

X
M
R
Z
Y

Name Dipesh Sah.

Standard 5 Section 'B' Roll No. 092

Subject Artificial Intelligence Lab.

Lab - 1

1) Vacuum Cleaner.

Algorithm :

1. Initialize goal state :

- Set goal_state = {'A': '0', 'B': '0'}
- (0 = clean, 1 = dirty)
- Initialize cost = 0.

2. Input the vacuum location :

- Get vacuum's location :
location_input = 'A' or 'B'.
- Get status of the current room :
status_input = '0' or '1'.
- Get status of the other room :
status_input_complement = '0' or '1'.

3. Vacuum in Room A :

- If location_input == 'A':
 - If status_input == '1' (A is dirty):
 - Clean A, set goal_state['A'] = '0', increment cost.
 - If status_input_complement == '1'
(B is dirty):
 - Move to B, clean it.
 - update goal_state['B'] = '0'
 - increment cost.

4. Vacuum in Room B :

- If location_input == 'B':
 - If status_input == '1' (B is dirty):

- Clean B, set goal-state [‘B’] = ‘0’,
- increment cost.
- If status-input-complement == ‘1’
(A is dirty):
 - Move to A, clean it
 - update goal-state [‘A’] = ‘0’
 - increment cost.

5. Output:

- print goal-state and cost as the performance measure.

Output: Initial condition: {‘A’: ‘1’, ‘B’: ‘1’}

Vacuum is placed in location A.

Location A is dirty.

Cost for Cleaning A: 1

Location A has been cleaned.

Location B is Dirty.

Moving right to location B.

Cost for moving Right: 2

Cost for Cleaning: 3

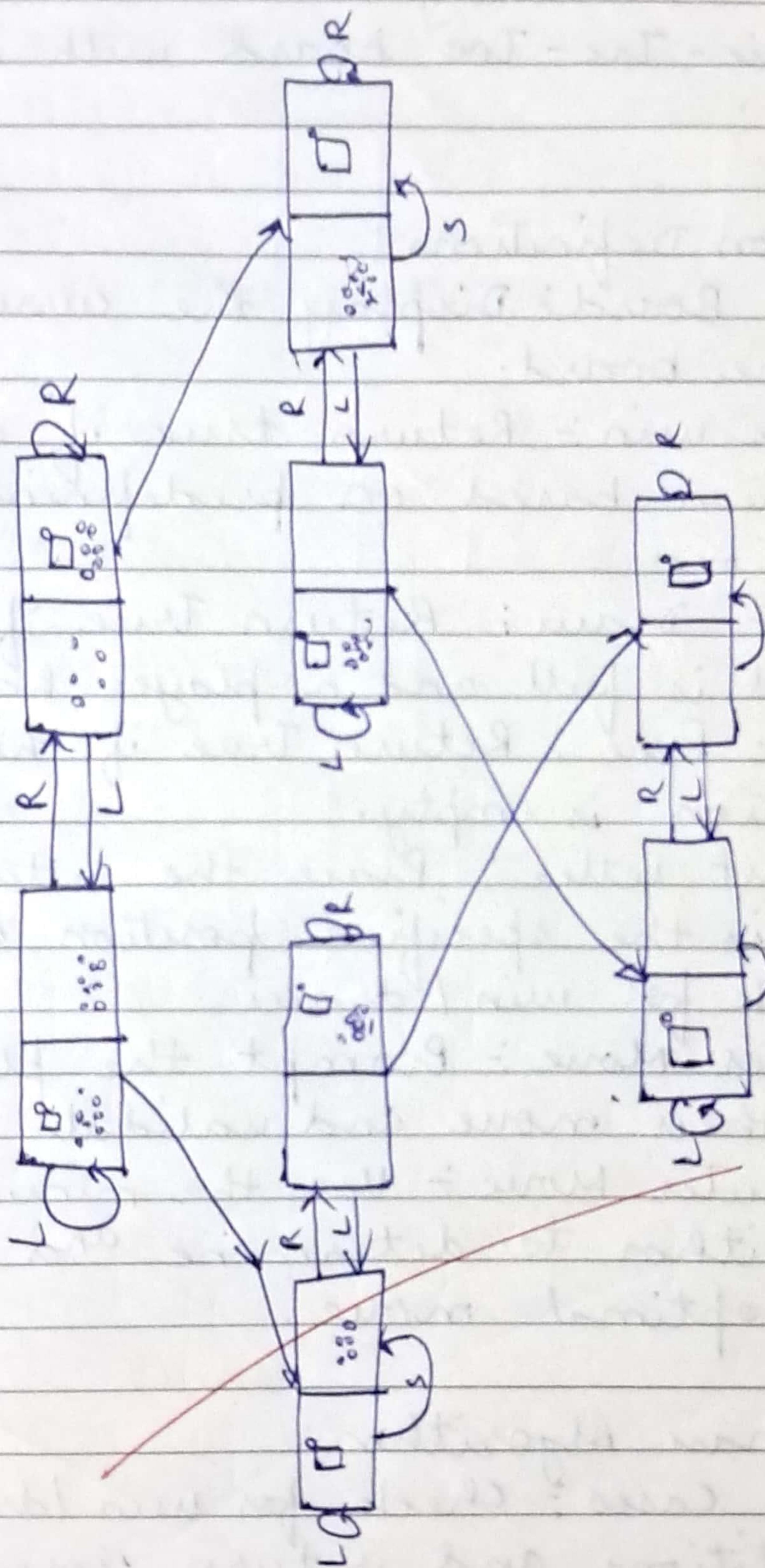
Location B has been cleaned.

