

## **Symbiosis Institute of Technology**

## **Department of Computer Science and Engineering**

## **Academic Year 2025-26**

## Design Analysis of Algorithm-Lab

## Batch 2023-27 - Sem V

Lab Assignment No:- 5	
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Batch	2023-27
Class	CSE- A3
Academic Year & Semester	2025-26 TY, 5 <sup>th</sup> sem
<b>Date of Submission</b>	15 Sept 2025,
Title of Assignment:	You're packing a suitcase for a flight, and the airline allows a maximum weight of 15 kg. You have a set of items to pack, each with a weight and a utility value (how useful it is to you on the trip). But you can't take all items — so you need to choose the most valuable combination without exceeding 15 kg.
	Item Weight (kg) Value (usefulness) Laptop 3 9 Headphones 1 5 Jacket 5 10 Camera 4 7 Book 2 4 Shoes 6 6

#### Problem:

Choose the combination of items that gives maximum value but keeps the total weight ≤ 15 kg.

This is the 0/1 Knapsack Problem, because you can either take an item (1) or leave it (0) — no partial items.

You should be able to change the maximum allowed weight, weights of different items. You can also include new items, delete items given in the example.

#### **Theory: (Handwritten)**

#### Fractional Knapsack (Greedy)

- Idea: Sort by value/weight ratio.
- Time: O(n log n).
- Space: O(1).

#### 0/1 Knapsack (Dynamic Programming)

- Idea: Use DP table dp[i][w] = max value with first i items and capacity w.
- Time: O(nW).
- Space: O(nW).

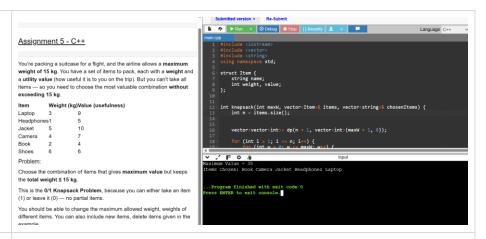
#### **Optimization:**

- Fractional: early break when bag is full.
- 0/1: Space optimization to O(W) possible.

#### Source code

```
if (items[i-1].weight <= w) {</pre>
         dp[i][w] = max(dp[i-1][w],
                  dp[i-1][w - items[i-1].weight] + items[i-1].value);
       } else {
         dp[i][w] = dp[i-1][w];
       }
    }
  }
  int w = maxW;
  for (int i = n; i > 0; i--) {
    if (dp[i][w] != dp[i-1][w]) {
       chosenItems.push_back(items[i-1].name);
       w -= items[i-1].weight;
    }
  }
  return dp[n][maxW];
}
int main() {
  int maxW = 15;
  vector<Item> items = {
    {"Laptop", 3, 9},
    {"Headphones", 1, 5},
    {"Jacket", 5, 10},
    {"Camera", 4, 7},
    {"Book", 2, 4},
    {"Shoes", 6, 6}
  };
  vector<string> chosenItems;
  int maxValue = knapsack(maxW, items, chosenItems);
  cout << "Maximum Value = " << maxValue << endl;</pre>
  cout << "Items chosen: ";</pre>
  for (auto &it : chosenItems) cout << it << " ";
  cout << endl;
  return 0;
```



