

# University of Alberta

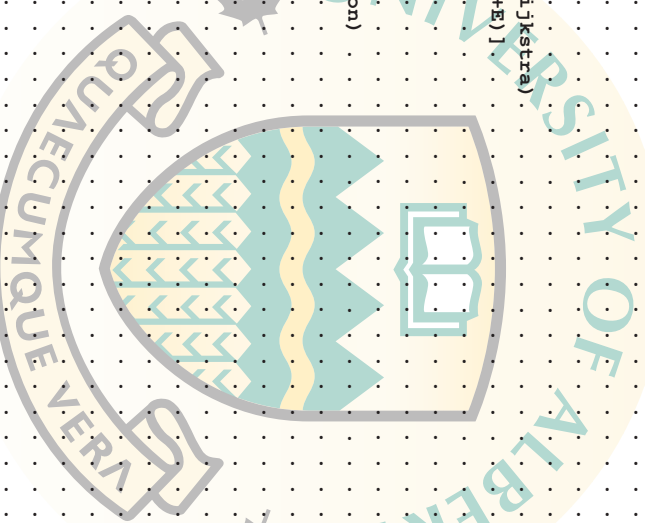
## 2007 ACM ICPC World Finals

## Code Archive



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```
/* Geometry: Complex Arithmetic -----*/
```

```
// These two values are used in most of the geometry algorithms.
```

```
double PI = 2*acos(0.0);
```

```
double EPS = 1E-8;
```

```
struct pol {
```

```
    double r, t;
```

```
    pol(double R = 0, double T = 0) : r(R), t(T) {}
```

```
};
```

```
struct point {
```

```
    double x, y;
```

```
    point(double X = 0, double Y = 0) : x(X), y(Y) {}
```

```
    point(const pol &P) : x(P.r*cos(P.t)), y(P.r*sin(P.t)) {}
```

```
    point conj() const { return point(x, -y); }
```

```
    double mag2() const { return x*x + y*y; }
```

```
    double mag() const { return sqrt(mag2()); }
```

```
    double arg() const { return atan2(y, x); }
```

```
    point operator-() const { return point(-x, -y); }
```

```
    point& operator+=(const point &a) { x += a.x; y += a.y; return *this; }
```

```
    point& operator-=(const point &s) { x -= s.x; y -= s.y; return *this; }
```

```
    point& operator*=(const point &m) {
```

```
        double tx = x*m.x - y*m.y, ty = x*m.y + y*m.x;
```

```
        x = tx; y = ty; return *this;
```

```
    }
```

```
    point& operator/=(const point &d) {
```

```
        double tx = y*d.y + x*d.x, ty = y*d.x - x*d.y, t = d.mag2();
```

```
        x = tx/t; y = ty/t; return *this;
```

```
    }
```

```
    bool operator<(const point &q) const {
```

```
        if (fabs(y-q.y) < EPS) return x < q.x;
```

```
        return y < q.y;
```

```
    }
```

```
    bool operator==(const point &q) const {
```

```
        return (fabs(x-q.x) < EPS) && (fabs(y-q.y) < EPS);
```

```
    }
```

```
    bool operator!=(const point &q) const { return !operator==(q); }
```

```
};
```

```
point operator+(point a, const point &b) { return a += b; }
```

```
point operator-(point a, const point &b) { return a -= b; }
```

```
point operator*(point a, const point &b) { return a *= b; }
```

```
point operator/(point a, const point &b) { return a /= b; }
```

```
.....
```

```
/* Geometry: Area of a polygon (positive <-> CCW orientation) -----*/
```

```
double areaPoly(vector<point> &p) {
```

```
    double sum = 0; int n = p.size();
```

```
    for (int i = n-1, j = 0; j < n; i = j++)
```

```
        sum += (p[i].conj()*p[j]).y;
```

```
    return sum/2;
```

```
}
```

```
.....
```

```
/* Geometry: Heron's formula for triangle area -----*/
```

```
// Given side lengths a, b, c, returns area or -1 if triangle is impossible
```

```
double area_heron(double a, double b, double c) {
```

```
    if (a < b) swap(a, b);
```

```
    if (a < c) swap(a, c);
```

```
    if (b < c) swap(b, c);
```

```
    if ((c-(a-b)) < 0) return -1;
```

```
    return sqrt((a+(b+c))*(c-(a-b))*(c+(a-b))*(a+(b-c)))/4.0;
```

```
}
```

```
.....
```

```
/* Geometry: Closest point on line segment a-b to point c -----*/
```

```
point closest_pt_lineseg(point a, point b, point c) {
```

```
    b -= a; c -= a; if (b == 0) return a;
```

```
    double d = (c/b).x;
```

```
    if (d < 0) d = 0; if (d > 1) d = 1;
```

```
    return a + d*b;
```

```
}
```

```
.....
```

```
/* Geometry: Rectangle in rectangle test -----*/
```

```
// Checks if rectangle of sides x,y fits inside one of sides X,Y
```

```
// Not tested with doubles but should work fine :)
```

```
// Code as written rejects rectangles that just touch.
```

```
bool rect_in_rect(int X, int Y, int x, int y) {
```

```
    if (Y > X) swap(Y, X);
```

```
    if (y > x) swap(y, x);
```

```
    double diagonal = sqrt(double(X)*X + double(Y)*Y);
```

```
    if (x < X && y < Y)
```

```
        return true;
```

```
    else if (y >= Y || x >= diagonal)
```

```
        return false;
```

```
    else {
```

```
        double w, theta, tMin = PI/4, tMax = PI/2;
```

```
        while (tMax - tMin > EPS) {
```

```
            theta = (tMax + tMin)/2.0;
```

```
            w = (Y-x*cos(theta))/sin(theta);
```

```
            if (w < 0 || x * sin(theta) + w * cos(theta) < Y)
```

```
                tMin = theta;
```

```

        else tMax = theta;
    }
    return (w > y);
}

.....

/* Geometry: Centroid of a simple polygon [O(N)] -----*/
// Points must be oriented (either CW or CCW), and non-convex is OK
point centroid(point p[], int n) {
    double sum = 0; point c;
    for (int i = n-1, j = 0; j < n; i = j++) {
        double area = (p[i].conj()*p[j]).y;
        sum += area; c += (p[i]+p[j])*area;
    }
    sum *= 3.0; c /= sum;
    return c;
}

.....

/* Geometry: Convex Hull -----*/
struct polar_cmp {
    point P0;
    polar_cmp(point p = 0) : P0(p) {}
    double turn(const point &p1, const point &p2) const {
        return ((p2-P0)*(p1-P0).conj()).y;
    }
    bool operator()(const point &p1, const point &p2) const {
        double d = turn(p1, p2);
        if (fabs(d) < EPS)
            return (p1-P0).mag2() < (p2-P0).mag2();
        else return d > 0;
    }
};
vector<point> convex_hull(vector<point> p) {
    sort(p.begin(), p.end());
    int n = unique(p.begin(), p.end()) - p.begin();
    sort(p.begin()+1, p.begin()+n, polar_cmp(p[0]));
    if (n <= 2) return vector<point>(p.begin(), p.begin()+n);
    vector<point> hull(p.begin(), p.begin()+2); int h = 2;
    for (int i = 2; i < n; ++i) {
        while ((h > 1) && (polar_cmp(hull[h-2], hull[h-1], p[i]) < EPS)) {
            hull.pop_back(); --h;
        }
        hull.push_back(p[i]); ++h;
    }
}

```

```

    }
    return hull;
}

.....

/* Geometry: Area of intersection of two circles -----*/
struct circle {
    point c; double r;
};
double CIArea(circle &a, circle &b) {
    double d = (b.c-a.c).mag();
    if (d <= (b.r - a.r)) return a.r*a.r*PI;
    if (d <= (a.r - b.r)) return b.r*b.r*PI;
    if (d >= a.r + b.r) return 0;
    double alpha = acos((a.r*a.r+d*d-b.r*b.r)/(2*a.r*d));
    double beta = acos((b.r*b.r+d*d-a.r*a.r)/(2*b.r*d));
    return a.r*a.r*(alpha-0.5*sin(2*alpha))+b.r*b.r*(beta-0.5*sin(2*beta));
}

.....