GOAL State:

{‘A’: ‘0’, ‘B’: ‘0’}

Performance Measurement: 3

State space Tree.



2) Tic-Tac-Toe:

Algorithm :

1. Initialize the board.

Create a dictionary (or list) to represent the Tic-Tac-Toe board with empty spaces.

2. Function Definitions :

- Print Board : Display the current state of the board.
- Check Win : Return true if any player has won based on predefining winning conditions.
- Check Draw : Return True if the board is full and no player has won.
- Space Free : Return True if the selected position is empty.
- Insert Letter : Place the letter ('X' or 'O') in the specified position and check for win / draw.
- Player Move : Prompt the player for their move and validate input.
- Computer Move : Use the minimax algorithm to determine and make the optimal move.

3. Minimax Algorithm :

- Base Cases : Check for win / draw conditions and return scores based on the outcome.
- Recursion : For maximizing (computer) and minimizing (player), evaluate

possible moves and return the best score.

4. Main Game Loop.

- While the game is ongoing:
 - If no win/draw, call computerMove.
 - check for win/draw conditions.
 - If no win/draw, call Player move.
 - check for win/draw conditions.

5. End Game :

- Display the final board and announce the result (win/draw).

Pseudocode :-

Initialize board

While true:

If not checkWin() & not checkDraw();
compMove()

If checkWin() or checkDraw();
break;

If not checkWin() & not checkDraw();
playerMove()

If checkWin() or checkDraw();
break

Display final board.

Announce result.

Output: X |

Enter position for O (1-9) : 3

X	O	X	O
		X	

Enter position for O (1-9) : 7.

