OREGON STATE UNIVERSITY

Project Group 21

CS 325: Project 3

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1 Problem 1

Part A:

i. Linear Problem Formulation

Let each route be represented as variables of the form "SourceDestination." For example a route from Plant 1 to Warehouse 1 would be "p1w1."

Objective function:

Minimize:

```
10~p1w1+15~p1w2+11~p2w1+8~p2w2+13~p3w1+8~p3w2+9~p3w3+14~p4w2+8~p4w3+5~w1r1+6~w1r2+7~w1r3+10~w1r4+12~w2r3+8~w2r4+10~w2r5+14~w2r6+14~w3r4+12~w3r5+12~w3r6+6~w3r7
```

Subject To:

```
Supply Constraints:
```

```
\begin{array}{l} p1w1 + p1w2 = 150 \\ p2w1 + p2w2 = 450 \\ p3w1 + p3w2 + p3w3 = 250 \\ p4w2 + p4w3 = 150 \end{array}
```

```
Demand Constraints:
```

```
w1r1 = 100

w1r2 = 150

w1r3 + w2r3 = 100

w1r4 + w2r4 + w3r4 = 200

w2r5 + w3r5 = 200

w2r6 + w3r6 = 150

w3r7 = 100
```

Balancing Constraints:

```
\begin{array}{l} p1w1+p2w1+p3w1-w1r1-w1r2-w1r3-w1r4<=0\\ p1w2+p2w2+p3w2+p4w2-w2r3-w2r4-w2r5-w2r6<=0\\ p3w3+p4w3-w3r4-w3r5-w3r6-w3r7<=0 \end{array}
```

Non-negativity Constraints:

p1w1, p1w2, p2w1, p2w2, p3w1, p3w2, p3w3, p4w2, p4w3, w1r1, w1r2, w1r3, w1r4, w2r3, w2r4, w2r5, w2r6, w3r4, w3r5, w3r6, w3r7 >= 0

ii. LINDO Code and Output

```
MIN 10 p1w1 + 15 p1w2 + 11 p2w1 + 8 p2w2 + 13 p3w1 + 8 p3w2 + 9 p3w3
 + 14 p4w2 + 8 p4w3 + 5 w1r1 + 6 w1r2 + 7 w1r3 + 10 w1r4 + 12 w2r3
  + 8  w2r4 + 10 w2r5 + 14 w2r6 + 14 w3r4 + 12 w3r5 + 12 w3r6 + 6 w3r7
   p1w1 + p1w2 = 150
   p2w1 + p2w2 = 450
   p3w1 + p3w2 + p3w3 = 250
   p4w2 + p4w3 = 150
   w1r1 = 100
   w1r2 = 150
   w1r3 + w2r3 = 100
   w1r4 + w2r4 + w3r4 = 200
   w2r5 + w3r5 = 200
   w2r6 + w3r6 = 150
   w3r7 = 100
   p1w1 + p2w1 + p3w1 - w1r1 - w1r2 - w1r3 - w1r4 <= 0
   p1w2 + p2w2 + p3w2 + p4w2 - w2r3 - w2r4 - w2r5 - w2r6 \le 0
   p3w3 + p4w3 - w3r4 - w3r5 - w3r6 - w3r7 <= 0
   p1w1 > 0
   p1w2 > 0
   p2w1 > 0
   p2w2 > 0
   p3w1 > 0
   p3w2 > 0
```

p3w3 > 0p4w2 > 0p4w3 > 0w1r1 > 0w1r2 > 0w1r3 > 0w1r4 > 0w2r3 > 0w2r4 > 0w2r5 > 0w2r6 > 0w3r4 > 0w3r5 > 0w3r6 > 0w3r7 > 0END

LP OPTIMUM FOUND AT STEP 13

OBJECTIVE FUNCTION VALUE

1) 17100.00

VARIABLE	VALUE	REDUCED COST
P1W1	150.000000	0.000000
P1W2	0.000000	8.000000
P2W1	200.000000	0.000000
P2W2	250.000000	0.000000
P3W1	0.000000	2.000000
P3W2	150.000000	0.000000
P3W3	100.000000	0.000000
P4W2	0.000000	7.000000
P4W3	150.000000	0.000000
W1R1	100.000000	0.000000
W1R2	150.000000	0.000000
W1R3	100.000000	0.000000
W1R4	0.000000	5.000000
W2R3	0.000000	2.000000
W2R4	200.000000	0.000000
W2R5	200.000000	0.000000
W2R6	0.000000	1.000000
W3R4	0.000000	7.000000
W3R5	0.000000	3.000000

