

ON TOUR | NEW YORK



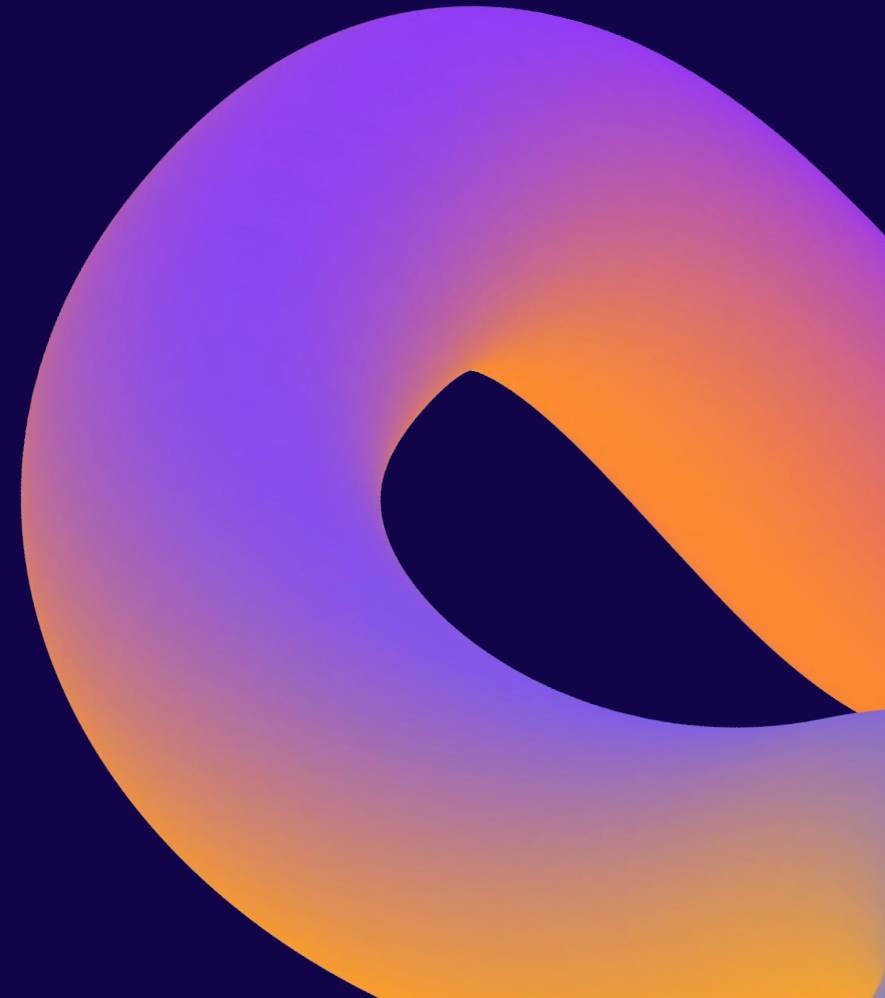
# Optimized Locking: Improving SQL Server Transaction Concurrency

**Deborah Melkin**

she\her

Data Engineer

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# Deborah Melkin

She/Her

Data Engineer  
Advisor360



- 25 years as a DBA
- Data Platform Women in Tech (WIT) Virtual UG, Co-leader
- WITspiration, Co-founder
- Redgate Community Ambassador
- Microsoft MVP, Data Platform
- In my spare time, I can usually be found doing something musical or something geeky with my husband, Andy, and our dog, Sebastian.



