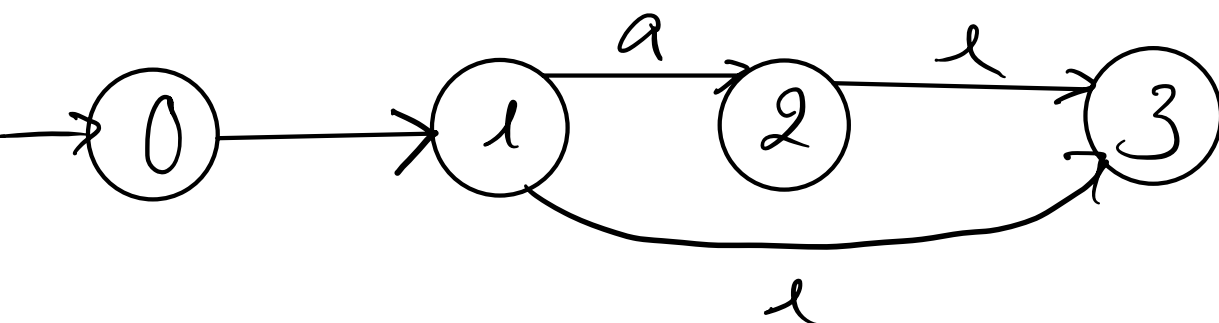


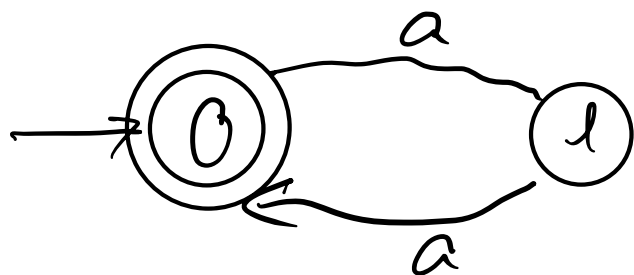
# TD2

Ex4:

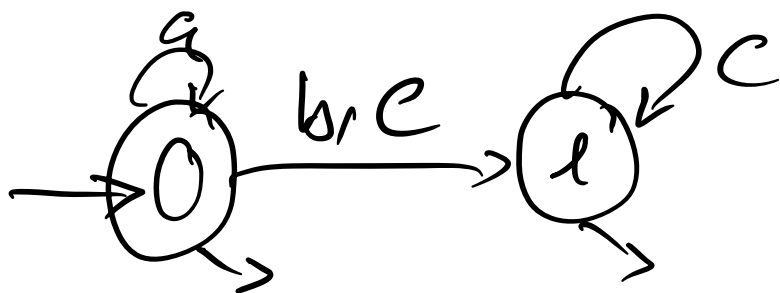
1.  $L = \{car, bar, or\}$



2.  $L = \{w \in \{a\}^* ; |w|_a \equiv 0 [2]\}$

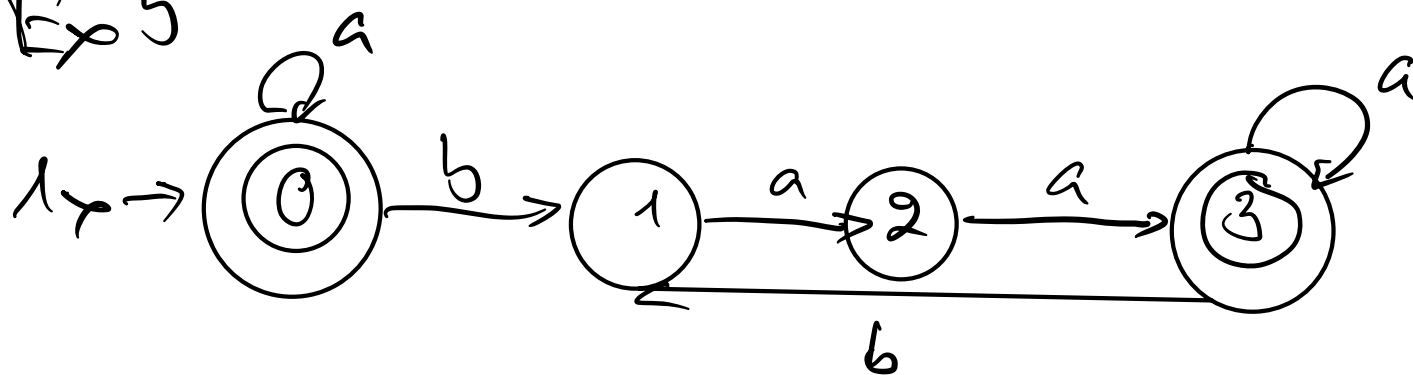


3.  $L = (a^*(b + \epsilon)c^*)$

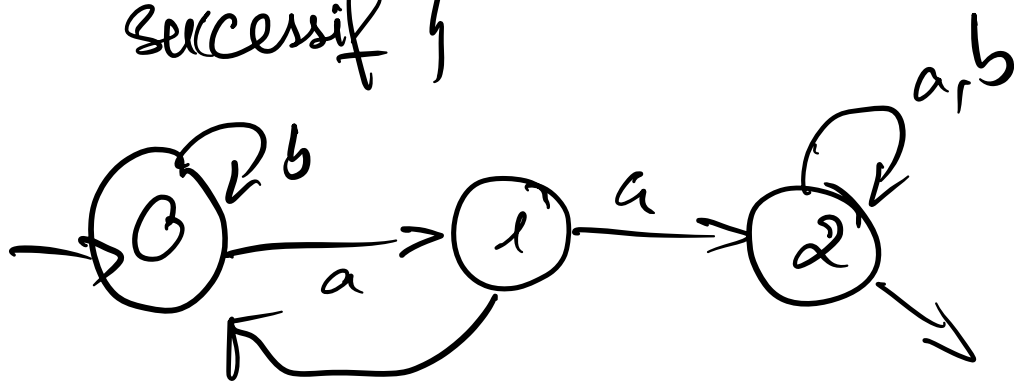


$$(a^*(\epsilon + b))c^* = (a^* + a^*b)c^* = a^*c^* + a^*bc^*$$

Ex 5

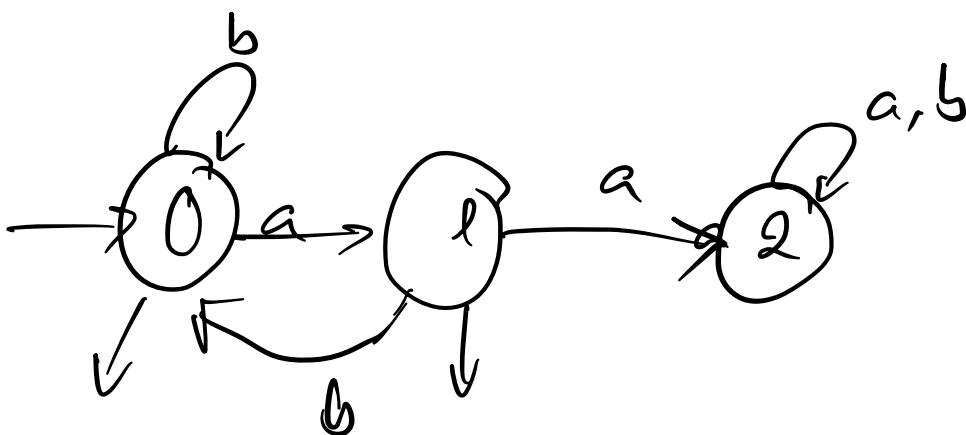


2,  $\mathcal{L}'_2 = \{u \in \{a,b\}^* \mid u \text{ contient deux } a \text{ successif}\}$



$\bar{\mathcal{L}} = \Sigma^* \setminus \mathcal{L}$  tout les alphabets de sigma privés de  $\mathcal{L}$

$\forall w \in \Sigma^* \quad u \in \bar{\mathcal{L}} \Rightarrow w \notin \mathcal{L}$



Algorithme pour compléter un automate 4

1<sup>o</sup> Compléter A

2<sup>o</sup> ~~Passer~~ les états acceptants et non acceptants dans A