

Verification and Explanation of Unfairness in Machine Learning

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Joint work with Bishwamittra Ghosh and Kuldeep S. Meel^b

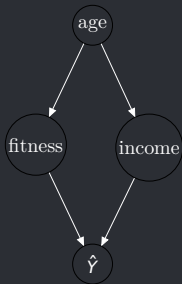
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(Un)Fairness in Machine Learning

Prediction of eligibility of health insurance

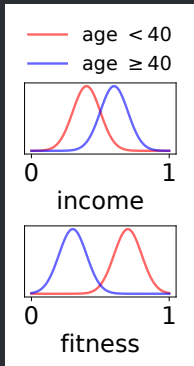
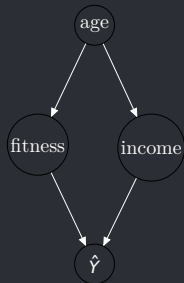
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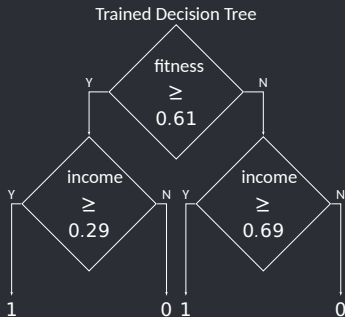
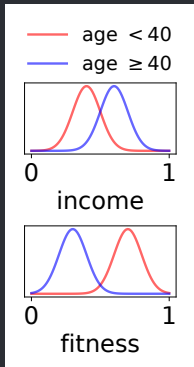
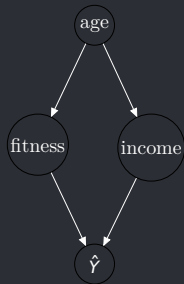
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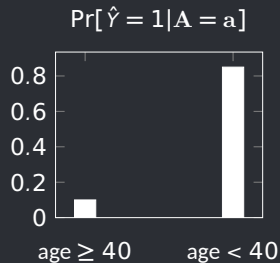
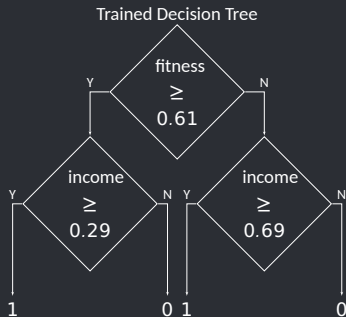
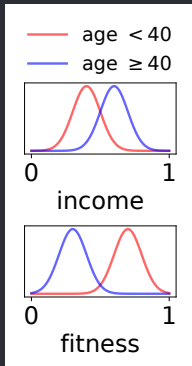
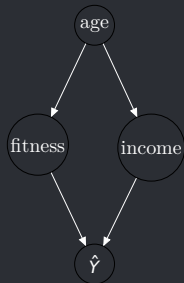
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Motivation

- Machine learning classifiers may become unfair to certain demographic groups
- Multiple fairness definitions and algorithms have been proposed to improve fairness
- What are still missing is scalable algorithms for verification and explanation of fairness
- Today, we focus on
 - **Fairness Verification:** A rigorous estimate of fairness of a classifier
 - **Fairness Explanation:** Identifying the source of unfairness of a classifier through the lens of input features

Outline

1. Motivation
2. Fairness Verification of Boolean Formulas
3. Fairness Verification of Linear Classifiers with Correlated Features
4. Fairness Explanation: A Model-agnostic Approach

Justicia: A Stochastic SAT Approach to Formally Verify Fairness [1]

Given

- a binary classifier $\mathcal{A} : (\mathbf{X}, \mathbf{A}) \rightarrow \hat{Y} \in \{0, 1\}$ and
- a probability distribution $(\mathbf{X}, \mathbf{A}, Y) \sim \mathcal{D}$,

verify whether \mathcal{A} achieves fairness w.r.t. \mathcal{D}

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$\Pr[\hat{Y} = 1 | \mathbf{A} = a]$ is called
the conditional PPV (Positive
Predictive Value)

