University of Alberta 2005 ACM ICPC World Finals Code Archive

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| ٠, | lden section Search | e. Generate | olex Method for | | ο. | (1) | Number Theory: Frimality testing with a sieve */ | Theory: Prime Factorization | Number of Divisors; | Primality | : Eul | ellaneous: Bit | Template: IO Reference | in lines of input until EOF | Number Theory: Converting between bases (arbitrary procision) | Theory: Chinese Postman Pr | _Theory: Maximur | н | y: Maximum Bipartite Matchir | : Binary Strings gene | | Dynamic Programming: Integer Parititoning | torics - Permutation i | ٠. | Digits in N! . | | Arithmetic: Binomial coefficient */ | č. | Cubic equation | : Simpson's Rule for | | ithmetic structures */ . | | 3D Geometry Primitives */ | bounding circle */ | n midpoints -> vertices (n | dearthed by three | | - e | (return | intersection | Hull (removes collinear/redunc | est point on line segment a-b to point $c */ \dots$ | of union of rectangles in $O(N^2)$ - include | Area of a polygon (positive : counterplockwise orientation) */ | |
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```
/* 2D Geometry */
struct point {
  double x, y;
  point(double X = 0, double Y = 0) : x(X), y(Y) {}
double SQR(double d) {
  return d*d;
double PI = 2*acos(0.0);
/* Area of a polygon (positive : counterclockwise orientation) */
double area poly(point p[], int n) {
  double sum = 0;
  for (int i = n-1, j = 0; j < n; i = j++)
    sum += p[i].x * p[j].y - p[i].y * p[j].x;
  return sum/2.0;
/* Area of union of rectangles in O(N^2) - includes custom structs */
// rect r[MAXN]; double ys[2*MAXN]; edge e[2*MAXN];
struct rect {
  double minx, miny, maxx, maxy;
  };
struct edge {
  double x, miny, maxy; char m;
  bool operator (const edge &a) const { return x < a.x; }
double area unionrect(rect r[], double ys[], edge e[], int n) {
  int flag; double curr, sum = 0, sx;
  for (int i = 0; i < n; ++i) {
    e[2*i].miny = e[2*i+1].miny = ys[2*i] = r[i].miny;
    e[2*i].maxy = e[2*i+1].maxy = ys[2*i+1] = r[i].maxy;
    e[2*i].x = r[i].minx; e[2*i].m = 1;
    e[2*i+1].x = r[i].maxx; e[2*i+1].m = -1;
  sort(ys, ys+2*n); sort(e, e+2*n);
  for (int i = 0; i < 2*n; ++i) {
   if (i) sum += (ys[i]-ys[i-1])*curr;
    curr = flaq = 0;
    for(int j = 0; j < 2*n; ++j)
      if ((e[j].miny \le ys[i]) \&\& (ys[i] < e[j].maxy)) {
        if (!flaq) sx = e[j].x;
        flaq += e[j].m;
        if (!flag) curr += e[i].x-sx;
  return sum;
/* Closest point on line segment a-b to point c */
// For infinite line a-b, remove clamping of dp
point closest pt lineseg(point a, point b, point c) {
  b.x -= a.x; \overline{b.y} -= a.y;
  if ((fabs(b.x) < EPS) && (fabs(b.y) < EPS)) return a;
  double dp = (b.x*(c.x-a.x) + b.y*(c.y-a.y))/(SQR(b.x)+SQR(b.y));
  if (dp > 1) dp = 1; if (dp < 0) dp = 0;
  return point(b.x*dp + a.x, b.y*dp + a.y);
/* Convex Hull (removes collinear/redundant points) */
// (reorders original data)
```

```
enum {CCW, CW, CL};
int cross prod(const point &p1, const point &p2, const point &p0) {
 double x1 = p1.x-p0.x, x2 = p2.x-p0.x;
  double y1 = p1.y-p0.y, y2 = p2.y-p0.y;
  double res = x1*y2 - x2*y1;
 if (fabs(res) < EPS) return CL:
 else if (res > 0.0) return CW;
 return CCW;
struct polar cmp {
 point P0;
  polar cmp(point p) : P0(p) {}
  bool operator()(const point &p1, const point &p2) const {
    int res = cross prod(p1, p2, P0);
    if (res == CW) return true;
    else if (res == CCW) return false:
    double x1 = p1.x-P0.x, x2 = p2.x-P0.x;
    double y1 = p1.y-P0.y, y2 = p2.y-P0.y;
    double d = (SQR(x1) + SQR(y1)) - (SQR(x2) + SQR(y2));
    if (fabs(d) < EPS) return false;
    else if (d < 0.0) return true;
    return false:
int convex hull(point poly[], point hull[], int n) {
 if (n < 1) return 0;
  int i, min = 0, h = 1;
  point *P0 = &hull[0]; *P0 = poly[0];
  for (i = 1; i < n; ++i)
   if ((poly[i].y < P0->y) ||
        ((fabs(poly[i].y - PO->y) < EPS) && (poly[i].x < PO->x))) 
      min = i; *P0 = poly[i];
  poly[min] = poly[0]; poly[0] = *P0;
  if (n == 1) return h;
  sort(poly+1, poly+n, polar cmp(*P0));
  for (i = 1; i < n; ++i)
    if ((fabs(poly[i].x - hull[0].x) > EPS) ||
        (fabs(poly[i].y - hull[0].y) > EPS))
      break:
  if (i == n) return h;
 hull[h++] = polv[i++];
  for (; i < n; i++) {
    while ((h > 1) \&\& (cross prod(poly[i], hull[h-1], hull[h-2]) != CCW))
    hull[h++] = poly[i];
  return h;
/* Area of intersection of two circles */
struct circle
 point o; double r;
double CIArea(circle &A, circle &B) {
 double dA, dB, tx, ty;
```

double d = sqrt(SQR(B.o.x-A.o.x)+SQR(B.o.y-A.o.y));

dA = tx = (SQR(d) + SQR(A.r) - SQR(B.r))/(d*2);

ty = sqrt(SQR(A.r) - SQR(tx)); dB = d - dA;

if $(d \ge A.r + B.r)$ return 0:

if ((d < EPS) || (d + A.r <= B.r) || (d + B.r <= A.r))
return SQR((B.r < A.r) ? B.r : A.r) * PI;

```
return SOR(A.r) *acos(dA/A.r) - dA*sgrt(SOR(A.r)-SOR(dA))
    + SOR(B.r)*acos(dB/B.r) - dB*sgrt(SOR(B.r)-SOR(dB));
/* Line segment a-b vs. c-d intersection (returns IP in p) */
// returns 1 if intersect, 0 if not, -1 if coincident
int intersect line(point a, point b, point c, point d, point &p) {
  double num1 = (a.y-c.y)*(d.x-c.x)-(a.x-c.x)*(d.y-c.y);
  double num2 = (a.y-c.y)*(b.x-a.x)-(a.x-c.x)*(b.y-a.y);
  double denom = (b.x-a.x)*(d.y-c.y)-(b.y-a.y)*(d.x-c.x);
  if (fabs(denom) >= EPS) {
    double r = num1 / denom, s = num2 / denom;
    if (0-EPS <= r && r <= 1+EPS &&
        0-EPS <= s && s <= 1+EPS) {
      p.x = a.x + r*(b.x - a.x); p.y = a.y + r*(b.y - a.y);
      return 1:
    return 0;
  if (fabs(num1) >= EPS) return 0;
  if (a.x > b.x || (a.x == b.x && a.y > b.y)) swap(a, b);
  if (c.x > d.x \mid | (c.x == d.x && c.v > d.v)) swap(c, d);
  if (a.x == b.x) {
    if (b.y == c.y) { p = b; return 1; }
    else if (a.y == d.y) { p = a; return 1; }
    else if (b.y < c.y \mid | d.y < a.y) return 0;
  else {
    if (b.x == c.x) \{ p = b; return 1; \}
    else if (a.x == d.x) \{ p = a; return 1; \}
    else if (b.x < c.x \mid | d.x < a.x) return 0;
  return -1;
/* Area of intersection of two general polygons (N^2) */
// 1 for clockwise ordering of points, -1 for counter-clockwise
bool operator<(const point &a, const point &b) {</pre>
  return (a.y+EPS < b.y) || ((a.y-EPS < b.y) && (a.x+EPS < b.x));
bool operator==(const point &a, const point &b) {
  return !(a < b) && !(b < a);
struct triangle {
  point p[3];
  };
//2D cross product of vectors a->b, c->d
double cross(const point &a, const point &b, const point &c, const point &d)
  return ((b.x - a.x)*(d.y - c.y) - (d.x - c.x)*(b.y - a.y));
// This function classifies p as either being "left of" [-1],
// "right of" [+1] or "on" [0] the line a -> b.
int leftRight(const point &a, const point &b, const point &p){
  double res = cross(a, b, a, p);
  if (res > EPS) return -1;
  else if (res < -EPS) return 1;
  return 0:
// This function returns non-zero if point b in the sequence a->b->c
// is a concave point or zero if it is convex.
```