/* Geometry: Area of union of rectangles [O(N^2)] -----*/
// Rectangle sides are parallel to the x & y axes
// May be desirable to add a constructor to 'rect' to ensure that the
// coordinates are properly sorted
struct rect {
    double minx, miny, maxx, maxy;
};
struct edge {
    double x, miny, maxy;
    char m;
    bool operator<(const edge &e) const {
        return x < e.x;
    }
};
double area_unionrect(vector<rect> R) {
    int n = R.size();
    vector<double> ys(2*n);
    vector<edge> e(2*n);
    for (int i = 0; i < n; ++i) {
        e[2*i].miny = e[2*i+1].miny = ys[2*i] = r[i].miny;
        e[2*i].maxy = e[2*i+1].maxy = ys[2*i+1] = r[i].maxy;
        e[2*i].x = r[i].minx;
        e[2*i].m = 1;
        e[2*i+1].x = r[i].maxx;
        e[2*i+1].m = -1;
    }
}

```

```

    }
    sort(ys.begin(), ys.end());
    sort(e.begin(), e.end());
    double sum = 0, cur = 0;
    for (int i = 0; i < 2*n; ++i) {
        if (i) sum += (ys[i]-ys[i-1])*cur;
        int flag = 0; double sx = cur = 0;
        for (int j = 0; j < 2*n; ++j) {
            if (e[j].miny <= ys[i] && ys[i] < e[j].maxy) {
                if (!flag) sx = e[j].x;
                flag += e[j].m;
                if (!flag) curr += e[j].x-sx;
            }
        }
    }
    return sum;
}

.....

/* Geometry: Line segment a-b vs. c-d intersection (IF returned in p) -----*/
// returns 1 if intersect, 0 if not, -1 if coincident
int intersect_line(point a, point b, point c, point d, point &p) {
    double num1 = ((a-c)*(d-c).conj()).y, num2 = ((a-c)*(b-a).conj()).y;
    double denom = ((d-c)*(b-a).conj()).y;
    if (fabs(denom) > EPS) {
        double r = num1/denom, s = num2/denom;
        if ((0 <= r) && (r <= 1) && (0 <= s) && (s <= 1)) {
            p = a+r*(b-a);
            return 1;
        }
        return 0;
    }
    if (fabs(num1) > EPS) return 0;
    if (b < a) swap(a, b); if (d < c) swap(c, d);
    if (a.y == b.y) {
        if (b.x == c.x) { p = b; return 1; }
        else if (a.x == d.x) { p = a; return 1; }
        else if ((b.x < c.x) || (d.x < a.x)) return 0;
    }
    else {
        if (b.y == c.y) { p = b; return 1; }
        else if (a.y == d.y) { p = a; return 1; }
        else if ((b.y < c.y) || (d.y < a.y)) return 0;
    }
    return -1;
}

```

```

/* Geometry: Area of intersection of two general polygons [O(N^2)] -----*/
int ORDER = -1; // CCW ordering, 1 for CW
struct triangle {
    point p[3];
};
double cross(point a, point b, point c, point d) {
    d -= c; b -= a;
    return (d*b.conj()).y;
}
int leftRight(const point &a, const point &b, const point &p) {
    // -1: p left of a->b, +1: p right of a->b, 0: p on a->b
    double d = cross(a, b, a, p);
    if (d > EPS) return -1;
    if (d < -EPS) return 1;
    return 0;
}
bool isConcave(point &a, point &b, point &c) {
    // tests if b in a->b->c is concave/flat
    return ORDER*leftRight(a, b, c) <= 0;
}
bool isInsideTriangle(point &a, point &b, point &c, point &p) {
    int r1 = leftRight(a,b,p), r2 = leftRight(b,c,p), r3 = leftRight(c,a,p);
    return (ORDER*r1 >= 0) && (ORDER*r2 >= 0) && (ORDER*r3 >= 0);
}
vector<triangle> triangulate(vector<point> &orig) {
    // Accepts a vector of n ordered vertices, returns triangulation.
    // No triangles if n < 3.
    vector<triangle> T;
    if (orig.size() < 3) return T;
    list<point> P(orig.begin(), orig.end());
    list<point>::iterator a, b, c, q;
    for (a = b = P.begin(), c = ++b, ++c; c != P.end(); a = b, c = ++b, ++c)
        if (!isConcave(*a, *b, *c)) {
            q = P.begin(); if (q == a) { ++q; ++q; ++q; }
            while ((q != P.end()) && !isInsideTriangle(*a, *b, *c, *q)) {
                ++q; if (q == a) { ++q; ++q; ++q; }
            }
            if (q == P.end()) {
                triangle t; t.p[0] = *a; t.p[1] = *b; t.p[2] = *c; T.push_back(t);
                P.erase(b); b = a;
                if (b != P.begin()) --b;
            }
        }
    return T;
}
bool isectLineSegs(point &a, point &b, point &c, point &d, point &p) {

```

```

// Finds intersection p of segments a-b and c-d (returns 0 if none/inf)
double n1 = cross(c, d, c, a), n2 = -cross(a, b, a, c);
double dn = cross(a, b, c, d);
if (fabs(dn) > EPS) {
    double r = n1/dn, s = n2/dn;
    if ((0 <= r) && (r <= 1) && (0 <= s) && (s <= 1)) {
        p = a+r*(b-a);
        return true;
    }
}
return false;
}

struct radialLessThan {
    point P0;
    radialLessThan(point p = 0) : P0(p) {}
    bool operator()(const point &a, const point &b) const {
        return (ORDER == leftRight(P0, a, b));
    }
};

double isectAreaTriangles(triangle &a, triangle &b) {
    vector<point> P;
    point p; triangle T[2] = {a, b};
    for (int r = 1, t = 0; t < 2; r = t++)
        for (int i = 2, j = 0; j < 3; i = j++) {
            if (isInsideTriangle(T[r].p[0], T[r].p[1], T[r].p[2], T[t].p[i]))
                P.push_back(T[t].p[i]);
            for (int u = 2, v = 0; v < 3; u = v++)
                if (isectLineSegs(T[t].p[i], T[t].p[j], T[r].p[u], T[r].p[v], p))
                    P.push_back(p);
        }
    if (P.empty()) return 0;
    sort(P.begin(), P.end());
    vector<point> U; unique_copy(P.begin(), P.end(), back_inserter(U));
    if (U.size() >= 3) {
        sort(++U.begin(), U.end(), radialLessThan(U[0]));
        return areaPoly(U);
    }
    return 0;
}

double isectAreaGpoly(vector<point> &P, vector<point> &Q) {
    vector<triangle> S = triangulate(P), T = triangulate(Q);
    double area = 0;
    for (vector<triangle>::iterator s = S.begin(); s != S.end(); ++s)
        for (vector<triangle>::iterator t = T.begin(); t != T.end(); ++t)
            area += isectAreaTriangles(*s, *t);
    return -ORDER*area;
}

```

```

/* Geometry: Point in polygon -----*/
bool pt_in_poly(vector<point> &p, const point &a) {
    int n = p.size(); bool inside = false;
    for (int i = 0, j = n-1; i < n; j = i++) {
        if ((a-p[i]).mag()+ (a-p[j]).mag()- (p[i]-p[j]).mag() < EPS)
            return true; // Boundary case (pt on edge), you may want false here
        if (((p[i].y<=a.y) && (a.y<p[j].y)) || ((p[j].y<=a.y) && (a.y<p[i].y)))
            if (a.x-p[i].x < (p[j].x-p[i].x)*(a.y-p[i].y) / (p[j].y-p[i].y))
                inside = !inside;
    }
    return inside;
}

// Geometry: Polygon midpoints -> vertices (n odd) -----*/
vector<point> midpts2vert(vector<point> &midpts) {
    int n = midpts.size(); vector<point> poly(n);
    poly[0] = midpts[0];
    for (int i = 1; i < n-1; i += 2) {
        poly[0].x += midpts[i+1].x - midpts[i].x;
        poly[0].y += midpts[i+1].y - midpts[i].y;
    }
    for (int i = 1; i < n; i++) {
        poly[i].x = 2.0*midpts[i-1].x - poly[i-1].x;
        poly[i].y = 2.0*midpts[i-1].y - poly[i-1].y;
    }
    return poly;
}

