

2022 Summer Mid Set A	<p>1a</p> <p>Calculate the time complexity of the following function:</p> <div data-bbox="491 300 1249 553" data-label="Text"> <pre>worstCase(n): int i, j, k, a, b, sum for (i = 0; i < n; i = i + 3) for (j = n; j >= 1; j = j / 5) for (k = 1; k <= n; k = k * 5) sum = a + b</pre> </div> <p>Ans : $O(n \log_5^2 n)$</p>
2022 Summer Mid Set B	<p>1a</p> <p>Calculate the time complexity of the following function:</p> <div data-bbox="481 799 1256 1052" data-label="Text"> <pre>worstCase(n): int i, j, k, a, b, sum for (i = 1; i < n; i = i + 3) for (j = n; j >= 1; j = j - 2) for (k = n; k >= 1; k = k / 4) sum = a + b</pre> </div> <p>Ans : $O(n^2 \log_4 n)$</p>
2022 Fall Mid Set A	<p>1a</p> <p>Finding_Worst_Case(n):</p> <div data-bbox="276 1301 1398 1910" data-label="Text"> <pre>int i,j,k,m,multi,a,b,c for(i = n; i >= 1; i = i / 7){ for(j = 1; j <= n; j = j + 3) { for(k=1; k<=40 ; k=k+1){ multi=a*b } for(m=n ; m>=1 ; m=m-5){ multi=multi*c } } }</pre> </div> <p>First for loop : $\log_7 n$ Second for loop : $n/3$ Third for loop : 40 Fourth for loop : $n/5$</p>

	<p>Total = $(\log_7 n * n/3 * (40+n/5))$</p> <p>Ans : $O(n^2 \log_7 n)$</p>
2022 Fall Mid Set B	<p>1a</p> <p>Finding_Worst_Case(n):</p> <pre> int i,j,k,m,multi,a,b,c for(i = n; i >= 1; i = i - 4){ for(j = 1; j <= n; j = j * 3) { for(k=1; k<=20 ; k=k+1){ multi=a*b } for(m=n ; m>=1 ; m=m/5){ multi=multi*c } } } </pre> <p>First for loop : $n/4$</p> <p>Second for loop : $\log_3 n$</p> <p>Third for loop : 20</p> <p>Fourth for loop : $\log_5 n$</p> <p>Total = $(n/4 * \log_3 n * (20 + \log_5 n))$</p> <p>Ans: $O(n(\log n)^2)$</p>
2023 Spring Mid Set A	<p>1a. In the primary scholarship exam in Bangladesh, four lakh ($n=4,00,000$) students take part but only the top 50 students are given an award. Write the asymptotic time complexity to give the awards. Assume that each award is given in a constant time.</p> <p>Ans: $O(1)$ / Constant</p> <p>1b. Write the asymptotic time complexity of the following function.</p> <div style="border: 1px solid black; padding: 5px; margin-top: 10px;"> <p>1. def contains_duplicates(elements):</p> </div>

	<pre> 2. for outer in range(len(elements)): 3. for inner in range(len(elements)): 4. if outer == inner: 5. continue 6. 7. if elements[outer] == elements[inner]: 8. return True 9. 10. return False </pre> <p>Ans: $O(n^2)$</p>
2023 Spring Mid Set B	<p>1a. In the primary scholarship exam in Bangladesh, two lakh ($n=2,00,000$) students take part but only the top 25 students are given an award. Write the asymptotic time complexity to give the awards. Assume that each award is given in a constant time.</p> <p>Ans : $O(1)$ / Constant.</p> <p>1b. Write the asymptotic time complexity of the following function.</p> <pre> 1. def cumulative_sum(elem): 2. for outer in range(len(elem)): 3. for inner in range(outer+1,len(elem)): 4. elem[outer]= elem[outer] + elem[inner] 5. 6. return elem </pre> <p>Ans: $O(n^2)$.</p>
2023 Summer Mid Set A	<p>1a. Find the time-complexity of the following task in terms of number of students.</p> <p>You are given a student attendance sheet. Each student has a unique integer ID. You have to count the number of students having an even number as ID. The list is sorted but the IDs are not necessarily consecutive. So you check each ID one by one.</p> <p>Ans: If number of students is N, then $O(N)$</p> <p>1b. Write the asymptotic time complexity of the following code snippet. Show your works/reasoning.</p> <pre> 1. for i in range (1,n) </pre>