W3R6	150.000000	0.000000
W3R7	100.000000	0.000000
ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.000000	-10.000000
3)	0.000000	-11.000000
4)	0.000000	-11.000000
5)	0.000000	-10.000000
6)	0.000000	-5.000000
7)	0.000000	-6.000000
8)	0.000000	-7.000000
9)	0.00000	-5.000000
10)	0.00000	-7.000000
11)	0.00000	-10.000000
12)	0.00000	-4.000000
13)	0.00000	0.000000
14)	0.00000	3.000000
15)	0.00000	2.000000
16)	150.000000	0.000000
17)	0.00000	0.000000
18)	200.000000	0.000000
19)	250.000000	0.000000
20)	0.00000	0.000000
21)	150.000000	0.000000
22)	100.000000	0.000000
23)	0.00000	0.000000
24)	150.000000	0.000000
25)	100.00000	0.000000
26)	150.000000	0.000000
27)	100.00000	0.000000
28)	0.00000	0.000000
29)	0.00000	0.000000
30)	200.000000	0.000000
31)	200.000000	0.000000
32)	0.000000	0.000000
33)	0.000000	0.000000
34)	0.00000	0.000000
35)	150.000000	0.000000
36)	100.000000	0.000000

iii. Optimal Shipping Routes and Minimum Cost

The optimal shipping routes and quantity of refrigerators per route are:

Route	Refrigerators
P1W1	150
P2W1	200
P2W2	250
P3W2	150
P3W3	100
P4W3	150
W1R1	100
W1R2	150
W1R3	100
W2R4	200
W2R5	200
W3R6	150
W3R7	100

The optimal minimum cost is \$17,100

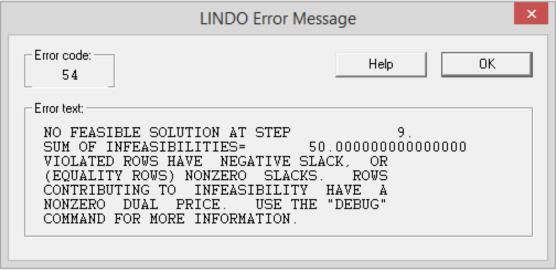
Part B:

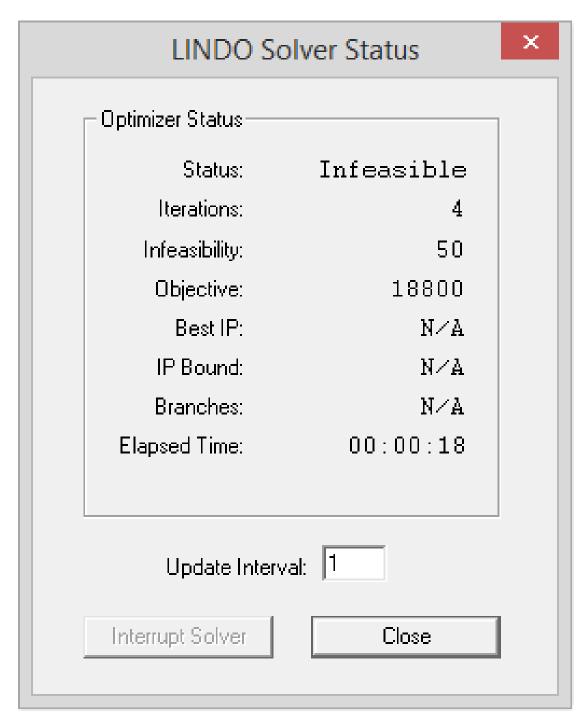
To remove Warehouse 2 and all its associated routes, we remove all variables in our equation that contain "w2" in it. The resulting LINDO code is as follows:

```
MIN 10 plw1 + 11 p2w1 + 13 p3w1 + 9 p3w3 + 8 p4w3 + 5 w1r1 + 6 w1r2
   + 7 w1r3 + 10 w1r4 + 14 w3r4 + 12 w3r5 + 12 w3r6 + 6 w3r7
ST
   p1w1 = 150
   p2w1 = 450
   p3w1 + p3w3 = 250
   p4w3 = 150
   w1r1 = 100
   w1r2 = 150
   w1r3 = 100
   w1r4 + w3r4 = 200
   w3r5 = 200
   w3r6 = 150
   w3r7 = 100
   p1w1 + p2w1 + p3w1 - w1r1 - w1r2 - w1r3 - w1r4 <= 0
   p3w3 + p4w3 - w3r4 - w3r5 - w3r6 - w3r7 <= 0
   p1w1 > 0
   p2w1 > 0
   p3w1 > 0
```

```
p3w3 > 0
p4w3 > 0
w1r1 > 0
w1r2 > 0
w1r3 > 0
w1r4 > 0
w3r4 > 0
w3r5 > 0
w3r6 > 0
w3r7 > 0
```