/* Geometry: 3D Primitives -----*/
struct point3 {
    double x, y, z;
    point3(double X=0, double Y=0, double Z=0) : x(X), y(Y), z(Z) {}
    point3 operator+(point3 p) { return point3(x + p.x, y + p.y, z + p.z); }
    point3 operator*(double k) { return point3(k*x, k*y, k*z); }
    point3 operator-(point3 p) { return *this + (p*-1.0); }
    point3 operator/(double k) { return *this*(1.0/k); }
    double mag2() { return x*x + y*y + z*z; }
    double mag() { return sqrt(mag2()); }
    point3 norm() { return *this/this->mag(); }
};

double dot(point3 a, point3 b) {
    return a.x*b.x + a.y*b.y + a.z*b.z;
}

```

```

point3 cross(point3 a, point3 b) {
    return point3(a.y*b.z - b.y*a.z, b.x*a.z - a.x*b.z, a.x*b.y - b.x*a.y);
}

struct line {
    point3 a, b;
    line(point3 A=point3(), point3 B=point3()) : a(A), b(B) {}
    // Direction unit vector a -> b
    point3 dir() { return (b - a).norm(); }
};

point3 cpoint_iline(line u, point3 p) {
    // Closest point on an infinite line u to a given point p
    point3 ud = u.dir();
    return u.a - ud*dot(u.a - p, ud);
}

double dist_ilines(line u, line v) {
    // Shortest distance between two infinite lines u and v
    return dot(v.a - u.a, cross(u.dir(), v.dir()).norm());
}

point3 cpoint_ilines(line u, line v) {
    // Finds the closest point on infinite line u to infinite line v.
    // Note: if (uv*uv - uu*vv) is zero then the lines are parallel and such a
    // single closest point does not exist. Check for this if needed.
    point3 ud = u.dir(); point3 vd = v.dir();
    double uu = dot(ud, ud), vv = dot(vd, vd), uv = dot(ud, vd);
    double t = dot(u.a, ud) - dot(v.a, ud); t *= vv;
    t -= uv*(dot(u.a, vd) - dot(v.a, vd));
    t /= (uv*uv - uu*vv);
    return u.a + ud*t;
}

point3 cpoint_lineseg(line u, point3 p) {
    // Closest point on a line segment u to a given point p
    point3 ud = u.b - u.a; double s = dot(u.a - p, ud)/ud.mag2();
    if (s < -1.0) return u.b;
    if (s > 0.0) return u.a;
    return u.a - ud*s;
}

struct plane {
    point3 n, p;
    plane(point3 ni = point3(), point3 pi = point3()) : n(ni), p(pi) {}
    plane(point3 a, point3 b, point3 c) : n(cross(b-a, c-a).norm()), p(a) {}
    //Value of d for the equation ax + by + cz + d = 0
    double d() { return -dot(n, p); }
};

point3 cpoint_plane(plane u, point3 p) {
    //Closest point on a plane u to a given point p
    return p - u.n*(dot(u.n, p) + u.d());
}

```

```

point3 iline_isect_plane(plane u, line v) {
    //Point of intersection between an infinite line v and a plane u.
    //Note: if dot(u.n, vd) == 0 then the line and plane do not intersect at
    //a single point. Check for this case if it is needed.
    point3 vd = v.dir();
    return v.a - vd*((dot(u.n, v.a) + u.d())/dot(u.n, vd));
}

line isect_planes(plane u, plane v) {
    //Infinite line of intersection between two planes u and v.
    //Note: if dot(v.n, uvu) == 0 then the planes do not intersect at a line.
    //Check for this case if it is needed.
    point3 o = u.n*-u.d(), uv = cross(u.n, v.n);
    point3 uvu = cross(uv, u.n);
    point3 a = o - uvu*((dot(v.n, o) + v.d())/(dot(v.n, uvu)*uvu.mag2()));
    return line(a, a + uv);
}

// Geometry: Great Circle distance (lat[-90,90], long[-180,180])-----*/
double greatcircle(double lt1, double lo1, double lt2, double lo2, double r) {
    double a = PI*(lt1/180.0), b = PI*(lt2/180.0);
    double c = PI*((lo2-lo1)/180.0);
    return r*acos(sin(a)*sin(b) + cos(a)*cos(b)*cos(c));
}

// Geometry: Circle described by three points -----*/
bool circle(point p1, point p2, point p3, point &center, double &r) {
    double G = 2*((p2-p1).conj()*(p3-p2)).y;
    if (fabs(G) < EPS) return false;
    center = p1*(p3.mag2()-p2.mag2());
    center += p2*(p1.mag2()-p3.mag2());
    center += p3*(p2.mag2()-p1.mag2());
    center /= point(0, G); r = (p1-center).mag();
    return true;
}

// Arithmetic: Discrete Logarithm solver [O(sqrt(P))] -----*/
// Given prime P, B, and N, finds least L such that B^L == N (mod P)
typedef unsigned int UI;
typedef unsigned long long ULL;
map<UI,UI> M;
UI times(UI a, UI b, UI m) {

```



```
Result s;
if(z <= 0) {
    s.n = 3; theta = acos(r/sqrt(q*q*q))
    s.x[0] = sq*cos(theta/3.0) - a1/3.0
    s.x[1] = sq*cos((theta+2.0*PI)/3.0)
    s.x[2] = sq*cos((theta+4.0*PI)/3.0)
}
else {
    s.n = 1; s.x[0] = pow(sqrt(z)+fabs(
    s.x[0] += q/s.x[0]; s.x[0] *= (r <
    s.x[0] -= a1/3.0;
}
return s;
```



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```

t = 0;
for (i = 1; i < n; i++) {
    if (A[m[k]] > A[m[k+1]]) {
        m[k] = m[k+1]; start = m[k] = i;
    }
    if (A[i] < A[m[1]])
        i = m[1];
    if (A[i] < A[m[k]]) {
        m[k+1] = i; up = k;
    }
    if (low != up-1) {
        m[(low+up)/2] = i;
    }
}

```

```
int main() {
    // Prints entire (2^n) Gray code for an input number of bits.
    unsigned int i; int n;
    char s[512];
    scanf("%d", &n);
    for(i = 0; i < (1 << n); i++)
```



```

printf("%s\n", pbits(s, i^(i>>1), n));
}

.....

/* Generators: Catalan Numbers -----*/

long long int cat[33];
void getcat() {
    cat[0] = cat[1] = 1;
    for (int i = 2; i < 33; ++i)
        cat[i] = cat[i-1]*(4*i-6)/i;
}

.....

/* Generators: Binary Strings generator (cardinal order) -----*/

char bit[MAXN];
void recurse(int n, int curr, int left) {
    if(curr == n)
        Process(n);
    else {
        if(curr+left < n) {
            bit[curr] = 0; recurse(n, curr+1, left);
        }
        if(left) {
            bit[curr] = 1; recurse(n, curr+1, left-1);
        }
    }
}

void gen_bin_card(int n) {
    for(int i = 0; i <= n; i++) {
        printf("Cardinality %d:\n", i);
        recurse(n, 0, i);
    }
}

.....

/* Graph Theory: Maximum Bipartite Matching -----*/

/* How to use (sample at bottom):
For vertex i of set U:
match[i] = -1 means i is not matched
match[i] = x means the edge i->(x-|U|) is selected
*****
For simplicity, use addEdge(i,j,n) to add edges, where
0 <= i < |U| and 0 <= j < |V| and |U| = n.
If there is an edge from vertex i of U to vertex

```

```

j of V then: e[i][j+|U|] = e[j+|U|][i] = 1.
*****
- If |U| = n and |V| = m, then vertices are assumed
  to be from [0,n-1] in set U and [0,m-1] in set V.
- Remember that match[i]-n gives the edge from i, not just match[i]. */

const int MAXN 300 // How many vertices in U+V (in total)
char e[MAXN][MAXN]; // MODIFIED Adj. matrix (see note)
int match[MAXN], back[MAXN], q[MAXN], tail;
void addEdge(int x, int y, int n) {
    e[x][y+n] = e[y+n][x] = 1;
}

int find(int x, int n, int m) {
    int i, j, r;
    if(match[x] != -1) return 0;
    memset(back, -1, sizeof(back));
    for(q[i=0]=x, tail = 1; i < tail; i++)
        for(j = 0; j < n+m; j++) {
            if(!e[q[i]][j]) continue;
            if(match[j] != -1) {
                if(back[j] == -1) {
                    back[j] = q[i];
                    back[q[tail++]] = match[j] = j;
                }
            }
            else {
                match[match[q[i]] = j] = q[i];
                for(r = back[q[i]]; r != -1; r = back[back[r]])
                    match[match[r] = back[r]] = r;
                return 1;
            }
        }
    return 0;
}

void bipmatch(int n, int m) {
    memset(match, -1, sizeof(match));
    for(int i = 0; i < n+m; i++) if(find(i,n,m)) i = 0;
}

int main() {
    int n, m, esize, x, y;
    int i, count;
    // Read size of set U into n, size of set V into m
    while(scanf("%d %d", &n, &m) == 2) {
        memset(e, 0, sizeof(e)); // Clear edges
        scanf("%d", &esize); // get # of edges
        while(esize--) {
            scanf("%d %d", &x, &y); // add edges

```

```

    addEdge(x,y,n);                // Edges [0,n-1]->[0,m-1]
}
bipmatch(n, m);                    // Perform matching
for(count = i = 0; i < n; i++) {   // Print results
    if(match[i] != -1) {
        printf("%d->%d\n", i, match[i]-n);
        count++;
    }
}
printf("Matching size: %d\n", count);
}
return 0;
}

```



**/\* Graph Theory: Eulerian Graphs -----\*/**

/\* Before adding edges, call Init() to initialize all necessary data structures.