Statistical parity: \mathcal{A} satisfies ϵ -statistical parity if for $\epsilon \in [0, 1]$,

$$\max_a \Pr[\hat{Y} = 1 | \mathbf{A} = a] - \min_a \Pr[\hat{Y} = 1 | \mathbf{A} = a] \leq \epsilon$$

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Our Approach: Compute the maximum and minimum of $\Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$
by a reduction to stochastic SAT

Satisfiability (SAT) problem

A Recap

Given a Boolean formula ϕ in CNF (Conjunctive Normal Form) defined over Boolean variables \mathbf{X} , the SAT problem finds a satisfying assignment of \mathbf{X} that evaluates ϕ to true

$$\phi = (X_1 \vee \neg X_2) \wedge (\neg X_1 \vee X_2 \vee X_3) \wedge \neg X_1$$

- SAT solution: $X_1 = \text{false}$, $X_2 = \text{false}$, $X_3 = \text{true}$

Stochastic SAT (SSAT)

A Brief Introduction

An SSAT formula Φ has a prefix and a CNF formula ϕ

$$\Phi = \underbrace{q_1 X_1, \dots, q_n X_n}_{\text{prefix}}, \phi$$

- q_i is an universal (\forall), existential (\exists), or randomized \forall^{p_i} quantifier with $p_i = \Pr[X_i = \text{true}]$
- SSAT computes the probability of satisfaction $\Pr[\Phi]$

Stochastic SAT (SSAT)

The Semantics

Let X be the left-most variable in the prefix of Φ . The recursive semantics of a SSAT formula are

1. $\Pr[\text{true}] = 1, \Pr[\text{false}] = 0$
2. $\Pr[\Phi] = \max_X \{ \Pr[\Phi|_X], \Pr[\Phi|_{\neg X}] \}$ if X is existentially quantified (\exists)
3. $\Pr[\Phi] = \min_X \{ \Pr[\Phi|_X], \Pr[\Phi|_{\neg X}] \}$ if X is universally quantified (\forall)
4. $\Pr[\Phi] = p \Pr[\Phi|_X] + (1 - p) \Pr[\Phi|_{\neg X}]$ if X is randomized quantified (\forall^p)

Stochastic SAT (SSAT)

A Tale of Two Encodings

- **Existential-random SSAT formula**

$$\Phi_{\text{ER}} = \exists X_2, \exists X_3, \forall^{0.25} X_1, (X_1 \vee \neg X_2) \wedge (\neg X_1 \vee X_2 \vee X_3) \wedge \neg X_1$$

- $\Pr[\Phi_{\text{ER}}] = 0.75$
- Optimal assignment (maximization): $X_2 = \text{false}, X_3 = \text{false}$

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- **Universal-random SSAT formula**

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- $\Pr[\Phi_{\text{UR}}] = 0$
- Optimal assignment (minimization): $X_2 = \text{true}, X_3 = \text{false}$

Justicia: Fairness Verification with SSAT

Consider

- features $\mathbf{X} \cup \mathbf{A}$ are Boolean
- predicted class \hat{Y} is a CNF formula $\phi_{\hat{Y}}$ defined on $\mathbf{X} \cup \mathbf{A}$

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Two Steps to Justicia

1. Computing $\max_{\mathbf{a}} \Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$, is equivalent to solving

$$\Phi_{\text{ER}} \triangleq \underbrace{\exists A_1, \dots, \exists A_n}_{\text{sensitive features}}, \underbrace{\forall^{p_1} X_1, \dots, \forall^{p_m} X_m}_{\text{non-sensitive features}}, \phi_{\hat{Y}}.$$

2. For computing $\min_{\mathbf{a}} \Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$, we substitute \exists with \forall for sensitive features, and observe $\Pr[\Phi_{\text{UR}}] = 1 - \Pr[\Phi_{\text{ER}}(\neg \phi_{\hat{Y}})]$.

