


- `malloc(0x40)` (第 18 次)
- `ExtendFrontEndUsageData` 及增加更大的 `BlocksIndex` (`0x80-0x400`)，並設置對應的 bitmap
- 並在該對應的 `FrontEndHeapUsageData` 寫上對應的 index，此時可 enable LFH 範圍變為 (`idx: 0-0x400`)
- 建立並初始化 `FrontEndHeap` (`mmap`)
- 初始化 `SegmentInfoArrays[idx]`
- 在 `SegmentInfoArrays[BucketIndex]` 填上 `segmentInfo`

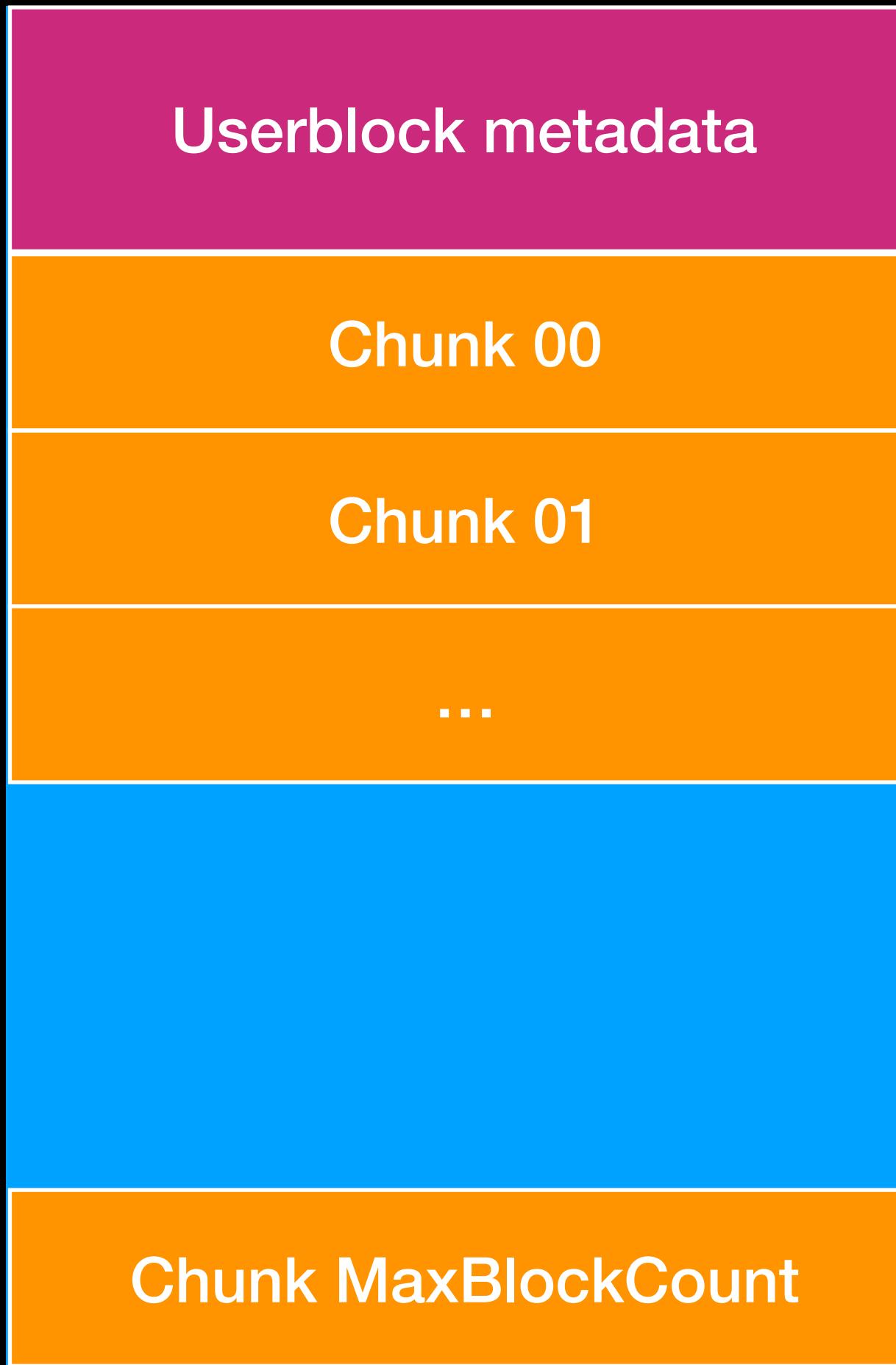
# Windows memory allocator



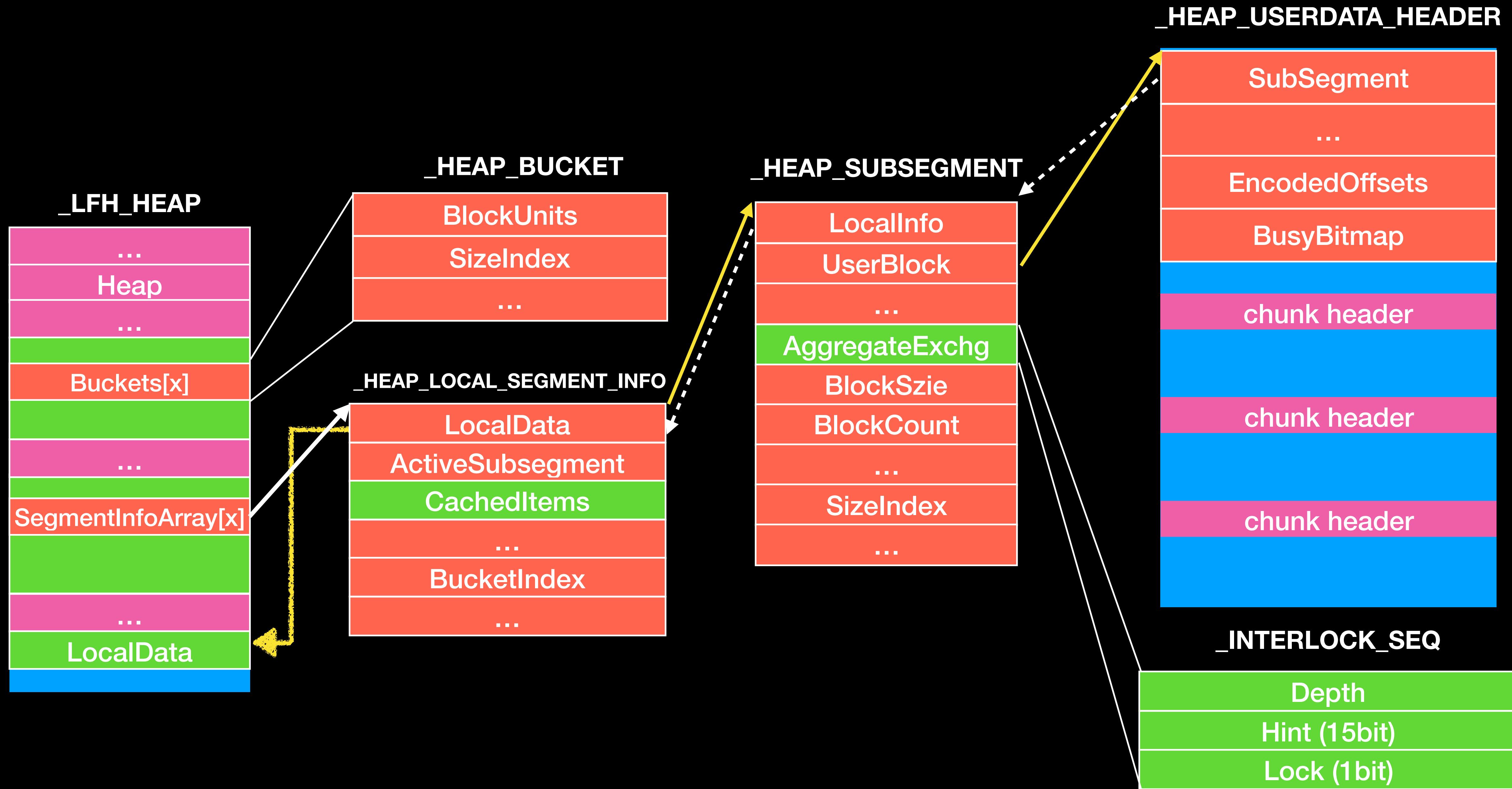
# Windows memory allocator



# Windows memory allocator

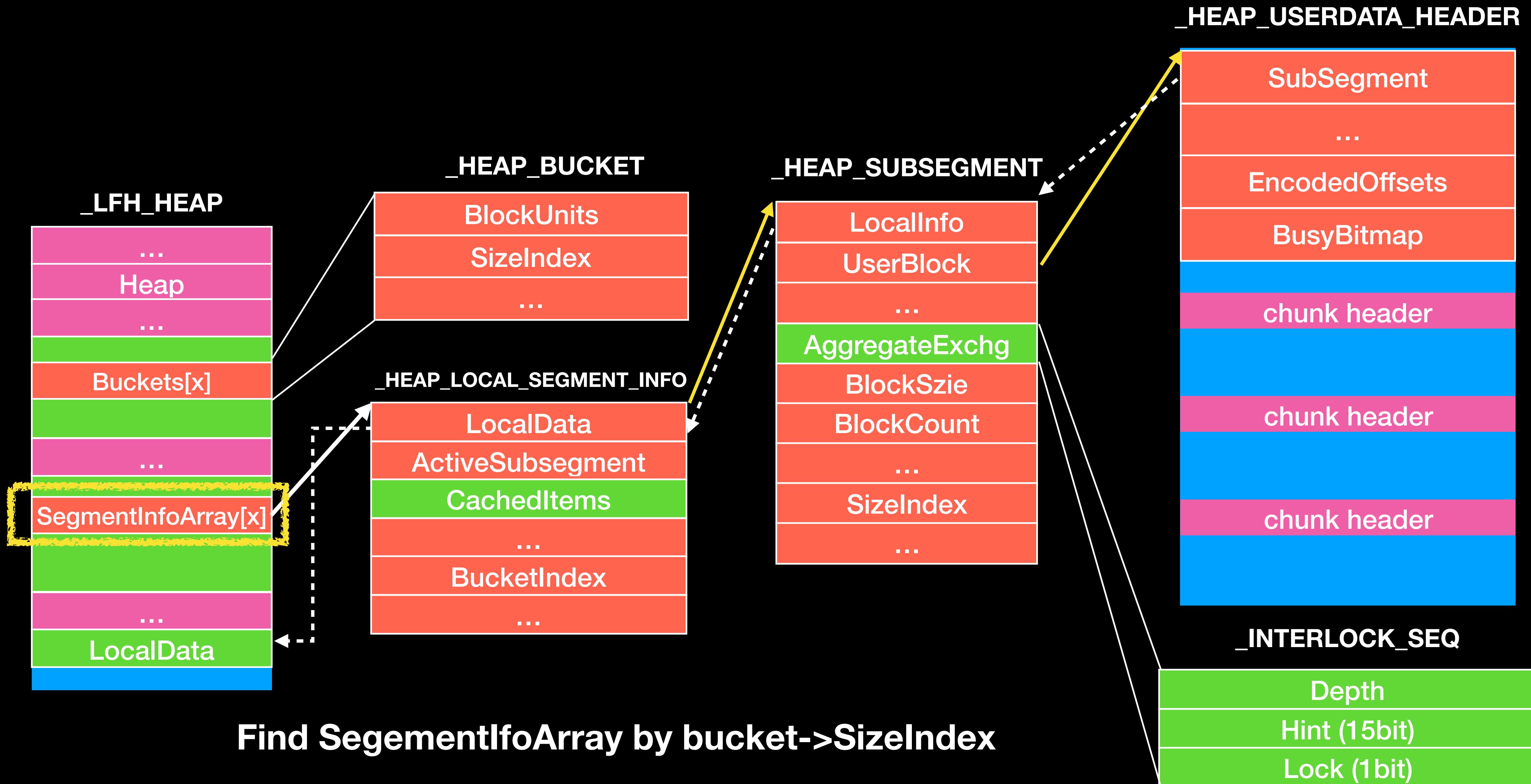


