Announcements

- 1. Problem Set 1 out today, due Thursday, Feb 12 (start of class)
- 2. Mid-term exam on Thursday, Feb 26 (in class)

3. Callback

Basic Idea

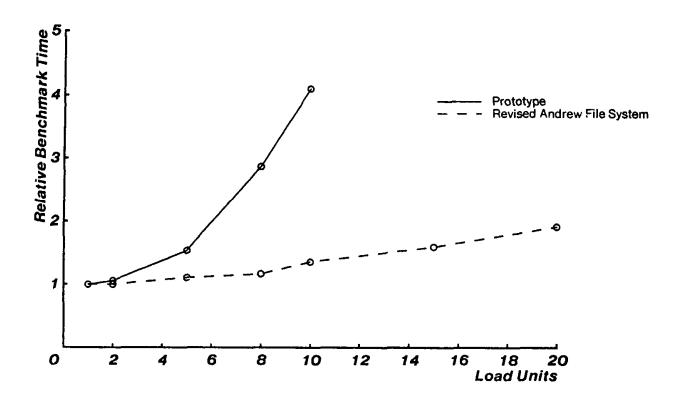
- targeted notification of caching sites
- master copy tracks sites with cached data
- typically done at coarse granularity (e.g. entire file) can be made to work with byte ranges
- on update, all sites with cached data notified ("callback")

Original Use

AFS-2 (circa 1985)

Advantages

- excellent scalability for Unix workloads
- zero network traffic for read of cached-valid objects
- precursor to caching for disconnected operation
- biases read performance in favor of write-performance



- sizable state on server
- complexity of tracking cached state on clients
- silence ambiguous for client network failure → lost callbacks periodic "keepalive" probes data could be stale between probes
- NAT networks with masquerading firewalls

4. Leases

Basic Idea

Caching site obtains finite-duration control from master copy

duration is called "lease period", typically few seconds multiple sites can obtain read lease; only one can get write lease

leases have to be renewed, else control expires with lease

Limiting cases

- lease duration = 0 → Check on Use
- lease duration = ∞ \rightarrow callback (Targeted Notification)

Original Use

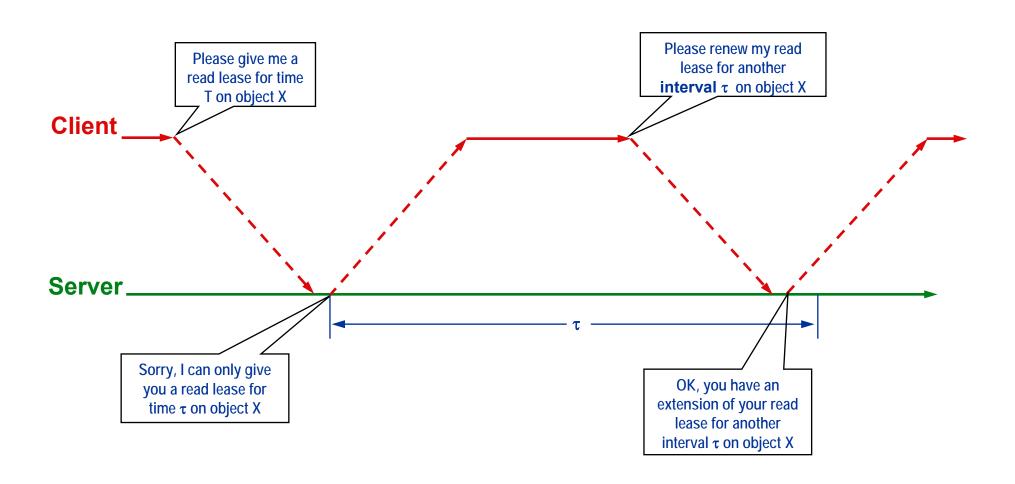
V System (circa 1985) [Gray89]

Advantages

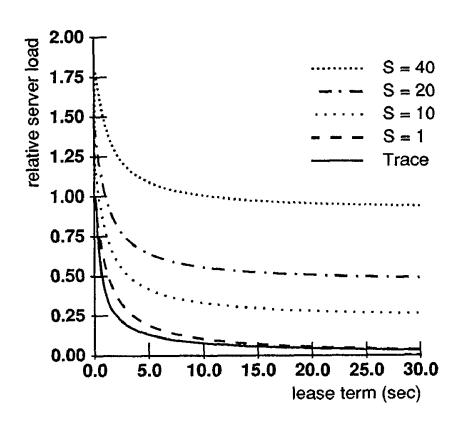
- generalizes the check on use and callback schemes
- lease duration can be tuned to adapt to mutation rate
 Lease duration is a clean tuning knob for design flexibility
- conceptually simple yet flexible
- time becomes a hidden communication channel