```
int isConcave(point &a, point &b, point &c) {
   return (ORDER*leftRight(a, b, c) <= 0);
 // This function returns non-zero if point p is located on or
  // inside the triangle <a b c>.
  int isInsideTriangle(point &a, point &b, point &c, point &p) {
   int r1 = leftRight(a,b,p), r2 = leftRight(b,c,p), r3 = leftRight(c,a,p);
   return ((ORDER*r1 >= 0) && (ORDER*r2 >= 0) && (ORDER*r3 >= 0));
 // Takes in a list of n ordered points forming the polygon P, and returns
  // a vector of n-2 triangles in T. P is modified during the
  // triangulation. (n >= 3)
  void triangulate(list<point> &P, vector<triangle> &T) {
   list<point>::iterator a, b, c, q; triangle t;
   T.clear();
   if (P.size() < 3) return;
   for (a=b=P.begin(), c=++b, ++c; c != P.end(); a=b, c=++b, ++c) {
      if (!isConcave(*a, *b, *c)) {
        for (q = P.begin(); q != P.end(); q++) {
          if (q == a) \{ ++q; ++q; continue; \}
          if (isInsideTriangle(*a, *b, *c, *q)) break;
        if (g == P.end()) {
                t.p[0] = *a; t.p[1] = *b; t.p[2] = *c; T.push_back(t);
                P.erase(b); b = a;
                if (b != P.begin()) b--;
  //Finds the pt of intersection between line segments a->b and c->d
  //Returns 1 if there is one point of intersection, stored in p.
  //Returns 0 if there is no point of intersection, or infinitely many
  int isectLineSeqs(point &a, point &b, point &c, point &d, point &p) {
   double r, s, dn, n1, n2;
   n1 = cross(c, d, c, a);
   n2 = -cross(a, b, a, c);
    dn = cross(a, b, c, d):
   if ((dn > EPS) || (dn < -EPS)) {
      r = n1/dn; s = n2/dn;
      if ((-EPS < r) && (r < 1+EPS) &&
                (-EPS < s) && (s < 1+EPS)) {
        p.x = a.x + r*(b.x - a.x);
        p.y = a.y + r*(b.y - a.y);
        return 1;
    return 0;
  double areaPoly(vector<point> &P) {
   double area = 0.0;
   vector<point>::iterator p, q;
   for (p = P.end()-1, q = P.begin(); q != P.end(); p = q++)
     area += p->x*q->y - p->y*q->x;
   return area/2.0;
  point P0;
  bool radialLessThan(const point &a, const point &b) {
   return (ORDER == leftRight(P0, a, b));
3 double isectAreaTriangles(triangle &a, triangle &b) {
   vector<point> P; vector<point>::iterator s, e, ne;
```

```
point p; triangle T[2] = { a, b };
  int i, j, r, t, u, v; double area = 0;
  P.clear():
  for(r=1, t=0; t < 2; r = t++)
    for(i = 2, j = 0; j < 3; i = j++) {
      if(isInsideTriangle(T[r].p[0],T[r].p[1],T[r].p[2],T[t].p[i]))
        P.push back(T[t].p[i]);
      for (u = 2, v = 0; v < 3; u = v++)
        if(isectLineSegs(T[t].p[i],T[t].p[j],T[r].p[u],T[r].p[v],p))
                 P.push back(p);
  if (!P.empty()) {
    s = P.begin(); e = P.end();
    sort(s, e); ne = unique(s, e);
    P.erase(ne, e);
    if (P.size() >= 3) {
      P0 = P[0]; sort(s+1, ne, radialLessThan);
      area = areaPoly(P);
  return area;
double isectAreaGpoly(list<point> &P, list<point> &Q) {
  double area = 0.0;
  vector<triangle> S, T; vector<triangle>::iterator s, t;
  triangulate(P, S); triangulate(Q, T);
  for (s = S.begin(); s != S.end(); s++)
    for (t = T.begin(); t != T.end(); t++)
      area += isectAreaTriangles(*s, *t);
  return -ORDER*area;
/* Point in polygon */
#define BOUNDARY 1
double dist2d(point a, point b) {
  return sqrt(SQR(a.x-b.x) + SQR(a.y-b.y));
int pt in poly(point p[], int n, point a) {
  int i, j, c = 0;
  for (i = 0, j = n-1; i < n; j = i++) {
    if (dist2d(p[i],a)+dist2d(p[j],a)-dist2d(p[i],p[j]) < EPS)
      return BOUNDARY:
    if (((p[i].y <= a.y) \&\& (a.y < p[j].y)) || ((p[j].y <= a.y) \&\& (a.y < p[i].y)))
      if (a.x-p[i].x < (p[j].x-p[i].x) * (a.y-p[i].y) / (p[j].y-p[i].y))
        C = !C:
  return c;
/* Rotate P around a point O counterclockwise */
point rotate 2d(point p, point o, double theta) {
  double m[2][2]; point r;
  m[0][0] = m[1][1] = cos(theta);
  m[0][1] = -\sin(\tanh a); m[1][0] = -m[0][1];
  p.x -= o.x; p.y -= o.y;
  r.x = m[0][0] * p.x + m[0][1] * p.y + o.x;
  r.y = m[1][0] * p.x + m[1][1] * p.y + o.y;
  if(fabs(r.x) < EPS) r.x = 0; if(fabs(r.y) < EPS) r.y = 0;
  return r;
```

```
int circle(point p1, point p2, point p3, point &center, double &r)
  double a,b,c,d,e,f,q;
  a = p2.x - p1.x; b = p2.y - p1.y;
  c = p3.x - p1.x; d = p3.y - p1.y;
  e = a*(p1.x + p2.x) + b*(p1.y + p2.y);
  f = c*(p1.x + p3.x) + d*(p1.y + p3.y);
  g = 2.0*(a*(p3.y - p2.y) - b*(p3.x - p2.x));
  if (fabs(q) < EPS) return 0;
  center.x = (d*e - b*f) / g; center.y = (a*f - c*e) / g;
  r = sqrt(SQR(p1.x-center.x)+SQR(p1.y-center.y));
  return 1;
/* Polygon midpoints -> vertices (n odd) */
void midpts2vert(point midpts[], point poly[], int n) {
  poly[0] = midpts[0];
  for (int i = 1; i < n; i += 2) {
    poly[0].x += midpts[i+1].x - midpts[i].x;
    poly[0].y += midpts[i+1].y - midpts[i].y;
  for (int i = 1; i < n; i++) {
    poly[i].x = 2.0*midpts[i-1].x - poly[i-1].x;
    poly[i].y = 2.0*midpts[i-1].y - poly[i-1].y;
/* Minimum bounding circle */
struct circle {
 double x, y, r;
bool inside(point p, circle c)
 return SQR(p.x-c.x)+SQR(p.y-c.y) <= SQR(c.r);
circle Circle1(point p) {
  circle c; c.x = p.x; c.y = p.y; c.r = 0;
  return c;
circle Circle2(point p1, point p2) {
  circle c; c.x = 0.5*(p1.x + p2.x); c.y = 0.5*(p1.y + p2.y);
  c.r = 0.5*sqrt(SQR(p1.x-p2.x)+SQR(p1.y-p2.y));
  return c;
circle Circle3(point p1, point p2, point p3) {
  circle res; double a,b,c,d,e,f,g;
  a = p2.x - p1.x; b = p2.y - p1.y;
  c = p3.x - p1.x; d = p3.y - p1.y;
  e = (p2.x + p1.x)*a + (p2.y + p1.y)*b;
  f = (p3.x + p1.x)*c + (p3.y + p1.y)*d;
  g = 2.0*(a*(p3.y - p2.y) - b*(p3.x - p2.x));
  if (fabs(q) < EPS) {
    res.x = res.y = res.r = DBL MAX; // cfloat
    return res;
  res.x = (d*e - b*f) / q; res.y = (a*f - c*e) / q;
  res.r = sqrt(SQR((p1.x-res.x))+SQR((p1.y-res.y)));
  return res;
circle min circle(point p[], int n) {
  int i, j, k; point t; circle c = Circle1(p[0]);
  for(i = 0; i < n; i++) {
    j = rand() % n; k = rand() % n;
    t = p[j]; p[j] = p[k]; p[k] = t;
```

```
for(i = 1; i < n; i++) if(!inside(p[i], c)) {
    c = Circle1(p[i]);
    for(j = 0; j < i; j++) if(!inside(p[j], c)) {
      c = Circle2(p[i],p[j]);
      for (k = 0; k < j; k++) if (!inside(p[k], c))
        c = Circle3(p[i],p[j],p[k]);
  return c;
/* 3D Geometry Primitives */
struct point {
  double x, y, z;
  point(double X=0, double Y=0, double Z=0) : x(X), y(Y), z(Z) {}
  point operator+(point p) { return point(x + p.x, y + p.y, z + p.z); }
  point operator*(double k) { return point(k*x, k*y, k*z); }
  point operator-(point p) { return *this + (p*-1.0); }
  point operator/(double k) { return *this*(1.0/k); }
  double mag2() { return x*x + y*y + z*z; }
  double mag()
                 { return sqrt(maq2()); }
  point norm() { return *this/this->mag(); }
double dot(point a, point b) {
  return a.x*b.x + a.y*b.y + a.z*b.z;
point cross(point a, point b) {
  return point(a.y*b.z - b.y*a.z, b.x*a.z - a.x*b.z, a.x*b.y - b.x*a.y);
struct line {
  point a, b;
  line(point A=point(), point B=point()) : a(A), b(B) {}
  // Direction unit vector a -> b
  point dir() { return (b - a).norm(); }
  };
// Closest point on an infinite line u to a given point p
point cpoint iline(line u, point p) {
  point ud = u.dir();
  return u.a - ud*dot(u.a - p, ud);
// Shortest distance between two infinite lines u and v
double dist ilines(line u, line v) {
  return dot(v.a - u.a, cross(u.dir(), v.dir()).norm());
// Finds the closest point on infinite line u to infinite line v.
// Note: if (uv*uv - uu*vv) is zero then the lines are parallel and such a
// single closest point does not exist. Check for this if needed.
point cpoint ilines(line u, line v) {
  point ud = u.dir(); point vd = v.dir();
  double uu = dot(ud, ud), vv = dot(vd, vd), uv = dot(ud, vd);
  double t = dot(u.a, ud) - dot(v.a, ud); t *= vv;
  t -= uv*(dot(u.a, vd) - dot(v.a, vd));
  t /= (uv*uv - uu*vv);
  return u.a + ud*t;
// Closest point on a line segment u to a given point p
point cpoint lineseq(line u, point p) {
  point ud = u.b - u.a; double s = dot(u.a - p, ud)/ud.mag2();
  if (s < -1.0) return u.b;
```

