



## 第二章 广义函数的卷积

### 函数与函数之间的卷积

定义	设 $f(x), g(x) \in C(\mathbb{R}^n)$ 且至少一个具有紧支集, 定义卷积 $(f * g)(x) = \int_{\mathbb{R}^n} f(y)g(x-y)dy$
交换	证明: $f, g \in C(\mathbb{R})$ 且至少有一个具有紧支集 $(f * g)(x) = \int_{\mathbb{R}} f(y)g(x-y)dy \xrightarrow{t=x-y} \int_{\mathbb{R}} f(x-t)g(t)dt = (g * f)(x)$
结合	$f, g, h \in C(\mathbb{R})$ 且至少有两个具有紧支集 $(f * (g * h))(x) = \int_{\mathbb{R}} \int_{\mathbb{R}} f(s)g(t)h(x-s-t)dt ds$ $((f * g) * h)(x) = \int_{\mathbb{R}} \int_{\mathbb{R}} f(s)g(t-s)h(x-t)dt ds \xrightarrow{u=t-s} \int_{\mathbb{R}} \int_{\mathbb{R}} f(s)g(u)h(x-s-u)du ds$
微分	微分可以跨过卷积 设 $f(x) \in C(\mathbb{R}^n)$ 且至具有紧支集, 即 $f \in C_0(\mathbb{R}^n), g \in C^k(\mathbb{R}^n)$ , 则 $f * g \in C^k(\mathbb{R}^n)$ , 有 $\partial_x^\alpha (f * g) = f * (\partial_x^\alpha g),  \alpha  \leq k$ 按微分定义算 简单 证明: 不妨设 $x \in \mathbb{R}, \text{supp } f = K$ 是紧集 $\begin{aligned} \frac{d(f * g)(x)}{dx} &= \lim_{h \rightarrow \infty} \frac{(f * g)(x+h) - (f * g)(x)}{h} \\ &= \lim_{h \rightarrow \infty} \frac{1}{h} \left( \int_{\mathbb{R}} f(y)g(x+h-y) - f(y)g(x-y)dy \right) \\ &= \lim_{h \rightarrow \infty} \int_{\mathbb{R}} f(y) \frac{1}{h} (f(x+h-y) - g(x-y))dy \\ &\stackrel{LDT}{=} \int_{\mathbb{R}} f(y) \frac{dg(x-y)}{dx} dy \\ &= \left( f * \frac{dg}{dx} \right)(x) \end{aligned}$ 于是有 $\partial_x^\alpha (f * g) = f * \partial_x^\alpha g$
支集	设 $f * g$ 有意义, $\text{supp}(f * g) \subset \text{supp } f + \text{supp } g = \{x+y; x \in \text{supp } f, y \in \text{supp } g\}$ 考虑 $x, y$ 的邻域和 $\text{supp } g$ 不交是套路, 证明支集的性质都这么考虑 要证: $\forall x \notin \text{supp } f + \text{supp } g$ , 都有 $x \notin \text{supp}(f * g)$ $\forall x \notin \text{supp } f + \text{supp } g$ 为闭集, $\exists x$ 的邻域 $V_x, s.t. V_x \cap (\text{supp } f + \text{supp } g) = \emptyset$ $\forall y \in \text{supp } f$ , 有 $x-y \notin \text{supp } g$ 为闭集, 从而 $\exists x-y$ 的邻域 $V_{x-y}, s.t. V_{x-y} \cap \text{supp } g = \emptyset$ , 即 $\int_{V_{x-y}} = 0$ 从而 $(f * g)(x) = \int_{\text{supp } f} f(y)g(x-y)dy = 0$
定理	卷积的性质很好, 跟正则性高的那边走 设 $f \in L^1(\mathbb{R}^n), g \in L^p(\mathbb{R}^n), 1 < p < \infty$ , 则 $f * g \in L^p(\mathbb{R}^n)$ 且 $\ f * g\ _{L^p(\mathbb{R}^n)} \leq \ f\ _{L^1} \ g\ _{L^p}$

### 广义函数与函数之间的卷积

定义	设 $f(x) \in \mathcal{D}'(\mathbb{R}^n), g(x) \in C_0^\infty(\mathbb{R}^n)$ , 定义 $(f * g)(x) = \langle f(y), g(x-y) \rangle$
例子	$\delta$ 在广义函数与函数的卷积运算中相当于单位元 $\delta(x) \in \mathcal{D}'(\mathbb{R}^n), f \in C^\infty(\mathbb{R}^n), (\delta * f)(x) = \langle \delta(y), f(x-y) \rangle = f(x-0) = f(x)$
