

UNIT - 1

→ Data Structures

Data are simply values or set of values. Data values refer to a single unit of values an entity is something that has certain attributes or property which may be assigned values. These values may be numeric or non-numeric.

Data structure is logical or mathematical ^{model} ~~order~~ of particular organization of data is called as choice of particular data model depends on the particular application.

Types of data structures

→ 1. Array

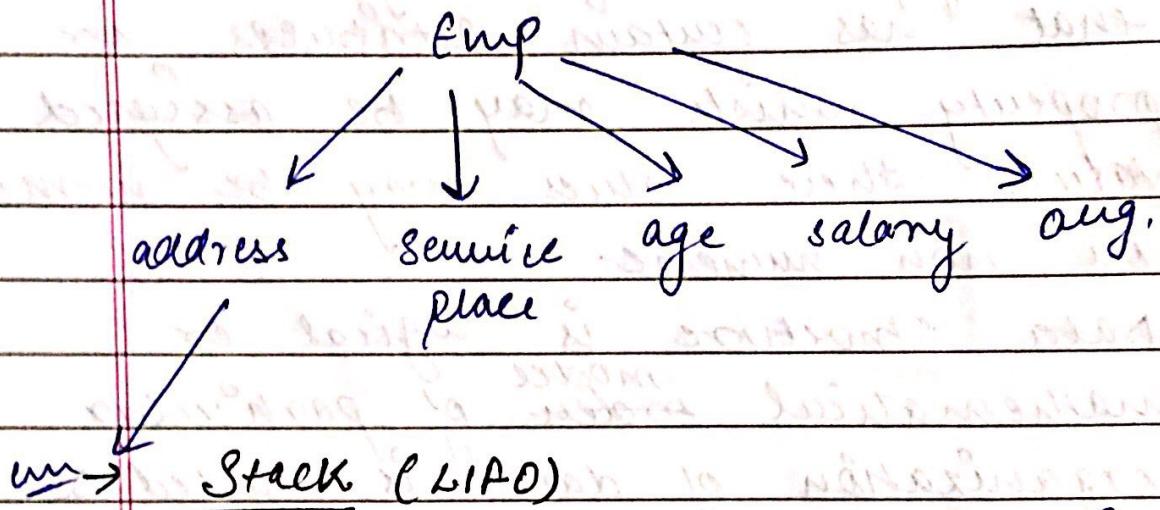
The simplest type of data structure is a linear array. List of finite no. of similar data elements sequence by a set of consecutive no's.

→ 2. Linked list

Linked list contains 2 parts in one node one part contain information and other contain linked information.

→ Tree

Tree is hierarchical relationship b/w various elements of the datastructure which reflects this relationship is called rooted tree or simple tree.



It is also called last in first out (LIFO) system is a linear list which insertion and deletion take place at only one end called top.

→ Queue (FIFO)

It is a first in first out (FIFO) system.

It is a linear list which insertion can take place only at one end of the list called rear. at the deletion can take place only at other end called front.

Data structure operations

4 operation play a major role for all data structure.

→ Traversing

Accessing each record exactly once so that certain item in the record may be proceed.

→ Searching

Finding the location of the record with the given key-value.

→ Inserting

Adding a new record to the data structure is called inserting.

→ Deleting

Removing a record from a data structure.

→ Sorting

Arranging the records in some logical order.

→ Merging

Combining the records of 2 different file into a single file.

ARRAY

UB → upper bound

LB → lower bound.

$$n = UB - LB + 1$$

* ALGORITHM - 1

Traversing of Array.

Here ~~add~~ LA is a linear array with lower bound LB and upper bound UB, apply operation PROCESS to each element.

Step 1 : [initialize counter] set $K := LB$

Step 2 : Repeat step 3 & 4 while $K \leq UB$

Step 3 : [visit element] apply PROCESS to $LA[K]$

Step 4 : [increment counter] set $K := K + 1$

[End of step 2 loop]

Step 5 : Exit

ALGORITHM - 2

Inserting an element in linear array
Insert (LA, n, k, value)

Here LA is linear array ~~subset~~
~~($LA_1, \dots, k, \dots, LA_n$)~~ ~~here~~ with UB & LB,
apply PROCESS to each element.

Step 1 : [Initialize counter] set $J := N$

Step 2 : Repeat step 3 & 4 while $J \geq k$

Step 3: [Move jth element downward] set set
 $LA[j+1] = LA[j]$

Step 4: [decrease counter] set $J := J - 1$
[End of loop 2]

Step 5: [insert an element] set $LA[k] := \text{value}$

Step 6: [reset N] set $N := N + 1$

Step 7: Exit

* ALGORITHM-3

Deletion an element in array at given location.

