



Angr & Pwn

vulnerability detection using symbolic execution

by malweisse



About me

Andrea Fioraldi [@andreafioraldi, @malweisse on IRC]

22 years old, student of Engineering in Computer Science.

Interested in binary analysis (symbolic execution, reversing and other pretty stuffs) and binary exploitation (do u know what is double free?).

Capturing flags with TheRomanXpl0it and mHACKeroni.

Not so much skilled trumpeter, mountain bike and trekking lover, Dragon Ball fanboy, homebrewer.





Angr crash course

Angr and the bugs



Concrete execution

- Single flow of instructions (path) at time
- Can't evaluate program behaviour

```
int foo(int a, int b)
       int c = 77:
       if(a + b == 42) {
            c = c - b:
       else {
            c = a - c;
10
       if(c == 38) {
           puts("Well done.");
14
       else {
15
           puts("Try again.");
16
17 }
```



- Explore all possible paths in a program
- Evaluate how inputs affect the choice of a path
- Use symbolic values as inputs
- Logical expression in function of the symbolic inputs (path constraints and symbolic storage)
- Execution forked on branches
- SMT solver to evaluate that expressions

```
int foo(int a. int b)
        int c = 77:
       if(a + b == 42) {
            c = c - b:
        else {
            c = a - c;
10
       if(c == 38) {
12
            puts("Well done.");
13
14
        else {
15
            puts("Try again.");
16
17
```



```
    [Path #1] Formula: X + Y = 42,
    Storage: c = 77
```

• [Path #2] Formula: ¬(X + Y = 42), Storage: c = 77

Target: line 12

```
int foo(int a, int b)
       int c = 77:
       if(a + b == 42) {
           c = c - b;
8
       else {
9
           c = a - c;
10
       if(c == 38) {
           puts("Well done.");
       else {
15
           puts("Try again.");
16
17 }
```



- [Path #1.1] Formula: X + Y = 42 ∧ c =
 38, Storage: c = 77 Y
- [Path #1.2] Formula: X + Y = 42 ∧ c!=
 38, Storage: c = 77 Y
- [Path #2.1] Formula: ¬(X + Y = 42) ∧ c
 = 38, Storage: c = X − 77
- [Path #2.2] Formula: ¬(X + Y = 42) ∧ c
 != 38, Storage: c = X − 77

The interesting paths are 1.1 and 2.1 because they reached our target

```
int foo(int a. int b)
       int c = 77:
       if(a + b == 42) {
            c = c - b:
        else {
            c = a - c;
10
11
       if(c == 38) {
12
            puts("Well done.");
13
14
        else {
15
            puts("Try again.");
16
17 }
```



[Path #1.1] Formula: X + Y = 42 ∧ c =
 38, Storage: c = 77 - Y

```
solve(X + Y = 42 \land 77 - Y = 38)
```

$$X = 3, Y = 39$$

```
int foo(int a, int b)
       int c = 77:
       if(a + b == 42) {
           c = c - b;
       else {
           c = a - c;
10
       if(c == 38) {
12
           puts("Well done.");
13
14
       else {
15
           puts("Try again.");
16
17 }
```



Symbolic execution (what about?)

- Symbolic pointers dereference?
- Loops?
- Exponential increase of the number of paths?
- Environment interaction?
- Non-linear constraints?





Angr and the bugs



Angr WTF

Angr is a binary analysis framework.

Binary loader, emulator, symbolic executor.

http://angr.io/

python3 -m pip install angr





Angr modules

- CLE, a multi-format binary loader with an intuitive API;
- archinfo, a collection of classes that contain architecture-specific information;
- PyVEX, a wrapper around Valgrind's VEX IR lifter, used to make the analyses architecture-agnostic;
- *Claripy*, the angr data backend, a wrapper around the Z3 solver and an interface to abstract concrete and symbolic values handling;



Loading a binary

```
import angr
# load the example
project = angr.Project("./foo")
```



Simulation

```
# start a new SimulationManager
simgr = project.factory.simulation_manager()
# step
simgr.step()
# step until it branches
simgr.run(until=lambda sm: len(sm.active) != 1)
# check the states that are still active
print (simgr.active)
```



States

```
state = simgr.active[0]

# a state has plugins, representing registers, memory, etc
print (state.regs.rax)
print (state.memory.load(state.regs.rsp, 8))
# one of the plugins represents the system state
print (state.posix.fd)
# files are backed by a memory region
print (state.posix.fd[0].read_data(8)) #return (value, size)
```



Solver

```
state = simgr.active[0]
# SMT solver
print (state.solver)
addr = state.regs.rsp + 0x100
# each value in angr is represented as an expression tree
v = state.memory.load(addr, 8) + 0x10
print (v)
print (v.op) # add
print (v.args) # (other bitvector, 0x10)
# add state constraints
state.add constraints(v < 0x2a)</pre>
state.add constraints(v > 0x1a)
# concretize a value
print (hex(state.solver.eval(v)))
```



Simulation 2

```
# create a state starting on foo()
initial state = project.factory.blank state()
initial state.regs.rip = 0x4005c7 # foo address
# create symbolic args
a = initial state.solver.BVS("sym a", 64)
b = initial state.solver.BVS("sym b", 64)
initial state.regs.rdi = a
initial state.regs.rsi = b
# start a new SimulationManager based on initial state
simgr = project.factory.simulation manager(initial state)
# explore using conditions
simgr.explore(find=0x400600, avoid=0x40060e) # well done, try again
print (simgr, simgr.found) # found stash
# conretize found inputs
print (a.args[0], "=", simgr.found[0].solver.eval(a))
print (b.args[0], "=", simgr.found[0].solver.eval(b))
```



Default Stashes

- active, states that are "live"
- errored, states that errored out (actual angr bugs)
- **found**, states that reached the "find" condition
- avoided, states that hit the "avoid" condition
- deadended, states that produced no successors
- unconstrained, states that jump to unconstrained symbolic values



Symbolic execution

Angr crash course

Angr and the bugs



Stack Buffer Overflow

What is the purpose of a stack buffer overflow?

Control the program counter value overwriting the return address on the stack



Stack Buffer Overflow

DEMO



Thank you!



QUESTIONS?



References



USEFUL PAPERS:

- R. Baldoni, E. Coppa, D. C. D'Elia, C. Demetrescu, and I. Finocchi, "A survey of symbolic execution techniques,"
 ACM Comput. Surv., vol. 51, no. 3, 2018
- Y. Shoshitaishvili, R. Wang, C. Salls, N. Stephens, M. Polino, A. Dutcher, J. Grosen, S. Feng, C. Hauser, C. Kruegel, and G. Vigna, "SoK: (State of) The Art of War: Offensive Techniques in Binary Analysis," in IEEE Symposium on Security and Privacy, 2016.

BRIEF SUMMARY (the chapter 1 of my thesis):

• https://www.researchgate.net/publication/327655380 Symbolic Execution and Debugging Synchronization

MORE ABOUT ANGR

- https://docs.angr.io/
- https://www.blackhat.com/docs/us-15/materials/us-15-Kruegel-Using-Static-Binary-Analysis-To-Find-Vulnerabiliti
 es-And-Backdoors-In-Firmware.pdf