交换	$\mathcal{D}'$ 广义函数 $* C_0^\infty$ 函数, 由 $\langle \cdot, \cdot \rangle$ 定义, 没有交换
结合	积分号可以放到试验函数上来, 即积分号可以和广义函数作用换序 设 $\Omega \subset \mathbb{R}^n, \omega \subset \mathbb{R}^m$ 为开集, $\varphi \in C_0^\infty(\Omega \times \omega)$ 且 $\text{supp } \varphi \subset K_1 \times K_2$ , 这里 $K_1 \subset \Omega, K_2 \subset \omega$ 为紧集, 如果 $u \in \mathcal{D}'(\Omega)$ , 则 $\int \langle u, \varphi(\cdot, y) \rangle dy = \left\langle u, \int \varphi(\cdot, y) dy \right\rangle$ 证明思路: $\langle f(x), \varphi(x, y) \rangle_x \in C_0^\infty(W)$ , 从而 $S(y) := \langle f(y), \varphi(x, y) \rangle$ 可积 将 $\int \langle f(x), \varphi(x, y) \rangle_x dy$ 写成求和的极限, 即有: 当 $h$ 充分小时 $\begin{aligned} \sum_{k \in \mathbb{Z}} \langle f(x), \varphi(x, kh) \rangle h^m &\xrightarrow{h \rightarrow 0} \int \langle f(x), \varphi(x, y) \rangle dy \\ &= \left\langle f(x), \sum_{k \in \mathbb{Z}} \varphi(x, kh) h^m \right\rangle = \left\langle f(x), \int \varphi(x, y) dy \right\rangle \end{aligned}$
定理	若 $\varphi, \psi \in C_0^\infty(\mathbb{R}^n)$ , 且 $f \in \mathcal{D}'(\mathbb{R}^n)$ , 则 $(f * \psi) * \varphi = f * (\psi * \varphi)$ $\begin{aligned} (f * \psi) * \varphi &= \int (f * \psi)(y) \varphi(x-y) dy \\ &= \int \langle f(z), \psi(y-z) \rangle_z \varphi(x-y) dy \\ &= \int \langle f(z), \psi(y-z) \varphi(x-y) \rangle_z dy \\ &\stackrel{\text{引理}}{=} \left\langle f(z), \int \psi(y-z) \varphi(x-y) dy \right\rangle_z \\ &\stackrel{y-z=w}{=} \left\langle f(z), \int \psi(w) \varphi(x-z-w) dw \right\rangle \\ &= \langle f(z), (\psi * \varphi)(x-z) \rangle \\ &= f * (\psi * \varphi) \end{aligned}$
微分	微分号可以放到试验函数上来, 即积分号可以和广义函数作用换序 设 $\omega \subset \mathbb{R}^n$ 为开集, $\varphi(x, y) \in C^\infty(\Omega \times \omega)$ 且有紧集 $K \subset \Omega$ 满足当 $x \notin K$ 时有 $\varphi(x, y) = 0$ 对任意 $y \in \omega$ 成立. 如果 $u \in \mathcal{D}'(\Omega)$ , 则函数 $y \mapsto \langle u, \varphi(\cdot, y) \rangle$ 是一 $C^\infty$ 函数, 且 $\partial_y^\alpha \langle u, \varphi(\cdot, y) \rangle = \langle u, \partial_y^\alpha \varphi(\cdot, y) \rangle$ 证明: 令 $F(x) = \langle f(y), \varphi(x, y) \rangle$ , 下证 $\partial_x^\alpha F(x) = \langle f(y), \partial_x^\alpha \varphi(x, y) \rangle$ 当 $ h  = \sqrt{h_1^2 + \dots + h_n^2}$ 充分小时, $x \in U_x, x+h \in U_x, F(x+h) - F(x) = \langle f(y), \varphi(x+h) - \varphi(x, y) \rangle$ 其中, $\varphi(x+h) - \varphi(x, y) = \sum_{j=1}^n h_j \partial_{x_j} \varphi(x, y) + R(x, y, h), R(x, y, h) \in o( h ^2)$ 