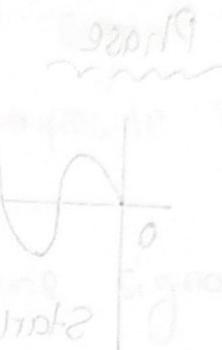
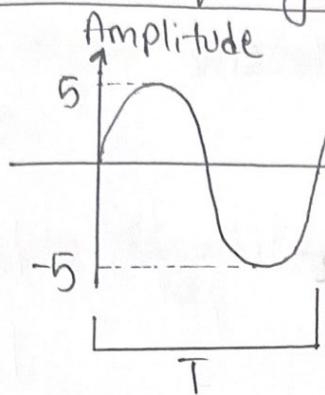


Q

Period Frequency Phase Wave



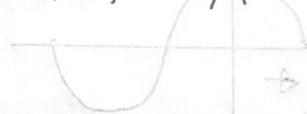
Peak Amplitude = 5

Peak to Peak = $|5 + (-5)| = 10$

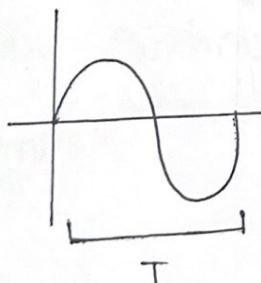
Time Period = To complete one cycle.

Frequency = 1 second how many cycle.

$$T = 1/f \text{ or, } f = 1/T$$



Wavelength



• Signal (WTF single) (W time lag be

= wavelength.

• Propagation speed x Period

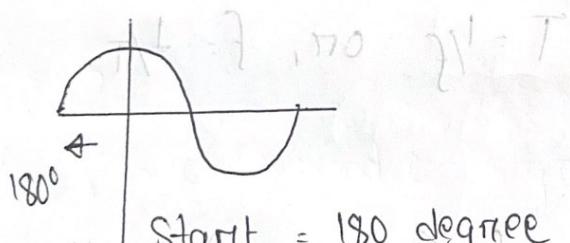
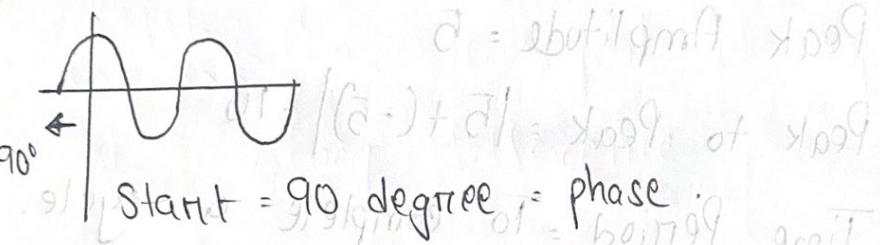
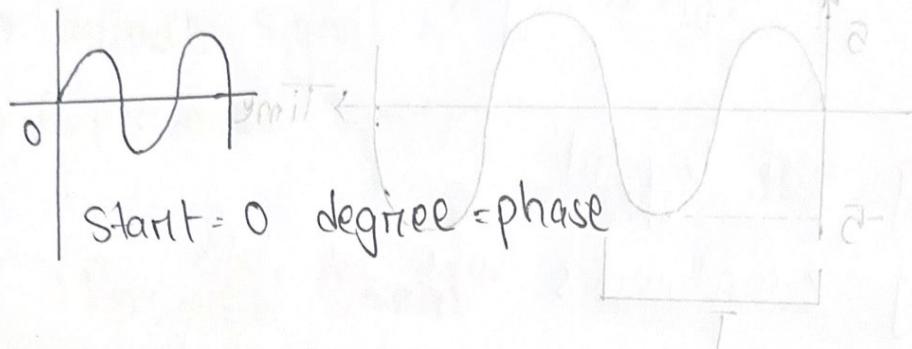
$$(1) \text{ Speed} = CT = C/f$$

Ex: In a vacuum, light propagated with a Speed

of $3 \times 10^8 \text{ m/s}$. Red light frequency $4 \times 10^{14} \text{ Hz}$. Wavelength

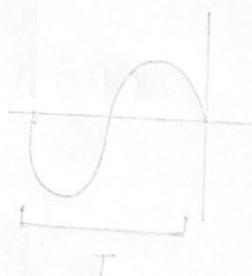
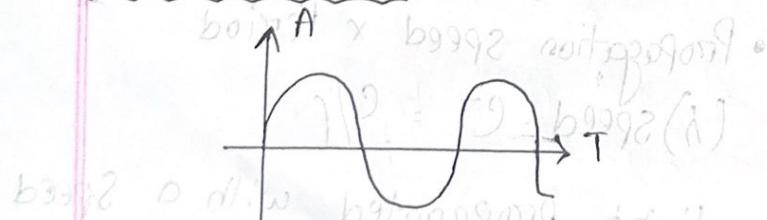
$$\text{of Red Light} = \frac{3 \times 10^8}{4 \times 10^{14}} = 750 \times 10^{-9}$$

Phase

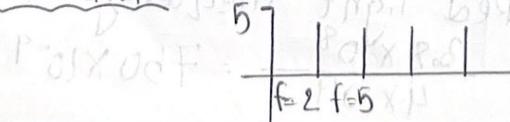


Start = 180 degree = phase

Time Domain



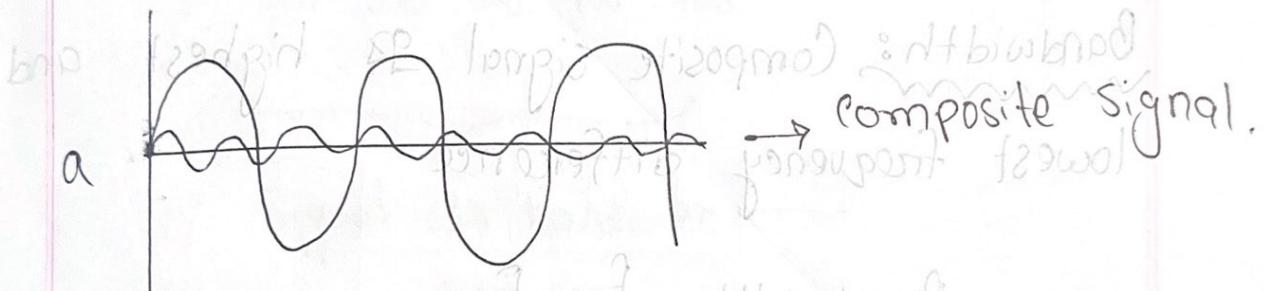
Frequency Domain



Peak amplitude
and frequency
representation.

