

CPU SCHEDULING.

A process execution consists of a cycle of CPU execution & I/O execution.

Every process begins with CPU burst that may be followed by I/O burst then another CPU burst & I/O burst & so on eventually in the last will end up on CPU burst.

CPU Bound -: Process requiring most of the time on CPU.

I/O Bound -: Process requiring most of the time on I/O devices & peripherals.

CPU Scheduling

Non-Preemptive

Preemptive

- When a process completes its execution
- When a process leaves CPU to perform some I/O operation or to wait for an event
- if the process enters in the ready state either from a new waiting state & high priority process
- if process switch from running to ready state because the quantum expires.

CPU Scheduling Terminology :-

① Burst Time | Execution Time | Running Time

Time a process requires for running on the CPU.

② Waiting Time

Time spent by a process in ready state waiting for CPU.

$$W.T = T.A.T - B.T.$$

③ Arrival Time

When a process enters ready state

④ Execution Time

When process completes execution & exit from System

⑤ Turn Around Time

Total time spent by a process in the system

$$\boxed{T.A.T. = E.T. - A.T. \\ = B.T. + W.T.}$$

⑥ Response Time :-

Time between a process enters ready queue and get scheduled on the CPU for the first time.

C.P.U. Scheduling Criteria

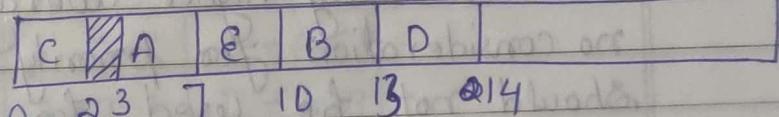
- Avg. Waiting Time
- Response Time (Avg.)
- CPU Utilization
- Throughput

D) First Come first Serve (FCFS)

- It assigns CPU to the process that arrives first
- Easy to understand & can easily be implemented using Queues
- Always non-preemptive in nature

Pid	A.T.	B.T.
A	3	4
B	5	3
C	0	2
D	5	1
E	4	3

Gantt Chart - !



$$\text{Now, } TAT = \text{E.T.} - \text{A.T.}$$

	TAT
A	$7-3=4$
B	$13-5=8$
C	$2-0=2$
D	$14-5=9$
E	$10-4=6$

$$\text{Now, W.T.}$$

$$\text{TAT} = \text{W.T.} + \text{B.T.}$$

$$\text{W.T.} \neq \text{B.T.} - \text{TAT}$$

$$\text{W.T.} = \text{TAT} - \text{B.T.}$$

$$\text{W.T.}$$

$$A \quad 4-4=0$$

$$B \quad 8-3=5$$

$$C \quad 2-2=0$$

$$D \quad 9-1=8$$

$$E \quad 6-3=3$$

$$\therefore \text{Avg. TAT} = \frac{4+8+2+9+6}{5}$$

$$= 29/5 = 5.8.$$

$$\& \text{Avg WT} = \frac{0+5+0+8+3}{5}$$

$$= 16/5 = 3.2$$

Convoy Effect - :

Smaller process have to wait for long time for bigger process to release CPU.

Eg - :	P	WT	B.T.	WT	
	P ₁	0	100	0	= 49.5 // Heavier to deal
	P ₂	1	1	99	with

Advantages - : Simple, easy to use, easy to understand, easy to implement, must be used for background process whose execution is not urgent.

Disadvantage - : Suffer from Convoy effect, normally higher average waiting time, no consideration for priority or Burst time. Should not be used for interactive systems.

Convoy → Starvation → Deadlock

Starvation occurs when the processor is biased. So, FCFS does not have starvation but it does have Convoy effect.

S_{HORTEST} J_OB F_{IRST}

→ SJF Non-Preemptive

→ Shortest Remaining time First Preemptive.

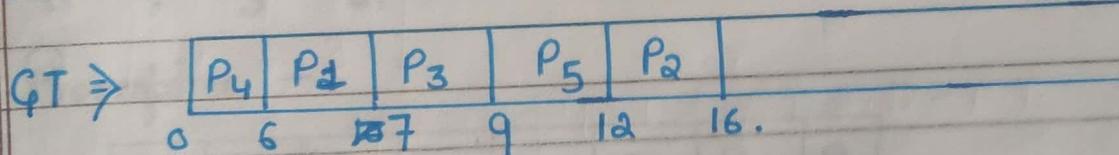
Out of all available process, CPU is assigned to the process having smallest Burst time. (No priority / seniority).

If there is a tie, FCFS is used.

Preemptive approach (SRTF) is optimal as it guarantees minimal average waiting time.