【多元的 Taylor 展开】 $\Rightarrow F(x+h) - F(x) = \left\langle f(y), \sum_{j=1}^n h_j \partial_{x_j} \varphi(x, y) + R(x, y, h) \right\rangle = \left\langle f(y), \sum_{j=1}^n h_j \partial_{x_j} \varphi(x, y) \right\rangle + \langle f(y), R(x, y, h) \rangle$ 其中, $ \langle f(y), R(x, y, h) \rangle  \leq c \sum_{ \alpha  \leq k} \sup_K  \partial_y^\alpha R(x, y, h)  = o( h ^2)$ , 而 $\frac{o( h ^2)}{ h } \xrightarrow{ h  \rightarrow 0} 0$ 证明: $F(x) = \langle f(y), \varphi(x, y) \rangle$ , 不妨设 $h = (0, \dots, 0, h_j, 0, \dots, 0),  h  =  h_j $ $\begin{aligned} \partial_{x_j} F(x) &= \lim_{h_j \rightarrow 0} \frac{F(x+h) - F(x)}{h_j} \\ &= \lim_{ h  \rightarrow 0} \left( \frac{1}{h_j} \sum_{j=1}^n h_j \langle f(y), \partial_{x_j} \varphi(x, y) \rangle + o( h ) \right) = \langle f(y), \partial_{x_j} \varphi(x, y) \rangle \end{aligned}$ 从而当 $f' \in \mathcal{D}'(\mathbb{R}^n), \varphi \in C_0^\infty(\mathbb{R}^n)$ 时, $\partial_x^\alpha (f * \varphi)(x) = f * (\partial_x^\alpha \varphi)(x), \forall \alpha \in \mathbb{N}^n$
定理	若 $f \in \mathcal{D}'(\mathbb{R}^n), \varphi \in C_0^\infty(\mathbb{R}^n)$ (或 $f \in \mathcal{D}'(\mathbb{R}^n), \varphi \in C^\infty(\mathbb{R}^n)$ ), 则 $\partial^\alpha (f * \varphi) = f * \partial^\alpha \varphi = (\partial^\alpha f) * \varphi$ $f * \partial_x^\alpha \varphi = \langle f(y), \partial_x^\alpha \varphi(x-y) \rangle = \langle f(y), (-1)^{ \alpha } \partial_y^\alpha \varphi(x-y) \rangle = (-1)^{ \alpha } \langle \partial_y^\alpha f(y), \varphi(x-y) \rangle = (\partial^\alpha f) * \varphi$
支集	$f \in \mathcal{D}', \varphi \in C_0^\infty \Rightarrow \text{supp}(f * \varphi) \subset \text{supp } f + \text{supp } \varphi$ 证明: 即证 $\forall x \in \text{supp } f + \text{supp } \varphi$ , 有 $x \notin \text{supp}(f * \varphi)$ 由于 $\text{supp } f + \text{supp } \varphi$ 是闭集, $\exists x$ 的邻域 $V_x, s.t. V_x \cap (\text{supp } f + \text{supp } \varphi) = \emptyset$ $\forall y \in \text{supp } \varphi, x-y \notin \text{supp } f$ , 则 $\exists x-y$ 的邻域 $V_{x-y}, s.t. V_{x-y} \cap \text{supp } f = \emptyset$ , 即 $\varphi _{V_{x-y}} = 0$ $\Rightarrow f * \varphi _{V_x} = \langle f(y), \varphi(x-y) \rangle = 0$
