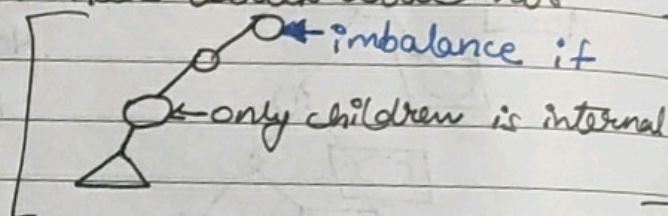


Homework 9.1Ans \Rightarrow Proof by contradiction

Let there exist a non ~~empty~~ empty AVL tree T such that $LR(T) > \frac{1}{2}$.

i.e. no. of nodes that are only children $> \frac{\text{no. of nodes}}{2}$.

Now we know that only child node is always a leaf. (Otherwise its grand parent node ~~would~~ would not be ~~imbalanced~~ balanced.)



Let number of only child $= k$.

Every node would have a parent ~~and every~~ i.e. there will be at least $k + k = 2k$ number of nodes.

~~the~~ $\frac{\text{only child nodes}}{\text{total nodes}} < \frac{1}{2}$ which contradicts the assumption.

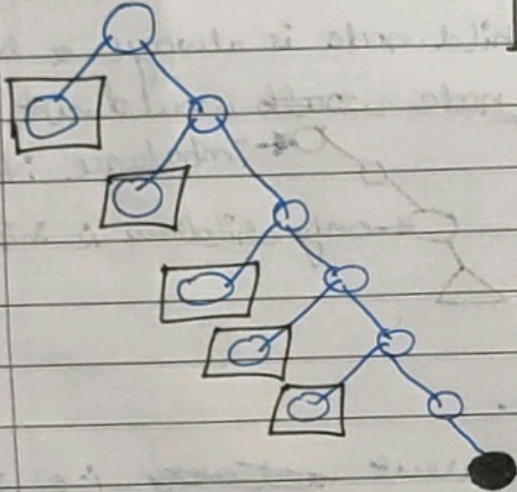
Hence for every non empty AVL tree T , $LR(T) \leq \frac{1}{2}$.

Hence Proved.

Homework 9.2

Ans \Rightarrow No. This is ~~not~~ true for all trees.

Consider the following tree as a counterexample



$\boxed{0} \Rightarrow \text{leaf node}$

\Rightarrow only child node

Here, if $n \geq 3$ is number of nodes,
then

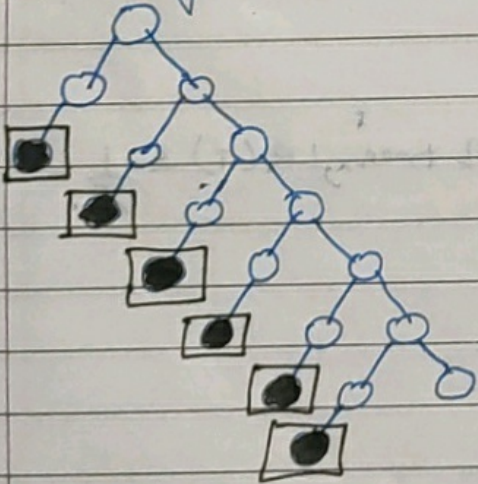
height = $\frac{n}{2}$ whereas

$$LR(T) = \frac{1.51}{2} \text{ [for all } n \geq 3]$$

Therefore, above graph has $LR(T) \leq \frac{1}{2}$ but height is not $O(\log n)$

Homework 9.3

Proof by counter example

$$\underline{\underline{Ans = 2}}$$


pb

In such trees, ~~let~~ ^{n} if ~~leaf~~ is the count for only-children, then the height of tree is $n+1$.

Therefore, $\text{height}(T) \neq O(\log n)$ ~~is not~~

① \Rightarrow leaf node



- \Rightarrow only child node