Upon execution of the code, we are met with these solution infeasibility notices:





There is no feasible solution here because Warehouse 3's total possible supply (400) is not capable of supporting the combined demand of R4, R5, R6, and R7 (650). On the other hand, Warehouse 1 is in fact capable of supporting R1, R2, R3, and R4 with its 850 supply to 550 demand ratio, but that still does not solve the larger supply chain deficit and feasibility issue. Overall, we need Warehouse 2 to supplement the supply to

R3, R4, R5, and R6 and sufficiently "spread the load" on the warehouses.

Part C:

To simulate the limitation of 100 shipments (inbound and outbound) on Warehouse 2, we add two new constraints to our code for Part A. The constraints limit the inbound and outbound shipments to being less than or equal to 100.

LINDO Code and Output:

```
MIN 10 plw1 + 15 plw2 + 11 p2w1 + 8 p2w2 + 13 p3w1 + 8 p3w2 + 9 p3w3
  + 14 p4w2 + 8 p4w3 + 5 w1r1 + 6 w1r2 + 7 w1r3 + 10 w1r4 + 12 w2r3
  + 8  w2r4 + 10  w2r5 + 14  w2r6 + 14  w3r4 + 12  w3r5 + 12  w3r6 + 6  w3r7
ST
   p1w1 + p1w2 = 150
   p2w1 + p2w2 = 450
   p3w1 + p3w2 + p3w3 = 250
   p4w2 + p4w3 = 150
   w1r1 = 100
   w1r2 = 150
   w1r3 + w2r3 = 100
   w1r4 + w2r4 + w3r4 = 200
   w2r5 + w3r5 = 200
   w2r6 + w3r6 = 150
   w3r7 = 100
   p1w1 + p2w1 + p3w1 - w1r1 - w1r2 - w1r3 - w1r4 <= 0
   p1w2 + p2w2 + p3w2 + p4w2 - w2r3 - w2r4 - w2r5 - w2r6 <= 0
   p3w3 + p4w3 - w3r4 - w3r5 - w3r6 - w3r7 <= 0
   p1w1 > 0
   p1w2 > 0
   p2w1 > 0
   p2w2 > 0
   p3w1 > 0
   p3w2 > 0
   p3w3 > 0
   p4w2 > 0
   p4w3 > 0
   w1r1 > 0
   w1r2 > 0
   w1r3 > 0
   w1r4 > 0
   w2r3 > 0
   w2r4 > 0
   w2r5 > 0
```

```
w2r6 > 0
w3r4 > 0
w3r5 > 0
w3r6 > 0
w3r7 > 0
p1w2 + p2w2 + p3w2 + p4w2 <= 100
w2r3 + w2r4 + w2r5 + w2r6 <= 100
END</pre>
```

LP OPTIMUM FOUND AT STEP 7

OBJECTIVE FUNCTION VALUE

1) 18300.00

4)

VARIABLE	VALUE	REDUCED COST
P1W1	150.00000	0.000000
P1W2	0.00000	8.000000
P2W1	350.00000	0.000000
P2W2	100.00000	0.000000
P3W1	0.00000	4.000000
P3W2	0.00000	2.000000
P3W3	250.000000	0.000000
P4W2	0.00000	9.000000
P4W3	150.000000	0.000000
W1R1	100.00000	0.000000
W1R2	150.00000	0.000000
W1R3	100.00000	0.000000
W1R4	150.000000	0.000000
W2R3	0.00000	7.000000
W2R4	50.00000	0.000000
W2R5	50.00000	0.000000
W2R6	0.00000	4.000000
W3R4	0.00000	4.000000
W3R5	150.00000	0.000000
W3R6	150.000000	0.000000
W3R7	100.00000	0.000000
ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.00000	-20.000000
3)	0.00000	-21.000000