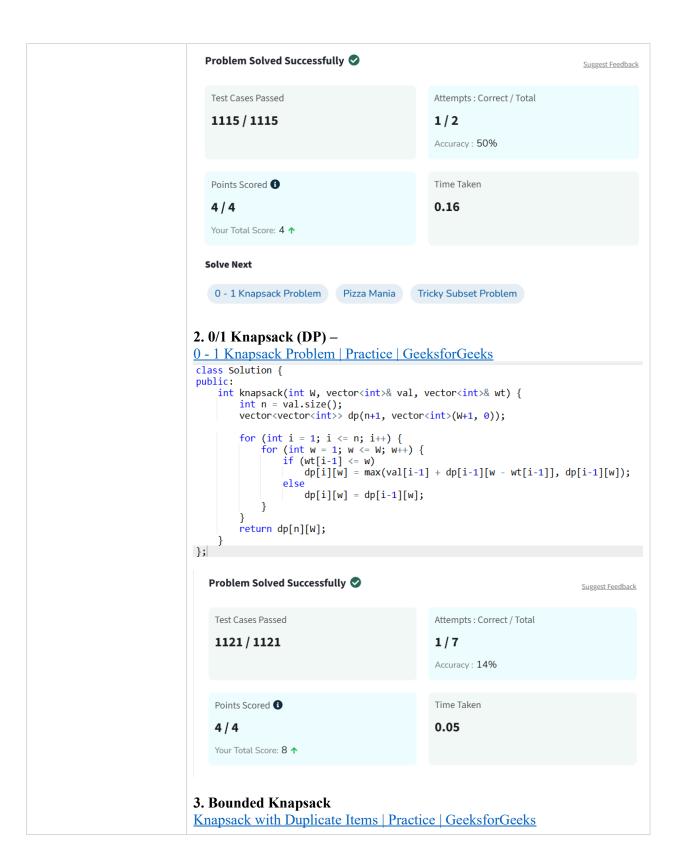


# Problems Solved from GeeksforGeeks

#### 1. Fractional Knapsack (Greedy) –

Fractional Knapsack | Practice | GeeksforGeeks

```
1 - class Solution {
 2 public:
          double fractionalKnapsack(vector<int>& val, vector<int>& wt, int capacity) {
3 +
4
               int n = val.size();
5
               // Pair each item as {ratio, index}
6
               vector<pair<double, int>> ratio;
for (int i = 0; i < n; i++) {
   double r = (double)val[i] / wt[i];</pre>
8 -
9
10
                    ratio.push_back({r, i});
11
12
               // Sort by ratio (highest first)
13
               sort(ratio.rbegin(), ratio.rend());
14
15
               double totalValue = 0.0;
16
17
18 -
               for (int i = 0; i < n; i++) {
19
                    int idx = ratio[i].second;
20
                    if (wt[idx] <= capacity) {
   // Take the whole item</pre>
21 -
22
                         totalValue += val[idx];
capacity -= wt[idx];
23
24
                    } else {
    // Take fraction of item
25 -
26
                         totalValue += (double)val[idx] * ((double)capacity / wt[idx]);
break; // bag is full
27
28
29
30
31
32
               return totalValue;
33
34 };
```



```
1 * class Solution {
  2 public:
 3 =
          int knapSack(vector<int>& val, vector<int>& wt, int capacity) {
 4
               int n = val.size();
 5
               vector<int> dp(capacity + 1, 0);
  6
 7 =
               for (int w = 0; w <= capacity; w++) {</pre>
 8 -
                   for (int i = 0; i < n; i++) {
                        if (wt[i] <= w) {
    dp[w] = max(dp[w], val[i] + dp[w - wt[i]]);</pre>
 9 =
 10
 11
                   }
 12
 13
 14
               return dp[capacity];
 15
 16 };
17
  Problem Solved Successfully
                                                                             Suggest Feedback
   Test Cases Passed
                                               Attempts: Correct / Total
   1111 / 1111
                                               1/1
                                               Accuracy: 100%
   Points Scored 1
                                               Time Taken
   4/4
                                               0.02
   Your Total Score: 12 ^
4. Subset Sum Problem
\underline{ \begin{array}{c|c} \textbf{Subset Sum Problem} \mid \textbf{Practice} \mid \textbf{GeeksforGeeks} \\ \textbf{class Solution} & \{ \end{array} }
public:
     bool isSubsetSum(vector<int>& arr, int sum) {
          int n = arr.size();
          vector<vector<bool>>> dp(n+1, vector<bool>(sum+1, false));
          for (int i = 0; i \le n; i++)
               dp[i][0] = true;
for (int i = 1; i <= n; i++) {
               for (int s = 1; s \le sum; s++) {
                    if (arr[i-1] <= s)
                        dp[i][s] = dp[i-1][s] || dp[i-1][s - arr[i-1]];
                        dp[i][s] = dp[i-1][s];
          return dp[n][sum];
     }
};
```