Non-Preemptive Example :-

Pid	AT	BT	TAT	W.T.
P ₁	3	1	7-3=4	3=4-1
P ₂	10	4	16-1=15	9=15-6
P ₃	4	2	9-4=5	3=5-2
P ₄	0	6	6-0=6	0=6-6
P ₅	2	3	12-2=10	7=10-3.



Preemptive Example - :-

P ₄	P ₂	P ₂	P ₁	P ₂	P ₃	P ₅	P ₄	TAT	W.T.
0	1	2	3	4	6	8	11	16.	0
*	Minimal average waiting time among all the scheduling approach.							5	1
								4	2
								16	10
								9	6.

Advantage - :

- SRTF guarantees minimal average time.
 - Provide a standard for other algos in terms of average time
 - Compared to FCFC better average response time.

Dissadvantage - :-

- Cannot be implemented, bcoz there is no way to know the burst time of a process
 - Process with longer CPU B.T. requirement will go into starvation
 - No idea of priority, process with large B.T. have poor response time.

PRIORITY SCHEDULING.

- CPU is allocated to the process which possess, the highest priorities
- Tie breaks by FCFS.
- No importance for A.T. or B.T.
- Supports both preemptive & non-preemptive approach.

P-id	B.T.	A.T.	Priority.	TAT	WT.
P ₁	4	0	2	4	0
P ₂	3	1	3	14	11
P ₃	1	2	4	10	9
P ₄	5	3	5	6	1
P ₅	2	4	5.	7	5

N.P. →

P ₁	P ₄	P ₅	P ₃	P ₂
0	4	9	11	12

15

Preemptive - :-

P ₁	P ₂	P ₃	P ₄	P ₄	P ₅	P ₂	P ₁
0	1	2	3	4	8	10	12

15.

TAT	B.T.
15	11
11	8
1	0
5	0
6	4

Advantage :-

- Provide a facility of priority specially for system process
- Allow to run important process first even if it is a user process

Disadvantage :-

- Process with smaller priority may starve
- No idea of Response or waiting time

* Ageing :-

Is a technique of gradually increasing the priority of process that wait in the system for long time.

9	9	29	5	9	3	9	9
31	61	01	8	4	8	6	10

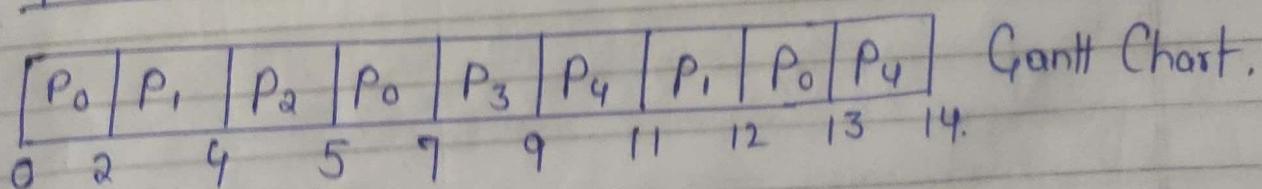
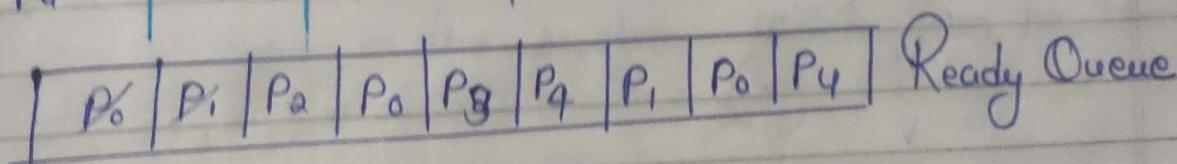
ROUND ROBIN SCHEDULING

- This algo is designed for time sharing systems, where it is not necessary to complete one process & then start another, but to be responsive & divide the time of CPU among the process in Ready State.
- Here, ready queue will be treated as circular queue.
- We fix a time quantum up to which a process can hold CPU in one go, within which either a process terminates or process must release the CPU & de-enter in the circular queue & wait for the next chance.

Always, Preemptive in nature.

$$TQ = 2.$$

P_id	A.T.	B.T.	TAT	WT
P ₀	0	5	13	8
P ₁	1	3	11	8
P ₂	2	1	3	2
P ₃	3	2	6	4
P ₄	4	3	7	4



Advantages - :

- Performs best in terms of Average response time
- Works well in time sharing System, client server architecture & interactive system
- Kind of SJF

Disadvantages - :

- Longer process may starve
- Performance depends highly on time quanta
- No idea of priority.

TU	TAT	W.A	T.A
2	8	2	0
2	11	2	1
3	2	1	6
3	3	0	2
5	5	2	2