Data communication এ আমরা single signal বিষ্ণু
ডাই ক্যাটে মার্কস না, আমাদের composite signal
নাস্তা।

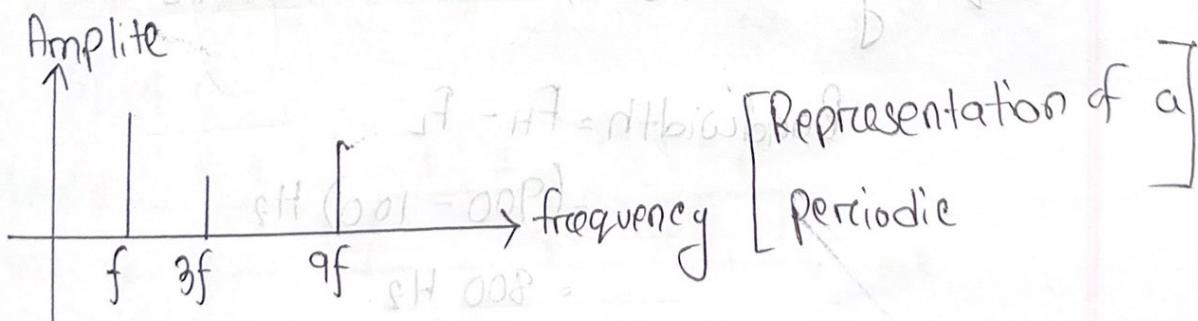
Composite Signal \rightarrow many simple sine signal.



Composite signal নির্দিষ্ট periodie $\frac{2\pi}{f}$ তাৰিখ
decomposition হয়ে মনে discrete frequency পাই,

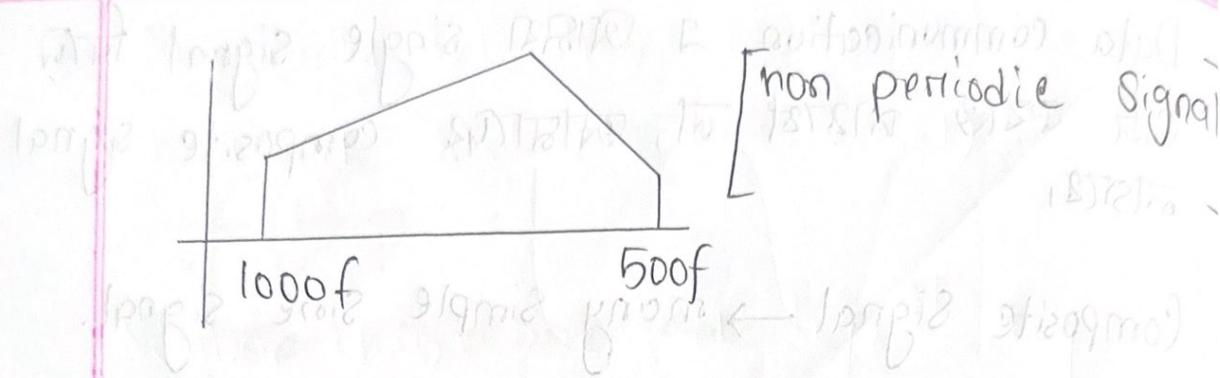
Composite signal (aperiodic) \rightarrow decomposition হয়।

মনে continuous signe frequency f



periodic signal

FO output



Bandwidth: Composite signal is highest and lowest frequency difference

$$\text{Bandwidth} = f_H - f_L$$

Ex: 3

Composite periodic signal

5 sine signal

frequency = 100, 300, 500, 700, 900 Hz.

To calculate Bandwidth = $f_H - f_L$

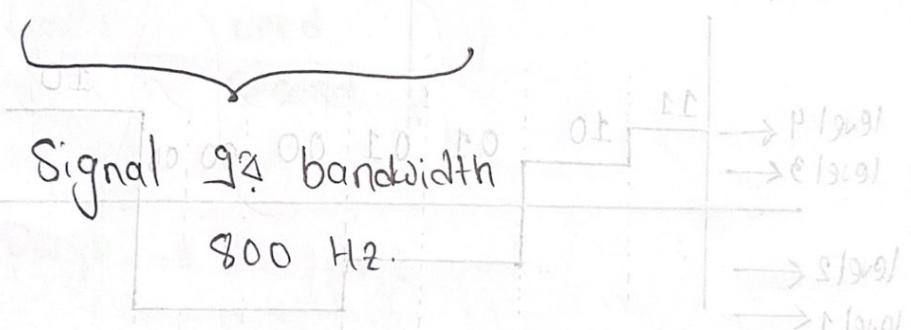
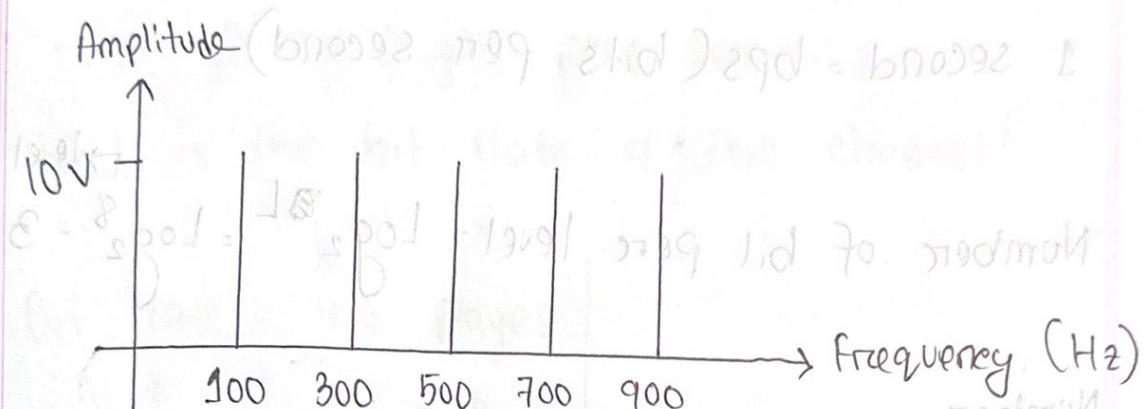
$$(900 - 100) \text{ Hz}$$

$$= 800 \text{ Hz.}$$

And draw the spectrum assuming all component

(time)

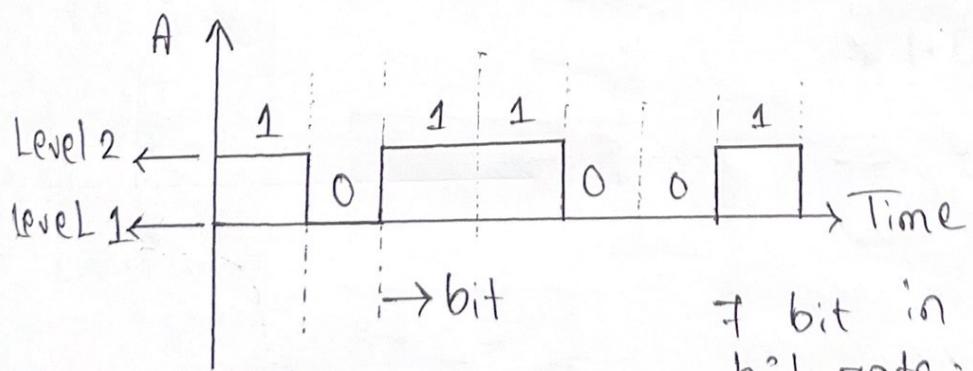
have a maximum amplitude of 10V .



We basically work with digital signal.

