Self Balancing Binary Tree

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Special Assignment: 2CSOE52 Data Structure

Self Balancing Binary Trees

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1 Introduction

Self balancing binary trees are height balanced tree which keeps its height as minimum as possible when a insertion or deletion operations are performed. Typically the height of the tree is kept in range of $\log(n)$ so that all operation cost $O(\log(n))$ average.

This project compares the three major self-balancing BST and determines the specific application of each of the algorithms and its possible real world application. The three self-balancing BST are;

- Red Black Tree
- AVL Tree
- Splay Tree

2 Red Black Trees

It is BST with an extra attribute for each node. The extra attribute is the colour. The colour of the node can either be black or red. So a structure of red black tree could be as shown;

```
1 struct RB_BST_node {
2    int data;
3    int color;
4    struct RB_BST_node* parent;
5    struct RB_BST_node* right;
6    struct RB_BST_node* left;
7    };
```

The properties of the Red-Black trees are;

- Root note is always black
- Every node is either red of black
- Every leaf which is NULL is black
- A red node cannot have red children

• Every simple path from node to descendant(NULL) contains the same number of black nodes.

2.1 Algorithm

Search Algorithm

```
search (tree, value)

Step #1 :

If tree -> data = value OR tree = NULL

Return tree

Else If value < data

Return search(tree -> left , value)

Else

Return search(tree -> righ , value)

END
```

Insert Algorithm

```
If tree = NULL

Insert the new value as root with color Black

Else

Insert new value as leaf with color red

If parent of new is black

exit

else

Check the color of parent sibling

if color of parent sibling = black

perform suitable rotation

else

change the color and recheck

END
```

2.2 Results

The code for the Red Black Tree implementation is in Appendix 5.1.

Figure 1: Result for RB tree

2.3 Application

The Major application of the Red Black tree is in Multi-Set, Multi-Map, Map and Set. A recent article also explores the use of Red-Black tree in Wireless sensor network in order the achieve substantial efficiency in network life time and energy efficiency.

3 AVL Trees

AVL tree is a self-balancing Binary Search Tree (BST) where the difference between heights of left and right subtrees cannot be more than one for all nodes.

Structure of AVL node

```
Struct AVLNode

int data;

struct AVLNode *left, *right;

int balfactor;

};
```

The important feature of a AVL tree is that the difference between the depth of right and left subtrees cannot be more than one. In order to maintain this feature, an implementation of an AVL will include an algorithm to rebalance the tree when adding an additional element would

upset this feature.

3.1 Algorithm

Insert Algorithm

```
Insert the element using BST insertion

Check for the balance factor of each node

If balance factor = [-1,0,1]

Proceed to next operation

Else

The tree is imbalanced, perform suitable rotations to balance

END
```

3.2 Results

The code for the AVL Tree is in Appendix 5.2

```
(base) Dehits-MacBook-Air:~ dehittrivedi$ /var/folders/m1/zy3xgbxs18gbs8sw35g67jhw0000gn/T/18BEC119_DSSA_B ; exit;
Insert elements

Number of key: 4
Enter key: 1
Enter key: 2
Enter key: 3
Enter key: 3
Enter key: 4
Preorder traversal of the constructed AVL tree is
2 1 3 4
-------
The elapsed time is 0.000459 secondslogout
```

Figure 2: Result for AVL tree

3.3 Application

The major use of AVL tree in database transactions and in memory sorts of set and dictionaries.

4 Conclusion

This project helps to understand different types of self balancing binary tree. The implementation of 2 common self-balancing BST was demonstrated in the project. The splay tree are also

Self-balancing BST however due to its complexity and lack of resources its implementation was not covered in the project.

The AVL and Red-Black both of them are complex algorithm however this algorithm help to achieve better time and space complexity. This algorithm brings the average complexity to $O(\log_2(n))$.

