Contents

1 Geometry - 2D Primitives

1.1 Basics

```cpp
#pragma once
double eps = 1e-6;
typedef complex<double> point;
struct circle {
  point c;
  double r;
  circle(point c, double r) : c(c), r(r) {}
  circle() {}
};
double cross(const point &a, const point &b) { return imag(conj(a) * b
↪  ); }
double dot(const point &a, const point &b) { return real(conj(a) * b);
↪  }
```

1.2 Area of intersection of two circles

```
#include "Basics.cpp"
double circ_inter_area(circle &a, circle &b) {
  double d = abs(b.c - a.c);
  if(d <= b.r - a.r) return a.r * a.r * M_PI;
  if(d <= a.r - b.r) return b.r * b.r * M_PI;
  if(d >= a.r + b.r) return 0;
  double alpha = acos((a.r * a.r + d * d - b.r * b.r) / (2 * a.r * d));
  double beta = acos((b.r * b.r + d * d - a.r * a.r) / (2 * b.r * d));
  return a.r * a.r * (alpha - 0.5 * sin(2 * alpha)) + b.r * b.r * (beta
  ↪  - 0.5 * sin(2 * beta));
}
```

1.3 Points of intersection of two circles

```
#include "Basics.cpp"
// Intersects two circles and intersection points are in 'inter'
// -1-> outside, 0-> inside, 1-> tangent, 2-> 2 intersections
int circ_circ_inter(circle &a, circle &b, vector<point> &inter) {
  double d2 = norm(b.c - a.c), rS = a.r + b.r, rD = a.r - b.r;
  if(d2 > rS * rS) return -1;
  if(d2 < rD * rD) return 0;
  double ca = 0.5 * (1 + rS * rD / d2);
  point z = point(ca, sqrt(a.r * a.r / d2 - ca * ca));
  inter.push_back(a.c + (b.c - a.c) * z);
  if(abs(z.imag()) > eps) inter.push_back(a.c + (b.c - a.c) * conj(z));
  return inter.size();
}
```

1.4 Line-circle intersection

```
#include "Basics.cpp"
// Intersects (infinite) line a-b with circle c
// Intersection points are in 'inter'
// 0 -> no intersection, 1 -> tangent, 2 -> two intersections
int line_circ_inter(point a, point b, circle c, vector<point> &inter) {
  c.c -= a;
  b -= a;
  point m = b * real(c.c / b);
  double d2 = norm(m - c.c);
  if(d2 > c.r * c.r) return 0;
  double l = sqrt((c.r * c.r - d2) / norm(b));
  inter.push_back(a + m + l * b);
  if(abs(l) > eps) inter.push_back(a + m - l * b);
  return inter.size();
}
```

1.5 Line-line intersection

```
#include "Basics.cpp"
// Intersects point of lines a-b and c-d
// -1->coincide,0->parallel,1->intersected(inter. point in 'p')
int line_line_inter(point a, point b, point c, point d, point &p) {
  if(abs(cross(b - a, d - c)) > eps) {
    p = cross(c - a, d - c) / cross(b - a, d - c) * (b - a) + a;
    return 1;
  }
  if(abs(cross(b - a, b - c)) > eps) return 0;
  return -1;
}
```

1.6 Segment-segment intersection

```
#include "Line-line intersection.cpp"
// Intersect of segments a-b and c-d
//   -2  -> not parallel and no intersection
//   -1  -> coincide with no common point
//    0  -> parallel and not coincide
//    1  -> intersected ('p' is intersection of segments)
//    2  -> coincide with common points ('p' is one of the end
//           points lying on both segments)
int seg_seg_inter(point a, point b, point c, point d, point &p) {
  int s = line_line_inter(a, b, c, d, p);
  if(s == 0) return 0;
  if(s == -1) {
    // '<-eps' excludes endpoints in the coincide case
    if(dot(a - c, a - d) < eps) {
      p = a;
      return 2;
    }
    if(dot(b - c, b - d) < eps) {
      p = b;
      return 2;
    }
    if(dot(c - a, c - b) < eps) {
      p = c;
      return 2;
    }
    return -1;
  }
  // '<-eps' excludes endpoints in intersected case
  if(dot(p - a, p - b) < eps && dot(p - c, p - d) < eps) return 1;
  return -2;
}
```

1.7 Parabola-line intersection

```cpp
#include "Basics.cpp"
// Find intersection of the line d-e and the parabola that
// is defined by point 'p' and line a-b
// Returns the number of intersections
// 'ans' has intersection points
int parabola_line_inter(point p, point a, point b, point d, point e,
↪  vector<point> &ans) {
  b = b - a;
  p /= b;
  a /= b;
  d /= b;
  e /= b;
  a -= p;
  d -= p;
  e -= p;
  point n = (e - d) * point(0, 1);
  double c = -dot(n, e);
  if(abs(n.imag()) < eps) {
    if(abs(a.imag()) > eps) {
      double x = -c / n.real();
      ans.push_back(point(x, a.imag() / 2 - x * x / (2 * a.imag())));
    }
  } else {
    double aa = 1;
    double bb = -2 * a.imag() * n.real() / n.imag();
    double cc = -2 * a.imag() * c / n.imag() - a.imag() * a.imag();
    double delta = bb * bb - 4 * aa * cc;
    if(delta > -eps) {
      if(delta < 0) delta = 0;
      delta = sqrt(delta);
      double x = (-bb + delta) / (2 * aa);
      ans.push_back(point(x, (-c - n.real() * x) / n.imag()));
      if(delta > eps) {
        double x = (-bb - delta) / (2 * aa);
        ans.push_back(point(x, (-c - n.real() * x) / n.imag()));
      }
    }
  }
  for(int i = 0; i < ans.size(); i++) ans[i] = (ans[i] + p) * b;
  return ans.size();
}
```

1.8 Circle described by 3 points

```cpp
#include "Basics.cpp"
// Returns whether they form a circle or not.
// 'center' and 'r' contain the circle if there is one
bool get_circle(point p1, point p2, point p3, point &center, double &r)
↪  {
  double g = 2 * imag(conj(p2 - p1) * (p3 - p2));
  if(abs(g) < eps) return false;
  center = p1 * (norm(p3) - norm(p2));
  center += p2 * (norm(p1) - norm(p3));
  center += p3 * (norm(p2) - norm(p1));
  center /= point(0, g);
  r = abs(p1 - center);
  return true;
}
```

1.9 Circle described by 3 lines

```cpp
#include "Line-line intersection.cpp"
// Returns number of circles that are tangent to all three lines
// 'cirs' has all possible circles with radius > 0
// It has zero circles when two of them are coincide
// It has two circles when only two of them are parallel
// It has four circles when they form a triangle. In this case
// first circle is incircle. Next circles are ex-circles tangent
// to edge a,b,c of triangle respectively.
int get_circle(point a1, point a2, point b1, point b2, point c1, point
↪  c2, vector<circle> &cirs) {
  point a, b, c;
  int sa = line_line_inter(a1, a2, b1, b2, c);
  int sb = line_line_inter(b1, b2, c1, c2, a);
  int sc = line_line_inter(c1, c2, a1, a2, b);
  if(sa == -1 || sb == -1 || sc == -1) return 0;
  if(sa + sb + sc == 0) return 0;
  if(sb == 0) {
    swap(a1, c1);
    swap(a2, c2);
  }
  if(sc == 0) {
    swap(b1, c1);
    swap(b2, c2);
  }
  sa = line_line_inter(a1, a2, b1, b2, c);
  line_line_inter(b1, b2, c1, c2, a);
  line_line_inter(c1, c2, a1, a2, b);
  if(sa == 0) {
    point v1 = polar(1.0, (arg(a2 - a1) + arg(a - b)) / 2) + b;
    point v2 = polar(1.0, (arg(a1 - a2) + arg(a - b)) / 2) + b;
    point v3 = polar(1.0, (arg(b2 - b1) + arg(a - b)) / 2) + a;
    point v4 = polar(1.0, (arg(b1 - b2) + arg(a - b)) / 2) + a;
```

```
    point p;
    if(line_line_inter(b, v1, a, v3, p) == 0) swap(v3, v4);
    line_line_inter(b, v1, a, v3, p);
    circle c1, c2;
    c1.c = p;
    line_line_inter(b, v2, a, v4, p);
    c2.c = p;
    c1.r = c2.r = abs(((a1 - b1) / (b2 - b1)).imag() * abs(b2 - b1)) /
    ↪  2;
    cirs.push_back(c1);
    cirs.push_back(c2);
  } else {
    if(abs(a - b) < eps) return 0;
    point bisec1[4][2];
    point bisec2[4][2];
    bisec1[0][0] = polar(1.0, (arg(c - a) + arg(b - a)) / 2);
    bisec1[0][1] = a;
    bisec2[0][0] = polar(1.0, (arg(c - b) + arg(a - b)) / 2);
    bisec2[0][1] = b;
    bisec1[1][0] = polar(1.0, (arg(c - a) + arg(b - a)) / 2);
    bisec1[1][1] = a;
    bisec2[1][0] = polar(1.0, (arg(c - b) + arg(b - a)) / 2);
    bisec2[1][1] = b;
    bisec1[2][0] = polar(1.0, (arg(a - b) + arg(c - b)) / 2);
    bisec1[2][1] = b;
    bisec2[2][0] = polar(1.0, (arg(a - c) + arg(c - b)) / 2);
    bisec2[2][1] = c;
    bisec1[3][0] = polar(1.0, (arg(b - c) + arg(a - c)) / 2);
    bisec1[3][1] = b;
    bisec2[3][0] = polar(1.0, (arg(b - a) + arg(a - c)) / 2);
    bisec2[3][1] = c;
    for(int i = 0; i < 4; i++) {
      point p;
      line_line_inter(bisec1[i][1], bisec1[i][1] + bisec1[i][0], bisec2
      ↪  [i][1], bisec2[i][1] + bisec2[i][0], p);
      circle c1;
      c1.c = p;
      c1.r = abs(((p - a) / (b - a)).imag()) * abs(b - a);
      cirs.push_back(c1);
    }
  }
  return cirs.size();
}
```

1.10 Circle described by 2 points and 1 line

```
#include "Parabola-line intersection.cpp"
// Returns number of circles that pass through point a and b and
// are tangent to the line c-d
// 'ans' has all possible circles with radius > 0
int get_circle(point a, point b, point c, point d, vector<circle> &ans)
↪  {
  point pa = (a + b) / 2.0;
  point pb = (b - a) * point(0, 1) + pa;
  vector<point> ta;
  parabola_line_inter(a, c, d, pa, pb, ta);
  for(point p : ta) ans.push_back(circle(p, abs(a - p)));
  return ans.size();
}
```

