

CSC212

# Data Structure



COMPUTER SCIENCE  
CITY COLLEGE OF NEW YORK

## Lecture 13

### Trees and Tree Traversals

Instructor: Prof. George Wolberg

Department of Computer Science

City College of New York

# Motivation

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- Linear structures
  - arrays
  - dynamic arrays
  - linked lists
- Nonlinear Structures
  - trees - Hierarchical Structures
  - Graphs
- Why???

# Application: Mailing Addresses

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George Wolberg, CS Dept, CCNY, New York, NY 10031, USA

6 billion = 6,000,000,000 people in the world

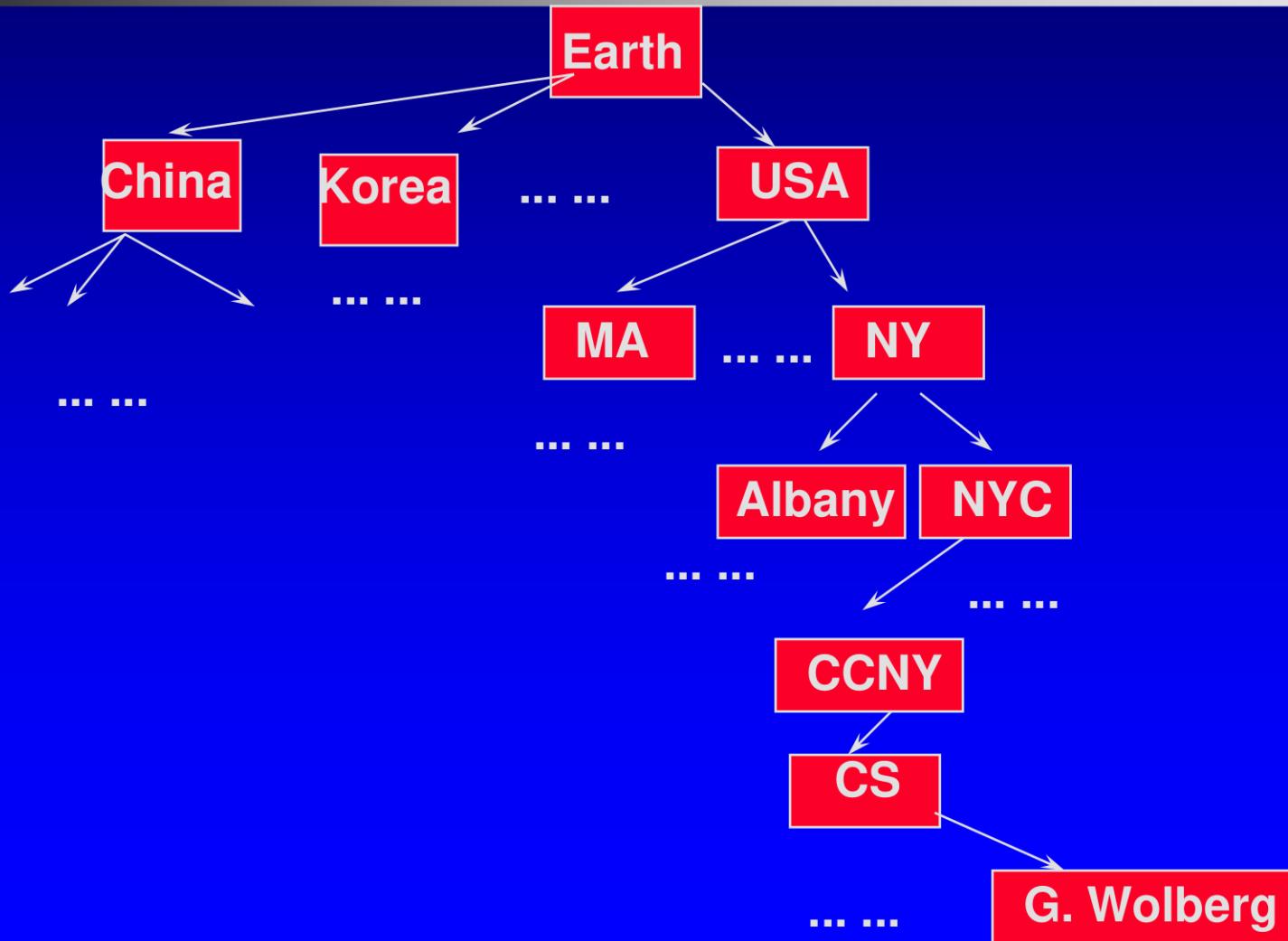
What kind of structure is the best for a postman to locate me?

Array ?

Linked list ?

Tree ?

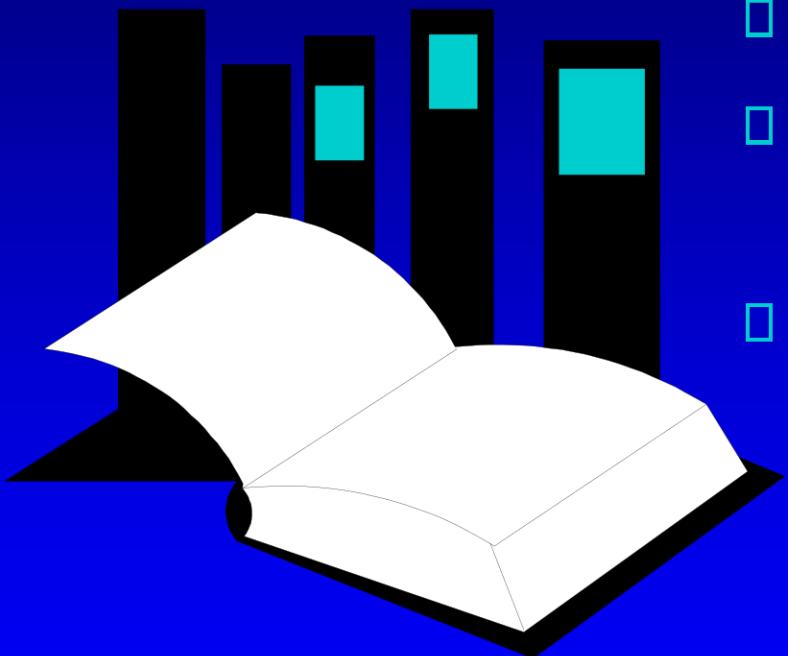
# A Tree for all the mailing addresses





# Trees and Binary Trees

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- Chapter 10 introduces **trees**.
- This presentation illustrates basic terminology for binary trees
- and focuses on
  - **Complete Binary Trees:** the simplest kind of trees
  - **Binary Tree Traversals:** any kind of binary trees

**Data Structures  
and Other Objects  
Using C++**

# Binary Trees

---

- A binary tree has **nodes**, similar to nodes in a linked list structure.
- **Data** of one sort or another may be stored at each node.
- But it is the **connections** between the nodes which characterize a binary tree.

# Binary Trees

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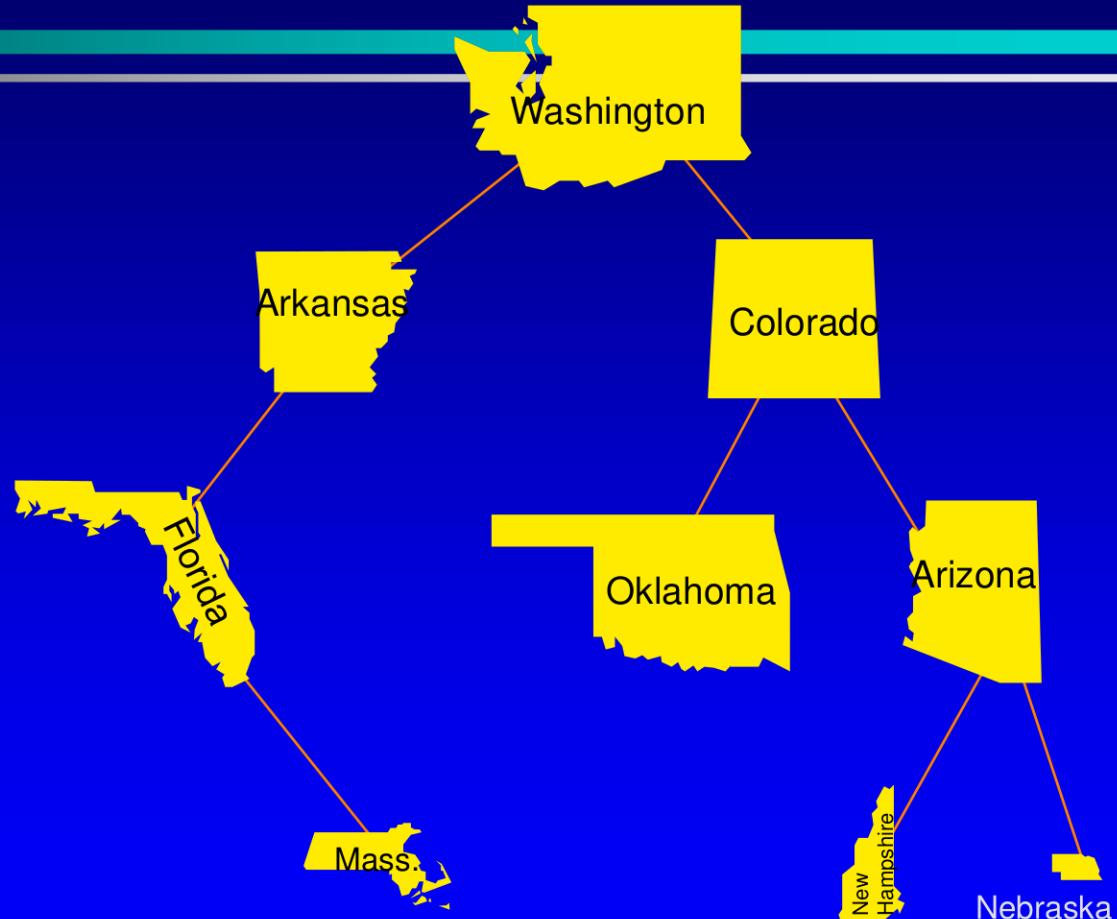
- A binary tree has **nodes**, similar to nodes in a linked list structure.
- **Data** of one sort or another may be stored at each node.
- But it is the **connections** between the nodes which characterize a binary tree.



An example can  
illustrate how the  
connections work

# A Binary Tree of States

In this example,  
the data  
contained at  
each node is  
one of the 50  
states.



# A Binary Tree of States

Each tree has a special node called its root, usually drawn at the top.



# A Binary Tree of States

Each tree has a special node called its root, usually drawn at the top.



# A Binary Tree of States

Each node is permitted to have two links to other nodes, called the left child and the right child.