Use the provided function addEdge(x,y,c) which adds c number of edges between x and y.

isEulerian(int n, int \*start, int \*end) returns:

- 0 if the graph is not Eulerian
- 1 if the graph has a Euler cycle
- 2 if the graph a path, from start to end

with n being the number of nodes in the graph

```

const int MAXN 105    // Number of nodes
const int MAXM 505    // Maximum number of edges
#define min(a,b) (((a)<(b))?(a):(b))
#define max(a,b) (((a)>(b))?(a):(b))
#define DEC(a,b) g[a][b]--;g[b][a]--;deg[a]--;deg[b]--
int sets[MAXN], deg[MAXN];
int g[MAXN][MAXN];
int seq[MAXM], seqsize;
// Uncomment if you need copy of graph
// int g2[MAXN][MAXN], deg2[MAXN];
int getRoot(int x) {
    if (sets[x] < 0) return x;
    return sets[x] = getRoot(sets[x]);
}
void Union(int a, int b) {
    int ra = getRoot(a), rb = getRoot(b);
    if (ra != rb) {
        sets[ra] += sets[rb];
        sets[rb] = ra;
    }
}

```

```

void Init() {
    memset(sets, -1, sizeof(sets));
    memset(g, 0, sizeof(g));
    memset(deg, 0, sizeof(deg));
}
void addEdge(int x, int y, int count) {
    g[x][y] += count; deg[x] += count;
    g[y][x] += count; deg[y] += count;
    Union(x,y);
}
int isEulerian(int n, int *start, int *end) {
    int odd = 0, i, count = 0, x;
    for (i = 0; i < n; i++)
        if (deg[i]) {
            x = i; count++;
        }
    if (sets[getRoot(x)] != -count) return 0;
    for (i = 0; i < n; i++) {
        if (deg[i] & 1) {
            odd++;
            if(odd == 1) *start = i;
            else if(odd == 2) *end = i;
            else return 0;
        }
    }
    return odd ? 2 : 1;
}
void getPath(int n, int start, int end) {
    int temp[MAXM], tsize = 1, i, j;
    temp[0] = start;
    while(1) {
        j = temp[tsize-1];
        for (i = 0; i < n; i++) {
            if (i == end) continue;
            if (g[i][j]) {
                temp[tsize++] = i;
                DEC(i,j); break;
            }
        }
        if (i == n) {
            if (g[end][j]) {
                temp[tsize++] = end;
                DEC(j,end);
            }
            break;
        }
    }
}

```

```

for (i = 0; i < tsize; i++)
    if (!deg[temp[i]])
        seq[seqsize++] = temp[i];
    else getPath(n, temp[i], temp[i]);
}
void buildPath(int n, int start, int end) {
    seqsize = 0;
    // Uncomment if you need copy of graph
    // memcpy(g, g2, sizeof(g));
    // memcpy(deg, deg2, sizeof(deg));
    getPath(n, start, end);
}
int main() {
    int i, x, y, start, end, n, m;
    while (scanf("%d %d", &n, &m) == 2) {
        Init();
        for (i = 0; i < m; i++) {
            scanf("%d %d", &x, &y);
            addEdge(x, y, 1);
        }
        // Uncomment if you need copy of graph
        // memcpy(g2, g, sizeof(g2));
        // memcpy(deg2, deg, sizeof(deg2));
        int res = isEulerian(n, &start, &end);
    }
}

```

**/\* Graph Theory: Maximum Flow in a directed graph -----\*/**

/\* - Multiple edges from u to v may be added. They are converted into a single edge with a capacity equal to their sum  
 - Vertices are assumed to be numbered from 0..n-1  
 - The graph is supplied as the number of nodes (n), the zero-based indexes of the source (s) and the sink (t), and a vector of edges u->v with capacity c (M). \*/

```

const int MAXN 200
struct Edge {
    //Edge u->v with capacity c
    int u, v, c;
};
int F[MAXN][MAXN]; //Flow of the graph
int maxFlow(int n, int s, int t, vector<Edge> &M) {
    int u, v, c, oh, min, df, flow, H[n], E[n], T[n], C[n][n];
    vector<Edge>::iterator m;
    list<int> N; list<int>::iterator cur;
    vector<int> R[n]; vector<int>::iterator r;

```

```

for (u = 0; u < n; u++) {
    E[u] = H[u] = T[u] = 0;
    R[u].clear();
    for (v = 0; v < n; v++)
        C[u][v] = F[u][v] = 0;
}
for (m = M.begin(); m != M.end(); m++) {
    u = m->u; v = m->v; c = m->c;
    if (c && !C[u][v] && !C[v][u]) {
        R[u].push_back(v);
        R[v].push_back(u);
    }
    C[u][v] += c;
}
H[s] = n;
for (r = R[s].begin(); r != R[s].end(); r++) {
    v = *r;
    F[s][v] = C[s][v]; F[v][s] = -C[s][v];
    E[v] = C[s][v]; E[s] -= C[s][v];
}
N.clear();
for (u = 0; u < n; u++)
    if ((u != s) && (u != t))
        N.push_back(u);
for (cur = N.begin(); cur != N.end(); cur++) {
    u = *cur; oh = H[u];
    while (E[u] > 0)
        if (T[u] >= (int)R[u].size()) {
            min = 10000000;
            for (r = R[u].begin(); r != R[u].end(); r++) {
                v = *r;
                if ((C[u][v] - F[u][v] > 0) && (H[v] < min))
                    min = H[v];
            }
            H[u] = 1 + min;
            T[u] = 0;
        }
    else {
        v = R[u][T[u]];
        if ((C[u][v] - F[u][v] > 0) && (H[u] == H[v]+1)) {
            df = C[u][v] - F[u][v];
            if (df > E[u])
                df = E[u];
            F[u][v] += df; F[v][u] = -F[u][v];
            E[u] -= df; E[v] += df;
        }
    }
}
else

```

```

        T[u]++;
    }
    if (H[u] > oh)
        N.splice(N.begin(), N, cur);
    }
    flow = 0;
    for (r = R[s].begin(); r != R[s].end(); r++)
        flow += F[s][*r];
    return flow;
}

