





From the design to the security assessment

12 January 2022 - Seminari De Cifris Athesis

Giada Sciarretta

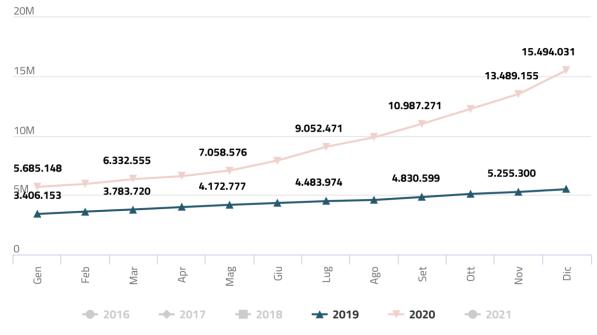


Digital Identity ...especially nowadays



20.269.399

05/05/2021



Andamento mensile delle identità SPID erogate (numero aggregato, totale dei gestori)



Social networks

More than 2 million new users



Digital Report 2021

Internet

We are connected for **over 6 hours a day**Over **1 million new users**



E-commerce

Spent 24% more than in 2019





Digital Identity Clusit Report 2021



La password de la sua carta Flash e stata inserita piu Per proteggere la sua carta abbiamo sospenso il acceso.

Per recuperare il acceso clicca su :

https://www.monetaonline.it/layout/03069/pop/code-ident 31357819513/

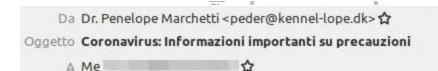
li cart

10% of serious attacks in 2021 are directly related to Covid-19

False e-mail dall'Organizzazione Mondiale della Sanità per il Coronavirus, occhio alla truffa

Nel testo della e-mail viene annunciato un documento con tutte le precauzioni per evitare il contagio da Coronavirus e si invita ad aprire il documento allegato

Increase of phishing, cloned sites, malware ...



Gentile Signore/Signora,

A causa del fatto che nella Sua zona sono documentati casi di infezione dal coronavirus, della Sanità ha preparato un documento che comprende tutte le precauzioni necessarie con coronavirus. Le consigliamo vivamente di leggere il documento allegato a questo messaggi

Distinti saluti, Dr. Penelope Marchetti (Organizzazione Mondiale della Sanità - Italia)





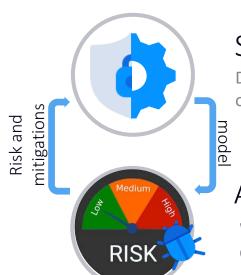
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FBK CyberSecurity Center – Security & Trust Our mission



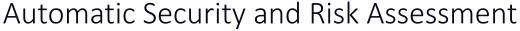


Need for secure solutions and analysis tools for the protection of our data



Security-by-design

Design of on-boarding, authentication and authorization solutions relying on innovative standards and considering the regulatory framework.



With the goal of automatically detecting possible threats and associating each successful attack with the related risk and mitigations.



Public Administration













Our experience DigiMat-Lab and F&C

Joint work with *Poligrafico e Zecca dello Stato Italiano* (IPZS, the **Italian Government Printing Office and Mint**).

- Shared laboratory *DigiMat Lab* (2017-2020);
- In-house company Futuro & Conoscenza from 2021.



mission

Design and analysis of innovative solutions based on CIE





Applications based on CIE 3.0 Index



Introduction to CIE 3.0

General features and real-world scenarios



CIE 3.0 and cryptography

A focus on the cryptographic features of CIE 3.0



What about security?

Our methodology to analyze protocols based on CIE 3.0





Carta d'Identità Elettronica – CIE 3.0 What is it?







A modern identification document: the ICAO MRTD application, containing the holder personal data, photo of the face and image of two fingerprints, is compliant with the ICAO specifications for travel documents

Replace paper-based version

CARDS ACTIVATED BY NOW 30/12/2021

25.768.740







Carta d'Identità Elettronica – CIE 3.0 Physical identity proofing



Personal data and photo of the owner are:

- printed on the plastic surface
 - visual security elements such as holograms and special inks prevent counterfeiting
- stored on the CIE chip can be accessed through NFC (Near Field Communication) interface
 - personal data cannot be accessed without the owner's knowledge: access key printed on the document (CAN – Card Access Number or MRZ – Machine Readable Zone).

Fingerprints may be processed only to control authorities with specific authorizations









Carta d'Identità Elettronica – CIE 3.0 Physical identity proofing



The use of eDocuments allow for



Automatic identity verification processes



Improving police checks security and efficacy

the reading and verification process of the personal and biometric data contained in the microchip verifies the authenticity of the document and the identity of the owner







ID.E.A – Identity Easy Access

An app to verify document's authenticity

Accessing data from CIE requires the MRZ, needed to derive the key for mutual authentication

Data can be accessed by interacting with the contactless chip via an NFC interface

.4 24% B 13:21

0104E

Carta d'Identità Elettronica – CIE 3.0 What is it?







proofing

A modern identification document: the ICAO MRTD application, containing the holder personal data, photo of the face and image of two fingerprints, is compliant with the ICAO specifications for travel documents



A tool for accessing services: the ECC (European Citizen Card) IAS application contains keys and X.509 certificates for secure access to online services

Online authentication





Carta d'Identità Elettronica – CIE 3.0 Online authentication

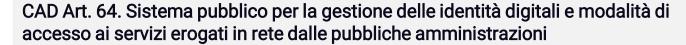




The CIE allows the citizen to authenticate securely to online services of institutions and public administrations



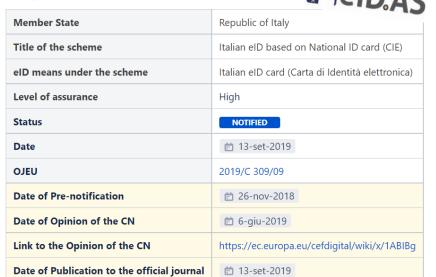
Descreto semplificazione (D.L. 76/2020) prevede l'equiparazione di SPID e CIE e indica il 28 febbraio 2021 quale data per lo switch off delle modalità diverse di identificazione per l'accesso ai servizi online delle pubbliche amministrazioni



2-quater. L'accesso ai servizi in rete erogati dalle pubbliche amministrazioni che richiedono identificazione informatica avviene tramite SPID, nonché tramite la carta di identità elettronica [...]