```
if (s > 0.0) return u.a;
 return u.a - ud*s;
// Planes
struct plane {
 point n, p;
 plane(point ni = point(), point pi = point()) : n(ni), p(pi) {}
 plane(point a, point b, point c): n(cross(b - a, c - a).norm()), p(a) {}
 //Value of d for the equation ax + by + cz + d = 0
 double d() { return -dot(n, p); }
//Closest point on a plane u to a given point p
point cpoint plane(plane u, point p) {
  return p - u.n*(dot(u.n, p) + u.d());
//Point of intersection between an infinite line v and a plane u.
//Note: if dot(u.n, vd) == 0 then the line and plane do not intersect at a
//single point. Check for this case if it is needed.
point iline isect plane(plane u, line v) {
 point vd = v.dir();
  return v.a - vd*((dot(u.n, v.a) + u.d())/dot(u.n, vd));
//Infinite line of intersection between two planes u and v.
//Note: if dot(v.n, uvu) == 0 then the two planes do not intersect at a line.
//Check for this case if it is needed.
line isect planes(plane u, plane v) {
 point o = u.n*-u.d(), uv = cross(u.n, v.n);
  point uvu = cross(uv, u.n);
  point a = o - uvu*((dot(v.n, o) + v.d())/(dot(v.n, uvu)*uvu.mag2()));
  return line(a, a + uv);
/* Great Circle distance (lat[-90,90], long[-180,180]) */
double greatcircle(double lt1, double lo1, double lt2, double lo2, double r) {
  double a = PI*(lt1/180.0), b = PI*(lt2/180.0);
  double c = PI*((lo2-lo1)/180.0);
  return r*acos(sin(a)*sin(b) + cos(a)*cos(b)*cos(c));
/* Complex Arithmetic structures */
struct rect {
  double x, y; rect(double X = 0, double Y = 0) : x(X), y(Y) {}
struct polar {
  double r, t; polar(double R = 0, double T = 0) : r(R), t(T) {}
  };
struct complex {
  double x, y, r, t;
  complex(double X = 0, double Y = 0) : x(X), y(Y) { setPol(); }
  complex(rect R) : x(R.x), y(R.y) { setPol(); }
  complex(polar P) : x(P.r*cos(P.t)), y(P.r*sin(P.t)) { setPol(); }
  void setPol() { r = sqrt(x*x + y*y); t = atan2(y, x); }
  complex conj() const { return complex(x, -y); }
  complex& operator+=(const complex &a) {
   x += a.x; y += a.y; setPol(); return *this;
  complex& operator = (const complex &a) {
    x -= a.x; y -= a.y; setPol(); return *this;
  complex& operator*=(const complex &a) {
    r *= a.r; t += a.t; x = r*cos(t); y = r*sin(t); setPol(); return *this;
```

```
complex& operator/=(const complex &a) {
    if (a == 0) return *this;
    r \neq a.r; t = a.t; x = r*cos(t); y = r*sin(t); setPol(); return *this;
  complex operator+(const complex &a) const {
    complex res = *this; return res += a;
  complex operator-(const complex &a) const {
    complex res = *this; return res -= a;
  complex operator*(const complex &a) const {
    complex res = *this; return res *= a;
  complex operator/(const complex &a) const {
    complex res = *this; return res /= a;
  bool operator == (const complex &a) const {
    return (fabs(x-a.x) < EPS) && (fabs(y-a.y) < EPS);
  bool operator!=(const complex &a) const {
    return !operator==(a);
  };
/* End Geometry */
/* Arithmetic: Discrete Logarithm solver
   Description: Given prime P, B, and N, finds the smallest
                 exponent L such that B^L == N \pmod{P} O(\operatorname{sgrt}(P)) */
map<UI,UI> M;
UL times (UL a, UL b, UL m) {
  return (ULL) a * b % m; }
UL power(UL val, UL power, UL m){
  UL res = 1, p;
  for(p = power; p; p=p>>1) \{
    if(p & 1) res = times(res, val, m);
    val = times(val, val, m);
  return res:
int discrete log(UI p, UI b, UI n){
  UL i, j, jump;
  M.clear();
  iump = (int)sqrt(p);
  for (i = 0; i < jump && i < p-1; i++){}
    M[power(b,i,p)] = i+1;
  for (i = 0; i < p-1; i+= jump)
    if (j = M[times(n,power(b,p-1-i,p),p)]) 
      return (i+j)%(p-1);
  return -1;
```

```
/* Arithmetic: Fast Exponentition */
LL fast exp(int b, int n){
 LL res = 1, x = b, p;
 for(p = n; p; p >>= 1, x *= x)
   if(p \& 1) res *= x;
 return res:
/* Arithmatic: Simpson's Rule for Numerical Intergration */
double Simpson(double a, double b, int k, double (*f)(double)){
 double dx, x, t; int i;
 dx = (b-a)/(2.0*k);
 t. = 0:
 for( i=0; i<k; i++ )
   t += (i==0 ? 1.0 : 2.0) * (*f)(a+2.0*i*dx);
   t += 4.0 * (*f)(a+(2.0*i+1.0)*dx);
 t += (*f)(b);
 return t * (b-a)/6.0/k;
/* Arithmetic: Cubic equation solver */
typedef struct{
                /* Number of solutions */
 int n;
 double x[3]; /* Solutions */
} Result:
double PI;
Result solve cubic(double a, double b, double c, double d) {
 Result s:
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double g = (a1*a1 - 3*a2)/9.0, sg = -2*sgrt(g);
 long double r = (2*a1*a1*a1 - 9*a1*a2 + 27*a3)/54.0;
 double z = r*r-a*a*a;
 double theta;
 if(z <= 0){
   s.n = 3:
    theta = acos(r/sqrt(q*q*q));
    s.x[0] = sq*cos(theta/3.0) - a1/3.0;
   s.x[1] = sq*cos((theta+2.0*PI)/3.0) - a1/3.0;
    s.x[2] = sq*cos((theta+4.0*PI)/3.0) - a1/3.0;
 } else {
   s.n = 1;
   s.x[0] = pow(sqrt(z)+fabs(r),1/3.0);
    s.x[0] += q/s.x[0];
   s.x[0] *= (r < 0) ? 1 : -1;
   s.x[0] -= a1/3.0;
 return s;
/* Miscellaneous: Coupons Problem
  Description: Coupons are given away in boxes of cereal. There are
                'm' different kinds of coupons (with equiprobable
                distribution). How many boxes of cereal would you
                have to buy, on average, to collect them all?
```

```
double ncoupons(int m) {
  double num = 0.0; int i;
  for (i = 1; i \le m; i++) num += m/(double) i;
  return num;
/* A related problem: If you buy 'n' boxes of cereal, what is the
   probability you get at least one of each of the 'm' coupons? */
double nways[100][100];
void make coupon table() {
  double fact = 1.0;
  int i, j;
  for (i = 1; i < 100; i++) {
   nways[i][1] = 1.0;
    for (j = 2; j < i; j++)
      nways[i][j] = j*(nways[i-1][j] + nways[i-1][j-1]);
    nwavs[i][i] = fact *= i;
double query table(int m, int n) {
  if (n < m) return 0.0;
  if (m == 0) return 1.0;
  if (n >= 100 \mid | m >= 100) exit(1);
  return nways[n][m]/pow(m,n);
/* Arithmetic: Binomial coefficient */
long double bin[MAXN+1][MAXN+1];
void getBinCoeff(){
  int i, k;
  for(k = 0; k \le MAXN; k++){}
    bin[k][0] = bin[k][k] = 1;
    for(i = 1; i < k; i++)
     bin[k][i] = bin[k-1][i-1]+bin[k-1][i];
/* Combinatorics: Digit Occurence count
   Description: Given a digit and a number N. return the number of
                times the digit occurs from 1..N. */
#include <stdio.h>
#include <string.h>
#include <math.h>
long long digit count(int digit, int max){
  long long res = 0;
  char buff[15];
  int i, count;
  if(max <= 0) return 0;
  /* Number of times "digit" occurs in the one's place */
  res += max/10 + ((max \% 10) >= digit ? 1 : 0);
```

```
/* Since we start from 1, if digit = 0, remove 1 since "0"
    doesn't count */
 if(digit == 0) res--;
 /* Get the number of occurences in max/10-1, and multiply this by
    10 since we can choose 10 possible last digits [0-9] */
 res += digit count(digit, max/10 - 1) * 10;
  /* The number of occurences in max/10 is equal to (1+max%10) * the
    number of times "digit" occurs in max/10 */
  sprintf(buff, "%d", max/10);
 for(i = 0, count = 0; i < strlen(buff); i++)
   if(buff[i] == digit+'0') count++;
 res += (1 + max%10) * count;
 return res:
/* Combinatorics: Digits in N!
  Description: Given N, computes the number of digits that N! will
               occupy in base B. */
long long fac digit(int n, int b) {
 double sum = 0; int i;
 for (i = 2; i \le n; i++) sum += log(i);
 return (long long) floor(1+sum/log(b)); /* don't use ceil! */
/* Combinatorics: Josephus Ring Survivor
  Description: Suppose that there are n people in a ring, [0..n-1].
               Count around the ring, starting from 0, and
               dismissing every m-th person.
*/
int survive[MAXN];
void josephus(int n, int m){
 int i:
 survive[1] = 0;
 for(i = 2; i <= n; i++)
   survive[i] = (survive[i-1]+(m%i))%i;
/* Combinatorics - Permutation index on distinct characters
  Description: Given a string formed of distinct characters,
                returns the index of the permutation from 0..N!-1.
  This does not work when characters can be the same for example: "aaba" */
int permdex (char *s){
 int i, j, size = strlen(s);
 int index = 0:
 for (i = 1; i < size; i++)
   for (j = i; j < size; j++)
     if (s[i-1] > s[j]) index ++;
   index *= size - i;
 return index;
* Dynamic Programming: Longest Ascending Subsequence */
```