I/O नियम बाज़।

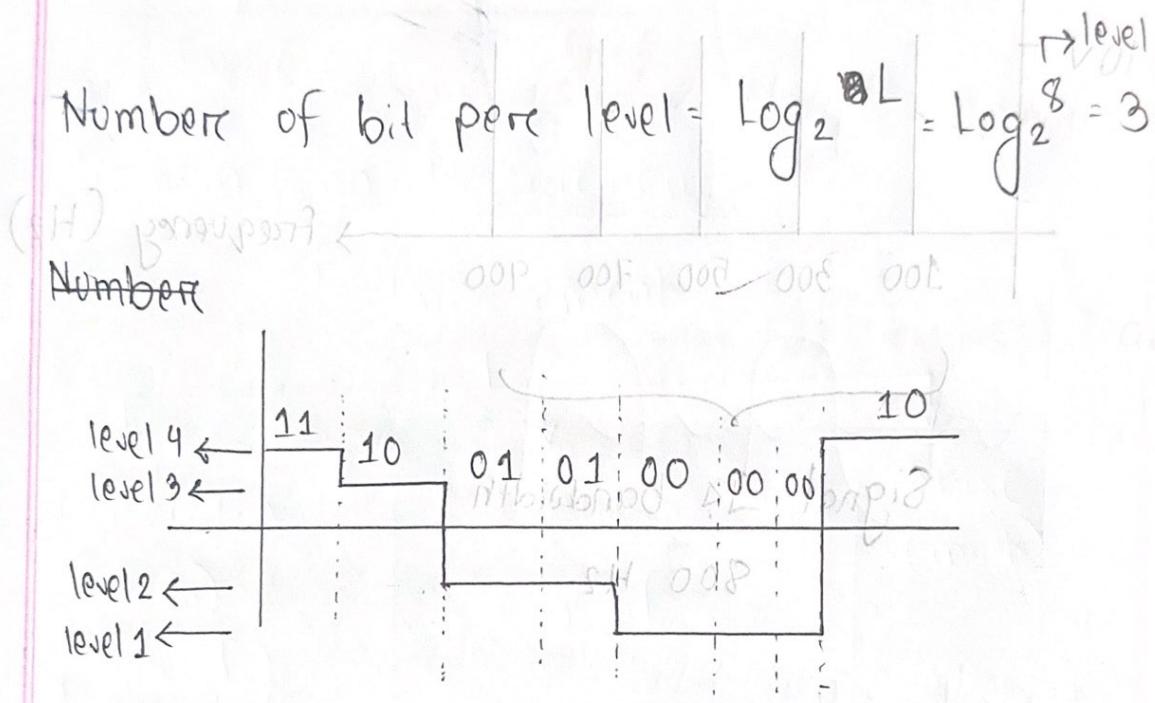
1 → positive voltage } एतो पिछे 2 level हो
0 → negative voltage } इसका signal होता है।



7 bit in 1 sec
bit rate = 7 bps

(send)

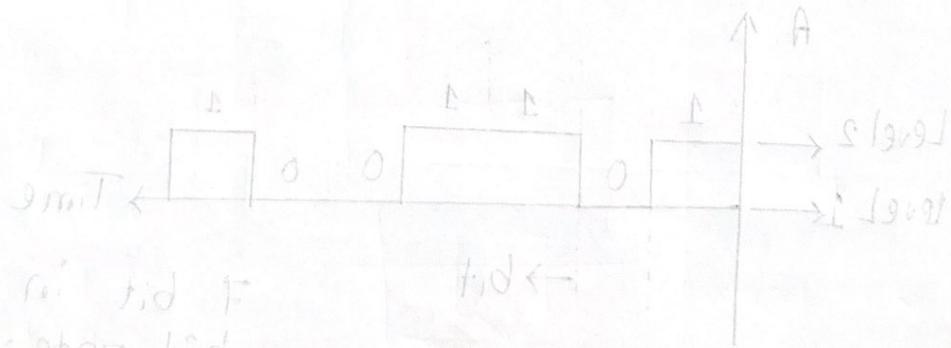
Number of bits that we can traverse in
1 second = bps (bits per second)



Number of bits per level = $\log_2 L$ = $\log_2 4 = 2$

16 bits → 1 second.

Bit rate per second : 16 bps.



Ex: 3.18

Assume you will download a document at the rate of 100 pages per second.
What is the bit rate of the channel?

$$\text{Bit rate} = 100 \text{ Pages}$$

↑
bps (unit)

↓
need
second.

$$\rightarrow 1 \text{ page} \rightarrow 24 \text{ line}$$

$$\rightarrow 1 \text{ line} \rightarrow 80 \text{ character}$$

$$\rightarrow 1 \text{ character} \rightarrow 8 \text{ bit}$$

$$\rightarrow 1 \text{ page} \rightarrow (24 \times 80 \times 8) = 15,360 \text{ bits}$$

$$\rightarrow 100 \text{ pages} \rightarrow (100 \times 24 \times 80 \times 8) = 15,360,000 \text{ bits.}$$