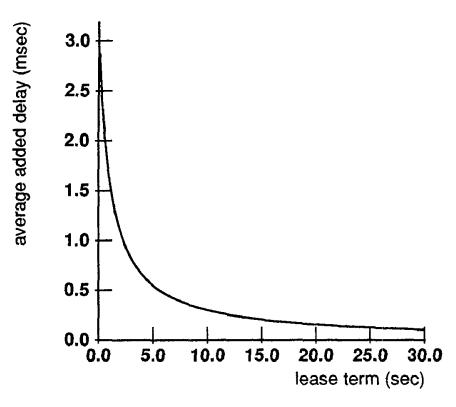
Lease Renewal

At most one write lease or multiple read leases at any one time



- lease-holder has total autonomy during lease; revocation?
- writers delayed while read lease holders complete their leases
- more traffic than callback (but less than check on use)
 keepalives for callback only one per server, not per lease





5. Skip Scary Parts

Basic Idea

- When write-sharing detected, turn off caching everywhere
 All references go directly master copy
- Resume caching when write-sharing ends

Original Use

• Sprite (circa 1987) (in conjunction with check on use)

Advantages

- Precise single-copy semantics (even at byte-level consistency)
- Excellent fallback position

 Exemplifies good engineering: "Handle average case well; worst case safely"

 Coda protocol achieves this differently: treats write-sharing as exception (conflict)
- Good adaptation of caching aggressiveness to workload

- Server maintains state
- Server aware of every use of data (open)

6. Faith-Based Caching

Basic Idea

- Blindly assume cached data is valid for a while
- Periodically check (based on time since last check)
- No communication needed during trust period

Original use

- Sun NFSv3 file system

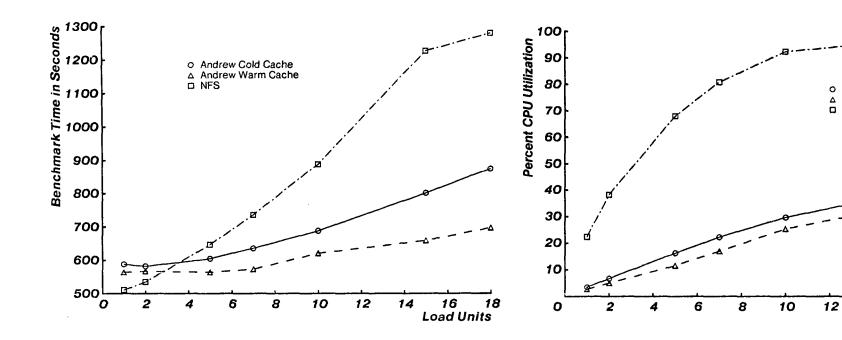
 cached file blocks assumed current for X seconds
 X = 3 for files, 30 for directories
- Small variant is a *TTL field* for each object used in web caching, gives content creator modicum of control

Advantages

- Simple implementation
- Server is stateless

Disadvantages

- User-visible inconsistencies sometimes seen (make)
- Blind faith sometimes misplaced!
- Not as efficient as callback-based schemes



16

Load Units

18

O Andrew Cold Cache

△ Andrew Warm Cache

14

□ NFS

7. Pass the Buck

Basic Idea

- Let the user trigger cache revalidation (hit "reload")
- Otherwise, all cached copies assumed valid forever
- Equivalent to infinite-TTL faith-based caching

Original use

Arose in the context of the Web

Advantages

- Trivial to implement, no server changes
- Avoids frivolous cache maintenance traffic

- Places burden on user
- User may be clueless about level of consistency needed
- Assumes existence of user: pain for write scripts/programs

Cache Consistency Strategies

- 1. Broadcast invalidations
- 2. Check on Use
- 3. Callbacks
- 4. Leases
- 5. Skip Scary Parts
- 6. Faith-based Caching
- 7. Pass the Buck

Many minor variants over the years, but these have withstood the test of time