0.000000 -19.000000

5)	0.000000	-18.000000
6)	0.000000	5.000000
7)	0.000000	4.000000
8)	0.000000	3.000000
9)	0.000000	0.000000
10)	0.000000	-2.000000
11)	0.000000	-2.000000
12)	0.000000	4.000000
13)	0.000000	10.000000
14)	0.000000	13.000000
15)	0.000000	10.000000
16)	150.000000	0.000000
17)	0.000000	0.000000
18)	350.000000	0.000000
19)	100.000000	0.000000
20)	0.000000	0.000000
21)	0.000000	0.000000
22)	250.000000	0.000000
23)	0.000000	0.000000
24)	150.000000	0.000000
25)	100.000000	0.000000
26)	150.000000	0.000000
27)	100.000000	0.000000
28)	150.000000	0.000000
29)	0.000000	0.000000
30)	50.000000	0.000000
31)	50.000000	0.00000
32)	0.000000	0.000000
33)	0.000000	0.000000
34)	150.000000	0.000000
35)	150.000000	0.000000
36)	100.000000	0.00000
37)	0.000000	0.00000
38)	0.000000	5.00000
•		

We see that this is a feasible solution, though it raises our optimal cost from \$17,100 to \$18,300.

The optimal shipping routes and quantity of refrigerators per route are:

7

Route	Refrigerators
P1W1	150
P2W1	350
P2W2	100

P3W3	250
P4W3	150
W1R1	100
W1R2	150
W1R3	100
W1R4	150
W2R4	50
W2R5	50
W3R5	150
W3R6	150
W3R7	100

Part D:

For some cost c_{ab} and some route x_{ab} where their subscripts represent the route from a to b, the objective function is:

$$Min(\sum_{p}\sum_{w}c_{pw}x_{pw}+\sum_{w}\sum_{r}c_{wr}x_{wr})$$

s.t.

```
(Supply Constraint) For each source node p, where s is supply: \sum_w x_{pw} = s_p (Demand Constraint) For each destination node r, where d is demand: \sum_w x_{wr} = d_r (Balancing Constraint) For each intermediate node w: \sum_p x_{pw} - \sum_r x_{wr} <= 0 (Non-negativity Constraint) For all nodes p and r: x_{pw}, x_{wr} >= 0
```

2 Problem 2

Part A:

i. Linear Problem Formulation

Let each ingredient be represented as variables using the first two letters of the ingredient's name. For example spinach is "sp."

Objective function:

Minimize:

$$21 \text{ to} + 16 \text{ le} + 40 \text{ sp} + 41 \text{ ca} + 41 \text{ su} + 120 \text{ sm} + 164 \text{ ch} + 884 \text{ oi}$$

Subject To:

Nutrition Constraints:

```
0.85 \text{ to} + 1.62 \text{ le} + 2.86 \text{ sp} + 0.93 \text{ ca} + 23.4 \text{ su} + 16 \text{ sm} + 9 \text{ ch} >= 15 0.33 \text{ to} + 0.20 \text{ le} + 0.39 \text{ sp} + 0.24 \text{ ca} + 48.7 \text{ su} + 5 \text{ sm} + 2.6 \text{ ch} + 100 \text{ oi} >= 2
```

```
0.33\ to+0.20\ le+0.39\ sp+0.24\ ca+48.7\ su+5\ sm+2.6\ ch+100\ oi<=8 4.64\ to+2.37\ le+3.63\ sp+9.58\ ca+15\ su+3\ sm+27\ ch>=4 9\ to+28\ le+65\ sp+69\ ca+3.8\ su+120\ sm+78\ ch<=200 le+sp-.4le-.4sp-.4to-.4ca-.4su-.4sm-.4ch-.4oi>=0
```

Non-negativity Constraints:

le, sp, to, ca, su, sm, ch, oi >= 0

ii. LINDO Code and Output

```
MIN 21 to + 16 le + 40 sp + 41 ca + 41 su + 120 sm + 164 ch + 884 oi
ST
     0.85 \text{ to} + 1.62 \text{ le} + 2.86 \text{ sp} + 0.93 \text{ ca} + 23.4 \text{ su} + 16 \text{ sm} + 9 \text{ ch} >= 15
     0.33 \text{ to} + 0.20 \text{ le} + 0.39 \text{ sp} + 0.24 \text{ ca} + 48.7 \text{ su} + 5 \text{ sm} + 2.6 \text{ ch} + 100 \text{ oi} > 0.00 \text{ sp}
     0.33 \text{ to} + 0.20 \text{ le} + 0.39 \text{ sp} + 0.24 \text{ ca} + 48.7 \text{ su} + 5 \text{ sm} + 2.6 \text{ ch} + 100 \text{ oi} < 0.00 \text{ ch}
     4.64 \text{ to} + 2.37 \text{ le} + 3.63 \text{ sp} + 9.58 \text{ ca} + 15 \text{ su} + 3 \text{ sm} + 27 \text{ ch} >= 4
     9 to + 28 le + 65 sp + 69 ca + 3.8 su + 120 sm + 78 ch <= 200
     le + sp - .4le - .4sp - .4to - .4ca - .4su - .4sm - .4ch - .4oi >= 0
    le > 0
     sp > 0
    to > 0
     ca > 0
    su > 0
     sm > 0
     ch > 0
     oi > 0
END
```

