

ANALYSIS OF ALGORITHMS

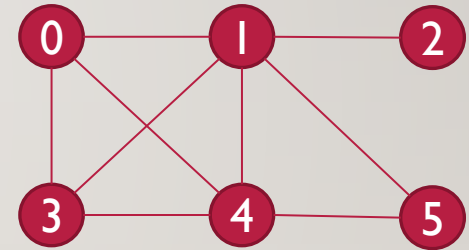
LECTURE 10 : SHORTEST PATH TREES

BASED ON SECTION 6.1



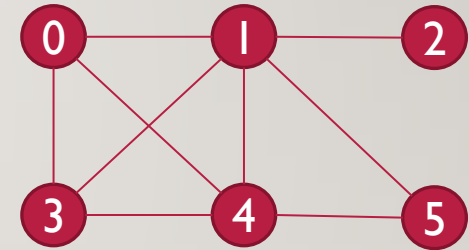
SHORTEST PATH

- Let's say you wanted to find the shortest path from some vertex to every other vertex
- What would you do?

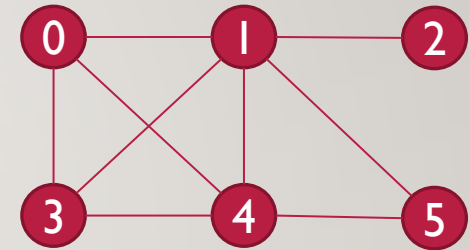
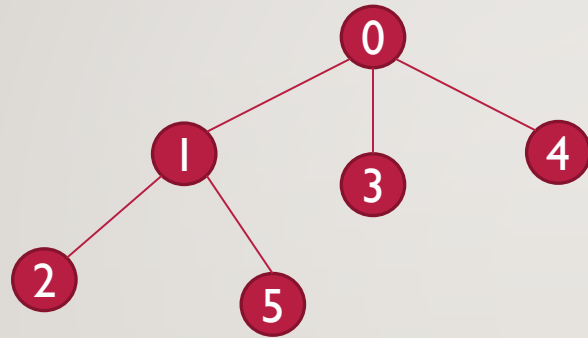


SHORTEST PATH

- Let's say you wanted to find the shortest path from some vertex to every other vertex
- What would you do?
- Breadth first search!
 - Because it tries to make shorter, wider trees, it is actually finding the shortest path from the root to anywhere else



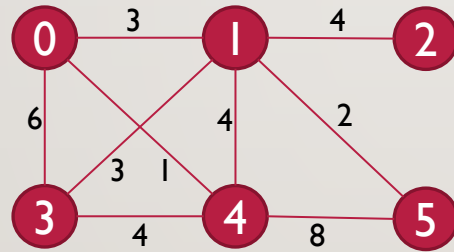
BREADTH FIRST SEARCH TREE



- Consider if there was a shorter path from 0 to 2
- That would have to be the edge 0,2
- But then that would have been in the tree
- The tree contains the shortest path from the root to any vertex

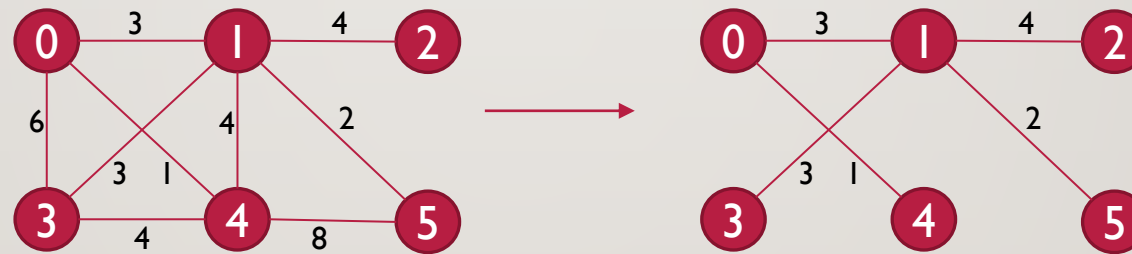
WEIGHTED GRAPHS

- What do we do if the graph is weighted?
- Minimum Weighted Spanning Tree?