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CO	ı y			



Regolamento (UE) n. 910/2014





Carta d'Identità Elettronica - CIE 3.0

Online authentication





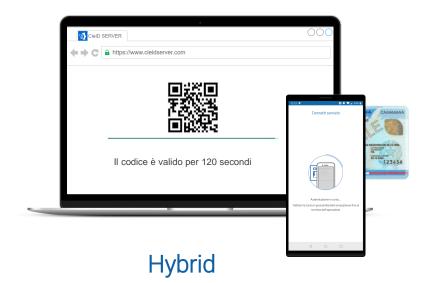
The CIE allows the citizen to authenticate securely to online services of institutions and public administrations



Total desktop



Total mobile







CIE 3.0 – Digital identity and more Index



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A focus on the cryptographic features of CIE 3.0



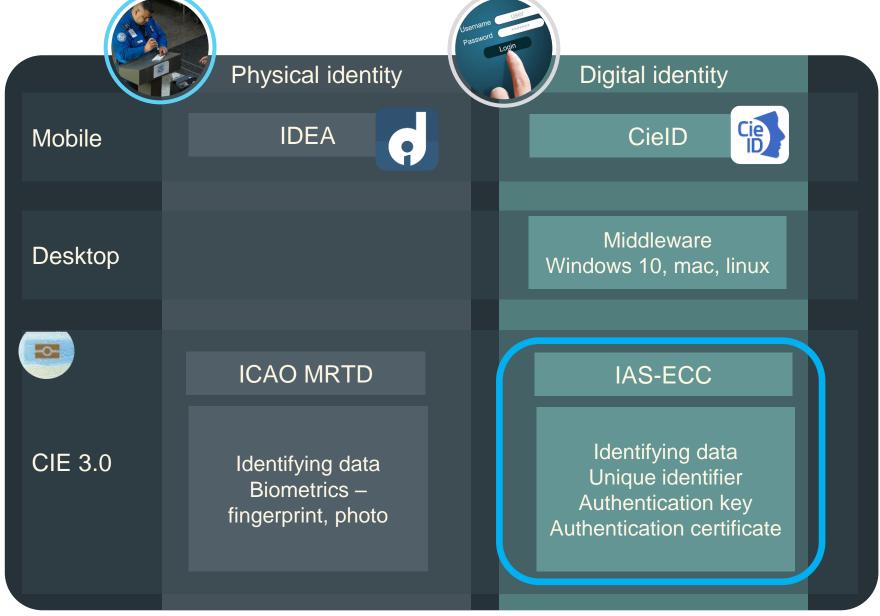
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CIE 3.0 Chip



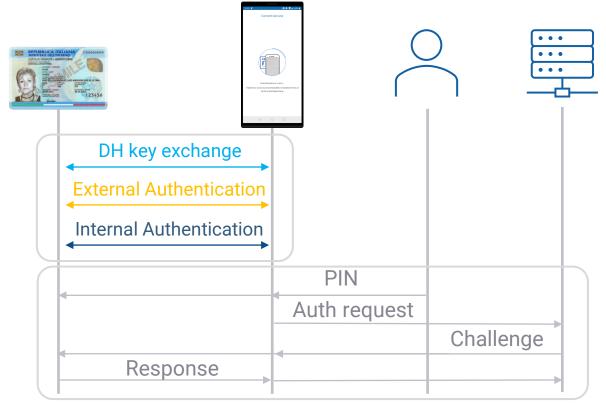




CIE 3.0 IAS-ECC

Online authentication by means of the IAS-ECC involves the following entities and consists two main steps:

- Creation of a secure channel:
 - The creation of a shared secret key between the CIE and the Terminal;
 - The mutual authentication of the peers
- 2. Challenge-Response auth protocol with the unlock of the secret key using the CIE PIN







CIE 3.0

IAS-ECC: DH key exchange



Diffie-Hellmann key exchange lets the two peers to share a secret key used to encrypt all the next exchanged messages \rightarrow protection from eavesdropping attacks





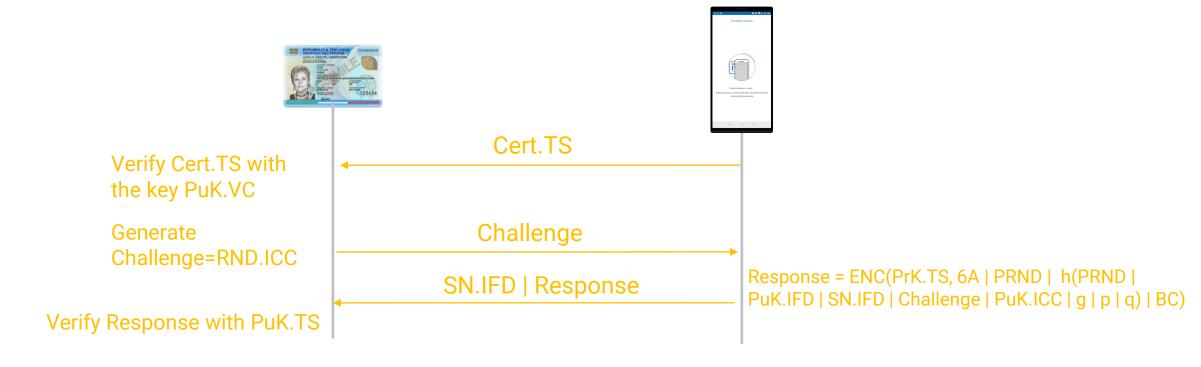


CIE 3.0

IAS-ECC: external authentication



Terminal authenticates with the CIE: is asked to prove its identity, by engaging a challenge-response based mechanism with the CIE