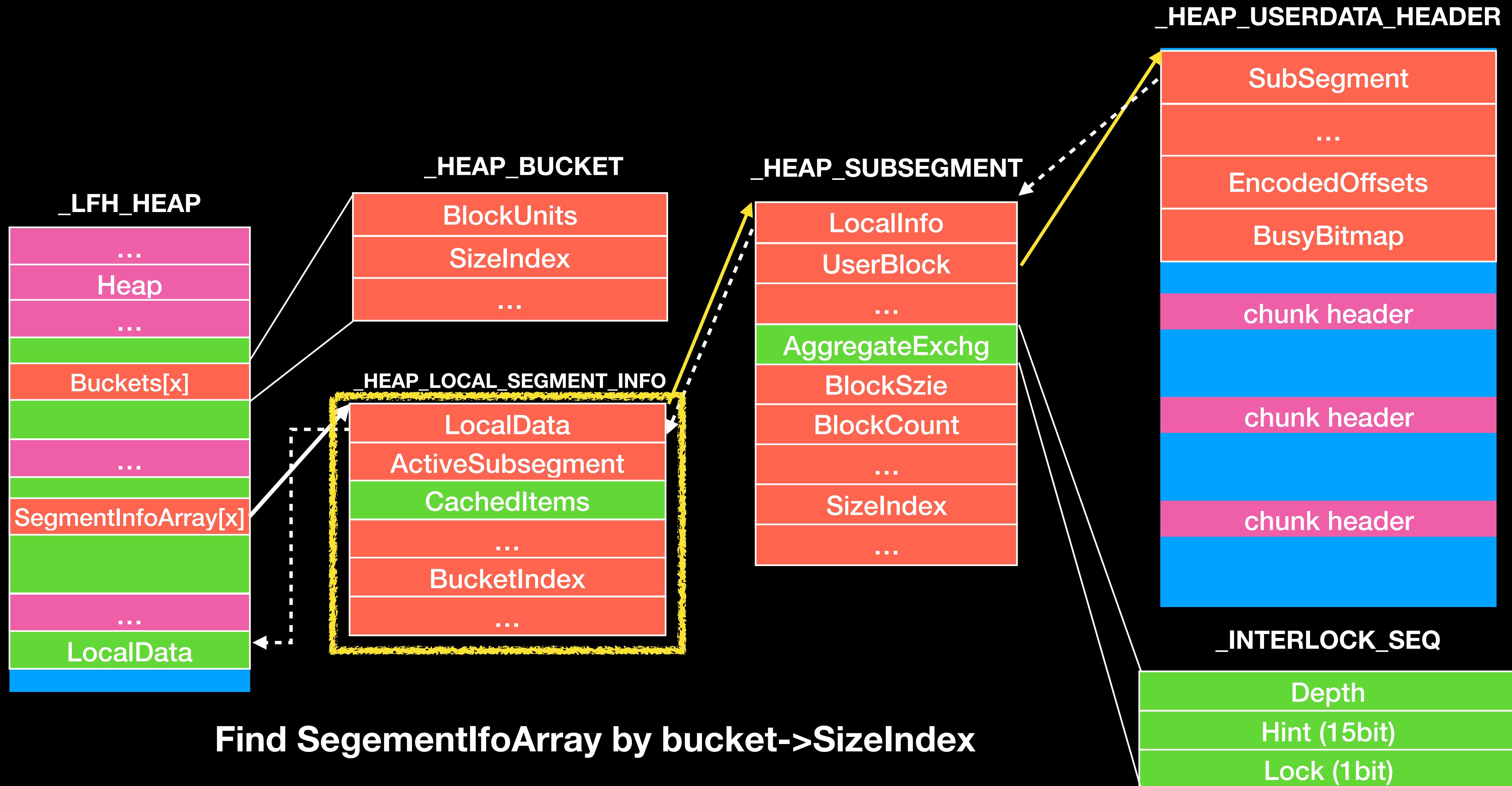
- LFH(初始化)
  - malloc(0x40) (第 19 次)
  - Allocate Userblock 並初始化
    - 設置對應的 chunk
    - 設置對應的 ActiviteSubsegment
    - 隨機返回其中的一個 chunk

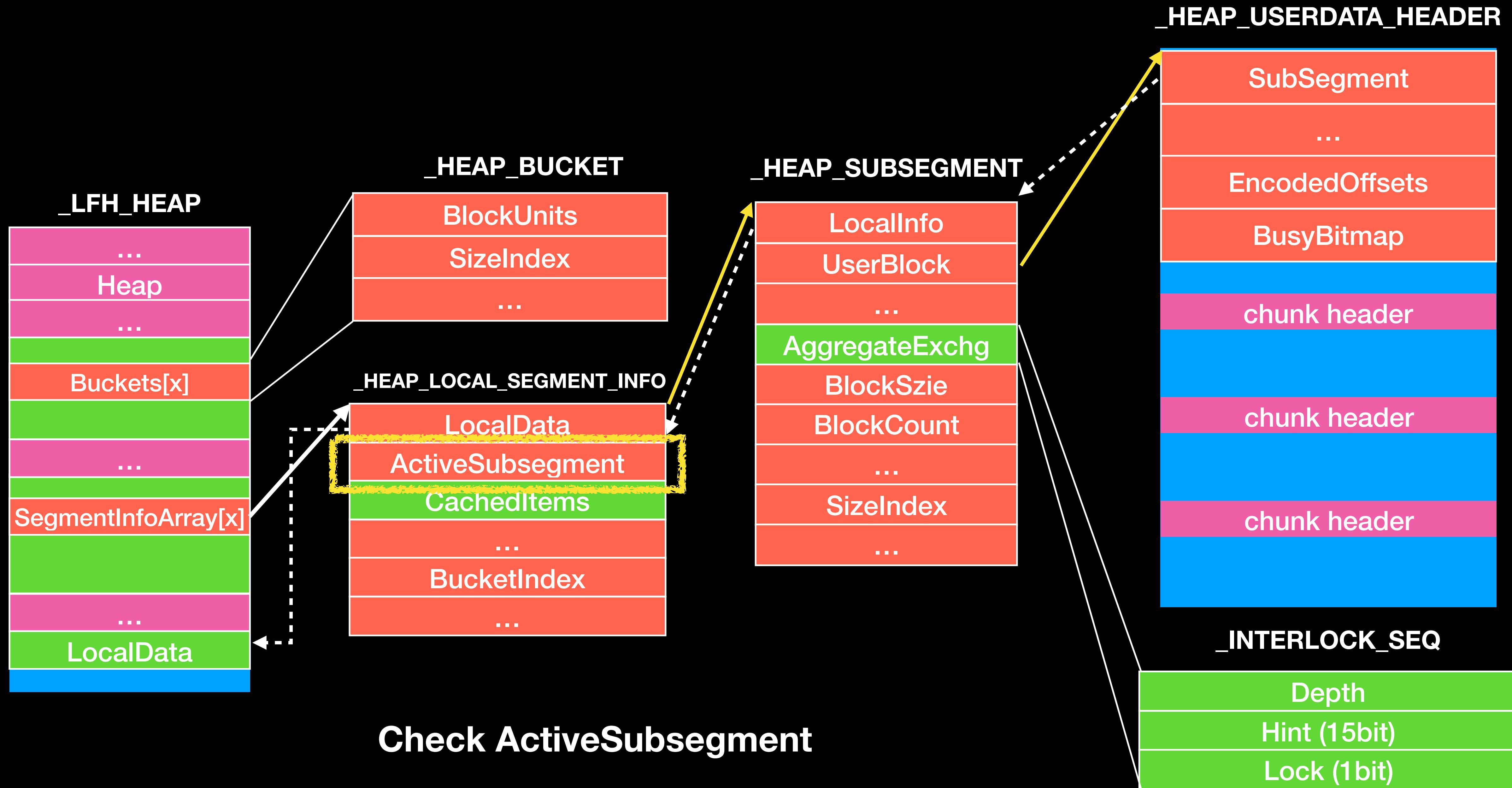


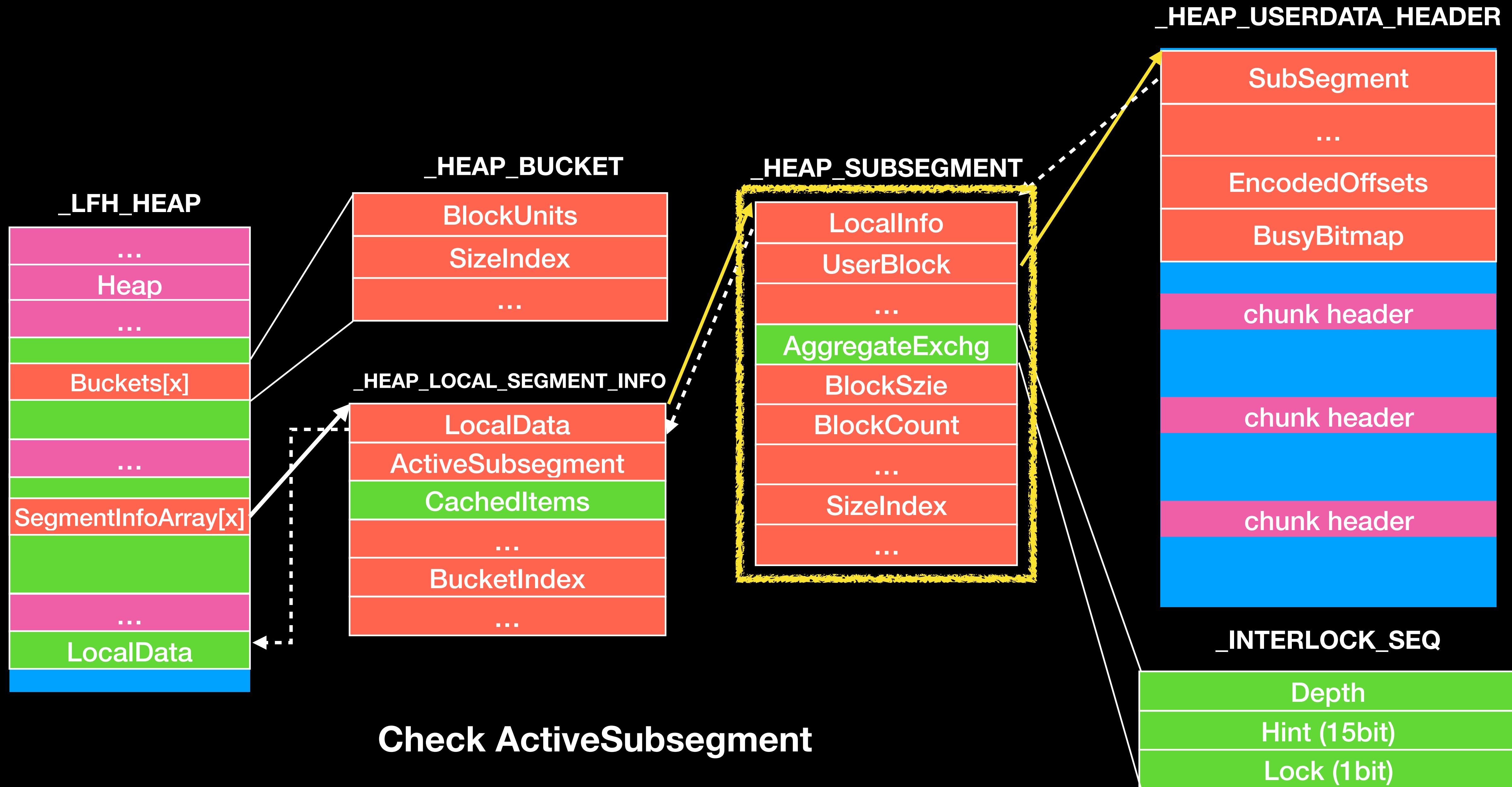
# Windows memory allocator

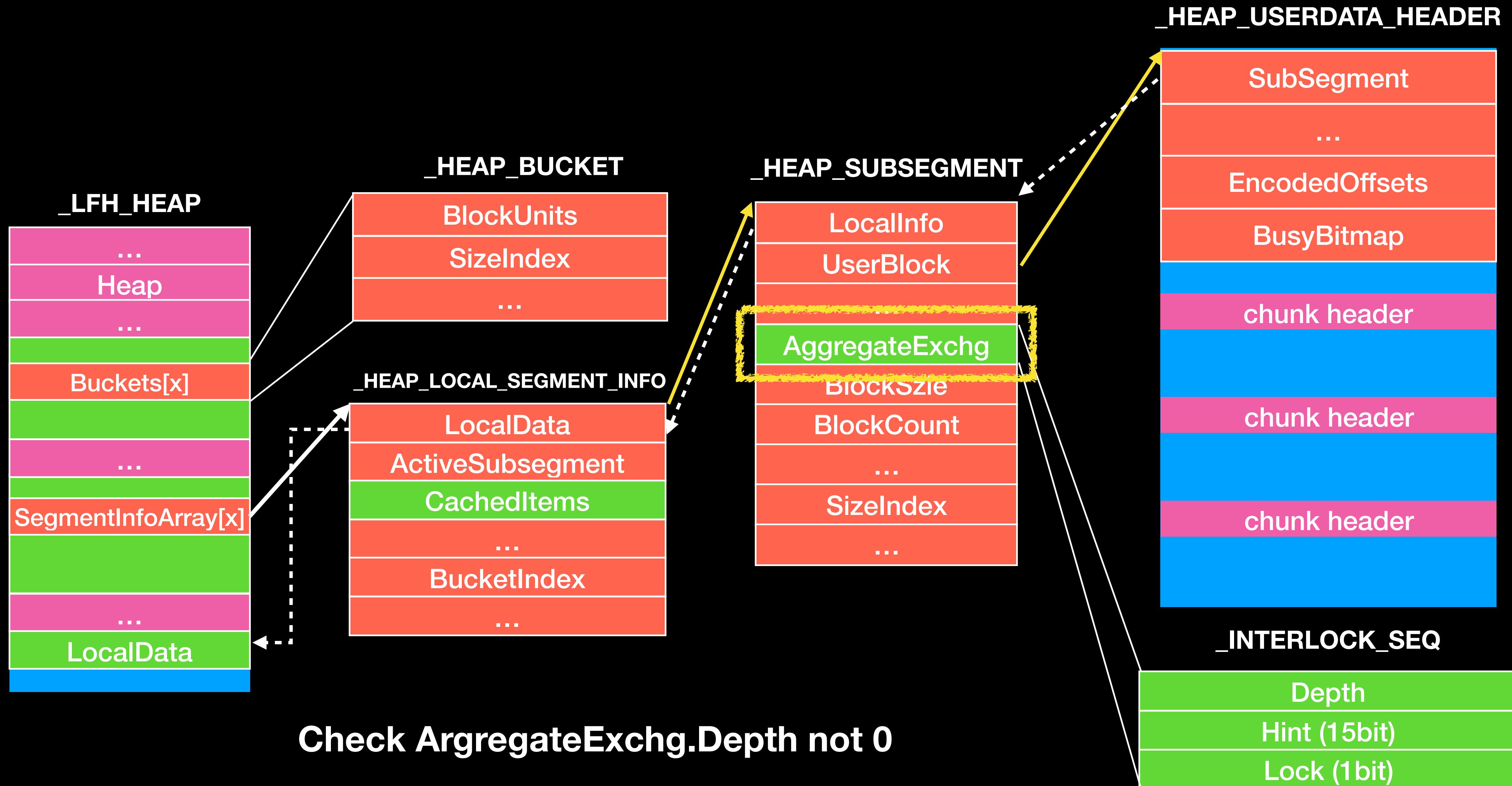
- LFH
  - Allocate (RtlpLowFragHeapAllocFromContext)
    - 先看看 ActiveSubsegment 中是否有可分配的 chunk
    - 從 ActiveSubsegment->depth 判斷
    - 如果沒有則會從 CachedItem 找，有找到的話會把 ActiveSubsegment 換成 CachedItem 中的 SubSegment