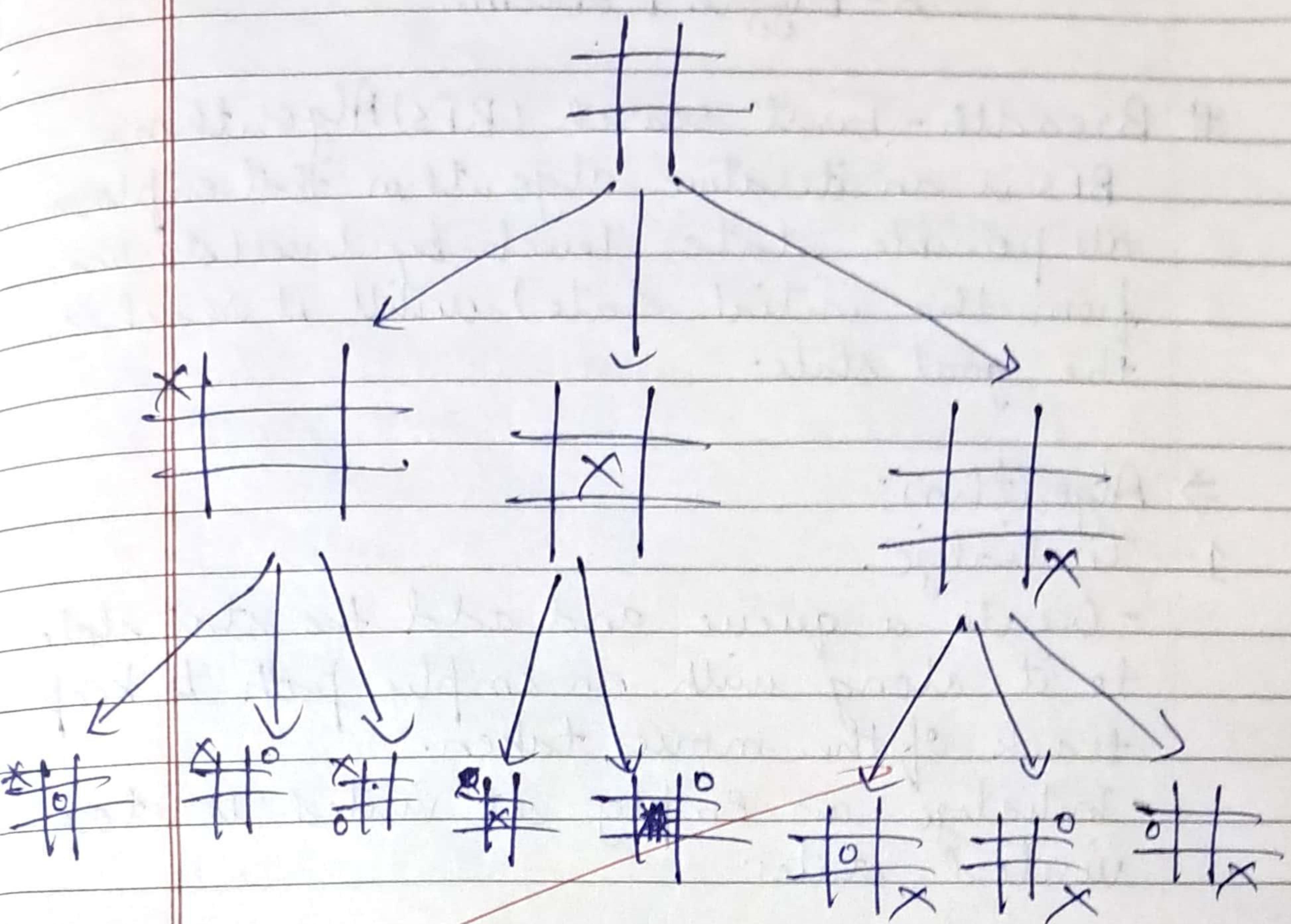
X	O	X	O
X		X	X
O		O	

Enter position for O (1-9) : 6

X	O	X	O
X	X	X	O
O		O	X

Bot wins!

- State Space Tree:



Q
11/10/24

Lab - 2

8 - Puzzle Problem.

Breadth - First Search (BFS) Algorithm.

BFS is an iterative algorithm that explores all possible states level by level (starting from the initial state) until it reaches the goal state.

⇒ Algorithm:-

1. Initialize.

- Create a queue and add the start state to it, along with an empty path to keep track of the moves taken.
- Initialize an empty set 'visited' to store visited states.

2. While queue is not empty:

- dequeue the front element from the queue (current state and the path).

3. Check goal.

- If the current state is the goal state, return the path (solution).

4. Mark current state as visited.

- add the current state to 'visited set'.

5. Generate successors:

- For each valid move from the current state:
 - apply the move to get a new state.
 - If the new state has not been visited:

- add the new state and the updated path to the queue.

6. Return Failure:

- If the queue is empty and no solution has been found, return failure (no solution).

→ Output:

Initial state:

$$\begin{bmatrix} 1, & 2, & 3 \\ 4, & 0, & 6 \\ 7, & 5, & 8 \end{bmatrix}$$

Solving using BFS:

Exploring state in BFS:

$$[1, 2, 3]$$

$$[4, 0, 6]$$

$$[7, 5, 8],$$

Exploring state in BFS:

$$[1, 0, 3]$$

$$[4, 2, 6]$$

$$[7, 5, 8]$$

↓ ↓ ↓

• BFS solution: [‘down’, ‘right’]

Exploring state in DFS:

$$[1, 2, 3]$$

$$[4, 5, 6]$$

$$[7, 8, 0]$$

8-puzzle problem using Depth-First Search Algorithm
DFS explores a branch as far as possible before backtracking when it reaches a dead end (or goal).