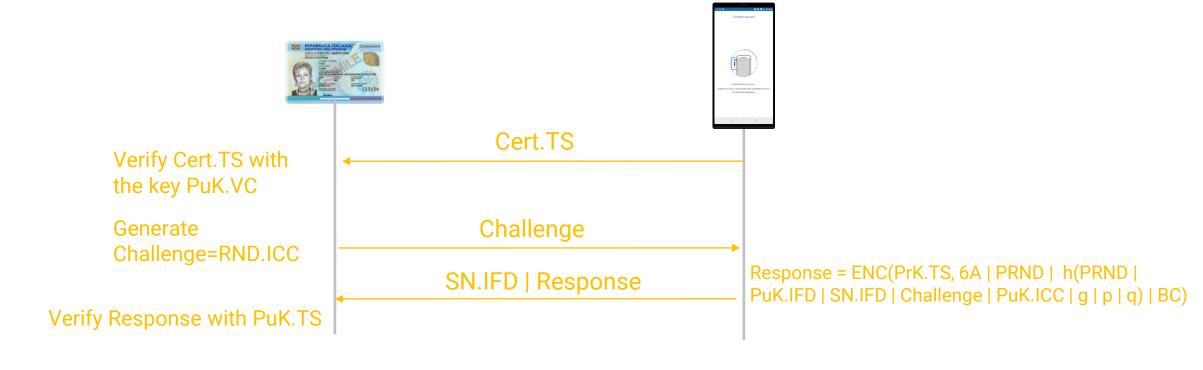
CIE 3.0 IAS-ECC: €

In the Italian implementation of the IAS-ECC, the PKI used to validate the digital certificates for the Terminals has been removed, with the goal of reducing the cost and complexity of the infrastructure and incentivize the diffusion of the CIE as online authentication mechanism.

Terminal authentic

This implies that the Cert.TS send by Terminal to the CIE, actually has not been issued by a CA. More in detail, the pair (PuK.CV, PrK.CV) is made publicly available to all the Terminals, so that they can produce a valid (from the CIE perspective) certificate.

