```
SELECT *
FROM COPY_EMP;

no rows selected
```

```
MERGE INTO copy_emp c
  USING employees e
  ON (c.employee_id = e.employee_id)
WHEN MATCHED THEN
  UPDATE SET
     ...
WHEN NOT MATCHED THEN
 INSERT VALUES...;
```

```
SELECT *
FROM COPY_EMP;

20 rows selected.
```

The condition `c.employee_id = e.employee_id` is evaluated. Because the `COPY_EMP` table is empty, the condition returns false: there are no matches. The logic falls into the `WHEN NOT MATCHED` clause, and the `MERGE` command inserts the rows of the `EMPLOYEES` table into the `COPY_EMP` table.

If rows existed in the `COPY_EMP` table and employee IDs matched in both tables (the `COPY_EMP` and `EMPLOYEES` tables), the existing rows in the `COPY_EMP` table would be updated to match the `EMPLOYEES` table.

Database Transactions

A database transaction consists of one of the following:

- **DML statements which constitute one consistent change to the data**
- **One DDL statement**
- **One DCL statement**

The Oracle Server ensures data consistency based on transactions. Transactions give you more flexibility and control when changing data, and they ensure data consistency in the event of user process failure or system failure.

Transactions consist of DML statements that make up one consistent change to the data. For example, a transfer of funds between two accounts should include the debit to one account and the credit to another account in the same amount. Both actions should either fail or succeed together; the credit should not be committed without the debit.

Transaction Types

Type	Description
Data manipulation language (DML)	Consists of any number of DML statements that the Oracle Server treats as a single entity or a logical unit of work
Data definition language (DDL)	Consists of only one DDL statement
Data control language (DCL)	Consists of only one DCL statement

- **Begin when the first DML SQL statement is executed**
- **End with one of the following events:**
  - **A COMMIT or ROLLBACK statement is issued**
  - **A DDL or DCL statement executes (automatic commit)**
  - **The user exits *i*SQL*Plus**
  - **The system crashes**

When Does a Transaction Start and End?

A transaction begins when the first DML statement is encountered and ends when one of the following occurs:

- A COMMIT or ROLLBACK statement is issued
- A DDL statement, such as CREATE, is issued
- A DCL statement is issued
- The user exits *i*SQL*Plus
- A machine fails or the system crashes

After one transaction ends, the next executable SQL statement automatically starts the next transaction.

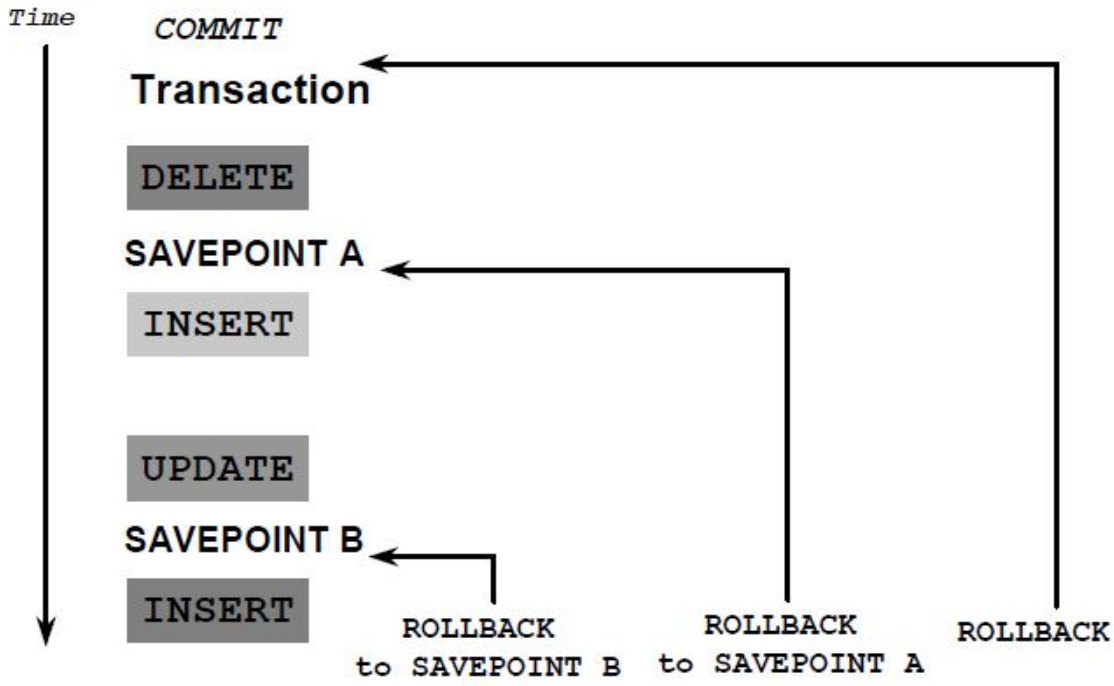
A DDL statement or a DCL statement is automatically committed and therefore implicitly ends a transaction.