**@dgmelkin.bsky.social**



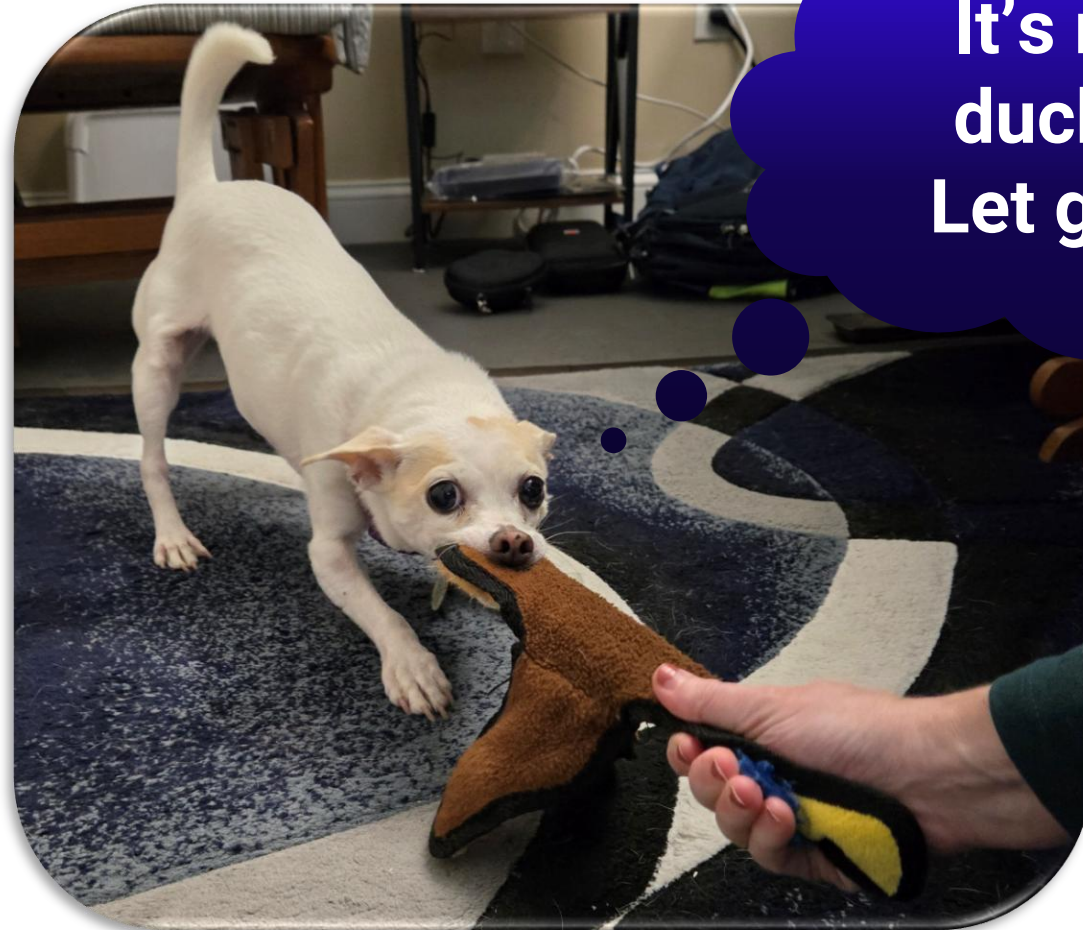
**github.com/DebtheDBA**



**DebtheDBA.wordpress.com**

# What's the problem we're trying to solve?

- Locking & Blocking!



It's my  
duck!!!  
Let go!!!

# Introducing Optimized Locking

- GA in Azure SQL DB in Feb 2023
- SQL Server 2025 Public Preview
- Writes not blocking Writes

# Agenda – How it works

- Components used by Optimized Locking
  - Read Committed Snapshot Isolation Level (RCSI)
  - Accelerated Database Recovery (ADR)
- Lock Escalation
- Optimized Locking
  - Transaction ID Locking
  - Lock After Qualification

**Warning:**  
**\* Internals Ahead \***  
**Level 300**



# Pessimistic vs Optimistic Locking

- **Pessimistic Locking** – Preventing users from modifying data in a way that affects other users
  - Read Uncommitted
  - Read Committed
  - Serializable
  - Repeatable Read

# Pessimistic vs Optimistic Locking

- **Optimistic Locking** – No locks when reading data but data is checked when modifying
  - Read Committed Snapshot Isolation
  - Snapshot Isolation
  - Optimized Locking

# Read Committed Snapshot Isolation Level (RCSI)

- Introduced with SQL Server 2005
- Writes and Reads don't block each other
- Two database options for setting this:
  - SET READ COMMITTED SNAPSHOT ISOLATION – sets default for the database
  - ALLOW SNAPSHOT ISOLATION LEVEL – allows individual sessions\queries to use SNAPSHOT isolation levels even if read committed snapshot isolation is not set.