	<div data-bbox="284 114 1460 293" data-label="Text"> <pre> 2. j= 1 3. while j < i*i 4. j= j+1 </pre> </div> <p>Ans:</p> $T(N) = 1^2 + 2^2 + 3^2 + \dots + n^2 = (n(n+1)(2n+1))/6$ $= O(n^3)$
2023 Summer Mid Set B	<p>1a. Find the time-complexity of the following task in terms of number of students.</p> <p>You are given a student attendance sheet. Each student has a unique integer ID. You have to count the number of students having an ID which is divisible by 3. The list is sorted but the IDs are not necessarily consecutive. So you check each ID one by one.</p> <p>Ans: If number of students is N, then O(N)</p> <p>1b. Write the asymptotic time complexity of the following code snippet. Show your works/reasoning.</p> <div data-bbox="284 958 1460 1189" data-label="Text"> <pre> 1. for i in range (1,n) 2. j= 1 3. while j*j < i 4. j= j+1 </pre> </div> <p>Ans:</p> <p>First for loop : $n-1 \approx n$</p> <p>Second while loop :</p> <p>Runs till $j*j \geq i$. Means j equals to \sqrt{i}.</p> <p>Number of iterations for each value of i is \sqrt{i}.</p> <p>So, the series becomes, $\sqrt{1} + \sqrt{2} + \sqrt{3} + \sqrt{4} + \sqrt{5} + \dots + \sqrt{n}$.</p> $\sqrt{1} + \sqrt{2} + \sqrt{3} + \sqrt{4} + \sqrt{5} + \dots + \sqrt{n} = \sum_{i=1}^n \sqrt{i} = \int_1^n \sqrt{x} dx = \frac{2}{3} * n^{3/2} \approx n^{3/2} \approx n^{1/2} * n^1 = \sqrt{n} * n$ <p>So, time complexity = $O(n\sqrt{n})$</p>
2023 Fall Mid Set A	<p>1a. Explain the time complexity of the following code snippet in regards of the Big-O notation:</p> <div data-bbox="284 1921 1460 2101" data-label="Text"> <pre> 1. for (i=0; i<n; i+=4) { 2. for (j=1; j<n; j*=2) { 3. for (k=0; k<30; k++) { </pre> </div>

```

4.                print("Am I still not 30?!!");
5.                }
6.                print("Why, God, why? We had a Deal!");
7.                for (m=n; m>0; m-=2) {
8.                    print("Could you BE more dramatic?");
9.                }
10.            }
11.}

```

Ans:

First Loop : $n/4$

2nd Loop : $\log_2 n$

3rd Loop : 30

4th Loop : $n/2$

$$n/4 * \log_2 n * (30 + n/2)$$

$$\approx n * \log_2 n * n$$

$$\approx n^2 * \log_2 n$$

1b. Consider the following functions.

$$f_1(n) = (\log n)^{2023}$$

$$f_2(n) = n^2 \log_n(n^n)$$

$$f_3(n) = n^3 + 7n^2$$

$$f_4(n) = 2.023^n$$

$$f_5(n) = n \log n$$

$$f_6(n) = n * \sqrt[3]{n^2}$$

Now do the followings:

- Write a correct asymptotic upper bound for each of the above.
- Sort the functions in ascending order of their growth rate, assuming n is significantly large. Just write the sorted order, no need to show any simulation.

Ans 1(b)(a): Yet to be covered!

And 1(b)(b): $f_1 < f_5 < f_6 < f_3 = f_2 < f_4$

2023
Fall
Mid
Set B

1a. Explain the time complexity of the following code snippet in regards of the Big-O notation:

```

1. for (i=0; i<n; i+=4) {
2.     for (j=1; j<n; j*=2) {
3.         for (k=0; k<20; k++) {

```

```

4.                print("Am I still not 30?!!");
5.                }
6.                print("Why, God, why? We had a Deal!");
7.                for (m=n; m>0; m-=4) {
8.                    print("Could you BE more dramatic?");
9.                }
10.            }
11.}

```

Ans:

First Loop : $n/4$
2nd Loop : $\log_2 n$
3rd Loop : 20
4th Loop : $n/4$

$n/4 * \log_2 n * (20 + n/4)$
 $\approx n * \log_2 n * n$
 $\approx n^2 * \log_2 n$

1b. Consider the following functions.

$$f_1(n) = (\log n)^{2000}$$

$$f_2(n) = n^3 \log_n(n^n)$$

$$f_3(n) = n^3 + 7n^2$$

$$f_4(n) = 4^n$$

$$f_5(n) = n \log n$$

$$f_6(n) = n * \sqrt{n}$$

Now do the followings:

- Write a correct asymptotic upper bound for each of the above.
- Sort the functions in ascending order of their growth rate, assuming n is significantly large. Just write the sorted order, no need to show any simulation.

Ans 1(b)(a): Yet to be covered!

And 1(b)(b): $f_1 < f_5 < f_6 < f_3 < f_2 < f_4$