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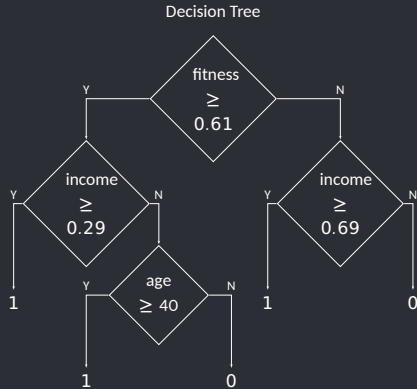
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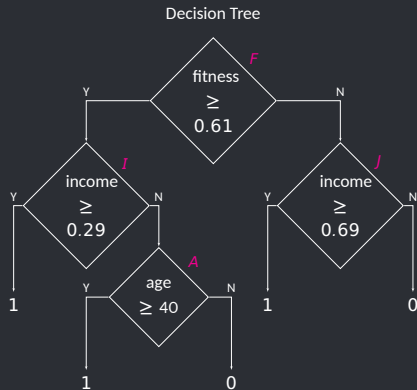
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Use an SSAT solver to solve the ER-SSAT problems [2].

An Illustration



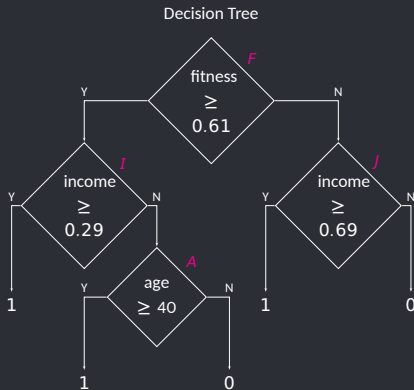
An Illustration



- CNF representation: $(\neg F \vee I \vee A) \wedge (F \vee J)$
- $\Pr[F] = 0.41, \Pr[I] = 0.93, \Pr[J] = 0.09$
- To compute $\max_a \Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$, we construct

$$\Phi_{\text{ER}} = \exists A, \mathbb{R}^{0.41}_F, \mathbb{R}^{0.93}_I, \mathbb{R}^{0.09}_J, (\neg F \vee I \vee A) \wedge (F \vee J)$$

An Illustration

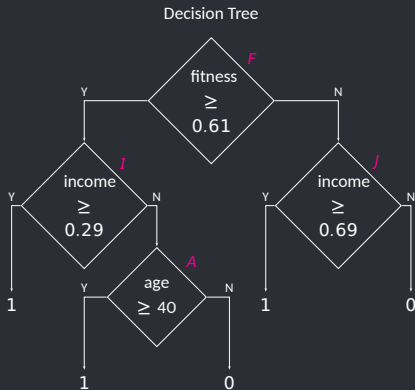


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- $\max_a \Pr[\hat{Y} = 1 | A = a] = \Pr[\Phi_{ER}] = 0.46$
- Similarly, $\min_a \Pr[\hat{Y} = 1 | A = a] = 0.43$
- Statistical parity is $0.46 - 0.43 = 0.03$

Theoretical Analysis

Pseudologarithmic Sample Complexity

Theorem (A PAC Bound for Justicia)

With probability $1 - \delta$, Justicia can estimate Statistical Parity (SP) up to a multiplicative error $2\epsilon_0$, i.e. $\widehat{SP} \leq 2\epsilon_0 SP$, if it has access to

$$k = O \left(\left(n + \ln \left(\frac{1}{\delta} \right) \right) \frac{\ln m}{\ln \epsilon_0} \right)$$

samples from the data-generating distribution.

Here, m and n are the number of variables with randomised and existential quantifiers respectively. Note that $\delta \in (0, 1)$ and $\epsilon_0 > 1$.