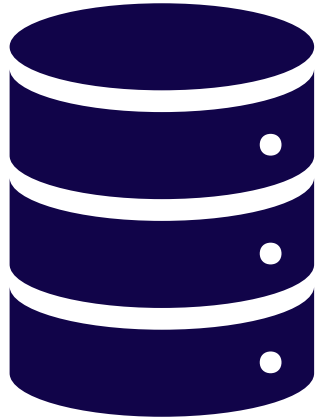


# Read Committed Snapshot Isolation Level (RCSI)

- Every row in the database has a Transaction ID (TID)
- Keeps copies of the previously committed versions of the records in the Version Store in TempDB
- Only reading the committed version of the data

BEGIN TRAN 1A

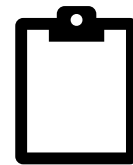
User DB



```
UPDATE Customer
SET State = 'MA'
WHERE FirstName = 'Sebastian'
```

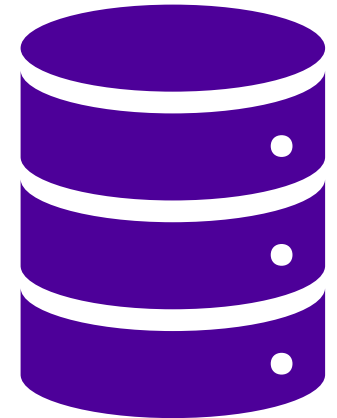


Sebastian	MA	Key	1A
-----------	----	-----	----



```
SELECT FirstName,
       LastName,
       State
FROM Customer
WHERE FirstName = 'Sebastian'
```

tempdb

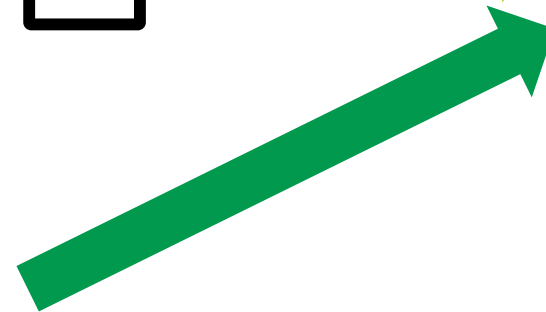


Version Store

Sebastian	AZ	1A	1
-----------	----	----	---

TID

Transaction Seq No.



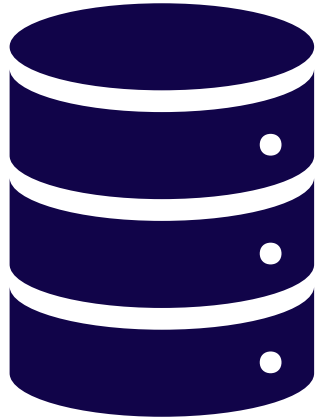
Read Committed  
Snapshot Isolation

# Accelerated Database Recovery (ADR)

- Introduced in SQL Server 2019
- Default for Azure SQL DB
  - Cannot be turned off
- Changes the way the transactions logs are read to be able to recover from long running transactions or just restore faster
- Uses a Persistent Version Store in Database rather than the Version Store in TempDB

BEGIN TRAN 1A

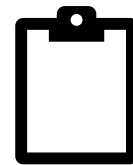
User DB



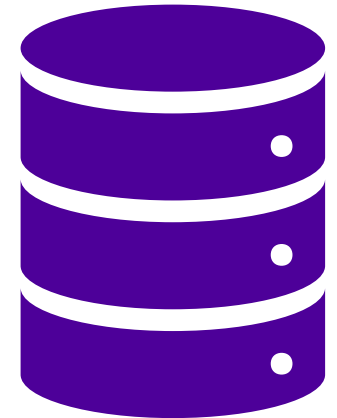
```
UPDATE Customer  
SET State = 'MA'  
WHERE FirstName = 'Sebastian'
```



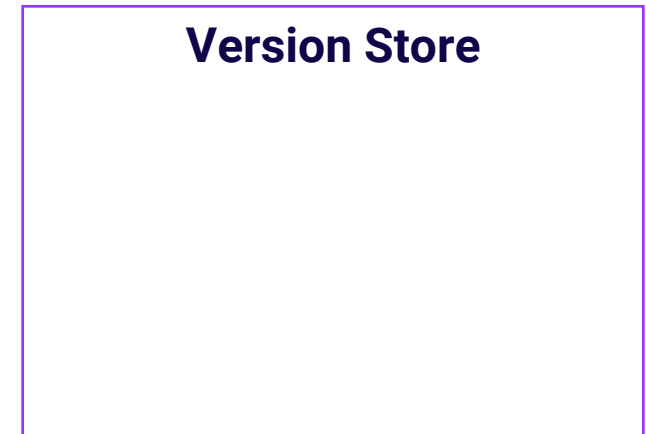
Sebastian	MA	 Key	1A
-----------	----	---	----



tempdb



Version Store



Persistent Version Store

Sebastian	AZ	1A	1
-----------	----	----	---



```
SELECT FirstName,  
        LastName,  
        State  
FROM Customer  
WHERE FirstName = 'Sebastian'
```