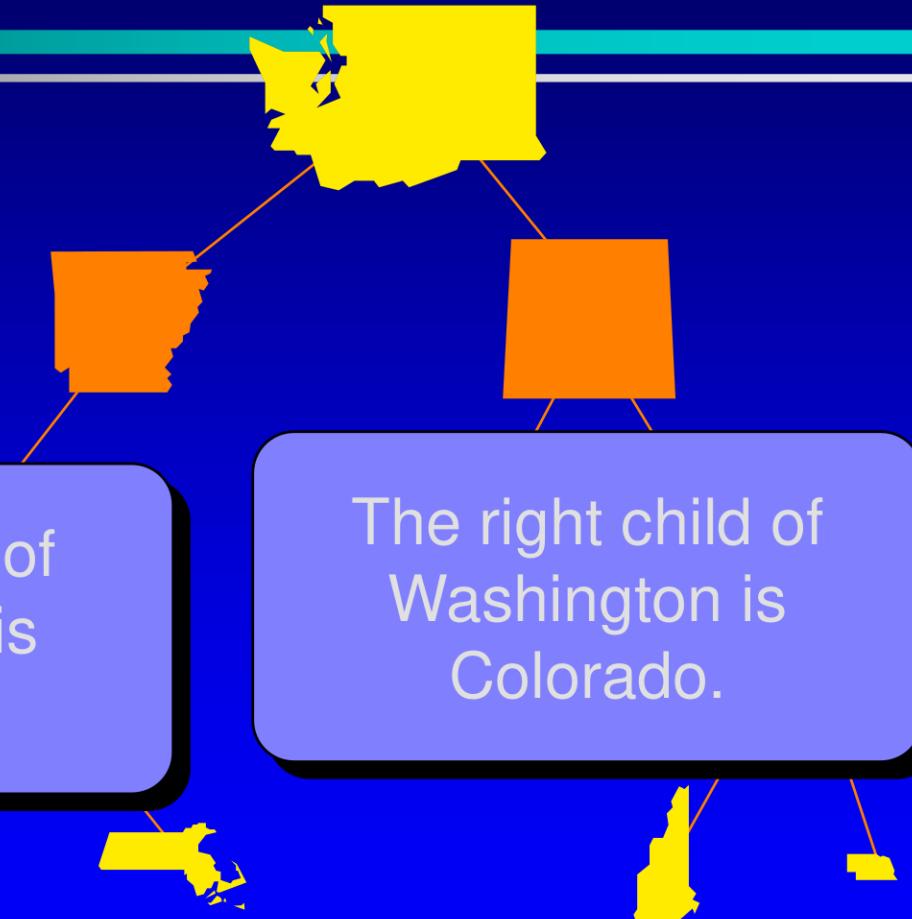
# A Binary Tree of States

Each node is permitted to have **two links** to other nodes, called the left child and the right child.



# A Binary Tree of States

Children are usually drawn below a node.



# A Binary Tree of States

Some nodes  
have only one  
child.

Arkansas has a  
left child, but no  
right child.



# A Quiz

Some nodes  
have only one  
child.

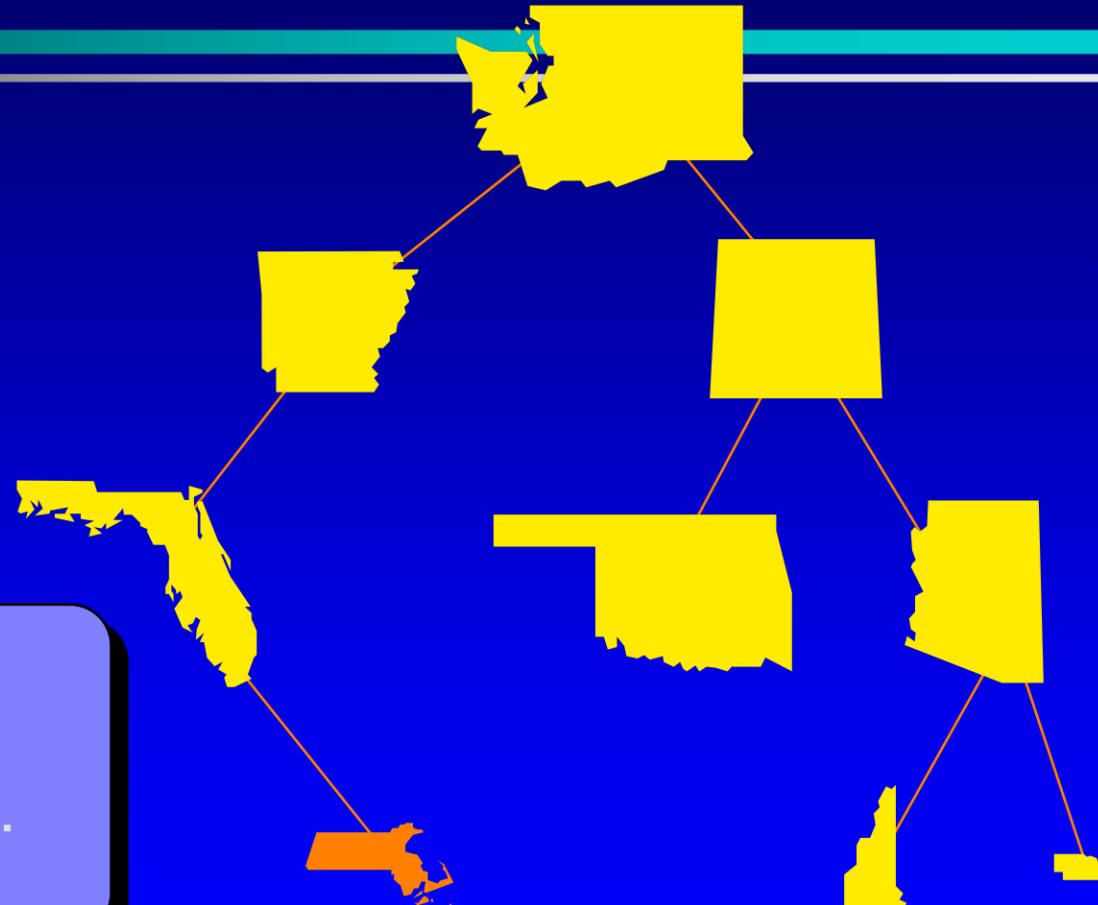
*Which node has  
only a right child?*



# A Quiz

Some nodes  
have only one  
child.

Florida has  
only a right child.



# A Binary Tree of States

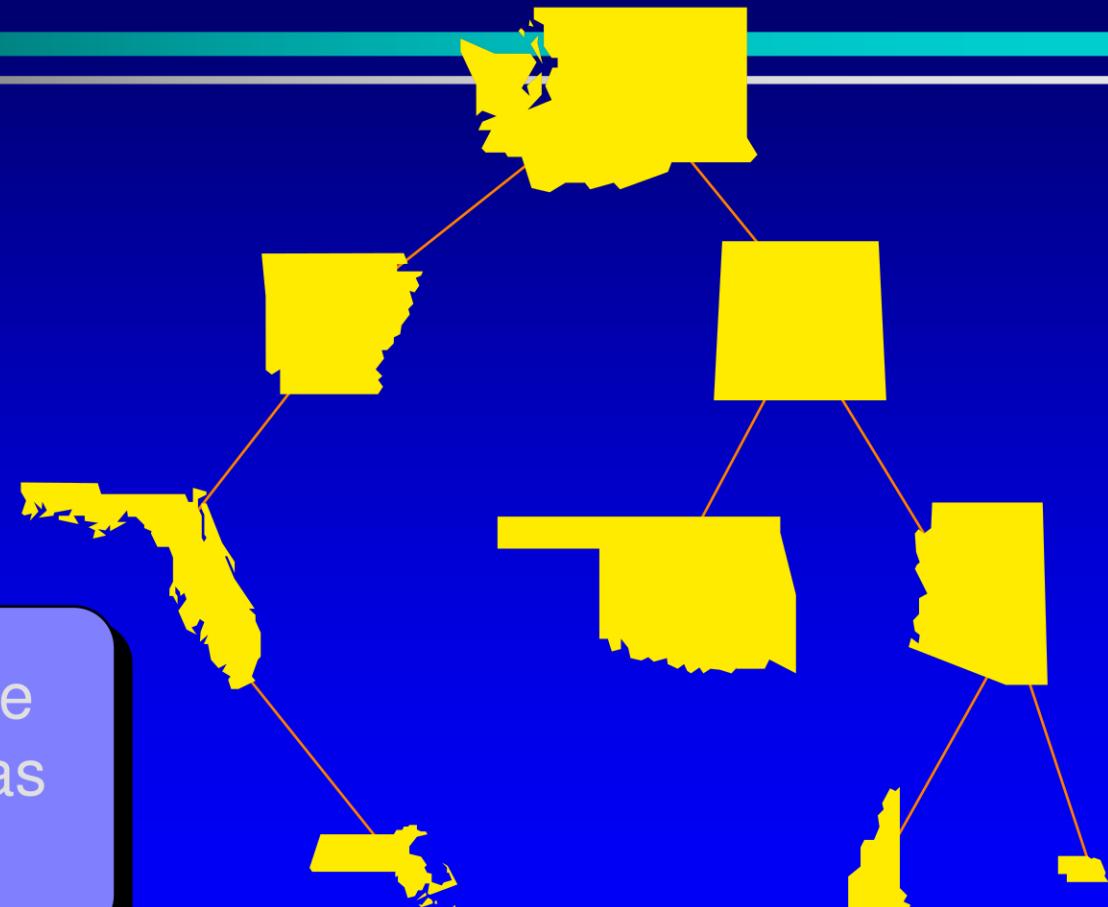
A node with no children is called a leaf.



# A Binary Tree of States

Each node is called the **parent** of its children.

Washington is the parent of Arkansas and Colorado.



# A Binary Tree of States

Two rules about parents:

- The root has no parent.
- Every other node has exactly one parent.



# A Binary Tree of States

Two nodes with  
the same parent  
are called  
siblings.

Arkansas  
and Colorado  
are siblings.



# Complete Binary Trees

---

A complete binary tree is a special kind of binary tree which will be useful to us.

# Complete Binary Trees

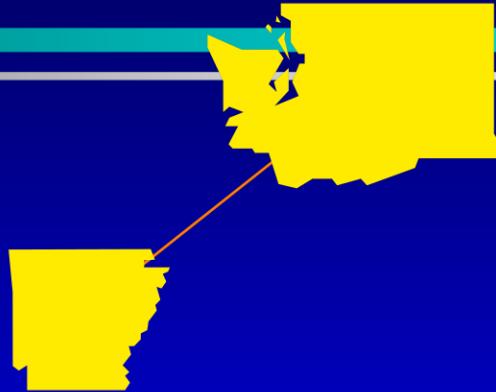


A complete binary tree is a special kind of binary tree which will be useful to us.

When a complete binary tree is built, its first node must be the root.

# Complete Binary Trees

The second node of a complete binary tree is always the left child of the root...



# Complete Binary Trees

The second node of a complete binary tree is always the left child of the root...

... and the third node is always the right child of the root.



# Complete Binary Trees

The next nodes must always fill the next level from left to right.



# Complete Binary Trees

The next nodes must always fill the next level from left to right.



# Complete Binary Trees

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# Complete Binary Trees

The next nodes must always fill the next level from left to right.