Experimental Analysis

Robustness and Compound Attribute Level Analysis

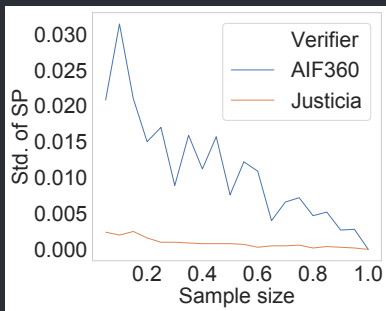


Figure: Robustness between probabilistic (Justicia) and dataset centric (AIF360 [3]) verifiers

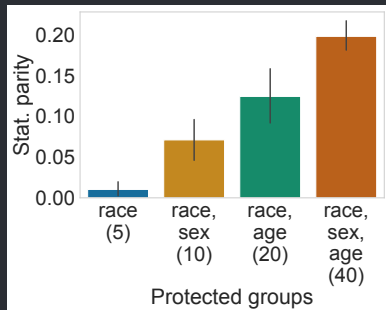


Figure: Verifying compound sensitive/protected groups with Justicia

Experimental Results

Faster than the Fastest

State-of-the-art probabilistic fairness verifiers

- FairSquare: computes weighted volume of logical programs using SMT reduction [4]
- VeriFair: probabilistic verification via sampling [5]

Dataset	FairSquare	VeriFair	Justicia
Ricci	4.8	5.3	0.1
Titanic	16	1.2	0.1
COMPAS	36.9	15.9	0.1
Adult	—	295.6	0.2

Table: Runtime of different verifiers in terms of execution time (in seconds) with decision tree classifiers. '—' refers to timeout.

Summary of Justicia

What Justicia can do?

- Justicia is a SSAT based probabilistic fairness verifier
- First method to verify compound sensitive groups
- More scalable in verifying decision trees and classifiers in Boolean formulas

What Justicia cannot do?

- Classifiers have to be expressed as Boolean formulas, which is computationally expensive even for linear classifiers
- Assumption of probabilistic independence of features leads to incorrect estimates

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FVGM: Algorithmic Fairness Verification with Graphical Models [6]

Fairness verification of Linear Classifiers

Challenges of earlier fairness verifiers

- **Scalability:** SSAT or SMT-based reduction of linear classifiers is computationally expensive
- **Accuracy:** Feature correlation is imprecisely modelled

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Proposed solutions

- **Scalability:** Novel stochastic subset-sum problem (S3P) based reduction
- **Accuracy:** Feature correlations represented as a Bayesian network

Linear Classifiers

Let

- w_{X_i} be the the weight/coefficient of non-sensitive feature X_i
- w_{A_j} be the the weight/coefficient of sensitive feature A_j
- τ is the offset parameter

The prediction of a binary linear classifier

$$\hat{Y} = \mathbb{1}\left[\sum_i w_{X_i} X_i + \sum_j w_{A_j} A_j \geq \tau\right].$$

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Our Approach: Compute the maximum and minimum of $\Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$
by a reduction to S3P

A Detour to Subset-sum Problem

- $\mathbf{B} \triangleq \{B_i\}_{i=1}^{|\mathbf{B}|}$ be a set of Boolean variables
- $w_i \in \mathbb{Z}$ be the weight of B_i
- a constant threshold $\tau \in \mathbb{Z}$

Given a constraint

$$\sum_{i=1}^{|\mathbf{B}|} w_i B_i = \tau$$

the subset-sum problem computes $\mathbf{b} \in \{0, 1\}^{|\mathbf{B}|}$ such that the constraint evaluates to true when \mathbf{B} is substituted with \mathbf{b}

Example:

- weights $\{-7, -3, -2, 9000, 5, 8\}$ and $\tau = 0$
- $\mathbf{b} = [0, 1, 1, 0, 1, 0]$ is the solution of the subset-sum problem, since $-3 - 2 + 5 = 0$

Stochastic Subset-sum Problem (S3P)

A Counting Analogue of the Subset-Sum Problem

S3P computes the *probability* of a subset of \mathbf{B} with sum of weights of non-zero variables to be at least τ . Formally,

$$s(\mathbf{B}, \tau) \triangleq \Pr \left[\sum_i w_i B_i \geq \tau \right] \in [0, 1].$$

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Similar to SSAT, we consider a quantifier $q_i \in \{\forall^{p_i}, \exists, \forall\}$ for each B_i in S3P

Stochastic Subset-sum Problem (S3P)

The Semantics

Let $\mathbf{B}[2 : n] \triangleq \{B_j\}_{j=2}^n$ be the subset of \mathbf{B} without the first variable B_1 .

