High-Dimensional Probability for Data Science

Based on the PhD working group lectures

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CONTENTS

LECTURE 1: CONCENTRATION INEQUALITIES

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The object of the first lectures is trying to characterize deviations of sums of random variables X_i w.r. to their expected value $\mathbb{E}$. These *concentration inequalities* take for instance the form of

$$\mathbb{P}(|S - \mu| > t) \leq \text{Bound},$$

where the bound is tighter than what we usually obtain using the standard inequalities that are presented in a first course in probability. In particular, we are <u>not</u> looking for asymptotic results as in the central limit theorem, but rather for estimates which are valid for any sample size N.

1.1 Hoeffding's inequality

Let us begin by recalling two standard inequalities which are going to be especially useful in the following sections.

Thm. 1 (Markov's inequality)

Let $X \geq 0$ be a random variable with finite expected value, $\mathbb{E}[X] < \infty$, then

$$\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}, \quad \text{for all } t > 0.$$

A straightforward consequence of Markov's inequality can be obtained by replacing the random variable X with $|X - \mu|$ and squaring both sides inside the probability operator, which yields the following inequality.

Corollary 1 (Chebyshev's inequality)

If X is a random variable with finite variance, $\mathbb{V}[X] < \infty$, then

$$\mathbb{P}(|X - \mu| \geq t) \leq \frac{\mathbb{V}[X]}{t^2}.$$

Remark   Many of the arguments that we make in this lecture will be based on the following trick: for any random variable X and for any $\lambda > 0$,

$$\mathbb{P}(X - \mu \geq t) = \mathbb{P}(e^{\lambda(X-\mu)} \leq e^{\lambda t}) \qquad \text{(monotone)}$$

$$\leq e^{-\lambda t}\mathbb{E}[e^{\lambda(X-\mu)}] \qquad \text{(Markov)}$$

Now, since it holds for any choice of $\lambda > 0$ we can obtain the tightest bound by optimizing w.r. to λ,

$$\mathbb{P}(X - \mu \geq t) \leq \inf_{\lambda>0} e^{-\lambda t}\mathbb{E}[e^{\lambda(X-\mu)}],$$

and since X is usually a sum of random variables, its characteristic function can be decomposed into a product and evaluated quite easily.

Thm. 2 (Hoeffding's inequality)

Let $X_1, \ldots, X_N$ be i.i.d Rademacher$(\frac{1}{2})$ random variables and $a_1, \ldots, a_N \in \mathbb{R}$, then for any $t > 0$ we have

$$\mathbb{P}\Big(\sum_{i=1}^{N} a_i X_i \geq t\Big) \leq \exp\Big(-\frac{t^2}{2\|a\|_2^2}\Big)$$

Sample size Unlike standard concentration inequalities based on the central limit theorem, this inequality gives an exact bound for any value of N.

Tightness Moreover, we can see that the tail behaviour, i.e. $\mathbb{P}(Y \geq t)$, is Gaussian-like in t, which means that this bound is extremely tight.

Proof.
Suppose that $\|a\|_2 = 1$, otherwise we can rescale t accordingly. For $\lambda > 0$, we have

$$\mathbb{P}\Big(\sum_{i=1}^{N} a_i X_i \geq t\Big) \overset{\text{Markov}}{\leq} e^{-\lambda t}\mathbb{E}[e^{\lambda \sum_{i=1}^{N} a_i X_i}]$$

$$= e^{-\lambda t}\prod_{i=1}^{N} \underbrace{\mathbb{E}[e^{\lambda a_i X_i}]}_{\frac{1}{2}e^{\lambda a_i} + \frac{1}{2}e^{-\lambda a_i}} \qquad \text{(Indep.)}$$

$$= e^{-\lambda t}\prod_{i=1}^{N} \cosh(\lambda a_i) \qquad (\tfrac{1}{2}e^x + \tfrac{1}{2}e^{-x} = \cosh(x))$$

$$\leq e^{-\lambda t}e^{\frac{\lambda^2}{2}\sum_{i=1}^{N} a_i^2} \qquad (\cosh(x) \leq e^{\frac{x^2}{2}}, \text{ see here})$$