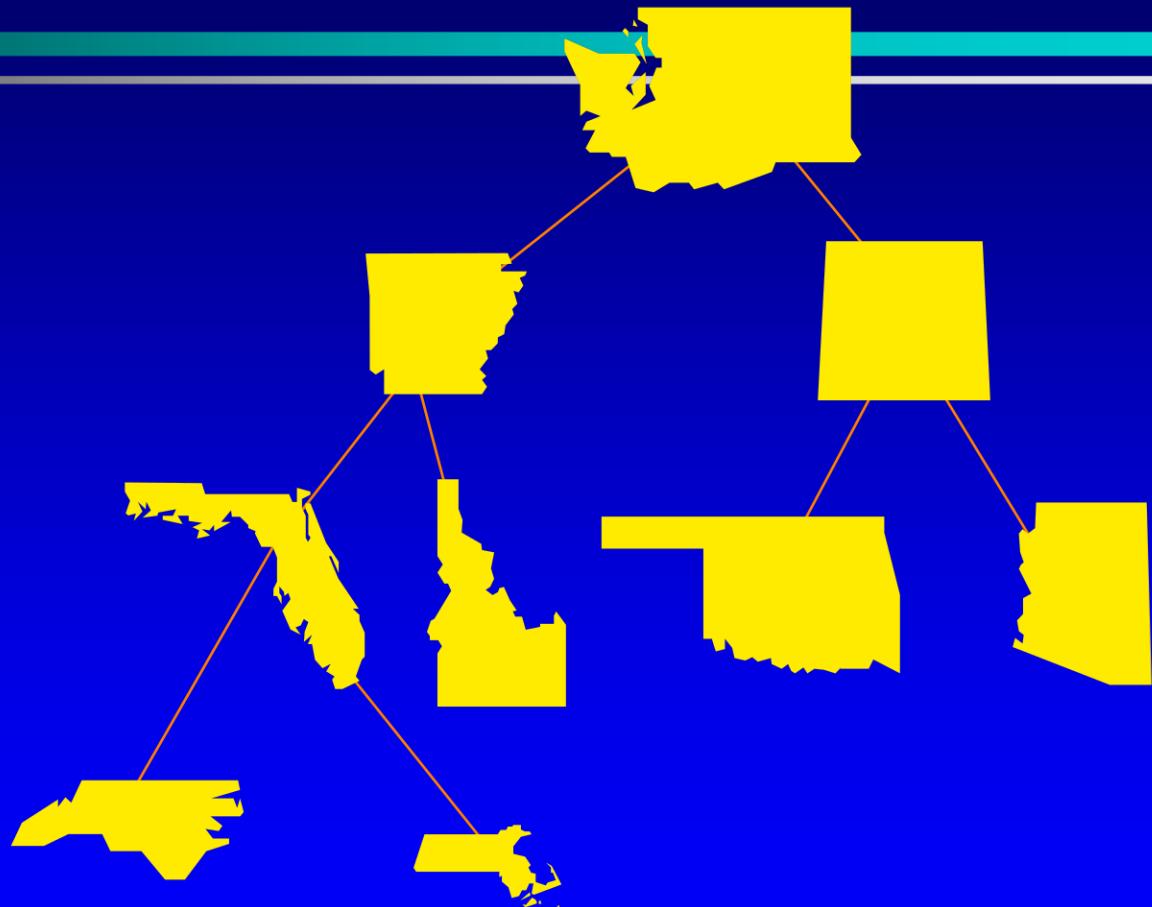
# Complete Binary Trees

The next nodes must always fill the next level from left to right.



# Complete Binary Trees

The next nodes must always fill the next level from left to right.



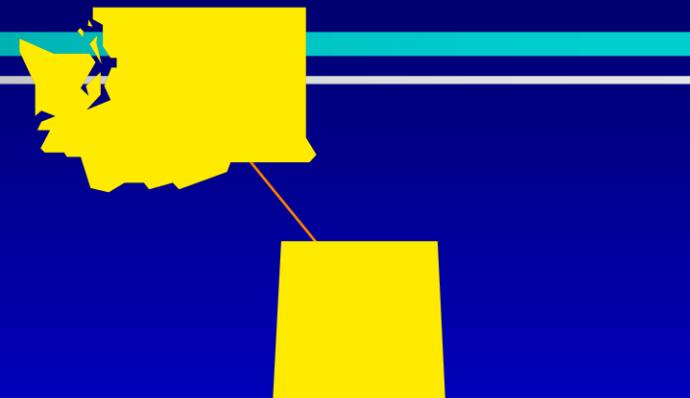
# Is This Complete?



# Is This Complete?



# Is This Complete?



# Is This Complete?



# Is This Complete?

---

Yes!

- It is called the empty tree, and it has no nodes, not even a root.

# Full Binary Trees

A full binary tree is a special kind of complete binary tree

FULL

When a full binary tree is built, its first node must be the root.

# Full Binary Trees

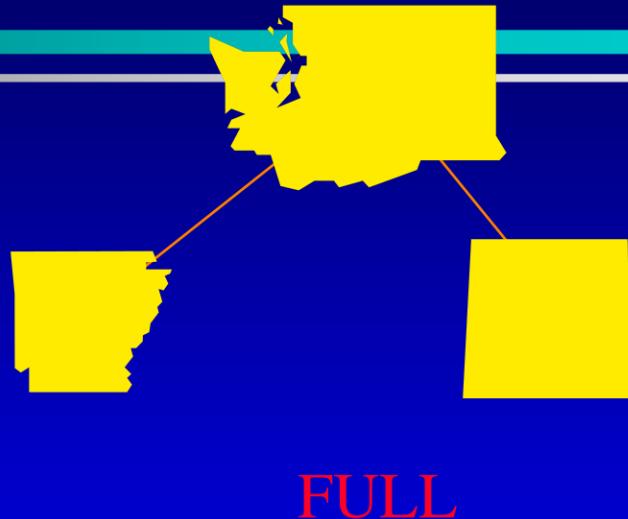
The second node of a full binary tree is always the left child of the root...



# Full Binary Trees

The second node of a full binary tree is always the left child of the root...

... and you MUST have the third node which always the right child of the root.



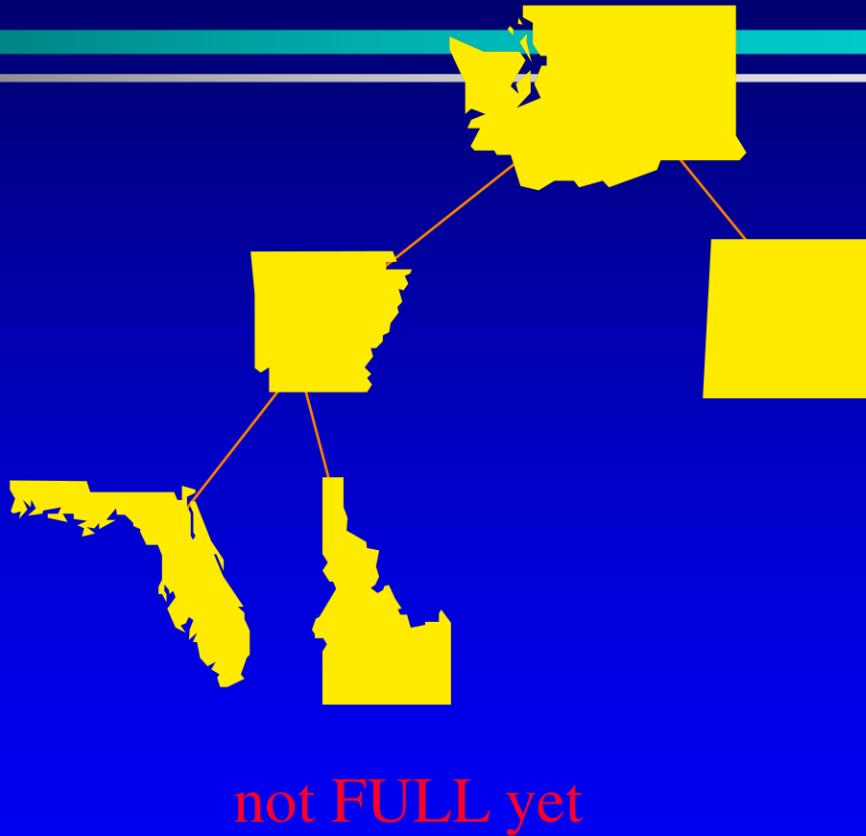
# Full Binary Trees

The next nodes must always fill the next level from left to right.



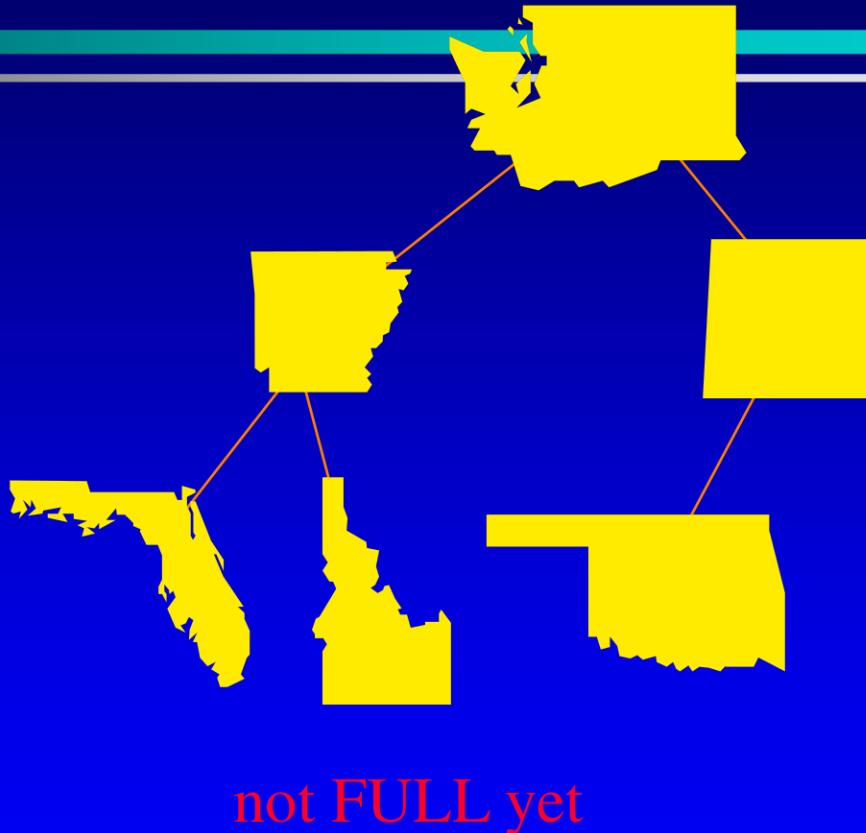
# Full Binary Trees

The next nodes must always fill the next level from left to right.



# Full Binary Trees

The next nodes must always fill the next level from left to right.



# Full Binary Trees

The next nodes must always fill the next level from **left to right**...until every leaf has the same depth (2)



FULL!