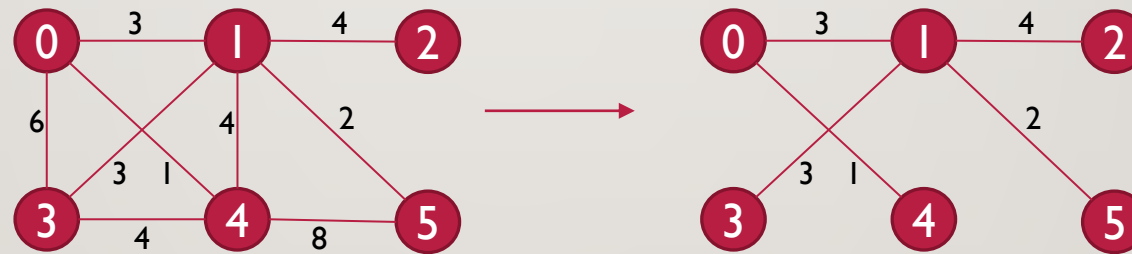
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WEIGHTED GRAPHS

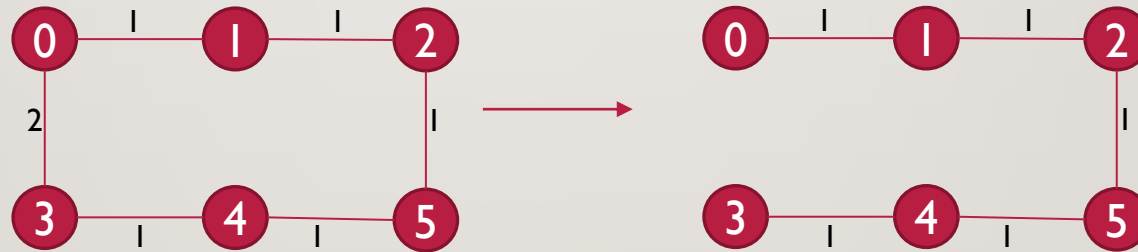
- What do we do if the graph is weighted?
- Minimum Weighted Spanning Tree?



- Finds the smallest total cost for a connected subtree, not the shortest path
- Consider the path between vertex 3 and vertex 4

WEIGHTED GRAPHS

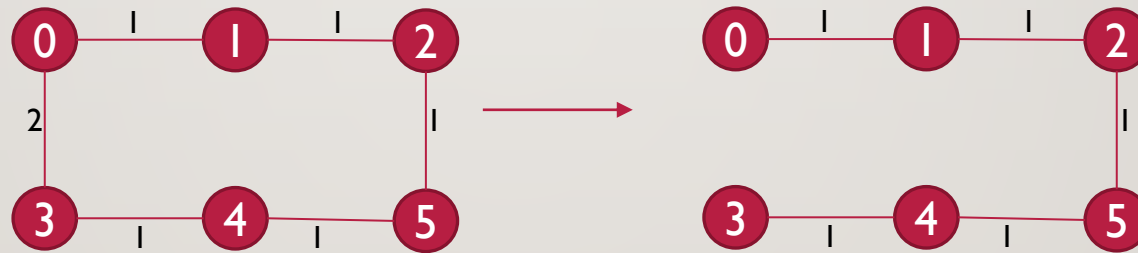
- What do we do if the graph is weighted?
- Minimum Weighted Spanning Tree?



- Finds the smallest total cost for a connected subtree, not the shortest path
- Consider the path between vertex 0 and vertex 3

WEIGHTED GRAPHS

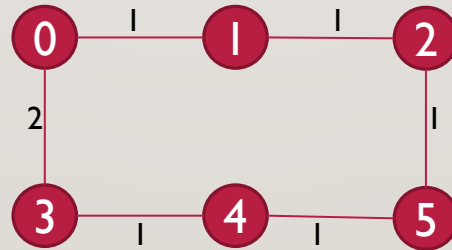
- What's wrong and how do we fix it?