$S(\mathbf{B}, \tau)$ is recursively defined as

$$S(\mathbf{B}, \tau) = \begin{cases} \mathbb{1}[\tau \leq 0], & \text{if } \mathbf{B} = \emptyset \\ S(\mathbf{B}[2 : n], \tau - \max\{w_1, 0\}), & \text{if } q_1 = \exists \\ S(\mathbf{B}[2 : n], \tau - \min\{w_1, 0\}), & \text{if } q_1 = \forall \\ p_1 \times S(\mathbf{B}[2 : n], \tau - w_1) + (1 - p_1) \times S(\mathbf{B}[2 : n], \tau), & \text{if } q_1 = \mathfrak{P}^{p_1} \end{cases}$$

Stochastic Subset-sum Problem (S3P)

Differences of S3P with SSAT

- Computation of \exists and \forall quantified variables is linear in S3P but exponential in SSAT.
- There is a pseudo-polynomial dynamic programming algorithm for S3P compared to the NP^{PP} -hardness of ER-SSAT and UR-SSAT.

FVGM: Fairness Verification of Linear Classifiers

1. Preprocess a linear classifier

- discretize each continuous feature X to a set of Boolean features \mathbf{B} using histogram
- if w is the weight of X and μ_i is the mean of feature values in the i -th bin, then the weight of $B_i \in \mathbf{B}$ is $w\mu_i$

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2. Learn a Bayesian network on discretized features¹
3. To compute $\max_a \Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$,
 - assign \exists quantifier to sensitive features \mathbf{A}
 - assign \forall quantifier to non-sensitive features \mathbf{X}
 - solve S3P problem

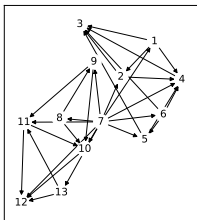
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4. To compute $\min_a \Pr[\hat{Y} = 1 | \mathbf{A} = \mathbf{a}]$, assign \forall quantifier to \mathbf{A} while keeping \exists quantifier on \mathbf{X}

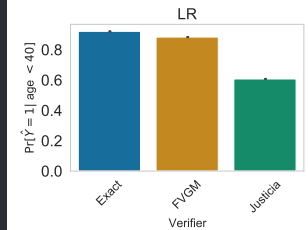
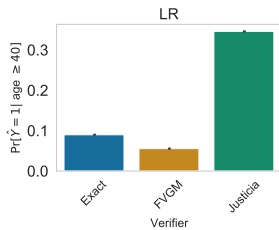
Experimental Analysis

Accuracy

- Sensitive features, $\mathbf{A} = \{\text{age}\}$
- Non-sensitive features, $\mathbf{X} = \{\text{health}, \text{income}\}$
- We discretize \mathbf{X} to Boolean features



1: $0.12 \leq \text{income} < 0.24$
2: $0.24 \leq \text{income} < 0.37$
3: $0.37 \leq \text{income} < 0.49$
4: $0.49 \leq \text{income} < 0.62$
5: $0.62 \leq \text{income} < 0.74$
6: $0.74 \leq \text{income} \leq 0.87$
7: $\text{age} \geq 40$
8: $-0.01 \leq \text{health} < 0.15$
9: $0.15 \leq \text{health} < 0.32$
10: $0.32 \leq \text{health} < 0.48$
11: $0.48 \leq \text{health} < 0.65$
12: $0.65 \leq \text{health} < 0.81$
13: $0.81 \leq \text{health} \leq 0.98$