$$= 1.536 \text{ Mbps.}$$



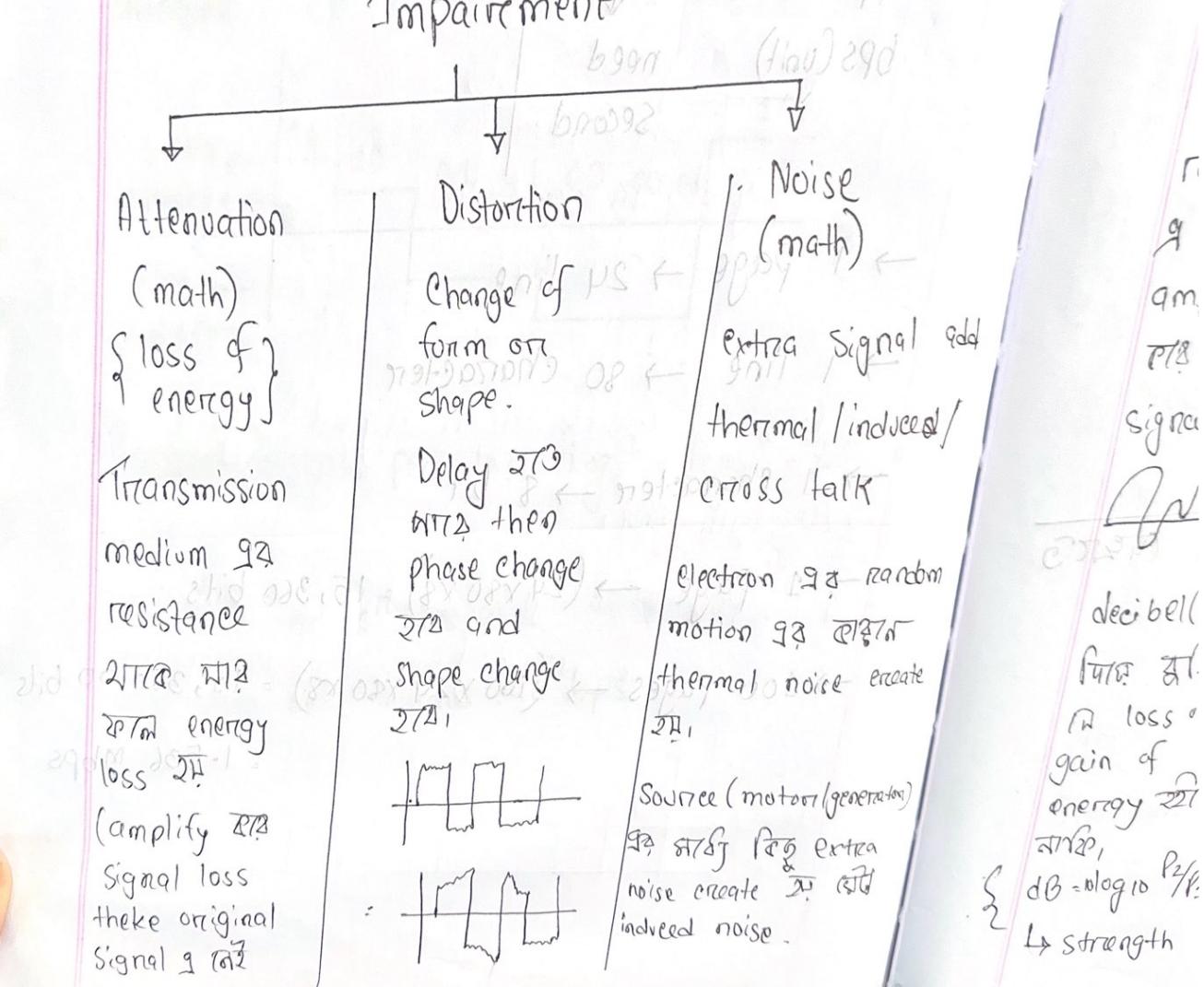
Bit Length

Bit length = propagation speed \times bit duration



Transmission Impairment

Impairment



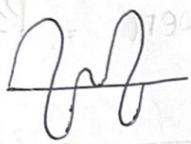
Signal

is shape

size

but size

cho to



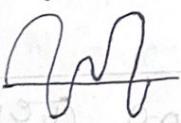
receiver end

is

amplifier use

P18 original

signal



decibel (dB)

dB

loss or

gain of

energy

dB

$$\{ \text{dB} = 10 \log_{10} \frac{P_2}{P_1} \}$$

strength

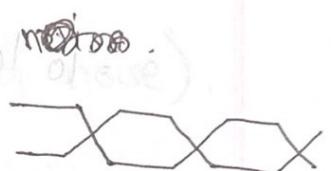
1 ft wire 53

5 ft 2 wires 100

effect

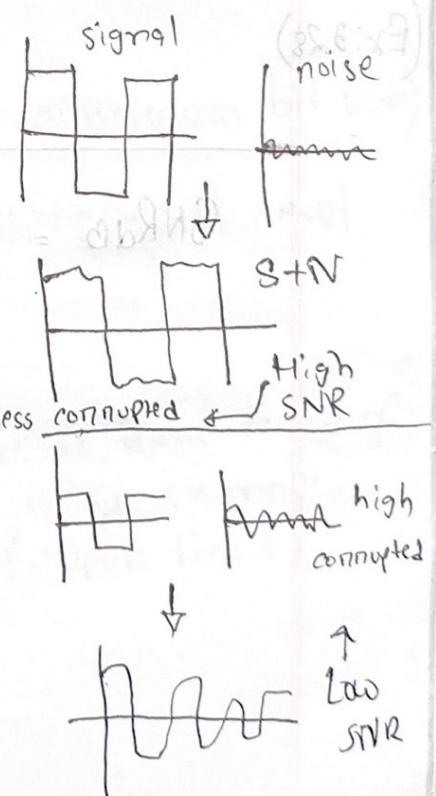
cross-talk

noise



Signal to Noise
Ratio (SNR).

$$\text{SNR} = \frac{\text{average signal power}}{\text{average noise power}}$$



(Ex 3.26) Signal \rightarrow end power reduced to one-half?

Calculate attenuation. or dB. or loss / Gain?

$$dB = 10 \log_{10} P_2/P_1$$

at -220173

$$= 10 \log_{10} \frac{0.5 P_1}{P_1}$$



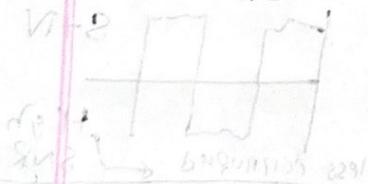
Previous power of P_1 .
end power = $P_1/2$
 P_1 → $P_2 = P_1/2$.

ratio of loss $= 3 \text{ dB}$

(3dB) \uparrow loss

(Ex 3.28), Level 2

$$SNR_{dB} = 10 \log_{10} SNR$$



Conversion करें

27A1

Data Rate Limit: $\text{Data rate} \leq \text{Bandwidth} \times \text{SNR}$

Data rate depends on 3 factors:

i) Bandwidth available

ii) Level of the signals

iii) Quality of the channel (level of noise).

Noiseless channel is theoretical and real life impossible.

Noiseless channel:

Nyquist Bit rate \rightarrow (theoretical maximum bit rate)

$\text{bitrate} = 2 \times \text{bandwidth} \times \log_2 \frac{\text{signal level}}{\text{noise level}}$

Noise Channel:

Shannon capacity \rightarrow (theoretical highest data rate of a noisy channel)

$$C = \text{bandwidth} \times \log_2 (1 + \text{SNR})$$

Ex: 3.41 Channel \rightarrow 1 MHz bandwidth
 $\text{SNR} = 63$ & no整形 short poll
 bitrate and signal level.

$$C = \text{bandwidth} \times \log_2(1 + \text{SNR})$$

$$= 10^6 \times \log_2(1 + 63)$$

$$\text{bitrate} = 6 \text{ Mbps}$$

$\left[\begin{array}{l} \text{Upper limit } 6 \text{ Mbps} \\ \text{Lower limit } 4 \text{ Mbps} \end{array} \right]$

\hookrightarrow Lower limit

$$\text{bitrate} = 2 \times \text{Bandwidth} \times \log_2 L \rightarrow \text{level}$$

$$6 \text{ Mbps} \leftarrow 4 \text{ Mbps} = 2 \times 10^6 \times \log_2 L$$

maths $\log_2 4 = ?$

$\log_2 4 = 2$ (Ans) $\log_2 4 = 2$

$2^2 = 4$

Ans,

3.6

Performance

~~Bandwidth~~ \rightarrow ~~Latency~~ \rightarrow ~~Throughput~~

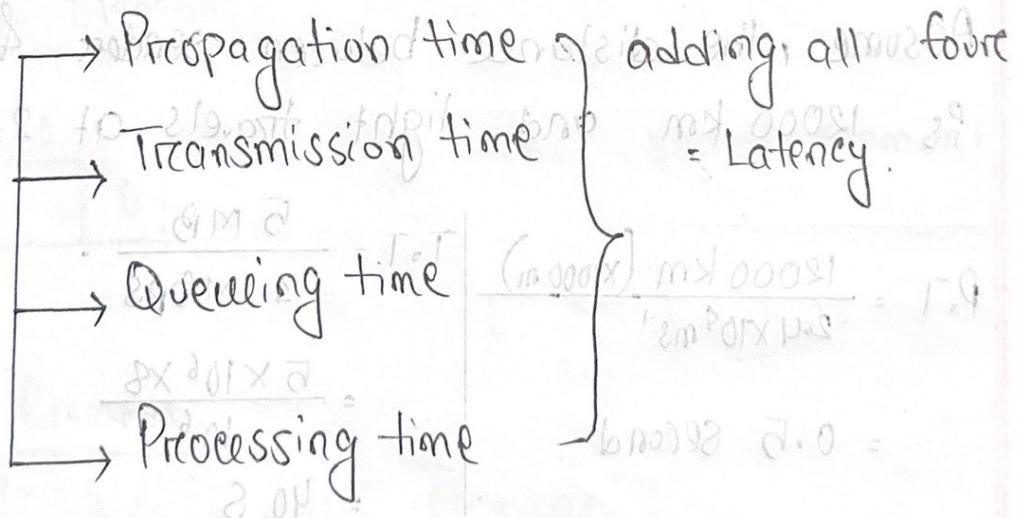
~~Bandwidth~~ \rightarrow $H_2 \uparrow$ $BPS \uparrow$ } Proportional to each other.
~~Bandwidth~~ \rightarrow $BPS \uparrow$ } Theoretical value.