⇒ Algorithm:

1. Initialize:

- Create a stack and add the start state to it, along with an empty path to keep track of the moves taken.
- Initialize an empty set 'visited' to store visited states.

2. While stack is not empty:

- pop the top element from the stack (the current state and the path).

3. Check goal:

- If the current state is the goal, return the path (solution).

4. Mark current state as visited:

- add the current state to the 'visited set'.

5. Generate successors:

- For each valid move from the current state:
 - apply the move to get a new state.
 - If the new state has not been visited:
 - push the new state and the updated path to the stack.

6. Return Failure:

- If the stack is empty and no solution has been found, return failure (no solution exists).

⇒ Output:

Solving using DFS:

Exploring state in DFS:

[1, 0, 3]

[4, 2, 6]

[7, 5, 8]

: : :

Exploring state in DFS:

[2, 3, 6]

[8, 0, 5]

[4, 1, 7]

: : :

: : :

: : :

: : :

,

Exploring state in DFS:

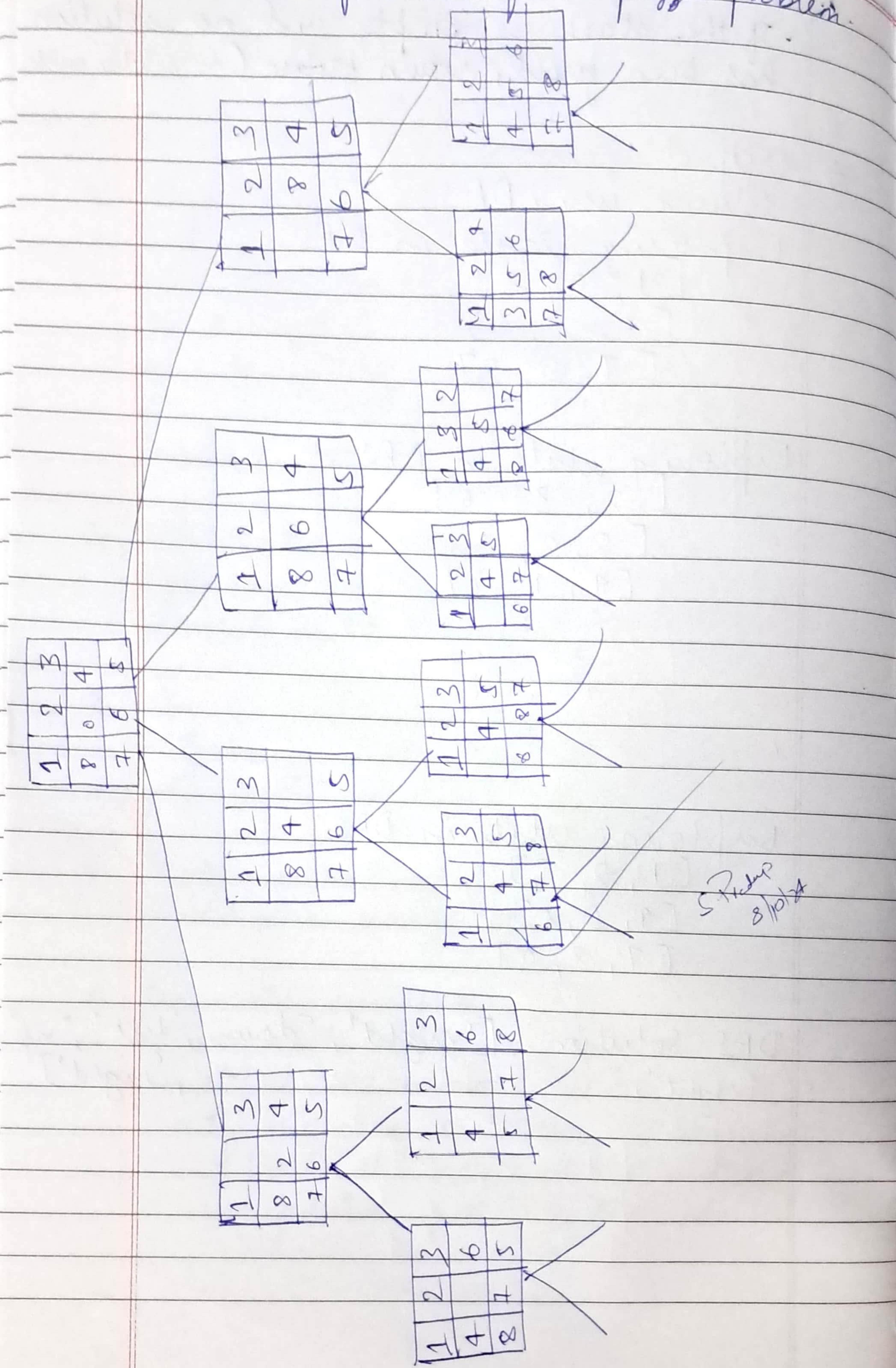
[1, 2, 3]