1.11 Circle described by 2 lines and 1 point

```
#include "Line-line intersection.cpp"
#include "Parabola-line intersection.cpp"
// Returns number of circles that pass through point p and are
// tangent to the lines a-b and c-d
// 'ans' has all possible circles with radius greater than zero
int get_circle(point p, point a, point b, point c, point d, vector<
↪  circle> &ans) {
  point inter;
  int st = line_line_inter(a, b, c, d, inter);
  if(st == -1) return 0;
  d -= c;
  b -= a;
  vector<point> ta;
  if(st == 0) {
    point pa = point(0, imag((a - c) / d) / 2) * d + c;
    point pb = b + pa;
    parabola_line_inter(p, a, a + b, pa, pb, ta);
  } else {
    if(abs(inter - p) > eps) {
      point bi;
      bi = polar(1.0, (arg(b) + arg(d)) / 2) + inter;
      vector<point> temp;
      parabola_line_inter(p, a, a + b, inter, bi, temp);
      ta.insert(ta.end(), temp.begin(), temp.end());
      temp.clear();
      bi = polar(1.0, (arg(b) + arg(d) + M_PI) / 2) + inter;
      parabola_line_inter(p, a, a + b, inter, bi, temp);
      ta.insert(ta.end(), temp.begin(), temp.end());
    }
  }
  for(point pt : ta) ans.push_back(circle(pt, abs(p - pt)));
  return ans.size();
}
```

2 Geometry - 2D Misc

2.1 Heron's formula for triangle area

```
// Given side lengths a, b, c, returns area or -1 if triangle is
// impossible
double area_heron(double a, double b, double c) {
  if(a < b) swap(a, b);
  if(a < c) swap(a, c);
  if(b < c) swap(b, c);
  if(a > b + c) return -1;
  return sqrt((a + b + c) * (c - a + b) * (c + a - b) * (a + b - c) /
   ↪ 16.0);
}
```

2.2 Rectangle in rectangle test

```
// Can rectangle with dims x*y fit inside box with dims w*h?
// Returns true for a "tight fit", if false is desired then swap
// strictness of inequalities.
bool rect_in_rect(double x, double y, double w, double h) {
  if(x > y) swap(x, y);
  if(w > h) swap(w, h);
  if(w < x) return false;
  if(y <= h) return true;
  double a = y * y - x * x;
  double b = x * h - y * w;
  double c = x * w - y * h;
  return a * a <= b * b + c * c;
}
```

2.3 Centroid and area of a simple polygon $O(N)$

```
#include "Basics.cpp"
// Points must be oriented (CW or CCW), and non-convex is OK
// Returns (nan,nan) if area of polygon is zero
point centroid(vector<point> p) {
  int n = p.size(); // should be at least 1
  double area = 0;
  point c(0, 0); // Not required for area of polygon
  for(int i = n - 1, j = 0; j < n; i = j++) {
    double a = cross(p[i], p[j]) / 2;
    area += a;
    c += (p[i] + p[j]) * (a / 3);
  }
  c /= area;
  return c; // or return 'area' for the area of polygon
}
```

2.4 Point in polygon $O(N)$

```
#include "Basics.cpp"
// outside -> 0, inside -> 1, on the border -> 2
int pt_in_poly(const vector<point> &p, const point &a) {
  int n = p.size();
  int inside = false;
  for(int i = 0, j = n - 1; i < n; j = i++) {
    if(abs(cross(a - p[i], a - p[j])) < eps && dot(a - p[i], a - p[j])
     ↪ < eps) return 2;
    if((imag(p[i]) <= imag(a) && imag(a) < imag(p[j])) || (imag(p[j])
     ↪ <= imag(a) && imag(a) < imag(p[i])))
      if(real(a) - real(p[i]) < (real(p[j]) - real(p[i])) * (imag(a) -
       ↪ imag(p[i])) / (imag(p[j]) - imag(p[i]))) inside = !inside;
  }
  return inside;
}
```

2.5 Convex-hull $O(N \log N)$

```
#include "Basics.cpp"
// Assumes pts.size()>0 and returns ccw convex hull with no
// 3 collinear points and with duplicated left most side node
int comp(const point &a, const point &b) {
  if(abs(a.real() - b.real()) > eps) return a.real() < b.real();
  if(abs(a.imag() - b.imag()) > eps) return a.imag() < b.imag();
  return 0;
}
inline vector<point> convexhull(vector<point> &pts) {
  sort(pts.begin(), pts.end(), comp);
  vector<point> lower, upper;
  for(int i = 0; i < (int)pts.size(); i++) {
    // <-eps include all points on border
    while(lower.size() >= 2 && cross(lower.back() - lower[lower.size()
     ↪ - 2], pts[i] - lower.back()) < eps) lower.pop_back();
    // >eps include all points on border
    while(upper.size() >= 2 && cross(upper.back() - upper[upper.size()
     ↪ - 2], pts[i] - upper.back()) > -eps) upper.pop_back();
    lower.push_back(pts[i]);
    upper.push_back(pts[i]);
  }
  lower.insert(lower.end(), upper.rbegin() + 1, upper.rend());
  return lower;
}
```

3 Geometry - 3D

3.1 Primitives

```
const double eps = 1e-6;
struct point3 {
  double x, y, z;
  point3(double x = 0, double y = 0, double z = 0) : x(x), y(y), z(z) {}
  point3 operator+(point3 p) const { return point3(x + p.x, y + p.y, z + p.z); }
  point3 operator*(double k) const { return point3(k * x, k * y, k * z); }
  point3 operator-(point3 p) const { return *this + (p * -1.0); }
  point3 operator/(double k) const { return *this * (1.0 / k); }
  double norm() { return x * x + y * y + z * z; }
  double abs() { return sqrt(norm()); }
  point3 normalize() { return *this / this->abs(); }
};
// dot product
double dot(point3 a, point3 b) { return a.x * b.x + a.y * b.y + a.z * b.z; }
// cross product
point3 cross(point3 a, point3 b) { return point3(a.y * b.z - b.y * a.z, b.x * a.z - a.x * b.z, a.x * b.y - b.x * a.y); }
struct line {
  point3 a, b;
  line(point3 A = point3(), point3 B = point3()) : a(A), b(B) {}
  // Direction unit vector a -> b
  point3 dir() { return (b - a).normalize(); }
};
// Returns closest point on an infinite line u to the point p
point3 cpoint_iline(line u, point3 p) {
  point3 ud = u.dir();
  return u.a - ud * dot(u.a - p, ud);
}
// Returns Shortest distance between two infinite lines u and v
double dist_ilines(line u, line v) { return dot(v.a - u.a, cross(u.dir(), v.dir()).normalize()); }
// Finds the closest point on infinite line u to infinite line v
// Note: if (uv*uv - uu*vv) is zero then the lines are parallel
// and such a single closest point does not exist. Check for
// this if needed.
point3 cpoint_ilines(line u, line v) {
  point3 ud = u.dir();
  point3 vd = v.dir();
  double uu = dot(ud, ud), vv = dot(vd, vd), uv = dot(ud, vd);
  double t = dot(u.a, ud) - dot(v.a, ud);
  t *= vv;
  t -= uv * (dot(u.a, vd) - dot(v.a, vd));
  t /= (uv * uv - uu * vv);
  return u.a + ud * t;
}
// Closest point on a line segment u to a given point p
point3 cpoint_lineseg(line u, point3 p) {
  point3 ud = u.b - u.a;
  double s = dot(u.a - p, ud) / ud.norm();
  if(s < -1.0) return u.b;
  if(s > 0.0) return u.a;
  return u.a - ud * s;
}
struct plane {
  point3 n, p;
  plane(point3 ni = point3(), point3 pi = point3()) : n(ni), p(pi) {}
  plane(point3 a, point3 b, point3 c) : n(cross(b - a, c - a).normalize()), p(a) {}
  // Value of d for the equation ax + by + cz + d = 0
  double d() { return -dot(n, p); }
};
// Closest point on a plane u to a given point p
point3 cpoint_plane(plane u, point3 p) { return p - u.n * (dot(u.n, p) + u.d()); }
// Point of intersection of an infinite line v and a plane u.
// Note: if dot(u.n, vd) == 0 then the line and plane do not
// intersect at a single point. Check for this if needed.
point3 iline_isect_plane(plane u, line v) {
  point3 vd = v.dir();
  return v.a - vd * ((dot(u.n, v.a) + u.d()) / dot(u.n, vd));
}
// Infinite line of intersection between two planes u and v.
// Note: if dot(v.n, uvu) == 0 then the planes do not intersect
// at a line. Check for this case if it is needed.
line isect_planes(plane u, plane v) {
  point3 o = u.n * -u.d(), uv = cross(u.n, v.n);
  point3 uvu = cross(uv, u.n);
  point3 a = o - uvu * ((dot(v.n, o) + v.d()) / (dot(v.n, uvu) * uvu.norm()));
  return line(a, a + uv);
}
// Returns great circle distance (lat[-90,90], long[-180,180])
double greatcircle(double lt1, double lo1, double lt2, double lo2, double r) {
  double a = M_PI * (lt1 / 180.0), b = M_PI * (lt2 / 180.0);
  double c = M_PI * ((lo2 - lo1) / 180.0);
```

```
  return r * acos(sin(a) * sin(b) + cos(a) * cos(b) * cos(c));
}
// Rotates point p around directed line a->b with angle 'theta'
point3 rotate(point3 a, point3 b, point3 p, double theta) {
  point3 o = cpoint_iline(line(a, b), p);
  point3 perp = cross(b - a, p - o);
  return o + perp * sin(theta) + (p - o) * cos(theta);
}
```

