Un jugement d'évaluation s'écrit sous la forme $\gamma \vdash [e, m_1] \downarrow [v, m_2]$.

Règles

$$\begin{array}{c} \gamma \vdash [entier\,,\,m] \Downarrow [entier\,,\,m] \\ \hline x \in \gamma \quad \gamma(x) = \langle e\,,\, \gamma_{def} \rangle \quad \gamma_{def} \vdash [e\,,\,m_1] \Downarrow [v\,,\,m_2] \\ \hline \gamma \vdash [x\,,\,m_1] \Downarrow [v\,,\,m_2] \\ \hline \gamma \vdash [e_2\,,\,m_1] \Downarrow [v_2\,,\,m_2] \quad \gamma \vdash [e_1\,,\,m_2] \Downarrow [v_1\,,\,m_3] \quad v_1 \times v_2 \in dom \ op \quad v = v_1 \ op \ v_2 \\ \hline \gamma \vdash [e_1 \ op \ e_2\,,\,m_1] \Downarrow [v\,,\,m_3] \\ \hline \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v\,,\,m'] \quad v \in dom \ op \quad v' = op \ v \\ \hline \gamma \vdash [op \ e\,,\,m] \Downarrow [v',\,m'] \\ \hline \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [true\,,\,m_2] \quad \gamma \vdash [e_2\,,\,m_2] \Downarrow [v\,,\,m_3] \\ \hline \gamma \vdash [if \ e_1 \ then \ e_2 \ else \ e_3\,,\,m_1] \Downarrow [v\,,\,m_3] \\ \hline \gamma \vdash [if \ e_1 \ then \ e_2 \ else \ e_3\,,\,m_1] \Downarrow [v\,,\,m_3] \\ \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \quad \gamma \colon \{x \mapsto v_1\} \vdash [e_2\,,\,m_2] \Downarrow [v\,,\,m_3] \\ \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \quad \gamma \colon \{x \mapsto v_1\} \vdash [e_2\,,\,m_2] \Downarrow [v\,,\,m_3] \\ \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \quad \gamma \vdash [e_1\,,\,m_1] \Downarrow [v\,,\,m_3] \\ \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \quad \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_3] \\ \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \quad \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_3] \\ \hline \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \quad \gamma \vdash [e_1\,,\,m_1] \Downarrow [v_1\,,\,m_2] \\ \hline \gamma \vdash [e_1\,,\,m_2] \Downarrow [\langle \operatorname{fun} \ x \to e_3\,,\,\gamma_{def}\,\rangle\,,\,m_3] \quad \gamma_{def} \colon \{x \mapsto v_2\} \vdash [e_3\,,\,m_3] \Downarrow [v\,,\,m_4] \\ \hline \gamma \vdash [(e_1\,)\ e_2\,,\,m_1] \Downarrow [v\,,\,m_4] \\ \hline \gamma \vdash [(e_1\,)\ e_2\,,\,m_1] \Downarrow [v\,,\,m_4] \\ \hline \gamma \vdash [(e_1\,)\ e_2\,,\,m_1] \Downarrow [v\,,\,m_2] \\ \hline \gamma \vdash [\operatorname{letrec} \ x = e_1\ \operatorname{in} \ e_1\,,\,\gamma\}\} \vdash [e_2\,,\,m_1] \Downarrow [v\,,\,m_2] \\ \hline \gamma \vdash [\operatorname{letrec} \ x = e_1\ \operatorname{in} \ e_2\,,\,m_1] \Downarrow [v\,,\,m_2] \\ \hline \gamma \vdash [\operatorname{letrec} \ x = e_1\ \operatorname{in} \ e_2\,,\,m_1] \Downarrow [v\,,\,m_2] \\ \hline \end{array}$$

Justifications

Opération binaire

```
# let x = ref 1 in ((x:=2;(!x))+(!x));;
- : int = 3
# let x = ref 1 in ((!x)+(x:=2;(!x)));;
- : int = 4
```

Appel de fonction

```
# let x = ref 1 in (x:=2; fun y -> y+(!x)) (!x) ;;
- : int = 3
# let x = ref 1 in (fun y -> y +(!x)) (x:=2; (!x));;
- : int = 4
```