

# Knights, Knaves, and Logical Reasoning Mechanising the Laws of Thought

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<sup>&</sup>lt;sup>1</sup>Special thanks to Francis Southern

## Introduction

**Thinking** 

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Formalising

⇓

Modelling

⇓

Computing

## **Thinking**

#### A Puzzle

You are on a strange island where people are divided into

- · Knights always saying the truth
- · Knaves always saying lies

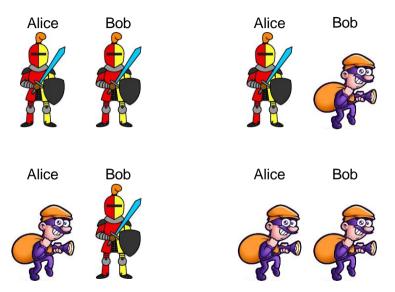
You meet two natives of the island Alice and Bob, and ask them

"Are you knights or knaves?"

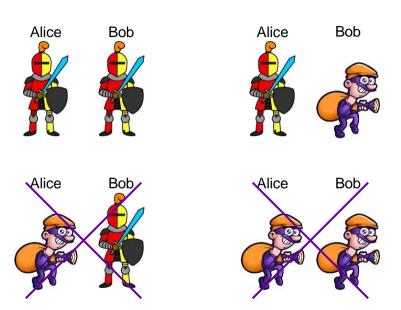
Alice answers "At least one of us is a knave"

What are Alice and Bob?

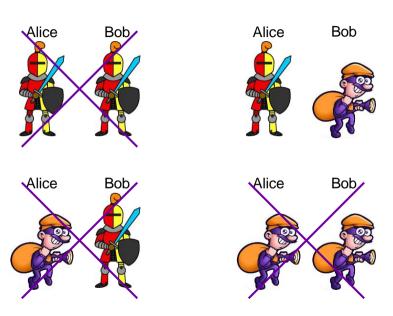
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## Formalising

## Formalising Correct Reasoning

A: Socrates is a man

B: All men are mortal

C: All men are Socrates C: Socrates is mortal

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Woody Allen - Love and Death Aristotle

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Linguistic, philosophical, or mathematical approaches to formalisation

Today: Propositional Logic

## **Propositions**

An expression which is either true or false.

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Proposition test: Is it true that...?

- $\cdot 2 + 2 = 5$
- Manchester
- Grass is green
- · We're in Manchester
- What's your name?
- It's raining

## Not $-\neg$ , And -&

#### Not

p	$\neg p$
F	Т
Т	F

It's not raining

Grass is not green.

## Not $-\neg$ , And -&

#### Not

p	$\neg p$
F	Т
Т	F

It's not raining

Grass is not green.

#### And

p	q	p & q
F	F	F
F	Т	F
Т	F	F
Т	Т	Т

Grass is green and it's raining.

We're in Manchester and we're in France.

## Or $- \mid$ , Implication (If, then) $- \rightarrow$

Or

p	q	$p \mid q$
F	F	F
F	Т	Т
Т	F	Т
Т	Т	Т

Take an aspirin or lie down.

You can have milk or sugar in your tea.

## Or $- \mid$ , Implication (If, then) $- \rightarrow$

#### Or

p	q	$p \mid q$
F	F	F
F	Т	Т
Т	F	Т
Т	Т	Т

Take an aspirin or lie down.

You can have milk or sugar in your tea.

#### **Implication**

p	q	$p \rightarrow q$
F	F	Т
F	Т	Т
Т	F	F
Т	Т	Т

If you get 90% on this assignment, then you'll pass the course.

If you're late, then you'll give me a fiver.

## Biimplication (If and only if) – ↔

#### Biimplication

p	q	$p \leftrightarrow q$
F	F	Т
F	Т	F
Т	F	F
Т	Т	Т

I'll buy you a new wallet if (and only if) you need one.

He studies if (and only if) he can.

p	q	r	(p & q)	$(p \& q) \rightarrow r$
F	F	F		
F	F	Т		
F	Т	F		
F	Т	Т		
Т	F	F		
Т	F	Т		
Т	Т	F		
Т	Т	Т		

p	q	r	(p & q)	$(p \& q) \rightarrow r$
F	F	F	F	
F	F	Т	F	
F	Т	F	F	
F	Т	Т	F	
Т	F	F	F	
Т	F	Т	F	
Т	Т	F		
Т	Т	Т		

p	q	r	(p & q)	$(p \& q) \rightarrow r$
F	F	F	F	
F	F	Т	F	
F	Т	F	F	
F	Т	Т	F	
Т	F	F	F	
Т	F	Т	F	
Т	Т	F	Т	
Т	Т	Т	Т	

p	q	r	(p & q)	$(p \& q) \to r$
F	F	F	F	Т
F	F	Т	F	Т
F	Т	F	F	Т
F	Т	Т	F	Т
Т	F	F	F	Т
Т	F	Т	F	Т
Т	Т	F	Т	
Т	Т	Т	Т	

p	q	r	(p & q)	$(p \& q) \to r$
F	F	F	F	Т
F	F	Т	F	Т
F	Т	F	F	Т
F	Т	Т	F	Т
Т	F	F	F	Т
Т	F	Т	F	Т
Т	Т	F	Т	F
Т	Т	Т	Т	

q	r	(p & q)	$(p \& q) \to r$		
F	F	F	Т		
F	Т	F	Т		
Т	F	F	Т		
Т	Т	F	Т		
F	F	F	Т		
F	Т	F	Т		
Т	F	Т	F		
Т	Т	Т	Т		
	F F T T F F	F F T T F F T F F F F T	F F F F F F F F F F F F T F T F		