```



**/\* Graph Theory: Chinese Postman Problem -----\*/**

// The maximum # of vertices solvable is roughly 20

```
#define MAXN 20
```

```
#define DISCONNECT -1
```

```
int g[MAXN][MAXN]; // Adj matrix (keep lowest cost if multiedge)
```

```
int deg[MAXN]; // Degree count
```

```
int A[MAXN+1]; // Used by perfect matching generator
```

```
int sum; // Sum of costs
```

```
int odd, best;
```

```
void floyd(int n) {
```

```
    int i, j, k;
```

```
    for(k = 0; k < n; k++)
```

```
        for(i = 0; i < n; i++)
```

```
            for(j = 0; j < n; j++)
```

```
                if(g[i][k] != -1 && g[k][j] != -1) {
```

```
                    int temp = g[i][k] + g[k][j];
```

```
                    if(g[i][j] == -1 || g[i][j] > temp)
```

```
                        g[i][j] = temp;
```

```
                }
```

```
    for(i = 0; i < n; i++) g[i][i] = 0;
```

```
}
```

```
void checkSum() {
```

```
    int i, temp;
```

```
    for(i = temp = 0; i < odd/2; i++)
```

```
        temp += g[A[2*i]][A[2*i+1]];
```

```
    if(best == -1 || best > temp) best = temp;
```

```
}
```

```
void perfmach(int x) {
```

```
    int i, t;
```

```
    if(x == 2) checkSum();
```

```
    else {
```

```
        perfmach(x-2);
```

```
        for(i = x-3; i >= 0; i--) {
```

```
            t = A[i]; A[i] = A[x-2];
```

```
            A[x-2] = t; perfmach(x-2);
```

```
        }
```

```
    t = A[x-2];
```

```
    for(i = x-2; i >= 1; i--) A[i] = A[i-1];
```

```
    A[0] = t;
```

```
}
```

```
int postman(int n) {
```

```
    int i; floyd(n);
```

```
    for(odd = i = 0; i < n; i++)
```

```
        if(deg[i]%2) A[odd++] = i;
```

```
    if(!odd) return sum;
```

```
    best = -1;
```

```
    perfmach(odd);
```

```
    return sum+best;
```

```
}
```

```
int main() {
```

```
    int i, u, v, c, n, m;
```

```
    while(scanf("%d %d", &n, &m) == 2){
```

```
        // Clear graph and degree count
```

```
        memset(g, -1, sizeof(g));
```

```
        memset(deg, 0, sizeof(deg));
```

```
        for(sum = i = 0; i < m; i++) {
```

```
            scanf("%d %d %d", &u, &v, &c);
```

```
            u--; v--; deg[u]++; deg[v]++;
```

```
            if(g[u][v] == -1 || g[u][v] > c) g[u][v] = c;
```

```
            if(g[v][u] == -1 || g[v][u] > c) g[v][u] = c;
```

```
            sum += c;
```

```
        }
```

```
        printf("Best cost: %d\n", postman(n));
```

```
    }
```

```
}
```



**/\* Graph Theory: Strongly Connected Components -----\*/**

```
#define VI vector<int>
```

```
#define MAXN 1000
```

```
VI g[MAXN], curr;
```

```
vector< VI > scc;
```

```
int dfsnum[MAXN], low[MAXN], id;
```

```
char done[MAXN];
```

```
void visit(int x) {
```

```
    curr.push_back(x);
```

```
    dfsnum[x] = low[x] = id++;
```

```
    for(size_t i = 0; i < g[x].size(); i++)
```

```

if(dfsnum[g[x][i]] == -1){
    visit(g[x][i]);
    low[x] <?= low[g[x][i]];
}
else if(!done[g[x][i]])
    low[x] <?= dfsnum[g[x][i]];
if(low[x] == dfsnum[x]) {
    VI c; int y;
    do {
        done[y = curr[curr.size()-1]] = 1;
        c.push_back(y);
        curr.pop_back();
    }
    while(y != x);
    scc.push_back(c);
}
}

void strong_conn(int n) {
    memset(dfsnum, -1, n*sizeof(int));
    memset(done, 0, sizeof(done));
    scc.clear(); curr.clear();
    for(int i = id = 0; i < n; i++)
        if(dfsnum[i] == -1) visit(i);
}

```



#### /\* Graph Theory: Min Cost Max Flow (Edmonds-Karp & Dijkstra) -----\*/

/\* Takes a directed graph where each edge has a capacity ('cap') and a cost per unit of flow ('cost') and returns a maximum flow network of minimal cost ('fcost') from s to t. USE THIS CODE FOR (MODERATELY) DENSE GRAPHS; FOR VERY SPARSE GRAPHS, USE mcmf4 (next)

##### PARAMETERS:

- cap (global): adjacency matrix where cap[u][v] is the capacity of the edge u->v. cap[u][v] is 0 for non-existent edges.
- cost (global): a matrix where cost[u][v] is the cost per unit of flow along the edge u->v. If cap[u][v] == 0, cost[u][v] is ignored. ALL COSTS MUST BE NON-NEGATIVE!
- n: the number of vertices ([0, n-1] are considered as vertices).
- s: source vertex.
- t: sink.

##### RETURNS:

- the flow
- the total cost through 'fcost'
- fnet contains the flow network. Careful: both fnet[u][v] and fnet[v][u] could be positive. Take the difference.

##### COMPLEXITY:

```

- Worst case: O(n^2*flow <? n^3*fcost) */
// Watch for commas when typing this in!
#define NN 1024 // the maximum number of vertices + 1
int cap[NN][NN]; // adjacency matrix (fill this up)
int cost[NN][NN]; // cost per unit of flow matrix (fill this up)
int fnet[NN][NN], adj[NN][NN], deg[NN]; // flow network and adjacency list
int par[NN], d[NN]; // par[source] = source;
int pi[NN]; // Labelling function
#define CLR(a, x) memset(a, x, sizeof(a))
#define Inf (INT_MAX/2)
#define Pot(u,v) (d[u] + pi[u] - pi[v])
bool dijkstra(int n, int s, int t) {
    // Dijkstra's using non-negative edge weights (cost + potential)
    for (int i = 0; i < n; i++)
        d[i] = Inf, par[i] = -1;
    d[s] = 0; par[s] = -n - 1;
    while (1) {
        int u = -1, bestD = Inf;
        for (int i = 0; i < n; i++)
            if (par[i] < 0 && d[i] < bestD)
                bestD = d[u = i];
        if(bestD == Inf) break;
        par[u] = -par[u] - 1;
        for (int i = 0; i < deg[u]; i++) {
            int v = adj[u][i];
            if (par[v] >= 0) continue;
            if (fnet[v][u] && d[v] > Pot(u,v) - cost[v][u])
                d[v] = Pot(u,v) - cost[v][u], par[v] = -u-1;
            if (fnet[u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v])
                d[v] = Pot(u,v) + cost[u][v], par[v] = -u - 1;
        }
    }
    for (int i = 0; i < n; i++)
        if (pi[i] < Inf)
            pi[i] += d[i];
    return par[t] >= 0;
}

#undef Pot
int mcmf3(int n, int s, int t, int &fcost) {
    CLR(deg, 0); CLR(fnet, 0); CLR(pi, 0);
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            if (cap[i][j] || cap[j][i])
                adj[i][deg[i]++] = j;
    int flow = fcost = 0;
    while (dijkstra(n, s, t)) {
        int bot = INT_MAX;

```

```

for (int v = t, u = par[v]; v != s; u = par[v = u])
    bot <?= fnet[v][u] ? fnet[v][u] : (cap[u][v] - fnet[u][v]);
for (int v = t, u = par[v]; v != s; u = par[v = u])
    if (fnet[v][u]) {
        fnet[v][u] -= bot; fcost -= bot * cost[v][u];
    }
    else {
        fnet[u][v] += bot; fcost += bot * cost[u][v];
    }
    flow += bot;
}
return flow;
}

int main() {
    int numV; cin >> numV;
    memset(cap, 0, sizeof(cap));
    int m, a, b, c, cp, s, t;
    cin >> m >> s >> t;
    // fill up cap with existing capacities.
    // if the edge u->v has capacity 6, set cap[u][v] = 6.
    // for each cap[u][v] > 0, set cost[u][v] to the
    // cost per unit of flow along the edge i->v
    // Uncomment the commented statements if caps/costs are bidirectional
    for (int i=0; i<m; i++) {
        cin >> a >> b >> cp >> c;
        cost[a][b] = c; // cost[b][a] = c;
        cap[a][b] = cp; // cap[b][a] = cp;
    }

    int fcost, flow = mcmf3(numV, s, t, fcost);
    cout << "flow: " << flow << endl;
    cout << "cost: " << fcost << endl;
}

.....

/* Graph Theory: Min Cost Max Flow (Edmonds-Karp & fast heap Dijkstra) -----*/

// Same as above, but better for sparse graphs
#define NN 1024 // the maximum number of vertices + 1
int cap[NN][NN]; // adjacency matrix (fill this up)
int cost[NN][NN]; // cost per unit of flow matrix (fill this up)
int fnet[NN][NN], adj[NN][NN], deg[NN]; // flow network and adjacency list
int par[NN], d[NN], q[NN], inq[NN], qs; // Dijkstra's variables
int pi[NN]; // Labelling function
#define CLR(a, x) memset(a, x, sizeof(a))
#define Inf (INT_MAX/2)
#define BUBL { \
    t = q[i]; q[i] = q[j]; q[j] = t; \