DELETE (LA, N, K, VALUE)

Hence LA is linear array K is location with LB & LB , apply process to each element.

Step 1: Value = $LA[K]$

Step 2: Repeat for $J = K + 1$ to $N - 1$
Set $LA[J] = LA[J + 1]$

[End of loop]

Step 3: [Reset the number of element in LA]
Set $: N = N - 1$

Step 4: Exit

Single dimension

Formulae to find index value from given array.

$$\text{LOC}(LACK) = \text{Base}(LA) + w(K-1)$$

Two dimension

Column major

$$\text{LOC}(A[i,j,k]) = \text{Base}(LA) + w[M(k-1) + (j-1)]$$

values
index is

Row major

$$\text{LOC}(A[i,j,k]) = \text{Base}(LA) + w[N(j-1) + (k-1)]$$

LA = linear array

Base = Base address of array

k = Given location

w = storage type

M, N \rightarrow row, column

J, K \rightarrow given locations

w \rightarrow Storage type

Base \rightarrow Base address of LA.

If M, N are not directly given then apply

$$A [u_1 \dots l_1, u_2 \dots l_2]$$

$$M = u_1 - l_1 + 1 \quad N = u_2 - l_2 + 1$$

$$M = 4 \quad N = 6$$

$A \in [4 \dots 7, -1 \dots 3]$

$w = 2 \text{ bytes}$

$\text{base} = 100$

address calculation for $A = 6$ and 2

$A[6, 2]$

$$M = 84$$

$$N = 5$$

$$w = 2$$

$$\begin{aligned} \text{Loc}(A[6, 2]) &= 100(\text{LA}) + 2(4(2-1) + (6-1)) \\ &= 100(\text{LA}) + 18 \\ &= 118 \end{aligned}$$

$$\begin{aligned} \text{Loc}(A[6, 2]) &= 100 + 2[5(6-1) + (2-1)] \\ &= 100 + 1 \\ &= \end{aligned}$$

column major

$$\text{Loc}(A[j, k]) = \text{Base(LA)} + w[M(k-u_1) + (j-u_2)]$$

row major

$$\text{Loc}(A[j, k]) = \text{Base(LA)} + w[N(j-l_1) + (k-l_2)]$$

* Recursion

#include <stdio.h>

main()

{

 Put a, fact ;

 printf ("Enter any number");

 scanf ("%d", &a);

 fact = rec(a);

 printf ("fact value %d", fact);

}

rec(x)

{

 Put x;

{

 Putf;

 if (x == 1)

 return(1);

 else

 f = x * rec(x-1);

 return(f);

}

g. Write a program in c for traversing an array.

```
#include <stdio.h>
void main()
{
    int LA[ ] = {2, 4, 6, 8, 9};
    int i, n=5;
    printf("Enter 5 elements in an array");
    for (i=0; i<n; i++)
    {
        printf(" LA[%d]=%d", i, LA[i]);
    }
}
```

g.

g. Write a program in c to insert an element in an array.

```
#include <stdio.h>
```

```
void main()
```

```
{
```

```
int array [MAX] = {1, 2, 4, 5};
```

```
int n=4
```

```
int i=0
```

```
int index=2
```

```
int value=3
```

```
printf("Enter array");
```

```
for (i=0; i<n; i++)
```

```
{
```

```
    printf(" array LA[%d]=%d", i, LA[i]);
```

```
for (i=N; i>=pIndex; i--)  
    {  
        CA[i+1] = CA[i]  
    }
```

$CA[pIndex] = \text{value}$,
 $N++$;

print ("Printing array after insertion")

```
for (i=0; i<N; i++)
```

```
{  
    print (CA[y.d] = y.d, E, CA[i]);  
}
```

return 0;

```
}
```

no comment on a code representing an array

the array is represented by an array of elements

each element is a character

the array is a character array

STACK : LIFO

insertion: PUSH

deletion: POP

→ Algorithm 1.4 [Push operation on stack]
PUSH [STACK, TOP, MAXSTK, VALUE]

Step 1: [stack already filled]
if $TOP = MAXSTK$ then

write OVERFLOW and return

Step 2: Set $TOP := TOP + 1$ [increase top by 1]

Step 3: Set $STACK [TOP] = VALUE$

Step 4: Return

→ Algorithm 1.5 [Pop operation on stack]
POP [STACK, TOP, VALUE]

Step 1: [stack already empty]
if $TOP = \text{other}$

write UNDERFLOW and return

Step 2: Set $VALUE = STACK [TOP]$

Step 3: Set $TOP = TOP - 1$

Step 4: Return

~~* Evaluation of postfix expression~~

→ Algorithm 1.056

(This algorithm find the value for given arithmetic expression P written in postfix notation)

$$\boxed{3 \ 5 \ 6 \ +}$$

Step1: Add right parenthesis ")" at the end of P.