- The MWST algorithm is greedy. It only cares what the cheapest edge right now is
- We need to keep track of the cheapest edge **cumulatively**

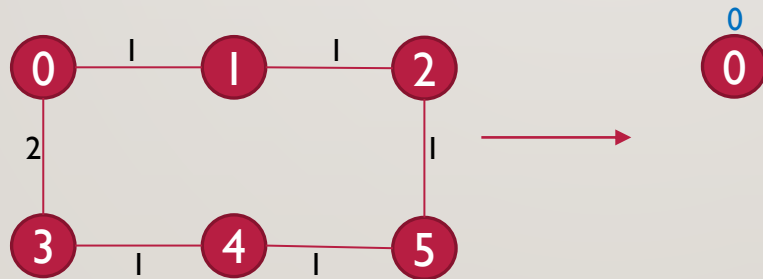
DIJKSTRA'S ALGORITHM

- Pretty much the same as the Prim's algorithm (The minimum weighted spanning tree algorithm we covered in class)
- But now instead of just picking the edge with the lowest cost, we pick the edge with the lowest cumulative cost.
- That means we need to keep track of the cumulative cost to each vertex



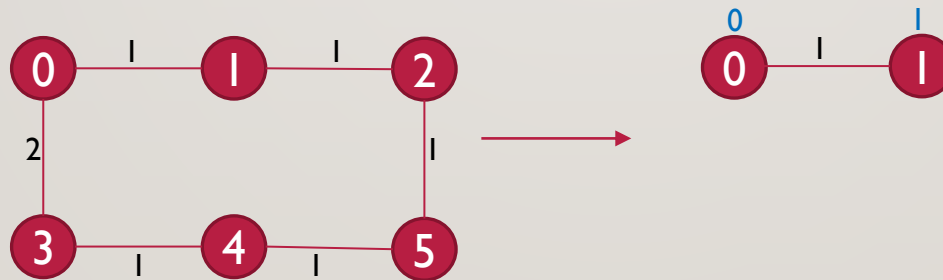
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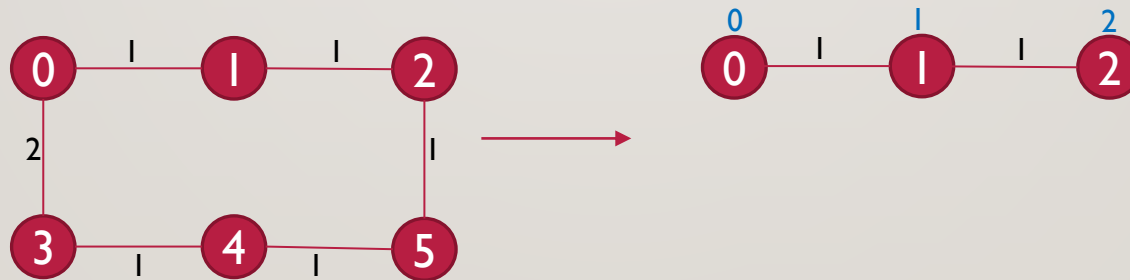
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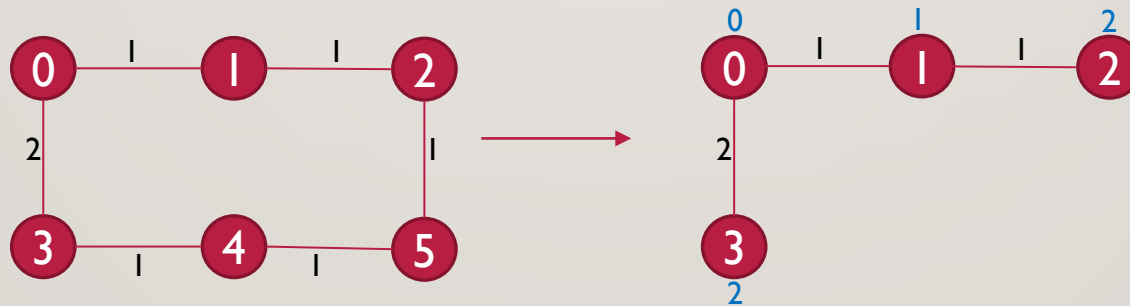
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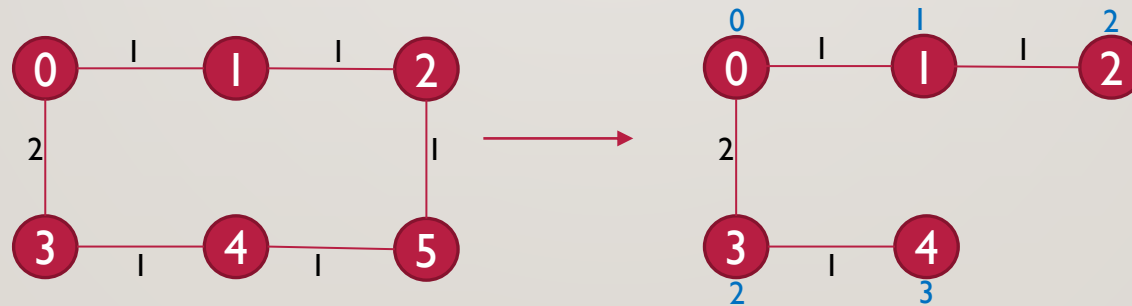
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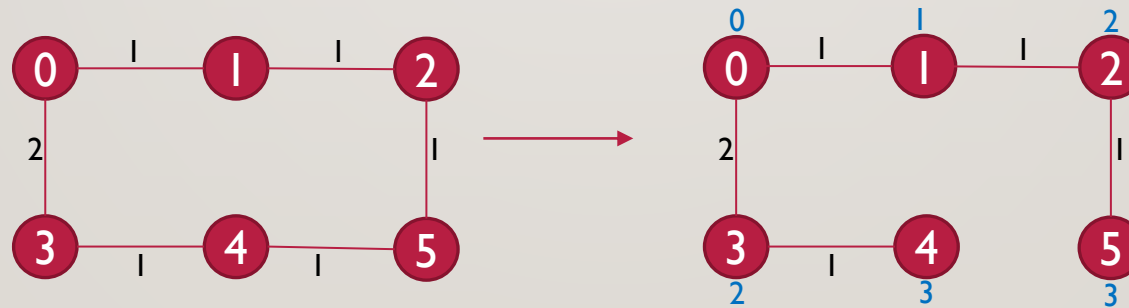
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DIJKSTRA'S ALGORITHM