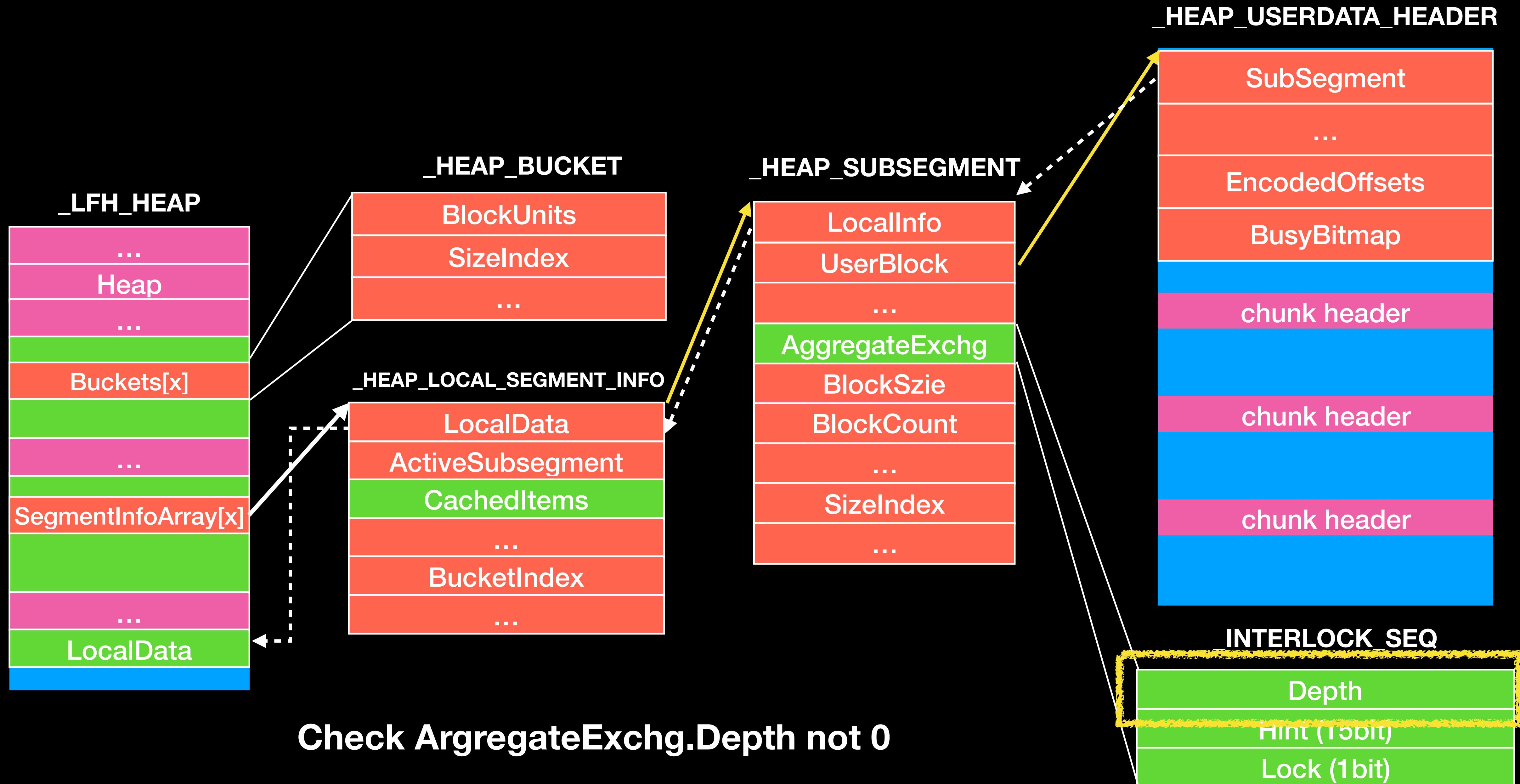


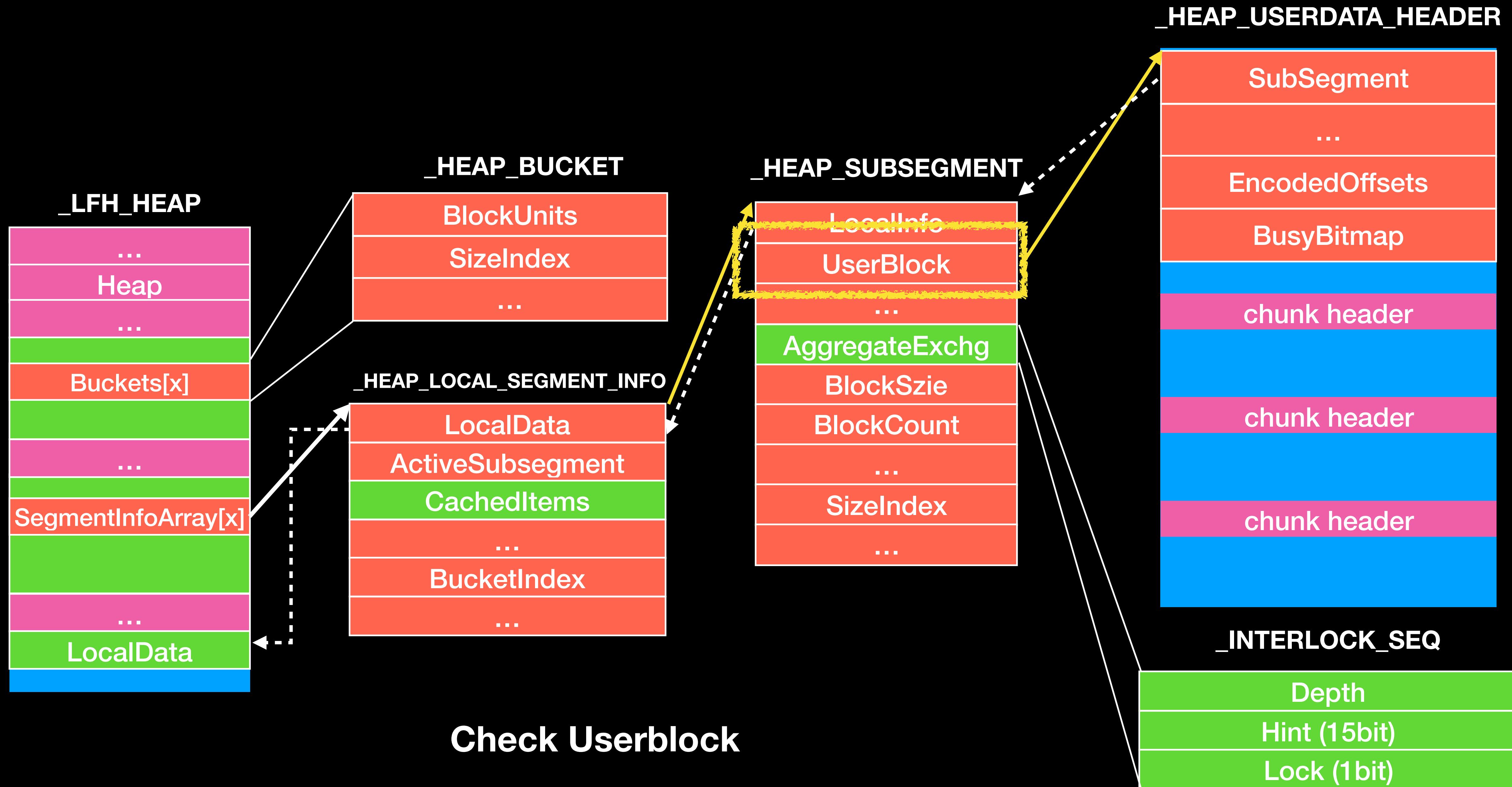












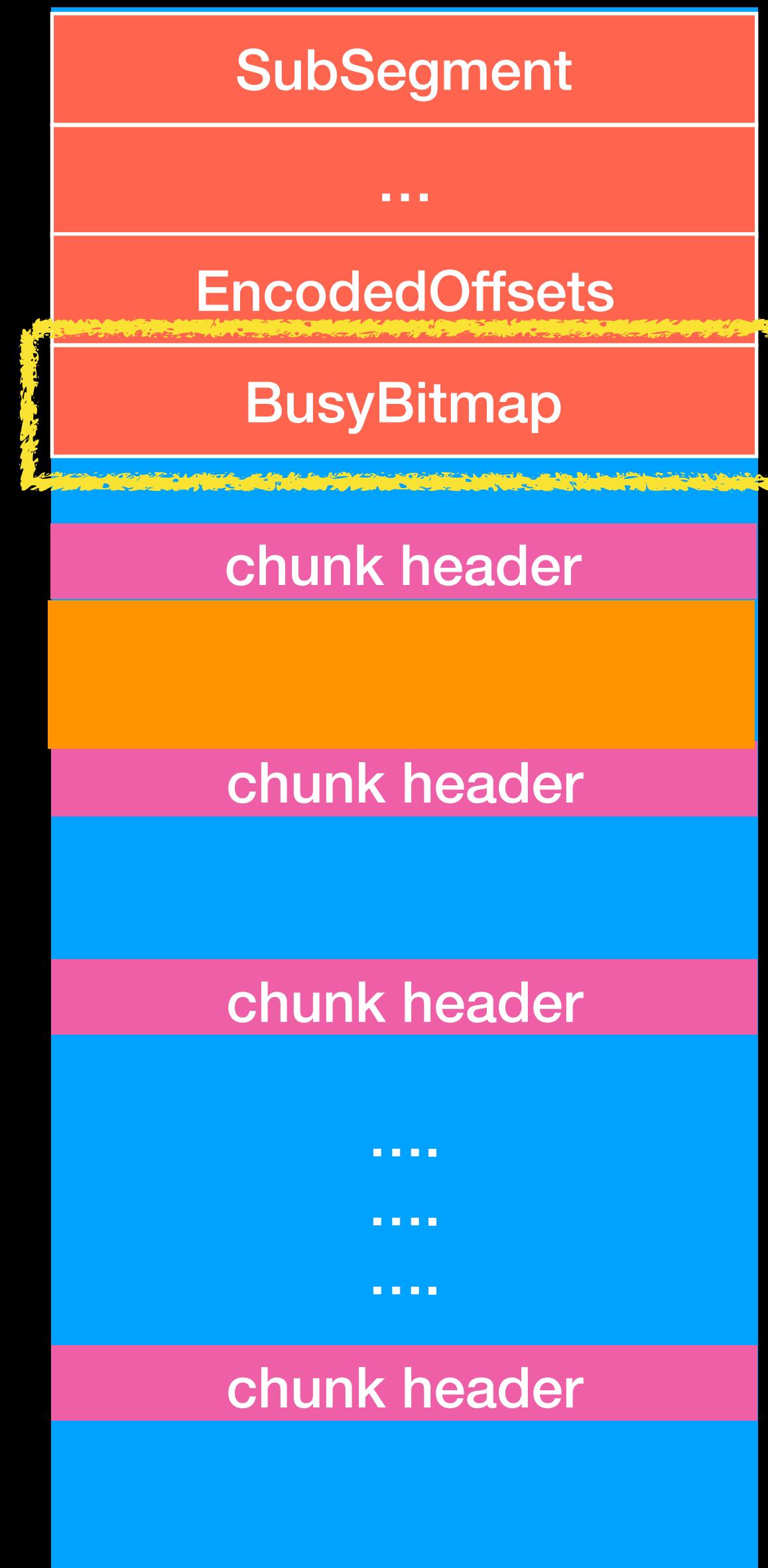
# Windows memory allocator

- LFH
  - Allocate (RtlpLowFragHeapAllocFromContext)
    - 取得 RtlpLowFragHeapRandomData[x] 上的直
    - 下一次會從 RtlpLowFragHeapRandomData[x+1] 取
    - x 為 1 byte , 當下次輪回 x 則會  $x = \text{rand()} \% 256$  繼續取
    - RtlpLowFragHeapRandomData 為一個 random 256 byte 的陣列
      - 範圍為 0x0 - 0x7f