$$\therefore \text{Bandwidth} = F_H - F_L \quad \left\{ \begin{array}{l} (\text{Highest} - \text{Lowest}) \text{ frequency} \\ \text{grit fesid} \end{array} \right.$$

~~Throughput~~: How fast data is transmitted.

$$BW = 10 \text{ Mbps}$$

$$\text{Throughput} = 7 \text{ Mbps}$$

Latency:

$$\left\{ \begin{array}{l} \text{Propagation} \leftarrow \text{distance} \times \text{propagation speed} \\ \text{Transmission} \leftarrow \text{frame size} \times \text{bit rate} \\ \text{Queuing} \leftarrow \text{average queue length} \times \text{processing time per frame} \\ \text{Processing} \leftarrow \text{fixed overhead} \end{array} \right\} \left\{ \begin{array}{l} \text{propagation speed} = \text{light speed} / \text{medium refractive index} \\ \text{bit rate} = \text{BW} / \text{overhead} \\ \text{processing time per frame} = \text{frame size} / \text{bit rate} \end{array} \right\}$$

11. Oct 19, 2019

8.9. output

$$\rightarrow \text{Propagation time} = \frac{\text{Distance}}{\text{Propagation Speed}}$$

$$\rightarrow \text{Transmission time} = \frac{\text{Message Size}}{\text{Bandwidth}}$$

$$\rightarrow \text{Queuing time} = \text{Packet Line}$$

$$\rightarrow \text{Processing time} = \text{Processor (Slow/fast)}$$

Ex-27 What are the Propagation time & transmission time for 5 Mbps message if the bandwidth of the network is 1 Mbps?

Assume, the distance between sender & Receiver is 12000 km and light travels at $2.4 \times 10^8 \text{ ms}^{-1}$.

$$P.T = \frac{12000 \text{ km} \times 1000 \text{ m}}{2.4 \times 10^8 \text{ ms}^{-1}}$$
$$= 0.5 \text{ second}$$
$$T.T = \frac{5 \text{ MB}}{1 \text{ MBps}}$$
$$= \frac{5 \times 10^6 \times 8}{10^6}$$
$$= 40 \text{ S}$$

$$\left\{ \begin{array}{l} 3.28, 3.18, \log_2 L \\ \text{imp} \end{array} \right\} \left\{ \begin{array}{l} \text{SNR} \& \text{SNR DB} \rightarrow \text{conversion} \\ 1 \text{ byte} \rightarrow 8 \text{ bits} \\ 1 \text{ character} \rightarrow 8 \text{ bits} \end{array} \right\}$$

(ii) formats Chapter 04



Digital Transmission

Digital to Digital Transmission

Analog to Digital Transmission



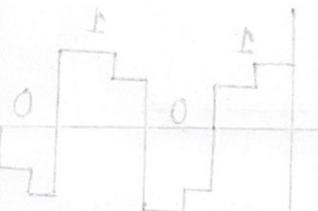
1/S = 3

D-D

→ Line coding

→ Block coding

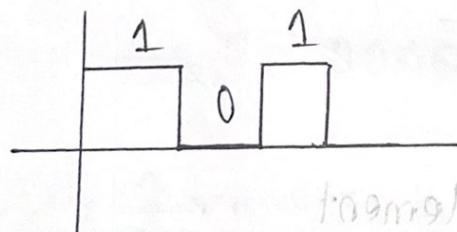
→ Scrambling



S/E = 3

■

Signal Element;



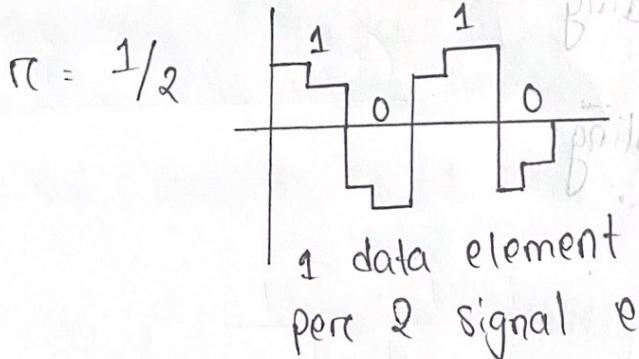
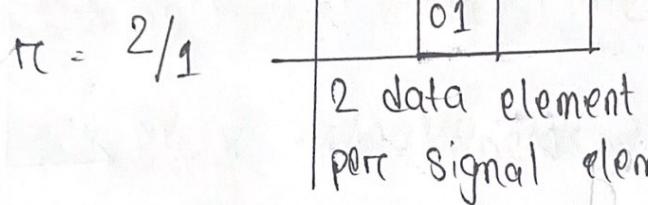
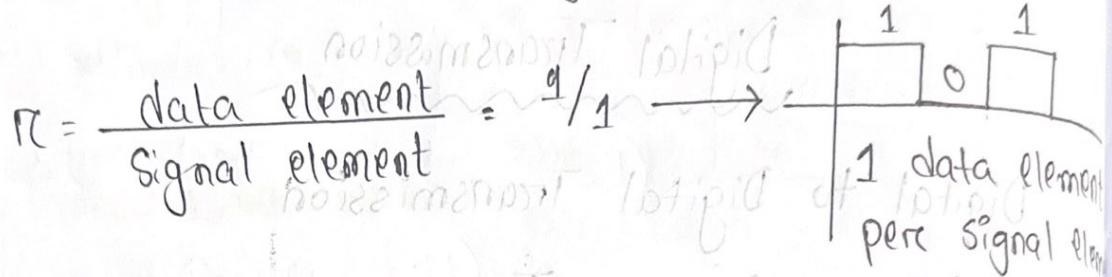
— = signal element

■

Data Element;

1/0 = data element

One data element per signal element (π)



Signal Rate:

- 1 second \rightarrow signal element

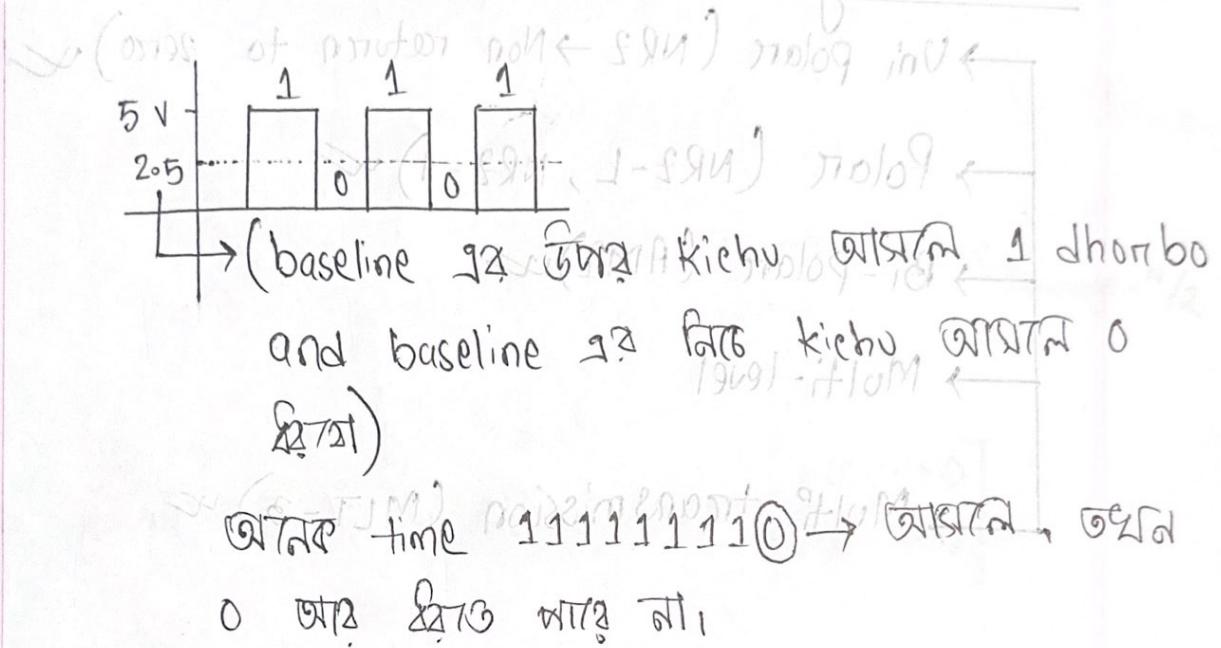
$$S = \frac{N}{\pi} \rightarrow$$

Signal rate \leftarrow Data rate
Pulse rate \leftarrow Proportion of signal element
Bandwidth \leftarrow data element

■ Data Rate: Same as bit rate.

Perc second \rightarrow how many bit.

Baseline Wandering:

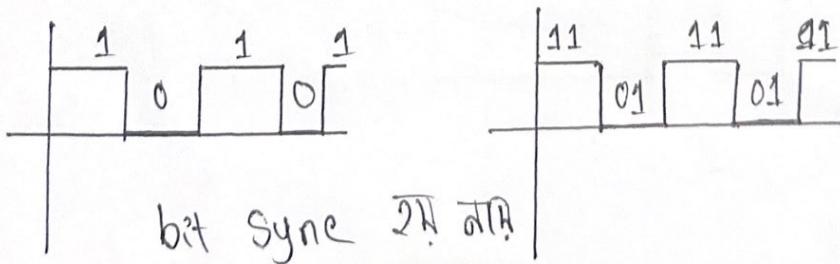


• DC Component:

frequency আনন্দ করে মান, voltage করে মান, যা
একই frequency আগলে ৫২৭৩ পাই না,

00000000001 → ৫২৭৩ পাই না

• Self Synchronization: Sender and Receiver এবং clock same time কৰে নি থাঁকলে Sync কৰে নি,

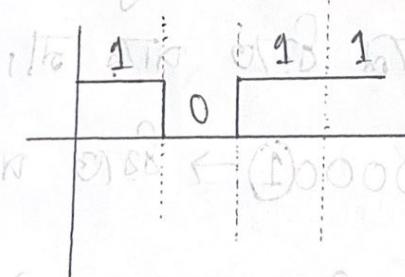


四

Line Coding

- Uni-polar (NRZ → Non return to zero)
 - Polar (NRZ-L, NRZ-I)
 - Bi-polar (AMI)
 - Multi-level
 - Multi transmission (MLT-3)

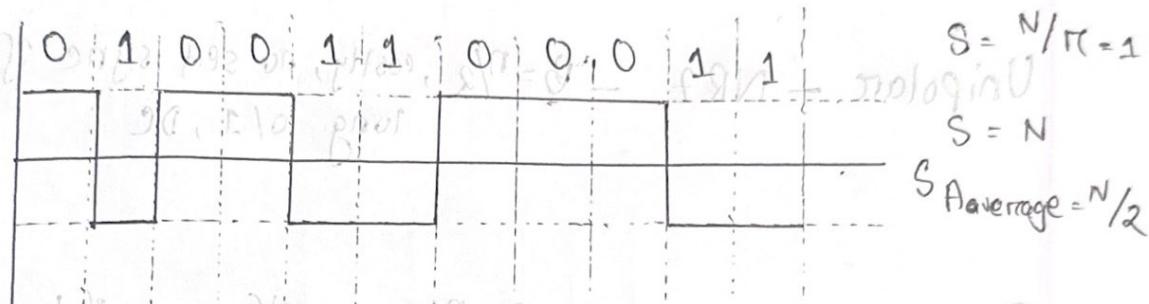
• Uni Polar (NRZ):



Lecture 09

Date: 29/02/2024

↳ NRZ-L & NRZ-I (Polar)

