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Hit and Run and Stuff

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6 1 Hit-and-Run

7 The Hit-and-Run algorithm used for sampling in a convex body K , was introduced by
8 Smith in 1984 [5]. The mixing time is known to be $\mathcal{O}^*(n^2 R^2/r^2)$, where R and r are
9 the radii of the inscribed and circumscribed ball of K respectively [1, 3]. In the case of
10 the muscles of a limb, we are interested in the polygon P that is given by the set of all
11 possible activations $\mathbf{a} \in \mathbb{R}^n$ that satisfy

$$\mathbf{f} = A\mathbf{a}, \mathbf{a} \in [0, 1]^n,$$

12 where $\mathbf{f} \in \mathbb{R}^m$ is a fixed force vector and $A = J^{-T} R F_m \in \mathbb{R}^{m \times n}$. P is bounded by the
13 unit n -cube since all variables a_i , $i \in [n]$ are bounded by 0 and 1 from below, above
14 respectively.

Consider the following 1×3 example.

$$1 = \frac{10}{3}a_1 - \frac{53}{15}a_2 + 2a_3$$

$$a_1, a_2, a_3 \in [0, 1],$$

15 the set of feasible activations is given by the shaded set in figure ??.

16 **FIGURE 1**

17 The Hit-and-Run walk on P is defined as follows (it works analogously for any convex
18 body).

- 19 1. Find a given starting point \mathbf{p} of P .
- 20 2. Generate a random direction through \mathbf{p} (uniformly at random over all directions).
- 21 3. Choose the next point of the sampling algorithm uniformly at random from the
22 segment of the line in P .
- 23 4. Repeat from (b) the above steps with the new point as the starting point.

24 **FIGURE 2:** *Figures in 2D or 3D? Better in 3D to have the overview.*

- 25 1. *figure inner point*
- 26 2. *figure choice of direction*
- 27 3. *figure line segment*
- 28 4. *figure choice of new point*
- 29 5. *figure of distribution after some sampling*

30 The implementation of this algorithm is straight forward except for the choice of the
 31 random direction. How do we sample uniformly at random (u.a.r.) from all directions
 32 in P ? Suppose that \mathbf{q} is a direction in P and $p \in P$. Then by definition of P , \mathbf{q} must
 33 satisfy $\mathbf{f} = A(\mathbf{p} + \mathbf{q})$. Since $\mathbf{p} \in P$, we know that $\mathbf{f} = A\mathbf{p}$ and therefore

$$\mathbf{f} = A(\mathbf{p} + \mathbf{q}) = \mathbf{f} + A\mathbf{q}$$

34 and hence

$$A\mathbf{q} = 0.$$

35 **FIGURE 3:** q must be parallel to the plane given by $1 = \frac{10}{3}a_1 - \frac{53}{15}a_2 + 2a_3$.

36 We therefore need to choose directions uniformly at random from all directions in
 37 the vectorspace

$$V = \{\mathbf{q} \in \mathbb{R}^n | A\mathbf{q} = 0\}.$$

38 As shown by Marsaglia this can be done as follows [4].

- 39 1. Find an orthonormal basis $b_1, \dots, b_r \in \mathbb{R}^n$ of $A\mathbf{q} = 0$.
- 40 2. Choose $(\lambda_1, \dots, \lambda_r) \in \mathcal{N}(0, 1)^n$ (from the Gaussian distribution).
- 41 3. $\sum_{i=1}^r \lambda_i b_i$ is a u.a.r. direction.

42 A basis of a vectorspace V is a minimal set of vectors that generate V , and it is
 43 orthonormal if the vectors are pairwise orthogonal (perpendicular) and have unit length.
 44 Using basic linear algebra one can find a basis for $V = \{A\mathbf{q} = 0\}$ and orthogonalize it
 45 with the well known Gram-Schmidt method (for details see e.g. [2]). Note that in order
 46 to get the desired u.a.r. distribution the basis needs to be orthonormal. For the limb
 47 case we can safely assume that the rows of A are linearly independent and hence the
 48 number of basis vectors is $n - m$.

49 **FIGURE 4**

- 50 1. *Some basis*
- 51 2. *orthonormal basis*

52 References

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