5 Appendix

5.1 CODE FOR RED-BLACK TREE

```
1 /*
2 Data Structure Special assignment
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4 RED - BLACK TREE
6 */
8 #include <stdio.h>
9 #include <time.h>
#include <stdlib.h>
12 #define RED 0
13 #define BLACK 1
15 struct node{
int key;
int color;
struct node *parent;
struct node *left;
20 struct node *right;
21 };
23 struct node *ROOT;
24 struct node *NILL;
```

```
26 void left_rotate(struct node *x);
27 void right_rotate(struct node *x);
28 void tree_print(struct node *x);
29 void red_black_insert(int key);
30 void red_black_insert_fixup(struct node *z);
31 struct node *tree_search(int key);
32 struct node *tree_minimum(struct node *x);
void red_black_transplant(struct node *u, struct node *v);
34 void red_black_delete(struct node *z);
35 void red_black_delete_fixup(struct node *x);
37 int main(){
   double time_spent = 0.0;
  clock_t begin = clock();
  NILL = malloc(sizeof(struct node));
  NILL->color = BLACK;
  ROOT = NILL;
  printf("Insert elements\n\n");
  int i, key;
  printf("Number of key: ");
  scanf("%d", &i);
  while(i--){
   printf("Enter key: ");
   scanf("%d", &key);
   red_black_insert(key);
54
55
   printf("Tree\n\n");
  tree_print(ROOT);
  printf("\n");
```

```
oprintf("Enter a key to be searhced\n\n");
         printf("Enter key: ");
         scanf("%d", &key);
         printf((tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL) ? "NILL \n" : "%p \n", tree\_search(key) == NILL \n", tree\_search(key) == NILL \n", tree\_search(key) ==
                  ));
          printf("minimum test\n\n");
          printf("MIN: %d\n", (tree_minimum(ROOT))->key);
          printf("Tree \n\n");
          tree_print(ROOT);
          printf("\n");
             printf("\n");
            clock_t end = clock();
             time_spent += (double)(end - begin) / CLOCKS_PER_SEC;
             printf("----\n");
              printf("The elapsed time is %f seconds", time_spent);
78 return 0;
79 }
81 void tree_print(struct node *x){
82 if(x != NILL){
           tree_print(x->left);
         printf("%d\t", x->key);
         tree_print(x->right);
86 }
87 }
89 struct node *tree_search(int key){
90 struct node *x;
92 x = ROOT;
93 while(x != NILL && x->key != key){
```

```
if(key < x->key){
    x = x -> left;
    }
    else{
   x = x->right;
   }
100
102 return x;
103 }
struct node *tree_minimum(struct node *x){
while(x->left != NILL){
  x = x -> left;
108 }
109 return x;
110 }
void red_black_insert(int key){
struct node *z, *x, *y;
z = malloc(sizeof(struct node));
115
z \rightarrow key = key;
z \rightarrow color = RED;
118 z->left = NILL;
z->right = NILL;
120
x = ROOT;
   y = NILL;
123
124 /*
   * Go through the tree untill a leaf(NILL) is reached. y is used for
    * track of the last non-NILL node which will be z's parent.
127 */
```

```
while(x != NILL){
     y = x;
129
    if(z\rightarrow key \leftarrow x\rightarrow key){
     x = x \rightarrow left;
131
     }
132
     else{
     x = x->right;
134
     }
135
136
137
   if(y == NILL){
    ROOT = z;
139
140
   else if(z\rightarrow key <= y\rightarrow key){
141
    y \rightarrow left = z;
142
   }
143
   else{
144
    y \rightarrow right = z;
145
   }
146
147
   z->parent = y;
149
   red_black_insert_fixup(z);
150
151 }
153 void red_black_insert_fixup(struct node *z){
   while(z->parent->color == RED){
154
155
     /* z's parent is left child of z's grand parent*/
     if(z->parent == z->parent->parent->left){
157
158
159
      /* z's grand parent's right child is RED */
      if(z->parent->right->color == RED){
160
       z->parent->color = BLACK;
161
       z->parent->parent->right->color = BLACK;
162
```

```
z->parent->parent->color = RED;
      z = z->parent->parent;
164
166
     /* z's grand parent's right child is not RED */
167
     else{
169
      /* z is z's parent's right child */
      if(z == z->parent->right){
171
       z = z->parent;
172
       left_rotate(z);
      }
174
      z->parent->color = BLACK;
176
      z->parent->parent->color = RED;
177
      right_rotate(z->parent->parent);
     }
179
    }
180
181
    /* z's parent is z's grand parent's right child */
182
    else{
183
184
     /* z's left uncle or z's grand parent's left child is also RED */
185
     if(z->parent->left->color == RED){
      z->parent->color = BLACK;
187
      z->parent->parent->left->color = BLACK;
      z->parent->parent->color = RED;
189
      z = z->parent->parent;
192
     /* z's left uncle is not RED */
193
194
     else{
      /* z is z's parents left child */
195
      if(z == z->parent->left){
       z = z - parent;
197
```

```
right_rotate(z);
199
      z->parent->color = BLACK;
201
      z->parent->parent->color = RED;
202
      left_rotate(z->parent->parent);
     }
204
    }
205
206
207
   ROOT -> color = BLACK;
209 }
void left_rotate(struct node *x){
   struct node *y;
   /* Make y's left child x's right child */
   y = x->right;
   x \rightarrow right = y \rightarrow left;
   if(y->left != NILL){
   y->left->parent = x;
   }
219
   /* Make x's parent y's parent and y, x's parent's child */
   y->parent = x->parent;
222
   if(y->parent == NILL){
   ROOT = y;
224
225
   else if(x == x->parent->left){
    x->parent->left = y;
227
   }
228
229
   else{
   x->parent->right = y;
230
232
```

```
233 /* Make x, y's left child & y, x's parent */
y - > left = x;
x - parent = y;
236 }
237
238 void right_rotate(struct node *x){
   struct node *y;
   /* Make y's right child x's left child */
241
y = x - > left;
   x \rightarrow left = y \rightarrow right;
244 if (y->right != NILL) {
   y->right->parent = x;
246
247
   /* Make x's parent y's parent and y, x's parent's child */
   y->parent = x->parent;
   if(y->parent == NILL){
   ROOT = y;
252
   else if(x == x->parent->left){
   x->parent->left = y;
254
   }
255
   else{
   x->parent->right = y;
257
   }
258
   y \rightarrow right = x;
x - parent = y;
261 }
```