Step2: Scan P from left to right and repeat step3
for each element of P until the ")" is encountered.

Step3: If operand is encountered push it on Stack.

Step4: If an operator \otimes is encountered, then

a) remove the two top element of stack,
where A is top element & B is the
next top element

b) evaluate $B \otimes A$

c) replace the result (b) back on stack

Step5: SET VALUE equal to the top of stack

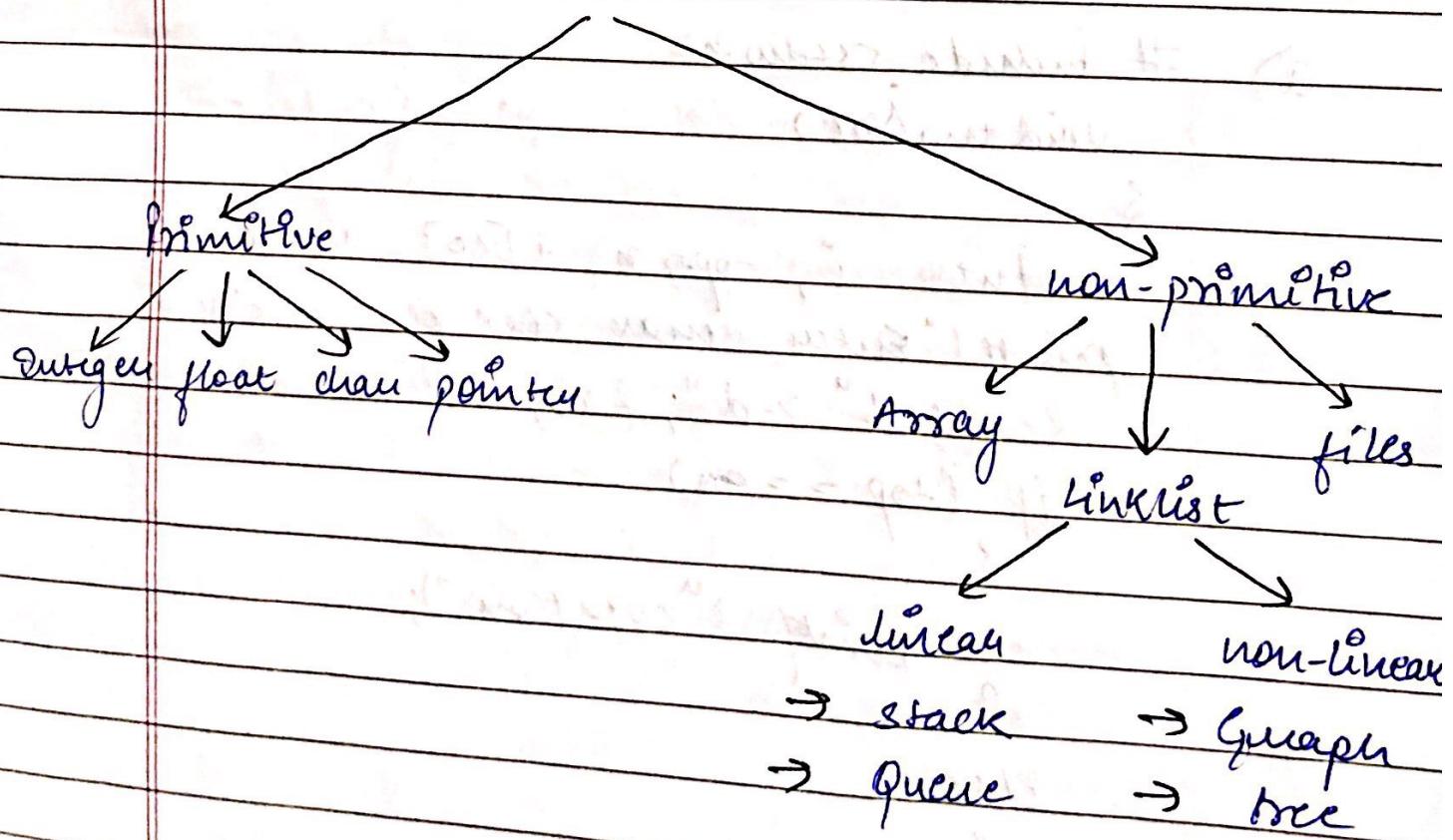
Step6: EXIT.

$$3, 5, 6, +, *, 7, 1$$

$$P = 5, 6, 2, +, *, 12, 4, /, -,)$$

1	S	15					
2	6	5	6				
3	2	5	6	2			
4	+	5	8				
5	*	40					
6	12	40	12				
7	4	40	12	4			
8	1	40	3				
9	-	37					
10)	Ans 37					

Data Structure.