```
int asc seg(int *A, int n, int *S){
  int *m, *seq, i, k, low, up, mid, start;
  m = malloc((n+1) * sizeof(int));
  seg = malloc(n * sizeof(int));
  /* assert(m && seq); */
  for (i = 0; i < n; i++) seq[i] = -1;
  m[1] = start = 0;
  for (k = i = 1; i < n; i++) {
    if (A[i] >= A[m[k]]) {
      seq[i] = m[k++];
      start = m[k] = i;
    } else if (A[i] < A[m[1]]) {</pre>
     m[1] = i;
    } else {
      /* assert(A[m[1]] <= A[c] && A[c] < A[m[k]]); */
      low = 1
      up = k;
      while (low != up-1) {
        mid = (low+up)/2;
        if(A[m[mid]] <= A[i]) low = mid:
        else up = mid;
      seq[i] = m[low];
     m[up] = i;
  for (i = k-1; i >= 0; i--)
    S[i] = A[start]:
    start = seq[start];
  free(m); free(seq);
  return k;
int sasc seq(int *A, int n, int *S){
  int *m, *seq, i, k, low, up, mid, start;
  m = malloc((n+1) * sizeof(int));
  seq = malloc(n * sizeof(int));
  /* assert(m && seq); */
  for (i = 0; i < n; i++) seq[i] = -1;
  m[1] = start = 0;
  for (k = i = 1; i < n; i++) {
   if (A[i] > A[m[k]]) {
      seq[i] = m[k++];
      start = m[k] = i;
    } else if (A[i] < A[m[1]]) {</pre>
      m[1] = i;
    } else if (A[i] < A[m[k]]) {</pre>
     low = 1;
      up = k;
      while (low != up-1) {
        /* assert(A[m[h]] <= A[c] && A[c] < A[m[j]]); */
        mid = (low+up)/2;
        if(A[m[mid]] <= A[i]) low = mid;</pre>
        else up = mid;
      if (A[i] > A[m[low]]) {
        seq[i] = m[low];
```

```
m[up] = i;
   for (i = k-1; i >= 0; i--)
     S[i] = A[start];
     start = seq[start];
   free(m); free(seq);
   return k:
  /* Dynamic Programming: Integer Parititoning
    Description: Template for calculating the number of ways of
                  partitioning the integer N into M parts.
    Notes:
                  A partition of a number N is a representation of
                  N as the sum of positive integers
                  e.g. 5 = 1+1+1+1+1
                  The number of ways of partitioning an integer N
                  into M parts is equal to the number of ways of
                  partitioning the number N with the largest element
                  being of size M. This is best seen with a Ferres-
                  Young diagram:
                  Suppose N = 8, M = 3:
                  4 = * * * *
                  3 = * * *
                  1 = *
                      3 2 2 1
                  By transposition from rows to columns, this equality
                  can be seen.
                  P(N, M) = P(N-1, M-1) + P(N-M, M)
                  P(0, M) = P(N, 0) = 0
                  P(N, 1) = 1
  #include <stdio.h>
  #include <string.h>
  #define MAXN 300
  #define ULL unsigned long long
  ULL A[MAXN+1][MAXN+1];
  void Build(){
   int i, j;
   memset(A, 0, sizeof(A));
   A[0][0] = 1;
   for(i = 1; i <= MAXN; i++){
     A[i][1] = 1;
     for(j = 2; j \le i; j++)
       A[i][j] = A[i-1][j-1] + A[i-j][j];
  /* Generator: Catalan Numbers */
8 long long int cat[33];
void getcat() {
```

```
int i;
  cat[0] = cat[1] = 1;
  for (i = 2; i < 33; i++)
    cat[i] = cat[i-1]*(4*i-6)/i;
/* Generators: Binary Strings generator - (ordered by cardinality) */
char bit[MAXN];
void recurse(int n, int curr, int left){
  if(curr == n){
    Process(n);
  } else {
    if(curr+left < n){
     bit[curr] = 0:
     recurse(n, curr+1, left);
    if(left){
     bit[curr] = 1;
      recurse(n, curr+1, left-1);
void gen bin card(int n){
  int i:
  for(i = 0; i \le n; i++)
    printf("Cardinality %d:\n", i);
    recurse(n, 0, i);
/* Graph Theory: Maximum Bipartite Matching
                For vertex i of set U:
                match[i] = -1 means i is not matched
                match[i] = x means the edge i -> (x-|U|) is selected
                For simplicity, use addEdge(i,j,n) to add edges, where
                0 \le i \le |U| and 0 \le j \le |V| and |U| = n.
                If there is an edge from vertex i of U to vertex
                j 	ext{ of } V 	ext{ then: } e[i][j+|U|] = e[j+|U|][i] = 1.
   Notes:
                - If |U| = n and |V| = m, then vertices are assumed
                  to be from [0,n-1] in set U and [0,m-1] in set V.
                - Remember that match[i] -n gives the edge from i,
                  not just match[i].
                - This code is roughly 2 times slower than the old
                  code since it doesn't try multiple BFS paths at
                  once, however, it's about 4 times shorter... */
                          /* How many vertices in U+V (in total) */
#define MAXN 300
                          /* MODIFIED Adj. matrix (see note) */
char e[MAXN][MAXN];
int match[MAXN], back[MAXN], q[MAXN], tail;
void addEdge(int x, int y, int n){
  e[x][y+n] = e[y+n][x] = 1;
```

```
int find(int x, int n, int m){
  int i, i, r:
  if(match[x] != -1) return 0;
  memset(back, -1, sizeof(back));
  for(q[i=0]=x, tail = 1; i < tail; i++)
    for(j = 0; j < n+m; j++)
      if(!e[q[i]][i]) continue;
      if(match[j] != -1){
        if(back[j] == -1){
          back[j] = q[i];
          back[q[tail++] = match[j]] = j;
      } else {
        match[match[q[i]] = j] = q[i];
        for(r = back[q[i]]; r != -1; r = back[back[r]])
          match[match[r] = back[r]] = r;
       return 1:
 return 0;
void bipmatch(int n, int m){
 int i;
 memset(match, -1, sizeof(match));
 for(i = 0; i < n+m; i++) if(find(i,n,m)) i = 0;
int main(){
 int n, m, esize, x, y;
  int i, count;
  /* Read size of set U into n, size of set V into m */
  while(scanf("%d %d", &n, &m) == 2){
                                      /* Clear edges */
    memset(e, 0, sizeof(e));
    scanf("%d", &esize);
                                       /* get # of edges */
    while(esize--){
      scanf("%d %d", &x, &y);
                                       /* add edges */
                                       /* Edges [0,n-1]->[0,m-1] */
      addEdge(x,y,n);
    bipmatch(n, m);
                                       /* Perform matching */
    for(count = i = 0; i < n; i++){ /* Print results */
      if(match[i] != -1){
       printf("%d->%d\n", i, match[i]-n);
        count++;
    printf("Matching size: %d\n", count);
 return 0:
/* Graph Theory: Eulerian Graphs
                necessary data structures.
```

raph Theory: Eulerian Graphs

Before adding edges, call Init() to initialize all necessary data structures.

Use the provided function addEdge(x,y,c) which adds c number of edges between x and y.

isEulerian(int n, int *start, int *end) returns:

```
0 if the graph is not Eulerian
                   1 if the graph has a Euler cycle
                   2 if the graph a path, from start to end
                with n being the number of nodes in the graph */
#define MAXN 105
                     /* Number of nodes */
#define MAXM 505
                     /* Maximum number of edges */
#define min(a,b) (((a)<(b))?(a):(b))
#define \max(a,b) (((a)>(b))?(a):(b))
#define DEC(a,b) q[a][b]--;q[b][a]--;deq[a]--;deq[b]--
int sets[MAXN], deg[MAXN];
int q[MAXN][MAXN];
int seq[MAXM], seqsize;
/* Uncomment if you need copy of graph
   int g2[MAXN][MAXN], deg2[MAXN];
int getRoot(int x){
  if(sets[x] < 0) return x;
  return sets[x] = getRoot(sets[x]);
void Union(int a, int b){
  int ra = getRoot(a), rb = getRoot(b);
  if(ra != rb){
    sets[ra] += sets[rb];
    sets[rb] = ra;
void Init(){
  memset(sets, -1, sizeof(sets));
  memset(q, 0, sizeof(q));
  memset(deg, 0, sizeof(deg));
void addEdge(int x, int y, int count){
  g[x][y] += count; deg[x] += count;
  g[y][x] += count; deg[y] += count;
  Union(x,y);
int isEulerian(int n, int *start, int *end){
  int odd = 0, i, count = 0, x;
  /* Check if graph is connected. If all vertices
     are quaranteed to be used then use this:
     if(sets[getRoot(0)] != -n) return 0;
     Otherwise, count only vertices used like this: */
  for(i = 0; i < n; i++)
    if(deg[i]){
     x = i; count++;
  if(sets[getRoot(x)] != -count) return 0;
  for(i = 0; i < n; i++){
   if(deg[i]%2){
     odd++;
```