3.2 3D Convex-hull $O(N^2)$

```
#include "Primitives.cpp"
// vector<hullFinder::hullFace> hull=hullFinder(pts).findHull();
// 'hull' will have triangular faces of convex-hull of the given
// points 'pts'. Some of them might be co-planar.
// There are O(pts.size()) of those disjoint triangles that
// cover all surface of convex hull
// Each element of hull is a hullFace which has indices of three
// vertices of a triangle
bool operator==(const point3 &p, const point3 &q) { return abs(p.x - q.
↪  x) < eps && abs(p.y - q.y) < eps && abs(p.z - q.z) < eps; }
point3 triNormal(const point3 &a, const point3 &b, const point3 &c) {
↪  return cross(a, b) + cross(b, c) + cross(c, a); }
class hullFinder {
  const vector<point3> &pts;
public:
  hullFinder(const vector<point3> &pts_) : pts(pts_), halfE(pts.size(),
  ↪  -1) {}
  struct hullFace {
    int u, v, w;
    point3 n;
    hullFace(int u_, int v_, int w_, const point3 &n_) : u(u_), v(v_),
    ↪  w(w_), n(n_) {}
  };
  vector<hullFinder::hullFace> findHull() {
    vector<hullFace> hull;
    int n = pts.size();
    if(n < 4) return hull;
    int p3 = 2;
    point3 tNorm;
    while(p3 < n && (tNorm = triNormal(pts[0], pts[1], pts[p3])) ==
    ↪  point3()) ++p3;
    int p4 = p3 + 1;
    while(p4 < n && abs(dot(tNorm, pts[p4] - pts[0])) < eps) ++p4;
    if(p4 >= n) return hull;
    edges.clear();
    edges.push_front(hullEdge(0, 1));
    setF1(edges.front(), p3);
    setF2(edges.front(), p3);
    edges.push_front(hullEdge(1, p3));
    setF1(edges.front(), 0);
    setF2(edges.front(), 0);
    edges.push_front(hullEdge(p3, 0));
    setF1(edges.front(), 1);
    setF2(edges.front(), 1);
    addPt(p4);
    for(int i = 2; i < n; ++i)
      if(i != p3 && i != p4) addPt(i);
    for(list<hullEdge>::const_iterator e = edges.begin(); e != edges.
    ↪  end(); ++e) {
      if(e->u < e->v && e->u < e->f1)
        hull.push_back(hullFace(e->u, e->v, e->f1, e->n1));
      else if(e->v < e->u && e->v < e->f2)
        hull.push_back(hullFace(e->v, e->u, e->f2, e->n2));
    }
    return hull;
  }
private:
  struct hullEdge {
    int u, v, f1, f2;
    point3 n1, n2;
    hullEdge(int u_, int v_) : u(u_), v(v_), f1(-1), f2(-1) {}
  };
  list<hullEdge> edges;
  vector<int> halfE;
  void setF1(hullEdge &e, int f1) {
    e.f1 = f1;
    e.n1 = triNormal(pts[e.u], pts[e.v], pts[e.f1]);
  }
  void setF2(hullEdge &e, int f2) {
    e.f2 = f2;
    e.n2 = triNormal(pts[e.v], pts[e.u], pts[e.f2]);
  }
  void addPt(int i) {
    for(list<hullEdge>::iterator e = edges.begin(); e != edges.end();
    ↪  ++e) {
      bool v1 = dot(pts[i] - pts[e->u], e->n1) > eps;
      bool v2 = dot(pts[i] - pts[e->u], e->n2) > eps;
      if(v1 && v2)
        e = --edges.erase(e);
      else if(v1) {
        setF1(*e, i);
        addCone(e->u, e->v, i);
      } else if(v2) {
```

```
        setF2(*e, i);
        addCone(e->v, e->u, i);
      }
    }
  }
  void addCone(int u, int v, int apex) {
    if(halfE[v] != -1) {
      edges.push_front(hullEdge(v, apex));
      setF1(edges.front(), u);
      setF2(edges.front(), halfE[v]);
      halfE[v] = -1;
    } else
      halfE[v] = u;
    if(halfE[u] != -1) {
      edges.push_front(hullEdge(apex, u));
      setF1(edges.front(), v);
      setF2(edges.front(), halfE[u]);
      halfE[u] = -1;
    } else
      halfE[u] = v;
  }
};
```

4 Combinatorics

4.1 (Un)Ranking of K-combination out of N $O(K \log N)$

```
const int maxn = 100;
const int maxk = 10;
// combination[i][j] = j!/(i!*(j-i)!)
long long combination[maxk][maxn];
long long cumsum[maxk][maxn];
void initialize() { //~O(nk)
  memset(combination, 0, sizeof combination);
  for(int i = 0; i < maxn; i++) combination[0][i] = 1;
  for(int i = 1; i < maxk; i++)
    for(int j = 1; j < maxn; j++) combination[i][j] = combination[i][j
    ↪  - 1] + combination[i - 1][j - 1];
  for(int i = 0; i < maxk; i++) cumsum[i][0] = combination[i][0];
  for(int i = 0; i < maxk; i++)
    for(int j = 1; j < maxn; j++) cumsum[i][j] = cumsum[i][j - 1] +
    ↪  combination[i][j];
}
// Returns rank of the given combination 'c' of n objects.
long long rank_comb(int n, vector<int> c) {
  long long ans = 0;
  int prev = -1;
  sort(c.begin(), c.end()); // comment this if it is sorted
  for(int i = 0; i < c.size(); i++) {
    ans += cumsum[c.size() - i - 1][n - prev - 2] - cumsum[c.size() - i
    ↪  - 1][n - c[i] - 1];
    prev = c[i];
  }
  return ans;
}
struct comp {
  long long base;
  comp(long long base) : base(base) {}
  int operator()(const long long &a, const long long &val) { return (
  ↪  base - a) > val; }
};
// Returns k-combination of rank 'r' of n objects
vector<int> unrank_comb(int n, int k, long long r) {
  vector<int> c;
  int prev = -1;
  for(int i = 0; i < k; i++) {
    long long base = cumsum[k - i - 1][n - prev - 2];
    prev = n - 1 - (lower_bound(cumsum[k - i - 1], cumsum[k - i - 1] +
    ↪  n - prev - 1, r, comp(base)) - cumsum[k - i - 1]);
    r -= base - cumsum[k - i - 1][n - prev - 1];
    c.push_back(prev);
  }
  return c;
}
```

4.2 (Un)Ranking of K-permutation out of N $O(K)$

```
void rec_unrank_perm(int n, int k, long long r, vector<int> &id, vector
↪  <int> &pi) {
  if(k > 0) {
    swap(id[n - 1], id[r % n]);
    rec_unrank_perm(n - 1, k - 1, r / n, id, pi);
    pi.push_back(id[n - 1]);
    swap(id[n - 1], id[r % n]);
  }
}
// Returns a k-permutation corresponds to rank 'r' of n objects.
// 'id' should be a full identity permutation of size at least n
// and it remains the same at the end of the function
vector<int> unrank_perm(int n, int k, long long r, vector<int> &id) {
  vector<int> ans;
  rec_unrank_perm(n, k, r, id, ans);
  return ans;
}
```

```
long long rec_rank_perm(int n, int k, vector<int> &pirev, vector<int> & pi) {
  if(k == 0) return 0;
  int s = pi[k - 1];
  swap(pi[k - 1], pi[pirev[n - 1] - (n - k)]);
  swap(pirev[s], pirev[n - 1]);
  long long ans = s + n * rec_rank_perm(n - 1, k - 1, pirev, pi);
  swap(pirev[s], pirev[n - 1]);
  swap(pi[k - 1], pi[pirev[n - 1] - (n - k)]);
  return ans;
}
// Returns rank of the k-permutaion 'pi' of n objects.
// 'id' should be a full identity permutation of size at least n
// and it remains the same at the end of the function
long long rank_perm(int n, vector<int> &id, vector<int> pi) {
  for(int i = 0; i < pi.size(); i++) id[pi[i]] = i + n - pi.size();
  long long ans = rec_rank_perm(n, pi.size(), id, pi);
  for(int v : pi) id[v] = v;
  return ans;
}
```

4.3 Digit occurrence count $O(\log n)$

```
// Given digit d and value N, returns # of times d occurs from 1..N
long long digit_count(int digit, int N) {
  long long res = 0;
  char buff[15];
  int i, count;
  if(N <= 0) return 0;
  res += N / 10 + ((N % 10) >= digit ? 1 : 0);
  if(digit == 0) res--;
  res += digit_count(digit, N / 10 - 1) * 10;
  sprintf(buff, "%d", N / 10);
  for(i = 0, count = 0; i < strlen(buff); i++)
    if(buff[i] == digit + '0') count++;
  res += (1 + N % 10) * count;
  return res;
}
```

4.4 Josephus Ring Survivor

```
/* Josephus Ring Survivor (n people, dismiss every m'th) */
const int MaxN = 1000;
int survive[MaxN];
void josephus(int n, int m) {
  survive[1] = 0;
  for(int i = 2; i <= n; i++) survive[i] = (survive[i - 1] + (m % i)) % i;
}
```

4.5 Derangement

```
// combinatorial: derangement
// count the number of permutations of n elements, such that no element app
// math formula: derange(n)=ceil(factorial(n)/e+0.5), probability of derang
const int maxN = 21;
long long derange[maxN];
long long cal_derange(int n) {
  derange[0] = 1;
  derange[1] = 0;
  for(int i = 2; i <= n; i++) derange[i] = (i - 1) * (derange[i - 1] + derange[i - 2]);
  return derange[n];
}
```

5 Graph Theory

5.1 Fast flow $O(V^2E)$

```
// find_flow returns max flow from s to t in an n-vertex graph.
// Use add_edge to add edges (directed/undirected) to the graph.
// Call clear_flow() before each testcase.
const int maxn = 1000;
int c[maxn][maxn];
vector<int> adj[maxn];
int par[maxn];
int dcount[maxn + maxn];
int dist[maxn];
void add_edge(int a, int b, int cap, int rev_cap = 0) {
  c[a][b] += cap;
  c[b][a] += rev_cap;
  adj[a].push_back(b);
  adj[b].push_back(a);
}
void clear_flow() {
  memset(c, 0, sizeof c);
  memset(dcount, 0, sizeof dcount);
  for(int i = 0; i < maxn; ++i) adj[i].clear();
}
int advance(int v) {
  for(int w : adj[v]) {
    if(c[v][w] > 0 && dist[v] == dist[w] + 1) {
      par[w] = v;
      return w;
    }
  }
```

```
  return -1;
}
int retreat(int v) {
  int old = dist[v];
  --dcount[dist[v]];
  for(int w : adj[v]) {
    if(c[v][w] > 0) dist[v] = min(dist[v], dist[w]);
  }
  ++dist[v];
  ++dcount[dist[v]];
  if(dcount[old] == 0) return -1;
  return par[v];
}
int augment(int s, int t) {
  int delta = c[par[t]][t];
  for(int v = t; v != s; v = par[v]) delta = min(delta, c[par[v]][v]);
  for(int v = t; v != s; v = par[v]) {
    c[par[v]][v] -= delta;
    c[v][par[v]] += delta;
  }
  return delta;
}
queue<int> q;
void bfs(int v) {
  memset(dist, -1, sizeof dist);
  while(!q.empty()) q.pop();
  q.push(v);
  dist[v] = 0;
  ++dcount[dist[v]];
  while(!q.empty()) {
    v = q.front();
    q.pop();
    for(int w : adj[v]) {
      if(c[w][v] > 0 && dist[w] == -1) {
        dist[w] = dist[v] + 1;
        ++dcount[dist[w]];
        q.push(w);
      }
    }
  }
}
int find_flow(int n, int s, int t) {
  bfs(t);
  int v = s;
  par[s] = s;
  int ans = 0;
  while(v != -1 && dist[s] < n) {
    int newv = advance(v);
    if(newv != -1)
      v = newv;
    else
      v = retreat(v);
    if(v == t) {
      v = s;
      ans += augment(s, t);
    }
  }
  return ans;
}
```