challenge-response paseu mechanism with the cit





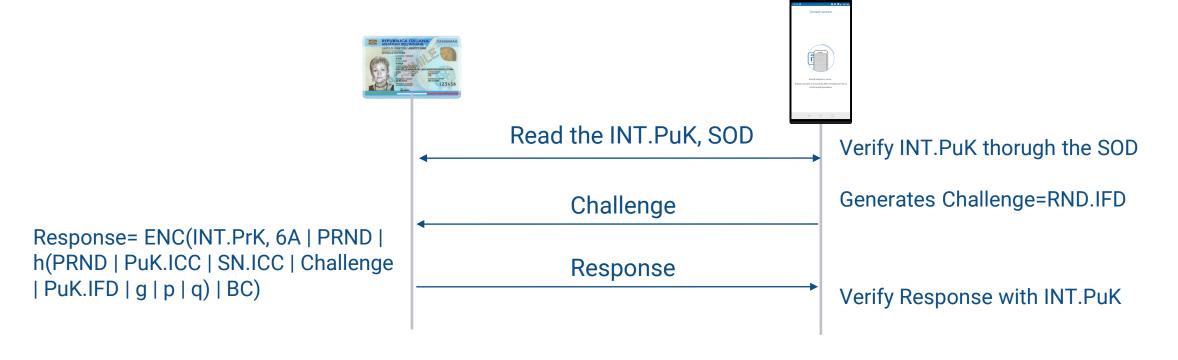


CIE 3.0

IAS-ECC: internal authentication

The CIE authenticates with the Terminal



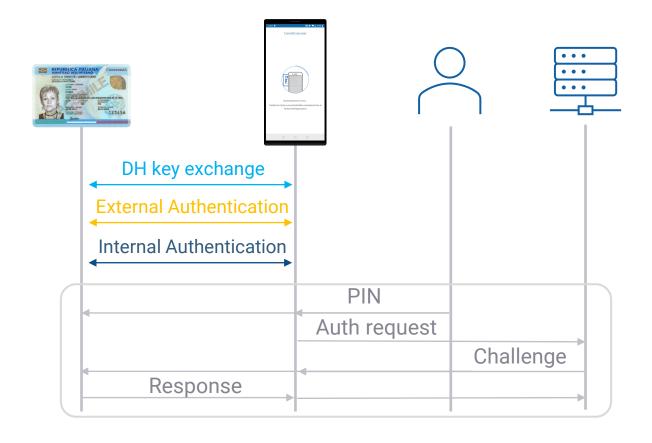






CIE 3.0

IAS-ECC: CR auth protocol



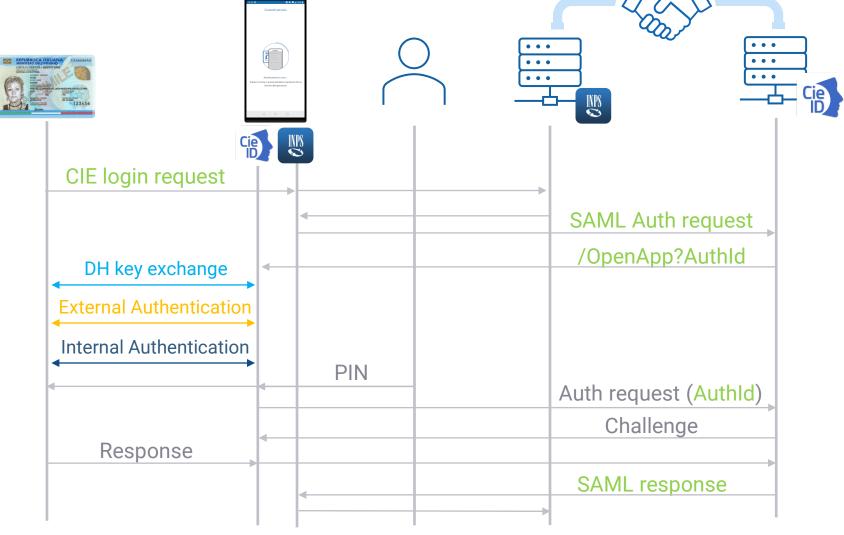




"Entra con CIE 3.0" federation

Cie Entra con CIE

SAML 2.0







CIE 3.0 - Digital identity and more Index



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DigiMat-Lab and F&C Methodology

Functional and security requirements



Security-by-design

Design and implementation of the protocol, based on state-of-the-art techniques



Automatic Security and Risk Assessment

With the goal of automatically detecting possible threats and associating each successful attack with the related risk and mitigations.





CIE 3.0

Online authentication



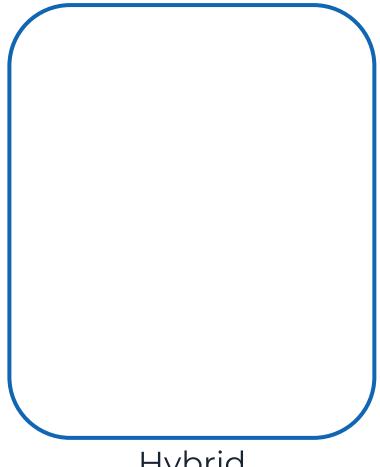








Total-mobile



Hybrid

CIE 3.0

Online authentication











Total-mobile



Hybrid



CIE 3.0 – Hybrid scenario Requirements

Requirements

- 1. Passwordless solution
- 2. SAML 2.0
- 3. No push notifications
- 4. No CIE ID SERVER databases
- 5. No preliminary phase to bind smartphone-desktop

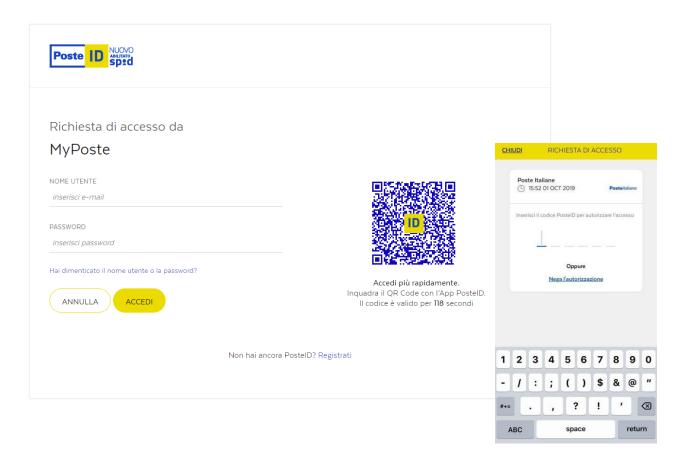
Final choice: QR codes







CIE 3.0 - Hybrid scenario State-of-the-art













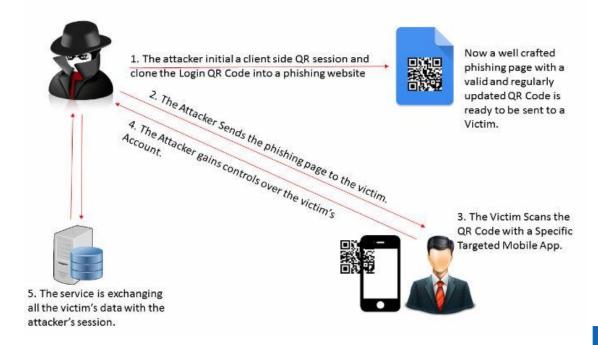
CIE 3.0 – Hybrid scenario Known attack

QRLJacking (Quick Response Code Login Jacking) is a social engineering attack vector capable of session hijacking affecting all applications that rely on "Login with QR code" feature as a secure way to login into accounts. [OWASP]

In a nutshell victim scans the attacker's QR code results of session hijacking.