Now, if we want to find the optimal bound, $\lambda_{\text{opt}} = \inf_{\lambda > 0} e^{-\lambda t + \frac{\lambda^2}{2}\|a\|_2^2}$, we first notice that the function inside the exponent is parabolic in λ,

$$f(\lambda) = -\lambda t + \frac{\lambda^2}{2}\|a\|_2^2 \overset{\text{parabola}}{\Longrightarrow} \lambda_{\text{opt}} = \frac{t}{\|a\|_2^2} \implies f(\lambda_{\text{opt}}) = -\frac{t^2}{2\|a\|_2^2}.$$

Therefore, by substituting the optimal λ we obtain the proof of Hoeffding's inequality,

$$\mathbb{P}\Big(\sum_{i=1}^{N} a_i X_i \geq t\Big) \leq e^{-\frac{t^2}{2\|a\|_2^2}}.$$

$\square$

Exercise Restate Hoeffding's inequality for $X_1, \ldots, X_N \overset{\text{iid}}{\sim} \text{Ber}(\frac{1}{2})$, using the fact that $Z_i = 2X_i - 1$ with $Z_i \sim \text{Rademacher}(\frac{1}{2})$.

Exercise Use Hoeffding's inequality for Bernoulli random variables to prove that by tossing a coin N times we have the exact bound

$$\mathbb{P}\Big(\text{at least } \frac{3}{4} \text{ heads}\Big) \leq e^{-N/8}.$$

Remark   We can get a double bound from the above 2 by using $\mathbb{P}(|S| \geq t) \leq \mathbb{P}(S \geq t) + \mathbb{P}(-S \geq t)$, and observing that the Rademacher r.v. is symmetric $S = -S$. Therefore, both bounds are equal and the following two-sided inequality can be stated.

> **Thm. 3 (Two-sided Hoeffding's inequality)**
>
> Let $X_1, \dots, X_N$ be i.i.d Rademacher r.v.'s, then for all $t \geq 0$ and for all $a \in \mathbb{R}^N$,
>
> $$\mathbb{P}\Big(\Big|\sum_{i=1}^{N} a_i X_i\Big| \geq t\Big) \leq 2\exp\Big(-\frac{t^2}{2\|a\|_2^2}\Big).$$

We now turn to the more general problem of bounded random variables, which include as a special case the setting of Bernoulli r.v.'s with varying parameter p_i.

> **Thm. 4 (Hoeffding's inequality for bounded r.v.'s)**
>
> Let $X_1, X_2, \dots, X_N$ be independent but not identically distributed r.v.'s, such that $X_i \in [m_i, M_i]$ and $\mathbb{E}[X_i] < \infty$. Then, for all $t \geq 0$ the following inequality holds,
>
> $$\mathbb{P}\Big(\sum_{i=1}^{N}(X_i - \mathbb{E}[X_i]) \geq t\Big) \leq \exp\Big(-\frac{2t^2}{\sum_{i=1}^{N}(M_i - m_i)^2}\Big).$$

Proof.
(Exercise 2.2.7 in the book) The difficult part is achieving the constant 2 in the numerator, therefore we start with a different constant and then use a trick to get it. Let $\lambda > 0$, then by the same argument as before we can write

$$\mathbb{P}\Big(\sum_{i=1}^{N}(X_i - \mathbb{E}[X_i]) \geq t\Big) \leq e^{-\lambda t}\mathbb{E}[e^{\lambda \sum_i X_i - \mathbb{E}[X_i]}]$$

$$= e^{-\lambda t}\prod_i \mathbb{E}[e^{\lambda(X_i - \mathbb{E}[X_i])}]$$

$$\leq e^{-\lambda t + \sum_i \lambda(M_i - m_i)}$$

This is not as easy to optimize as before since we don't have a quadratic form, therefore we need a subtle trick to transform it into a more easily handled problem.