Q Evaluate the postfix expression

1. 9

9

2. 10

9	10
---	----

3. 3

9	10	3
---	----	---

4. *

9	30
---	----

5. 2

9	30	2
---	----	---

6. /

9	15
---	----

7. +

24

8.)

Q WAP in c for push and pop operation of stack.

a) #include <stdio.h>
void main()

{
 int i, top = -1, a[10], value;
 printf("Enter max size of stack");
 scanf("%d", &n);
 if (top == n)
 {

} printf("overflow");

else

{

top = top + 1;

printf("Enter the value");
 scanf("%d", &value);

a[top] = value;

3

for (i=0; i< top; i++)

3

printf("%d", a[i]);

3.

Pop

b) #include <stdio.h>

void main()

{

int c, top=0, n, a[10], value;

printf("Enter max size of stack");

scanf("%d", &n);

if (top == n)

{

printf("underflow");

}

else

{

top = top - 1;

printf("Entry the value");

scanf("%d", &value);

*top = value = a[top];

{

for (i=0; i< top; i++)

{

printf("%d", a[i]);

$+AB \rightarrow$ Prefix
 $A+B \rightarrow$ infix
 $AB+ \rightarrow$ postfix.

classmate

Date _____

Page _____

Conversion of suffix to post fix rotation

$$\text{infix} = (E) \quad Q = (A+B)$$

Step 1: Push "(" onto stack and ")" to the end of expression (Q).

Step 2: Scan Q from left to right and repeat Step 3 to 6 for each element of Q until the stack is empty.

Step 4: If left parenthesis is encountered push it onto stack.

Step 5: If operand is encountered, add it to

a) If an operator \otimes is encountered then repeatedly pop from stack and add to P each operator on the top of stack which has the same precedence or higher precedence than \otimes .

b) Add \otimes to stack.

Step 6: If right parenthesis is encountered then a) repeatedly pop from stack and add to P each operator until left parenthesis is encountered

b) remove the left parenthesis

Step 7: Exit

S-No	Stack	Expression (postfix)
C	CC	A
A	CC	AB
+	CC +	AB+
B	CC +	AB+
)	(+	AB*
*	(* +	AB*
C	(* +)	AB+C
)		A+C*

$(A+B)/(C-E)$

SNo	Stack	Expression (postfix)
A C	CC	-
A	CC	A
+	CC +	A
B	CC +	AB
)	(+	AB +
/	C /	AB +
C	C/C	AB + C
C	C/C	AB + C
-	C/C -	AB + CD
D	C/C -	AB + CD -
)	C/C -	AB + CD -
)		AB + CD - /

$A * B + C / D$

for power
to hybrid

SNO

A

*

B

+

C

/

D

)

stack

C

~~AB~~ * C *

~~CD~~ * C *

~~CD~~ + C +

C +

C + /

C + /

C + /

Expression/Postfix

A

A *

A * B

A * B *

A B * C

A B * C D

(AB * CD) /

$5 * (6 + 2) - (12 / 4)$

SNO

5

*

6

+

2

)

-

C

12

/

4

)

)

stack

C

C *

CC *

CC *

CC +

CC +

C

C -

CC -

CC -

CC - /

CC - /

CC - /

Exp.

5

5

5

5, 6

5, 6 *

5, 6 * 2

5, 6 * 2 +

5, 6 * 2 +

5, 6 * 2 + 12

5, 6 * 2 + 12

5, 6 * 2 + 12

5, 6 * 2 + 12

5, 6 * 2 + 12

$$(A + (B * C) - (D * E * F) * G) * H)$$

SNO

Stack

Exp

C

CC

A

CC@

+

CC +

C

CC + C

B

CC + C

*

CC + C *

C

CC + C *

)

CC +

A, B, C *

-

CC -

A B C * +

C

CC - C

A B C * +

D

CC - C

A B C * + D

/

CC - C /

A B C * + D

E

CC - C /

A B C * + D E

T

CC - C / T

A B C * + D E

F

CC - C / T

A B C * + D E F

)

CC -

A B C * + D E F / T

*

CC *

A B C * + D E F / T *

G

CC *

A B C * + D E F / T * G

)

C

A B C * + D E F / T * G

*

C *

A B C * + D E F / T * G * -

H

C *

A B C * + D E F / T * G * - H *

)