Experimental Analysis

Scalability

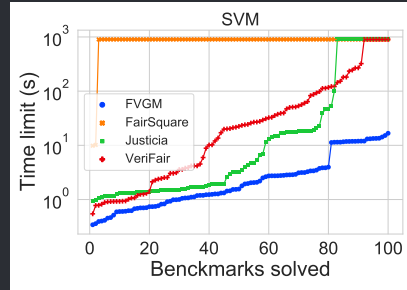
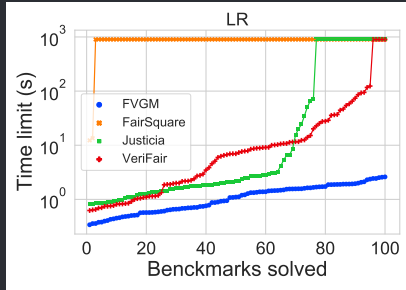


Figure: A cactus plot to present the scalability of different fairness verifiers on Linear Regression (LR) classifiers and Support Vector Machine (SVM)

Summary of FVGM

- FVGM is an efficient fairness verification framework for linear classifiers based on a novel stochastic subset-sum problem (S3P).
- FVGM is the first method to include feature correlations using a Bayesian network.
- FVGM demonstrates higher *scalability* and higher *accuracy* in comparison with earlier fairness verifiers.

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Fairness Explanation

- Identification of the source of unfairness is important to take affirmative actions
- Data contains bias and classifiers trained on the data inherit the bias.

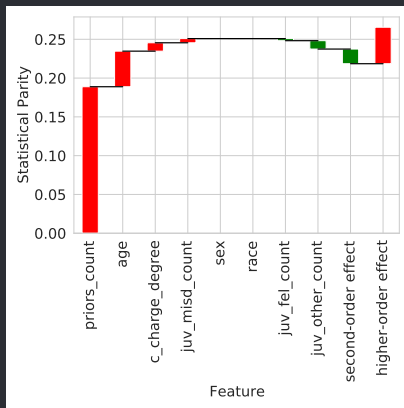


Figure: Explaining statistical parity in COMPAS recidivism prediction dataset for the feature 'sex'

Computing the Fairness Explanations

A Model-agnostic Approach

Observations

- Fairness, particularly group fairness, is a global property of the classifier.
- Fairness computation is equivalent to computing *the sensitivity of the classifier* w.r.t. different sensitive groups

Our approach: Extend global sensitivity analysis techniques from functional analysis to classification for explaining fairness.

FairXplain: Key Ideas

Idea 1

Statistical parity can be computed using the difference between variance of outcomes for sensitive groups

If $p_a \triangleq \max_a \Pr[\hat{Y} = 1 | A = a]$ and $p_{a'} \triangleq \min_{a'} \Pr[\hat{Y} = 1 | A = a']$,

$$\begin{aligned} \text{Statistical Parity} &= \frac{\text{Var}[\hat{Y} = 1 | A = a] - \text{Var}[\hat{Y} = 1 | A = a']}{1 - (p_a + p_{a'})} \\ &= \frac{\sum_{i=1}^n \overbrace{(v_i^{(a)} - v_i^{(a')})}^{\text{1-th order}} + \sum_{i < j}^n \overbrace{(v_{ij}^{(a)} - v_{ij}^{(a')})}^{\text{2-th order}} + \dots + \overbrace{(v_{12\dots n}^{(a)} - v_{12\dots n}^{(a')})}^{\text{n-th order}}}{1 - (p_a + p_{a'})} \end{aligned}$$

$$v_i^{(a)} = \text{Var}_{X_i}[E_{X_{\sim i}}[\hat{Y} = 1 | X_i, A = a]], \quad v_{ij}^{(a)} = \text{Var}_{X_{ij}}[E_{X_{\sim ij}}[\hat{Y} = 1 | X_i, X_j, A = a]] - v_i^{(a)} - v_j^{(a)}$$

FairXplain: Key Ideas

Idea 2

If we can decompose the variance in terms of the basis functions of the classifier, we can decompose the first and higher order variances as the variances of these decompositions.