Attack Flow

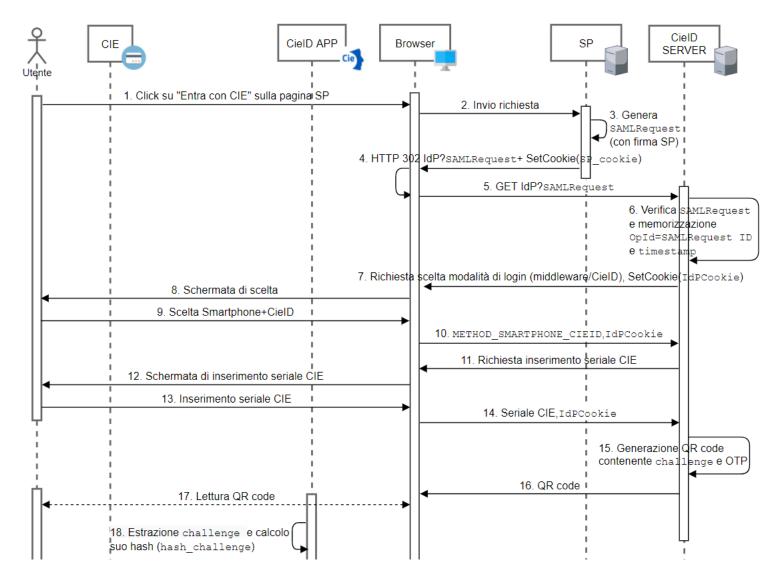
- The attacker initial a client side QR session and clone the Login QR Code into a phishing website.
- The Attacker Sends the phishing page to the victim. (refer to <u>QRLJacking real life attack vectors</u>)
- 3. The Victim Scans the QR Code with a Specific Targeted Mobile App.
- 4. The Attacker gains controls over the victim's Account.
- The service is exchanging all the victim's data with the attacker's session





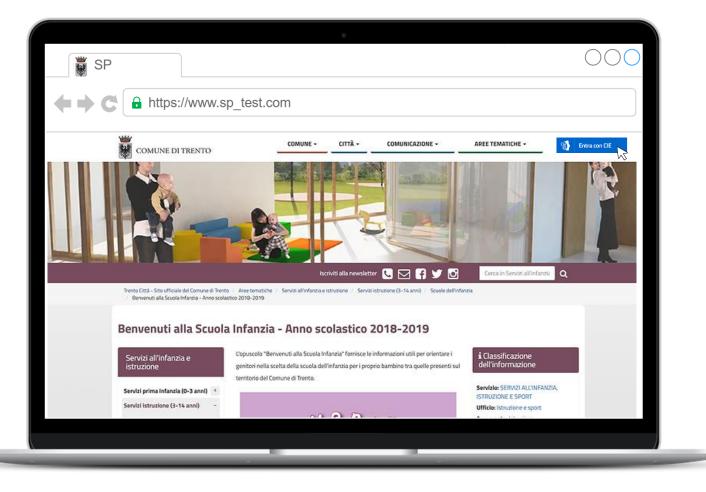


CIE 3.0 – Hybrid scenario MSC . . .