LP OPTIMUM FOUND AT STEP

OBJECTIVE FUNCTION VALUE

1) 103.3371

VARIABLE	VALUE	REDUCED COST
TO	0.000000	16.520000
LE	0.561163	0.00000
SP	0.000000	13.916369
CA	0.000000	35.535095
SU	0.084189	0.00000
SM	0.757556	0.00000
СН	0.000000	96.583527
OI	0.000000	1212.744385

ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.00000	-8.632957
3)	6.00000	0.000000
4)	0.00000	3.269654
5)	0.865455	0.000000
6)	93.060745	0.000000
7)	0.00000	-4.447570
8)	0.561163	0.000000
9)	0.00000	0.000000
10)	0.00000	0.000000
11)	0.00000	0.000000
12)	0.084189	0.000000
13)	0.757556	0.000000
14)	0.00000	0.000000
15)	0.00000	0.000000

The minimum amount of calories a salad can contain while meeting the nutrition constraints is 103.3371 calories.

The optimal ingredient mix is:

Ingredient	Amount
LE	0.561163
SU	0.084189
SM	0.757556

iii.

The cost of this salad is: $(0.561163*.75) + (0.084189*.45) + (0.757556*2.15) = $2.0875 \approx 2.09 .

Part B:

i. Linear Problem Formulation

Let each ingredient be represented as variables using the first two letters of the ingredient's name. For example spinach is "sp."

Objective function:

Minimize:

```
1\ {\rm to}\ +\ .75\ {\rm le}\ +\ .5\ {\rm sp}\ +\ .5\ {\rm ca}\ +\ .45\ {\rm su}\ +\ 2.15\ {\rm sm}\ +\ .95\ {\rm ch}\ +\ 2\ {\rm oi}
```

Subject To:

```
Nutrition Constraints:
```

```
\begin{array}{l} 0.85\ {\rm to}+1.62\ {\rm le}+2.86\ {\rm sp}+0.93\ {\rm ca}+23.4\ {\rm su}+16\ {\rm sm}+9\ {\rm ch}>=15\\ 0.33\ {\rm to}+0.20\ {\rm le}+0.39\ {\rm sp}+0.24\ {\rm ca}+48.7\ {\rm su}+5\ {\rm sm}+2.6\ {\rm ch}+100\ {\rm oi}>=2\\ 0.33\ {\rm to}+0.20\ {\rm le}+0.39\ {\rm sp}+0.24\ {\rm ca}+48.7\ {\rm su}+5\ {\rm sm}+2.6\ {\rm ch}+100\ {\rm oi}<=8\\ 4.64\ {\rm to}+2.37\ {\rm le}+3.63\ {\rm sp}+9.58\ {\rm ca}+15\ {\rm su}+3\ {\rm sm}+27\ {\rm ch}>=4\\ 9\ {\rm to}+28\ {\rm le}+65\ {\rm sp}+69\ {\rm ca}+3.8\ {\rm su}+120\ {\rm sm}+78\ {\rm ch}<=200\\ {\rm le}+{\rm sp}-.4{\rm le}-.4{\rm sp}-.4{\rm to}-.4{\rm ca}-.4{\rm su}-.4{\rm sm}-.4{\rm ch}-.4{\rm ci}>=0 \end{array}
```

Non-negativity Constraints:

le, sp, to, ca, su, sm, ch, oi ≥ 0

ii. LINDO Code and Output

```
MIN 1 to + .75 le + .5 sp + .5 ca + .45 su + 2.15 sm + .95 ch + 2 oi
ST
    0.85 \text{ to} + 1.62 \text{ le} + 2.86 \text{ sp} + 0.93 \text{ ca} + 23.4 \text{ su} + 16 \text{ sm} + 9 \text{ ch} >= 15
    0.33 \text{ to} + 0.20 \text{ le} + 0.39 \text{ sp} + 0.24 \text{ ca} + 48.7 \text{ su} + 5 \text{ sm} + 2.6 \text{ ch} + 100 \text{ oi} > 0.00 \text{ sp}
    0.33 \text{ to} + 0.20 \text{ le} + 0.39 \text{ sp} + 0.24 \text{ ca} + 48.7 \text{ su} + 5 \text{ sm} + 2.6 \text{ ch} + 100 \text{ oi} < 0.00 \text{ sp}
    4.64 \text{ to} + 2.37 \text{ le} + 3.63 \text{ sp} + 9.58 \text{ ca} + 15 \text{ su} + 3 \text{ sm} + 27 \text{ ch} >= 4
    9 to + 28 le + 65 sp + 69 ca + 3.8 su + 120 sm + 78 ch <= 200
    le + sp - .4le - .4sp - .4to - .4ca - .4su - .4sm - .4ch - .4oi >= 0
    le > 0
    sp > 0
    to > 0
    ca > 0
    su > 0
    sm > 0
    ch > 0
    oi > 0
END
```