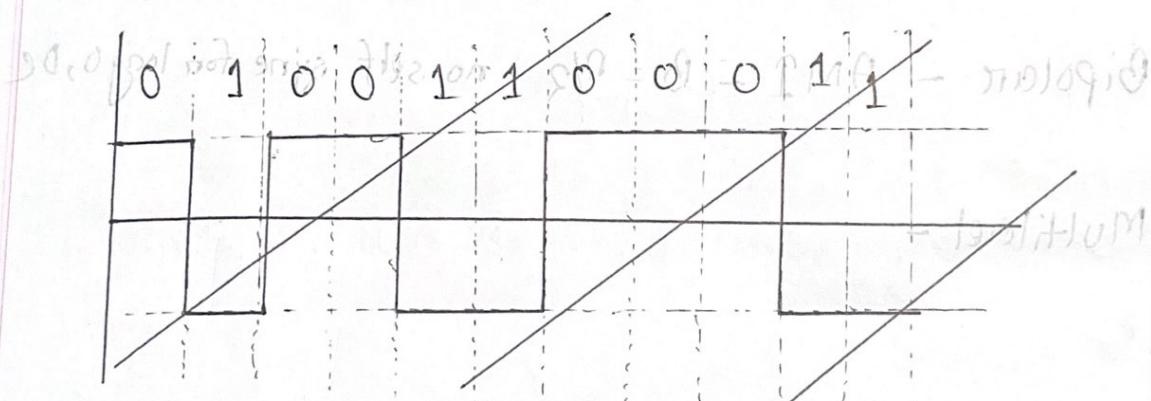


$\pi = \frac{\text{data element}}{\text{Signal element}} = \frac{1}{1} = 1$

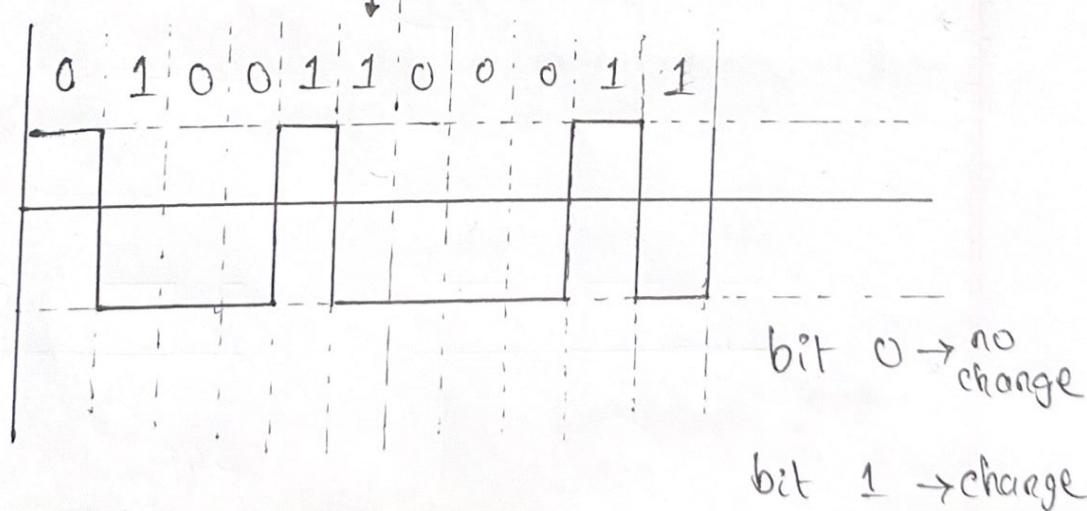
$\pi = 1$ for NRZ-L → level

$[+ve \rightarrow 0]$

$[-ve \rightarrow 1]$



NRZ-L → NRZ-I → Inversion



unipolar - NRZ

PO: 0.571721

Summary of line coding schemes (Table) 4.1

Unipolar - NRZ - $B = \frac{n}{2}$, costly, no self sync if long 0/1, DC



Polar - NRZ-L - $B = \frac{n}{2}$, no self sync if long 0/1, DC

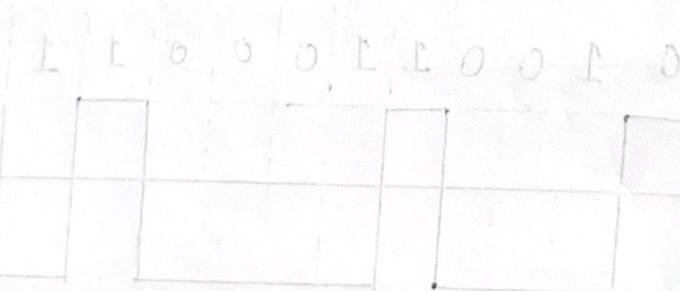
$\left[\begin{array}{l} \text{0 or 1} \\ \text{L} \leftarrow \text{NRZ-I} \end{array} \right] - B = \frac{n}{2}$ no self sync for long 0, DC

Bipolar - AMI - $B = \frac{n}{2}$ no self sync for long 0, DC

Multilevel -



Multi transmission - MLT-3 - $B = \frac{n}{3}$ no self sync for long 0.



0.571721

0.571721

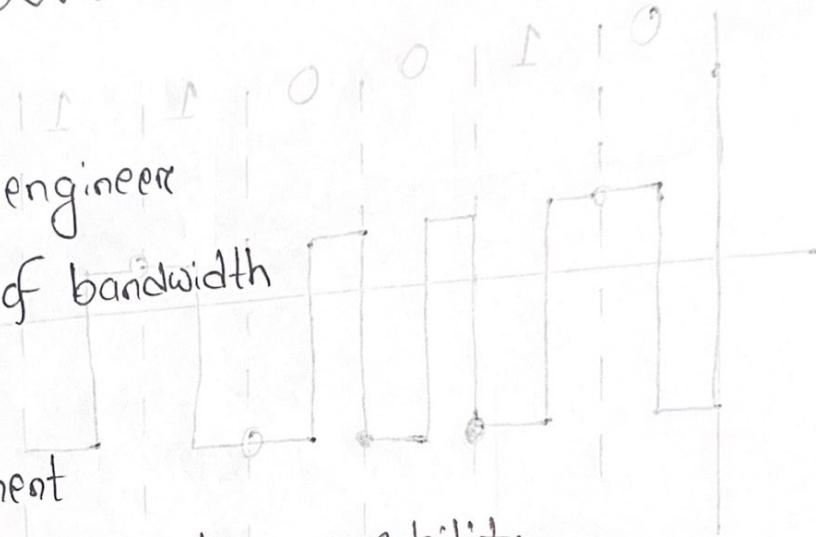
NRZ - Pros and Cons

Pros

- Easy to engineer
- good use of bandwidth

Cons

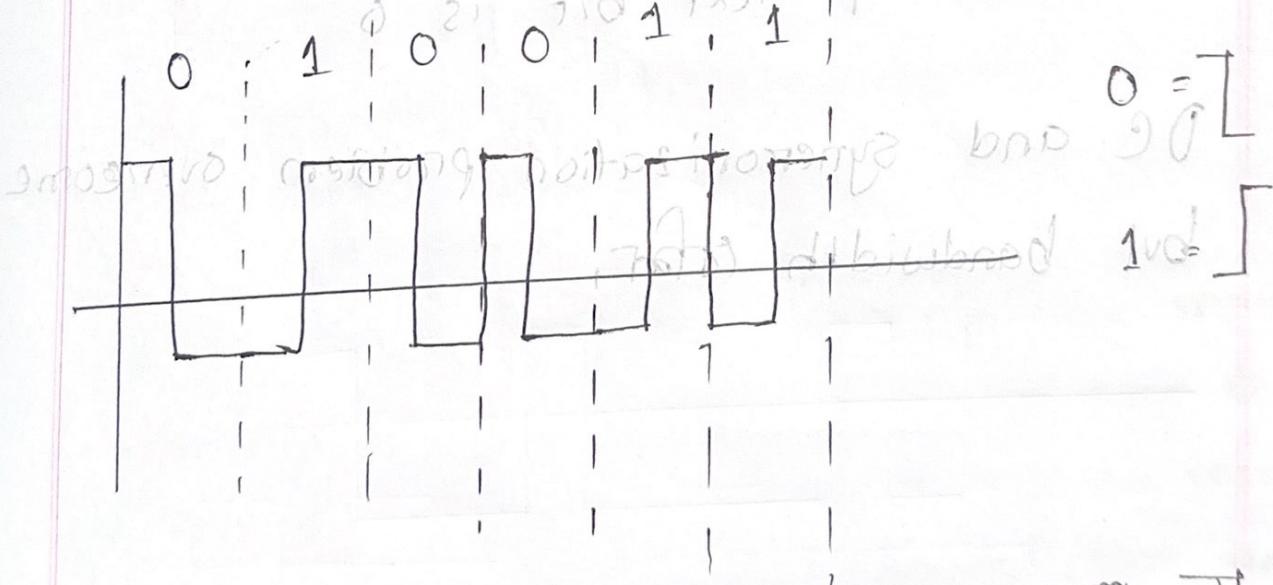
- DC component
- lack of synchronization capability.



