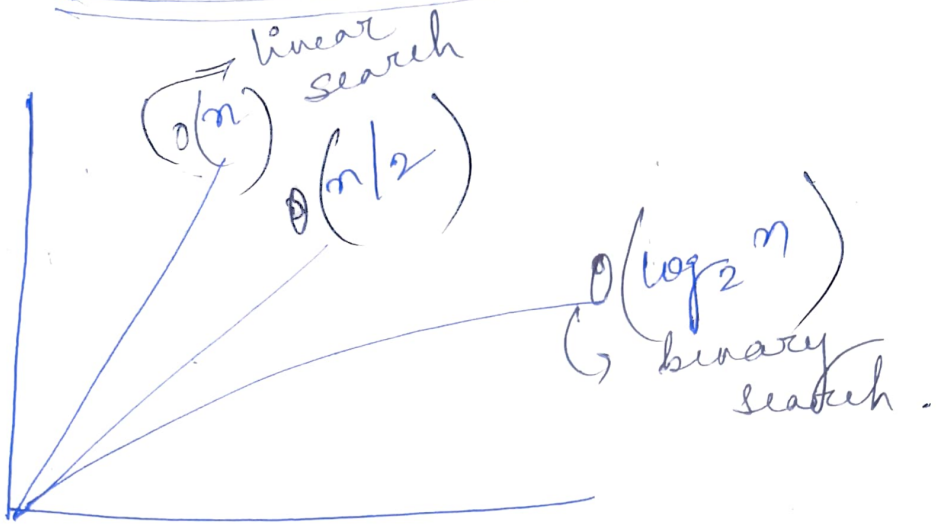
03/07/20

Lecture - 3 Algorithms



where n is the no. of steps.

- Computer scientists determine algorithms or how well they are by a terminology known as Big O.

→ Big O means or the order of.
Efficiency of algorithm is determined with the help of Big O. It is just an approximation of how fast or slow the code is.

→ Running time is how much time it takes for your program or your algorithm to run. How many steps or seconds etc..

~~Linear Search~~ $O(n)$

Running time of some algorithms

Better in terms of time req.

$O(n^2)$	— Bubble sort, selection sort.
$O(n \log n)$	merge sort
$O(n)$	— Linear search.
$O(\log n)$	— Binary search.
$O(1)$	

⇒ Big O is essentially an upper bound on how much time an algorithm might take. Usually refers to the worst cases.

where as,

Ω is referred to the best cases, opposite of Big O. It refers to the lower bound on the running-time.

$\Omega(n^2)$	— bubble sort., selection sort
$\Omega(n \log n)$	— merge sort.
$\Omega(n)$	— bubble sort
$\Omega(\log n)$	
$\Omega(1)$	— Linear search, Binary search

- $\Omega(n)$ & $O(n)$
  - counting anything.
  - It takes the same no. of steps in the best case & worst case.

- One should look for the worst case.

- → In C, strings are arrays. So, can't compare to other string using ==. Use 'strcmp'. It returns 0 if two strings are the same.

- → We can define custom data types in C. using typedef.

→ Struct is a container where we can put multiple data types.

Bubble Sort Pseudocode

```
Repeat n-1 times
    for i from 0 to n-2
        if ith & i+1'th element out of order
            swap them
```

Running time

$$(n-1)(n-1)$$
$$= n^2 - 2n + 1$$
$$= O(n^2)$$

Pseudocode of Selection sort

n times ← for i from 0 to n-1

n times ← find smallest item between ith item & last item

swap smallest item with ith item.

$$n + (n-1) + (n-2) + \ldots + 1$$

$$= \frac{n(n+1)}{2} = \frac{n^2+1}{2} = \frac{n^2}{2} + \frac{1}{2}$$

$$= O(n^2)$$

Improved bubble sort

Repeat until no swaps

for i from 0 to n-2

if ith & i'th elements out of order

swap them.

→ $\Omega(n)$

→ Merge Sort divides the problem in half each time, so running time is $O(\log n) + n$ → steps to look at each element once.

Θ Notation

To describe the running times of algorithms if the upper bound and the lower bound is the same..

Ex — Merge sort $\Theta(n \log n)$

selection sort $\Theta(n^2)$