LP OPTIMUM FOUND AT STEP

OBJECTIVE FUNCTION VALUE

1) 1.554133

VARIABLE	VALUE	REDUCED COST
TO	0.000000	1.002081
LE	0.00000	0.402912

SP CA SU SM CH	0.832298 0.000000 0.096083 0.000000 1.152364 0.000000	0.000000 0.486914 0.000000 0.405609 0.000000 7.281258
ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.00000	-0.131261
3)	6.000000	0.000000
4)	0.000000	0.051847
5)	31.576324	0.000000
6)	55.651089	0.000000
7)	0.00000	-0.241358
8)	0.000000	0.000000
9)	0.832298	0.000000
10)	0.00000	0.000000
11)	0.00000	0.000000
12)	0.096083	0.000000
13)	0.00000	0.000000
14)	1.152364	0.000000
15)	0.00000	0.000000

The minimum cost of a salad while meeting the nutrition constraints is $\$1.554133 \approx \1.55

The optimal ingredient mix is:

Ingredient	Amount
SP	0.832298
SU	0.096083
СН	1.152364

iii.

The caloric content of this salad is: (0.832298*40) + (0.096083*585) + (1.152364*164) = 278.488 calories.

Part C:

i.

ii.

iii.

3 Problem 3

Part A:

i. General Problem Statement

Objective:

```
Min \sum_{i=1}^{n} b_i, for some b_i representing |y_i - (a_1x_i + a_0)| for each i = 1...n s.t.
For each i = 1...n, b_i - (y_i - (a_1x_i + a_0)) >= 0
For each i = 1...n, b_i + (y_i - (a_1x_i + a_0)) >= 0
```

ii. LINDO Code and Output

```
MIN b1 + b2 + b3 + b4 + b5 + b6 + b7
ST
   b1 + 1 a1 + a0 >= 5
   b1 - 1 a1 - a0 >= -5
   b2 + 1 a1 + a0 >= 3
   b2 - 1 a1 - a0 >= -3
   b3 + 2 a1 + a0 >= 13
   b3 - 2 a1 - a0 >= -13
   b4 + 3 a1 + a0 >= 8
   b4 - 3 a1 - a0 >= -8
   b5 + 4 a1 + a0 >= 10
   b5 - 4 a1 - a0 >= -10
   b6 + 5 a1 + a0 >= 14
   b6 - 5 a1 - a0 >= -14
   b7 + 6 a1 + a0 >= 18
   b7 - 6 a1 - a0 >= -18
   b1 > 0
   b2 > 0
   b3 > 0
   b4 > 0
   b5 > 0
   b6 > 0
   b7 > 0
END
```