```
if(odd == 1) *start = i;
      else if(odd == 2) *end = i:
      else return 0;
 return odd ? 2 : 1;
void getPath(int n, int start, int end){
  int temp[MAXM], tsize = 1, i, j;
  temp[0] = start;
 while(1){
    j = temp[tsize-1];
    for(i = 0; i < n; i++)
      if(i == end) continue;
      if(g[i][j]){
       temp[tsize++] = i;
       DEC(i,j);
       break;
    if(i == n){
      if(g[end][j]){
       temp[tsize++] = end;
       DEC(j,end);
      break;
  for(i = 0; i < tsize; i++)
    if(!deg[temp[i]]) seq[seqsize++] = temp[i];
    else getPath(n, temp[i], temp[i]);
void buildPath(int n, int start, int end){
 segsize = 0;
 /* Uncomment if you need copy of graph
     memcpy(q, q2, sizeof(q));
     memcpy(deg, deg2, sizeof(deg));
 getPath(n, start, end);
int main(){
 int i, x,y,start,end, n, m;
  while(scanf("%d %d", &n, &m) == 2){
    Init();
    for(i = 0; i < m; i++){
      scanf("%d %d", &x, &y);
      addEdge(x,y,1);
    /* Uncomment if you need copy of graph
       memcpy(q2, q, sizeof(q2));
       memcpy(deg2, deg, sizeof(deg2));
    switch(isEulerian(n, &start, &end)){
  return 0;
```



```
/* Graph Theory: Maximum Flow in a directed graph
        - Multiple edges from u to v may be added. They are converted into a
          single edge with a capacity equal to their sum
        - Vertices are assumed to be numbered from 0..n-1
        - The graph is supplied as the number of nodes (n), the zero-based
          indexes of the source (s) and the sink (t), and a vector of edges
u->v
          with capacity c (M).
*/
#include <cstdio>
#include <vector>
#include <list>
using namespace std;
#define MAXN 200
//Edge u->v with capacity c
struct Edge {
  int u, v, c;
};
int F[MAXN][MAXN]; //Flow of the graph
int maxFlow(int n, int s, int t, vector<Edge> &M)
  int u, v, c, oh, min, df, flow, H[n], E[n], T[n], C[n][n];
  vector<Edge>::iterator m;
  list<int> N:
  list<int>::iterator cur;
  vector<int> R[n];
  vector<int>::iterator r;
  for (u = 0; u < n; u++) {
    E[u] = H[u] = T[u] = 0;
    R[u].clear();
    for (v = 0; v < n; v++)
      C[u][v] = F[u][v] = 0;
  for (m = M.begin(); m != M.end(); m++) {
    u = m->u;
    v = m->v;
    C = m->C;
    if (c && !C[u][v] && !C[v][u]) {
      R[u].push back(v);
      R[v].push back(u);
    C[u][v] += c;
  H[s] = n;
  for (r = R[s].begin(); r != R[s].end(); r++) {
    v = *r;
    F[s][v] = C[s][v];
    F[v][s] = -C[s][v];
    E[v] = C[s][v];
    E[s] -= C[s][v];
```

```
N.clear();
    for (u = 0; u < n; u++)
      if ((u != s) && (u != t))
        N.push back(u);
    for (cur = N.begin(); cur != N.end(); cur++) {
      u = *cur;
      oh = H[u];
      while (E[u] > 0)
        if (T[u] >= (int)R[u].size()) {
          min = 10000000;
           for (r = R[u].begin(); r != R[u].end(); r++) {
            if ((C[u][v] - F[u][v] > 0) \&\& (H[v] < min))
               min = H[v];
          H[u] = 1 + min;
          T[u] = 0;
         else {
          v = R[u][T[u]];
          if ((C[u][v] - F[u][v] > 0) \&\& (H[u] == H[v]+1)) {
            df = C[u][v] - F[u][v];
            if (df > E[u])
              df = E[u];
             F[u][v] += df;
             F[v][u] = -F[u][v];
             E[u] -= df;
            E[v] += df;
           else
            T[u]++;
      if (H[u] > oh)
        N.splice(N.begin(), N, cur);
    flow = 0:
    for (r = R[s].begin(); r != R[s].end(); r++)
      flow += F[s][*r];
    return flow;
  /* Graph Theory: Chinese Postman Problem
                   - The maximum # of vertices solvable is roughly 20
  #define MAXN 20
  #define DISCONNECT -1
                           /* Adj matrix (keep lowest cost multiedge) */
  int g[MAXN][MAXN];
  int deg[MAXN];
                           /* Degree count */
  int A[MAXN+1];
                          /* Used by perfect matching generator */
                           /* Sum of costs */
  int sum;
11 int odd;
int best;
```

```
void flovd(int n){
  int i, j, k;
  for(k = 0; k < n; k++) for(i = 0; i < n; i++) for(j = 0; j < n; j++)
    if(q[i][k] != -1 \&\& q[k][j] != -1){
      int temp = q[i][k] + q[k][j];
      if(g[i][j] == -1 || g[i][j] > temp)
        g[i][j] = temp;
  for(i = 0; i < n; i++) q[i][i] = 0;
void checkSum(){
  int i, temp;
  for(i = temp = 0; i < odd/2; i++)
   temp += q[A[2*i]][A[2*i+1]];
  if(best == -1 || best > temp) best = temp;
void perfmatch(int x){
  int i, t:
  if(x == 2) checkSum();
  else {
    perfmatch(x-2);
    for(i = x-3; i >= 0; i--){
      t = A[i];
      A[i] = A[x-2];
      A[x-2] = t:
      perfmatch(x-2);
    t = A[x-2];
    for(i = x-2; i >= 1; i--) A[i] = A[i-1];
    A[0] = t;
int postman(int n){
  int i;
  flovd(n):
  for(odd = i = 0; i < n; i++)
   if(deg[i]%2) A[odd++] = i;
  if(!odd) return sum;
  best = -1:
  perfmatch(odd);
  return sum+best;
int main(){
  int i, u, v, c, n, m;
  while(scanf("%d %d", &n, &m) == 2){
    /* Clear graph and degree count */
    memset(g, -1, sizeof(g));
    memset(deg, 0, sizeof(deg));
    for(sum = i = 0; i < m; i++){
      scanf("%d %d %d", &u, &v, &c);
      u--; v--;
```

```
deg[u]++; deg[v]++;
        if(g[u][v] == -1 || g[u][v] > c) g[u][v] = c;
        if(g[v][u] == -1 || g[v][u] > c) g[v][u] = c;
        sum += C;
      printf("Best cost: %d\n", postman(n));
    return 0;
  /* Graph Theory: Strongly Connected Components */
  #define VI vector<int>
  #define MAXN 1000
  VI g[MAXN], curr;
  vector< VI > scc;
  int dfsnum[MAXN], low[MAXN], id;
  char done[MAXN]:
  void visit(int x){
    curr.push back(x);
    dfsnum[x] = low[x] = id++;
    for(size t i = 0; i < g[x].size(); i++)
      if(dfsnum[g[x][i]] == -1)
        visit(q[x][i]);
        low[x] <?= low[g[x][i]];
      } else if(!done[q[x][i]])
        low[x] <?= dfsnum[q[x][i]];
    if(low[x] == dfsnum[x])
      VI c; int y;
        done[y = curr[curr.size()-1]] = 1;
        c.push back(y);
        curr.pop back();
      \} while(y != x);
      scc.push back(c);
  void strong conn(int n){
    memset(dfsnum, -1, n*sizeof(int));
    memset(done, 0, sizeof(done));
    scc.clear(); curr.clear();
    for(int i = id = 0; i < n; i++)
      if(dfsnum[i] == -1) visit(i);
  /* Number Theory: Converting between bases (arbitrary precision)
     Description: Given a starting base b1, and a target base b2, */
  import java.math.*;
  import java.jo.*;
  import java.util.*;
  class base convert{
    // invalid is the string that is returned if the N is not valid
    static String invalid = new String("Number is not valid");
private static String convert _ base(int base1, int base2,
                                        String n, String key) {
```

```
int i, x;
    String n2 = "", n3 = "";
    BigInteger
       a = BigInteger.ZERO,
        b1 = BigInteger.valueOf(base1),
       b2 = BiqInteger.valueOf(base2);
    for(i = 0; i < n.length(); i++){
        a = a.multiply(b1);
        x = key.indexOf(n.charAt(i));
        if(x == -1 || x >= base1) return invalid;
        a = a.add(BigInteger.valueOf(x));
    while (a.signum() == 1)
        BigInteger r[] = a.divideAndRemainder(b2);
        n2 += key.charAt(r[1].intValue());
       a = r[0];
    for(i = n2.length()-1; i >= 0; i--) n3 += n2.charAt(i);
    if(n3.length() == 0) n3 += '0';
    return n3;
  public static void main(String[] args){
try{
    String line, n;
    int tnum, base1, base2;
    StringTokenizer st;
    // key is the base system that you may change as needed
    String kev = new
       String("0123456789ABCDEFGHIJKLMNOPORSTUVWXYZabcdefghijklmnopgrstu-
vwxyz");
    // Standard IO
    BufferedReader in = new BufferedReader(new InputStreamReader(System.in));
    PrintStream out = System.out;
    // File IO
    // BufferedReader in = new BufferedReader(new FileReader("prob1.dat"));
   // PrintWriter out = new BufferedWriter(new FileWriter("prob1.out"));
   line = in.readLine();
                                         // Get number of test cases
    st = new StringTokenizer(line);
    tnum = Integer.parseInt(st.nextToken());
    for(int t = 0; t < tnum; t++){
        line = in.readLine();
        st = new StringTokenizer(line);
        base1 = Integer.parseInt(st.nextToken());
        base2 = Integer.parseInt(st.nextToken());
        n = st.nextToken();
        String result = convert base(base1, base2, n, key2);
        out.println(result);
} catch(Exception e){
    System.err.println(e.toString());
}}}
/* Java Template: BigInteger Reference
   _____
```

Description: This document is a reference for the use of the