Advantages of COMMIT and ROLLBACK Statements

With COMMIT and ROLLBACK statements, you can:

- **Ensure data consistency**
- **Preview data changes before making changes permanent**
- **Group logically related operations**

Controlling Transactions



Explicit Transaction Control Statements

You can control the logic of transactions by using the COMMIT, SAVEPOINT, and ROLLBACK statements.

Statement	Description
COMMIT	Ends the current transaction by making all pending data changes permanent
SAVEPOINT *name*	Marks a savepoint within the current transaction
ROLLBACK	ROLLBACK ends the current transaction by discarding all pending data changes
ROLLBACK TO *SAVEPOINT name*	ROLLBACK TO SAVEPOINT rolls back the current transaction to the specified savepoint, thereby discarding any changes and or savepoints created after the savepoint to which you are rolling back. If you omit the TO SAVEPOINT clause, the ROLLBACK statement rolls back the entire transaction. As savepoints are logical, there is no way to list the savepoints you have created.

Note: SAVEPOINT is not ANSI standard SQL.

Rolling Back Changes to a Marker

- Create a marker in a current transaction by using the SAVEPOINT statement.
- Roll back to that marker by using the ROLLBACK TO SAVEPOINT statement.

```
UPDATE...
SAVEPOINT update_done;
Savepoint created.
INSERT...
ROLLBACK TO update_done;
Rollback complete.
```

Rolling Back Changes to a Savepoint

You can create a marker in the current transaction by using the SAVEPOINT statement which divides the transaction into smaller sections. You can then discard pending changes up to that marker by using the ROLLBACK TO SAVEPOINT statement.

If you create a second savepoint with the same name as an earlier savepoint, the earlier savepoint is deleted.

Implicit Transaction Processing

- An automatic commit occurs under the following circumstances:
  - DDL statement is issued
  - DCL statement is issued
  - Normal exit from *i*SQL*Plus, without explicitly issuing COMMIT or ROLLBACK statements
- An automatic rollback occurs under an abnormal termination of *i*SQL*Plus or a system failure.

Status	Circumstances
Automatic commit	DDL statement or DCL statement is issued. *i*SQL*Plus exited normally, without explicitly issuing COMMIT or ROLLBACK commands.
Automatic rollback	Abnormal termination of *i*SQL*Plus or system failure.

Note: A third command is available in *i*SQL*Plus. The AUTOCOMMIT command can be toggled on or off. If set to on, each individual DML statement is committed as soon as it is executed. You cannot roll back the changes. If set to off, the COMMIT statement can still be issued explicitly. Also, the COMMIT statement is issued when a DDL statement is issued or when you exit from *i*SQL*Plus.

System Failures

When a transaction is interrupted by a system failure, the entire transaction is automatically rolled back. This prevents the error from causing unwanted changes to the data and returns the tables to their state at the time of the last commit. In this way, the Oracle Server protects the integrity of the tables.

From *i*SQL*Plus, a normal exit from the session is accomplished by clicking the Exit button. With SQL*Plus, a normal exit is accomplished by typing the command EXIT at the prompt. Closing the window is interpreted as an abnormal exit.

State of the Data

Before COMMIT or ROLLBACK

- The previous state of the data can be recovered.
- The current user can review the results of the DML operations by using the SELECT statement.
- Other users *cannot* view the results of the DML statements by the current user.
- The affected rows are *locked*; other users cannot change the data within the affected rows.

Committing Changes

Every data change made during the transaction is temporary until the transaction is committed.

State of the data before COMMIT or ROLLBACK statements are issued:

- Data manipulation operations primarily affect the database buffer; therefore, the previous state of the data can be recovered.
- The current user can review the results of the data manipulation operations by querying the tables.
- Other users cannot view the results of the data manipulation operations made by the current user. The Oracle Server institutes read consistency to ensure that each user sees data as it existed at the last commit.
- The affected rows are locked; other users cannot change the data in the affected rows.