5.2 CODE FOR AVL TREE

```
/*
2 Data Structure Special assignment
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```

```
4 AVL TREE : Insertion
5 */
7 #include < stdio.h>
8 #include <time.h>
9 #include < stdlib.h>
11 struct Node
int key;
    struct Node *left;
    struct Node *right;
int height;
17 };
int max(int a, int b);
20 int height(struct Node *N)
21 {
   if (N == NULL)
        return 0;
return N->height;
25 }
27 int max(int a, int b)
28 {
   return (a > b)? a : b;
30 }
32 struct Node* newNode(int key)
33 {
     struct Node* node = (struct Node*)
35
                        malloc(sizeof(struct Node));
     node->key = key;
     node->left = NULL;
    node->right = NULL;
```

```
node->height = 1;
      return(node);
40
41 }
42
43 struct Node *rightRotate(struct Node *y)
44 {
      struct Node *x = y->left;
45
      struct Node *T2 = x->right;
47
      x \rightarrow right = y;
48
      y \rightarrow left = T2;
50
      y->height = max(height(y->left), height(y->right))+1;
      x->height = max(height(x->left), height(x->right))+1;
52
53
      return x;
55 }
56 struct Node *leftRotate(struct Node *x)
      struct Node *y = x->right;
      struct Node *T2 = y->left;
60
      y \rightarrow left = x;
      x \rightarrow right = T2;
62
63
      x->height = max(height(x->left), height(x->right))+1;
      y->height = max(height(y->left), height(y->right))+1;
65
      return y;
67
68 }
70 int getBalance(struct Node *N)
71 {
      if (N == NULL)
          return 0;
```

```
return height(N->left) - height(N->right);
75 }
77 struct Node* insert(struct Node* node, int key)
78 {
      if (node == NULL)
          return(newNode(key));
82
      if (key < node->key)
83
          node->left = insert(node->left, key);
      else if (key > node->key)
85
          node->right = insert(node->right, key);
      else
87
          return node;
      node->height = 1 + max(height(node->left),
90
                               height(node->right));
92
      int balance = getBalance(node);
93
      if (balance > 1 && key < node->left->key)
95
          return rightRotate(node);
      if (balance < -1 && key > node->right->key)
98
          return leftRotate(node);
100
      if (balance > 1 && key > node->left->key)
101
      {
          node->left = leftRotate(node->left);
103
          return rightRotate(node);
      }
105
106
      if (balance < -1 && key < node->right->key)
      {
108
```

```
node->right = rightRotate(node->right);
           return leftRotate(node);
110
      }
112
      return node;
113
114 }
115
void preOrder(struct Node *root)
117 {
      if(root != NULL)
      {
           printf("%d ", root->key);
120
           preOrder(root->left);
           preOrder(root->right);
      }
123
124 }
126 int main()
127 {
    double time_spent = 0.0;
128
    clock_t begin = clock();
    struct Node *root = NULL;
130
131
    printf("Insert elements\n\n");
    int i, key;
134
   printf("Number of key: ");
135
   scanf("%d", &i);
   while(i--){
    printf("Enter key: ");
138
    scanf("%d", &key);
    root = insert(root, key);
140
   }
141
   root = insert(root, 10);
```

```
root = insert(root, 20);
    root = insert(root, 30);
145
    root = insert(root, 40);
147
    root = insert(root, 50);
    root = insert(root, 25);
148
    */
149
    \verb|printf("Preorder traversal of the constructed AVL"|\\
150
            " tree is \n");
    preOrder(root);
152
    printf("\n");
153
    clock_t end = clock();
    time_spent += (double)(end - begin) / CLOCKS_PER_SEC;
155
    printf("----\n");
    printf("The elapsed time is %f seconds", time_spent);
157
    return 0;
158
159 }
```