Trick   In order to replace "$\cosh x \leq e^{x^2/2}$" we can use the following trick: Let Y be a r.v. with $\mathbb{E}[Y] = 0$ (our case of $X - \mathbb{E}[X]$) and $Y \in [a, b]$, then for all $\lambda \in \mathbb{R}$,

$$\mathbb{E}[e^{\lambda Y}] \leq e^{\lambda^2 \frac{(b-a)^2}{2}}.$$

This is based on a symmetrization of Y by introducing another independent random variable $Y' \overset{\mathrm{d}}{=} Y$ and $Z \sim \text{Rademacher}(\frac{1}{2})$ from which we have $\mathbb{E}[e^{-\lambda Y'}] \overset{\text{Jens.}}{\leq} e^{-\lambda \mathbb{E}[Y]} = 1$, therefore

$$\mathbb{E}[e^{\lambda Y}] \leq \mathbb{E}[e^{\lambda Y}]\cdot\mathbb{E}[^{-\lambda Y'}] = \mathbb{E}[e^{\lambda(Y - Y')}] = \mathbb{E}[e^{\lambda Z(Y - Y')}] = \mathbb{E}[\cosh(\lambda(Y - Y'))] \leq \mathbb{E}[e^{\lambda^2 \frac{(Y - Y')^2}{2}}] = e^{\frac{\lambda^2(b-a)^2}{2}}.$$

Using this trick, we can optimize the equation using

$$\mathbb{P}\Big(\sum_{i=1}^{N}(X_i - \mathbb{E}[X_i]) \geq t\Big) \leq e^{-\lambda t} \prod_i e^{\lambda^2 \frac{(M_i - m_i)^2}{2}}$$

$$= \exp\Big(-\lambda t + \frac{\lambda^2}{2}\sum_i \frac{(M_i - m_i)^2}{2}\Big).$$

We can optimize with $\lambda > 0$ and get the minimum with a different constant than 2. Finding this other minimum requires more work.

$\square$

Example (Book 2.2.9 – Boosting a randomized algorithm)

We have an algorithm that gives the right answer out of two classes with a probability $\frac{1}{2} + \delta$, with $\delta > 0$. We run this algorithm N (odd) times and take the majority vote to get the final classification.

Problem   Find the minimal N such that $\mathbb{P}(\text{correct answer}) \geq 1 - \varepsilon$ for $\varepsilon \in (0,1)$ fixed.

Solution   Consider the following r.v. $X_1, \ldots, X_N$ be the indicator of the wrong answer

$$X_i = \begin{cases} 1 & \text{if } i^{\text{th}} \text{ run is wrong} \\ 0 & otherwise \end{cases}$$

then, using thm. 4 with $t = N\delta$, $M_i = 1$ and $m_i = 0$ we can bound the probability of wrong answer as

$$\mathbb{P}\Big(X_1 + \ldots + X_N \geq \frac{N}{2}\Big) = \mathbb{P}\Big(\sum_{i=1}^{N}(X_i - (\frac{1}{2} - \delta)) \geq N\delta\Big) \overset{4}{\leq} \exp\Big(-\frac{2N^2\delta^2}{N}\Big).$$

Therefore, in order to have the required bounded probability we need

$$-2N\delta^2 \leq \log\varepsilon \iff \boxed{N \geq \frac{1}{2\delta^2}\log\frac{1}{\varepsilon}}.$$

1.2   Chernoff's inequality

Consider the last Hoeffding's inequality (thm. 4), then for a sum of random variables we can write the Gaussian tail using the CLT as approximately

$$\mathbb{P}(|Z| \geq t) \leq 2e^{-\frac{t^2}{2}}.$$

Chernoff's inequality is useful in regimes of sums in order to prove a bound that is again independent from the central limit theorem. The following theorem is a merged result of Theorem 2.3.1, Exercise 2.3.2 and Exercise 2.3.5 in the book.