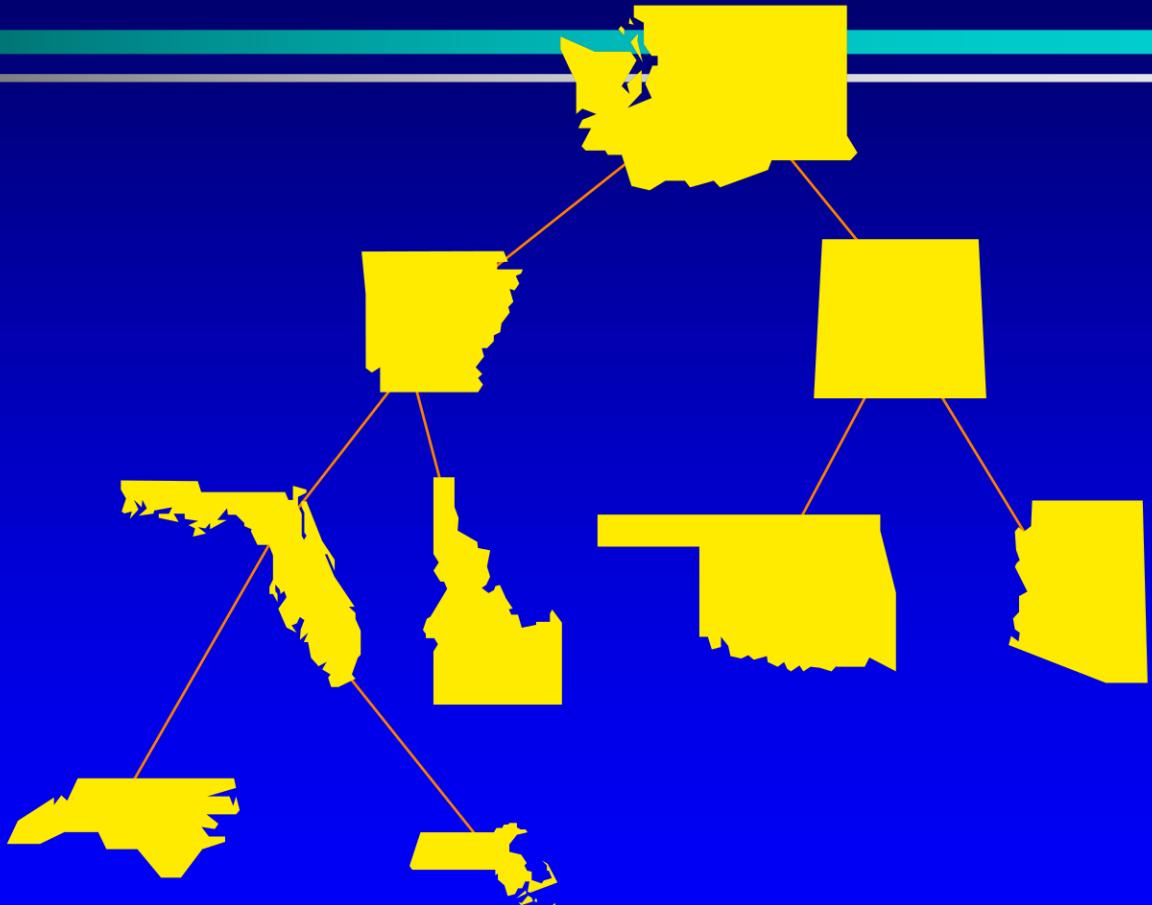
# Full Binary Trees

The next nodes must always fill the next level from left to right.



# Full Binary Trees

The next nodes must always fill the next level from left to right.



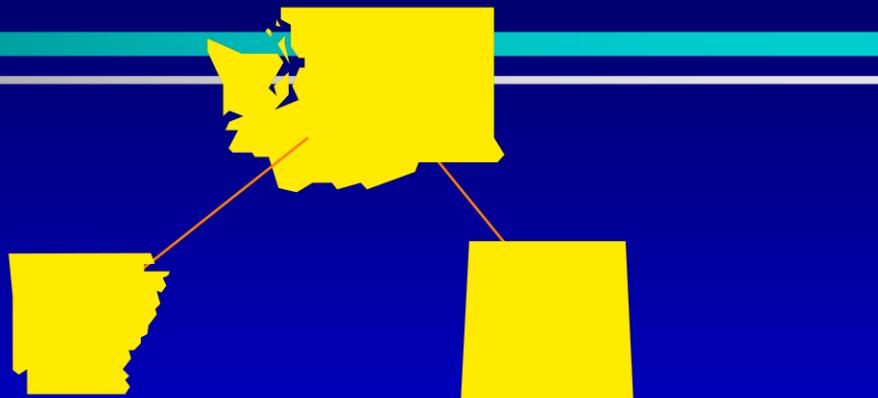
# Is This Full?



# Is This Full?



# Is This Full?



# Is This Full?



# Is This Full?

---

Yes!

- It is called the empty tree, and it has no nodes, not even a root.

# Implementing a Complete Binary Tree

- We will store the data from the nodes in a partially-filled array.

3

An integer to keep track of how many nodes are in the tree



An array of data

We don't care what's in this part of the array.

# Implementing a Complete Binary Tree Using an Array

- We will store the date from the nodes in a partially-filled array.

3

An integer to keep track of how many nodes are in the tree



Read Section 10.2 to see details of how the entries are stored.

An array of ~~dates~~ integers will store the data. Care what's in this part of the array.

# Implementing a Complete Binary Tree Using an Array

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- Root is at component [0]
- Parent of node in [i] is at  $[(i-1)/2]$
- Children (if exist) of node [i] is at  $[2i+1]$  and  $[2i+2]$
- Total node number
  - $2^0+2^1+2^2+\dots+2^{d-1}+r, r \leq 2^d$ , d is the depth



# Binary Tree Summary

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- Binary trees contain nodes.
- Each node may have a left child and a right child.
- If you start from any node and move upward, you will eventually reach the root.
- Every node except the root has one parent. The root has no parent.
- Complete binary trees require the nodes to fill in each level from left-to-right before starting the next level.

# Binary Tree Basics

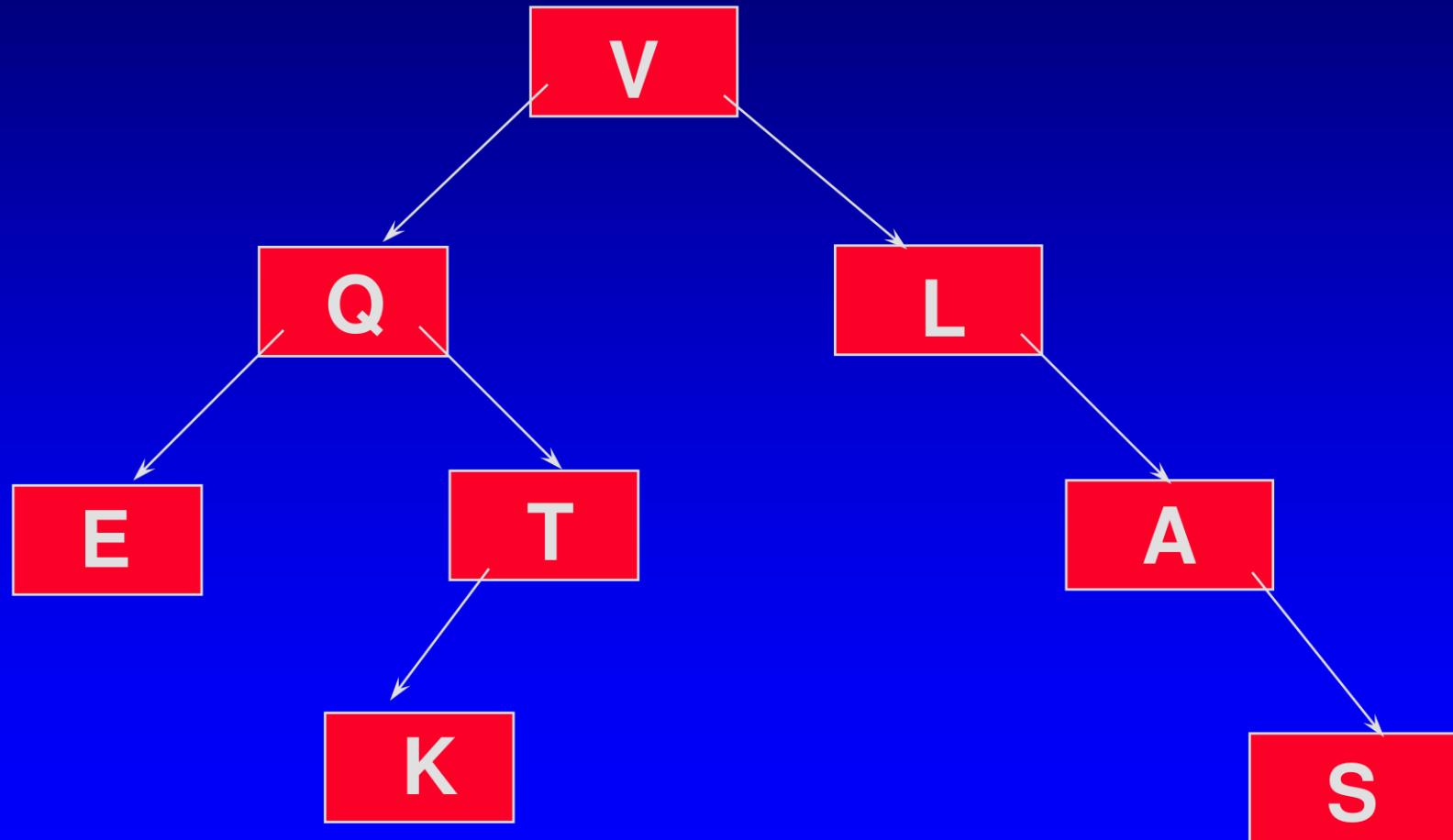
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**A binary tree is a structure in which:**

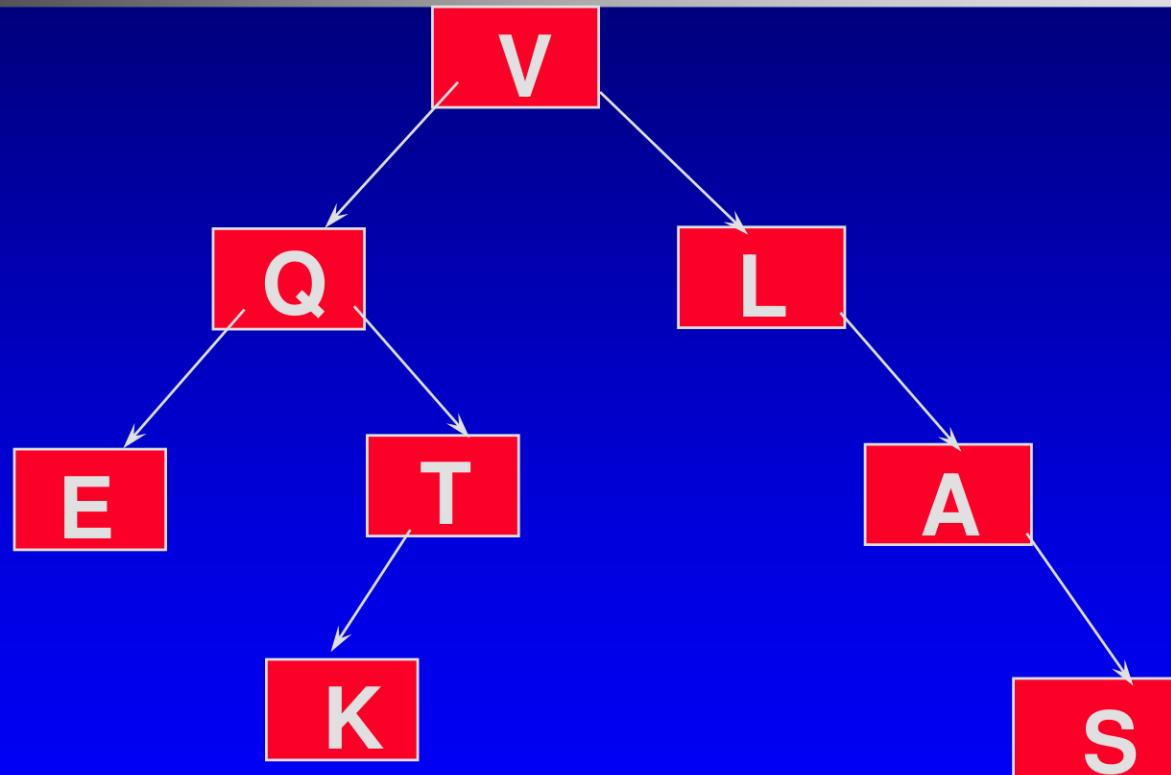
**Each node can have at most two children, and in which a unique path exists from the root to every other node.**

