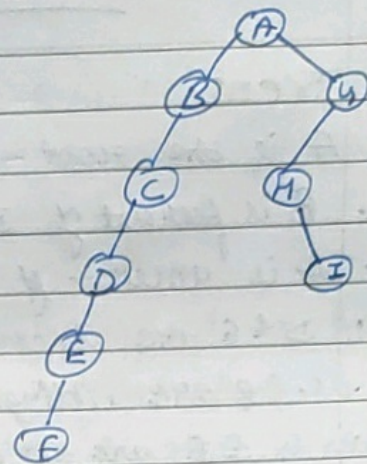


Homework 6.1



1. A is the root of the tree : True
2. F is the root of the tree : False
3. The leaves of the tree are : F and I
4. The ancestors of I are : H, G and A
5. H is a sibling of C : False
6. The height of the tree is : 5
7. The children of G are : H
8. It is a binary tree : True
9. Degrees of A, C and F are : A: 2, C: 1, F: 0
10. The root of right subtree of A is : G

Homework 6.2

Note the following points when arithmetic expression is expressed in form of binary trees.

- All the numerical values are present at the leaves.
- None of the node has exactly one child. Having ~~one~~ exactly one child is meaningless when we have operators that require 2 operands.

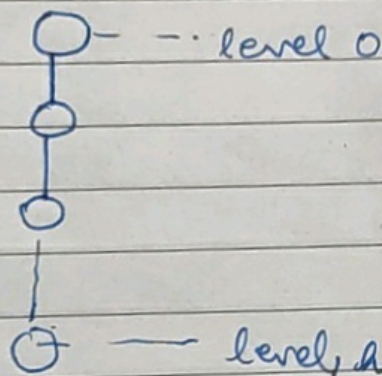
Note:- If the enum `OpType` had operators like factorial, square-root etc. that require single operand, node can have exactly one child.

```

Algorithm solve(Node node)
    if (node.type == NUM)
        return node.number
    switch (node.operation)
        case PLUS
            return solve(node.left) + solve(node.right)
        case MINUS
            return solve(node.left) - solve(node.right)
        case MULT
            return solve(node.left) * solve(node.right)
        case DIV
            return solve(node.left) / solve(node.right)
    
```


Homework 6.3

To make a binary tree having minimum nodes, ~~of~~ of given height, we will make a tree such that every node contributes to increment in height.



$$\text{total nodes} = h + 1$$

- This tree somewhat resembles singly linked list.

~~return node.parent~~

Homework 6.5

```
public boolean isRoot(Node node) {  
    return node.parent == NULL;  
}
```

```
public boolean isLeaf(Node node) {  
    return node.left == NULL && node.right == NULL;  
}
```

```
public boolean isInternal(Node node) {  
    return ! isLeaf(node);  
}
```

Homework 6.4

@ Distinct binary trees of height $n-1$ with n nodes:

This is one of the extreme case when every node except leaf node has exactly one child.

Putting node on left and right would generate a new binary tree.

1 node is root node. Remaining $n-1$ nodes have 2 choices each.

Total possible trees: $2 \times 2 \dots n-1$ times

$$= \underline{\underline{2^{n-1}}}$$

Ho

⑥ Number of possible binary trees with n nodes =

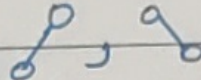
1. 1 if $n=1$
2. $n \times (\text{Number of possible binary trees with } n-1 \text{ nodes})$ if $n \neq 1$

Verification

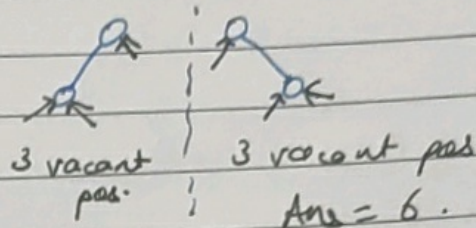
$n=1$



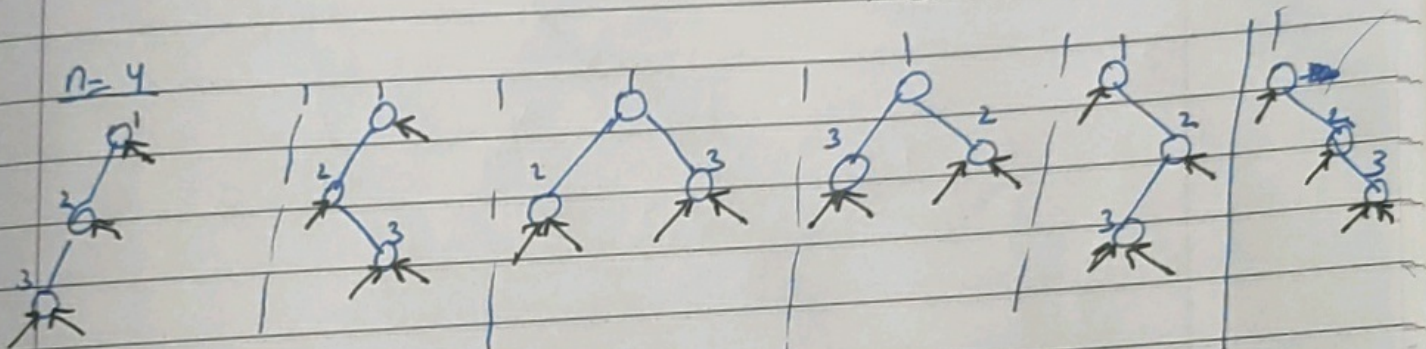
$n=2$



$n=3$



$n=4$



Total 4 possible trees for each case made by $n=3$.