# Windows memory allocator

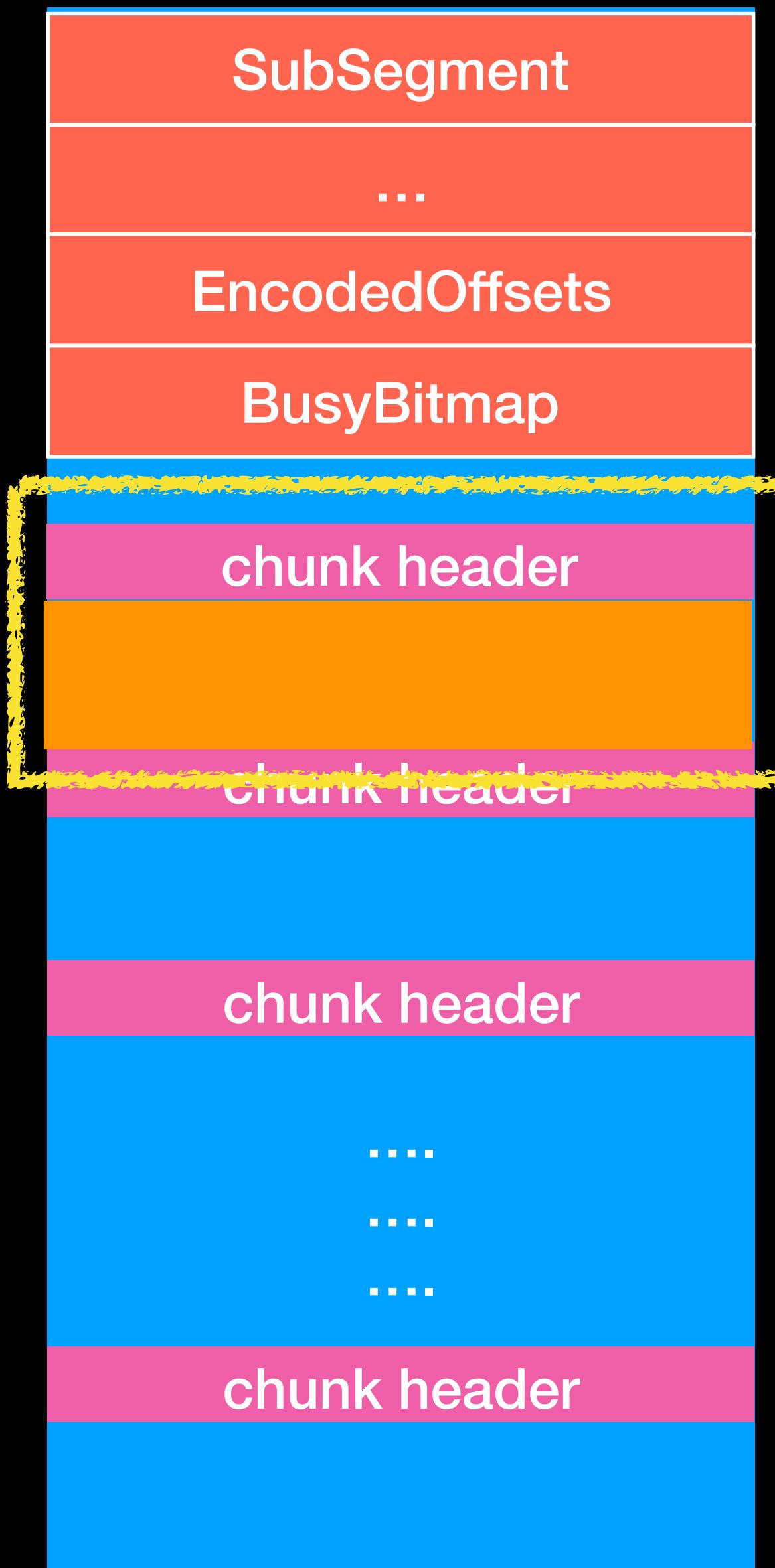
- LFH
  - Allocate (RtlpLowFragHeapAllocFromContext)
    - 最獲得的 index 為
      - RtlpLowFragHeapRandomData[x]\*maxidx >> 7
      - 如果 collision 則往後取最近的
    - 檢查 (unused byte & 0x3f) !=0 (表示 chunk 是 freed)
    - 最後設置 index 及 unused byte 就返回給使用者

## \_HEAP\_USERDATA\_HEADER



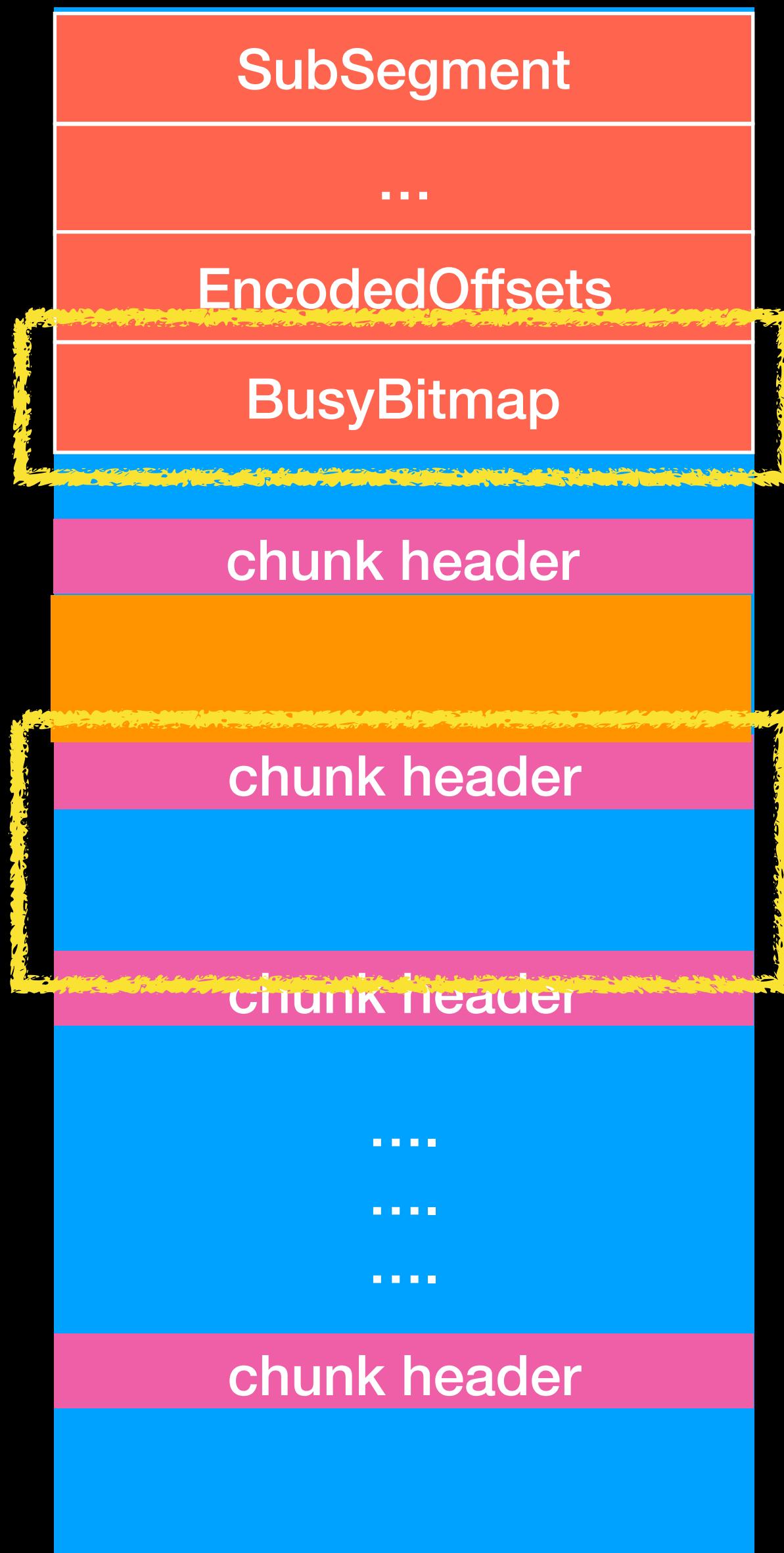
- Get an index
  - `RtlpLowFragHeapRandomData[x]*maxidx >> 7`
  - 檢查該 index 對映到的 bitmap 是否為 0
    - 是 0 則返回對映的 bitmap
    - 非 0 擇取最近的下一個