**The two children of a node are called the **left child** and the **right child**, if they exist.**

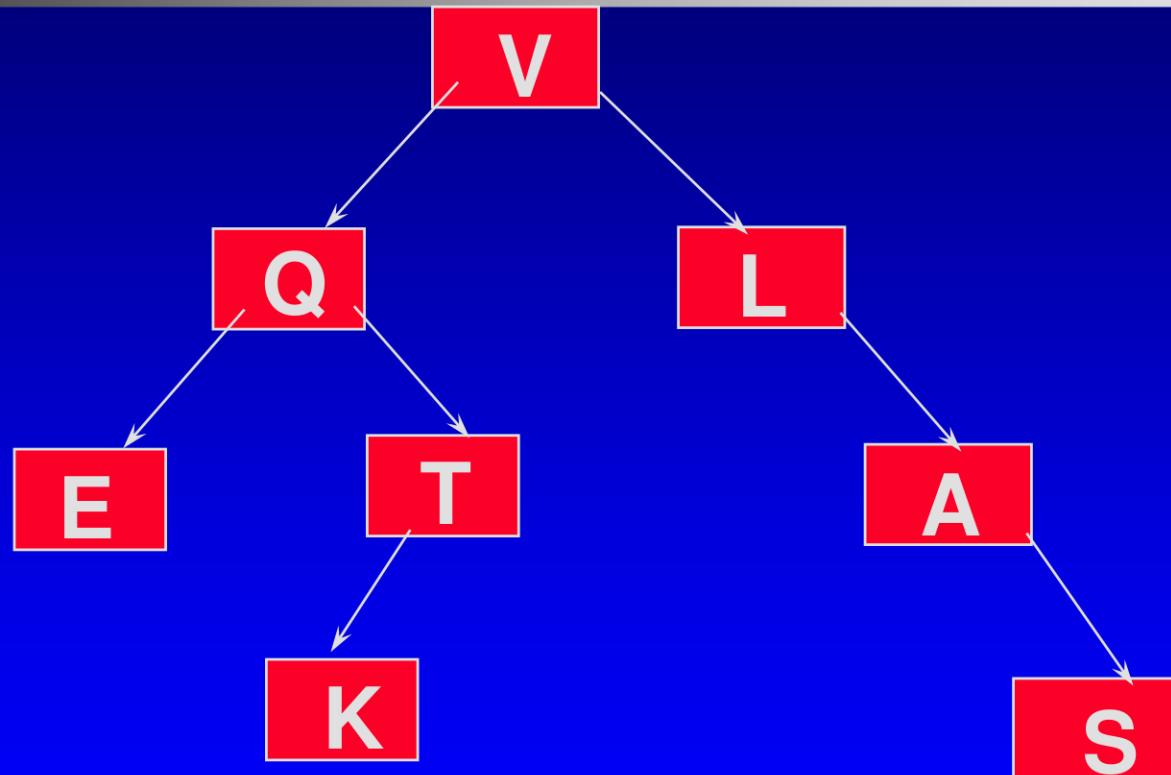
# A Binary Tree Exercise



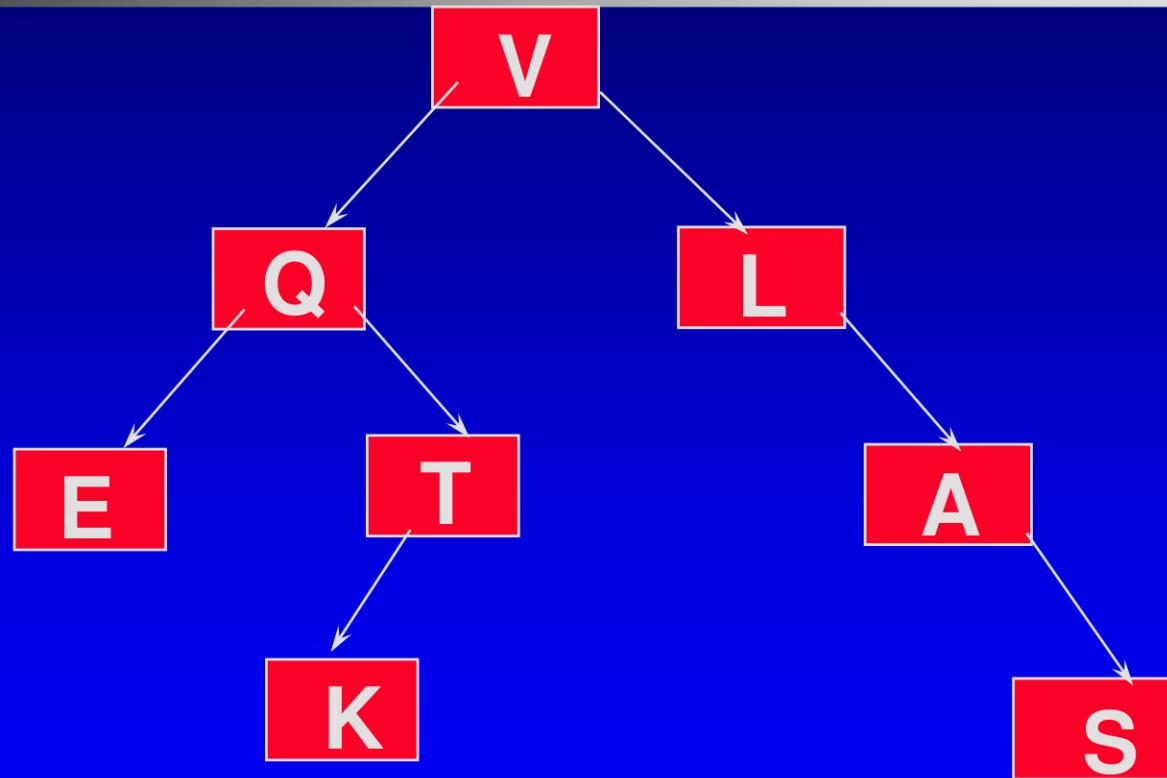
# How many leaf nodes?



# How many descendants of Q?

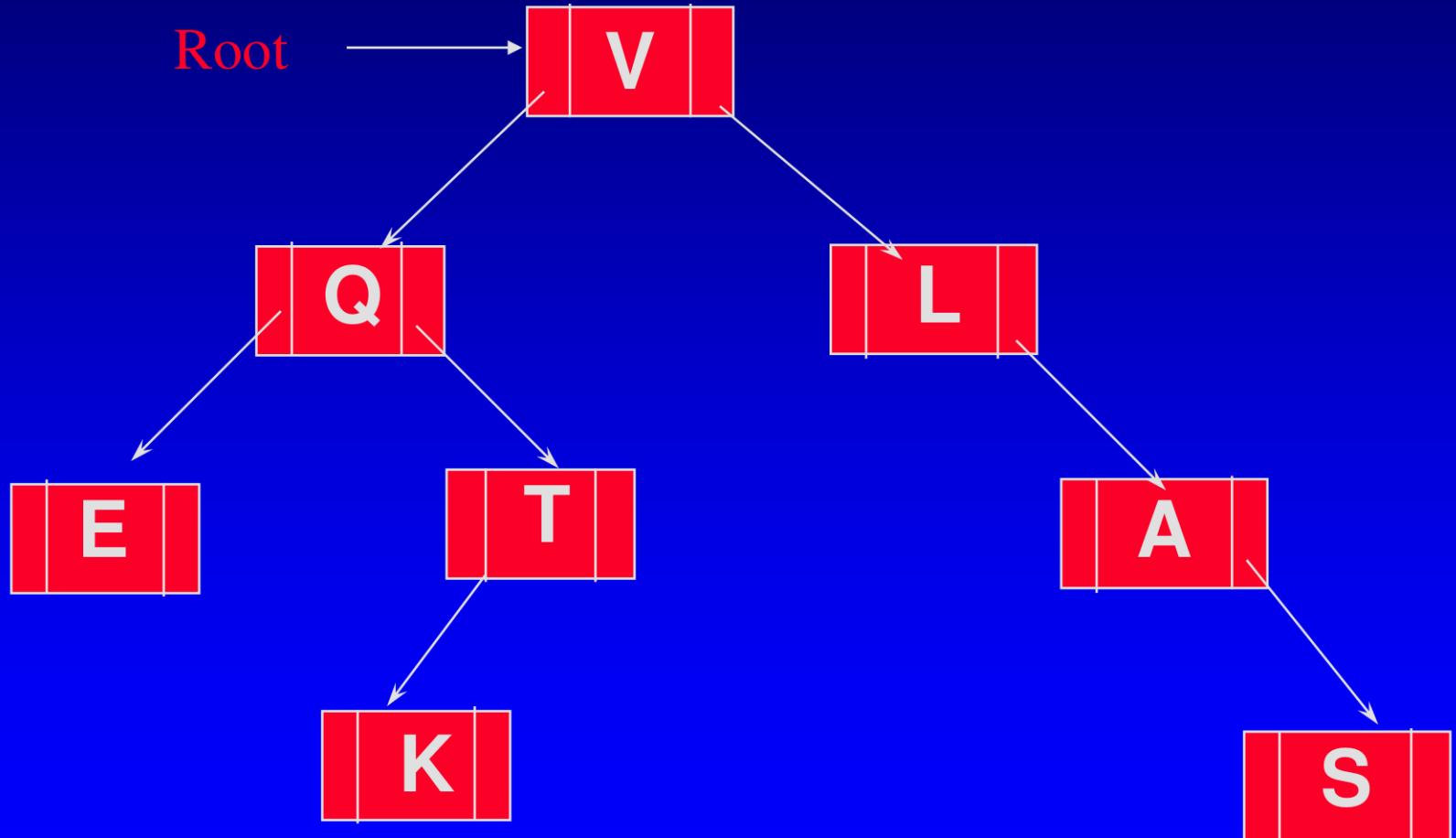


# How many ancestors of K?



Question: How to implement a general binary tree ?

# Implementing a Binary Tree with a Class for Nodes



# Binary Tree Nodes

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- Each node of a binary tree is stored in an object of a new `binary_tree_node` class that we are going to define
- Each node contains data as well as pointers to its children (nodes)
- An entire tree is represented as a pointer to the root node

# binary\_tree\_node Class

bintree

## □ variables

## □

```
template <class Item>
class binary_tree_node
{
public:
    .....
private:
    Item data_field;
    binary_tree_node *left_field;
    binary_tree_node *right_field;
};
```