[4, 5, 6]

[7, 8, 0]

DFS Solution: ['right', 'down', 'left', 'up',
'right', ..., ..., ..., ..., 'right']

State Space Tree for 8-puzzle problem



Lab-3

Algorithm 1: A* search for 8-puzzle with Misplaced Tiles Heuristic.

1. Initialize

Open list $\leftarrow \{\text{start node}\}$

Closed list $\leftarrow \emptyset$

$g(\text{start}) \leftarrow 0$

$h(\text{start}) \leftarrow \text{number of misplaced tiles}$

$f(\text{start}) \leftarrow g(\text{start}) + h(\text{start})$

2. Loop until Open list is empty

- Select the node n from the Open list with the lowest $f(n)$.

- Remove n from the Open list and add it to the Closed list.

3. For each neighbour of node n .

- If the neighbour is in the Closed list, skip it.

- Calculate $g(\text{neighbour}) \leftarrow g(n) + 1$ (depth of node)

- Calculate $h(\text{neighbour}) \leftarrow \text{number of misplaced tiles}$

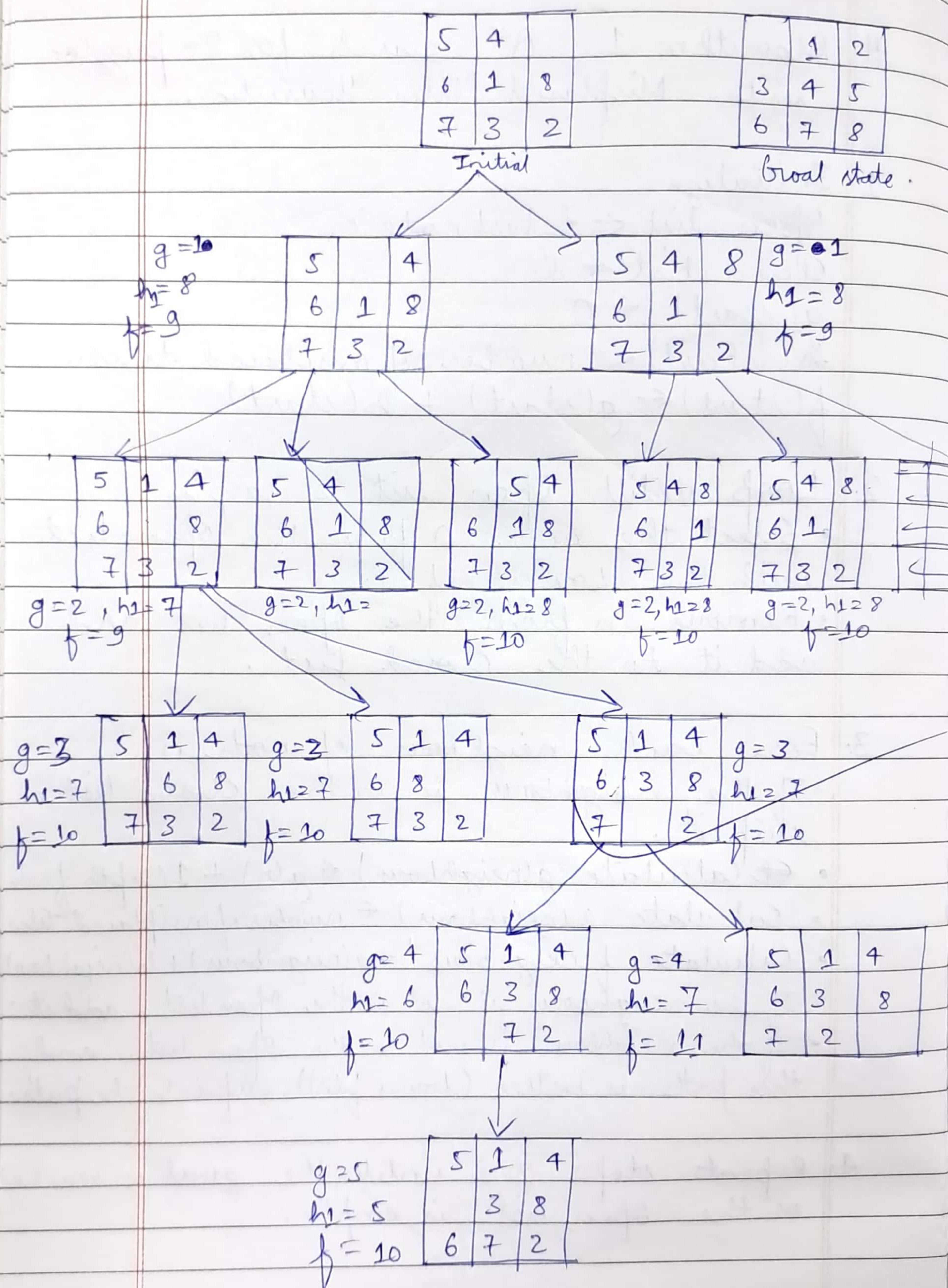
- Calculate $f(\text{neighbour}) \leftarrow g(\text{neighbour}) + h(\text{neighbour})$