5.2 Flow and negative flow

```
const int inf = (int)1e9;
const int maxn = 300;
int x[maxn][maxn], m;
int c[maxn][maxn], n;
int f[maxn][maxn];
int flow_k, flow_t, mark[maxn];
int dfs(int v, int m) {
  if(v == flow_t) return m;
  for(int i = 0, x; i < n; ++i)
    if(c[v][i] - f[v][i] >= flow_k && !mark[i]++)
      if(x = dfs(i, min(m, c[v][i] - f[v][i]))) return (f[i][v] = -(f[v
      ↪ ][i] += x)), x;
  return 0;
}
// Input: n(# of vertices),s(source),t(sink),c[n][n](capacities)
// Finds flow from i to j (i.e. f[i][j]) in the maximum flow
// where f[i][j]=-f[j][i]
// Requirements: f[i][j] should be filled with initial flow
// before calling the function and c[i][j] >= f[i][j]
void flow(int s, int t) {
  int flow_ans = 0;
  flow_t = t;
  flow_k = 1;
  for(int i = 0; i < n; ++i)
    for(int j = 0; j < n; ++j)
      for(; flow_k < c[i][j]; flow_k *= 2)
        ;
  for(; flow_k; flow_k /= 2) {
    memset(mark, 0, sizeof mark);
    for(; dfs(s, inf);) memset(mark, 0, sizeof mark);
  }
}
```

```
// Input: m(# of vertices), x[m][m](capacities)
// Finds f[i][j] in a circular flow satisfying x[i][j]
// If you have a real sink and source set x[sink][source]=inf
// x[i][j]<0 means capacity of i->j is zero and a flow of at
// least abs(x[i][j]) should go from j to i.
// If you have two capacities for i->j and j->i and some
// min flow for at least one of them you should resolve this
// before calling the function by filling some flow in f[i][j]
// and f[j][i]
// Returns false when can't satisfy x and returns false when
// x[i][j] and x[j][i] are both negative. Check this if needed
bool negative_flow() {
  for(int i = 0; i < m; ++i)
    for(int j = 0; j < m; ++j) {
      if(x[i][j] < 0) {
        if(x[j][i] < 0) return false;
        continue;
      }
      if(x[j][i] >= 0) {
        c[i][j] = x[i][j];
        continue;
      }
      c[i][j] = x[i][j] + x[j][i];
      c[j][i] = 0;
      c[i][m + 1] -= x[j][i];
      c[m][j] -= x[j][i];
      if(c[i][j] < 0) return false;
    }
  n = m + 2;
  flow(n - 2, n - 1);
  for(int i = 0; i < m; ++i)
    if(c[m][i] != f[m][i]) return false;
  for(int i = 0; i < m; ++i)
    for(int j = 0; j < m; ++j)
      if(x[i][j] < 0) {
        f[i][j] += x[i][j];
        f[j][i] -= x[i][j];
      }
  return true;
}
```

5.3 Min cost max flow

```
// Input (zero based, non-negative edges):
// n = |V|, e = |E|, s = source, t = sink
// cost[v][u] = cost for each unit of flow from v to u
// cap[v][u] = capacity
// Output of mcf():
// Flow contains the flow value
// Cost contains the minimum cost
// f[n][n] contains the flow
const int maxn = 300;
const int inf = 1e9;
int cap[maxn][maxn], cost[maxn][maxn], f[maxn][maxn];
int p[maxn], d[maxn], mark[maxn], pi[maxn];
int n, s, t, Flow, Cost;
int pot(int u, int v) { return d[u] + pi[u] - pi[v]; }
int dijkstra() {
  memset(mark, 0, sizeof mark);
  memset(p, -1, sizeof p);
  for(int i = 0; i <= n; i++) d[i] = inf;
  d[s] = 0;
  // Doesn't use a priority queue due to it not really improving
  // the algorithm - will still be O(n^2)
  while(1) {
    int u = n;
    for(int i = 0; i < n; i++)
      if(!mark[i] && d[i] < d[u]) u = i;
    if(u == n) break;
    mark[u] = 1;
    for(int v = 0; v < n; v++) {
      if(!mark[v] && f[v][u] && d[v] > pot(u, v) - cost[v][u]) {
        d[v] = pot(u, v) - cost[v][u];
        p[v] = u;
      }
      if(!mark[v] && f[u][v] < cap[u][v] && d[v] > pot(u, v) + cost[u][
      ↪ v]) {
        d[v] = pot(u, v) + cost[u][v];
        p[v] = u;
      }
    }
  }
  for(int i = 0; i < n; i++)
    if(pi[i] < inf) pi[i] += d[i];
  return mark[t];
}
void mcf() {
  memset(f, 0, sizeof f);
  memset(pi, 0, sizeof pi);
  Flow = Cost = 0;
  while(dijkstra()) {
    int min = inf;
    for(int x = t; x != s; x = p[x])
      if(f[x][p[x]])
```

```
      min = std::min(f[x][p[x]], min);
    else
      min = std::min(cap[p[x]][x] - f[p[x]][x], min);
  for(int x = t; x != s; x = p[x])
    if(f[x][p[x]]) {
      f[x][p[x]] -= min;
      Cost -= min * cost[x][p[x]];
    } else {
      f[p[x]][x] += min;
      Cost += min * cost[p[x]][x];
    }
  Flow += min;
 }
}
```

5.4 2-Sat & strongly connected component $O(V + E)$

```
// Vertices are numbered 0..n-1 for true states.
// False state of the variable i is i+n (i.e. other(i))
// For SCC 'n', 'adj' and 'adjrev' need to be filled.
// For 2-Sat set 'n' and use add_edge
// 0<=val[i]<=1 is the value for binary variable i in 2-Sat
// 0<=group[i]<2*n is the scc number of vertex i.
const int maxn = 1000;
int n;
vector<int> adj[maxn * 2];
vector<int> adjrev[maxn * 2];
int val[maxn];
int marker, dfst, dfstime[maxn * 2], dfsorder[maxn * 2];
int group[maxn * 2];
// For 2SAT Only
inline int other(int v) { return v < n ? v + n : v - n; }
inline int var(int v) { return v < n ? v : v - n; }
inline int type(int v) { return v < n ? 1 : 0; }
void satclear() {
  for(int i = 0; i < maxn + maxn; i++) {
    adj[i].resize(0);
    adjrev[i].resize(0);
  }
}
void dfs(int v) {
  if(dfstime[v] != -1) return;
  dfstime[v] = -2;
  int deg = adjrev[v].size();
  for(int i = 0; i < deg; i++) dfs(adjrev[v][i]);
  dfstime[v] = dfst++;
}
void dfsn(int v) {
  if(group[v] != -1) return;
  group[v] = marker;
  int deg = adj[v].size();
  for(int i = 0; i < deg; i++) dfsn(adj[v][i]);
}
// For 2SAT Only
void add_edge(int a, int b) {
  adj[other(a)].push_back(b);
  adjrev[b].push_back(other(a));
  adj[other(b)].push_back(a);
  adjrev[a].push_back(other(b));
}
bool solve() {
  dfst = 0;
  memset(dfstime, -1, sizeof dfstime);
  for(int i = 0; i < n + n; i++) dfs(i);
  memset(val, -1, sizeof val);
  for(int i = 0; i < n + n; i++) dfsorder[n + n - dfstime[i] - 1] = i;
  memset(group, -1, sizeof group);
  for(int i = 0; i < n + n; i++) {
    marker = i;
    dfsn(dfsorder[i]);
  }
  // For 2SAT Only
  for(int i = 0; i < n; i++) {
    if(group[i] == group[i + n]) return 0;
    val[i] = (group[i] > group[i + n]) ? 0 : 1;
  }
  return 1;
}
```

5.5 Bipartite matching, vertex cover, edge cover, disjoint set $O(VE)$

```
// Input:
//    n: size of part1, m: size of part2
//    a[i]: neighbours of i-th vertex of part1
//    b[i]: neighbours of i-th vertex of part2
const int maxn = 2020, maxm = 2020;
int n, m;
vector<int> a[maxn], b[maxm];
int matched[maxn], mark[maxm], mate[maxm];
bool dfs(int v) {
  if(v < 0) return 1;
  for(int to : a[v])
    if(!mark[to]++ && dfs(mate[to])) return matched[mate[to] = v] = 1;
  return 0;
}
```

```
void set_mark() {
  memset(matched, 0, sizeof matched);
  memset(mate, -1, sizeof mate);
  memset(mark, 0, sizeof mark);
  for(int i = 0; i < n; ++i)
    for(int to : a[i])
      if(mate[to] < 0) {
        matched[mate[to] = i] = 1;
        break;
      }
  for(int i = 0; i < n; ++i)
    if(!matched[i] && dfs(i)) memset(mark, 0, sizeof mark);
  for(int i = 0; i < n; ++i)
    if(!matched[i]) dfs(i);
}
// res.size(): size of matching
// res[i]: i-th edge of matching
// res[i].first is in part1, res[i].second is in part2
void matching(vector<pair<int, int>> &res) {
  set_mark();
  res.clear();
  for(int i = 0; i < m; ++i)
    if(mate[i] >= 0) res.push_back(make_pair(mate[i], i));
}
// p1: vertices in part1, p2: vertices in part2
// union of p1 and p2 cover the edges of the graph
void vertex_cover(vector<int> &p1, vector<int> &p2) {
  set_mark();
  p1.clear();
  p2.clear();
  for(int i = 0; i < m; ++i)
    if(mate[i] >= 0)
      if(mark[i])
        p2.push_back(i);
      else
        p1.push_back(mate[i]);
}
// p1: vertices in part1, p2: vertices in part2
// union of p1 and p2 is the largest disjoint set of the graph
void disjoint_set(vector<int> &p1, vector<int> &p2) {
  set_mark();
  p1.clear();
  p2.clear();
  for(int i = 0; i < m; ++i)
    if(mate[i] >= 0 && mark[i])
      p1.push_back(mate[i]);
    else
      p2.push_back(i);
  for(int i = 0; i < n; ++i)
    if(!matched[i]) p1.push_back(i);
}
// edges in res cover the vertices of the graph
// res[i].first is in part1, res[i].second is in part2
void edge_cover(vector<pair<int, int>> &res) {
  set_mark();
  res.clear();
  for(int i = 0; i < m; ++i)
    if(mate[i] >= 0)
      res.push_back(make_pair(mate[i], i));
    else if(b[i].size())
      res.push_back(make_pair(b[i][0], i));
  for(int i = 0; i < n; ++i)
    if(!matched[i] && a[i].size()) res.push_back(make_pair(i, a[i][0
    ↪ ]));
}
```