## \_HEAP\_USERDATA\_HEADER



- 算出來的 index 發現已用過
  - 會取下一個
  - 並設置對映的 bitmap

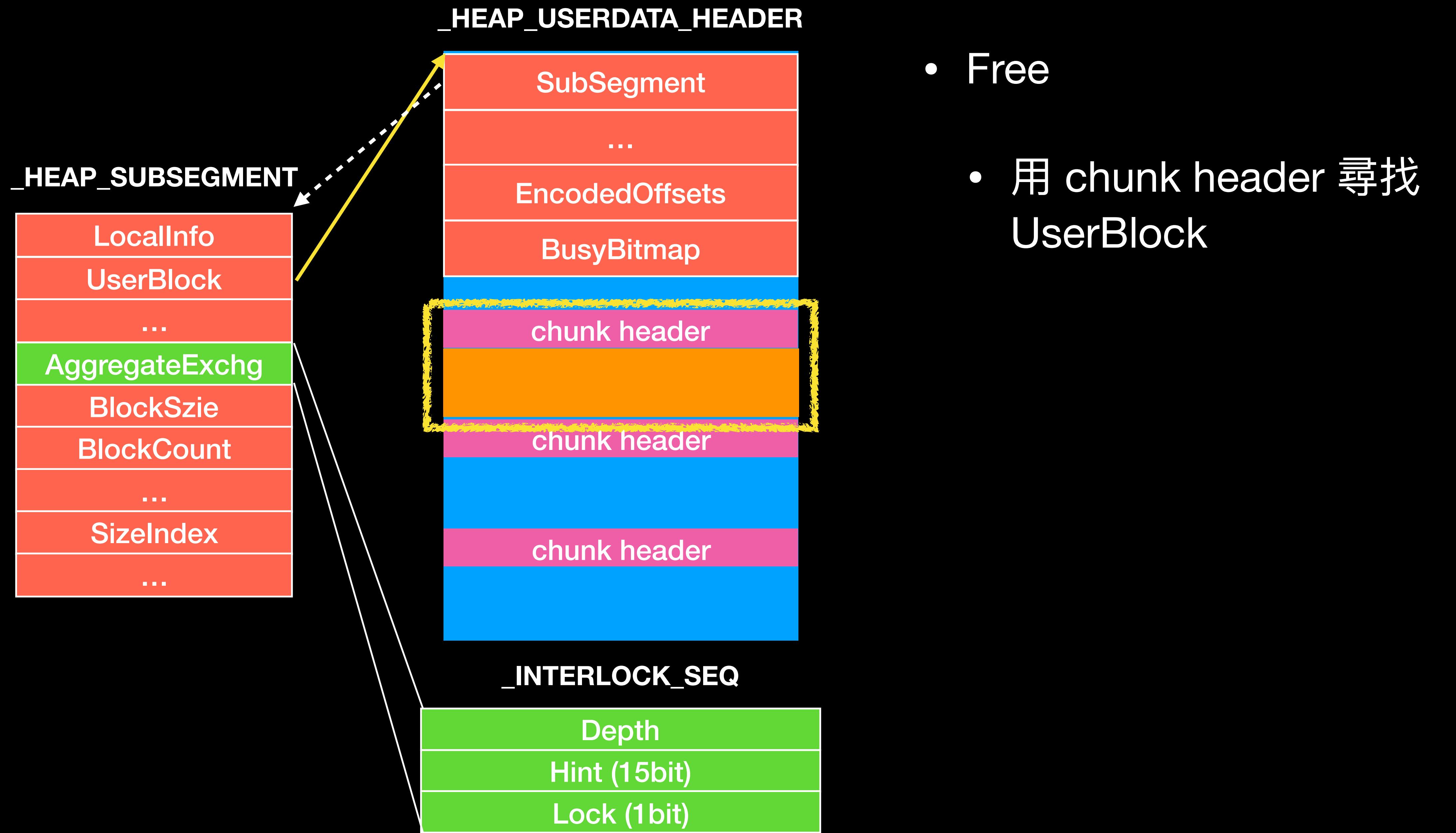
## \_HEAP\_USERDATA\_HEADER

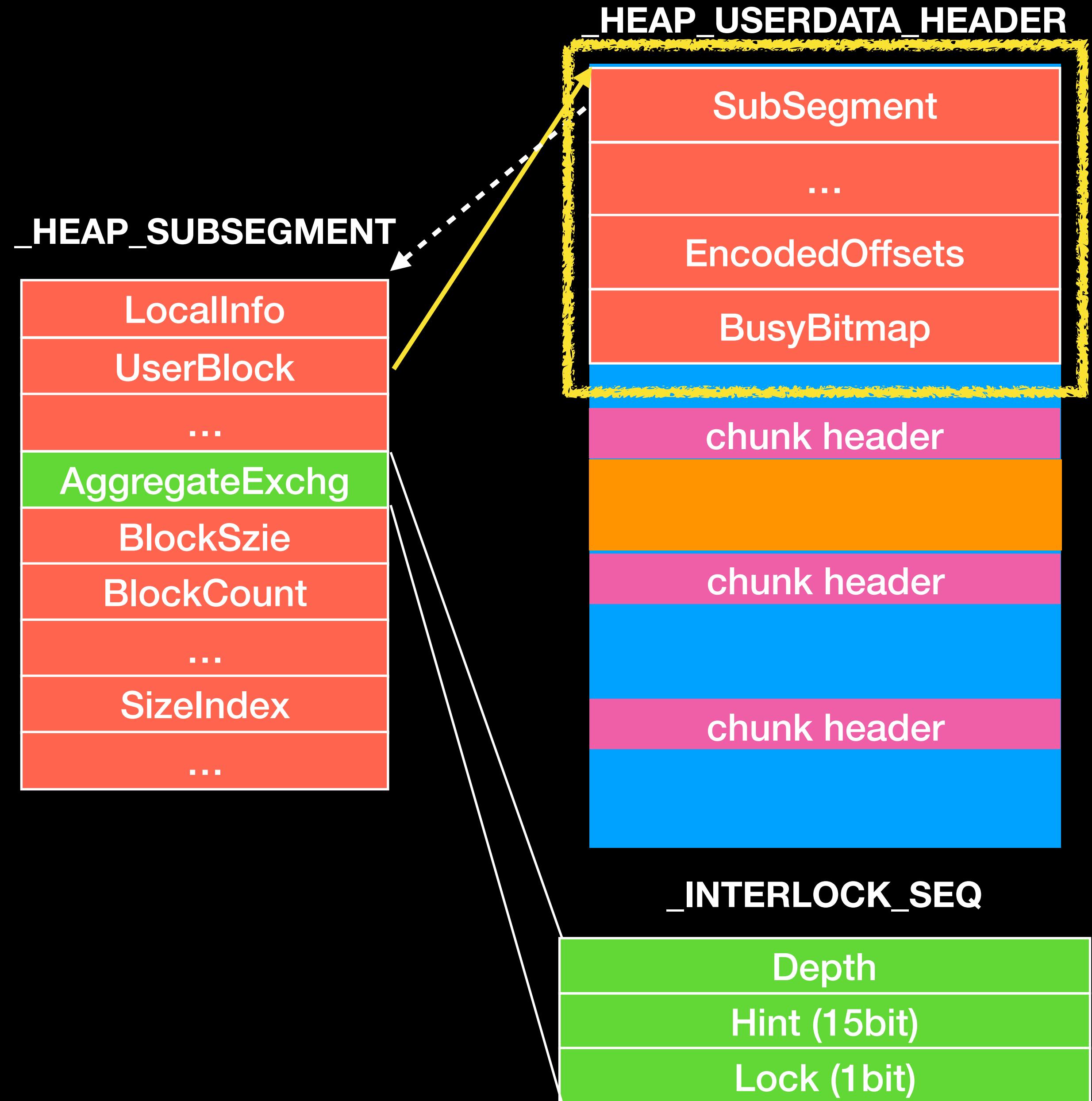


- 算出來的 index 發現已用過
  - 會取下一個
  - 並設置對映的 bitmap

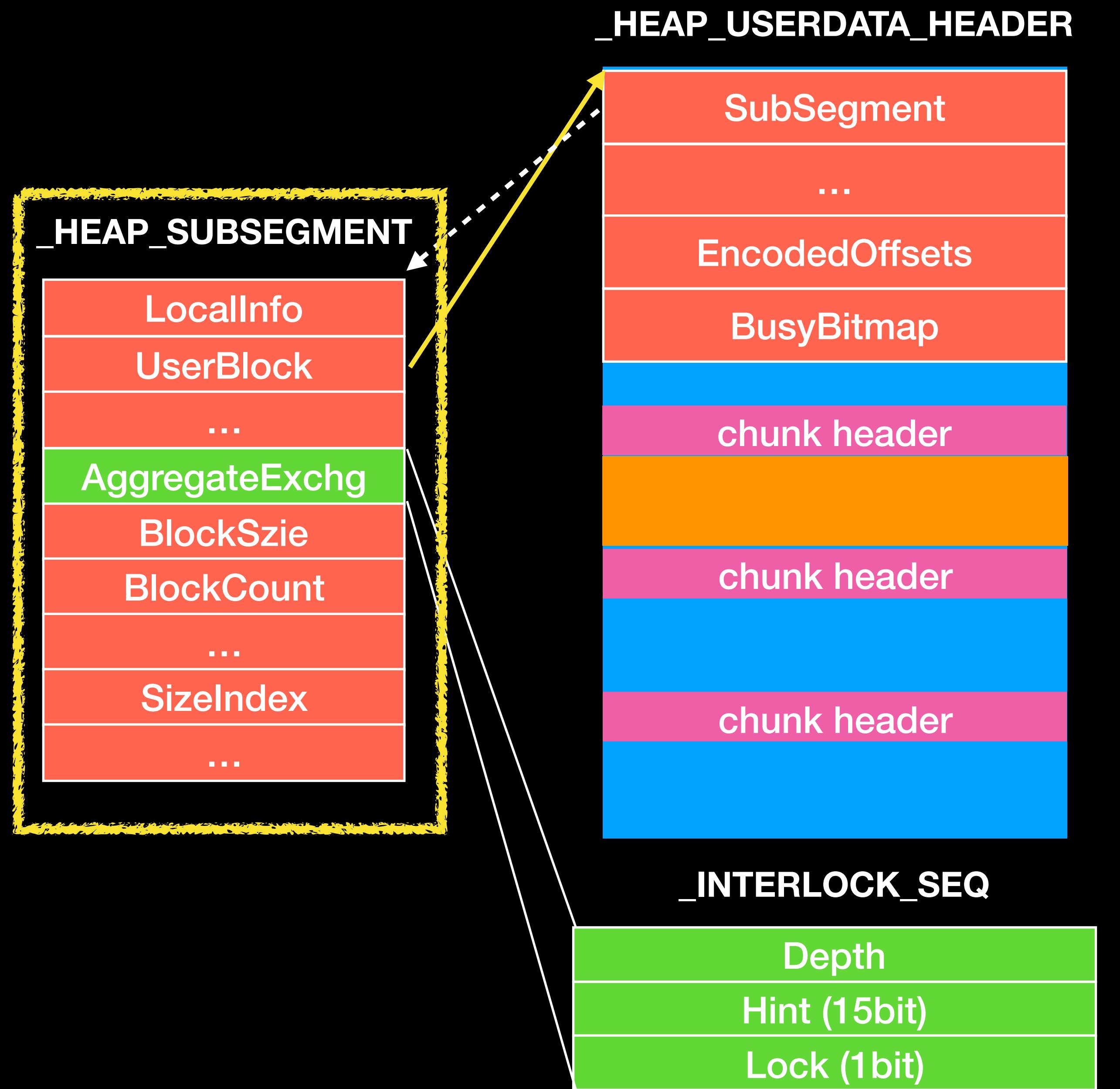
# Windows memory allocator

- LFH
  - Free (RtlFreeHeap)
    - Update unused byte in chunk header
    - Find the index of the Chunk and reset Userblock->BusyBitmap
    - Update ActiveSubsegment->AggregateExchg
    - 如果 Free 的 chunk 不屬於當前 ActiveSubsegment 則會看看能不能放進 cachedItems ，可以放就放進去

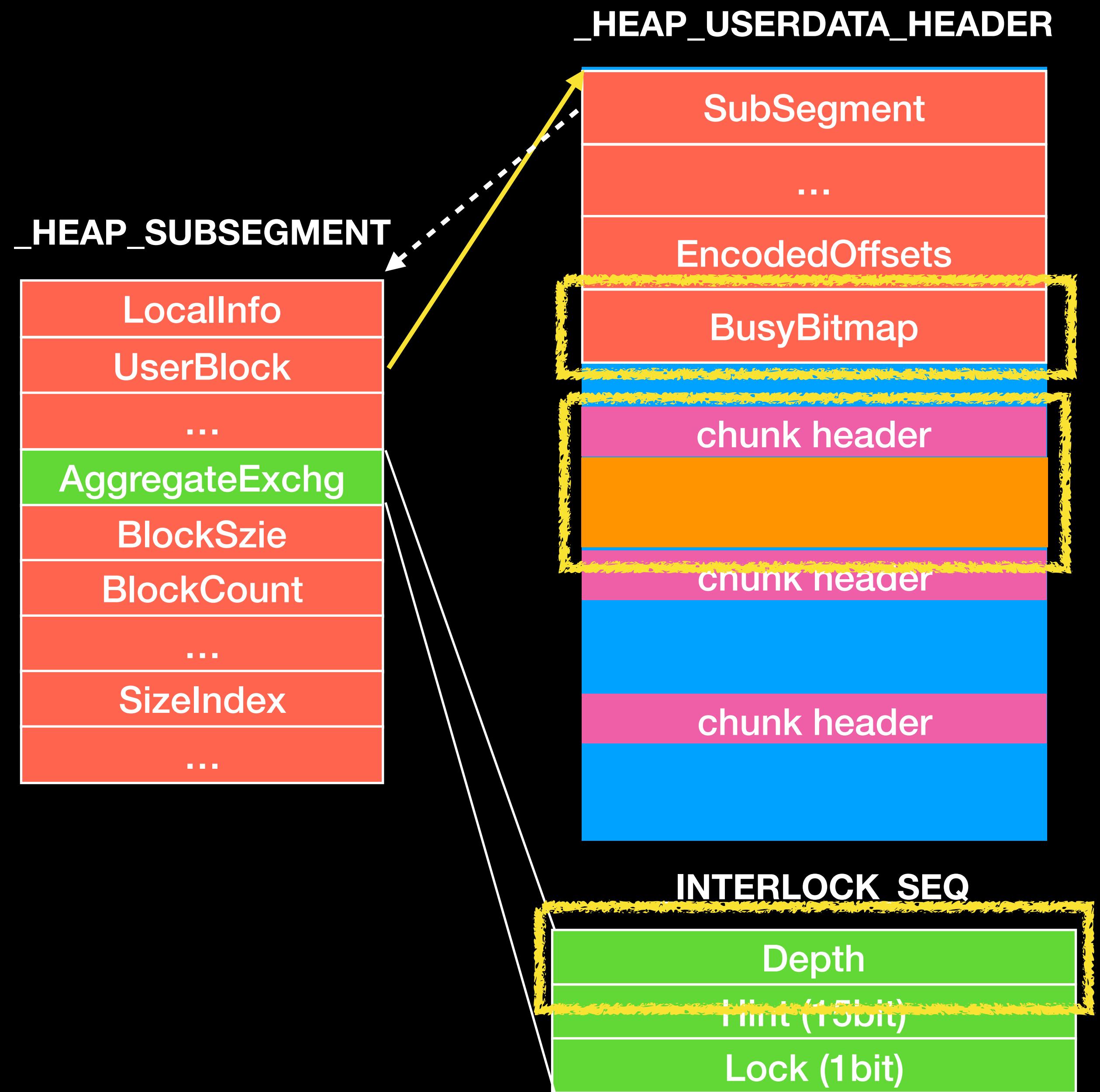




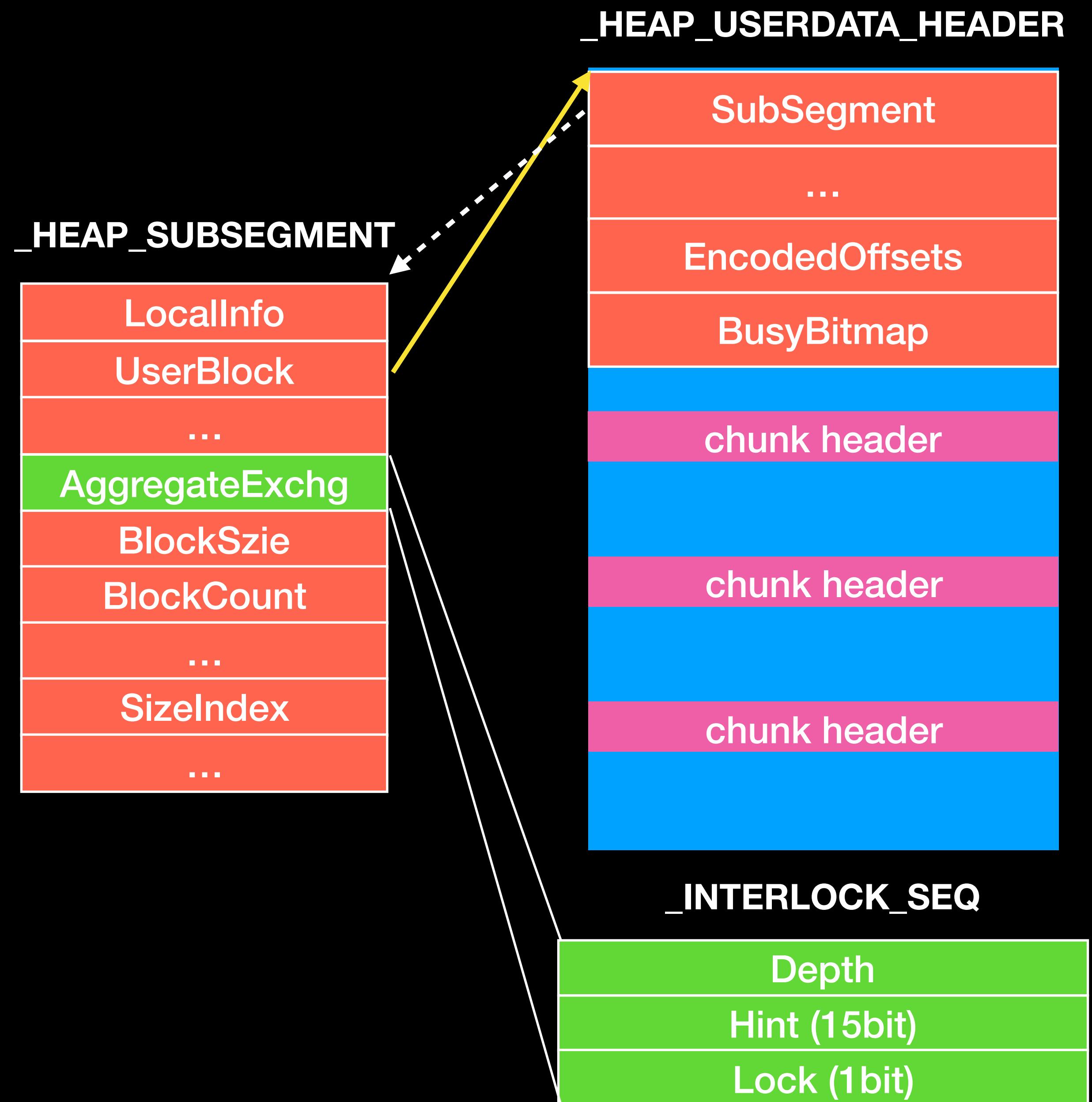
- Free
- 用 chunk header 尋找 UserBlock
- 找回對映的 SubSegment