State of the Data After COMMIT

- Data changes are made permanent in the database.
- The previous state of the data is permanently lost.
- All users can view the results.
- Locks on the affected rows are released; those rows are available for other users to manipulate.
- All savepoints are erased.

Make all pending changes permanent by using the COMMIT statement. Following a COMMIT statement:

- Data changes are written to the database.
- The previous state of the data is permanently lost.
- All users can view the results of the transaction.
- The locks on the affected rows are released; the rows are now available for other users to perform new data changes.
- All savepoints are erased.

Committing Data

- Make the changes.

```
DELETE FROM employees
WHERE  employee_id = 99999;
1 row deleted.

INSERT INTO departments
VALUES (290, 'Corporate Tax', NULL, 1700);
1 row inserted.
```

- Commit the changes.

```
COMMIT;
Commit complete.
```

The example in the slide deletes a row from the EMPLOYEES table and inserts a new row into the DEPARTMENTS table. It then makes the change permanent by issuing the COMMIT statement.

Example

Remove departments 290 and 300 in the DEPARTMENTS table, and update a row in the COPY_EMP table. Make the data change permanent.

```
DELETE FROM departments
WHERE  department_id IN (290, 300);
2 rows deleted.

UPDATE copy_emp
  SET department_id = 80
  WHERE employee_id = 206;
1 row updated.

COMMIT;
Commit Complete.
```

State of the Data After ROLLBACK

Discard all pending changes by using the ROLLBACK statement:

- Data changes are undone.
- Previous state of the data is restored.
- Locks on the affected rows are released.

```
DELETE FROM   copy_emp;
22 rows deleted.
ROLLBACK;
Rollback complete.
```

Rolling Back Changes

Discard all pending changes by using the ROLLBACK statement. Following a ROLLBACK statement:

- Data changes are undone.
- The previous state of the data is restored.
- The locks on the affected rows are released.

Example

While attempting to remove a record from the TEST table, you can accidentally empty the table. You can correct the mistake, reissue the proper statement, and make the data change permanent.

```
DELETE FROM  test;
25,000 rows deleted.

ROLLBACK;
Rollback complete.

DELETE FROM  test
WHERE        id = 100;
1 row deleted.

SELECT   *
FROM     test
WHERE    id = 100;
No rows selected.

COMMIT;
Commit complete.
```

Statement-Level Rollback

- If a single DML statement fails during execution, only that statement is rolled back.
- The Oracle Server implements an implicit savepoint.
- All other changes are retained.
- The user should terminate transactions explicitly by executing a COMMIT or ROLLBACK statement.

Statement-Level Rollbacks

Part of a transaction can be discarded by an implicit rollback if a statement execution error is detected. If a single DML statement fails during execution of a transaction, its effect is undone by a statement-level rollback, but the changes made by the previous DML statements in the transaction are not discarded. They can be committed or rolled back explicitly by the user.

The Oracle Server issues an implicit commit before and after any data definition language (DDL) statement. So, even if your DDL statement does not execute successfully, you cannot roll back the previous statement because the server issued a commit.

Terminate your transactions explicitly by executing a COMMIT or ROLLBACK statement.

Read Consistency

- **Read consistency guarantees a consistent view of the data at all times.**
- **Changes made by one user do not conflict with changes made by another user.**
- **Read consistency ensures that on the same data:**
  - **Readers do not wait for writers**
  - **Writers do not wait for readers**

Database users access the database in two ways:

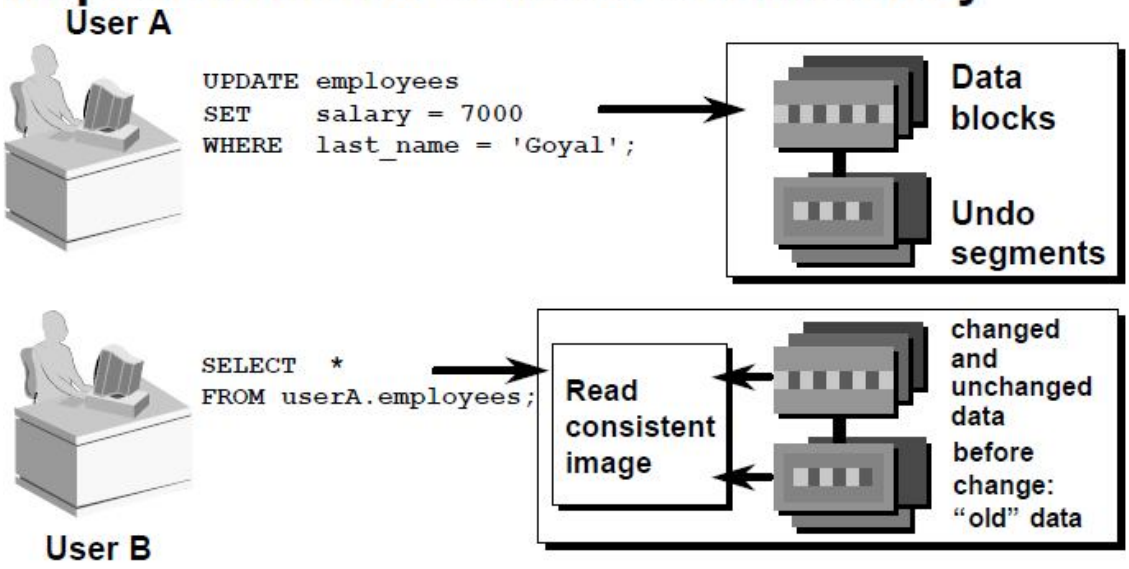
- Read operations (`SELECT` statement)
- Write operations (`INSERT`, `UPDATE`, and `DELETE` statements)

You need read consistency so that the following occur:

- The database reader and writer are ensured a consistent view of the data.
- Readers do not view data that is in the process of being changed.
- Writers are ensured that the changes to the database are done in a consistent way.
- Changes made by one writer do not disrupt or conflict with changes another writer is making.

The purpose of read consistency is to ensure that each user sees data as it existed at the last commit, before a DML operation started.

Implementation of Read Consistency



Implementation of Read Consistency

Read consistency is an automatic implementation. It keeps a partial copy of the database in undo segments.

When an insert, update, or delete operation is made to the database, the Oracle Server takes a copy of the data before it is changed and writes it to a undo segment.

All readers, except the one who issued the change, still see the database as it existed before the changes started; they view the undo segment's snapshot of the data.

Before changes are committed to the database, only the user who is modifying the data sees the database with the alterations; everyone else sees the snapshot in the undo segment. This guarantees that readers of the data read consistent data that is not currently undergoing change.

When a DML statement is committed, the change made to the database becomes visible to anyone executing a `SELECT` statement. The space occupied by the "old" data in the undo segment file is freed for reuse.

If the transaction is rolled back, the changes are undone:

- The original, older version, of the data in the undo segment is written back to the table.
- All users see the database as it existed before the transaction began.

Locking

In an Oracle database, locks:

- **Prevent destructive interaction between concurrent transactions**
- **Require no user action**
- **Use the lowest level of restrictiveness**
- **Are held for the duration of the transaction**
- **Are of two types: explicit locking and implicit locking**

Locks are mechanisms that prevent destructive interaction between transactions accessing the same resource, either a user object (such as tables or rows) or a system object not visible to users (such as shared data structures and data dictionary rows).

How the Oracle Database Locks Data

Locking is performed automatically and requires no user action. Implicit locking occurs for SQL statements as necessary, depending on the action requested. Implicit locking occurs for all SQL statements except `SELECT`.

The users can also lock data manually, which is called explicit locking.

Implicit Locking

- Two lock modes:
  - Exclusive: Locks out other users
  - Share: Allows other users to access the server
- High level of data concurrency:
  - DML: Table share, row exclusive
  - Queries: No locks required
  - DDL: Protects object definitions
- Locks held until commit or rollback

