# Statistical Computation

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### 1 Basics

## 1.1 Floating Point Representation

**Definition 1.1.** A *floating point number* is represented by three components: (S, F, E) where S is the sign of the number  $(\pm 1)$ , F is a fraction (lying between 0 and 1), E is an exponent. S, F, E are all represented as binary digits (bits). The *floating point representation* of x, fl(x) is

$$fl(x) = S \times F \times 2^E$$

**Note.** x and f(x) need not be the same, since f(x) is a binary approximation to x.

#### 1.1.1 Round-Off Error

Mathematical operations introduce further approximation errors

$$f(f(x)) = f(x + \varepsilon) \approx f(x) + \varepsilon f'(x)$$

and the goal is to make the round-off error |f(x) - f(f(f(x)))| as small as possible.

#### 1.1.2 Machine Epsilon

For a given real number x, we have

$$|f(x) - x| \le U|x|$$
 or  $f(x) = x(1+u), |u| \le U$ 

where U is **machine epsilon** or **machine unit**.

#### 1.1.3 Overflow and Underflow Error

**Definition 1.2.** If the result of a floating point operation exceeds the maximum possible floating point number  $x_{\text{max}}$ , then the value returned is Inf.

Note. Inf indicates an overflow error.

**Definition 1.3.** If the result of a floating point operation is undefined then NaN is returned.

**Definition 1.4.** An *underflow error* occurs when the result of a floating point calculation is smaller (in absolute value) than the smallest floating point number  $x_{\min}$ .

**Note.** There are two possible outcomes: an error is reported or an exact 0 is returned. The latter outcome may cause problems in subsequent computations.

Note. There are some ways to avoid overflow and underflow errors:

- 1. Use logarithmic scale: Changes multiplication/division into addition/subtraction, e.g., lgamma, lfactorial, lchoose.
  - 2. Use series expansions (e.g., Taylor series).

**Example 1.1.** For x close to 0,  $\frac{\exp(x)-1}{x} \approx 1$ . Naive computation of  $\frac{\exp(x)-1}{x}$  is problematic for x close to 0 due to possible round-off and underflow errors:

$$\frac{\mathrm{fl}(\exp(x) - 1)}{\mathrm{fl}(x)} \neq \frac{\exp(x) - 1}{x}$$

We solve the problem by using a series approximation, for  $|x| \leq \varepsilon$ ,

$$\frac{\exp(x) - 1}{x} = \frac{x + x^2/2 + x^3/6 + \dots}{x} = 1 + \frac{x}{2} + \frac{x^2}{6} + \dots$$

## 1.1.4 Catastrophic Cancellation

Suppose  $z_1 = g_1(x_1, \dots, x_n)$  and  $z_2 = g_2(x_1, \dots, x_n)$ . We want to compute  $y = z_1 - z_2$ . What we actually compute is

$$y^* = \mathrm{fl}(\mathrm{fl}(z_1) - \mathrm{fl}(z_2))$$

where  $f(z_1) = z_1(1 + u_1)$  and  $f(z_2) = z_2(1 + u_2)$ . We have

$$fl(z_1) - fl(z_2) = \underbrace{z_1 - z_2}_{y} + \underbrace{z_1 u_1 - z_2 u_2}_{error}$$

If  $z_1$  and  $z_2$  are large but  $y = z_1 - z_2$  is small then the magnitude of the error may be larger than the magnitude of y - **catastrophic cancellation**.