Thm. 5 (Chernoff's inequality)

Let $X_1, \dots, X_N$ be such that $X_i \overset{iid}{\sim} Bern(p_i)$ and consider the cumulative sum $S_N = \sum_i X_i$ with expected value $\mu = \mathbb{E}[S_N] = \sum_i p_i$. Then, the following inequalities hold:

$$\mathbb{P}(S_N \geq t) \leq e^{-\mu} \cdot \left(\frac{e\mu}{t}\right)^t \qquad for\ t > \mu,$$

$$\mathbb{P}(S_N \leq t) \leq e^{-\mu} \left(\frac{e\mu}{t}\right)^t \qquad for\ t < \mu,$$

"SMALL DEVIATIONS": $\mathbb{P}(|S_N - \mu| \geq \delta\mu) \leq 2e^{-C\mu\delta^2} \qquad for\ \delta \in (0, 1],$

where C is a universal constant (i.e. does not depend on the other quantities).

Proof.

1. The first step is always the same, let $\lambda > 0$ then

$$\mathbb{P}(S_N \geq t) = \mathbb{P}(e^{\lambda S_N} \geq e^{\lambda t}) \leq e^{-\lambda t}\mathbb{E}[e^{\lambda S_N}] = e^{-\lambda t}\prod_i \mathbb{E}[e^{\lambda X_i}]. \qquad (1)$$

   Now for a Bernoulli random variable, $\mathbb{E}[e^{\lambda X_i}] = (1 - p_i)e^0 + p_i e^\lambda = 1 + (e^\lambda - 1)p_i$, and we use the following identity:
$$1 + x \leq e^x \quad for\ all\ x > 0,$$

   to write

$$\mathbb{E}[e^{\lambda X_i}] = 1 + \overbrace{(e^\lambda - 1)p_i}^{x} \leq \exp\left((e^\lambda - 1)p_i\right).$$

   Going back to (1), we have the following bound for any $\lambda > 0$,

$$\mathbb{P}(S_N \geq t) \leq e^{-\lambda t}e^{(e^\lambda - 1)\sum_i p_i} = e^{-\lambda t + \mu(e^\lambda - 1)}.$$

   Again, by optimizing over λ we find that the tightest bound from (1) is given by

$$f(\lambda) = -\lambda t + \mu(e^\lambda - 1) \implies \lambda_{\text{opt}} = \operatorname*{argmin}_{\lambda > 0} f(\lambda) = \log\frac{t}{\mu},$$

   from which we obtain the first Chernoff bound,

$$\mathbb{P}(S_N \geq t) \leq e^{-\mu}\left(\frac{e\mu}{t}\right)^t.$$

2. For the second inequality, proceed as before using

$$\mathbb{P}(S_N \leq t) \overset{\lambda \geq 0}{=} \mathbb{P}(e^{-\lambda S_N} \geq e^{-\lambda t}).$$

3. We can obtain the bound on $\mathbb{P}(|S_N - \mu| \geq \delta\mu)$ by using the fact that

$$\mathbb{P}(|S_N - \mu| \geq \delta\mu) \leq \mathbb{P}(S_N - \mu \geq \delta\mu) + \mathbb{P}(S_N - \mu \leq -\delta\mu) \overset{(1),(2)}{\leq} \dots$$

$\square$

Thm. 6 (Poisson tail behaviour)

Let $Z \sim Pois(\gamma)$ with $\gamma > 0$, i.e. X has probability mass function $\mathbb{P}(X = k) = e^{-\gamma} \frac{\gamma^k}{k!}$, for $k = 0, 1, \ldots$. Then,

1. *For all $\delta \in (0, 1]$ thm. 5-3 holds*

$$\mathbb{P}(|Z - \gamma| \geq \delta\gamma) \leq 2e^{-C\lambda\delta^2}$$

2. *Let now $t > \gamma$, then the following bound holds*

$$\mathbb{P}(X \geq t) \leq e^{-\gamma} \left(\frac{e\gamma}{t}\right)^t \tag{A}$$

Remark   These bound are extremely useful in practical applications and is similar to Chernoff's bound (thm. 5), which works instead for a sum of Bernoulli variables.