# Lock Escalation

- Threshold at which SQL Server will convert lower granularity locks to a higher level in order to manage a small number of locks
  - 5,000 on a table or index for a single statement
  - Each additional 1,250 locks taken in a transaction

# Lock Granularity (from Microsoft)

Resource	Description
RID	A row identifier used to lock a single row within a heap.
KEY	A row lock to lock a single row in a B-tree index.
PAGE	An 8 kilobyte (KB) page in a database, such as data or index pages.
EXTENT	A contiguous group of eight pages, such as data or index pages.
HoBT	A heap or B-tree. A lock protecting a B-tree (index) or the heap data pages in a table that doesn't have a clustered index.
TABLE	The entire table, including all data and indexes.
FILE	A database file.
APPLICATION	An application-specified resource.
METADATA	Metadata locks.
ALLOCATION_UNIT	An allocation unit.
DATABASE	The entire database.
XACT	Transaction ID (TID) lock used in Optimized Locking

# Lock Escalation (cont'd)

- Row and key locks escalate to table locks, not page locks
- “Queries Decide the Isolation Level” (Erik Darling)

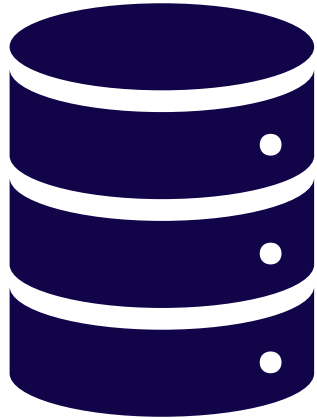
# Optimized Locking

- Changes to the locking mechanisms
  - Locks on the Transaction ID
  - Releases locks on other objects quickly
  - Minimizes lock escalation
- Accelerated Database Recovery is Required



BEGIN TRAN 1A

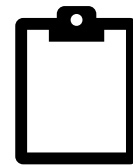
User DB



```
UPDATE Customer  
SET State = 'MA'  
WHERE FirstName = 'Sebastian'
```



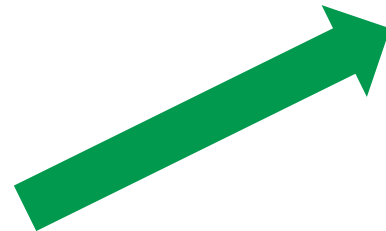
Sebastian	MA	Key	
-----------	----	-----	---



Persistent Version Store

Sebastian	AZ	1A	1
-----------	----	----	---

```
SELECT FirstName,  
        LastName,  
        State  
FROM Customer  
WHERE FirstName = 'Sebastian'
```



Optimized Locking -  
Reads

**Yeah, but that was just  
a SELECT statement...**

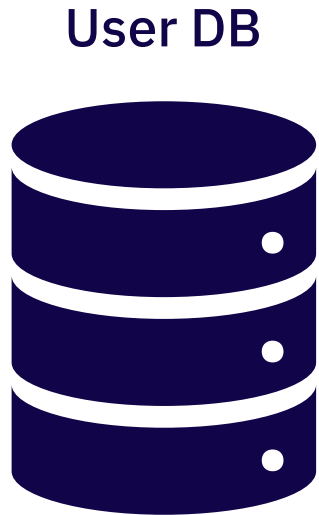
**Regular RCSI stuff.**

**Now what about writes?**



# Transaction ID Locking

- As the lock is held on the Transaction ID, the next transaction can see when rows aren't affected and update them without waiting.