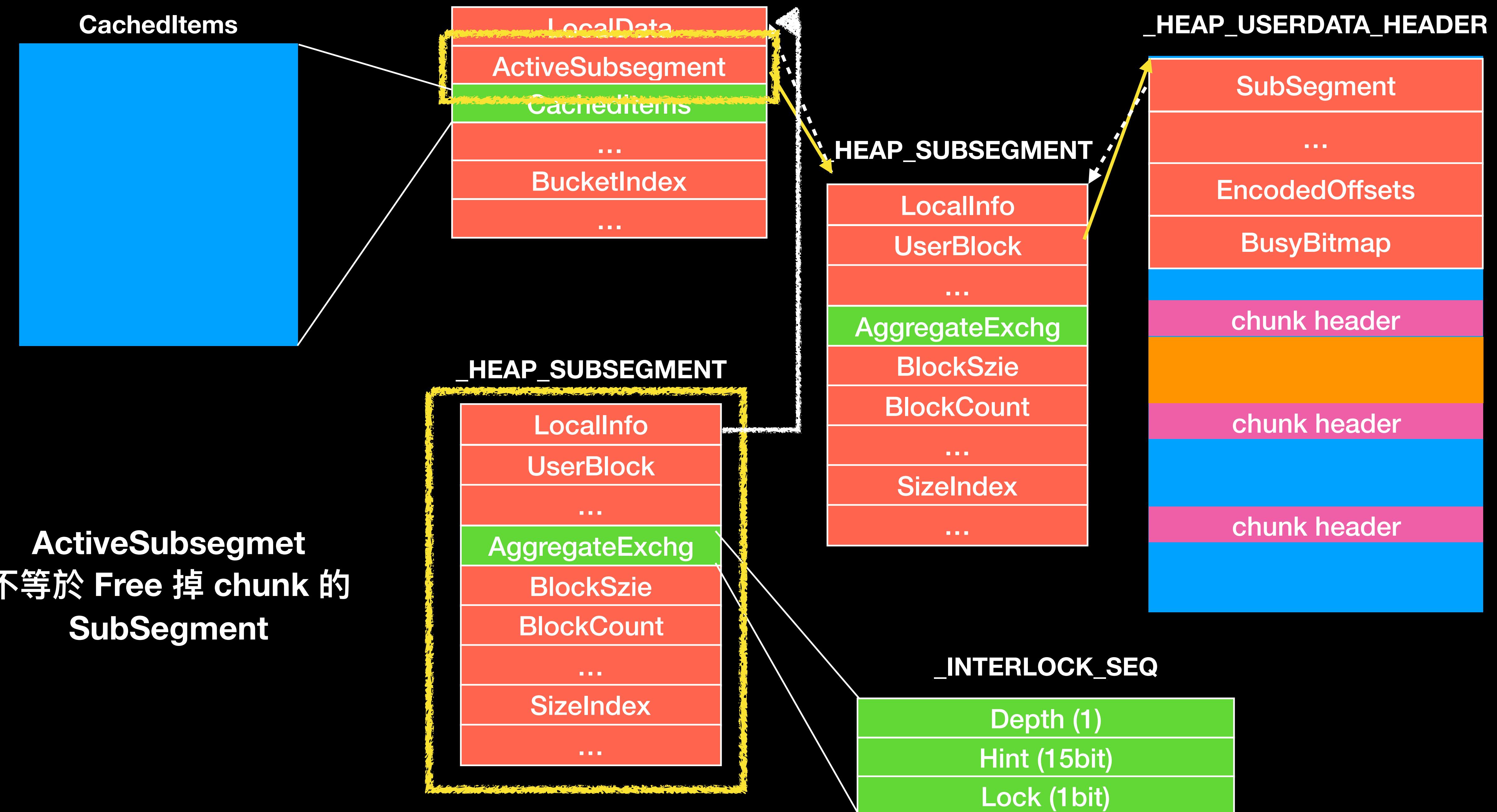
- Free
  - 用 chunk header 尋找 UserBlock
  - 找回對映的 SubSegment
  - 設置 unused byte = 0x80
  - 清除對映的 bitmap
  - Update AggregateExchg

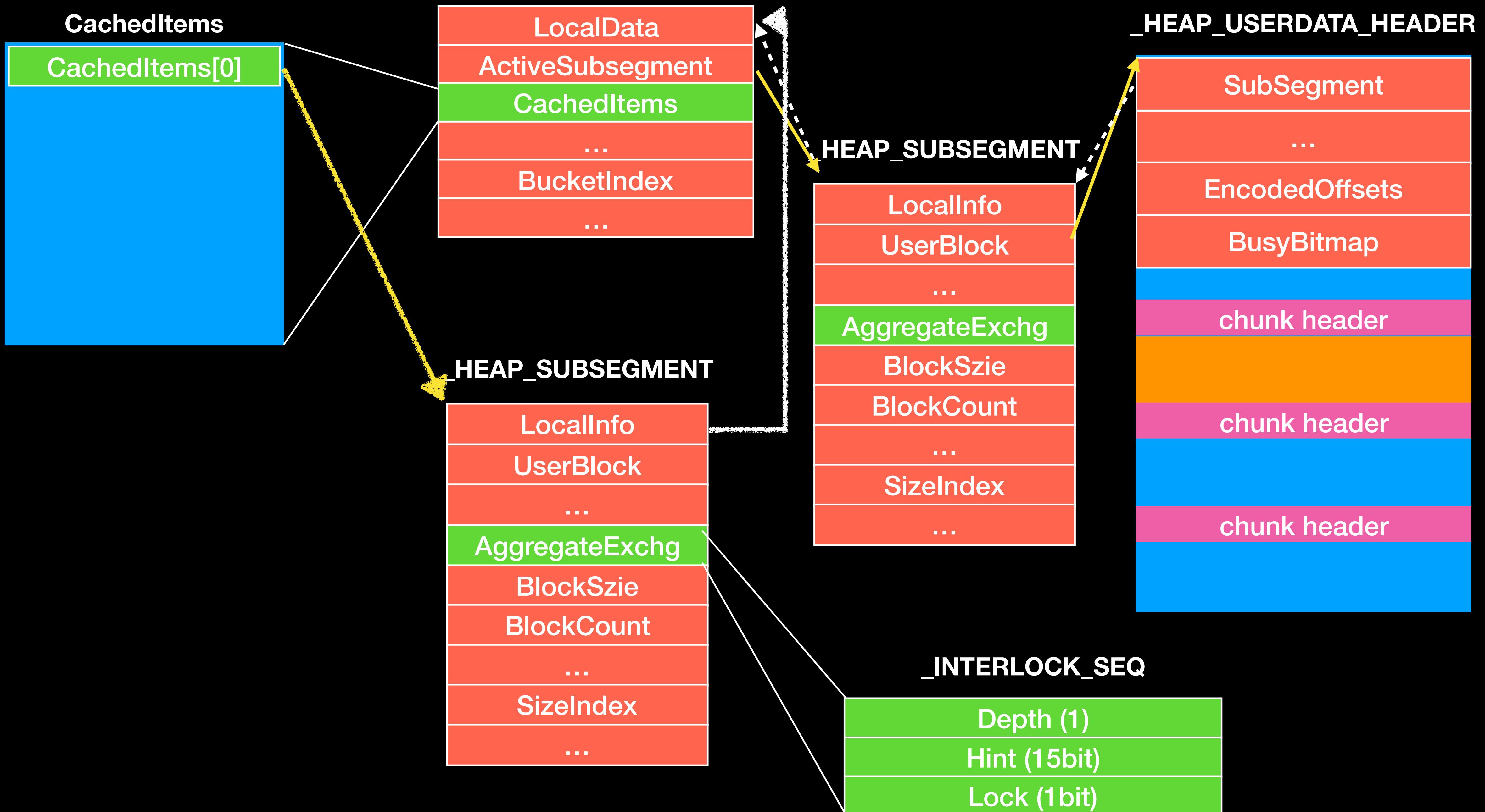


- Free
  - 用 chunk header 尋找 UserBlock
  - 找回對映的 SubSegment
  - 設置 unused byte = 0x80
  - 清除對映的 bitmap
  - Update AggregateExchg



- Free
  - 用 chunk header 尋找 UserBlock
  - 找回對映的 SubSegment
  - 設置 unused byte = 0x80
  - 清除對映的 bitmap
  - Update AggregateExchg



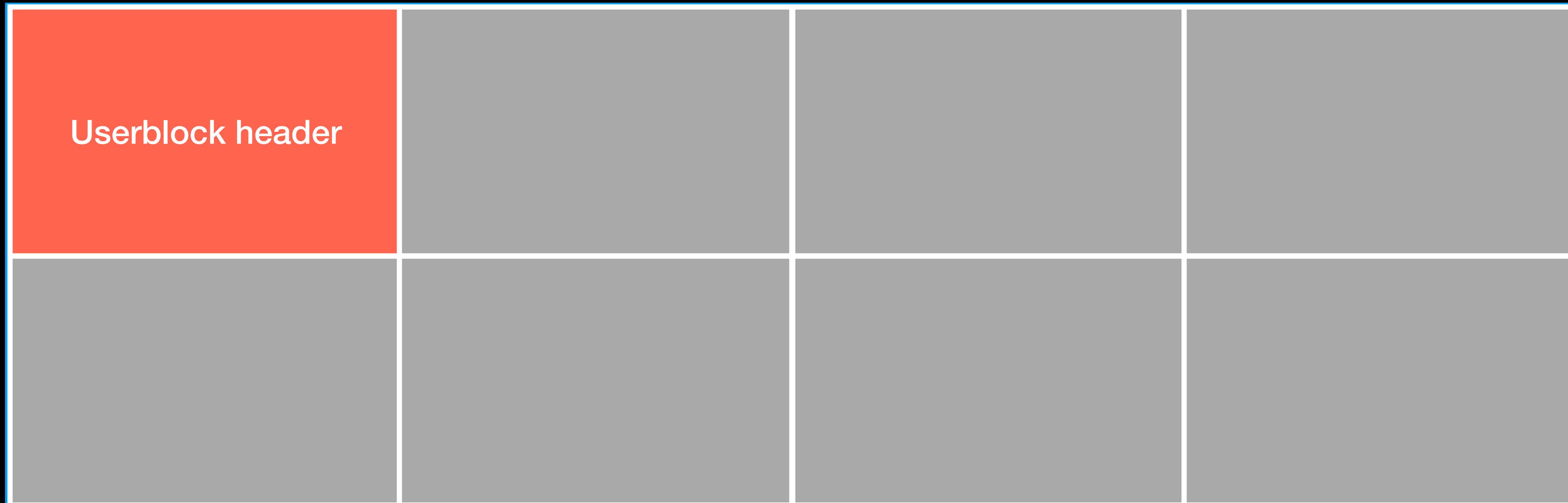


# LFH Exploitation

- Reuse attack
  - 假設我們有 Use after free 的漏洞
  - 但因為 LFH 的隨機性，導致無法預側下一塊 chunk 會在哪，使得我們很難利用
  - 這時可以填滿 Userblock 的方式，在 free 掉其中一塊，那麼下次該 chunk 必定會拿到同一塊，我們可以利用這個特性，拿到 overlap chunk 做進一步利用