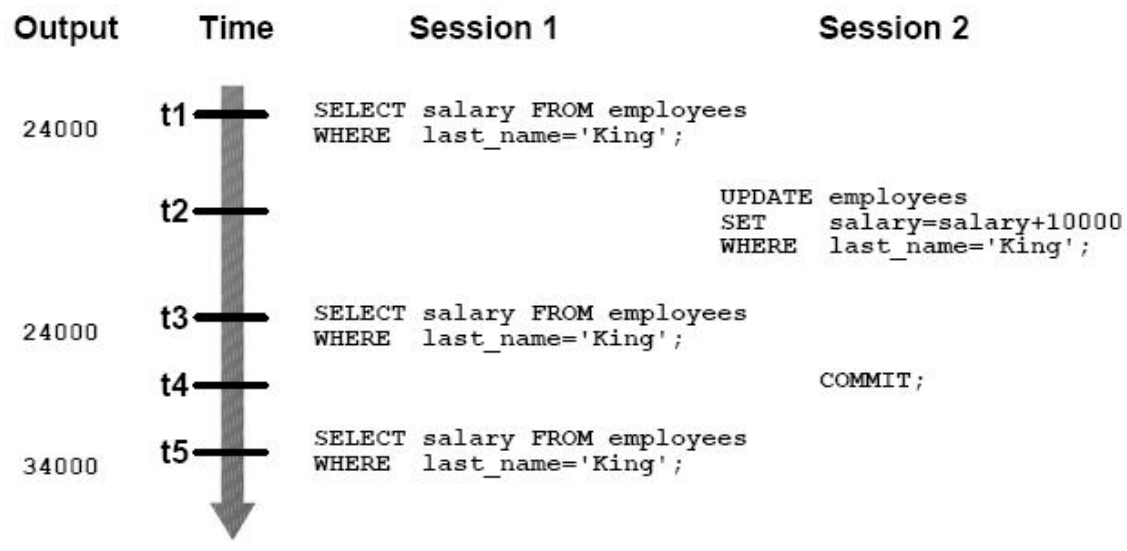
DML Locking

When performing data manipulation language (DML) operations, the Oracle Server provides data concurrency through DML locking. DML locks occur at two levels:

- A share lock is automatically obtained at the table level during DML operations. With share lock mode, several transactions can acquire share locks on the same resource.
- An exclusive lock is acquired automatically for each row modified by a DML statement. Exclusive locks prevent the row from being changed by other transactions until the transaction is committed or rolled back. This lock ensures that no other user can modify the same row at the same time and overwrite changes not yet committed by another user.

Note: DDL locks occur when you modify a database object such as a table.

Read Consistency Example



6. Write an `INSERT` statement in a text file named `loademp.sql` to load rows into the `MY_EMPLOYEE` table. Concatenate the first letter of the first name and the first seven characters of the last name to produce the user ID.
7. Populate the table with the next two rows of sample data by running the `INSERT` statement in the script that you created.
8. Confirm your additions to the table.

ID	LAST_NAME	FIRST_NAME	USERID	SALARY
1	Patel	Ralph	rpatel	895
2	Dancs	Betty	bdancs	860
3	Biri	Ben	bbiri	1100
4	Newman	Chad	cnewman	750

9. Make the data additions permanent.

Update and delete data in the `MY_EMPLOYEE` table.

10. Change the last name of employee 3 to Drexler.
11. Change the salary to 1000 for all employees with a salary less than 900.
12. Verify your changes to the table.

ID	LAST_NAME	FIRST_NAME	USERID	SALARY
1	Patel	Ralph	rpatel	1000
2	Dancs	Betty	bdancs	1000
3	Drexler	Ben	bbiri	1100
4	Newman	Chad	cnewman	1000

13. Delete Betty Dancs from the `MY_EMPLOYEE` table.
14. Confirm your changes to the table.

ID	LAST_NAME	FIRST_NAME	USERID	SALARY
1	Patel	Ralph	rpatel	1000
3	Drexler	Ben	bbiri	1100
4	Newman	Chad	cnewman	1000

15. Commit all pending changes.

Control data transaction to the `MY_EMPLOYEE` table.

16. Populate the table with the last row of sample data by modifying the statements in the script that you created in step 6. Run the statements in the script.
17. Confirm your addition to the table.

ID	LAST_NAME	FIRST_NAME	USERID	SALARY
1	Patel	Ralph	rpatel	1000
3	Drexler	Ben	bbiri	1100
4	Newman	Chad	cnewman	1000
5	Ropeburn	Audrey	aropebur	1550

18. Mark an intermediate point in the processing of the transaction.
19. Empty the entire table.
20. Confirm that the table is empty.
21. Discard the most recent `DELETE` operation without discarding the earlier `INSERT` operation.
22. Confirm that the new row is still intact.

ID	LAST_NAME	FIRST_NAME	USERID	SALARY
1	Patel	Ralph	rpatel	1000
3	Drexler	Ben	bbiri	1100
4	Newman	Chad	cnewman	1000
5	Ropeburn	Audrey	aropebur	1550

23. Make the data addition permanent.