```

```

    t = inq[q[i]]; inq[q[i]] = inq[q[j]]; inq[q[j]] = t; }
#define Pot(u,v) (d[u] + pi[u] - pi[v])
bool dijkstra(int n, int s, int t) {
    // Dijkstra's using non-negative edge weights (cost + potential)
    CLR(d, 0x3F); CLR(par, -1); CLR(inq, -1);
    d[s] = qs = 0;
    inq[q[qs++] = s] = 0;
    par[s] = n;
    while (qs) {
        int u = q[0]; inq[u] = -1;
        q[0] = q[--qs];
        if (qs) inq[q[0]] = 0;
        for (int i = 0, j = 2*i + 1, t; j < qs; i = j, j = 2*i + 1) {
            if (j + 1 < qs && d[q[j + 1]] < d[q[j]]) j++;
            if (d[q[j]] >= d[q[i]]) break;
            BUBL;
        }
        for (int k = 0, v = adj[u][k]; k < deg[u]; v = adj[u][++k]) {
            if (fnet[v][u] && d[v] > Pot(u,v) - cost[v][u])
                d[v] = Pot(u,v) - cost[v][par[v] = u];
            if (fnet[u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v])
                d[v] = Pot(u,v) + cost[par[v] = u][v];
            if (par[v] == u) {
                if (inq[v] < 0) { inq[q[qs] = v] = qs; qs++; }
                for (int i=inq[v], j=(i-1)/2, t; d[q[i]]<d[q[j]]; i=j, j=(i-1)/2)
                    BUBL;
            }
        }
    }

    for (int i = 0; i < n; i++)
        if (pi[i] < Inf)
            pi[i] += d[i];
    return par[t] >= 0;
}

#undef Pot
int mcmf4(int n, int s, int t, int &fcost) {
    CLR(deg, 0); CLR(fnet, 0); CLR(pi, 0);
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            if (cap[i][j] || cap[j][i])
                adj[i][deg[i]++] = j;
    int flow = fcost = 0;
    while (dijkstra(n,s,t)) {
        int bot = INT_MAX;
        for (int v = t, u = par[v]; v != s; u = par[v = u])
            bot <?= fnet[v][u] ? fnet[v][u] : (cap[u][v] - fnet[u][v]);
        for (int v = t, u = par[v]; v != s; u = par[v = u])

```

```

    if (fnet[v][u]) {
        fnet[v][u] -= bot; fcost -= bot * cost[v][u];
    }
    else {
        fnet[u][v] += bot; fcost += bot * cost[u][v];
    }
    flow += bot;
}
return flow;
}

```



**/\* Graph Theory: Articulation Points & Bridges (adj list) [O(V+E)] -----\*/**

/\* - array entry art[v] is true iff vertex v is an articulation point  
 - array entries bridge[i][0] and bridge[i][1] are the endpoints of a bridge  
 in the graph. If bridge (u,v) is represented in the array, (v,u) is not.  
 - 'bridges' is the number of bridges in the graph  
 - index vertices from 0 to n-1 \*/

```

#define MAX_N 200
#define min(a,b) (((a)<(b))?(a):(b))
// Why typedef struct, when struct will do? I didn't change this one.
typedef struct {
    int deg;
    int adj[MAX_N];
} Node;
Node alist[MAX_N];
bool art[MAX_N], seen[MAX_N];
int df_num[MAX_N], low[MAX_N], father[MAX_N], cnt;
int bridge[MAX_N*MAX_N][2], bridges;
void add_edge(int v1, int v2) {
    alist[v1].adj[alist[v1].deg++] = v2;
    alist[v2].adj[alist[v2].deg++] = v1;
}
void add_bridge(int v1, int v2) {
    bridge[bridges][0] = v1;
    bridge[bridges][1] = v2;
    ++bridges;
}
void clear() {
    for (int i = 0; i < MAX_N; ++i)
        alist[i].deg = 0;
}
void search(int v, bool root) {
    int w, child = 0;
    seen[v] = true;
    low[v] = df_num[v] = cnt++;

```

```

    for (int i = 0; i < alist[v].deg; ++i) {
        w = alist[v].adj[i];
        if (df_num[w] == -1) {
            father[w] = v; ++child;
            search(w, false);
            if (low[w] > df_num[v]) add_bridge(v, w);
            if (low[w] >= df_num[v] && !root)
                art[v] = true;
            low[v] = min(low[v], low[w]);
        }
        else if (w != father[v]) {
            low[v] = min(low[v], df_num[w]);
        }
    }
    if (root && child > 1) art[v] = true;
}
void articulate(int n) {
    int child = 0;
    for (int i = 0; i < n; ++i) {
        art[i] = false;
        df_num[i] = father[i] = -1;
    }
    cnt = bridges = 0;
    memset(seen, false, sizeof(seen));
    for (int i = 0; i < n; ++i)
        if (!seen[i])
            search(i, true);
}
int main() {
    int n, m, v1, v2, c = 0;
    while (true) {
        scanf("%d %d", &n, &m);
        if (!n && !m) break;
        clear();
        for (int i = 0; i < m; ++i) {
            scanf("%d %d", &v1, &v2);
            add_edge(v1 - 1, v2 - 1);
        }
        articulate(n);
        printf("Articulation Points:");
        for (int i = 0; i < n; ++i)
            if (art[i]) printf(" %d", i + 1);
        printf("\n");
        printf("Bridges:");
        for (int i = 0; i < bridges; ++i)
            printf(" (%d,%d)", bridge[i][0] + 1, bridge[i][1] + 1);
        printf("\n\n");
    }
}

```

```

}
}

```



**/\* Graph Theory: Maximum Weighted Bipartite Matching [O(n<sup>3</sup>)] -----\*/**

```

/* Given N workers and N jobs to complete, where each worker has a
certain compatibility (weight) to each job, find an assignment
(perfect matching) of workers to jobs which maximizes the
compatibility (weight).
- W is a 2 dimensional array where W[i][j] is the weight of
worker i doing job j.  Weights must be non-negative.  If
there is no weight assigned to a particular worker and job
pair, set it to zero.  If there is a different number of
workers than jobs, create dummy workers or jobs accordingly
with zero weight edges.
- M is a 1 dimensional array populated by the algorithm where
M[i] is the index of the job matched to worker i.
- This algorithm could be used on non-negative floating point
weights as well.
#define MAX_N 100 // Max number of workers/jobs
int W[MAX_N][MAX_N], U[MAX_N], V[MAX_N], Y[MAX_N]; // weight vars
int M[MAX_N], N[MAX_N], P[MAX_N], Q[MAX_N], R[MAX_N], S[MAX_N], T[MAX_N];
int Assign(int n) {
// Returns max weight, corresponding matching inside global M
int w, y; // weight vars
int i, j, m, p, q, s, t, v;
for (i = 0; i < n; i++) {
M[i] = N[i] = -1; U[i] = V[i] = 0;
for (j = 0; j < n; j++)
if (W[i][j] > U[i])
U[i] = W[i][j];
}
for (m = 0; m < n; m++) {
for (p = i = 0; i < n; i++) {
T[i] = 0; Y[i] = -1;
if (M[i] == -1) {
S[i] = 1; P[p++] = i;
}
else S[i] = 0;
}
while (1) {
for (q = s = 0; s < p; s++) {
i = P[s];
for (j = 0; j < n; j++)
if (!T[j]) {
y = U[i] + V[j] - W[i][j];

```

```

if (y == 0) {
R[j] = i;
if (N[j] == -1)
goto end_phase; // I hate goto's!
T[j] = 1; Q[q++] = j;
}
else if ((Y[j] == -1) || (y < Y[j])) {
Y[j] = y; R[j] = i;
}
}
}
if (q == 0) {
y = -1;
for (j = 0; j < n; j++)
if (!T[j] && ((y == -1) || (Y[j] < y)))
y = Y[j];
for (j = 0; j < n; j++) {
if (T[j])
V[j] += y;
if (S[j])
U[j] -= y;
}
for (j = 0; j < n; j++)
if (!T[j]) {
Y[j] -= y;
if (Y[j] == 0) {
if (N[j] == -1)
goto end_phase; // again!
T[j] = 1; Q[q++] = j;
}
}
}
for (p = t = 0; t < q; t++) {
i = N[Q[t]];
S[i] = 1; P[p++] = i;
}
}
end_phase:
i = R[j]; v = M[i];
M[i] = j; N[j] = i;
while (v != -1) {
j = v; i = R[j];
v = M[i];
M[i] = j; N[j] = i;
}
}
for (i = w = 0; i < n; i++)