Remark 2   If $p_i = \frac{\gamma}{N}$, then $S_N \approx Z \sim \text{Pois}(\gamma)$ for $N \gg 1$ and the rate of convergence is very fast, therefore this result could also be obtained as a limit. However, the above theorem is *exactly* valid.

Proof.

(Execise) Prove equation (A) using the basic trick $\mathbb{P}(X \geq t) \leq e^{-\lambda t}\mathbb{E}[e^{\lambda X}]$, which can be computed explicitly, and then optimize over $\lambda > 0$. Briefly comment on why this bound is optimal.

$\square$

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LECTURE 2: SUBGAUSSIAN RANDOM VARIABLES

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In this lecture we generalize Hoeffding's inequality to subgaussian random variables, which are a class of distributions that enjoy nice properties and are fundamental in the high-dimensional setting. We begin by recalling some properties of the Gaussian distribution

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Prop. 1 (Properties of the gaussian distribution)

Let $X \sim \mathcal{N}(0,1)$, then the following statements hold:

1. *We have a tail estimate for X given by*

$$\left(\frac{1}{t} - \frac{1}{t^3}\right) \frac{1}{\sqrt{2\pi}} e^{-\frac{t^2}{2}} \leq \mathbb{P}(X \geq t) \leq \frac{1}{t} \frac{1}{\sqrt{2\pi}} e^{-\frac{t^2}{2}}, \quad t > 0.$$

   This estimate in particular implies that

$$\mathbb{P}(X \geq t) \leq \frac{1}{\sqrt{2\pi}} e^{-\frac{t^2}{2}} \qquad t \geq 1,$$

$$\mathbb{P}(|X| \geq t) \leq 2 e^{-\frac{t^2}{2}} \qquad\qquad t \geq 0.$$

2. *Given $p \geq 1$, we have that*

$$\|X\|_{L^p} = \mathbb{E}[|X|^p]^{\frac{1}{p}} = \sqrt{2} \left(\frac{\Gamma(\frac{1+p}{2})}{\Gamma(\frac{1}{2})}\right)^{\frac{1}{p}}$$

3. *The moment-generating function of X is $\mathbb{E}[e^{\lambda X}] = e^{\frac{\lambda^2}{2}}$ for all $\lambda \in \mathbb{R}$.*

Corollary 2 (Bounded norm of a gaussian r.v.)

If $X \sim \mathcal{N}(0,1)$ there exists a $C > 0$ such that $\|X\|_{L^p} \leq C\sqrt{p}$ for all $p \geq 1$.

Proof.

Use Stirling's approximation for the Gamma function to obtain the bound.

$\square$

With these properties we can now discuss another class of random variables, which include the Gaussian distribution.

2.1 Space of subgaussian random variables

We begin the analysis of subgaussian random variables by stating a sequence of equivalent properties that turn out to be equivalent to each other.

Thm. 7 (Equivalence of properties for subgaussian r.v.'s)

Let X be a generic random variable, then the following properties are equivalent:

1. *(TAIL OF X) $\exists K_1 > 0$ such that $\mathbb{P}(|X| > t) \leq 2e^{-t^2/k_1^2}$ for all $t \geq 0$.*

2. *(MOMENTS OF X) $\exists k_2 > 0$ such that $\|X\|_{L^p} \leq k_2\sqrt{p}$ for all $p \geq 1$.*

3. *(MGF OF X^2) $\exists k_3 > 0$ such that $\mathbb{E}[e^{\lambda^2 X^2}] \leq e^{k_3^2 \lambda^2}$ for $|\lambda| \leq \frac{1}{k_3}$.*

4. *(MGF OF X^2) $\exists k_4 > 0$ such that $\mathbb{E}[e^{X^2/k_4^2}] \leq 2$.*

In addition, if $\mathbb{E}[X] = 0$ we can add another equivalent property:

5. *(MFG OF X) $\exists k_5 > 0$ such that $\mathbb{E}[e^{\lambda X}] \leq e^{k_5 \lambda^2}$ for all $\lambda \in \mathbb{R}$.*

Moreover, the above constants $k_1, \ldots, k_5$ differ by a constant factor, i.e. if one property holds then all properties hold and $\exists C_{ij} > 0$ such that

$$k_i \leq C_{ij} k_j \quad \text{for all } i, j \text{ and } C_{ij} \text{ does not depend on } X.$$

Proof.
Long and boring.

