Solution

- **Problem**: each process has a page table that maps all pages in its address space

- ✓ **Solution**: we only need to map the portion of the address space actually being used

- ➡ Use another level of indirection : **two-level page tables**

Two-Level Page Tables

Virtual addresses have three parts

- **a master page number** i.e the index in the master page table (a.k.a "page directory") that maps to a secondary page table

- **a secondary page number** i.e the index in the secondary page table (a.k.a "page table") that maps to the physical memory

- **an offset** that indicates where in physical page address is located