## functions

//retrievals

data

left

right

//set

set\_data

set\_left

set\_right

//boolean

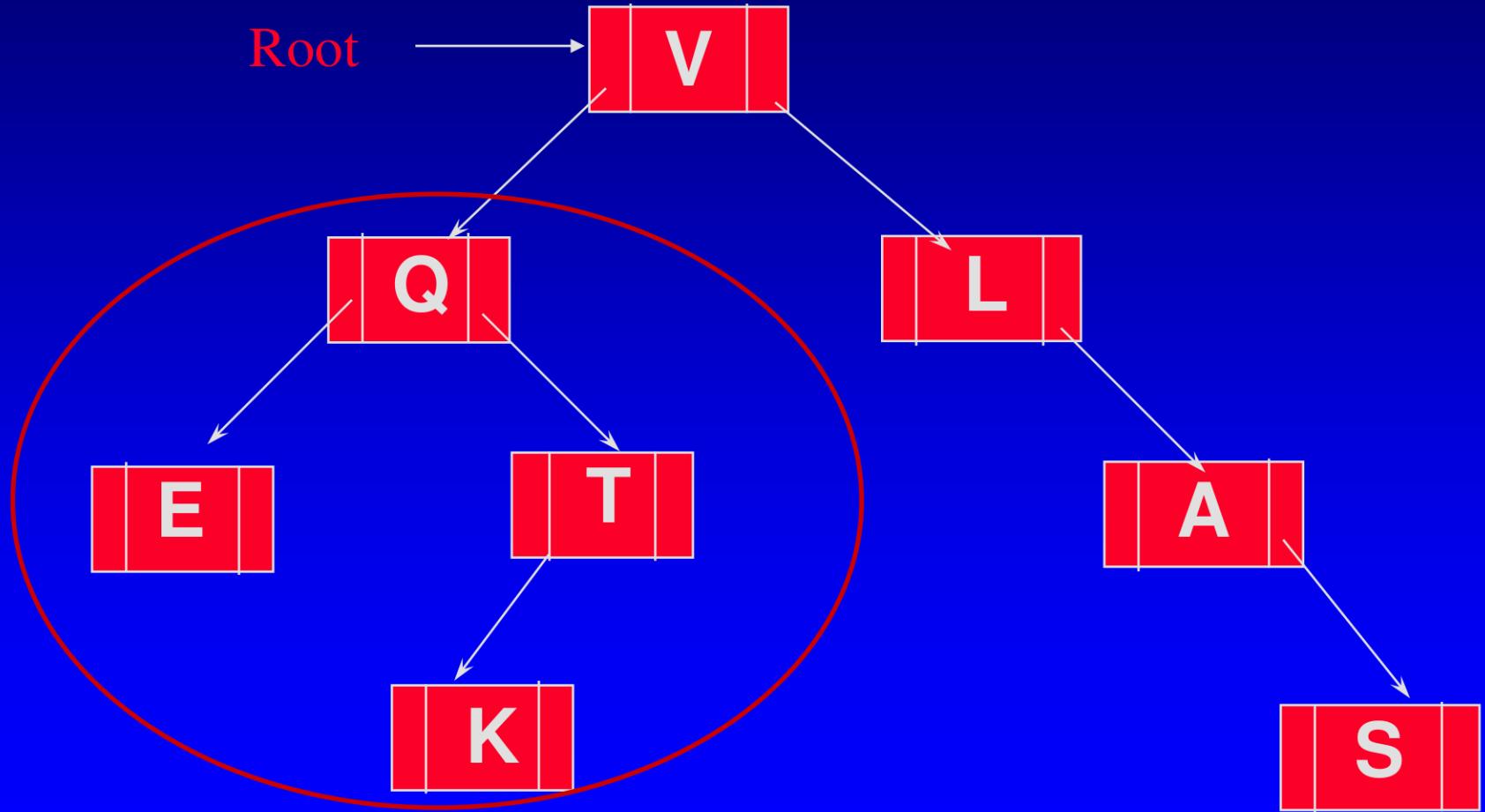
is\_leaf

# Creating and Manipulating Trees

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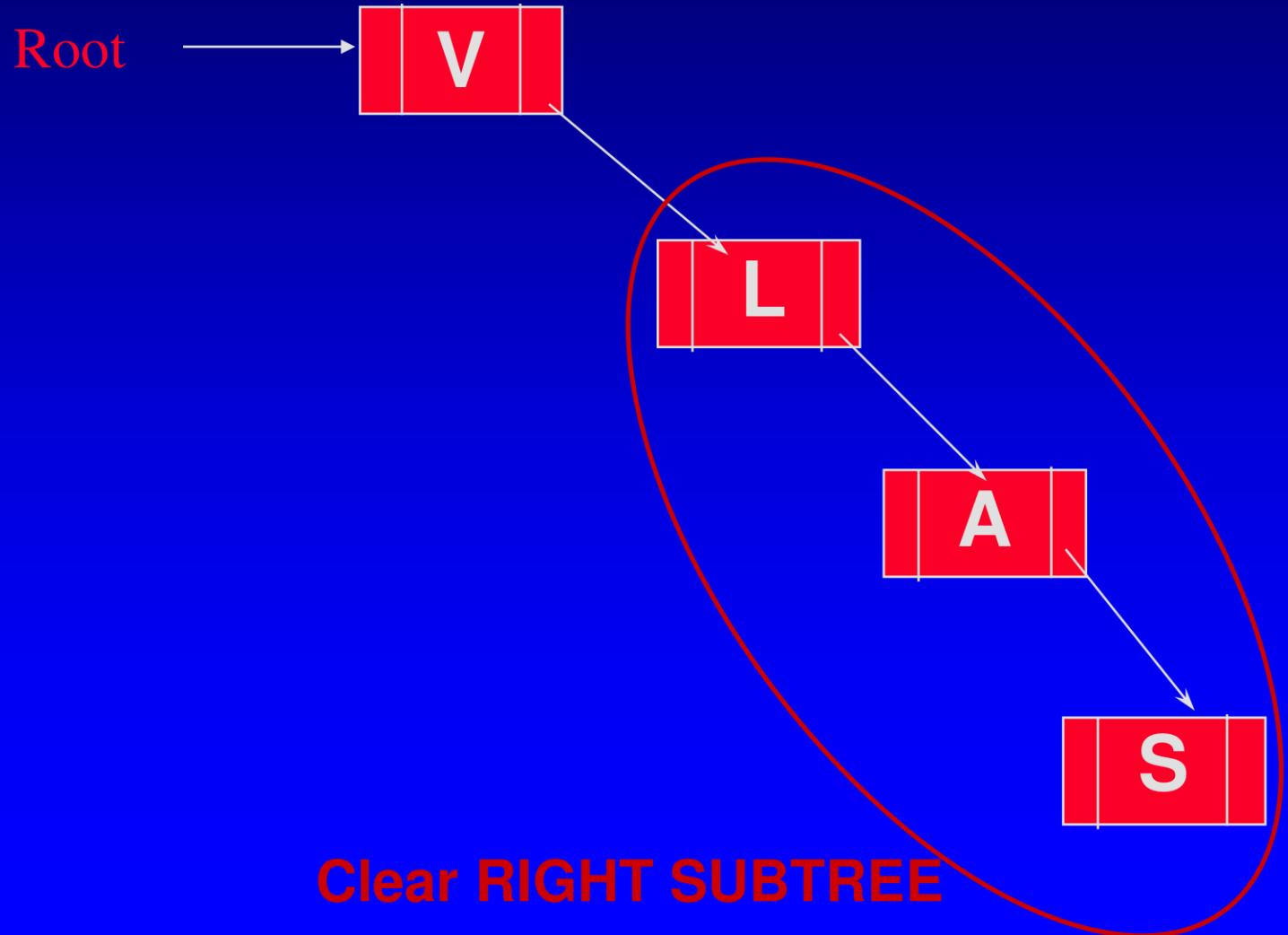
- Consider only two functions
  - Clearing a tree
    - Return nodes of a tree to the heap
  - Copying a tree
- The Implementation is easier than it seems
  - if we use recursive thinking

# Clearing a Tree

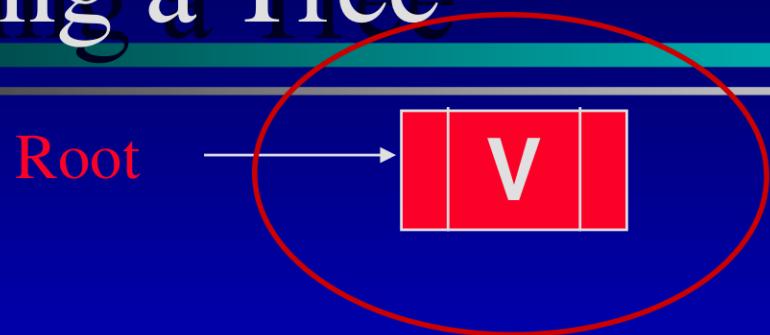


Clear LEFT SUBTREE

# Clearing a Tree



# Clearing a Tree



**Return root node to the heap**

# Clearing a Tree

---



**Set the root pointer to NULL**

# Clear a Tree

bintree

- key: recursive thinking

```
template <class Item>
void tree_clear(binary_tree_node<Item*>*& root_ptr)
// Library facilities used: cstdlib
{
    if (root_ptr != NULL)
    {
        tree_clear( root_ptr->left( ) ); // clear left sub_tree
        tree_clear( root_ptr->right( ) ); // clear right sub_tree
        delete root_ptr; // return root node to the heap
        root_ptr = NULL; // set root pointer to the null
    }
}
```

# Copy a Tree

bintree

- Can you implement the copy? (p 467)

```
template <class Item>
binary_tree_node<Item>* tree_copy(const binary_tree_node<Item>* root_ptr)
// Library facilities used: cstdlib
{
    binary_tree_node<Item> *l_ptr;
    binary_tree_node<Item> *r_ptr;

    if (root_ptr == NULL)
        return NULL;
    else
    {
        l_ptr = tree_copy( root_ptr->left( ) ); // copy the left sub_tree
        r_ptr = tree_copy( root_ptr->right( ) ); // copy the right sub_tree
        return
            new binary_tree_node<Item>( root_ptr->data( ), l_ptr, r_ptr );
    } // copy the root node and set the the root pointer
}
```

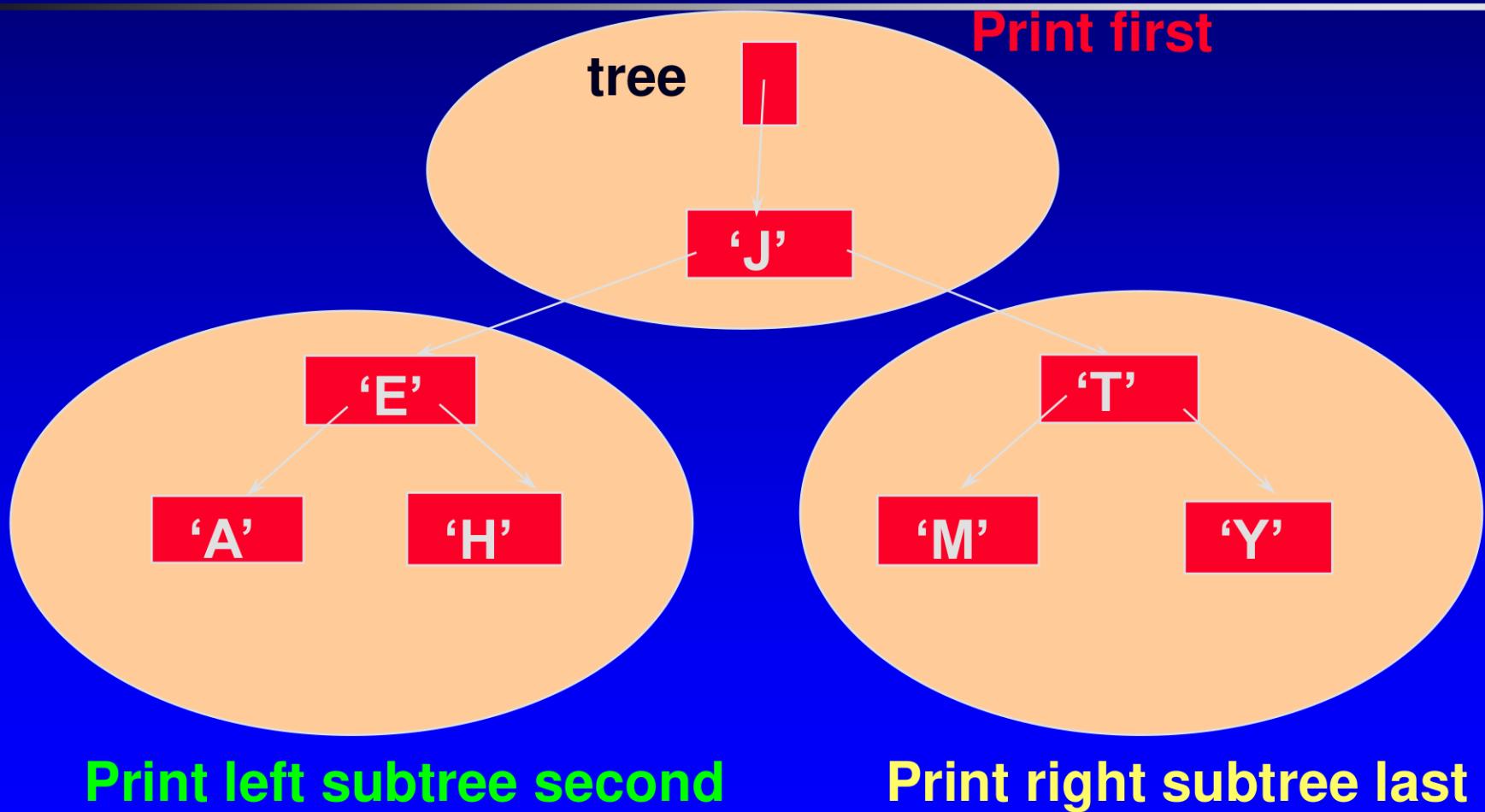
# Binary Tree Traversals

bintree

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- pre-order traversal
  - root (left sub\_tree) (right sub\_tree)
- in-order traversal
  - (left sub\_tree) root (right sub\_tree)
- post-order traversal
  - (left sub\_tree) (right sub\_tree) root
- backward in-order traversal
  - (right sub\_tree) root (left sub\_tree)