- Start off with the root vertex
- Repeatedly
 - Add the edge with the lowest cumulative cost from the original graph going from a vertex in the tree to one not in the tree
 - Stop when you have all the vertices from original tree

DIJKSTRA'S ALGORITHM

- Start off with the root vertex
- Repeatedly
 - Add the edge with the lowest cumulative cost from the original graph going from a vertex in the tree to one not in the tree
 - Stop when you have all the vertices from original tree
- Data structures needed
 - Parent array – stores the parent of each vertex
 - Cost array – keeps best known cost to each vertex
 - Marked array – stores whether a vertex is in the tree

Cost	0	1	2	3	4	5

Parent	0	1	2	3	4	5

Marked	0	1	2	3	4	5

DIJKSTRA'S ALGORITHM

- $\text{marked}[\text{root}] = \text{true}$, $\text{parent}[\text{root}] = \text{root}$, $\text{cost}[\text{root}] = 0$
- Set $\text{cost}[v] = \text{cost}(\text{root}, v)$ for all v
- Repeatedly
 - Find the unmarked vertex y with the lowest cost
 - Set $\text{marked}[y] = \text{true}$
 - For every unmarked neighbour z
 - The current cost of getting to z is in the cost array
 - By finding the edge (y, z) we've found another possible cost. In this case $\text{cost}(z) = \text{cost}(y) + \text{cost}(y, z)$
 - Update the cost and parent of z if it's lower than the current known cost or the cost is unknown
- Stop when all vertices are in the tree

Cost	0	1	2	3	4	5
	-1	-1	-1	-1	-1	-1

Parent	0	1	2	3	4	5
	-1	-1	-1	-1	-1	-1

Marked	0	1	2	3	4	5
	f	f	f	f	f	f

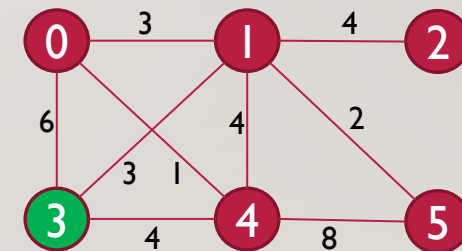
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	-1	-1	-1	-1	-1	-1