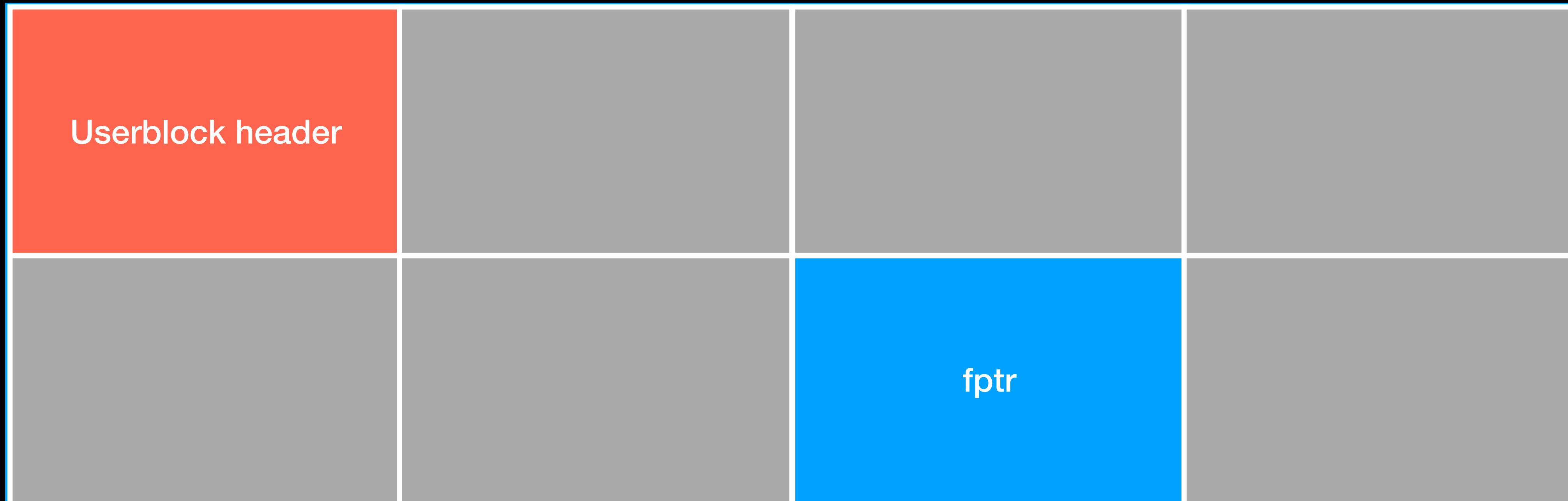
# LFH Exploitation

- normal case
  - `malloc(sizeof(A))`



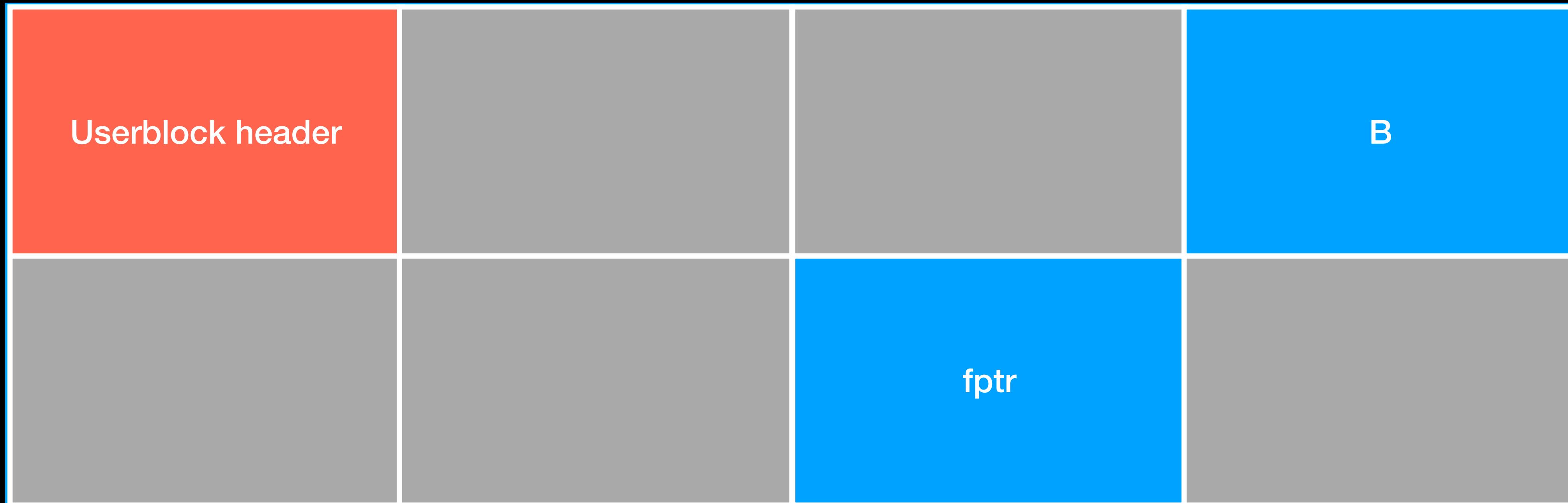
# LFH Exploitation

- normal case
  - `malloc(sizeof(B))`



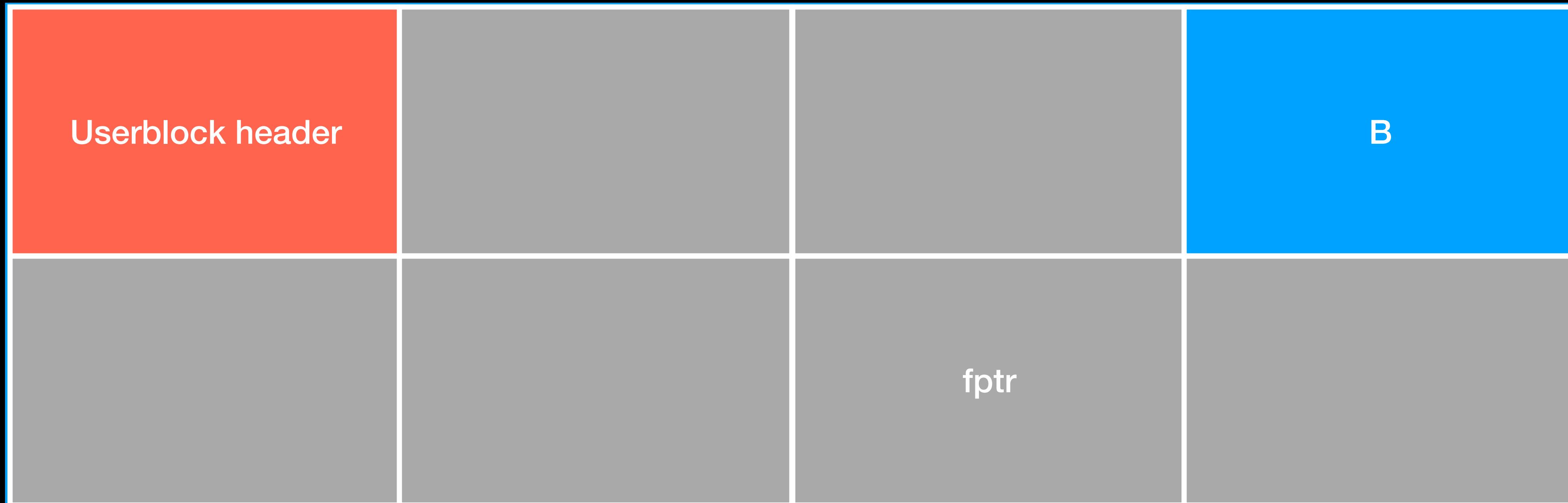
# LFH Exploitation

- normal case
  - free(A)



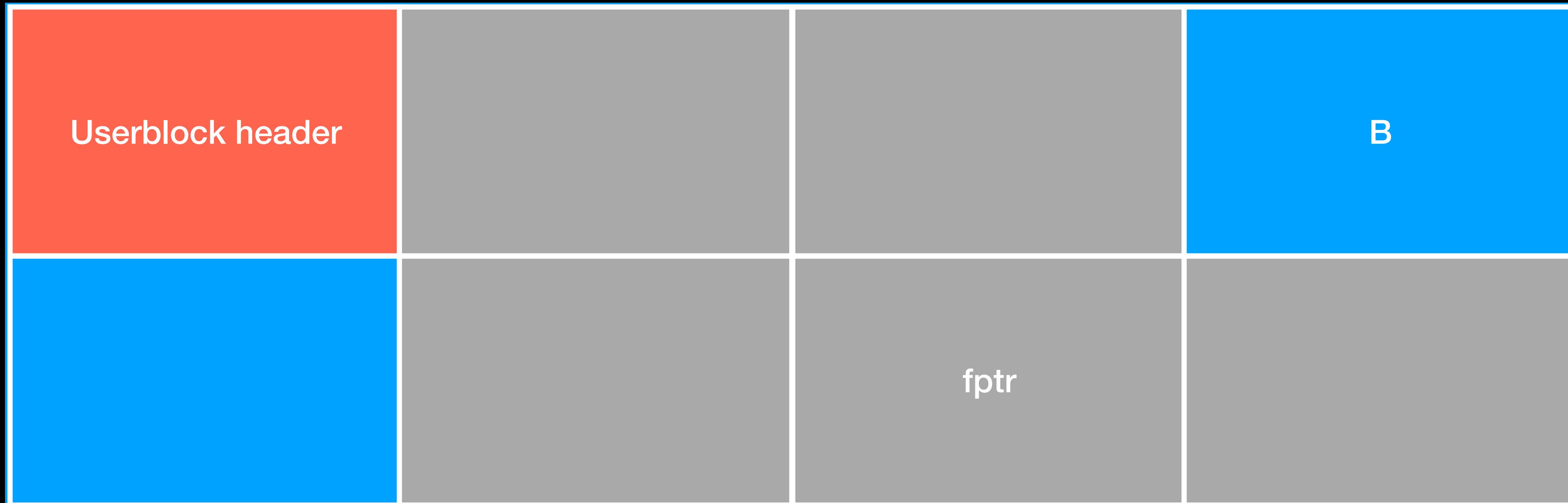
# LFH Exploitation

- normal case
  - `malloc(sizeof(A))`



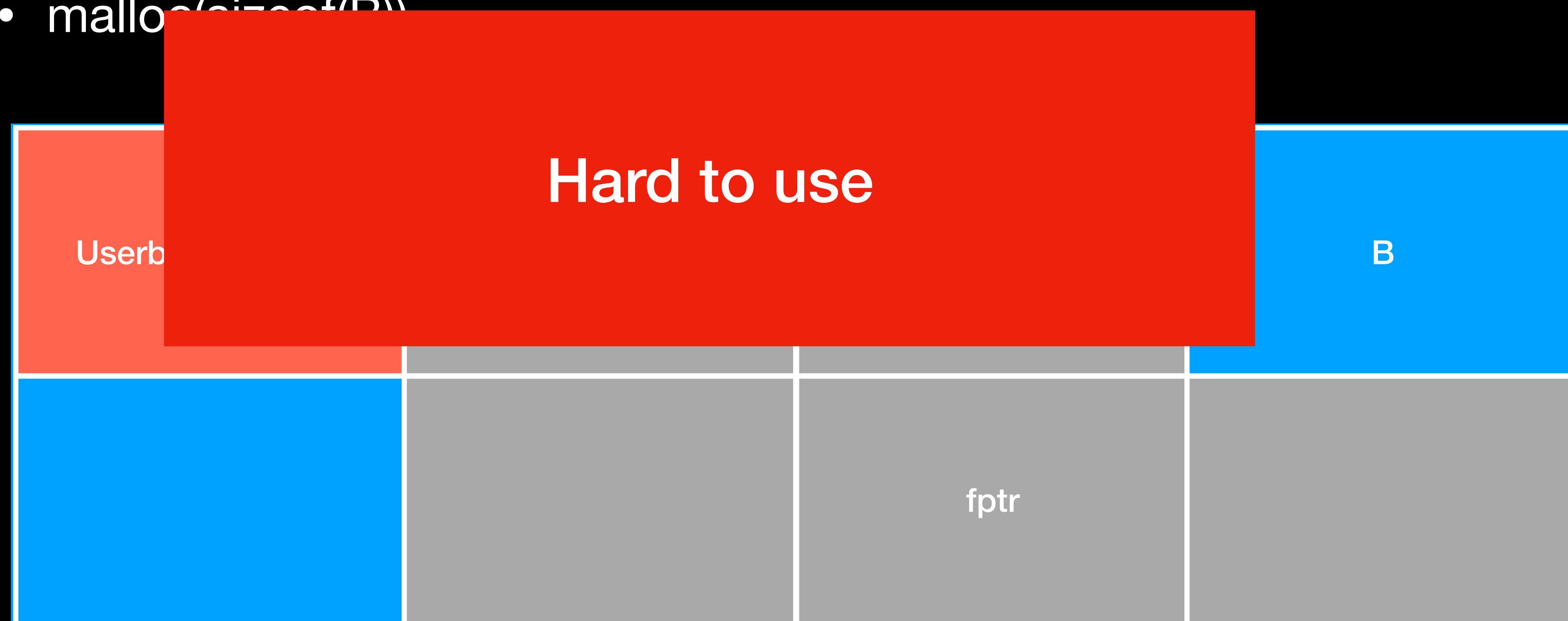
# LFH Exploitation

- normal case
  - `malloc(sizeof(B))`



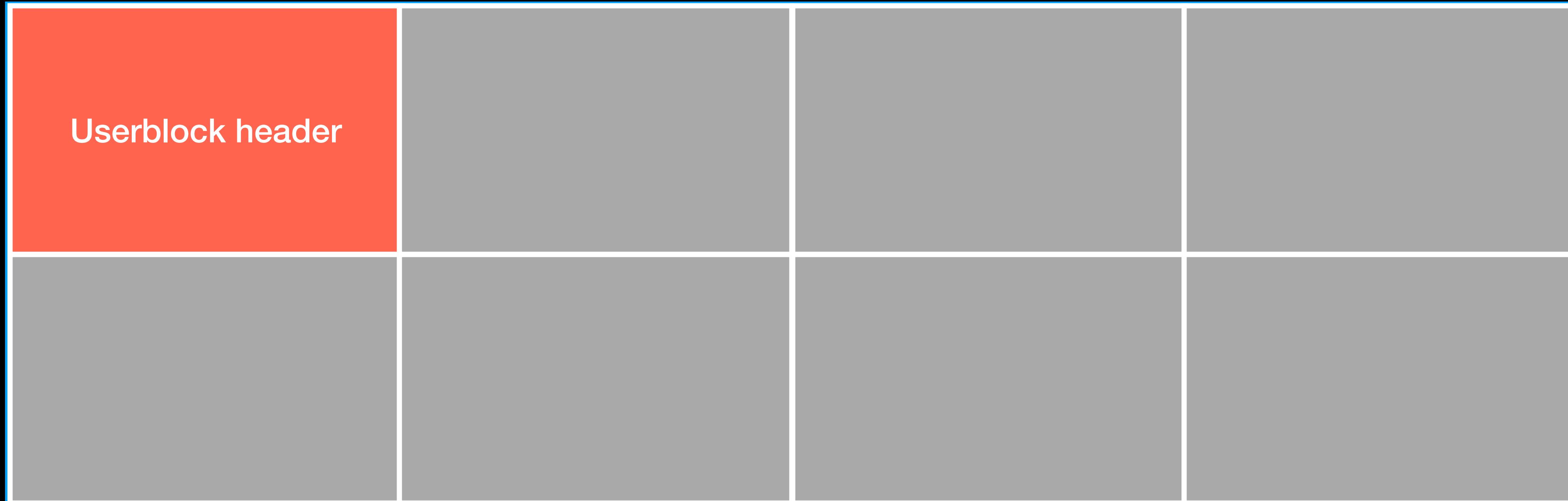
# LFH Exploitation

- normal case
  - malloc(sizeof(B))