```
BigInteger class in Java. It contains sample code
             that computes GCDs of pairs of integers.
Constants:
[1.2] BigInteger.ONE - The BigInteger constant one.
[1.2] BigInteger.ZERO - The BigInteger constant zero.
Creating BigIntegers
1. From Strings
   a) BigInteger(String val);
   b) BigInteger(String val, int radix);
2. From byte arrays
   a) BigInteger(byte[] val);
   b) BigInteger(int signum, byte[] magnitude)
3. From a long integer
   a) static BigInteger BigInteger.valueOf(long val)
Math operations:
                      --> C = A.add(B):
A + B = C
A - B = C
                     --> C = A.subtract(B);
                      --> C = A.multiplv(B);
A * B = C
A / B = C
                      --> C = A.divide(B);
                      --> C = A.remainder(B);
A % B = C
A % B = C where C > 0 --> C = A.mod(B);
A / B = O & A % B = R --> C = A.divideAndRemainder(B):
                                 (O = C[0], R = C[1])
A ^b = C
                      --> C = A.pow(B);
abs(A) = C
                      --> C = A.abs();
-(A) = C
                      --> C = A.negate();
gcd(A.B) = C
                      --> C = A.qcd(B);
                      --> C = A.modPow(B,M);
(A ^ B) % M
C = inverse of A mod M --> C = A.modInverse(M);
max(A,B) = C
                      --> C = A.max(B);
min(A,B) = C
                      --> C = A.min(B);
Bit Operations
~A = C
            (NOT)
                      --> C = A.not();
A \& B = C
            (AND)
                      --> C = A.and(B);
A \mid B = C
            (OR)
                      --> C = A.or(B);
A \stackrel{\star}{B} = C
                      --> C = A.xor(B);
            (XOR)
A \& \sim B = C (ANDNOT) --> C = A.andNot(B);
A \ll n = C (LSHIFT) --> C = A.shiftLeft(n);
A >> n = C (RSHIFT) --> C = A.shiftRight(n);
Clear n'th bit of A
                     --> C = A.clearBit(n);
                     --> C = A.setBit(n);
Set n'th bit of A
Flip n'th bit of A
                     --> C = A.flipBit(n);
Test n'th bit of A
                    --> C = A.testBit(n);
                      --> n = A.bitCount();
Bitcount of A = n
Bitlength of A = n
                    --> n = A.bitLength();
Lowest set bit of A --> n = A.getLowestSetBit();
```

```
or A.equals(B);
                       --> A.compareTo(B) == 1;
  A > B
                       --> A.signum() == -1;
  A < 0
  A == 0
                       --> A.signum() == 0;
  A > 0
                       --> A.signum() == 1;
  Conversion.
  double
                       --> A.doubleValue();
  float
                      --> A.floatValue();
  int
                       --> A.intValue();
                      --> A.longValue();
  long
                       --> A.toByteArray();
  byte[]
  String
                       --> A.toString();
  String (base b)
                      --> A.toString(b);
   -----*/
/* Reads in lines of input until EOF is encountered. For each line
  of input it will extract two integers and then print out their
  GCD. */
import java.math.*;
import java.io.*;
import java.util.*;
class BigIntegers {
   public static void main(String[] args) {
       BufferedReader in = new BufferedReader(
                            new InputStreamReader(System.in));
       String line;
       StringTokenizer st;
       BigInteger a;
       BigInteger b;
       try {
          while(true) {
              line = in.readLine();
              if(line == null) break;
              st = new StringTokenizer(line);
              a = new BigInteger(st.nextToken());
              b = new BigInteger(st.nextToken());
              System.out.println( a.gcd(b) );
       } catch(Exception e) {
          System.err.println(e.toString());
/* Java Template: IO Reference
   _____
  Description: This document is a reference for the use of the
              java for regular IO purposes. It covers stdin and
              stdout as well as file IO. It also shows how to use
              StringTokenizer for parsing.
```

--> A.compareTo(B) == -1;

--> A.compareTo(B) == 0

A == B

```
Author: Patrick Earl
  Date:
               Nov 14, 2002
  References: Java API Documentation
  Reliability: 0
*/
import java.util.*;
import java.io.*;
class IO {
    public static void main(String[] args) {
        try {
            BufferedReader in=new BufferedReader(
                                new FileReader("prob1.dat"));
            PrintWriter out=new PrintWriter(
                              new BufferedWriter(
                                new FileWriter("prob1.out")));
            /* For stdin/stdout IO, use: */
            PrintStream out = System.out;
            BufferedReader in = new BufferedReader(
                                 new InputStreamReader(System.in));
            String line;
            int num=0;
            StringTokenizer st;
            while(true) {
                /* Newlines are removed by readLine(). */
                line = in.readLine();
                if(line == null) break;
                num++;
                /* Print out line number. */
                out.println("Line #" + num);
                /* Split on whitespace */
                st = new StringTokenizer(line);
                while(st.hasMoreTokens()) {
                    out.print("Token: ");
                    out.println(st.nextToken());
                /* To split on something else, use:
                     st = new StringTokenizer(line, delim);
                   Or use this to change in the middle of parsing:
                     line = st.nextToken(delim);
            /* You must flush for files! */
            out.flush();
        } catch(Exception e) {
            System.err.println(e.toString());
```

```
/* Miscellaneous: Bit Count */
int bitcount(int a){
 int c = 0;
  while(a){
    c++; a &= a-1;
  } return c;
/* Number Theory: Euler Phi function */
int phi(int n){
  int i, count, res = 1;
  for(i = 2; i*i <= n; i++){}
    count = 0;
   while (n % i == 0)
     n /= i;
     count++;
    if(count > 0) res *= (pow(i, count)-pow(i, count-1));
  if(n > 1) res *= (n-1);
  return res:
/* Number Theory: Primality Testing */
int isPrime(int x){
  int i;
  if( x == 1 ) return ONEPRIME;
  if(x == 2) return 1;
  if( x % 2 == 0) return 0;
  for(i = 3; i*i <= x; i+=2)
   if(x \% i == 0) return 0;
  return 1:
/* Number Theory: Number of Divisors; O(sqrt(N)) */
#include <stdio.h>
int num divisors(int n){
  int i, count, res = 1;
  for(i = 2; i*i <= n; i++){}
   count = 0;
   while(!(n%i)){
     n /= i;
     count++;
    if(count) res *= (count+1);
  if(n > 1) res *= 2;
  return res;
/* Number Theory: Prime Factorization */
```

```
int primes[MAXP]; int psize;
  void getPrimes(){
    int i, j, isprime;
    psize = 0;
    primes[psize++] = 2;
    for(i = 3; i \le MAXN; i+= 2)
      for(isprime = j = 1; j < psize; j++){
        if(i % primes[j] == 0){
          isprime = 0;
          break;
         if(1.0*primes[j]*primes[j] > i) break;
      if(isprime) primes[psize++] = i;
  typedef struct{
    int size;
    int f[32]:
  } Factors:
  Factors getPFactor(int n){
    Factors x:
    int i;
    x.size = 0;
    for(i = 0; i < psize; i++){}
      while(n % primes[i] == 0){
        x.f[x.size++] = primes[i];
        n /= primes[i];
      if(1.0*primes[i]*primes[i] > n) break;
    if(n > 1)
      x.f[x.size++] = n;
    return x;
  /* Number Theory: Primality testing with a sieve */
  #define TEST(f,x) (*(f+(x)/16)&(1<<(((x)%16L)/2)))
  #define SET(f,x) *(f+(x)/16)|=1<<(((x)%16L)/2)
  #define ONEPRIME 0 /* whether or not 1 is considered to be prime */
  #define UL unsigned long
  #define UC unsigned char
  UC *primes = NULL;
  UL getPrimes(UL maxn) {
    UL x, y, psize=1;
    primes = calloc(((maxn)>>4)+1L, sizeof(UC));
    for (x = 3; x*x <= maxn; x+=2)
      if (!TEST(primes, x))
        for (y = x*x; y \le maxn; y += x << 1) SET (primes, y);
15 /* Comment out if you don't need # of primes <= maxn */
```

```
for(x = 3; x \leq maxn; x+=2)
    if(!TEST(primes, x)) psize++:
  return psize;
/* Returns whether or not a given POSITIVE number if prime. */
int isPrime(UL x){
  if(x == 1) return ONEPRIME;
  if(x == 2) return 1:
  if(x % 2 == 0) return 0;
  return (!TEST(primes, x));
/* Number Theory: Sum of divisors O(sqrt(N)) */
LL sum divisors(LL n){
  int i, count; LL res = 1;
  for(i = 2; i*i <= n; i++){}
    count = 0;
    while(n \% i == 0){
      n /= i:
      count++;
    if(count) res *= ((pow(i, count+1)-1)/(i-1));
  if(n > 1) res *= ((pow(n, 2)-1)/(n-1));
  return res;
/* Chinese Remainder Theorem (cra.c)
 * Author: Howard Cheng
 * Reference:
 * Geddes, K.O., Czapor, S.R., and Labahn, G. Algorithms for Computer
   Algebra, Kluwer Academic Publishers, 1992, p. 180
 * Given n relatively prime modular in m[0], ..., m[n-1], and right-hand
 * sides a[0], ..., a[n-1], the routine solves for the unique solution
 * in the range 0 \le x \le m[0]*m[1]*...*m[n-1] such that x = a[i] \mod m[i]
 * for all 0 <= i < n. The algorithm used is Garner's algorithm, which
 * is not the same as the one usually used in number theory textbooks.
 * It is assumed that m[i] are positive and pairwise relatively prime.
 * a[i] can be any integer.
 * If the system of equations is
 * x = a[0] \mod m[0]
 * \quad x = a[1] \mod m[1]
 * then a[i] should be reduced mod m[i] first.
 * Also, if 0 \ll a[i] \ll m[i] for all i, then the answer will fall
 * in the range 0 \le x \le m[0] * m[1] * ... * m[n-1].
 * Added: 5 January 2000
 * Confirmed: Matthew McNaughton (mcnaught@cs.ualberta.ca)
#include <stdio.h>
#include <stdlib.h>
```