广义函数的正则化	$\forall f \in \mathcal{D}'(\mathbb{R}^n)$ , 都可由一列 $C^\infty(\Omega)$ 函数在 $\mathcal{D}'$ 的意义下去逼近 证明: 设 $f \in \mathcal{D}'(\mathbb{R}^n), f_\varepsilon = f * \Phi_\varepsilon \in C^\infty(\mathbb{R}^n)$ , 要证: $\forall \varphi \in C_0^\infty(\mathbb{R}^n), \langle f_\varepsilon, \varphi \rangle \xrightarrow{\varepsilon \rightarrow 0} \langle f, \varphi \rangle$ 关注到 $\langle f_\varepsilon * \tilde{\varphi} \rangle(0) = \langle f_\varepsilon(y), \tilde{\varphi}(x-y) \rangle(0) = \langle f_\varepsilon(y), \varphi(y-x) \rangle _{x=0} = \langle f_\varepsilon(y), \varphi(y) \rangle$ 从而 $\langle f_\varepsilon, \varphi \rangle = \langle f_\varepsilon * \tilde{\varphi} \rangle(0) = [(f * \Phi_\varepsilon) * \tilde{\varphi}](0) \stackrel{\text{结合律}}{=} [f * (\Phi_\varepsilon * \tilde{\varphi})](0) = \left\langle f * (\Phi_\varepsilon * \tilde{\varphi})^\vee \right\rangle \xrightarrow{\varepsilon \rightarrow 0} \langle f, \varphi \rangle$ 若 $S \in \mathcal{D}'(\Omega)$ , 则必可找到一串 $S_j \in \mathcal{D}'(\Omega), j = 1, 2, \dots, s.t. S_j \rightarrow S$ (于 $\mathcal{D}'(\Omega)$ 中) 证明略

### 广义函数与广义函数之间的卷积

定义	设 $f(x) \in \mathcal{D}', g \in \mathcal{D}'$ , 则可定义 $f$ 和 $g$ 的卷积 $f * g \in \mathcal{D}'$ 为: $\langle f * g, \varphi \rangle = \langle f(x), \langle g(y), \varphi(x+y) \rangle \rangle, \forall \varphi \in \mathcal{D}(\mathbb{R}^n)$ 注意: $f$ 与 $g$ 中至少有一个有紧支集, 对于两个一般的 $f, g \in \mathcal{D}'(\mathbb{R}^n)$ , 不能一般地定义其卷积
结合	2020考过 设 $f, g, h \in \mathcal{D}'(\mathbb{R}^n)$ , 且其中至少有两个具紧支集, 则 $(f * g) * h = f * (g * h)$ 证明: $\forall \varphi \in C_0^\infty(\mathbb{R}^n), \langle f * (g * h), \varphi \rangle = \langle f(x), \langle (g * h)(y), \varphi(x+y) \rangle \rangle$ $= \langle f(x), \langle g(y), \langle h(z), \varphi(x+y+z) \rangle \rangle \rangle$ 而考虑 $\langle (f * g) * h, \varphi \rangle = \langle (f * g)(x), \langle h(z), \varphi(x+z) \rangle \rangle = \langle f(x), \langle g(y), \langle h(z), \varphi(x+y+z) \rangle \rangle \rangle$
性质	要把函数看成广义函数, 利用性质两层嵌套证明, 2023考过 定理: $f \in \mathcal{D}', g \in \mathcal{D}',$ 则 $f * g = g * f$ ; (1) $\forall \varphi, \psi \in C_0^\infty(\mathbb{R}^n)$ , 则 $((f * g) * \varphi) * \psi = ((g * f) * \varphi) * \psi$ 证明: $((f * g) * \varphi) * \psi \stackrel{\text{广义函数之间卷积的结合律}}{=} (f * (g * \varphi)) * \psi$ $\stackrel{\text{广义函数与函数的卷积结合律}}{=} f * ((g * \varphi) * \psi)$ $\stackrel{\text{函数之间卷积可交换}}{=} f * (\psi * (g * \varphi))$ $\stackrel{\text{广义函数与函数的卷积结合律}}{=} (f * \psi) * (g * \varphi)$ $\stackrel{\text{函数之间卷积可交换}}{=} (g * \varphi) * (f * \psi)$ $\stackrel{\text{广义函数与函数卷积的结合律}}{=} g * (\varphi * (f * \psi))$ $\stackrel{\text{函数与函数卷积的交换律}}{=} g * ((f * \psi) * \varphi)$ $\stackrel{\text{广义函数与函数卷积的结合律}}{=} g * (f * (\psi * \varphi))$ $\stackrel{\text{广义函数之间卷积的结合律}}{=} (g * f) * (\varphi * \psi)$ $\stackrel{\text{广义函数与函数卷积的结合律}}{=} ((g * f) * \varphi) * \psi$ (2) 则 $((f * g) * \varphi) * \tilde{\psi}(0) = ((g * f) * \varphi) * \tilde{\psi}(0)$ $\Rightarrow \langle (f * g) * \varphi, \tilde{\psi} \rangle = \langle (g * f) * \varphi, \tilde{\psi} \rangle$ $\Rightarrow (f * g) * \varphi = (g * f) * \varphi$ 又相等. 