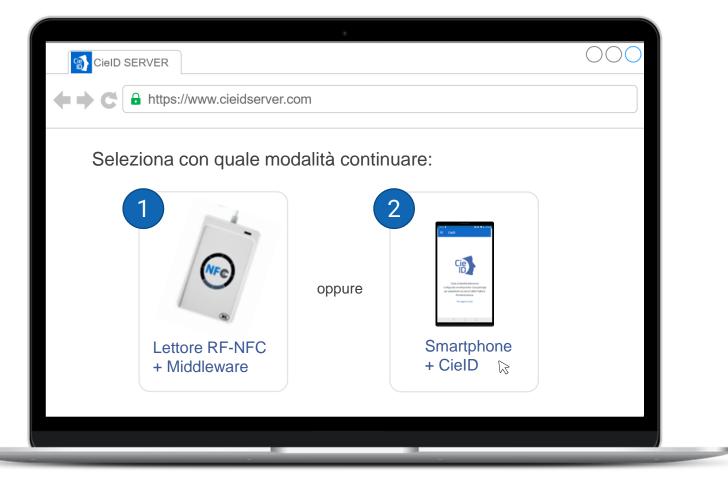






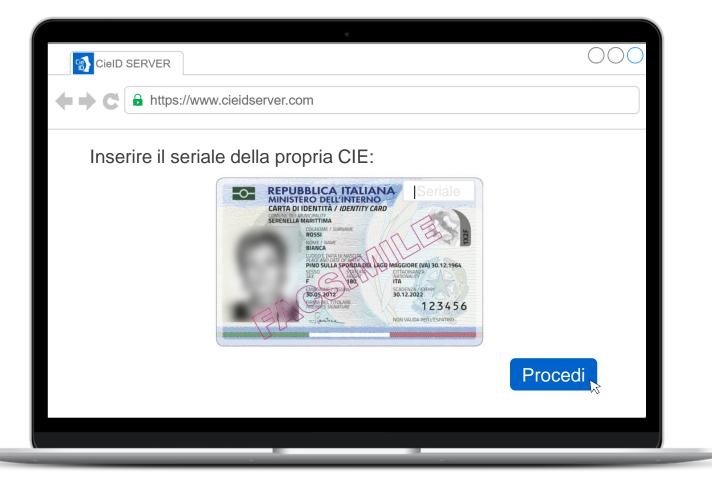






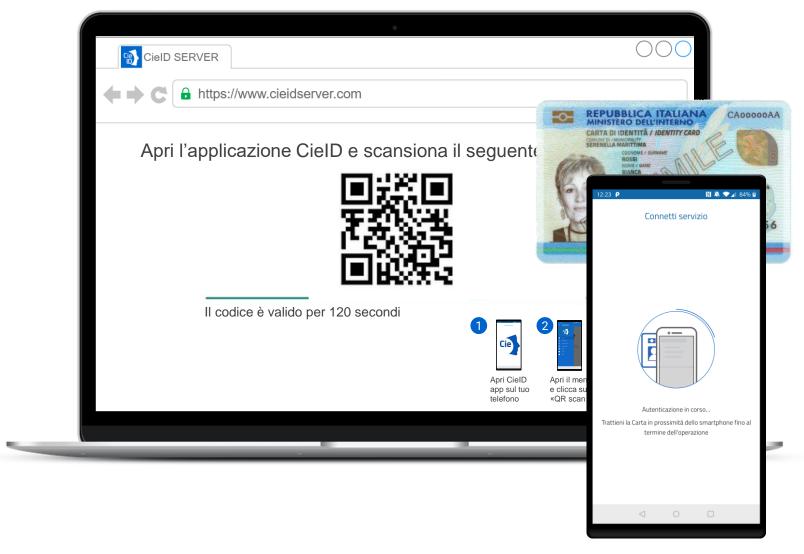






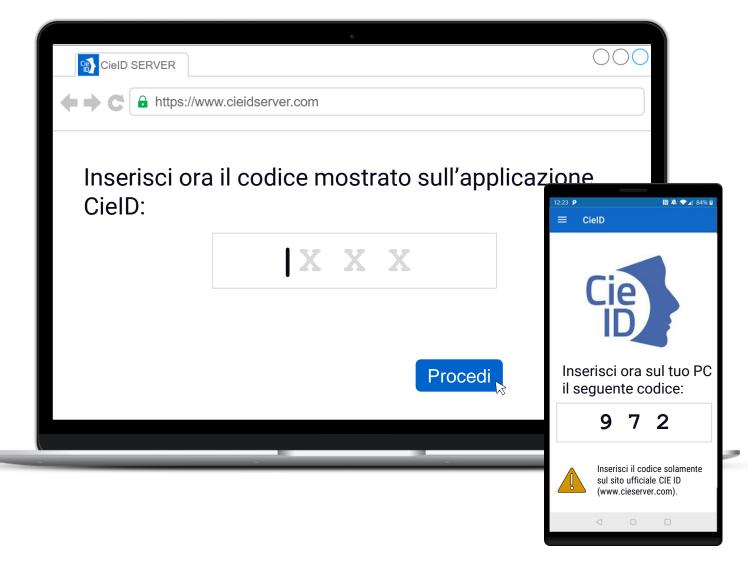






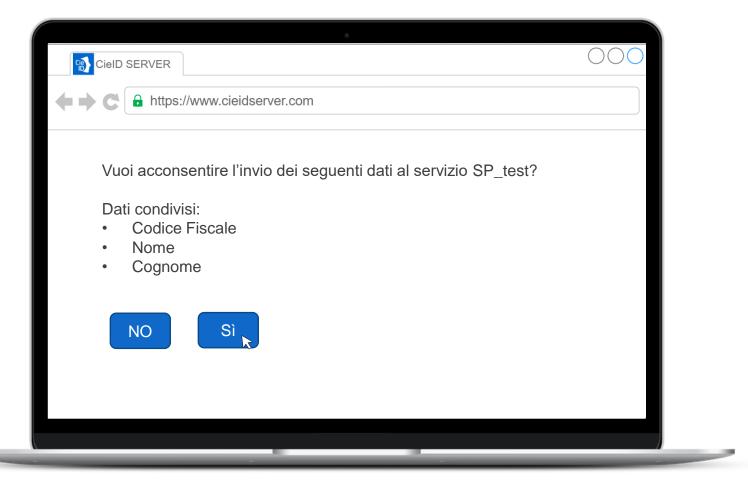








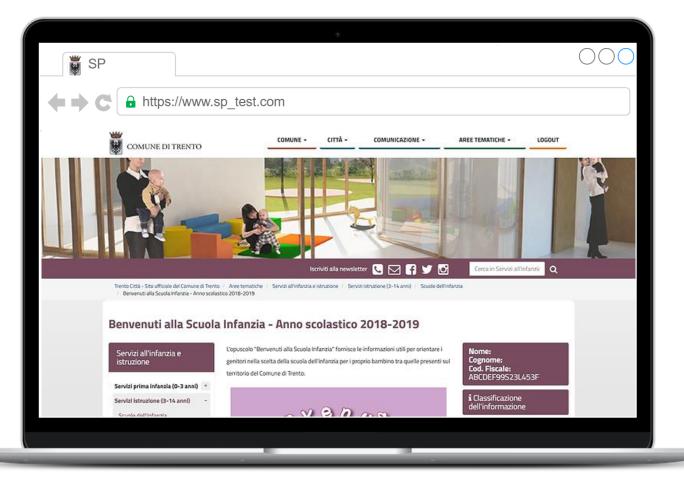








CIE 3.0 – Hybrid scenario Mockup







QRLJacking Mitigations

1. Require the input of a uniquely identifying information for the generation of the QR code



Inserire il seriale della propria CIE:







QRLJacking Mitigations

- 1. Require the input of a uniquely identifying information for the generation of the QR code
- 2. QR code with a time validity









| QRLJacking | Mitigations

- 1. Require the input of a uniquely identifying information for the generation of the QR code
- 2. QR code with a time validity
- 3. Advise users to verify the trustworthiness of the source by checking that the URL really belongs to CIE ID SERVER







QRLJacking Mitigations

- 1. Require the input of a uniquely identifying information for the generation of the QR code
- 2. QR code with a time validity
- 3. Advise users to verify the trustworthiness of the source by checking that the URL really belongs to CIE ID SERVER
- 4. At the end of the authentication procedure, display an OTP on the mobile device and require users to insert it in the personal computer's browser







QRLJacking Mitigations

- 1. Require the input of a uniquely identifying information for the generation of the QR code
- 2. QR code with a time validity
- 3. Advise users to verify the trustworthiness of the source by checking that the URL really belongs to CIE ID SERVER
- 4. At the end of the authentication procedure, display an OTP on the mobile device and require users to insert it in the personal computer's browser
- 5. Show the details of the on-going operation





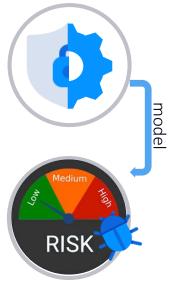


DigiMat-Lab and F&C Methodology



M. Pernpruner, R. Carbone, S. Ranise, G. Sciarretta. *The Good, the Bad and the (Not So) Ugly of Out-of-Band Authentication with eID Cards and Push Notifications: Design, Formal and Risk Analysis.* CODASPY '20

M. Pernpruner, G. Sciarretta, S. Ranise. *A Framework for Security and Risk Analysis of Enrollment Procedures: Application to Fully-Remote Solutions Based on eDocuments*. SECRYPT 2021.