# Preorder Traversal: J E A H T M Y

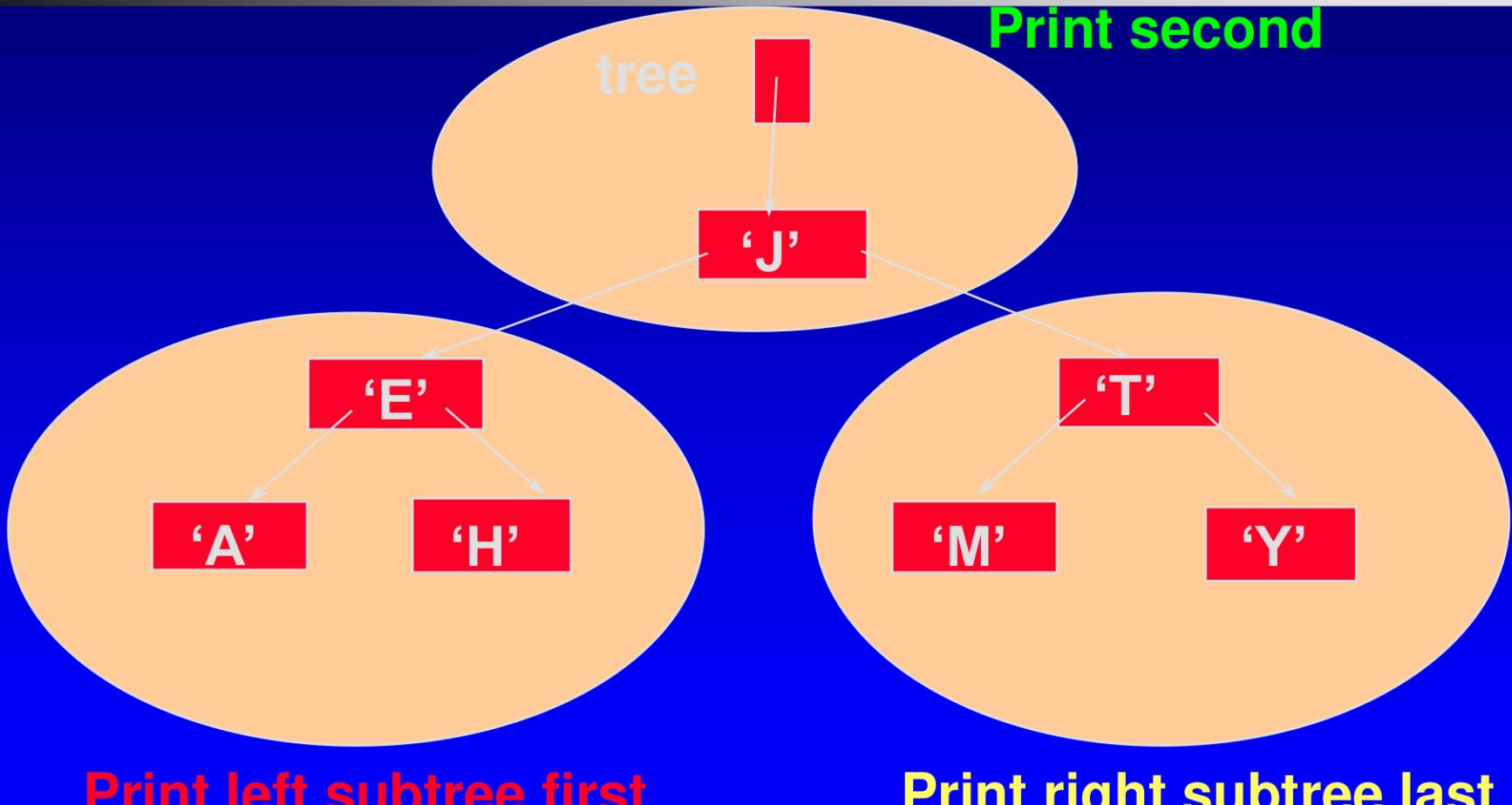


# Preorder Traversal

- Example: print the contents of each node

```
template <class Item>
void preorder_print(const binary_tree_node<Item>* node_ptr)
// Library facilities used: cstdlib, iostream
{
    if (node_ptr != NULL)
    {
        std::cout << node_ptr->data( ) << std::endl;
        preorder_print(node_ptr->left( ));
        preorder_print(node_ptr->right( ));
    }
}
```

# Inorder Traversal: A E H J M T Y

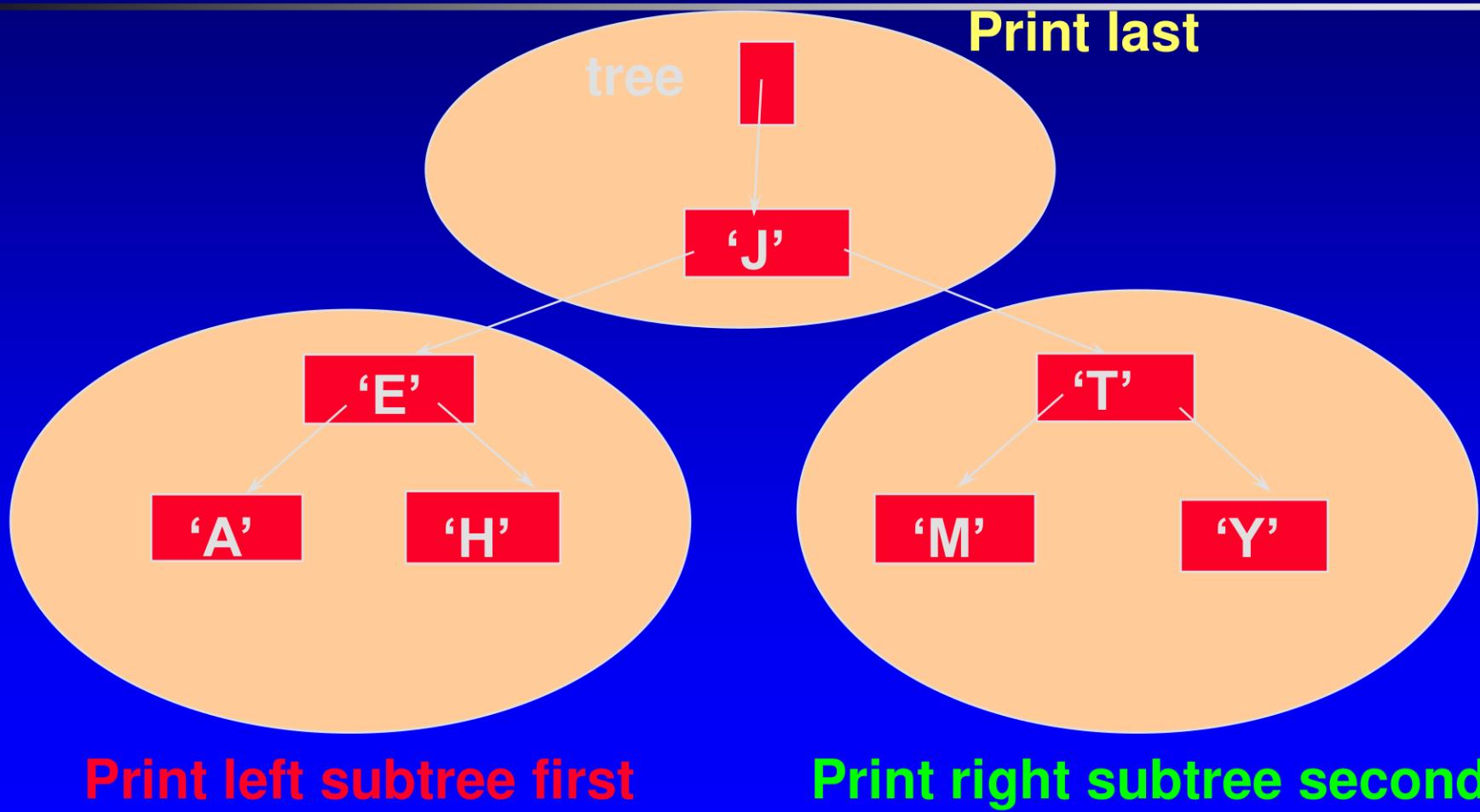


# Inorder Traversal

- Example: print the contents of each node

```
template <class Item>
void inorder_print(const binary_tree_node<Item>* node_ptr)
// Library facilities used: cstdlib, iostream
{
    if (node_ptr != NULL)
    {
        inorder_print(node_ptr->left( ));
        std::cout << node_ptr->data( ) << std::endl;
        inorder_print(node_ptr->right( ));
    }
}
```