5.6 Bipartite weighted matching $O(VE^2)$

```
// Input: n, m, w[n][m] (n <= m)
//        w[i][j] is the weight between the i-th vertex of part1
//        and the j-th vertex of part2. w[i][j] can be any
//        integer (including negative values)
// Output: res, size of res is n
const int inf = 1e7;
const int maxn = 200, maxm = 200;
int n, m, w[maxn][maxm], u[maxn], v[maxm];
int mark[maxn], mate[maxm], matched[maxn];
int dfs(int x) {
  if(x < 0) return 1;
  if(mark[x]++) return 0;
  for(int i = 0; i < m; i++)
    if(u[x] + v[i] - w[x][i] == 0)
      if(dfs(mate[i])) return matched[mate[i] = x] = 1;
  return 0;
}
void _2matching() {
  memset(mate, -1, sizeof mate);
  memset(mark, 0, sizeof mark);
  memset(matched, 0, sizeof matched);
  for(int i = 0; i < n; i++)
    for(int j = 0; j < m; j++)
      if(mate[j] < 0 && u[i] + v[j] - w[i][j] == 0) {
        matched[mate[j] = i] = 1;
        break;
```

```
    }
  for(int i = 0; i < n; i++)
    if(!matched[i])
      if(dfs(i)) memset(mark, 0, sizeof mark);
}
void wmatching(vector<pair<int, int>> &res) {
  for(int i = 0; i < m; i++) v[i] = 0;
  for(int i = 0; i < n; i++) {
    u[i] = -inf;
    for(int j = 0; j < m; j++) u[i] = max(u[i], w[i][j]);
  }
  memset(mate, -1, sizeof mate);
  memset(matched, 0, sizeof matched);
  int counter = 0;
  while(counter != n) {
    for(int flag = 1; flag;) {
      flag = 0;
      memset(mark, 0, sizeof mark);
      for(int i = 0; i < n; i++)
        if(!matched[i] && dfs(i)) {
          counter++;
          flag = 1;
          memset(mark, 0, sizeof mark);
        }
    }
    int epsilon = inf;
    for(int i = 0; i < n; i++)
      for(int j = 0; j < m; j++) {
        if(!mark[i]) continue;
        if(mate[j] >= 0)
          if(mark[mate[j]]) continue;
        epsilon = min(epsilon, u[i] + v[j] - w[i][j]);
      }
    for(int i = 0; i < n; i++)
      if(mark[i]) u[i] -= epsilon;
    for(int j = 0; j < m; j++)
      if(mate[j] >= 0)
        if(mark[mate[j]]) v[j] += epsilon;
  }
  res.clear();
  for(int i = 0; i < m; i++)
    if(mate[i] != -1) res.push_back(pair<int, int>(mate[i], i));
}
```

5.7 Cut edges and 2-edge-connected components $O(V+E)$

```
// input (zero based):
//       g[n] should be the adjacency list of the graph
//       g[i] is a vector of int
// output of cut_edge():
//       cut_edges is a vector of pair<int, int>
//       comp[comp_size] contains the 2 connected components
//       comp[i] is a vector of int
const int maxn = 1000;
typedef pair<int, int> edge;
vector<int> g[maxn];
int n, mark[maxn], d[maxn], jad[maxn];
vector<edge> cut_edges;
// for components only
vector<int> comp[maxn];
int comp_size;
vector<int> comp_stack;
void dfs(int x, int level) {
  mark[x] = 1;
  // for components only
  comp_stack.push_back(x);
  int t = 0;
  for(int u : g[x]) {
    if(!mark[u]) {
      jad[u] = d[u] = d[x] + 1;
      dfs(u, level + 1);
      jad[x] = std::min(jad[u], jad[x]);
      if(jad[u] == d[u]) {
        cut_edges.push_back(edge(u, x));
        // for components only
        while(comp_stack.back() != u) {
          comp[comp_size].push_back(comp_stack.back());
          comp_stack.pop_back();
        }
        comp[comp_size++].push_back(u);
        comp_stack.pop_back();
        //
      }
    } else {
      if(d[u] == d[x] - 1) t++;
      if(d[u] != d[x] - 1 || t != 1) jad[x] = std::min(d[u], jad[x]);
    }
  }
  // for components only
  if(level == 0) {
    while(comp_stack.size() > 0) {
      comp[comp_size].push_back(comp_stack.back());
      comp_stack.pop_back();
    }
```

```
    comp_size++;
  }
}
void cut_edge() {
  memset(mark, 0, sizeof mark);
  memset(d, 0, sizeof d);
  memset(jad, 0, sizeof jad);
  cut_edges.clear();
  // for components only
  for(int i = 0; i < maxn; i++) comp[i].clear();
  comp_stack.clear();
  comp_size = 0;
  for(int i = 0; i < n; i++)
    if(!mark[i]) dfs(i, 0);
}
```

5.8 Cut vertices and 2-connected components $O(V+E)$

```
// Input (zerobased):
//       g[n] should be the adjacency list of the graph
//       g[i] is a vector of int
// Output of cut_ver():
//       cut_vertex is a vector of int
//       comp[comp_size] contains the 2 connected components
//       comp[i] is a vector of int
const int maxn = 1000;
vector<int> g[maxn];
int d[maxn], mark[maxn], mark0[maxn], jad[maxn];
int n;
vector<int> cut_vertex;
// for components only
vector<int> comp[maxn];
int comp_size;
vector<int> comp_stack;
void dfs(int x, int level) {
  mark[x] = 1;
  // for components only
  comp_stack.push_back(x);
  for(int u : g[x]) {
    if(!mark[u]) {
      jad[u] = d[u] = d[x] + 1;
      dfs(u, level + 1);
      jad[x] = std::min(jad[u], jad[x]);
      if(jad[u] >= d[x] && d[x]) {
        cut_vertex.push_back(x);
        // for components only
        while(comp_stack.back() != u) {
          comp[comp_size].push_back(comp_stack.back());
          comp_stack.pop_back();
        }
        comp[comp_size].push_back(u);
        comp_stack.pop_back();
        comp[comp_size++].push_back(x);
      }
    } else if(d[u] != d[x] - 1)
      jad[x] = std::min(d[u], jad[x]);
  }
  // for components only
  if(level == 0) {
    while(comp_stack.size() > 0) {
      comp[comp_size].push_back(comp_stack.back());
      comp_stack.pop_back();
    }
    comp_size++;
  }
}
int dfs0(int x) {
  mark0[x] = 1;
  for(int to : g[x])
    if(!mark0[to]) return dfs0(to);
  return x;
}
void cut_ver() {
  memset(mark, 0, sizeof mark);
  memset(mark0, 0, sizeof mark0);
  memset(d, 0, sizeof d);
  memset(jad, 0, sizeof jad);
  // for components only
  for(int i = 0; i < maxn; i++) comp[i].clear();
  comp_stack.clear();
  comp_size = 0;
  cut_vertex.clear();
  for(int i = 0; i < n; i++)
    if(!mark[i]) dfs(dfs0(i), 0);
}
```

5.9 Dijkstra $O(E \log V)$

```
const int maxn = 1000; // Max # of vertices
int n; //# of vertices
vector<pair<int, int>> v[maxn]; // weighted adjacency list
int d[maxn]; // distance from source
struct comp {
  bool operator()(int a, int b) { return (d[a] != d[b]) ? d[a] < d[b] :
  ↪  a < b; }
```

```
};
set<int, comp> mark;
void dijkstra(int source) {
  memset(d, -1, sizeof d);
  d[source] = 0;
  mark.clear();
  for(int i = 0; i < n; ++i) mark.insert(i);
  while(mark.size()) {
    int x = *mark.rbegin();
    mark.erase(x);
    if(d[x] == -1) break;
    for(auto &it : v[x]) {
      if(d[it.first] == -1 || d[x] + it.second < d[it.first]) {
        mark.erase(it.first);
        d[it.first] = d[x] + it.second;
        mark.insert(it.first);
      }
    }
  }
}
```

5.10 Bellman ford with negative cycle detection

```
const int MaxN = 205;
int V;
struct Edge {
  int from, to, cost;
};
vector<Edge> allEdgesFromNode[MaxN];
// MUST be updated in update loop
int predecessor[MaxN];
// If the END of a path is in negative cycle, then no min cost path
bool inNegativeCycle[MaxN];
// Black - No cycle.
// Gray - Is in a cycle
// White - unknown.
const int White = 0, Gray = 1, Black = 2;
int color[MaxN];
// Determines if a node is contained in an infinite cycle
int ExpandPredecessor(int node) {
  if(color[node] != White) return color[node];
  color[node] = Gray;
  // Not part of a cycle at all
  if(predecessor[node] == -1) return color[node] = Black;
  int newColor = ExpandPredecessor(predecessor[node]);
  inNegativeCycle[node] = (newColor == Gray);
  return color[node] = newColor;
}
void ExpandNegativeCycle(int node) {
  inNegativeCycle[node] = true;
  for(Edge& e : allEdgesFromNode[node]) {
    if(!inNegativeCycle[e.to]) ExpandNegativeCycle(e.to);
  }
}
void FinishUpBellmanFord() {
  // Go along the predecessor graph
  for(int i = 0; i < V; ++i) color[i] = White;
  // Find all nodes that are part of a negative cycle
  for(int i = 0; i < V; ++i) ExpandPredecessor(i);
  // Now, expand from all nodes that are in a negative cycle
  // - they cause all children to become negative cycle nodes
  for(int i = 0; i < V; ++i)
    if(inNegativeCycle[i]) ExpandNegativeCycle(i);
}
```

5.11 Minimum spanning tree

```
#define MAXN 1000
#define MAXM 1000000
#define EPS 1e-8
int n;
struct Edge {
  int u, v; /* Edge between u, v with weight w */
  double w;
};
int sets[MAXN];
Edge edge[MAXM], treeedge[MAXN];
int numedge;
int getRoot(int x) {
  if(sets[x] < 0) return x;
  return sets[x] = getRoot(sets[x]);
}
void Union(int a, int b) {
  int ra = getRoot(a);
  int rb = getRoot(b);
  if(ra != rb) {
    sets[ra] += sets[rb];
    sets[rb] = ra;
  }
}
double mintree() {
  double weight = 0.0;
  int i, count;
  sort(edge, edge + numedge, [](auto a, auto b) { return a.w < b.w; });
  for(i = count = 0; count < n - 1; i++) {
```

```
    if(getRoot(edge[i].u) != getRoot(edge[i].v)) {
      Union(edge[i].u, edge[i].v);
      weight += edge[i].w;
      treeedge[count++] = edge[i];
    }
  }
  return weight;
}
```