```
#include <assert.h>
int gcd(int a, int b, int *s, int *t){
 int r, r1, r2, a1, a2, b1, b2, q;
 a1 = b2 = 1;
 a2 = b1 = 0;
 while (b) {
   /* assert(a1*A + a2*B == a); */
   q = a / b;
   r = a % b;
   r1 = a1 - q*b1;
   r2 = a2 - q*b2;
   a = b;
   a1 = b1:
   a2 = b2:
   b = r:
   b1 = r1;
   b2 = r2;
 *s = a1;
 *t = a2:
 /* assert(a >= 0); */
 return a;
int cra(int n, int *m, int *a){
 int x, i, k, prod, temp;
 int *gamma, *v;
 gamma = malloc(n*sizeof(int));
       = malloc(n*sizeof(int));
  /* assert(gamma && v); */
  /* compute inverses */
  for (k = 1; k < n; k++) {
   prod = m[0] % m[k];
   for (i = 1; i < k; i++) {
     prod = (prod * m[i]) % m[k];
   gcd(prod, m[k], gamma+k, &temp);
   qamma[k] %= m[k];
    if (qamma[k] < 0) {
      qamma[k] += m[k];
  /* compute coefficients */
 v[0] = a[0];
  for (k = 1; k < n; k++) {
    temp = v[k-1];
    for (i = k-2; i >= 0; i--)
      temp = (temp * m[i] + v[i]) % m[k];
     if (temp < 0) {
        temp += m[k];
    v[k] = ((a[k] - temp) * gamma[k]) % m[k];
    if (v[k] < 0) {
     v[k] += m[k];
```



```
/* convert from mixed-radix representation */
  x = v[n-1];
  for (k = n-2; k >= 0; k--)
   x = x * m[k] + v[k];
  free (gamma);
  free(v);
  return x:
int main(void) {
  int n, *m, *a, i, x;
  while (scanf(``%d'', &n) == 1 && n > 0) {
   m = malloc(n*sizeof(int));
   a = malloc(n*sizeof(int));
   assert(m && a);
   printf("Enter moduli:\n");
    for (i = 0; i < n; i++) {
      scanf("%d", m+i);
   printf("Enter right-hand side:\n");
    for (i = 0; i < n; i++) {
     scanf("%d", a+i);
   x = cra(n, m, a);
   printf("x = %d\n", x);
    for (i = 0; i < n; i++) {
     assert((x-a[i]) % m[i] == 0);
    free(m);
   free(a);
  return 0:
/* Extended Euclidean Algorithm */
int gcd(int a, int b, int *s, int *t)
  int r, r1, r2, a1, a2, b1, b2, q;
 int A = a;
  int B = b:
  /* unnecessary if a, b >= 0 */
  if (a < 0) {
   r = gcd(-a, b, s, t);
   *s *= -1;
   return r;
  if (b < 0) {
   r = gcd(a, -b, s, t);
    *t *= -1;
    return r;
  a1 = b2 = 1;
  a2 = b1 = 0;
```

```
while (b) {
      assert(a1*A + a2*B == a);
      q = a / b;
      r = a % b;
      r1 = a1 - q*b1;
      r2 = a2 - q*b2;
      a = b;
      a1 = b1;
      a2 = b2:
      b = r:
      b1 = r1;
      b2 = r2;
    *s = a1:
    *t = a2;
    assert(a >= 0);
    return a:
  /* Fast Exponentiation mod m */
  int fast exp(int b, int n, int m)
    int res = 1:
    int x = b;
    while (n > 0) {
     if (n & 0x01) {
        res = (res * x) % m;
      } else {
        n >>= 1;
        x = (x * x) % m;
    return res:
  // Simplex Method for Linear Programming
  // m - number of (less than) inequalities
  // n - number of variables
      C - (m+1) by (n+1) array of coefficients:
                       - objective function coefficients
 //
          row 0
          row 1:m
                       - less-than inequalities
          column 0:n-1 - inequality coefficients
          column n - inequality constants (0 for objective function)
 // X[n] - result variables
      return value - maximum value of objective function
                      (-inf for infeasible, inf for unbounded)
17<sub>#define MAXM 400</sub>
                    // leave one extra
```

```
#define MAXN 400 // leave one extra
#define EPS 1e-9
#define INF 1.0/0.0
double A[MAXM][MAXN];
int basis[MAXM], out[MAXN];
void pivot(int m, int n, int a, int b) {
   int i,j;
   for (i=0;i<=m;i++) if (i!=a) for (j=0;j<=n;j++) if (j!=b) {
      A[i][j] -= A[a][j] * A[i][b] / A[a][b];
   for (j=0;j<=n;j++) if (j!=b) A[a][j] /= A[a][b];
   for (i=0;i<=m;i++) if (i!=a) A[i][b] = -A[i][b]/A[a][b];
   A[a][b] = 1/A[a][b];
   i = basis[a];
   basis[a] = out[b];
   out[b] = i;
double simplex(int m, int n, double C[][MAXN], double X[])
   int i,j,ii,jj; // i,ii are row indexes; j,jj are column indexes
   for (i=1;i<=m;i++) for (j=0;j<=n;j++) A[i][j] = C[i][j];
   for (j=0;j<=n;j++) A[0][j] = -C[0][j];
   for (i=0;i<=m;i++) basis[i]=-i;
   for (i=0;i<=n;i++) out[i] = i;
   for(::) {
      for (i=ii=1;i<=m;i++) {
         if (A[i][n]<A[ii][n]
                  || (A[i][n] == A[ii][n] && basis[i] <basis[ii]))
              ii=i:
      if (A[ii][n] >= -EPS) break;
      for (j=jj=0;j< n;j++)
         if (A[ii][j]<A[ii][jj]-EPS
                  || (A[ii][j]<A[ii][j]-EPS && out[i]<out[j]))
      if (A[ii][jj] >= -EPS) return -INF;
      pivot(m,n,ii,ii);
   for(;;) {
      for (j=jj=0;j< n;j++)
         if (A[0][j]<A[0][jj]
                  || (A[0][j] == A[0][jj] && out[j] < out[jj]))
      if (A[0][jj] > -EPS) break;
      for (i=1, ii=0; i <= m; i++)
         if (A[i][ii]>EPS &&
              (!ii | A[i][n]/A[i][jj]<A[ii][n]/A[ii][jj]-EPS |
               (A[i][n]/A[i][jj]<A[ii][n]/A[ii][jj]+EPS
                   && basis[i]<basis[ii])))</pre>
              ii=i;
      if (A[ii][jj] <= EPS) return INF;
      pivot(m,n,ii,jj);
   for (j=0;j< n;j++) X[j] = 0;
   for (i=1;i<=m;i++) if (basis[i] >= 0) X[basis[i]] = A[i][n];
```

```
return A[0][n];
void print(int m, int n, char *msg) { // not used -- debug only
   int i,j;
   printf("%s\n",msq);
   for(i=0;i<=m;i++) {
      for (j=0;j<=m;j++) printf(" %10d",i==j);
      for (j=0;j<=n;j++) printf(" %10g",A[i][j]);
      printf("\n");
   for (i=0;i<=m;i++) printf(" %10d",basis[i]);
   for (j=0;j<n;j++) printf(" %10d",out[j]);
   printf("\n");
/* Gray code. Generates a b-bit gray code starting from 0. */
/* the i'th gray code is i^(i>>1). Magic. */
char *
pbits(char *s, int n, int b) {
 unsigned int i: char *t:
 for( i = 1 << (b-1); i != 0; i >>= 1 ) {
    *s++ = n&i ? '1' : '0';
  *s++ = '\0';
 return t;
/* Search: Golden section Search
  Description: Given an function f(x) with a single local minimum,
                a lower and upper bound on x, and a tolerance for
                convergence, this function finds the value of x
                The function is written globally as f(x)
   Notes:
                - watch out for -0.000 */
#include <stdio.h>
#define GOLD 0.381966
#define move(a,b,c)
                      x[a]=x[b];x[b]=x[c];fx[a]=fx[b];fx[b]=fx[c]
double f(double x){
 return x*x;
double golden (double xlow, double xhigh, double tol) {
 double x[4], fx[4], L;
 int iter = 0, left = 0, mini, i;
  fx[0] = f(x[0]=xlow);
 fx[3] = f(x[3]=xhigh);
  while(1){
   L = x[3]-x[0];
    if(!iter || left){
      x[1] = x[0] + GOLD*L;
      fx[1] = f(x[1]);
    if(!iter || !left){
      x[2] = x[3] - GOLD*L;
```

```
fx[2] = f(x[2]);
    for(mini = 0, i = 1; i < 4; i++)
     if(fx[i] < fx[mini]) mini = i;</pre>
    if(L < tol) break;</pre>
    if(mini < 2)
     left = 1;
     move(3,2,1);
    } else {
     left = 0;
     move(0,1,2);
    iter++;
  return x[mini];
/* Searching: Suffix array
   ______
  Description: Builds a suffix array of a string of N characters
   Complexity: O(N log N)
   Author:
               Howard Cheng
              Oct 30, 2003
  Date:
  References: Manber, U. and Myers, G. "Suffix Arrays: a New
               Method for On-line String Searches."
               SIAM Journal on Computing. 22(5) p. 935-948, 1993.
               T. Kasai, G. Lee, H. Arimura, S. Arikawa, and
               K. Park. "Linear-time Longest-common-prefix
               Computation in Suffix Arrays and Its Applications."
               Proc. 12th Annual Conference on Combinatorial
               Pattern Matching, LNCS 2089, p. 181-192, 2001
   Reliability: 1 (Spain 719 - Glass Beads)
              The build sarray routine takes in a string S of n
              characters (null-terminated), and constructs two
              arrays sarray and lcp. The properties are:
            - If p = sarray[i], then the suffix of str starting at
              p (i.e. S[p..n-1] is the i-th suffix when all the
              suffixes are sorted in lexicographical order
            - NOTE: the empty suffix is not included in this list.
                    so sarray[0] != n.
            - lcp[i] contains the length of the longest common
              prefix of the suffixes pointed to by sarray[i-1]
              and sarray[i]. lcp[0] is defined to be 0.
            - To see whether a pattern P occurs in str, you can
              look for it as the prefix of a suffix. This can be
              done with a binary search in O(|P| \log n) time.
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <limits.h>
```