~~Point for~~

```
#include <stdio.h>
```

```
#include <conio.h>
```

```
#define maxsize 30
```

```
void push();
```

```
void pop();
```

```
void display();
```

```
int stack [Maxsize], top;
```

```
void main()
```

```
clrscr();
```

```
int ch=1, top=-1;
```

```
while (ch)
```

```
    printf ("Push");
```

```
    printf ("Pop");
```

```
    printf ("Enter your choice");
```

```
    scanf ("%d", &ch);
```

```
    switch (ch)
```

```
        case 1: Push();
```

```
        break;
```

```
        case 2: Pop();
```

```
        break;
```

```
        default: exit(0);
```

```
}
```

```
}
```

```
getch();
```

Scanned with CamScanner

```
void push()
{
    int no;
    if (top == maxsize - 1)
        cout ("overflow");
    else
        top = top + 1;
    cout ("Enter the number");
    scanf ("%d", &no);
    stack [top] = no;

    cout ("Element in the stack");
    display();
}

void pop()
{
    int no;
    if (top == -1)
        cout ("Underflow");
    else
        top = top - 1;
    cout ("Element in stack");
    display();
}
```

void display()

{

int i;

for (i=0; i<=top; i++)

{

printf("%d", stack[i]);

}

(*push) = (*stack + top);
top++;

(*push) = (*stack + top);
top++;

(*) push = stack

top = 0;

(*) push = 0;

(*push) = (*stack + top);

Queue

front → deletion

$f \geq 0$

rear → insertion

$f \leq 0$

1	2	3	4	5	6	7
10	13	14	15	16	17	18

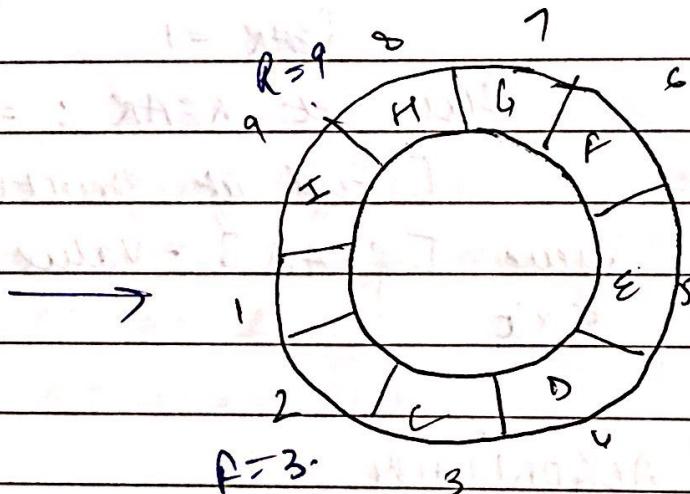
$f = 1$ $r = 2$ $l = 3$

$R = 7$

$R = 1$ $r = 2$ $l = 3$

$f = R = \text{Valid position}$

single element.



10 Six array (Linear array)

Insert \rightarrow 5, 6, 7, 8, 9

After insert

Delete \rightarrow 5, 6

Insert 15, 16, 17

Insert \rightarrow 10, 11, 12, 13, 14

Delete 10, 11

Delete \rightarrow 7, 8, 9

$f = 8$

$R = 13$

5	6	7	8	9	
---	---	---	---	---	--

$f = 6, 6$

$R = 6, 10$

ALGORITHM

QUEUE : QINSERT (QUEUE, N, FRONT, REAR, VALUE).

Step 1 : [Queue already filled]

if FRONT = 1 and REAR = N
Print "overflow".

Step 2 :

if front FRONT = NULL

[Queue is empty]

Set FRONT := 1, REAR := 1

else if REAR = N and FRONT ≠ 1

then

REAR = 1

else Set REAR := REAR + 1

[End if structure]

Step 3 : Queue [REAR] = value

Step 4 : Exit

ALGORITHM

QDELETE [QUEUE, N, FRONT, REAR, VALUE]

Step 1 : [Queue already empty]

if FRONT = NULL and REAR = NULL

Print "underflow".

Step 2 :

Value = Queue [FRONT]

Step 3 :

if FRONT = REAR

Set FRONT = NULL, REAR = NULL

else if FRONT = N then

SET FRONT = 1

else

SET FRONT = FRONT + 1

[End if structure]

Step 4: Exit.