## Modelling

 $k_A$  = Alice is a knight  $\neg k_A$  = Alice is a knave "Alice says X" is the same as  $k_A \longleftrightarrow X$ 

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- "at least one of us is a knave"
- · "I'm a knave or Bob is a knave"
- $\neg k_A \mid \neg k_B$

$$\Rightarrow k_A \longleftrightarrow (\neg k_A | \neg k_B)$$

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F	F	Т	Т	Т	
F	Т	Т	F	Т	
Т	F	F	Т	Т	
Т	Т	F	F	F	

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F	Т	Т	F	Т	
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F	Т	Т	F	Т	F
Т	F	F	Т	Т	
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F	Т	Т	F	Т	F
Т	F	F	Т	Т	Т
Т	Т	F	F	F	

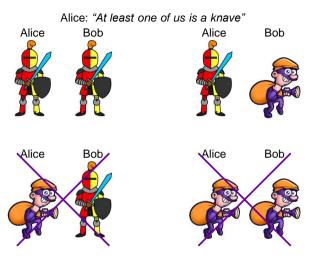
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$k_A$	$k_B$	$\neg k_A$	$\neg k_B$	$\neg k_A \mid \neg k_B$	$k_A \leftrightarrow (\neg k_A \mid \neg k_B)$
F	F	Т	Т	Т	F
F	Т	Т	F	Т	F
Т	F	F	Т	Т	Т
Т	Т	F	F	F	F

## From Solving to Modelling





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Alice: "At least one of us is a knave"

 $k_A$  = Alice is a knight

The trick: "Alice says X" is the same as  $k_A \leftrightarrow X$ 

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Alice says "At least one of us is a knave" =  $k_A \leftrightarrow (\neg k_A \mid \neg k_B)$ 



## From Solving to Modelling

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"At least one of us is a knave" =  $\neg k_A \mid \neg k_B$ 

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It can be (really) hard, but you only have to do it once!



## Modelling a Sudoku

			7			4	1	
		3		2				6
1		7	4			5	2	3
4		1	6				8	
	2	9		7		6	3	
	7				4	2		1
7	5	2			6	3		9
3				4		1		
	1	4			3			

What propositions do we need?

			7			4	1	
		3		2				6
1		7	4			5	2	3
4		1	6				8	
	2	9		7		6	3	
	7				4	2		1
7	5	2			6	3		9
3				4		1		
	1	4			3			

What propositions do we need?

Number n is in row i and column j

- number 7 is in row I and column 4
- number 2 is in row 6 and column 7

		7			4	1	
	3		2				6
	7	4			5	2	3
	1	6				8	
2	9		7		6	3	
7				4	2		1
5	2			6	3		9
			4		1		
1	4			3			
	7 5	7 1 2 9 7 5 2	3 7 4 1 6 2 9 7 5 2	3	3 2 2 7 4 7 4 7 7 7 7 7 7 7 7 7 7 7 7 7 7	3     2       7     4       1     6       2     9       7     4       2     9       7     4       2     4       3     4       4     1	3     2     5       7     4     5       1     6     6       2     9     7     6       3     3     4     2       5     2     6     3       4     4     1

What propositions do we need?

Number n is in row i and column j

- number 7 is in row I and column 4
- number 2 is in row 6 and column 7

- at least one number per cell  $(p_{1,1,4} | \dots | p_{9,1,4})$
- at most one number per cell  $(p_{7,1,4} \rightarrow \neg p_{1,1,4}, p_{7,1,4} \rightarrow \neg p_{2,1,4})$

			7			4	1	
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1		7	4			5	2	3
4		1	6				8	
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- no number can be repeated in a row

			7			4	1	
		3		2				6
1		7	4			5	2	3
4		1	6				8	
	2	9		7		6	3	
	7				4	2		1
7	5	2			6	3		9
3				4		1		
	1	4			3			

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- no number can be repeated in a row/column

			7			4	1	
		3		2				6
1		7	4			5	2	3
4		1	6				8	
	2	9		7		6	3	
	7				4	2		1
7	5	2			6	3		9
3				4		1		
	1	4			3			

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- no number can be repeated in a row/column/region

# Computing

# **Automating the Process**

#### Truth table

- mechanical
- time consuming (2<sup>n</sup> rows!)
- tedious

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- mechanical
- time consuming (2<sup>n</sup> rows!)
- tedious

## Let a computer do it for you!

- · ideal for mechanical tasks
- only needs an input formula
- more reliable than us
- · much faster than us
- the output is easily customisable

# **Automated Reasoning**

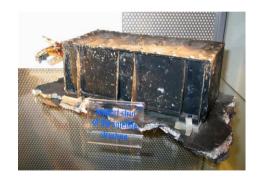
## Much more than solving puzzles!

- software and hardware verification
   Intel and Microsoft
- information management biomedical ontologies, Semantic Web, databases
- combinatorial reasoning constraint satisfaction, planning, scheduling
- Internet security
- theorem proving in mathematics

## Where Could Have Been Used

Ariane 5 rocket failure due to a software bug, cost \$370 million.





## Where Has Been Used

IDK / IDK-8072909

### To find and fix a bug in a widely used sorting algorithm!



## Where Has Been Used

To find and fix a bug in a widely used sorting algorithm!



Even Amazon and Facebook use automated reasoning techniques!

# **Automated Reasoning Competitions**

- The CADE ATP System Competition (CASC)
- OWL Reasoning Competition (ORE)
- SAT-Race



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You can bet on the winner!

## Do You Want to Know More?

Look at the references on the handout!