Marked	0	1	2	3	4	5
	f	f	f	f	f	f



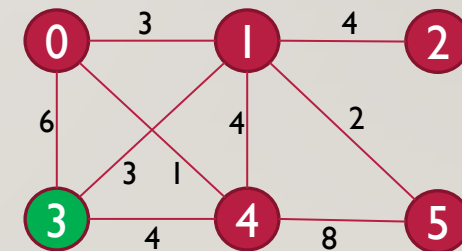
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Parent	0	1	2	3	4	5
	-1	-1	-1	3	-1	-1

Marked	0	1	2	3	4	5
	f	f	f	f	f	f



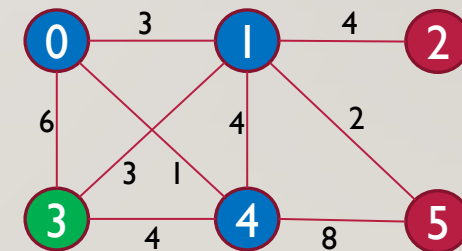
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Marked	0	1	2	3	4	5
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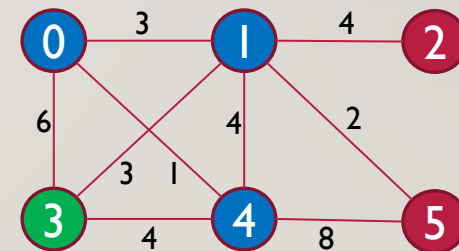
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Cost	0	1	2	3	4	5
	6	-1	-1	0	-1	-1

Parent	0	1	2	3	4	5
	3	-1	-1	3	-1	-1

Marked	0	1	2	3	4	5
	f	f	f	t	f	f



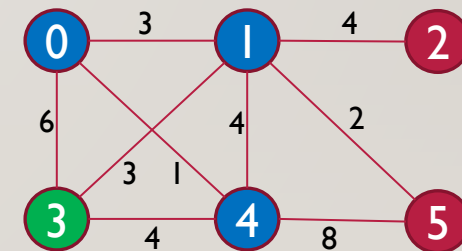
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Cost	0	1	2	3	4	5
	6	3	-1	0	-1	-1

Parent	0	1	2	3	4	5
	3	3	-1	3	-1	-1

Marked	0	1	2	3	4	5
	f	f	f	t	f	f



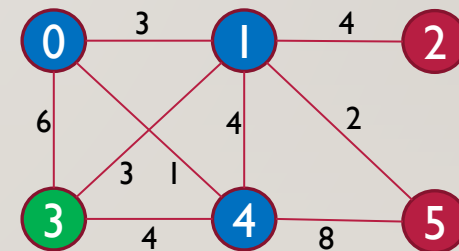
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	3	3	-1	3	3	-1

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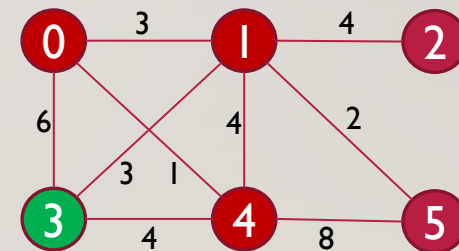
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	f	f	f	t	f	f



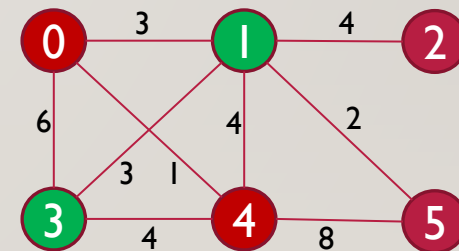
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	3	3	-1	3	3	-1

Marked	0	1	2	3	4	5
	f	t	f	t	f	f



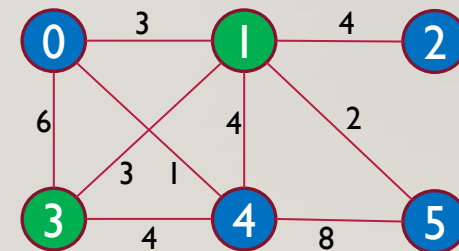
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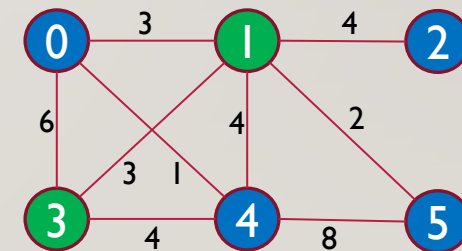
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Cost	0	1	2	3	4	5
	6	3	7	0	4	-1