5.12 Minimum weight Steiner tree $O(|V| * 3^{|S|} + |V|^3)$

```
// Given a weighted undirected graph G = (V, E) and a subset S of V,
// finds a minimum weight tree T whose vertices are a superset of S.
// NP-hard -- this is a pseudo-polynomial algorithm.
// Minimum stc[(1<<s)-1][v] (0 <= v < n) is weight of min. Steiner tree
// Minimum stc[i][v] (0 <= v < n) is weight of min. Steiner tree for
// the i'th subset of Steiner vertices
// S is the list of Steiner vertices, s = |S|
// d is the adjacency matrix (use infinities, not -1), and n = |V|
const int N = 32;
const int K = 8;
int d[N][N], n, S[K], s, stc[1 << K][N];
void steiner() {
  for(int k = 0; k < n; ++k)
    for(int i = 0; i < n; ++i)
      for(int j = 0; j < n; ++j) d[i][j] = min(d[i][k], d[i][k] + d[k][
      ↪  j]);
  for(int i = 1; i < (1 << s); ++i) {
    if(!(i & (i - 1))) {
      int u;
      for(int j = i, k = 0; j; u = S[k++], j >>= 1)
        ;
      for(int v = 0; v < n; ++v) stc[i][v] = d[v][u];
    } else
      for(int v = 0; v < n; ++v) {
        stc[i][v] = 0xffffff;
        for(int j = 1; j < i; ++j)
          if((j | i) == i) {
            int x1 = j, x2 = i & (~j);
            for(int w = 0; w < n; ++w) stc[i][v] = min(stc[i][v], d[v][
            ↪  w] + stc[x1][w] + stc[x2][w]);
          }
      }
  }
}
```

6 Number Theory

6.1 Sieve of Eratosthenes $O(N \log\log N)$

```
// Returns all prime numbers in [0,n]
const int maxn = 1000000;
int isnprime[maxn];
vector<int> sieve(int n) {
  memset(isnprime, 0, sizeof isnprime);
  isnprime[0] = isnprime[1] = 1;
  vector<int> ps;
  for(int i = 2; i < n; i++)
    if(!isnprime[i]) {
      ps.push_back(i);
      if(n / i >= i)
        for(int j = i * i; j <= n; j += i) isnprime[j] = 1;
    }
  return ps;
}
```

6.2 Chinese remaindering and ext. Euclidean $O(N \log \max(M_i))$

```
typedef long long int LLI;
LLI mod(LLI a, LLI m) { return ((a % m) + m) % m; }
// Assumes non-negative input. Returns d such that d=a*ss+b*tt
LLI gcdex(LLI a, LLI b, LLI &ss, LLI &tt) {
  if(b == 0) {
    ss = 1;
    tt = 0;
    return a;
  }
  LLI g = gcdex(b, a % b, tt, ss);
  tt = tt - (a / b) * ss;
  return g;
}
// Returns x such that 0<=x<lcm(m_0, ..., m_(n-1)) and
// x==a_i (mod m_i), if such an x exists. If x does not exist -1
// is returned.
LLI chinese_rem(vector<LLI> &a, vector<LLI> &m) {
  LLI g, s, t, a_tmp, m_tmp;
  a_tmp = mod(a[0], m[0]);
  m_tmp = m[0];
  for(int i = 1; i < a.size(); ++i) {
    g = gcdex(m_tmp, m[i], s, t);
    if((a_tmp - a[i]) % g) return -1;
    a_tmp = mod(a_tmp + (a[i] - a_tmp) / g * s * m_tmp, m_tmp / g * m[i
    ↪  ]);
    m_tmp = m[i] * m_tmp / gcdex(m[i], m_tmp, s, t);
```

```
  }
  return a_tmp;
}
```

6.3 Discrete logarithm solver $O(\sqrt{P})$

```
// Given prime P, B>0, and N, finds least L
// such that B^L==N (mod P)
// Returns -1, if no such L exist.
map<int, int> mow;
int times(int a, int b, int m) { return (long long)a * b % m; }
int power(int val, int power, int m) {
  int res = 1;
  for(int p = power; p; p >>= 1) {
    if(p & 1) res = times(res, val, m);
    val = times(val, val, m);
  }
  return res;
}
int discrete_log(int p, int b, int n) {
  int jump = sqrt(double(p));
  mow.clear();
  for(int i = 0; i < jump && i < p - 1; ++i) mow[power(b, i, p)] = i +
  ↪ 1;
  for(int i = 0, j; i < p - 1; i += jump)
    if(j = mow[times(n, power(b, p - 1 - i, p), p)]) return (i + j - 1)
    ↪ % (p - 1);
  return -1;
}
```

6.4 Euler phi $O(\sqrt{n})$

```
// The Euler Phi function returns the number of
// positive integers less than N that are relatively
// prime to N. O(sqrt(n))
int phi(int n) {
  int i, count, res = 1;
  for(i = 2; i * i <= n; i++) {
    count = 0;
    while(n % i == 0) {
      n /= i;
      count++;
    }
    if(count > 0) res *= (pow(i, count) - pow(i, count - 1));
  }
  if(n > 1) res *= (n - 1);
  return res;
}
```

6.5 Binomial coefficient

```
// binomial coefficient C(n,k)
long long Cdp[51][51]; // memset to -1
long long C(int n, int k) {
  if(Cdp[n][k] != -1) return Cdp[n][k];
  return Cdp[n][k] = (k == 0 || k == n ? 1 : C(n - 1, k - 1) + C(n - 1,
  ↪ k));
}
```

7 String

7.1 Manacher's algorithm $O(N)$

```
// Returns half of length of largest palindrome centered at
// every position in the string
// Add \0 at start, end, and middle to handle palindromes between
// characters + get length of largest palindrome at each index.
// Then, int characterBefore = i / 2; int lenToStart = P[i] / 2;
// int lenToEnd = lenToStart - (i % 2 == 0);
vector<int> manacher(string s) {
  vector<int> ans(s.size(), 0);
  int maxi = 0;
  for(int i = 1; i < s.size(); i++) {
    int k = 0;
    if(maxi + ans[maxi] >= i) k = min(ans[maxi] + maxi - i, ans[2 *
    ↪ maxi - i]);
    for(; s[i + k] == s[i - k] && i - k >= 0 && i + k < s.size(); k++)
      ;
    ans[i] = k - 1;
    if(i + ans[i] > maxi + ans[maxi]) maxi = i;
  }
  return ans;
}
```

7.2 KMP string matching $O(N + M)$

```
// Given strings t and p, return the indices of t where p occurs
// as a substring
vector<int> compute_prefix(string s) {
  vector<int> pi(s.size(), -1);
  int k = -1;
  for(int i = 1; i < s.size(); i++) {
    while(k >= 0 && s[k + 1] != s[i]) k = pi[k];
    if(s[k + 1] == s[i]) k++;
    pi[i] = k;
  }
```

```
  return pi;
}
vector<int> kmp_match(string t, string p) {
  vector<int> pi = compute_prefix(p);
  vector<int> shifts;
  int m = -1;
  for(int i = 0; i < t.size(); i++) {
    while(m > -1 && p[m + 1] != t[i]) m = pi[m];
    if(p[m + 1] == t[i]) m++;
    if(m == p.size() - 1) {
      shifts.push_back(i + 1 - p.size());
      m = pi[m];
    }
  }
  return shifts;
}
```

7.3 Suffix array $O(NlogN)$

```
// Calculate the suffix array for a string. Includes code for
// LCP and how to code up string matching range.
typedef pair<int, int> ii;

const int MaxN = 100010;
char T[MaxN];
int N;
int SA[MaxN], tempSA[MaxN]; // SA[i] = index of suffix i in string
int RA[MaxN], tempRA[MaxN]; // Rank of i in T
int c[MaxN];
void radixSort(int k) {
    int i, maxi = max(300, N);
    memset(c, 0, sizeof c);
    for (i = 0; i < N; ++i)
        c[i + k < N ? RA[i + k] : 0]++; // TODO: Mod for circular
    int sum = 0;
    for (i = 0; i < maxi; ++i) {
        int t = c[i]; c[i] = sum; sum += t;
    }
    for (i = 0; i < N; ++i) {
        // TODO: Mod for circular
        int indexToC = SA[i] + k < N ? RA[SA[i] + k] : 0;
        tempSA[c[indexToC]++] = SA[i];
    }
    for (i = 0; i < N; ++i) SA[i] = tempSA[i];
}

void constructSA() {
    int i;
    for (i = 0; i < N; ++i) RA[i] = T[i];
    for (i = 0; i < N; ++i) SA[i] = i;
    for (int k = 1; k < N; k <<= 1) {
        radixSort(k); radixSort(0);
        int r = 0;
        tempRA[SA[0]] = r;
        for (i = 1; i < N; ++i) {
            tempRA[SA[i]] =
                (RA[SA[i]] == RA[SA[i - 1]] &&
                 RA[(SA[i] + k) %N] == RA[(SA[i-1] + k) %N]) ? r : ++r;
        }
        for (i = 0; i < N; ++i) RA[i] = tempRA[i];
        if (RA[SA[N - 1]] == N - 1) break;
    }
}

// Returns inclusive set of all matches of P[0:pLen] into SA.
// Returns -1, -1 if no matches exist.
char P[MaxN];
ii StringMatching(int pLen) { // O(|P|log|N|)
    int low = 0, high = N-1;
    while (low < high) {
        int mid = (low + high) / 2;
        int result = strncmp(T + SA[mid], P, pLen);
        if (result >= 0) high = mid;
        else low = mid + 1;
    }
    if (strncmp(T + SA[low], P, pLen) != 0) return ii(-1, -1);
    ii ans; ans.first = low;
    low = 0; high = N - 1;
    while (low < high) {
        int mid = (low + high) / 2;
        int result = strncmp(T + SA[mid], P, pLen);
        if (result > 0) high = mid;
        else low = mid + 1;
    }
    if (strncmp(T + SA[high], P, pLen) != 0) --high;
    ans.second = high;
    return ans;
}


int LCP[MaxN];
↪  // LCP[i] = prefix size that SA[i] has in common with SA[i-1]
void ComputeLCP() {
    int Phi[MaxN], PLCP[MaxN], i, L;
```

```
    Phi[SA[0]] = -1;
    for (i = 1; i < N; ++i) Phi[SA[i]] = SA[i-1];
    for (L = i = 0; i < N; ++i) {
        if (Phi[i] == -1) { PLCP[i] = 0; continue; }
        while (T[i + L] == T[Phi[i] + L]) ++L;
        PLCP[i] = L;
        L = max(L - 1, 0);
    }
    for (int i = 0; i < N; ++i) LCP[i] = PLCP[SA[i]];
}
```

7.4 Trie

```
// This Trie is designed for returning the longest
// substring that appears more than once in a string
struct TrieNode {
  int cnt;
  map<char, TrieNode *> child;
  TrieNode() { cnt = 0; }
};
// Need to initialize each time after calling deleteTrie
TrieNode *root = NULL;
string resultString = "";
int resultCnt = 0;
// start is the start point in s
void insertTrie(string &s, int start) {
  if(root == NULL) root = new TrieNode;
  TrieNode *current = root;
  for(int i = start; i < s.size(); i++) {
    if(current->child[s[i]] == NULL) current->child[s[i]] = new
    ↪ TrieNode;
    current->child[s[i]]->cnt++;
    current = current->child[s[i]];
  }
}
// string tmp="";findLongest(root,tmp);
void findLongest(TrieNode *current, string &s) {
  for(auto c : current->child)
    if(c.second->cnt > 1) {
      s.push_back(c.first);
      findLongest(c.second, s);
      s.pop_back();
    }
  if(s.size() > resultString.size()) {
    resultString = s;
    resultCnt = current->cnt;
  }
}
void deleteTrie(TrieNode *current) {
  if(current == NULL) return;
  for(auto c : current->child) deleteTrie(c.second);
  delete current;
}
```