$\square$

Remark   5. really needs that $\mathbb{E}[X] = 0$, otherwise it does not work independently of X.

Given the usefulness of these bounds, it's important to isolate the class of r.v.'s that share these properties.

Def. (Subgaussian r.v.)

A r.v. X is called ***subgaussian*** if it satisfies one of the equivalent properties in thm. 7.

Thm. 8 (Subgaussian random variables form a vector space)

The set of subgaussian random variables is a vector space, which means that

$$X, Y \text{ subgaussian} \implies X + Y \text{ is subgaussian}$$

$$X \text{ subgaussian} \implies \alpha X \text{ is subgaussian}$$

Proof.
We aim to prove that $\mathbb{E}[e^{\frac{(X+Y)^2}{(a+b)^2}}] \leq 2$, we can consider

$$\frac{X+Y}{a+b} = \frac{a}{a+b}\frac{X}{a} + \frac{b}{b+a}\frac{Y}{b},$$

use the fact that e^{x^2} is convex to conclude that

$$e^{\frac{(x+y)^2}{(a+b)^2}} \leq \frac{a}{a+b}e^{\frac{x^2}{a^2}} + \frac{b}{a+b}e^{\frac{y^2}{b^2}}.$$

$\square$

> **Def. (Subgaussian norm)**
>
> Let X be a subgaussian r.v., then we define the **_subgaussian norm of X_** as
>
> $$\|X\|_{\psi_2} := \inf \left\{ t > 0 : \mathbb{E}[e^{X^2/t^2}] \leq 2 \right\}.$$

Remark   Take $t = k_4$ and we see that the set over which the inf is taken is never empty.

Remark 2   By dominated convergence this infimum is a minimum.

> **Prop. 2 (Subgaussian norm is indeed a norm)**
>
> $\|\cdot\|_{\psi_2}$ _is a norm on the space of subgaussian r.v.'s._

Proof.
Everything is simple, except for the triangle inequality which is not straightforward.

$\square$

Since we have that the optimal constant for property _4._ is given by the subgaussian norm $\|X\|_{\psi_2}$, then we have the following updated set of inequalities in terms of $k_4 = \|X\|_{\psi_2}^2$:

1. $\mathbb{P}(|X| > t) \leq 2e^{-\frac{Ct^2}{\|X\|_{\psi_2}^2}}$ for all $t \geq 0$.

2. $\|X\|_{L^p} \leq C\|X\|_{\psi_2}\sqrt{p}$ for all $p \geq 1$.

3. $\mathbb{E}[e^{\frac{X^2}{\|X\|_{\psi_2}^2}}] \leq 2$.

4. If $\mathbb{E}[X] = 0$, then $\mathbb{E}[e^{\lambda X}] \leq e^{C\lambda^2\|X\|_{\psi_2}^2}$.

> **Prop. 3 (Bounded r.v.'s are subgaussian)**
>
> _If X is a bounded random variable then X is subgaussian._

Proof.
$\|X\|_{\psi_2} \leq \dfrac{\|X\|_\infty}{\log 2}.$

$\square$

Non-subgaussian r.v.'s   Poisson, exponential, Pareto, Cauchy, . . .
For subgaussian random variables we have something similar to the property of Gaussian random variables

Prop. 4 (Sums of subgaussians)

Let $X_1, \ldots, X_N$ be i.i.d subgaussian random variables with $\mathbb{E}[X_i] = 0$ for all i. Then, $\sum_{i=1}^{N} X_i$ is subgaussian and

$$\Big\| \sum_{i=1}^{N} X_i \Big\|_{\psi_2}^2 \leq C \sum_{i=1}^{N} \|X_i\|_{\psi^2}^2.$$