```
#include <assert.h>
#define MAXN 100000
int bucket[CHAR MAX-CHAR MIN+1];
int prm[MAXN], count[MAXN];
char bh[MAXN+1];
void build sarray(char *str, int* sarray, int *lcp){
 int n, a, c, d, e, f, h, i, j, x;
 n = strlen(str);
 /* sort the suffixes by first character */
 memset(bucket, -1, sizeof(bucket));
 for (i = 0; i < n; i++) {
   j = str[i] - CHAR MIN;
   prm[i] = bucket[j];
   bucket[j] = i;
  for (a = c = 0; a \le CHAR MAX - CHAR MIN; a++) {
   for (i = bucket[a]; i != -1; i = i) {
     j = prm[i];
     prm[i] = c;
     bh[c++] = (i == bucket[a]);
 bh[n] = 1;
  for (i = 0; i < n; i++)
   sarray[prm[i]] = i;
  /* inductive sort */
 x = 0;
 for (h = 1; h < n; h *= 2)
   for (i = 0; i < n; i++) {
     if (bh[i] & 1) {
       x = i:
       count[x] = 0;
     prm[sarray[i]] = x;
   d = n - h;
   e = prm[d];
   prm[d] = e + count[e];
    count[e]++;
   bh[prm[d]] |= 2;
   i = 0:
     for (j = i; (j == i || !(bh[j] & 1)) &  j < n; j++) {
       d = sarray[j] - h;
       if (d >= 0) {
         e = prm[d];
         prm[d] = e + count[e];
         count[e]++;
         bh[prm[d]] |= 2;
```

```
for (j = i; (j == i || !(bh[j] & 1)) &  j < n; j++) {
       d = sarray[j] - h;
       if (d >= 0 \&\& bh[prm[d]] \& 2) {
         for (e = prm[d]+1; bh[e] == 2; e++);
         for (f = prm[d]+1; f < e; f++) {
           bh[f] \&= 1;
        = j;
    for (i = 0; i < n; i++) {
     sarray[prm[i]] = i;
     if (bh[i] == 2) {
       bh[i] = 3;
  for (i = 0; i < n; i++) {
   e = prm[i];
   if (e > 0) {
     j = sarray[e-1];
     while (str[i+h] == str[j+h]) {
       h++;
     lcp[e] = h;
     if (h > 0) {
       h--;
  lcp[0] = 0;
int main(){
  char S[MAXN]; int sarray[MAXN], lcp[MAXN], i;
  char T[MAXN];
  int n, j;
  while (scanf("%s", S) == 1) {
   n = strlen(S);
   for(i = 0; i < n; i++) S[n+i] = S[i];
   S[n+n] = 0;
   build sarray(S, sarray, lcp);
   for (i = 0; S[i]; i++)
     if(sarray[i] < n){</pre>
        printf("%3d: %2d [%d]\n", i, lcp[i], n);
        for(j = 0; j < n; j++){
         printf("%c", S[sarray[i]+j]);
       printf("\n");
  return 0;
```



```
/* Graph Theory: Articulation Points in a Graph (Adj List version)
   ______
  Description: An articulation point in a undirected graph is a
               vertex which disconnects the graph when removed.
               This routine takes a graph represented by an
               adjacency list, and finds all articulation points in
               the graph.
               If b is an articulation point, then there exists
               two distinct vertices a and c, such that b is on
               EVERY path from a to c.
               - An array called ART[] contains 1 if node i is an
                 articulation point and 0 otherwise.
               - use addEdge(int x, int y) to create a undirected
                 edge between node x and y
               - use clearList() to clear out all elements of the
                 adiacency list
  Complexity: O(V+E)
  Author.
               Gilbert Lee
               Oct 23, 2002 (mod Jan 16, 2003)
  Date:
  References: Algorithms in C, Robert Sedgewick pg.440-441
  Reliability: 0
               - Nodes must be indexed from 0..N-1.
  Notes:
               - Edges should not be added more than once to
                 prevent array overflow
               - On graphs which are not connected, nodes which
                 disconnect components they are part of are
                 considered articulation points.
#include <stdio.h>
#include <string.h>
#define MAXN 200
\#define min(a,b) (((a)<(b))?(a):(b))
typedef struct{
 int deg:
 int adi[MAXN];
} Node;
Node alist[MAXN];
char ART[MAXN], val[MAXN];
int id;
void addEdge(int x, int y) {
 alist[x].adj[alist[x].deq++] = y;
 alist[y].adj[alist[y].deg++] = x;
void clearList(){
 memset(alist, 0, sizeof(alist));
int visit(int x, int root){
```

```
int i, y, m, res, child = 0;
 res = val[x] = ++id;
 for (i = 0; i < alist[x].deg; i++)
   y = alist[x].adj[i];
   if(!val[v]){
     if (root && ++child > 1) ART[x] = 1;
     m = visit(v, 0);
     res = min(res, m);
    if(m \ge val[x] \&\& !root) ART[x] = 1;
     res = min(val[y], res);
 return res;
void articulate(int n){
 int i;
 memset(ART, 0, sizeof(ART));
 memset(val, 0, sizeof(val));
 for (id = i = 0; i < n; i++)
   if(!val[i]) visit(i, 1);
int main(){
 int i, n, m, x, y, found;
 /* Read in number of vertices, number of edges */
 while (scanf("%d %d", &n, &m) == 2)
   /* Read in edge between node x and node y */
   for(i = 0; i < m; i++) {
     scanf("%d %d", &x, &y);
     addEdge(x,y);
   /* Find articulation points */
   articulate(n);
   for (found = i = 0; i < n; i++)
     if(ART[i]){
       printf("Node %d is an articulation point\n", i);
       found = 1:
   if(!found) printf("No articulation points\n");
   clearList();
 return 0;
/* Graph Theory: Maximum Weighted Bipartite Matching
  Combinatorics: Assignment Problem
   _____
  Description: Given N workers and N jobs to complete, where each worker has a
               certain compatibility (weight) to each job, find an assignment
               (perfect matching) of workers to jobs which maximizes the
               compatibility (weight).
```

```
Complexity: O(n^3), where n is the number of workers or jobs.
  Author:
               Jason Klaus
  Date:
               February 18, 2004
  References: www.cs.umd.edu/class/fall2003/cmsc651/lec07.ps
             - W is a 2 dimensional array where W[i][i] is the weight of
                 worker i doing job j. Weights must be non-negative. If
                 there is no weight assigned to a particular worker and job
                 pair, set it to zero. If there is a different number of
                 workers than jobs, create dummy workers or jobs accordingly
                 with zero weight edges.
               - M is a 1 dimensional array populated by the algorithm where
                 M[i] is the index of the job matched to worker i.
               - This algorithm could be used on non-negative floating point
                 weights as well.
*/
#include <stdio.h>
/* Maximum number of workers/jobs */
#define MAX N 100
int W[MAX N] [MAX N], U[MAX N], V[MAX N], Y[MAX N]; /* <-- weight variables */
int M[MAX N], N[MAX N], P[MAX N], Q[MAX N], R[MAX N], S[MAX N], T[MAX N];
/st Returns the maximum weight, with the perfect matching stored in M. st/
int Assign(int n)
 int w, y; /* <-- weight variables */
 int i, j, m, p, q, s, t, v;
 for (i = 0; i < n; i++) {
   M[i] = N[i] = -1;
   U[i] = V[i] = 0;
   for (j = 0; j < n; j++)
     if (W[i][j] > U[i])
       U[i] = W[i][j];
 for (m = 0; m < n; m++)
   for (p = i = 0; i < n; i++) {
     T[i] = 0;
     Y[i] = -1;
     if (M[i] == -1) {
       S[i] = 1:
       P[p++] = i;
     else
       S[i] = 0;
    while (1) {
     for (q = s = 0; s < p; s++) {
       i = P[s];
```

```
for (j = 0; j < n; j++)
       if (!T[j]) {
         y = U[i] + V[j] - W[i][j];
         if (y == 0) {
           R[i] = i;
           if (N[i] == -1)
            goto end phase;
           T[j] = 1;
           Q[q++] = j;
         else if ((Y[j] == -1) | | (y < Y[j])) 
           Y[j] = y;
           R[j] = i;
   if (q == 0) {
     y = -1;
     for (j = 0; j < n; j++)
       if (!T[j] \&\& ((y == -1) || (Y[j] < y)))
         y = Y[i];
     for (j = 0; j < n; j++) {
       if (T[j])
         V[j] += y;
       if (S[j])
         U[j] -= y;
     for (j = 0; j < n; j++)
      if (!T[j]) {
         Y[j] -= y;
         if (Y[j] == 0) {
           if (N[j] == -1)
             goto end phase;
           T[j] = 1;
           Q[q++] = j;
    for (p = t = 0; t < q; t++) {
    i = N[O[t]];
     S[i] = 1:
     P[p++] = i;
end phase:
 i = R[i];
 v = M[i];
 M[i] = j;
 N[j] = i;
```

```
while (v != -1) {
    j = v;
    i = R[j];
     v = M[i];
     M[i] = j;
     N[j] = i;
 for (i = w = 0; i < n; i++)
   W += W[i][M[i]];
 return w;
int main()
 int w; /* <-- weight variables */
 int n, i, j;
 while ((scanf("%d", &n) == 1) && (n != 0)) {
   for (i = 0; i < n; i++) {
     for (j = 0; j < n; j++) {
       scanf("%d", &W[i][j]);
   w = Assign(n);
   printf("Optimum weight: %d\n", w);
   printf("Matchings:\n");
   for (i = 0; i < n; i++) {
     printf("%d matched to %d\n", i, M[i]);
 return 0;
```