Parent	0	1	2	3	4	5
	3	3	1	3	3	-1

Marked	0	1	2	3	4	5
	f	t	f	t	f	f



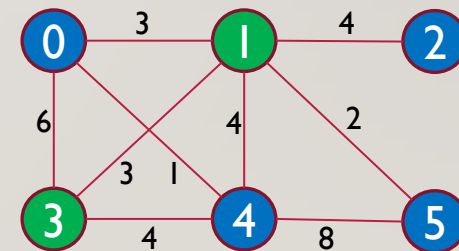
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	f	t	f	t	f	f



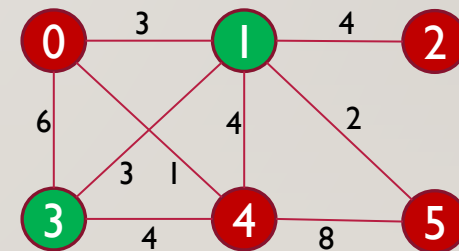
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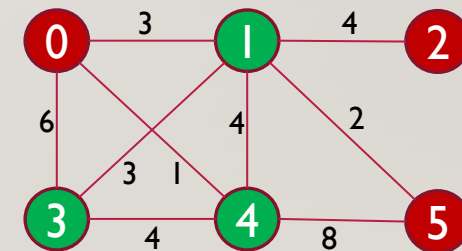
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	f	t	f	t	t	f



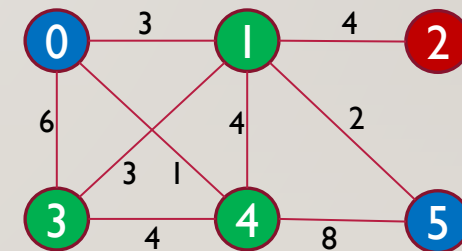
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 - By finding the edge (y, z) we've found another possible cost. In this case $\text{cost}(z) = \text{cost}(y) + \text{cost}(y, z)$
 - Update the cost and parent of z if it's lower than the current known cost or the cost is unknown
- Stop when all vertices are in the tree

Cost	0	1	2	3	4	5
	6	3	7	0	4	5

Parent	0	1	2	3	4	5
	3	3	1	3	3	1

Marked	0	1	2	3	4	5
	f	t	f	t	t	f



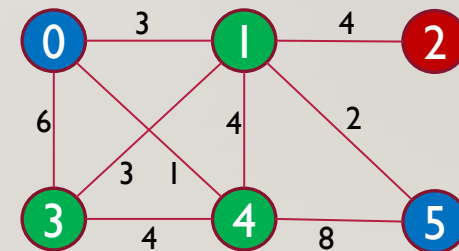
DIJKSTRA'S ALGORITHM

- $\text{parent}[\text{root}] = \text{root}$, $\text{cost}[\text{root}] = 0$
- Set $\text{cost}[v] = \text{cost}(\text{root}, v)$ for all v
- Repeatedly
 - Find the unmarked vertex y with the lowest cost
 - Set $\text{marked}[y] = \text{true}$
 - For every unmarked neighbour z
 - The current cost of getting to z is in the cost array
 - By finding the edge (y, z) we've found another possible cost. In this case $\text{cost}(z) = \text{cost}(y) + \text{cost}(y, z)$
 - Update the cost and parent of z if it's lower than the current known cost or the cost is unknown
- Stop when all vertices are in the tree

Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	f	t	f	t	t	f



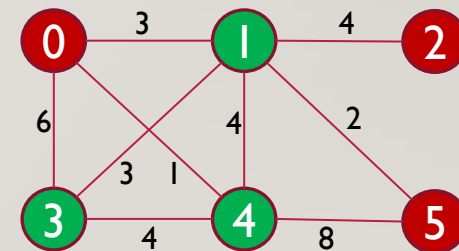
DIJKSTRA'S ALGORITHM

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Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	f	t	f	t	t	f