8 Misc

8.1 FFT $O(n \log n)$

```
/* Author: Zachary Friggstad, 2017

   The standard FFT of a sequence.
   Given n = 2^k complex numbers a[i], i=0...n-1 r, this will compute all n
   values \sum_i a[i]*zeta^{i+j} for j=0...n-1 in O(n * log(n)) time where
   zeta = exp(2 \pi i / n) is the standard n'th root of unity.

   Multiplication of polynomials.
   Given two polynomials f,g, will compute the product f*g
   in O(n * log(n)) time where n = max(deg(f), deg(g)).

   This is a numerical algorithm, in the sense calculations are performed
   with doubles and may be very slightly off, but I suspect it will not be
   issue in any contest problem that requires FFT.

   Use:
   fft(f, v, invert):
    - f is a vector of complex values of size 2^k for some k
    - v is a vector of complex values that stores the result
    - invert is a boolean indicating if we should compute the fft or invers

   multiply(f, g, res):
    - f, g are two vectors of complex values representing polynomials
      e.g. to represent x^3 - 3x + 7 the corresponding vector is {7, -3, 0,
      (each being a "complex" type) where 7 has index 0 and 1 has index 3 (
    - res is a complex vector that will store the result

   Note:
    - f and g do not need to have sizes being powers of 2
    - after, we will have res.size() = f.size() + g.size() - 2
    - if f and g have leading 0s, then so will res

   Reference:
   Introduction to Algorithms by Cormen, Lieserson, Rivest, and Stein
   This is the "iterative" version that does not use recursion.

   Reliability:
```

```
   kinversions - Open Kattis
   tiles - icpc.kattis.com (World Finals 2015)
   polymul2 - Open Kattis
   matchings - Open Kattis
*/

typedef double ld; // can always try long double if you are concerned
typedef complex<ld> cplx;
typedef vector<cplx> vc;
typedef vector<int> vi;
/* Compute the fft of f, store in v.
   invert will compute the inverse of the fft.
   f.size() *MUST* be a power of 2
   In particular, the regular fft (invert == false) will not normalize by
   1/f.size() but reverse fft (invert == true) will normalize. This is the easiest
   approach for convolution/polynomial multiplication.
*/
void fft(const vc& f, vc& v, bool invert) {
    int n = f.size();
    assert(n > 0 && (n&(n-1)) == 0);
    v.resize(n);
    for (int i = 0; i < n; ++i) {
        int r = 0, k = i;
        for (int j = 1; j < n; j <<= 1, r = (r<<1)|(k&1), k >>= 1);
        v[i] = f[r];
    }
    for (int m = 2; m <= n; m <<= 1) {
        int mm = m>>1;
        cplx zeta = polar<ld>(1, (invert?2:-2)*M_PI/m);
        for (int k = 0; k < n; k += m) {
            cplx om = 1;
            for (int j = 0; j < mm; ++j, om *= zeta) {
                cplx tl = v[k+j], th = om*v[k+j+mm];
                v[k+j] = tl+th;
                v[k+j+mm] = tl-th;
            }
        }
    }
    if (invert) for (auto& z : v) z /= ld(n);
}

/* Multiply polynomials f and g. Equivalently, compute the convolution of
   the sequences f[0], ..., f[df] and g[0], ..., g[dg]).

   IMPORTANT: if f is the zero polynomial, should still have a 0 entry (i.e.
   f.size() > 0 should always hold). Same for g.
   res - holds the results: the coefficients from res[0] to res[df+dg].

   Can use f = g (reference to same vector) safely.
   f,g are not constant because they are padded with 0s, but then are rever
   to their original form again.
*/
void multiply(vc& f, vc& g, vc& res) {
    int df = f.size()-1, dg = g.size()-1;

    assert(df >= 0 && dg >= 0);

    int n = df+dg+1;
    while (n&(n-1)) ++n;

    f.resize(n,0);
    g.resize(n,0);
    vc tmp;

    fft(f, tmp, false);
    fft(g, res, false);

    for (int i = 0; i < n; ++i) tmp[i] *= res[i];

    fft(tmp, res, true);

    f.resize(df+1);
    g.resize(dg+1);
    res.resize(df+dg+1);
}
```

Given a function $f : \{0, \ldots, n-1\} \to C$ into the complex numbers, we can write the function uniquely as $f(x) = \sum_{k=0}^{n-1} a_k z^k$, where $z = e^{2\pi \frac{i}{n}}$. The function fft takes an input function f (as a vector of complex numbers), and returns the values $a_0, \ldots, a_{n-1}$ in a vector. Conversely, the inverseffft function takes the values $a_0, \ldots, a_{n-1}$, and returns the function f. The Fourier transform is useful whenever we want to compute a convolution, which is a function $h : \{0, \ldots, n-1\} \to C$ defined in terms of two other functions f and g by $h(m) = \sum_{k=0}^{n-1} f(m)g(k-m)$, where the values of m and $k-m$ are considered modulo n. If we write $f = \sum a_k z^k, g = \sum b_k z^k$, and $h = \sum c_k z^k$, then we find $c_k = a_k * z_k$, so that we can compute $h = inverse fft(fft(f) * fft(g))$ in $O(n \log(n))$ time, rather than the naive $O(n^2)$ time. Complexity: $O(n \log(n))$, where n is the size of the domain of the input function. Notes: Watch out for standard floating point inaccuracies, etc. This code only works if your function is defined on a domain of length $[n]$, where n is a power of 2.

8.2 Longest ascending subsequence $O(n \log n)$

```
typedef pair<int, int> pii;
int comp(const pii &a, const pii &b) {
```

```
  if(a.first != b.first) return a.first < b.first;
  return a.second < b.second;
  ↪ // return 0 to find strictly ascending subsequence
}
vector<int> lis(const vector<int> &in) {
  vector<pii> l;
  vector<int> par(in.size(), -1);
  for(int i = 0; i < in.size(); i++) {
    int ind = lower_bound(l.begin(), l.end(), pii(in[i], i), comp) - l.
    ↪ begin();
    if(ind == l.size()) l.push_back(pii(0, 0));
    l[ind] = pii(in[i], i);
    if(ind != 0) par[i] = l[ind - 1].second;
  }
  vector<int> ans;
  int ind = l.back().second;
  while(ind != -1) {
    ans.push_back(in[ind]);
    ind = par[ind];
  }
  reverse(ans.begin(), ans.end());
  return ans;
}
```

8.3 Simplex

```
// m - number of (less than) inequalities
// n - number of variables
// c - (m+1) by (n+1) array of coefficients:
//   row 0       - objective function coefficients
//   row 1:m     - less-than inequalities
//   column 0:n-1 - inequality coefficients
//   column n    - inequality constants (0 for obj. function)
// x[n] - result variables
// Returns value - maximum value of objective function
//   (-inf for infeasible, inf for unbounded)
const int maxm = 400; // leave one extra
const int maxn = 400; // leave one extra
const double eps = 1e-9;
const double inf = 1.0 / 0.0;
double ine[maxm][maxn];
int basis[maxm], out[maxn];
void pivot(int m, int n, int a, int b) {
  int i, j;
  for(i = 0; i <= m; i++)
    if(i != a)
      for(j = 0; j <= n; j++)
        if(j != b) ine[i][j] -= ine[a][j] * ine[i][b] / ine[a][b];
  for(j = 0; j <= n; j++)
    if(j != b) ine[a][j] /= ine[a][b];
  for(i = 0; i <= m; i++)
    if(i != a) ine[i][b] = -ine[i][b] / ine[a][b];
  ine[a][b] = 1 / ine[a][b];
  i = basis[a];
  basis[a] = out[b];
  out[b] = i;
}
double simplex(int m, int n, double c[][maxn], double x[]) {
  int i, j, ii, jj;
  for(i = 1; i <= m; i++)
    for(j = 0; j <= n; j++) ine[i][j] = c[i][j];
  for(j = 0; j <= n; j++) ine[0][j] = -c[0][j];
  for(i = 0; i <= m; i++) basis[i] = -i;
  for(j = 0; j <= n; j++) out[j] = j;
  for(;;) {
    for(i = ii = 1; i <= m; i++)
      if(ine[i][n] < ine[ii][n] || (ine[i][n] == ine[ii][n] && basis[i]
      ↪ < basis[ii])) ii = i;
    if(ine[ii][n] >= -eps) break;
    for(j = jj = 0; j < n; j++)
      if(ine[ii][j] < ine[ii][jj] - eps || (ine[ii][j] < ine[ii][jj] -
      ↪ eps && out[i] < out[j])) jj = j;
    if(ine[ii][jj] >= -eps) return -inf;
    pivot(m, n, ii, jj);
  }
  for(;;) {
    for(j = jj = 0; j < n; j++)
      if(ine[0][j] < ine[0][jj] || (ine[0][j] == ine[0][jj] && out[j] <
      ↪ out[jj])) jj = j;
    if(ine[0][jj] > -eps) break;
    for(i = 1, ii = 0; i <= m; i++)
      if(ine[i][jj] > eps && (!ii || (ine[i][n] / ine[i][jj] < ine[ii][
      ↪ n] / ine[ii][jj] - eps) || (ine[i][n] / ine[i][jj] < ine[ii][
      ↪ n] / ine[ii][jj] + eps && basis[i] < basis[ii]))) ii = i;
    if(ine[ii][jj] <= eps) return inf;
    pivot(m, n, ii, jj);
  }
  for(j = 0; j < n; j++) x[j] = 0;
  for(i = 1; i <= m; i++)
    if(basis[i] >= 0) x[basis[i]] = ine[i][n];
  return ine[0][n];
}
```

8.4 Segment tree $O(\log n)$

```cpp
const int maxn = 1 << 20; // must be a power of 2
long long seg[2 * maxn];
// Add the value 'val' to the index 'num'
void add(int num, long long val) {
  num += maxn;
  while(num > 0) {
    seg[num] += val;
    num >>= 1;
  }
}
// returns sum of the elements in range [0,num]
long long get(int num) {
  num += maxn;
  long long ans = 0;
  ans = seg[num]; // Comment this to change the range to [0,num)
  while(num > 0) {
    if(num & 1) { ans += seg[num & (~1)]; }
    num >>= 1;
  }
  return ans;
}
```