```



```

        w += W[i][M[i]];
    return w;
}

int main() {
    int w; // weight var
    int n, i, j;
    while ((scanf("%d", &n) == 1) && (n != 0)) {
        for (i = 0; i < n; i++)
            for (j = 0; j < n; j++)
                scanf("%d", &W[i][j]);
        w = Assign(n);
        printf("Optimum weight: %d\n", w);
        printf("Matchings:\n");
        for (i = 0; i < n; i++)
            printf("%d matched to %d\n", i, M[i]);
    }
}

```



**/\* Linear Programming: Simplex Method-----\*/**

```

/* m - number of (less than) inequalities
   n - number of variables
   C - (m+1) by (n+1) array of coefficients:
       row 0      - objective function coefficients
       row 1:m    - less-than inequalities
       column 0:n-1 - inequality coefficients
       column n   - inequality constants (0 for objective function)
   X[n] - result variables
   return value - maximum value of objective function
                 (-inf for infeasible, inf for unbounded) */

#define MAXM 400 // leave one extra
#define MAXN 400 // leave one extra
#define EPS 1e-9
#define INF 1.0/0.0
double A[MAXM][MAXN];
int basis[MAXM], out[MAXN];
void pivot(int m, int n, int a, int b) {
    int i, j;
    for (i=0; i<=m; i++)
        if (i!=a)
            for (j=0; j<=n; j++)
                if (j!=b)
                    A[i][j] -= A[a][j] * A[i][b] / A[a][b];
    for (j=0; j<=n; j++)
        if (j!=b) A[a][j] /= A[a][b];
    for (i=0; i<=m; i++)

```

```

        if (i!=a) A[i][b] = -A[i][b]/A[a][b];
    A[a][b] = 1/A[a][b];
    i = basis[a];
    basis[a] = out[b];
    out[b] = i;
}

double simplex(int m, int n, double C[][MAXN], double X[]) {
    int i, j, ii, jj; // i, ii are row indexes; j, jj are column indexes
    for (i=1; i<=m; i++)
        for (j=0; j<=n; j++)
            A[i][j] = C[i][j];
    for (j=0; j<=n; j++)
        A[0][j] = -C[0][j];
    for (i=0; i<=m; i++)
        basis[i] = -i;
    for (j=0; j<=n; j++)
        out[j] = j;
    for (;;) {
        for (i=ii=1; i<=m; i++)
            if (A[i][n]<A[ii][n] || (A[i][n]==A[ii][n] && basis[i]<basis[ii]))
                ii=i;
        if (A[ii][n] >= -EPS) break;
        for (j=jj=0; j<=n; j++)
            if (A[ii][j]<A[jj][j]-EPS || (A[ii][j]<A[jj][j]-EPS && out[i]<out[j]))
                jj=j;
        if (A[ii][jj] >= -EPS) return -INF;
        pivot(m, n, ii, jj);
    }
    for (;;) {
        for (j=jj=0; j<=n; j++)
            if (A[0][j]<A[jj][j] || (A[0][j]==A[jj][j] && out[j]<out[jj]))
                jj=j;
        if (A[0][jj] > -EPS) break;
        for (i=1, ii=0; i<=m; i++)
            if (A[i][jj]>EPS &&
                (!ii || A[i][n]/A[i][jj]<A[ii][n]/A[ii][jj]-EPS ||
                 (A[i][n]/A[i][jj]<A[ii][n]/A[ii][jj]+EPS
                  && basis[i]<basis[ii])))
                ii=i;
        if (A[ii][jj] <= EPS) return INF;
        pivot(m, n, ii, jj);
    }
    for (j=0; j<=n; j++)
        X[j] = 0;
    for (i=1; i<=m; i++)
        if (basis[i] >= 0)
            X[basis[i]] = A[i][n];

```

A horizontal row of 25 solid green dots.

```
/* Description: This document is a reference for the use of the
               java for regular IO purposes.  It covers stdin and
               stdout as well as file IO.  It also shows how to use
               StringTokenizer for parsing. */
```

```
while (st.hasMoreTokens()) {
    out.print("Token: ");
    out.println(st.nextToken());
}

// To split on something else,
// st = new StringTokenizer(line, delim);
// Or use this to change in the future
// line = st.nextToken(delim);
}

// You must flush for files!
out.flush();
}
```

A decorative horizontal line composed of 20 small green circles spaced evenly across the width of the page.

```
/* Description: This document is a reference for the use of the
               BigInteger class in Java. It contains sample code
               that computes GCDs of pairs of integers.
```

.....

## Creating BigIntegers

## 1. From Strings

## 2. From byte arrays

3. From a long integer

Math operations:

-----

```
A ^ b = C      --> C = A.pow(B);
abs(A) = C      --> C = A.abs();
-(A) = C        --> C = A.negate()
```

## Bit Operations

-----

```
~A = C      (NOT)      --> C = A.not();
A & B = C    (AND)     --> C = A.and(B);
```



```

int tnum, base1, base2;
StringTokenizer st;
// key is the base system that you may change as needed
String key = new
String("0123456789ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz");
// Standard IO
BufferedReader in = new BufferedReader(new InputStreamReader(System.in));
PrintStream out = System.out;
// File IO
// BufferedReader in = new BufferedReader(new FileReader("probl.dat"));
// PrintWriter out = new BufferedWriter(new FileWriter("probl.out"));
line = in.readLine(); // Get number of test cases
st = new StringTokenizer(line);
tnum = Integer.parseInt(st.nextToken());
for (int t = 0; t < tnum; t++) {
    line = in.readLine();
    st = new StringTokenizer(line);
    base1 = Integer.parseInt(st.nextToken());
    base2 = Integer.parseInt(st.nextToken());
    n = st.nextToken();
    String result = convert_base(base1, base2, n, key2);
    out.println(result);
}
catch (Exception e) {
    System.err.println(e.toString());
}
}

/* Number Theory: Primality Testing -----*/

bool isPrime(int x) {
    if(x == 1) return ONEPRIME;
    if(x == 2) return true;
    if(!(x & 1)) return false;
    for(int i = 3; i*i <= x; i += 2)
        if (x % i) return false;
    return true;
}

/* Number Theory: Number of Divisors [O(sqrt(N))] -----*/

int num_divisors(int n) {
    int i, count, res = 1;

```

```

for(i = 2; i*i <= n; i++) {
    count = 0;
    while(!(n%i)) {
        n /= i; count++;
    }
    if(count) res *= (count+1);
}
if(n > 1) res *= 2;
return res;
}

/* Number Theory: Prime Factorization -----*/

int primes[MAXP]; int psize;
void getPrimes() {
    int i, j, isprime;
    psize = 0; primes[psize++] = 2;
    for (i = 3; i <= MAXN; i += 2) {
        for (isprime = j = 1; j < psize; j++) {
            if (i % primes[j] == 0) {
                isprime = 0;
                break;
            }
            if (1.0*primes[j]*primes[j] > i) break;
        }
        if(isprime) primes[psize++] = i;
    }
}

struct Factors {
    int size;
    int f[32];
};

Factors getPFactor(int n) {
    Factors x;
    int i;
    x.size = 0;
    for (i = 0; i < psize; i++) {
        while (n % primes[i] == 0) {
            x.f[x.size++] = primes[i];
            n /= primes[i];
        }
        if(1.0*primes[i]*primes[i] > n) break;
    }
    if(n > 1)
        x.f[x.size++] = n;
    return x;
}

```

```

}

.....

/* Number Theory: Primality testing with a sieve -----*/

// Consider using typedefs and functions instead of defines...
#define TEST(f,x)  ((f+(x)/16)&(1<<(((x)%16L)/2)))
#define SET(f,x)   ((f+(x)/16)|=1<<(((x)%16L)/2))
#define ONEPRIME 0 // whether or not 1 is considered to be prime
#define UL unsigned long
#define UC unsigned char
UC *primes = NULL;
UL getPrimes(UL maxn) {
    UL x, y, psize=1;
    primes = calloc(((maxn)>>4)+1L, sizeof(UC));
    for (x = 3; x*x <= maxn; x+=2)
        if (!TEST(primes, x))
            for (y = x*x; y <= maxn; y += x<<1) SET (primes, y);
    // Comment out if you don't need # of primes <= maxn
    for(x = 3; x <= maxn; x+=2)
        if(!TEST(primes, x)) psize++;
    return psize;
}

int isPrime(UL x) {
    // Returns whether or not a given POSITIVE number is prime
    if(x == 1) return ONEPRIME;
    if(x == 2) return 1;
    if(x % 2 == 0) return 0;
    return (!TEST(primes, x));
}

.....

/* Number Theory: Sum of divisors [O(sqrt(N))] -----*/

typedef long long int LL;
LL sum_divisors(LL n) {
    int i, count; LL res = 1;
    for (i = 2; i*i <= n; i++) {
        count = 0;
        while (n % i == 0) {
            n /= i; count++;
        }
        if (count) res *= (pow(i, count+1)-1)/(i-1);
    }
    if(n > 1) res *= (pow(n, 2)-1)/(n-1);
    return res;
}