LP OPTIMUM FOUND AT STEP

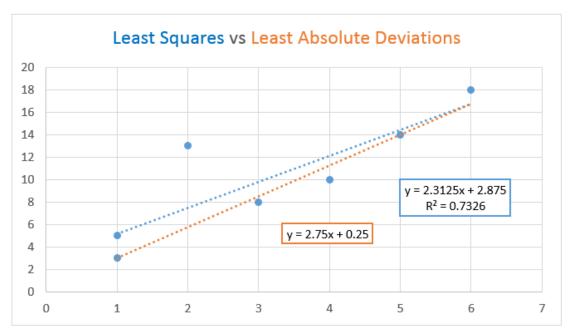
OBJECTIVE FUNCTION VALUE

1) 12.25000

VARIABLE	VALUE	REDUCED COST	
В1	2.00000	0.000000	
В2	0.00000	0.250000	
В3	7.250000	0.000000	
В4	0.500000	0.000000	
В5	1.250000	0.000000	
В6	0.00000	0.750000	
В7	1.250000	0.000000	
A1	2.750000	0.000000	
A0	0.250000	0.000000	
ROW	SLACK OR SURPLUS	DUAL PRICES	
2)	0.00000	-1.000000	
3)	4.00000	0.000000	
4)	0.00000	0.000000	
5)	0.00000	-0.750000	
6)	0.00000	-1.000000	
7)	14.500000	0.000000	
8)	1.000000	0.000000	
9)	0.00000	-1.000000	
10)	2.500000	0.000000	
11)	0.00000	-1.000000	
12)	0.00000	0.000000	
13)	0.00000	-0.250000	
14)	0.00000	-1.000000	
15)	2.500000	0.000000	
16)	2.000000	0.000000	
17)	0.00000	0.000000	
18)	7.250000	0.000000	
19)	0.500000	0.000000	
20)	1.250000	0.000000	
21)	0.000000	0.000000	
22)	1.250000	0.000000	

The sum of the absolute deviations was 12.25.

iii. Graph



The LAD approach creates a line that has a steeper slope and lesser y-intercept than the least squares line. The line more tightly follows the path of the majority of the data points and seems to be less affected by the outlier at x=2.

Part B:

i. General Problem Statement

Objective:

Min b, for some b representing $Max|y_i-(a_1x_i+a_0)|$ s.t. For each $i=1...n,\ b+a_1x_i+a_0>=y_i$ For each $i=1...n,\ b-a_1x_i-a_0>=-y_i$

ii. LINDO Code and Output

```
MIN b
ST

b + 1 a1 + a0 >= 5
b - 1 a1 - a0 >= -5
b + 1 a1 - a0 >= 3
b - 1 a1 - a0 >= -3
b + 2 a1 + a0 >= 13
b - 2 a1 - a0 >= -13
b + 3 a1 + a0 >= 8
b - 3 a1 - a0 >= -8
b + 4 a1 + a0 >= 10
```

LP OPTIMUM FOUND AT STEP

OBJECTIVE FUNCTION VALUE

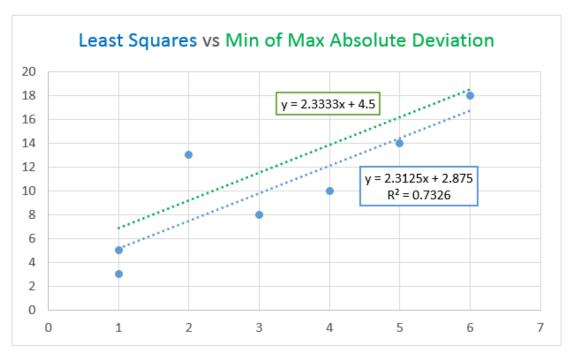
1) 3.833333

VARIABLE B A1 A0	VALUE 3.833333 2.333333 4.500000	REDUCED COST 0.000000 0.000000 0.000000
ROW 2) 3) 4) 5) 6) 7) 8) 9) 10) 11) 12) 13) 14) 15)	SLACK OR SURPLUS 5.666667 2.000000 7.666667 0.000000 7.666667 7.333333 0.333333 7.666667 0.000000 6.000000 1.666667 4.333333 3.333333	DUAL PRICES
16)	3.833333	0.000000

NO. ITERATIONS= 3

The min of the max absolute deviations was 3.833333.

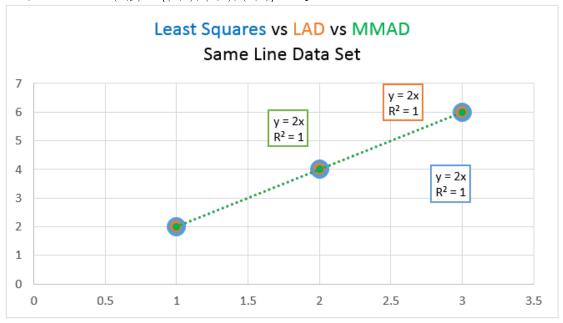
iii. Graph



The MMAD approach gives a line that appears to be more affected by the outlier at x = 2 with a much higher y-intercept than the LAD and least squares approaches. It's slope however, is only slightly steeper than the least squares line.

iv. Same Line Data Set

Yes, the data set (x,y) = [(1,2), (2,4), (3,6)] computes the same line for all three methods.