Moreover, since $\| \cdot \|_{\psi_2}^2$ is a norm, we also have the bound

$$\Big\| \sum_{i=1}^{N} X_i \Big\|_{\psi_2}^2 \leq C \sum_{i=1}^{N} \|X_i\|_{\psi^2}.$$

Proof.
Since $\mathbb{E}[X_i] = 0$ then also $\mathbb{E}[\sum_i X_i] = 0$ and we use property $5.$ to show

$$\mathbb{E}[e^{\lambda \sum_i X_i}] \overset{5.}{\leq} \prod_i e^{C\lambda^2 \|X_i\|_{\psi_2}^2}$$

$$= e^{C\lambda^2 \sum_i \|X_i\|_{\psi_2}^2}.$$

Moreover, since the best constant is k_4 we have the norm.

$\square$

2.2 General Hoeffding's inequality

Subgaussian random variables are extremely useful since we have a general form of the Hoeffding's inequality without passing through Rademacher or boundedness.

Thm. 9 (General Hoeffding's inequality)

Let $X_1, \ldots, X_N$ be independent subgaussian random variables with $\mathbb{E}[X_i] = 0$ for all i. Then, for each $t \geq 0$ we have a tail estimate

$$\mathbb{P}\Big(|\sum_{i=1}^{N} X_i| \geq t \Big) \leq 2 \exp\left(-\frac{ct^2}{\sum_{i=1}^{N} \|X_i\|_{\psi_2}^2} \right).$$

Proof.
using the previous Prop. 4 and

$$\mathbb{P}(|X| > t) \leq 2e^{-\frac{Ct^2}{\|X\|_{\psi_2}^2}}, \quad \text{for all } t \geq 0.$$

Apply the above bound to the subgaussian r.v. $X := \sum_{i=1}^{N} X_i$ (by thm. 8) and use the bound on the norm given by Prop. 4.

$\square$

Corollary 3 (General Hoeffding's inequality 2)

Let $X_1, X_2, \ldots, X_n$ be independent subgaussian random variables with $\mathbb{E}[X_i] = 0$, and let $a_1, a_2, \ldots, a_n \in \mathbb{R}$. Then,

$$\mathbb{P}\Big(|\sum_{i=1}^{N} a_i X_i| \geq t\Big) \leq 2\exp\Big(-\frac{ct^2}{k^2\|a\|_2^2}\Big),$$

where $k = \max_i \|X_i\|_{\psi_2}^2$.

Proof.
Use again the same properties, recall that it is a norm and then bound using the maximum of a_i.
$\square$

Note   One can apply the theorem to $X_1, \ldots, X_N$ independent and subgaussian but have to replace X_i by $X_i - \mathbb{E}[X_i]$.

Recall that $\|X - \mathbb{E}[X]\|_{L^2} \leq \|X\|_{L^2}$. This does not hold for the subgaussian norm, however we do have a lemma in this direction.

Lemma 1 (Centering of a subgaussian r.v.)

Let X be subgaussian, then $X - \mathbb{E}[X]$ is subgaussian (vector space) and

$$\|X - \mathbb{E}[X]\|_{\psi_2} \leq C\|X\|_{\psi_2}.$$

Proof.
$\|\cdot\|_{\psi_2}$ is a norm, therefore

$$\|X - \mathbb{E}[X]\|_{\psi_2} \leq \|X\|_{\psi_2} + \|\mathbb{E}[X]\|_{\psi_2}$$

$$= \|X\|_{\psi_2} + |\mathbb{E}[X]| \cdot \|1\|_{\psi_2}$$

$$\leq \|X\|_{\psi_2} + \|X\|_{L^1} \cdot \|1\|_{\psi_2} \qquad (|\mathbb{E}[X]| \leq \mathbb{E}[|X|] = \|X\|_{L^1})$$

$$\leq \|X\|_{\psi_2} + C \cdot \|X\|_{\psi_2} \cdot \sqrt{1} \cdot \|1\|_{\psi_2} \qquad (\text{using } 2.)$$

$$\leq K\|X\|_{\psi_2}.$$

$\square$