Security-by-design

Design and implementation of the protocol, based on state-of-the-art techniques

Automatic Security and Risk Assessment

With the goal of automatically detecting possible threats and associating each successful attack with the related risk and mitigations.





Protocol analysis Security and Risk Assessment



Login as another's identity

Security Analysis: detects which attackers have success



Is able to steal a user device



is able to intercept the data typed on the device

Risk Analysis: allows us to classify the successful attacks based on their seriousness, and thus to prepare a mitigation plan accordingly



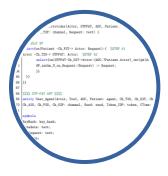


Security analysis 2 layers



Combinatorial Analysis

Relying on attackers' capabilities on the authentication factors, given that the violation of all of them necessarily results in the violation of the whole protocol. This analysis can be carried out quickly and returns all the explicit attacks (those deriving from the violation of all the authentication factors).



Formal Analysis

Relying on formal methods (a specification language and a model checker). It can be computationally expensive, but manages to find even more complex categories of attacks: implicit attacks manage to violate the protocol even without explicitly compromising all the authentication factors, as they manage to deceive the victim into implicitly compromising the remaining factors on their behalf.





• The security goal (SG) is the set of authentication factors that should not be compromised for the authentication procedure to be considered secure.

$$SG = \{ \blacksquare; PIN \}$$



Multi-Factor Cryptographic Device: «a hardware device that performs cryptographic operations using one or more protected cryptographic keys and requires activation through a second authentication factor»









• The security goal (SG) is the set of authentication factors that should not be compromised for the authentication procedure to be considered secure.

$$SG = \{ \blacksquare; PIN \}$$

• A threat model (TM) over the identification factors is a pair: $(\mathcal{A}TT; \mathcal{C})$

where:

- ATT is the set of considered attackers;
- C represents their capabilities.





Security analysis
Attackers

Shoulder Surfer: obtains secrets by looking at the user insering sensitive info



MA SE ES

Social Engineer: deceives people into revealing secret information or performing actions to their advantage



Malicious App: runs on the attacker's or the victim's mobile device

CIE thief: steals an CIE from its legitimate owner



Eavesdropping Software: intercepts the data typed on the device (e.g., keylogger)







can only compromise the PIN (in case it is written on paper)

can only compromise the PIN (by looking at the victim while typing it)

can only compromise the PIN (by deceiving the victim into revealing it)

Attackers Personal Computer Thief PCT **Mobile Device Thief MDT Card Thief** CT **Authenticator Duplicator** AD **Eavesdropping Software** ES Shoulder Surfer SS **Social Engineer** SE Man in the Browser MB Man in the Mobile MM

can only compromise the eDocument

can only compromise the PIN (by eavesdropping it while it is being typed)

can compromise the eDocument (indirectly, by deceiving the victim into interact with it), the PIN (by eavesdropping it while it is being typed)



= compromised







- An auth flow **violates** the security goal SG under the threat model TM = (ATT; C) iff there is an attacker (or a combination of them) in ATT that compromises all the identification factors contained in the SG associated to the flow.
- A subset $ATT \subseteq \mathcal{A}TT$ is minimal iff ATT violates SG and, for each $ATT' \subseteq ATT$, ATT' does not violate SG.

Attackers	****	2=	
Personal Computer Thief	PCT	<u></u>	
Mobile Device Thief	MDT	<u></u>	
Card Thief	CT	<u></u>	
Authenticator Duplicator	AD	-	<u></u>
Eavesdropping Software	ES	-	
Shoulder Surfer	SS	-	<u></u>
Social Engineer	SE	-	<u></u>
Man in the Browser	MB	<u> </u>	<u></u>
Man in the Mobile	ММ	a	• *

The security analysis problem for an auth flow under a threat model $T\mathcal{M} = (\mathcal{A}TT; \mathcal{C})$ is to find all (if any) minimal subsets $ATT \subseteq \mathcal{A}TT$ so that ATT violates $\mathcal{S}\mathcal{G}$.





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Successful explicit attackers
MM
CT+AD
CT+ES
CT+SS
CT+SE

The security analysis problem for an auth flow under a threat model $T\mathcal{M} = (\mathcal{A}TT; \mathcal{C})$ is to find all (if any) minimal subsets $ATT \subseteq \mathcal{A}TT$ so that ATT violates $\mathcal{S}\mathcal{G}$.