Program

QINSERT

```
#include <stdio.h>
```

```
int main()
```

```
{
```

```
    int N, QUEUE, FRONT, REAR, VALUE, NULL=1;  
    printf("Enter the size of queue");
```

Scans ("I. d", &N);

```
if (FRONT=1 && REAR=N)  
{
```

printf("overflow");

}

else

```
{
```

```
if (FRONT=NULL)
```

{

FRONT=1;

{ REAR=1;

```
else if (REAR=N && FRONT!=1)  
{
```

REAR=1;

& REAR=REAR+1;

```
    printf ("Enter the value");
    scanf ("%d", &value);
else Queue [REAR] = value;
    & printf ("%d", Queue [N]);
}
```

⇒ Program for Queue (Circular)

```
#include <stdio.h>
insertion #define max 5
int queue [MAX];
```

```
int front = -1;
```

```
int rear = -1;
```

```
Void main ()
```

```
{
```

```
void insert();
```

```
void del();
```

```
void display();
```

```
int choice;
```

```
while (1)
```

```
{
```

```
printf ("1. Insert");
```

```
printf ("2. Delete");
```

```
printf ("3. display");
```

```
printf ("4. Exit");
```

```
printf ("Enter the choice");
```

```
scanf ("%d", &choice);
```

```
switch (choice)
```

```
{
```

```
case 1: insert();
```

```
break;
```

```
case 2: del();
```

```
break;
```

```
case 3: display();
```

```
break;
```

case 4: exit();

break;

default : printf (" invalid choice");

3.7.2

void insert ()

{

int value;

if (front == 0 && rear == max - 1)

printf (" overflow");

return;

if (front == -1)

if

front = 0;

rear = 0;

else

use if (rear == max - 1 && front != 0)

rear = 0;

else

rear = rear + 1

printf (" Enter the value");

scanf ("%d", &value);

queue [rear] = value;

else

void display ()

{

Infix to prefix :

$2 * 3 / (2 - 1) + 5 * (4 - 1)$

SNO	Stack	Exp.
))	
)))	
1)))	
-)) -	1,
4)) -	1, 4
() +	1 4 -
*) *	1 4 -
5) *	1 4 - 5
+) +	1 4 - 5 *
)) +)	1 4 - 5 *
1) +) +)	1 4 - 5 *
-) +) -	1 4 - 5 * +
2) +) -	1 4 - 5 * 1 2
() +	1 4 - 5 * 1 2 -
/) + /	1 4 - 5 * 1 2 -
3) + /	1 4 - 5 * 1 2 - 3
*) + / *	1 4 - 5 * 1 2 - 3
2) + / *	1 4 - 5 * 1 2 - 3 2
(1 4 - 5 * 1 2 - 3 2 *

+ , / , + , 2 , 3 , - , 2 , 1 , * , 5 , - , 4 , 1

$$(A + (B * C) - (D / E \uparrow F) * G) * H$$

SNO	STACK	Expression
))	
H)	H
*) *	H
)) *	H
G) *	AG
*) *) *	AG
)) *) *	AG
F) *) *	HGF
\uparrow) *) * \uparrow	HGF
\varepsilon) *) * \uparrow	HGF\epsilon
/) *) *) /	HGFET
D) *) *) /	HGFETD
() *) *	HGFETD /
-) *) -	HGFETD / *
)) *) -	HGFETD / *
C) *) -	HGFETD / * C
*) *) -) *	HGFETD / * C
B) *) -) *	HGFETD / * C B
() *) -	HGFETD / * C B *
+) *) - +	HGFETD / * C B *
A) *) - +	HGFETD / * C B * A
() *	HGFETD / * C B * A
(HGFETD / * C B * A + - *

* - + A * B C * / D \uparrow E F G H

Deletion

void delete ()

{

int value;

if [front == 0 & rear == 0]

print ("underflow");

return;

if [front == ~~rear~~ ~~front == 0~~]

{

front = 0;

rear = 0;

}

else if [front == max - 1]

{

front = 0;

}

else

{

front = front + 1;

}

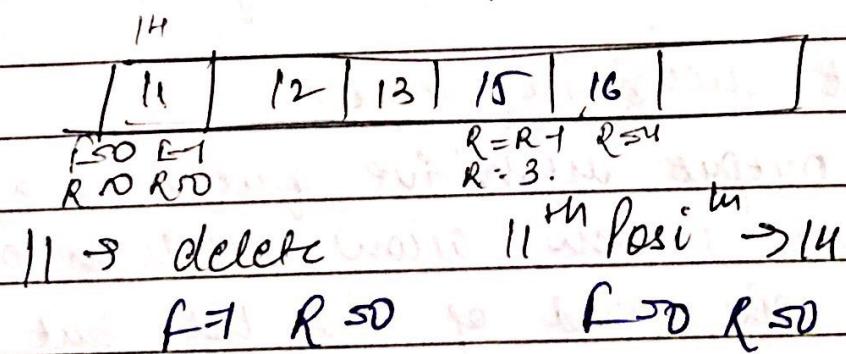
queue [front] = value;

Priority Queue

Priority queue is a collection of elements such that each element has been assigned a priority and the order in which elements are deleted and process comes from the following rules :-

- (1) An element of higher priority is processed before element of lower priority
- (2) 2 elements with same priority are processed acc. to the order in which they were added in the queue.