LAD Method:

```
MIN b1 + b2 + b3
ST

b1 + 1 a1 + a0 >= 2
b1 - 1 a1 - a0 >= -2
b2 + 2 a1 + a0 >= 4
b2 - 2 a1 - a0 >= -4
b3 + 3 a1 + a0 >= 6
b3 - 3 a1 - a0 >= -6
b1 > 0
b2 > 0
b3 > 0

END
```

LP OPTIMUM FOUND AT STEP 1

OBJECTIVE FUNCTION VALUE

1) 0.000000E+00

VARIABLE B1	VALUE 0.000000	REDUCED COST 1.000000
В2	0.00000	1.000000
В3	0.00000	1.000000
A1	2.00000	0.000000
AO	0.000000	0.00000
ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.00000	0.000000
3)	0.00000	0.000000
4)	0.00000	0.000000
5)	0.00000	0.000000
6)	0.00000	0.000000
7)	0.00000	0.000000
8)	0.00000	0.000000
9)	0.00000	0.000000
10)	0.00000	0.000000

NO. ITERATIONS= 1

MMAD Method:

```
MIN b
ST

b + 1 a1 + a0 >= 2
b - 1 a1 - a0 >= -2
b + 2 a1 + a0 >= 4
b - 2 a1 - a0 >= -4
b + 3 a1 + a0 >= 6
b - 3 a1 - a0 >= -6
b > 0

END
```

LP OPTIMUM FOUND AT STEP

OBJECTIVE FUNCTION VALUE

1) 0.000000E+00

VALUE	REDUCED COST
0.00000	1.000000
2.000000	0.000000
0.00000	0.000000
SLACK OR SURPLUS	DUAL PRICES
0.00000	0.000000
0.00000	0.000000
0.00000	0.000000
0.00000	0.000000
0.00000	0.000000
0.00000	0.000000
0.00000	0.000000
	0.000000 2.000000 0.000000 0.000000 0.000000 0.000000

NO. ITERATIONS= 1

Least Squares Line was produced in Excel (see graph above) rather than LINDO.

Part C: Objective:

 $Min \sum_{i=1}^{n} b_i$, for some b_i representing $|y_i - (a_2x_{2i} + a_1x_{1i} + a_0)|$ for each i = 1...n s.t.

For each i = 1...n, $b_i - (y_i - (a_2x_{2i} + a_1x_{1i} + a_0)) >= 0$ For each i = 1...n, $b_i + (y_i - (a_2x_{2i} + a_1x_{1i} + a_0)) >= 0$

```
MIN b1 + b2 + b3 + b4 + b5 + b6
ST
  b1 + 1 a2 + 1 a1 + a0 >= 5
   b1 - 1 a2 - 1 a1 - a0 >= -5
   b2 + 2 a2 + 1 a1 + a0 >= 9
  b2 - 2 a2 - 1 a1 - a0 >= -9
  b3 + 2 a2 + 2 a1 + a0 >= 12
  b3 - 2 a2 - 2 a1 - a0 >= -12
  b4 + 1 a2 + 0 a1 + a0 >= 3
  b4 - 1 a2 - 0 a1 - a0 >= -3
  b5 + 0 a2 + 0 a1 + a0 >= 0
  b5 - 0 a2 - 0 a1 - a0 >= 0
  b6 + 3 a2 + 1 a1 + a0 >= 11
  b6 - 3 a2 - 1 a1 - a0 >= -11
  b1 > 0
  b2 > 0
  b3 > 0
  b4 > 0
  b5 > 0
  b6 > 0
END
```

LP OPTIMUM FOUND AT STEP 8

3)

4)

OBJECTIVE FUNCTION VALUE

1) 2.000000

VARIA	ABLE	7	/ALUE	REDUCE	D CC	ST
	В1		1.000000	0.	0000	00
	В2		0.000000	0.	0000	00
	В3		0.000000	0.	5000	00
	В4		0.000000	0.	0000	00
	В5		0.000000	0.	5000	00
	В6		1.000000	0.	0000	00
	A2		3.000000	0.	0000	00
	A1		3.000000	0.	0000	00
	ΑO		0.000000	0.	0000	00
	ROW	SLACK	OR SURPLUS	DUAL	PRIC	ES
	2)		2.000000	0.	0000	00

0.000000

0.00000

-1.000000

-1.000000

5)	0.000000	0.000000
6)	0.000000	-0.500000
7)	0.000000	0.000000
8)	0.000000	-1.000000
9)	0.000000	0.000000
10)	0.000000	0.000000
11)	0.000000	-0.500000
12)	2.000000	0.000000
13)	0.000000	-1.000000
14)	1.000000	0.000000
15)	0.000000	0.000000
16)	0.000000	0.000000
17)	0.000000	0.000000
18)	0.000000	0.000000
19)	1.00000	0.000000

The sum of the absolute deviations was 2. The regression model is: $y = 3x_2 + 3x_1 + 0$