

# PARALLEL MACHINE SCHEDULING

## 1.introduction

- generalization of the single machine
- special case of flow shop
- decomposition procedures for multistage systems

## 2. the makespan without preemption: $P_m || C_{max}$

- **LPT(longest process time first) rule** : yield a good bound, not optimal ([page95](#)).
- **LFJ(least flexible job fist) rule**: optimal for  $P_m | p_j = 1, M_j | C_{max}$  ([page103](#)).

## 3. the makespan with preemption

- **An optimal algorithm: more practical in practise** ([page106](#)).
- **LRPT largest remaining process time fist**: optimal rule for discrete time and continues time  $P_m | r_j, prmp | C_{max}$  (可以指定一个整数的时间间隔来切换, 得到方案后再重新整理, 使方案更合理)
- **LRPT-FM largest remaining process time fist to the fast machine**: optimal rule for discrete time and continues time  $Q_m | r_j, prmp | C_{max}$  (可以指定一个整数的时间间隔来切换, 得到方案后再重新整理, 使方案更合理)

## 4. the total completion time without preemptions

- **SPT shortest process time first rule** : optimal for both  $P_m || \sum C_j$  and  $1 || \sum C_j$
- **WSPT weight shortest process time first rule** : optimal for  $1 || \sum w_j C_j$ , but not optimal for  $P_m || \sum w_j C_j$
- **LFJ(least flexible job fist) rule**: optimal for  $P_m || \sum C_j$

## 5. the total completion time with preemptions

- **SPT shortest process time first rule:** optimal for  $P_m | prmp | \sum C_j$
- **SRPT-FM shortest remian process time with fast machine first rule:** optimal for  $Q_m | prmp | \sum C_j$