

Statistical Computation

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1 Basics

1.1 Floating Point Representation

Definition 1.1. A *floating point number* is represented by three components: (S, F, E) where S is the sign of the number (± 1), F is a fraction (lying between 0 and 1), E is an exponent. S, F, E are all represented as binary digits (bits). The *floating point representation* of x , $\text{fl}(x)$ is

$$\text{fl}(x) = S \times F \times 2^E$$

Note. x and $\text{fl}(x)$ need not be the same, since $\text{fl}(x)$ is a binary approximation to x , and there are only a finite number of floating point numbers.

1.1.1 Round-Off Error

Mathematical operations introduce further approximation errors

$$f(\text{fl}(x)) = f(x + \varepsilon) \approx f(x) + \varepsilon f'(x)$$

and the goal is to make the round-off error $|f(x) - \text{fl}(f(\text{fl}(x)))|$ as small as possible.

1.1.2 Machine Epsilon and Other Constants

For a given real number x , we have

$$|\text{fl}(x) - x| \leq U|x| \text{ or } \text{fl}(x) = x(1 + u), |u| \leq U$$

where U is *machine epsilon* or *machine unit*. U is machine dependent but very small. In R, $U = 2^{-52} = 2.220 \times 10^{-16}$.

Other machine dependent constants include:

1. The minimum and maximum positive floating point numbers: $x_{\min} = 2^{-1022} = 2.225 \times 10^{-308}$ and $x_{\max} = 2^{1024} - 1 = 1.798 \times 10^{308}$.
2. The maximum integer: $2147383647 = 2^{31} - 1$.

1.1.3 Overflow and Underflow Error

Definition 1.2. If the result of a floating point operation exceeds x_{\max} , then the value returned is `Inf`.

Note. `Inf` indicates an *overflow error*.

Definition 1.3. If the result of a floating point operation is undefined then `NaN` is returned.

Definition 1.4. An *underflow error* occurs when the result of a floating point calculation is smaller (in absolute value) than x_{\min} .

Note. There are two possible outcomes (for underflow error): an error is reported or an exact 0 is returned. The latter outcome may cause problems in subsequent computations (e.g., division by 0).

Note. There are some ways to avoid overflow and underflow errors:

1. Use logarithmic scale: Changes multiplication/division into addition/subtraction, e.g., `lgamma`, `lfactorial`, `lchoose`.
2. Use series expansions (e.g., Taylor series).

Example 1.1. For x close to 0, $\frac{\exp(x) - 1}{x} \approx 1$. Naive computation of $\frac{\exp(x) - 1}{x}$ is problematic for x close to 0 due to possible round-off and underflow errors:

$$\frac{\text{fl}(\exp(x) - 1)}{\text{fl}(x)} \neq \frac{\exp(x) - 1}{x}$$

We solve the problem by using a series approximation, for $|x| \leq \varepsilon$,

$$\frac{\exp(x) - 1}{x} = \frac{x + x^2/2 + x^3/6 + \cdots}{x} = 1 + \frac{x}{2} + \frac{x^2}{6} + \cdots$$

1.1.4 Catastrophic Cancellation

Suppose $z_1 = g_1(x_1, \dots, x_n)$ and $z_2 = g_2(x_1, \dots, x_n)$. We want to compute $y = z_1 - z_2$. What we actually compute is

$$y^* = \text{fl}(\text{fl}(z_1) - \text{fl}(z_2))$$

where $\text{fl}(z_1) = z_1(1 + u_1)$ and $\text{fl}(z_2) = z_2(1 + u_2)$. We have

$$\text{fl}(z_1) - \text{fl}(z_2) = \underbrace{z_1 - z_2}_y + \underbrace{z_1 u_1 - z_2 u_2}_{\text{error}}$$

If z_1 and z_2 are large but $y = z_1 - z_2$ is small then the magnitude of the error may be larger than the magnitude of y - ***catastrophic cancellation***.