Double ended Queue



After 13^{th} 15, 16
 $F = 0$ $R = 4$

16 Delete
 $R = 3$ $F = 0$.

Double ended queue

- (1) Input restrictive Queue
- (2) Output restrictive queue

Dequeue or Double ended queue

Dequeue is a linear list in which elements can be added or removed at either end but not in the middle.

There are 2 variation in queue.

- (1) Input restrictive Queue

This queue allows insertion at only one end of the list but deletion at both end of the list.

- (2) Output restrictive Queue.

An output restrictive queue is a queue which allow deletion only at one end of the list but insertion at both end of the list.

Q

info

link.

del

start

[5]

1

50 50

82

2

10

484

3

40

8X

4

20

6

[10 3] []

5

50

2

6

60

87

7

30

83

60 → Link [67]

info

link

start

[5]

1

2

++

3

3

4

22

87

5

77

*

6

44

7

7

8

66

4X

DYNAMIC MEMORY ALLOCATION

Malloc (Memory allocation)

Dynamic memory allocation can be defined as a procedure in which the size of data structure is changed during the runtime.

`malloc` or memory allocation method is used to dynamically allocate the single large block of memory with specified size. It returns a pointer of type void which can be cast into a pointer of any form.

for dynamic: #include < stdlib.h >

`ptr = (cast-type *) malloc (byte-size)`

example:

`ptr = (Put *) malloc (100 * size of (int))`

```

#include <stdio.h>
#include <stdlib.h>

int main ()
{
    int *ptr;
    int n, i;
    n = 5;
    printf ("Enter number of elements %d", n);
    ptr = (int *) malloc (n * sizeof (int));
    if (ptr == NULL)
    {
        printf ("Memory not allocated");
        exit (0);
    }
    else
    {
        for (i = 0; i < n; i++)
        {
            ptr[i] = i + 1;
        }
        printf ("The elements of the array");
        for (i = 0; i < n; i++)
        {
            printf ("%d", ptr[i]);
        }
    }
    return 0;
}

```

alloc on contiguous Memory allocation.

alloc means ^{on} contiguous allocation
means is used to dynamically
allocate the specified no of blocks of
memory of specified type.
it initialize each block with a
default value 0.

Syntax :

$\text{ptr} = (\text{cast - type *}) \text{alloc}(n, \text{element - size})$

example :

$\text{ptr} = (\text{float *}) \text{alloc}(25, \text{size of float})$

program - same as Malloc.

Malloc \rightarrow alloc.

$\text{ptr} = (\text{int *}) \text{alloc}(5, \text{size of int})$

Free

free method is used to dynamically deallocate the memory the memory allocated using functions malloc or Calloc are not deallocated on their own. hence the free method is used whenever the dynamic memory allocation takes place. It helps to reduce the wastage of memory by freeing it.

free (ptr);

Realloc

ptr = & realloc (ptr, newsize)

for exam.

ptr = realloc (ptr, n * size of (int))

Realloc means reallocation method is used to dynamically change the memory allocation of previously assigned memory in other words if memory previously allocated with the help of malloc or Calloc is insufficient can be used to dynamically reallocate memory.

```

#include <stdio.h>
#include <stdlib.h>
int main()
{
    int *ptr;
    int n, i;
    n = 5;
    printf("Enter number of elements %d", n);
    ptr = (int *) malloc(n, sizeof(int));
    if (ptr == NULL)
    {
        printf("Memory not allocated");
        exit(0);
    }
    else
    {
        for (i = 0; i < n; i++)
        {
            ptr[i] = i + 1;
        }
        n = 10;
        printf("Enter the number of elements %d", n);
        ptr = realloc(ptr, n * sizeof(int));
        for (i = 5; i < n; i++)
        {
            ptr[i] = i + 1;
        }
        for (i = 0; i < n; i++)
        {
            printf("%d", ptr[i]);
        }
    }
}

```

func (P+8);
return 0;

}

Output - 12345

12345678910

++

--

+ with integer

- with integer

ptr1 + ptr2 ← not applicable

ptr1 - ptr2. ←

2-D array

pointers to pointers

Put arr [3][3]; → 0,1

arr →	2	3	4
arr+i →	5	6	7
arr+i+j →	8	9	1

* (* (arr+i)+j)

7 = * (* (arr+1)+2)

9 = * (* (arr+2)+1).

5 = * (* (arr+1)+0)

* (* arr) = 2.

* (* arr + 1) ⇒ * ((* arr + 0) + 1). = 3

⇒ * (* arr + 1) = 3

Link list

Traversing of link list

Traverse [PTR, INFO, LINK]

Step 1: PTR = START

Step 2: Repeat step 3 & 4 while PTR ≠ NULL.