则 $(f * g) * \tilde{\varphi}(0) = (g * f) * \tilde{\varphi}(0)$ $\Rightarrow \langle f * g, \varphi \rangle = \langle g * f, \varphi \rangle, \forall \varphi \in C_0^\infty(\mathbb{R}^n)$ . 则 $f * g = g * f$ 又相等 2020考过 证明: $\forall \varphi \in C_0^\infty(\mathbb{R}^n), \langle f * \delta, \varphi \rangle = \langle f(x), \langle \delta(y), \varphi(x+y) \rangle \rangle = \langle f(x), \varphi(x) \rangle$ $\Rightarrow f * \delta = f$ , 由交换性, $\delta * f = f$ 微分 $\begin{aligned} \partial^\alpha (f * g) &= (\partial^{\alpha_1} f) * (\partial^{\alpha_2} g), \forall \alpha = \alpha_1 + \alpha_2; \\ &\langle \partial^\alpha (f * g), \varphi \rangle, \forall \varphi \in C_0^\infty \\ &= (-1)^{ \alpha } \langle (f * g), \partial^\alpha \varphi \rangle \\ &= (-1)^{ \alpha } \langle f(x), \langle g(y), \partial^\alpha \varphi(x+y) \rangle \rangle \\ &= (-1)^{ \alpha } \langle f(x), \langle \partial^{\alpha_2} g(y), \partial^{\alpha_1} \varphi(x+y) \rangle \rangle \\ &= (-1)^{ \alpha } \langle f * \partial^{\alpha_2} g, \partial^{\alpha_1} \varphi \rangle \\ &= (-1)^{ \alpha } \langle \partial^{\alpha_2} g * f, \partial^{\alpha_1} \varphi \rangle \\ &= (-1)^{ \alpha } \langle \partial^{\alpha_2} g(x), \langle f(y), \partial^{\alpha_1} \varphi(x+y) \rangle \rangle \\ &= (\partial^{\alpha_1} f) * (\partial^{\alpha_2} g) \end{aligned}$ 2023伯苓考过 $\text{supp}(f * g) \subset \text{supp } f + \text{supp } g$ 证明思路还是类似的. 但用试验函数的支集 $\text{supp } \varphi$ 代替原先的点 $x$ 即证 $\forall \varphi \in C_0^\infty, \text{supp } \varphi \cap (\text{supp } f + \text{supp } g) = \emptyset$ , 一定有 $\varphi \cap \text{supp}(f * g) = \emptyset$ $\forall x \in \text{supp } f, \begin{cases} \forall y \in \text{supp } g \Rightarrow x+y \in \text{supp } f + \text{supp } g, \varphi(x+y) = 0, \\ \forall y \notin \text{supp } g, \langle g(y), \varphi(x+y) \rangle = 0, \end{cases}$ 于是 $\langle f * g, \varphi \rangle = \langle f(x), \langle g(y), \varphi(x+y) \rangle \rangle = 0$ 即 $\langle f * g, \varphi \rangle = 0, \Rightarrow \text{supp } \varphi \cap \text{supp}(f * g) = \emptyset$
线性	卷积运算关于每个因子是线性的.