# Postorder Traversal: A H E M Y T J



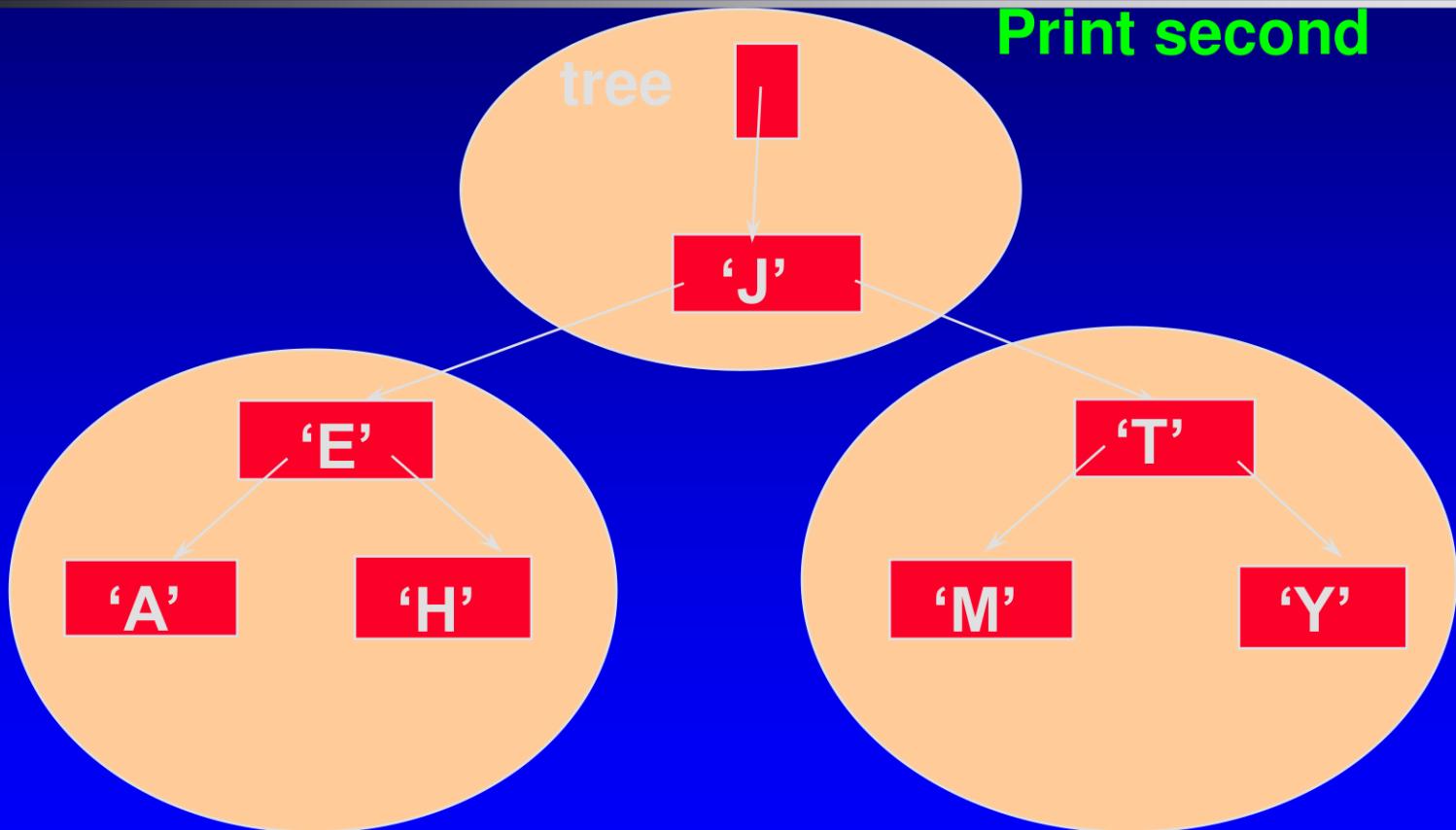
# Postorder Traversal

- Example: print the contents of each node

```
template <class Item>
void postorder_print(const binary_tree_node<Item>* node_ptr)
// Library facilities used: cstdlib, iostream
{
    if (node_ptr != NULL)
    {
        postorder_print(node_ptr->left( ));
        postorder_print(node_ptr->right( ));
        std::cout << node_ptr->data( ) << std::endl;
    }
}
```

# Backward Inorder Traversal:

Y T M I J H E A



Print second

Print left subtree last

Print right subtree first

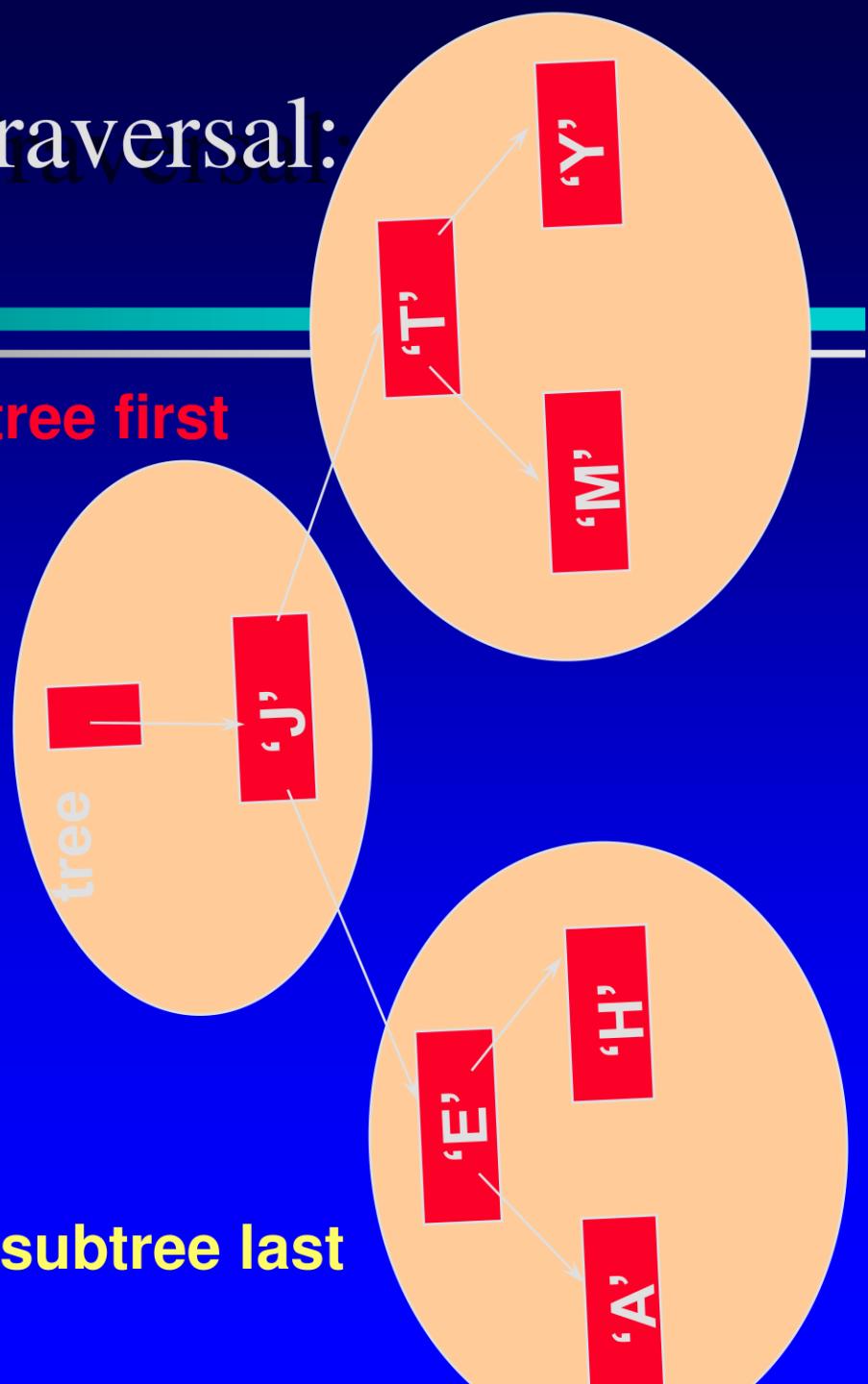
# Backward Inorder Traversal:

Y T M J H E A

**Print right subtree first**

**Print second**

**Print left subtree last**



# A Useful Backward Inorder Traversal

bintree

- Indent each number according its depth

```
template <class Item, class SizeType>
void print(binary_tree_node<Item>* node_ptr, SizeType depth)
// Library facilities used: iomanip, iostream, stdlib
{
    if (node_ptr != NULL)
    {
        print(node_ptr->right( ), depth+1);
        std::cout << std::setw(4*depth) << ""; // Indent 4*depth spaces.
        std::cout << node_ptr->data( ) << std::endl;
        print(node_ptr->left( ), depth+1);
    }
}
```

# A Challenging Question:

---

- For the traversals we have seen, the “processing” was simply printing the values of the node
- But we’d like to do any kind of processing
  - We can replace “cout” with some other form of “processing”
- But how about 1000 kinds?
  - Can template be helpful?
- Solution::::::> (pages 501 – 507)

# A parameter can be a function

---

- write one function capable of doing anything
- A parameter to a function may be a function. Such a parameter is declared by
  - the name of the function's return type (or void),
  - then the name of the parameter (i.e. the function),
  - and finally a pair of parentheses ().
  - Inside () is a list of parameter types of that parameter function
- Example
  - `int sum ( void f (int&, double), int i,...);`

# Preorder Traversal – print only

- Example: print the contents of each node

```
template <class Item>
void preorder_print(const binary_tree_node<Item>* node_ptr)
// Library facilities used: cstdlib, iostream
{
    if (node_ptr != NULL)
    {
        std::cout << node_ptr->data( ) << std::endl;
        preorder_print(node_ptr->left( ));
        preorder_print(node_ptr->right( ));
    }
}
```

# Preorder Traversal – general form

- A template function for tree traversals

```
template <class Item>
void preorder(void f(Item&), binary_tree_node<Item>* node_ptr)
// Library facilities used: cstdlib
{
    if (node_ptr != NULL)
    {
        f( node_ptr->data( ) ); // node_ptr->data() return reference !
        preorder(f, node_ptr->left( ));
        preorder(f, node_ptr->right( ));
    }
}
```

# Preorder Traversal – how to use

- Define a real function before calling

```
void printout(int & it)
    // Library facilities used: iostream
{
    std::cout << it << std::endl;
}
```

Can you print out all the node of a tree pointed by root ?

```
binary_tree_node<int> *root;
.....
preorder(printout, root);
```

Yes!!!

# Preorder Traversal – another functions

- Can define other functions...

```
void assign_default(int& it)
    // Library facilities used: iostream
{
    it = 0;
} // unfortunately template does not work here for function parameters
```

You can assign a default value to all the node of a tree pointed by root:

```
binary_tree_node<int> *root;
....
preorder(assign_default, root);
```

# Preorder Traversal – how to use

- Can the function-arguments be template?

```
template <class Item>
void printout(Item& it)
// Library facilities used: iostream
{
    std::cout << it << std::endl;
}
```

Can you print out all the nodes of a tree pointed by root ?

```
binary_tree_node<string> *root;
```

....

```
preorder(print_out, root);
```

X ! print\_out should have real types

# Preorder Traversal – how to use

- The function-arguments may be template if...

```
template <class Item>
void printout(Item& it)
// Library facilities used: iostream
{
    std::cout << it << std::endl;
}
```

Can you print out all the node of a tree pointed by root ?

```
binary_tree_node<string> *root;
....
preorder(print_out<string>, root);
```

But you may do the instantiation like this

# Preorder Traversal

## = a more general form

bintree

- An extremely general implementation (p 505)

```
template <class Process, class BTNode>
void preorder(Process f, BTNode* node_ptr)
// Note: BTNode may be a binary_tree_node or a const binary tree node.
// Process is the type of a function f that may be called with a single
// Item argument (using the Item type from the node),
// as determined by the actual f in the following.
// Library facilities used: cstdlib
{
    if (node_ptr != NULL)
    {
        f( node_ptr->data( ) );
        preorder(f, node_ptr->left( ));
        preorder(f, node_ptr->right( ));
    }
}
```

# Functions as Parameters

---

- We can define a template function  $X$  with functions as parameters – which are called *function parameters*
- A function parameter can be simply written as  $\text{Process } f()$  (where Process is a template), and the forms and number of parameters for  $f$  are determined by the actual call of  $f$  inside the template function  $X$
- The real function argument for  $f$  when calling the the template function  $X$  cannot be a template function, it must be instantiated in advance or right in the function call

# Summary

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- Tree, Binary Tree, Complete Binary Tree
  - child, parent, sibling, root, leaf, ancestor,...
- Array Representation for Complete Binary Tree
  - Difficult if not complete binary tree
- A Class of `binary_tree_node`
  - each node with two link fields
- Tree Traversals
  - recursive thinking makes things much easier
- A general Tree Traversal
  - A Function as a parameter of another function

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