Bipolar (Manchester and Differential Manchester)

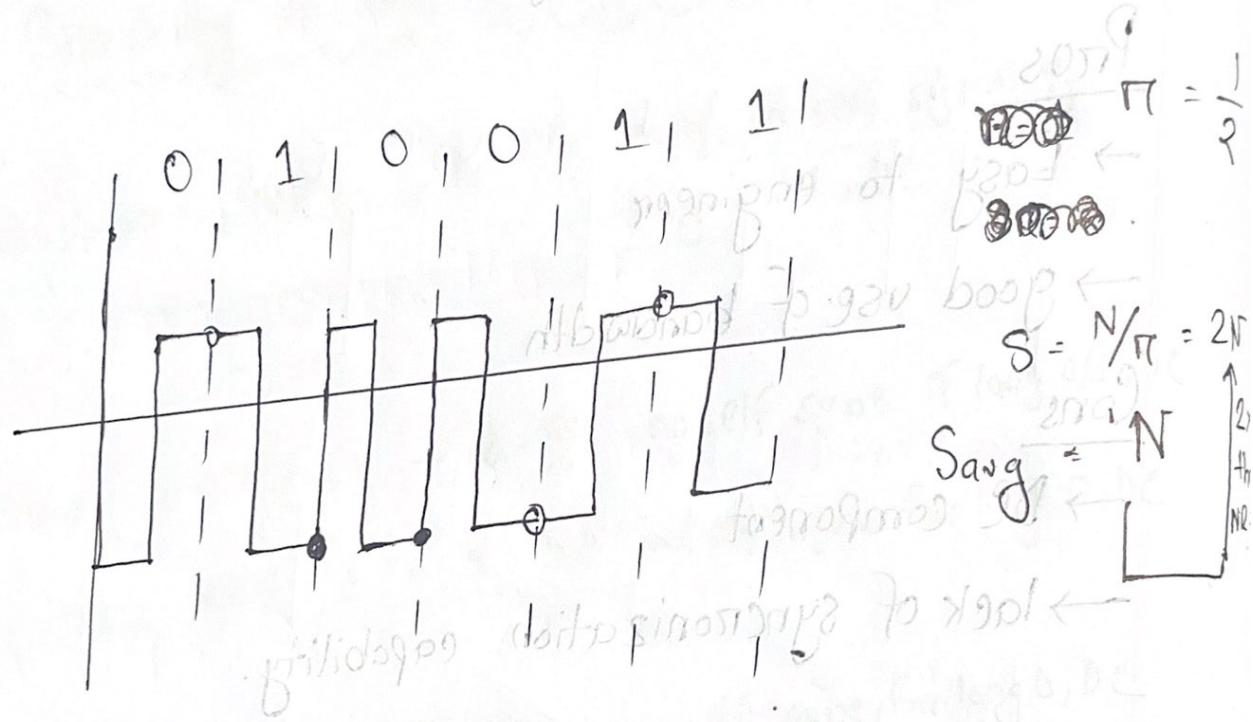
MAN

DMAN



Manchester

DC Problem 272



D-manchester

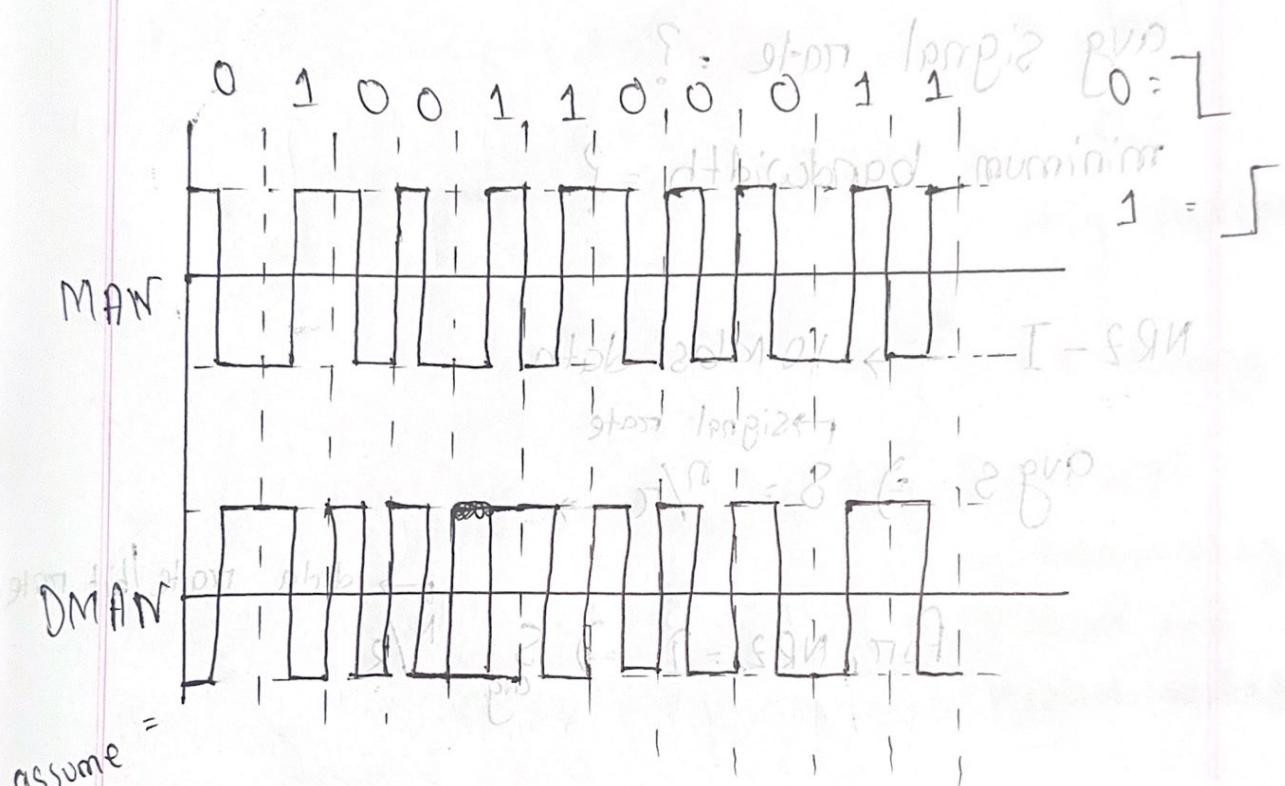
No inversion if next bit is 1

- Inversion if next bit is 0

DC and synchronization problem overcome

but bandwidth is lost

Biphase (MAN and DMAN)



Pros

- Sync on mid transition
- no dc
- error detection

Cons

- twice NRZ max
- require more b.w.
- at least one transition per bit time and possibly two.

Ex-4.4 math (NAMO bho NAM) geostiqid

avg signal rate = ?

minimum bandwidth = ?

NRZ-I → 10 Mbs data

$$\text{avg } S \Rightarrow S_{\text{avg}} = \frac{n}{\pi} \rightarrow \text{signal rate}$$

$$\text{For NRZ-I} \Rightarrow S_{\text{avg}} = \frac{N}{2} \rightarrow \text{data rate/bit rate}$$

$$\therefore S_{\text{avg}} = \frac{10}{2} = 500 \text{ kbps}$$

↳ noiseless or
minimum

$$\text{bandwidth} = S_{\text{avg}} = 500 \text{ Hz}$$

more bits ↪

minimum bit no. ↪

more power required ↪

9b or ↪

more noise ↑ ↪

noised noise ↪

less power

more bits ↓ ↪