8.5 Lazy segment tree

```cpp
const int MAX_N = 1024000;
struct Node {
  int value;
  // Action will need to be applied to all children
  // Will already have been applied to the node
  // EG: Increase for all numbers in range
  int action;
  bool hasAction;
};
const int NO_ACTION = 0; // TODO?
const int NEUTRAL_QUERY_VALUE = 0; // TODO?
int N;
Node allNodes[5 * MAX_N];
// After generating the segment tree, doesn't use
// the index specific array
int baseValue[MAX_N];
int SubQueryMerge(int lhs_val, int rhs_val) {
  return lhs_val + rhs_val; // TODO?
}
// Inclusive on both
void GenerateSegmentTree(int index, int nodeLeft, int nodeRight) {
  allNodes[index].action = NO_ACTION;
  if(nodeLeft == nodeRight) {
    allNodes[index].value = baseValue[nodeLeft]; // TODO?
    return;
  }
  int mid = (nodeLeft + nodeRight) / 2;
  GenerateSegmentTree(index * 2, nodeLeft, mid);
  GenerateSegmentTree(index * 2 + 1, mid + 1, nodeRight);
  allNodes[index].value = SubQueryMerge(allNodes[index * 2].value,
    allNodes[index * 2 + 1].value);
}
void AddLazyUpdateAction(int index, int action) {
  allNodes[index].action += action;
    // TODO - Handle multiple different lazy updates
  allNodes[index].hasAction = true;
}
void ApplyAndPushLazyUpdate(int index, int nodeStart, int nodeEnd) {
  if(!allNodes[index].hasAction) return;
  allNodes[index].value += allNodes[index].action; // TODO: Apply
  if(nodeStart != nodeEnd) {
    int middle = (nodeStart + nodeEnd) / 2;
    // Tell children about their lazy status
    AddLazyUpdateAction(index * 2, allNodes[index].action);
    AddLazyUpdateAction(index * 2 + 1, allNodes[index].action);
  }
  allNodes[index].action = NO_ACTION;
  allNodes[index].hasAction = false;
}
// Inclusive on both starts and ends
void ApplyLazyChange(int index, int nodeStart, int nodeEnd, int
  changeStart, int changeEnd, int action) {
  // Make sure the value is updated and moved to children
  ApplyAndPushLazyUpdate(index, nodeStart, nodeEnd);
  if(nodeEnd < changeStart || nodeStart > changeEnd) { return; }
  // This index is contained completely
  if(nodeStart >= changeStart && nodeEnd <= changeEnd) {
    // Add the update to this node, then apply
    // it so parent will get correct value.
    AddLazyUpdateAction(index, action);
    ApplyAndPushLazyUpdate(index, nodeStart, nodeEnd);
    return;
  }
  int middle = (nodeStart + nodeEnd) / 2;
  ApplyLazyChange(index * 2, nodeStart, middle, changeStart, changeEnd,
    action);
  ApplyLazyChange(index * 2 + 1, middle + 1, nodeEnd, changeStart,
    changeEnd, action);
```

```
  allNodes[index].value = SubQueryMerge(allNodes[index * 2].value,
   ↪  allNodes[index * 2 + 1].value);
}
// Inclusive on both starts and ends
int Query(int index, int nodeStart, int nodeEnd, int queryStart, int
↪  queryEnd) {
  if(nodeEnd < queryStart || nodeStart > queryEnd) return
   ↪  NEUTRAL_QUERY_VALUE;
  // Make sure the value is updated and moved to children
  ApplyAndPushLazyUpdate(index, nodeStart, nodeEnd);
  // This index is contained completely
  if(nodeStart >= queryStart && nodeEnd <= queryEnd) { return allNodes[
   ↪  index].value; }
  int middle = (nodeStart + nodeEnd) / 2;
  int count = SubQueryMerge(Query(index * 2, nodeStart, middle,
   ↪  queryStart, queryEnd), Query(index * 2 + 1, middle + 1, nodeEnd,
   ↪  queryStart, queryEnd));
  return count;
}
```

8.6 Equation solving $O(NM(N+M)$

```
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a > -eps); }
// m = number of equations, n = number of variables,
// a[m][n+1] = coefficients matrix
// Returns double ans[n] containing the solution, if there is no
// solution returns NULL
double *solve(double **a, int m, int n) {
  int cur = 0;
  for(int i = 0; i < n; ++i) {
    for(int j = cur; j < m; ++j)
      if(!zero(a[j][i])) {
        if(j != cur) swap(a[j], a[cur]);
        for(int sat = 0; sat < m; ++sat) {
          if(sat == cur) continue;
          double num = a[sat][i] / a[cur][i];
          for(int sot = 0; sot <= n; ++sot) a[sat][sot] -= a[cur][sot]
           ↪  * num;
        }
        cur++;
        break;
      }
  }
  for(int j = cur; j < m; ++j)
    if(!zero(a[j][n])) return NULL;
  double *ans = new double[n];
  for(int i = 0, sat = 0; i < n; ++i) {
    ans[i] = 0;
    if(sat < m && !zero(a[sat][i])) {
      ans[i] = a[sat][n] / a[sat][i];
      sat++;
    }
  }
  return ans;
}
```

8.7 Cubic equation solver

```
// Solves ax^3 + bx^2 + cx + d = 0
vector<double> solve_cubic(double a, double b, double c, double d) {
  long double a1 = b / a, a2 = c / a, a3 = d / a;
  long double q = (a1 * a1 - 3 * a2) / 9.0, sq = -2 * sqrt(q);
  long double r = (2 * a1 * a1 * a1 - 9 * a1 * a2 + 27 * a3) / 54.0;
  double z = r * r - q * q * q, theta;
  vector<double> res;
  res.clear();
  if(z <= 0) {
    theta = acos(r / sqrt(q * q * q));
    res.push_back(sq * cos(theta / 3.0) - a1 / 3.0);
    res.push_back(sq * cos((theta + 2.0 * M_PI) / 3.0) - a1 / 3.0);
    res.push_back(sq * cos((theta + 4.0 * M_PI) / 3.0) - a1 / 3.0);
    return res;
  }
  double v = pow(sqrt(z) + fabs(r), 1 / 3.0);
  v += q / v;
  v *= (r < 0) ? 1 : -1;
  v -= a1 / 3.0;
  res.push_back(v);
  return res;
}
```

8.8 Calendar

```
import java.io.BufferedReader;
import java.io.InputStreamReader;
import java.util.GregorianCalendar;
import java.util.Scanner;
class Main {
    public static void main (String args[]) {
        Scanner scanner = new Scanner(System.in);
        int T = scanner.nextInt();
        // Different option for input.
        BufferedReader inputHandler = new BufferedReader(new
         ↪  InputStreamReader(System.in));
        try {
```

```
            String line = inputHandler.readLine();
            // Do more
        } catch (Exception e) { // Better not reach here
        }
    }
    GregorianCalendar createCal(int year, int zeroIndexMonth, int day, int dayIncrease)
        GregorianCalendar cal = new GregorianCalendar(year, zeroIndexMonth, day);
        cal.add(Calendar.DAY_OF_MONTH, dayIncrease);
        return cal;
    }
}
```

8.9 Unbounded maximum stack size

```
/* Sometimes you want to recurse very deeply or
   allocate massive arrays on the stack. This code removes the limit
   on maximum stack size so that you can.
   Disclaimer:  This is usually not a good idea. You can solve any
   problem without increasing the max stack size. However, doing so
   can save you some time in a few cases. */
#include <sys/resource.h>
void remove_stack_limit() {
  struct rlimit rl;
  getrlimit(RLIMIT_STACK, &rl);
  rl.rlim_cur = RLIM_INFINITY;
  setrlimit(RLIMIT_STACK, &rl);
}
```

8.10 C++ IO Format

```
#include <bits/stdc++.h> // include everything
freopen("test.in","r",stdin);
freopen("test.out","w",stdout);
cout << fixed << setprecision(7) << M_PI << endl; // 3.1415927
cout << scientific << M_PI << endl; // 3.1415927e+000
int x=15, y=12094;
cout << setbase(10) << x << " " << y << endl; // 15 12094
cout << setbase(8)  << x << " " << y << endl; // 17 27476
cout << setbase(16) << x << " " << y << endl; // f 2f3e
x=5; y=9;
cout << setfill('0') << setw(2) << x << ":" << setw(2) << y <<
endl; // 05:09
printf ("%10d\n", 111); //       111
printf ("%010d\n", 111); //0000000111
printf ("%d %x %X %o\n", 200, 200, 200, 200); //200 c8 C8 310
printf ("%010.2f %e %E\n", 1213.1416, 3.1416, 3.1416);
//0001213.14 3.141600e+00 3.141600E+00
printf ("%*.*d\n",10, 5, 111); //     00111
printf ("%-*.*d\n",10, 5, 111); //00111
printf ("%+*.*d\n",10, 5, 111); //    +00111
char in[20]; int d;
scanf ("%s %*s %d",in,&d); //<- it's number 5
printf ("%s %d \n", in,d); //it's 5
```

8.11 Formulas

Pick's Theorem: $A = i + \frac{b}{2} - 1$ (A:area, i:interior, b:boundary points)

Catalan numbers: $C_n = \frac{1}{n+1}\binom{2n}{n} = \frac{4i-2}{i+1}C_{n-1} = \sum_{i=0}^{n-1} C_i C_{n-1-i}, C_0 = 1$

Triangle: $c^2 = a^2+b^2-2ab\cos(\theta_c)$, $s = \frac{1}{2}(a+b+c)$, $inradius = \sqrt{\frac{(s-a)(s-b)(s-c)}{s}}$, $exradii_a = \sqrt{\frac{s(s-b)(s-c)}{s-a}}$

Spherical cap: $V = \frac{\pi h}{6}(3a^2+h^2)$, $A = 2\pi rh$ (a:radius of base of cap, r: radius of sphere, h: height of cap)

8.12 Common bugs

```
* READ THE STATEMENT AGAIN. TELL YOUR TEAMMATE IF NECESSARY
* Double check spell of literals
* Graph: Multiple components, Multiple edges, Loops
* Geometry: Be careful about +pi,-pi
* Initialization: Use memset/clear(). Don't expect global
variables to be zero. Care about multiple tests.
* Precision and Range: Use long long if necessary. Use
BigInteger/BigDecimal
* Derive recursive formulas that use sum instead of
multiplication to avoid overflow.
* Small cases (n=0,1,negative)
* 0-based <=> 1-based
* Division by zero. Integer division a/(double)b
* Stack overflow (DFS on 1e5)
* Infinite loop?
* array bound check. maxn or x*maxn
* Don't use .size()-1 !
* \(int)-3 < (unsigned int) 2" is false!
* Check copy-pasted codes!
* Be careful about -0.0
* Remove debug info!
* Output format: Spaces at the end of line. Blank lines. View
the output in VIM if necessary
```

```
* Add eps to double before getting floor or round
* If setting a dp table to Inf when executing to avoid returning to
↪  this state, should look at using B/DFS instead!
```