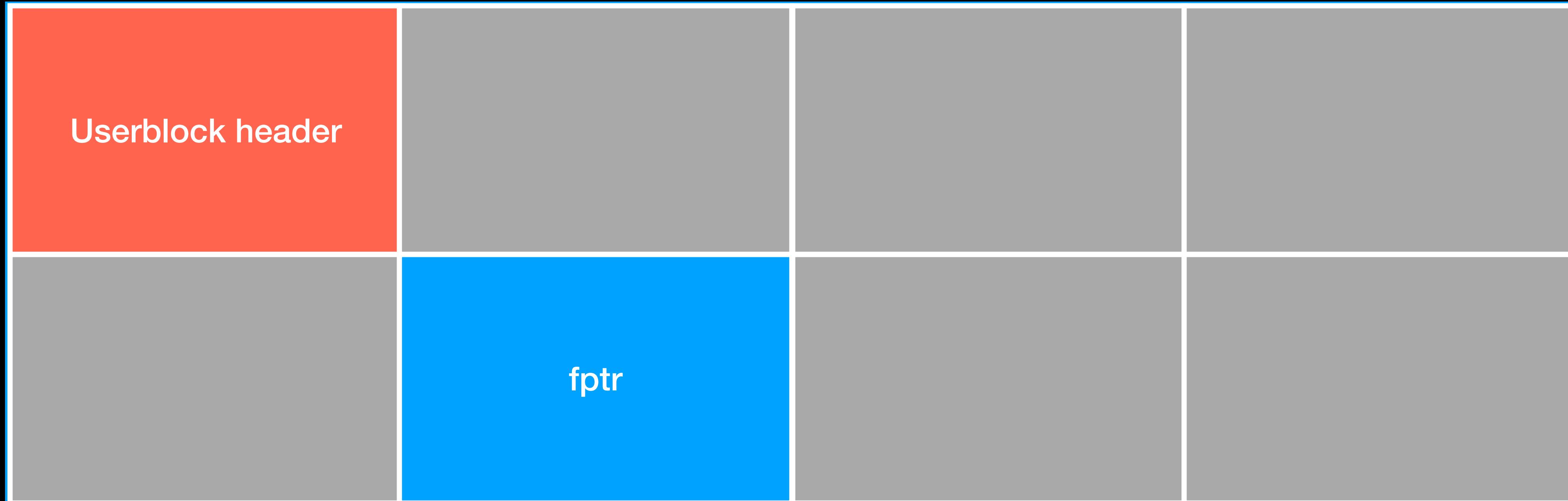
# LFH Exploitation

- Reuse attack
  - `malloc(sizeof(A))`



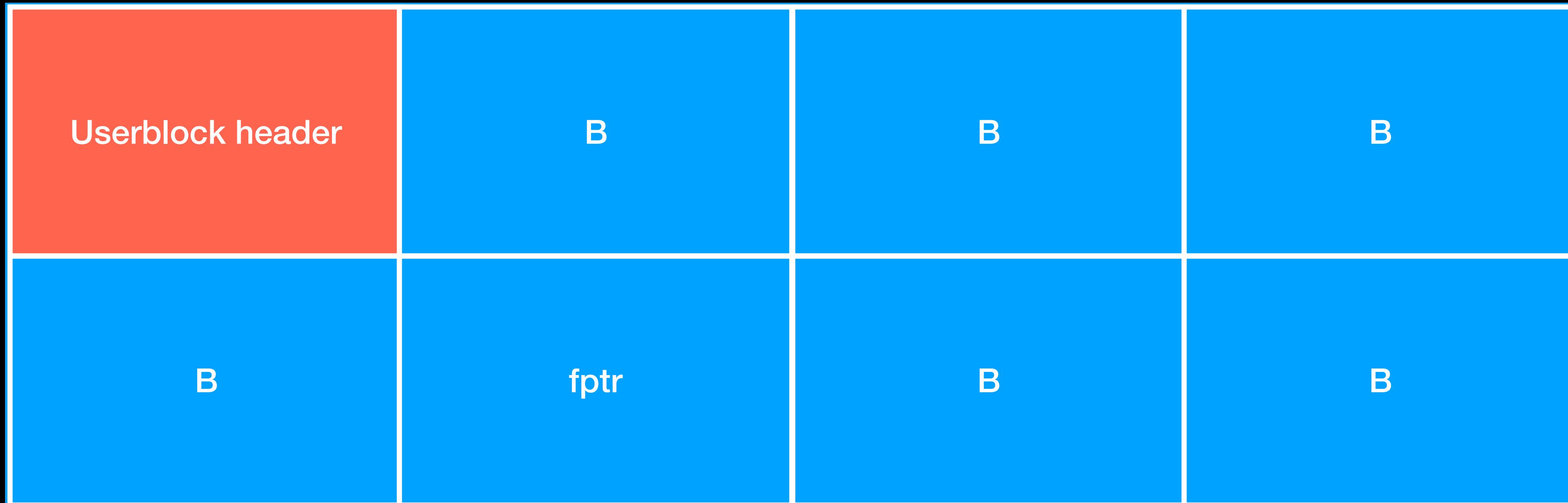
# LFH Exploitation

- Reuse attack
  - `malloc(sizeof(B)) x 6`



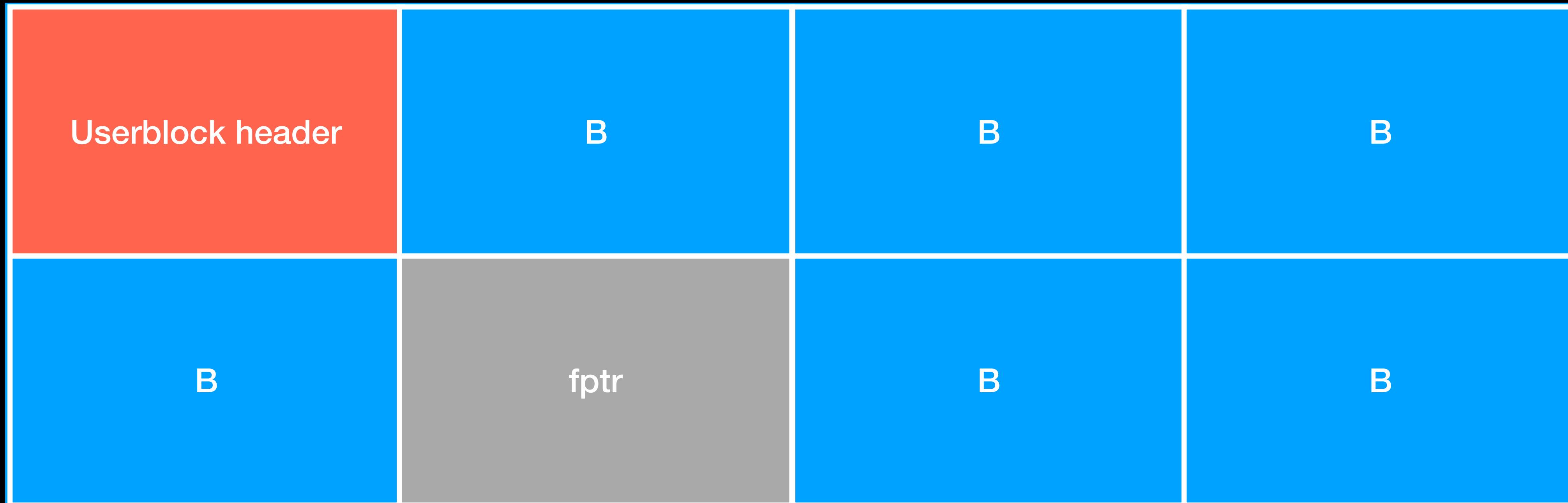
# LFH Exploitation

- Reuse attack
  - `free(A)`



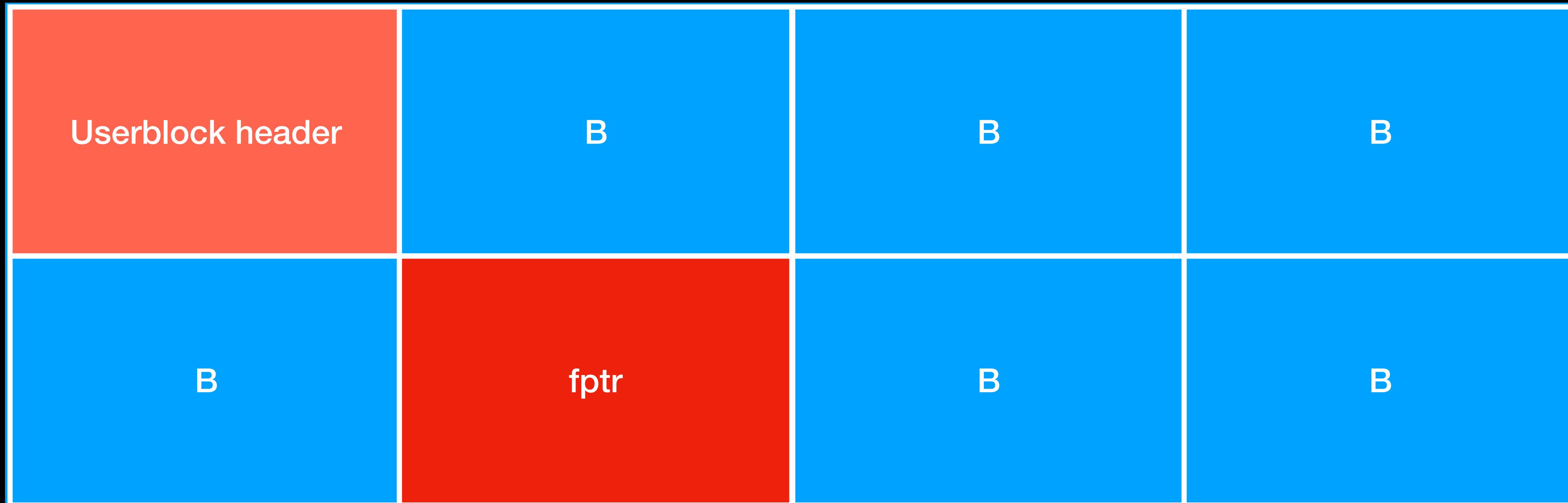
# LFH Exploitation

- Reuse attack
  - `malloc(sizeof(B))`



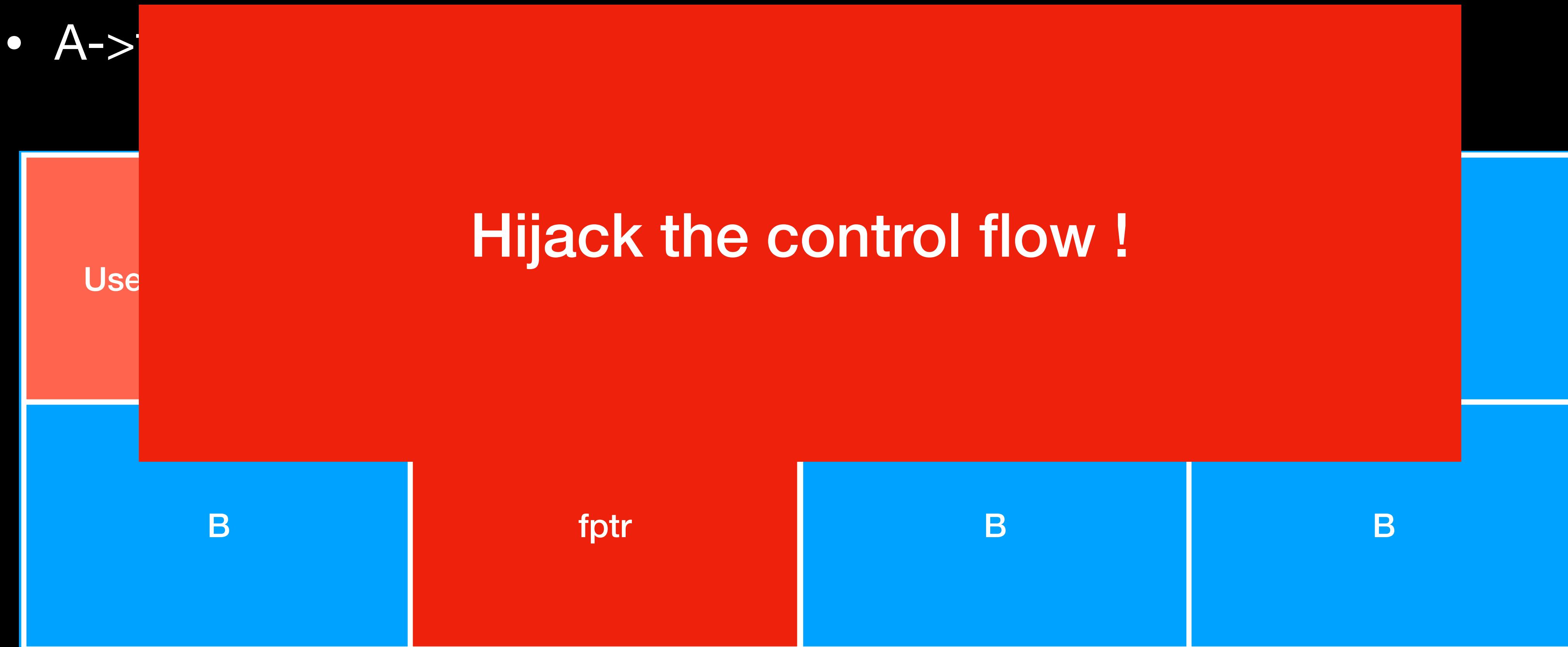
# LFH Exploitation

- Reuse attack
  - A->fptr



# LFH Exploitation

- Reuse attack



# Reference

- [https://github.com/saaramar/  
35C3 Modern Windows Userspace Exploitation](https://github.com/saaramar/35C3_Modern_Windows_Userspace_Exploitation)
- [http://illmatics.com/Understanding the LFH.pdf](http://illmatics.com/Understanding_the_LFH.pdf)
- [https://github.com/saaramar/Deterministic LFH](https://github.com/saaramar/Deterministic_LFH)

# Thank you for listening



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