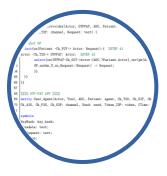


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Security analysis Formal analysis

SECURITY & TRUST **select** {

```
entity IdPServer(Actor, EICApp, User, SPServer, Browser, EIC: agent, Ch B2IdPS, Ch EICApp2IdPS, Ch IdPS2EICApp: channel) {
symbols
OpId: opid;
Request: userrequest;
OTP: otp;
IdPSessionCookie: cookie;
body { % of IdPServer
   select {
       on(Browser -Ch B2IdPS-> Actor: ?Request):{
             IdPSessionCookie := fresh();
             OpId := fresh();
             Actor -Ch IdPS2B-> Browser: Actor.IdPSessionCookie;
             select {
                 on(Browser -Ch B2IdPS-> Actor: User.IdPSessionCookie):{
                 Actor -Ch IdPS2B-> Browser: OpId.Actor.SPServer.User;
                 select {
                     on(EICApp -Ch EICApp2IdPS-> Actor: OpId.{OpId.Actor.SPServer.User} inv(pk(EIC))):{
                     OTP := fresh();
                     % iknows(OTP); % uncomment for SE
                     Actor -Ch IdPS2EICApp-> EICApp: OTP;
                     Actor -Ch IdPS2B-> Browser: Actor;
```

Security analysis Formal analysis

•	Formal Spec	ication				
Att.	Without Attacker	With Attacker				
PCT	(1) authentic_on(Ch_U2B, User); (2) userOwnComputer; ————————————————————————————————————	(3) weakly_authentic(Ch_U2B); (4) Use same channel ch_U2B in sessions				
MDT	(5) authentic_on(Ch_U2EICApp, User); (6) userOwnSmartphone; — — —	(7) weakly_authentic(Ch_U2EICApp); (8) Use same channel ch_U2EICApp in sessions				
CT	(9) authentic_on(Ch_U2EIC, User); (10) userOwnEIC; ————————————————————————————————————	(11) weakly_authentic(Ch_U2EIC); (12) Use same channel ch_U2EIC in sessions				
D	_ _	(13) iknows(PIN); (14) iknows(IdPCookie);				
ES, SS	<pre>(15) confidential_to(Ch_U2EICApp, EICApp); (16) confidential_to(Ch_EICApp2U, User); (17) confidential_to(Ch_U2B, Browser);</pre>	- - -				
SE	(18) confidential_to(Ch_U2B, Browser); (19) authentic_on(Ch_B2U, Browser); — — —	(20) iknows(PIN); (21) iknows(OTP);				
MB	_ _ _	(22) iknows(IdPCookie); (23) Replace browser with i in one session				
MM	<pre>(24) authentic_on(Ch_EICApp2U, EICApp); (25) authentic_on(Ch_EICApp2EIC, EICApp); (26) confidential_to(Ch_FCMSvc2EICApp, EICApp);</pre>	(27) Replace eicapp with i in one session				

User_authn_to_SP:(_) User *->> SPServer;



Implicit attacks
SE MB





Risk analysis OWASP Risk Rating Methodology

The **risk analysis problem** for an enrollment flow under a threat model $T\mathcal{M}=(\mathcal{A}TT;\mathcal{C})$ is to find the risk associated with all the minimal subsets of attackers violating SG.

A++		Likelihood							Impact					Diek
Att.	TD	0	AV	UI	SA	Aver.	Over.	LSP	AS	AD	AP	Aver.	Over.	Risk
MM	2	2	7	1	2	2.80	Low	9	8	3	2	5.50	Medium	Low

- 1. Assign a score (0-9) to each factor
- 2. Compute the average of likelihood and impact factors
- 3. Obtain the overall likelihood and impact
- 4. Compute the risk

ina impact									
•		Likelihood							
		Low	Medium	High					
	Low	Note	Low	Medium					
Impact	Medium	Low	Medium	High					
	High	Medium	High	Critical					





Low

Medium

High

v < 3

3 < v < 6

v < 9

Protocol analysis Final results

At the end of the analysis, we can know:

- a list of attackers that manage to compromise the protocol;
- an indication of the risk for each attacker.

	Attackers	Likelihood						Impact						Risk	
		TD	О	AV	UI	SA	Overall		Overall LSP AS AD AP Overall		verall	Overall			
	MM	2	2	7	1	2	2.80	Low	9	8	3	2	5.50	Medium	Low
explicit	CT+D	8	1	1	7	4	4.20	Medium	9	2	3	8	5.50	Medium	Medium
	CT+ES	5	0	1	4	2	2.40	Low	9	2	3	8	5.50	Medium	Low
	CT+SS	8	4	1	2	5	4.00	Medium	9	2	3	8	5.50	Medium	Medium
	CT+SE	4	2	1	4	3	2.80	Low	9	2	3	8	5.50	Medium	Low
implicit	SE	4	9	7	1	4	5.00	Medium	9	5	7	2	5.75	Medium	Medium
ппрпоп	MB	3	5	7	1	4	4.00	Medium	9	8	7	2	6.50	High	High





Protocol analysis The role of mitigations

- Mitigations can be specified by properly adjusting:
 - the attackers' capabilities (C);
 - the risk scores assigned to the likelihood and impact factors.



Require the input of a uniquely identifying information during the authentication protocols

SE has less probability of performing a QRLJacking



Implement root detection mechanisms on the mobile application

MM is less likely able to steal the PIN from the internal storage

Attackers	Likelihood	Impact	Risk
MM	Medium	High	High
CT+D	Medium	Medium	Medium
CT+ES	Low	Medium	Low
CT+SS	Medium	Medium	Medium
CT+SE	Medium	Medium	Medium
SE	High	High	Critical
MB	Medium	High	High





CIE 3.0

Real-world scenarios and future directions







Online authentication



Remote identity proofing



Electronic Signature



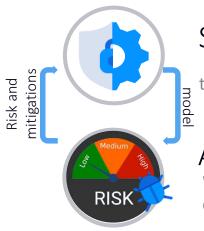
Advanced scenarios







DigiMat-Lab and F&C Methodology



Security-by-design

Design and implementation of the protocol, based on state-of-the-art techniques

Automatic Security and Risk Assessment

With the goal of automatically detecting possible threats and associating each successful attack with the related risk and mitigations.

It allows what-if analyses, by providing information on how specific mitigations affect the set of successful attackers and their risks





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https://stfbk.github.io/



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