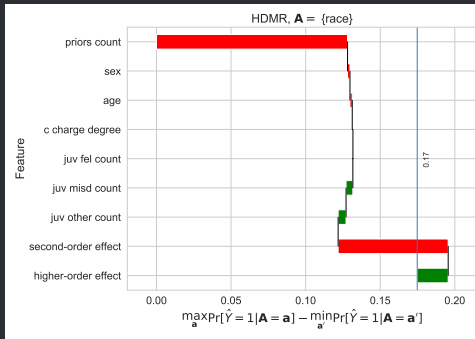
$$f_{\{i\}}(\mathbf{X}_{\{i\}}) \approx \sum_{r=-1}^{m+1} \alpha_r^i B_r(\mathbf{X}_{\{i\}})$$

$$f_{\{i,j\}}(\mathbf{X}_{\{i,j\}}) \approx \sum_{p=-1}^{m+1} \sum_{q=-1}^{m+1} \beta_{pq}^{ij} B_p(\mathbf{X}_{\{i\}}) B_q(\mathbf{X}_{\{j\}})$$

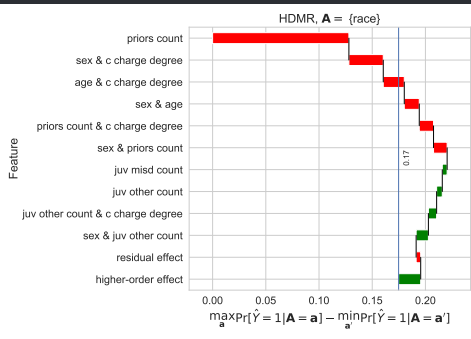
$$f_{\{i,j,k\}}(\mathbf{X}_{\{i,j,k\}}) \approx \sum_{p=-1}^{m+1} \sum_{q=-1}^{m+1} \sum_{r=-1}^{m+1} \gamma_{pqr}^{ijk} B_p(\mathbf{X}_{\{i\}}) B_q(\mathbf{X}_{\{j\}}) B_r(\mathbf{X}_{\{k\}})$$

Explaining Statistical Parity in COMPAS Dataset

Higher Order Effects are Decisive



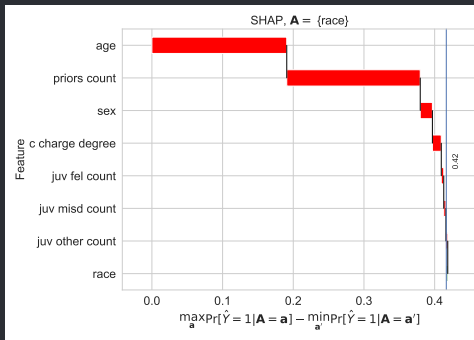
(a) FairXplain: First order effects



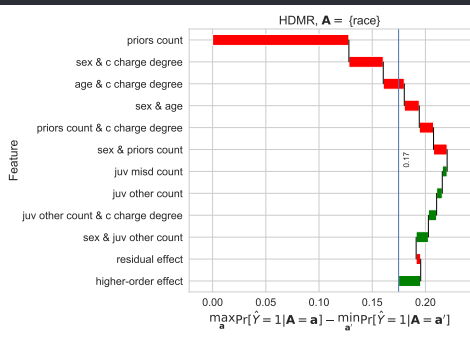
(b) FairXplain: First and second order effect

Explaining Statistical Parity in COMPAS Dataset

Local Explanations cannot Explain Unfairness



(c) Shapley Explanations



(d) FairXplain: First and second order effect

Conclusion

- **Fairness verification** and **explanation** are important problems in estimating the bias of classifiers and identifying the source of bias
- **Fairness verifiers**, Justicia and FVGM, improve upon existing fairness verifiers in terms of scalability and accuracy
- **Fairness explanation** shows the potential in identifying the effect of individual features or their interactions on the unfairness of the classifier. We currently focus in it.
- As a future work, we aim to design fairness enhancing algorithms relying on fairness verification and explanation

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Want to detect unfairness in your favourite classifier?

Use our Python library: “pip install justicia”