- If the neighbour is not in the Open list, add it.

- If the neighbour is not in the Open list, and this path is better (lower $f(n)$), update the path.

4. Repeat steps 2-3 until the goal is reached or the Open list is empty.

State Space Tree



Algorithm 2: A* Search for 8-puzzle with Manhattan Distance Heuristic.

1. Initialize

open list $\leftarrow \{\text{start node}\}$

closed list $\leftarrow \emptyset$

$g(\text{start}) \leftarrow 0$

$h(\text{start}) \leftarrow \text{Manhattan distance}$

$f(\text{start}) \leftarrow g(\text{start}) + h(\text{start})$

2. Loop until open list is empty.

- Select the node n from the open list with the lowest $f(n)$
- If n is the goal state/node, terminate and return the path.
- Remove n from the open list and add it to the closed list.

3. For each neighbour of node n ,

- If the neighbour is in the closed list, skip it.
- Calculate $g(\text{neighbour}) \leftarrow g(n) + 1$ (Depth of node)
- Calculate $h(\text{neighbour}) \leftarrow \text{Manhattan Distance}$.
- Calculate $f(\text{neighbour}) \leftarrow g(\text{neighbour}) + h(\text{neighbour})$
- If the neighbour is not in the open list, add it.
- If the neighbour is in the open list and this path is better (lower ($f(n)$)), update the path.

4. Repeat steps 2-3 until the goal is reached or the open list is empty.

State Space Tree

Initial

$g=0$	2	8	3		1	2	3
$h_2=4$	1		4		8		4
$f=4+0=4$	7	6	5		7	6	5

$g=1$	2	8	3	$g=1$	2	3	$g(n)=1$	2	8	3
$h(n)=1+1+2+1$ =5	1	4		$h(n)=3$	1	8	$h(n)=5$	1	4	
$f=6$	7	6	5	$f(n)=4$	7	6	$f(n)=6$	7	6	5

$g(n)=2$		2	3	$g(n)=2$	2	3	
$h(n)=2+1=2$	1	8	4	$h(n)=1+1+2+1$	1	8	4
$f(n)=4$	7	6	5	$f(n)=2+4=6$	7	6	5

1	2	3	$g(n)=3$
	8	4	$h(n)=1$
7	6	5	$f(n)=3+1=4$

1	2	3		1	2	3	
8		4		7	8	4	
7	6	5		6	5		

S.Ridhi
15/10/2024

Final state
goal reached.