Step 3: Apply process to INFO [PTR]

Step 4: Set PTR = LINK [PTR]
[End of while loop]

Step 5: Exit.

Insertion at the beginning

Insert [INFO, LINK, START, VALUE, AVAIL]

Step 1: if AVAIL = NULL, hint overflow.

Step 2: Set New = AVAIL

Step 3: AVAIL = LINK [AVAIL]

Step 4: INFO [New] = VALUE.

Step 5: Set LINK [New] = START.

Step 6: Set START = NEW

Step 7: Exit

creation:

```
#include <stdio.h>
#include <stdlib.h>
struct node
{
    Put info;
    struct node * link
};
```

```
struct node * first;
```

```
void main()
```

```
{ void Create();
```

```
Create(); void printf("data entered");
```

```
void creat()
```

```
struct node * ptr, * cpt.
```

```
char ch;
```

```
ptr = (struct node *) malloc (sizeof(struct node))
```

```
printf ("Input fixed first node");
```

```
scanf ("%d", & ptr->info);
```

```
first = ptr;
```

```
do
```

```
{
```

```
cpt = (struct node *) malloc (sizeof (struct node));
```

```
printf ("Input next node in link list");
```

```
scanf ("%d", & cpt->info);
```

```
ptr->link = cpt;
```

```
ptr = cpt;
```

```
printf ("Press 'Y/N' for more node");
```

ch = getch();

while (ch == 'y').

ptr → link = NULL;

}

void display()

struct node * ptr;

printf ("Element of link list");

ptr = first;

while (ptr != NULL)

{

printf (" %d ", ptr → info);

ptr = ptr → link;

}

}

→ Insert as a link is given location

Insert LOC (INFO, LINK, START, AVAIL, LOC, VALUE)

Step 1: If AVAIL = NULL Print "Overflow".

Step 2: Set NEW = Avail and Avail = Link[Avail]

Step 3: Set INFO[NEW] = Value

Step 4: If LOC = NULL then Insert as a first node
Set LINK[NEW] = Start and Start = NEW
else

Set LINK[NEW] = LINK[LOC]

LINK[LOC] = NEW

[End of structure]

Step 5: EXIT

Program

```
void Insert - beg()
{
    struct node *ptr;
    ptr = (struct node *) malloc (sizeof(struct node));
    if (ptr == NULL)
        printf ("Overflow");
    else
    {
        printf ("Input new node");
        scanf ("%d", &ptr->info);
        ptr->link = first;
```

~~X~~ Write an algorithm for insert a node in a linklist at the end.

Step 1: If $\text{AVAIL} = \text{NEW}$ print "Overflow".

Step 2: Set $\text{NEW} = \text{Avail}$ and $\text{Avail} = \text{LINK}[\text{Avail}]$

Step 3: Set $\text{INFO}[\text{NEW}] = \text{value};$

Step 4: If $\text{LOC} = \text{NULL}$ set $\text{LINK}[\text{NEW}] = \text{start}$ and $\text{start} = \text{NEW}$

else set $\text{LINK}[\text{NEW}] = \text{NULL}$

$\text{LINK}[\text{LOC}] = \text{NEW}$

(End if structure)

Exit



Deletion in a link list :-

DEL linklist(INFO, START, LINK, AVAIL, LOC, LOC

Step 1: If $\text{LOC} = \text{NULL}$ (i.e. i=4)

Set $\text{start} = \text{LINK}[\text{start}]$

else

set $\text{LINK}[\text{loc}] = \text{LINK}[\text{loc}]$

[End if structure]

Step 2: [From deleted node to Avail link]

Set $\text{LINK}[\text{loc}] = \text{Avail} - \text{Avail} - \text{loc}$

Step 3: Exit

This algorithm deletes the node ~~of~~ with
location loc.

LOC is the location of the node which
preceeds ~~at~~ n.

Program

insertion at given LOC

```
void insert_given_node()
```

```
{
```

```
struct node * ptr, * cpt;
```

```
int data;
```

```
ptr = (struct node *) malloc (size of (struct node));
```

```
if (ptr == NULL)
```

```
{
```

```
printf ("overflow");
```

```
return;
```

```
}
```

```
printf ("Input Newnode Information");
```

```
scanf ("%d", &ptr->info);
```

```
printf ("Input Information of node after  
which insertion will be made");
```

```
scanf ("%d", &data)
```

```
cpt = first;
```

```
while (cpt->info != data)
```

```
{
```

```
cpt = cpt->link;
```

$\text{ptr} \rightarrow \text{link} = \text{cpt} \rightarrow \text{link};$

$\text{cpt} \rightarrow \text{link} = \text{ptr};$

g -