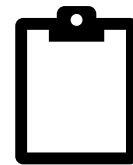


BEGIN TRAN 1A

UPDATE Customer  
SET State = 'MA'  
WHERE FirstName = 'Sebastian'



Sebastian	MA		
Deborah	MA		



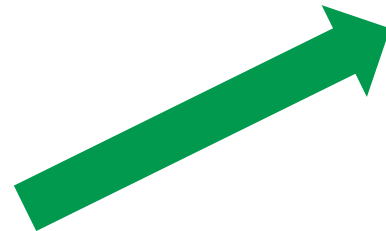
### Persistent Version Store

Sebastian	AZ	1A	1
Deborah	TN	2A	1



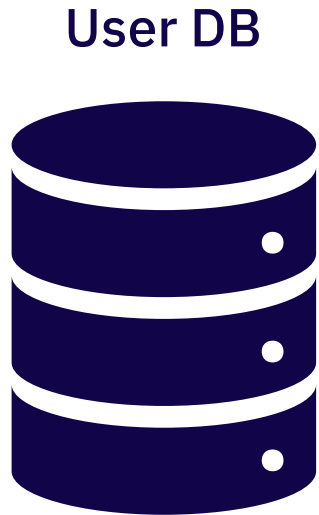
BEGIN TRAN 2A

UPDATE Customer  
SET State = 'MA'  
WHERE FirstName = 'Deborah'



# Lock After Qualification (LAQ)

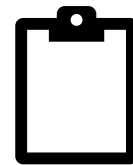
- RCSI is a requirement for this to work
- Looks to see if the previously committed version of the rows with TID locks would also be affected by the current transaction
  - If yes, transaction ***will wait*** for previous transaction to finish and include those columns as part of the change
  - If no, the transaction will update the rows that do match



BEGIN TRAN 1A

```
UPDATE Customer  
SET State = 'GA'  
WHERE FirstName = 'Sebastian'
```

COMMIT TRAN 1A



BEGIN TRAN 2A

```
UPDATE Customer  
SET State = 'MA'  
WHERE FirstName = 'Sebastian'
```

Persistent Version Store			
Sebastian	AZ	1A	1
Sebastian	GA	2A	1

# Demos!



# Be aware of

- Persistent Version Store
  - Make sure there is enough space available
    - When full, only reads can happen – no writes
  - Asynchronous Cleaner
    - Data remains until it's no longer needed
  - Keep Transactions Fast

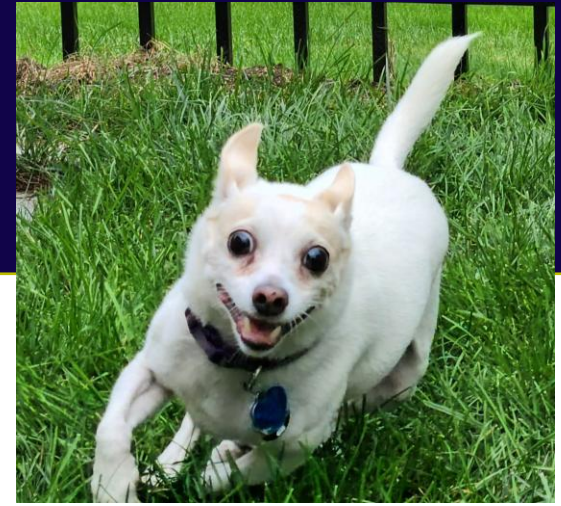


# Be aware of

- Lock After Qualification
  - Blocking and waits can still occur
  - Be aware of multiple transactions that may change data for the same set of records

# Recap

- Understand how RCSI & ADR are used for optimized locking
- Minimizes blocking but does not solve ALL situations
- Reduces the number of locks being held and the length they are held for
- Writes don't block Reads OR Writes



# Resources – Optimized Locking

- [SQL Server Transaction Locking and Row Versioning Guide](#)
- [Optimized Locking](#)
- [Bob Ward's Github - Optimized Locking Demos](#)
- [Isolation Level Locking \(Erik Darling\)](#)

# Resources - Accelerated Database Recovery

- [Soaring to New Heights with Accelerated Database Recovery \(John Morehouse\)](#)
- [Monitor and Troubleshoot Accelerated Database Recovery](#)
- [Why to Use Accelerated Database Recovery in SQL Server \(Luis Lema\)](#)
- [Constant Time Recovery in Azure SQL Database \(whitepaper\)](#)

# Your feedback is important to us



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# Thank you

Any Questions?

**Deborah Melkin**



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