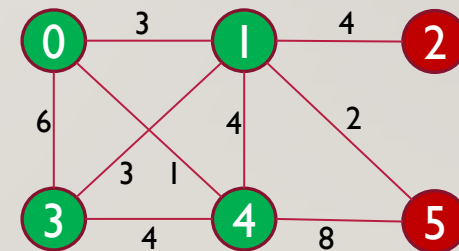
DIJKSTRA'S ALGORITHM

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Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	f	t	t	f



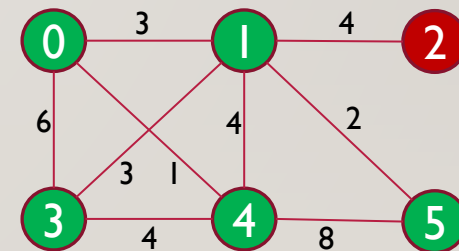
DIJKSTRA'S ALGORITHM

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	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	f	t	t	t



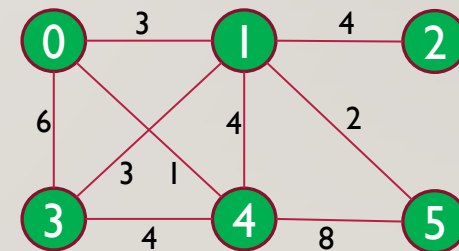
DIJKSTRA'S ALGORITHM

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- Stop when all vertices are in the tree

Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	t	t	t	t



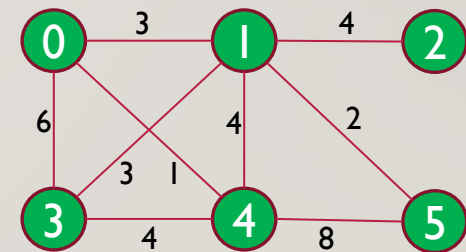
DIJKSTRA'S ALGORITHM

3

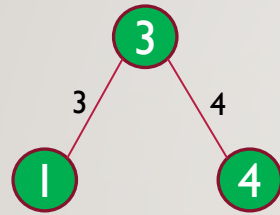
Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	t	t	t	t



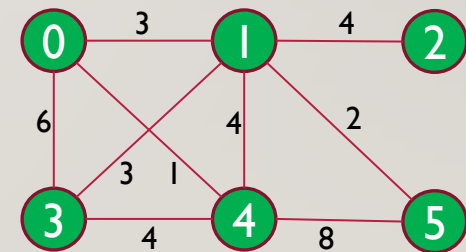
DIJKSTRA'S ALGORITHM



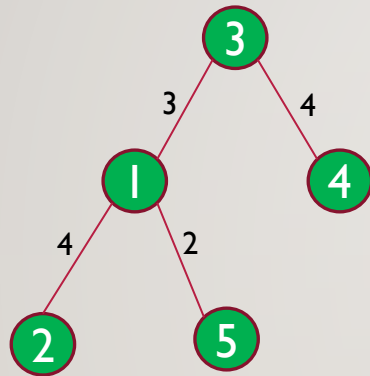
Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	t	t	t	t



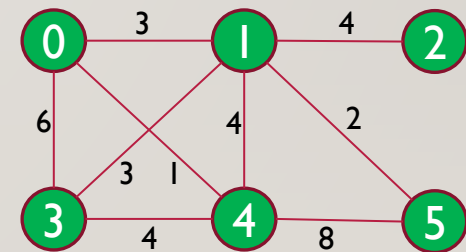
DIJKSTRA'S ALGORITHM



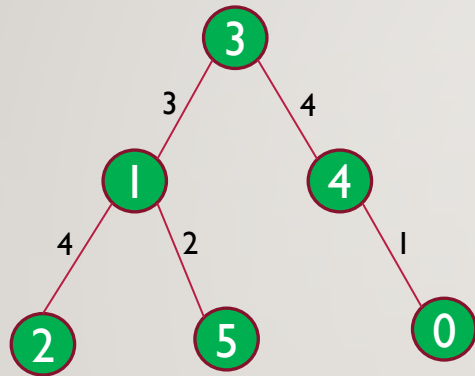
Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	t	t	t	t



DIJKSTRA'S ALGORITHM



Cost	0	1	2	3	4	5
	5	3	7	0	4	5

Parent	0	1	2	3	4	5
	4	3	1	3	3	1

Marked	0	1	2	3	4	5
	t	t	t	t	t	t