```

```

/* Number Theory: Chinese Remainder Theorem -----*/

/* Given n relatively prime modular in m[0], ..., m[n-1], and right-hand
sides a[0], ..., a[n-1], the routine solves for the unique solution
in the range 0 <= x < m[0]*m[1]*...*m[n-1] such that x = a[i] mod m[i]
for all 0 <= i < n. The algorithm used is Garner's algorithm, which
is not the same as the one usually used in number theory textbooks.

It is assumed that m[i] are positive and pairwise relatively prime.
a[i] can be any integer.

If the system of equations is
    x = a[0] mod m[0]
    x = a[1] mod m[1]
    ...
then a[i] should be reduced mod m[i] first.
Also, if 0 <= a[i] < m[i] for all i, then the answer will fall
in the range 0 <= x < m[0]*m[1]*...*m[n-1]. */

int gcd(int a, int b, int *s, int *t) {
    int r, r1, r2, a1, a2, b1, b2, q;
    a1 = b2 = 1;
    a2 = b1 = 0;
    while (b) {
        q = a / b; r = a % b;
        r1 = a1 - q*b1;
        r2 = a2 - q*b2;
        a = b; a1 = b1; a2 = b2;
        b = r; b1 = r1; b2 = r2;
    }
    *s = a1; *t = a2;
    return a;
}

int cra(int n, int *m, int *a) {
    int x, i, k, prod, temp;
    int *gamma, *v;
    gamma = malloc(n*sizeof(int));
    v      = malloc(n*sizeof(int));
    for (k = 1; k < n; k++) {
        prod = m[0] % m[k];
        for (i = 1; i < k; i++) {
            prod = (prod * m[i]) % m[k];
        }
        gcd(prod, m[k], gamma+k, &temp);
        gamma[k] %= m[k];
        if (gamma[k] < 0)
            gamma[k] += m[k];
    }
}

```

```

v[0] = a[0];
for (k = 1; k < n; k++) {
    temp = v[k-1];
    for (i = k-2; i >= 0; i--) {
        temp = (temp * m[i] + v[i]) % m[k];
        if (temp < 0)
            temp += m[k];
    }
    v[k] = ((a[k] - temp) * gamma[k]) % m[k];
    if (v[k] < 0)
        v[k] += m[k];
}
x = v[n-1];
for (k = n-2; k >= 0; k--)
    x = x * m[k] + v[k];
free(gamma); free(v);
return x;
}

int main(void) {
    int n, *m, *a, i, x;
    while (scanf("%d", &n) == 1 && n > 0) {
        m = malloc(n*sizeof(int));
        a = malloc(n*sizeof(int));
        printf("Enter moduli:\n");
        for (i = 0; i < n; i++)
            scanf("%d", m+i);
        printf("Enter right-hand side:\n");
        for (i = 0; i < n; i++)
            scanf("%d", a+i);
        x = cra(n, m, a);
        printf("x = %d\n", x);
        free(m); free(a);
    }
}

```

**/\* Number Theory: Extended Euclidean Algorithm -----\*/**

```

// Assumes non-negative input. Returns d s.t. d = a*x + b*y
// x,y passed in by reference, #include <algorithm> for swap function
int gcd(int a, int b, int &x, int &y) {
    x = 1; y = 0; int nx = 0, ny = 1;
    while (b) {
        int q = a/b;
        x -= q*nx; swap(x, nx);
        y -= q*ny; swap(y, ny);
        a -= q*b; swap(a, b);
    }
}

```

```

}
return a;
}

```



**/\* Number Theory: Generalized Chinese Remaindering -----\*/**

```

/* Given [a_0, ..., a_(n-1)] and [m_0, ..., m_(n-1)]
Computes 0 <= x < lcm(m_0, ..., m_(n-1)) such that
x == a_0 mod m_0, ..., x == a_(n-1) mod m_(n-1), if
such an x exists. True is returned iff such an
x exists. If x does not exist then the value at
the address of x will not be affected.
Complexity: O(n log(MAX(m_0, ..., m_(n-1)))) */

```

```

typedef long long int LLI;
LLI safe_mod(LLI a, LLI m) {
    if (a < 0) return (a + m + m * (-a/m)) % m;
    else return a % m;
}

LLI abs(LLI a) {
    return a < 0 ? -a : a;
}

LLI gcdex(LLI a, LLI b, LLI *ss, LLI *tt) {
    LLI q, r[150], s[150], t[150];
    int num = 2;
    r[0] = a; r[1] = b;
    s[0] = t[1] = 1;
    s[1] = t[0] = 0;
    while (r[num - 1]) {
        q = r[num - 2] / r[num - 1];
        r[num] = r[num - 2] % r[num - 1];
        s[num] = s[num - 2] - q * s[num - 1];
        t[num] = t[num - 2] - q * t[num - 1];
        ++num;
    }
    *ss = s[num - 2]; *tt = t[num - 2];
    return r[num - 2];
}

bool gen_chrem(LLI *a, LLI *m, int n, LLI *x) {
    LLI g, s, t, a_tmp, m_tmp;
    a_tmp = safe_mod(a[0], m[0]);
    m_tmp = m[0];
    for (int i = 1; i < n; ++i) {
        g = gcdex(m_tmp, m[i], &s, &t);
        if (abs(a_tmp - a[i]) % g) return false;
        a_tmp = safe_mod(a_tmp + (a[i] - a_tmp) / g * s * m_tmp, m_tmp/g*m[i]);
        m_tmp = m[i];
    }
}

```

```

    }
    x = a_tmp;
    return true;
}

int main() {
    int n; LLI a[20], m[20], x;
    while (true) {
        scanf("%lld", &n);
        if (!n) break;
        for (int i = 0; i < n; ++i)
            scanf("%lld %lld", &a[i], &m[i]);
        if (!gen_chrem(a, m, n, &x))
            printf("No solution.\n\n");
        else
            printf("X = %lld\n\n", x);
    }
}

```



**/\* Number Theory: Rational Reconstruction [O(log m)] -----\*/**

*/\* Description: Given integers m, g and k, computes integers 'num' and 'den' (if they exist) such that  $\text{num} \equiv g \cdot \text{den} \pmod{m}$  where  $|\text{num}| < k$  and  $0 < \text{den} < g/k$ . True is returned iff den is invertible mod m. This algorithm is useful if computations on rational numbers is to be used when the input and output numbers have small numerators and denominators but intermediate results can have very large numerators and denominators. To use in this fashion, reduce the input rationals modulo some number m (probably a prime), perform the operations modulo m and then use rational reconstruction to recover the results. m and k must be selected such that  $|\text{num}|, \text{den} < k$  and  $2 \cdot k \cdot k < m$  for all input and output rational numbers. \*/*

```

typedef long long int LLI;
int gcd_table(LLI a, LLI b, LLI *r, LLI *q, LLI *s, LLI *t) {
    int n = 2;
    assert(0 <= a && 0 < b);
    r[0] = a; r[1] = b;
    s[0] = t[1] = 1;
    s[1] = t[0] = 0;
    while (r[n - 1]) {
        r[n] = r[n - 2] % r[n - 1];
        q[n - 1] = r[n - 2] / r[n - 1];
        s[n] = s[n - 2] - s[n - 1] * q[n - 1];
        t[n] = t[n - 2] - t[n - 1] * q[n - 1];
        ++n;
    }
    return n;
}

```

```

LLI gcd(LLI a, LLI b) {
    if (a < 0) return gcd(-a, b);
    if (b < 0) return gcd(a, -b);
    if (!b) return a;
    return gcd(b, a % b);
}

bool rat_recon(LLI m, LLI g, LLI k, LLI *num, LLI *den) {
    int n, j;
    LLI r[200], q[200], s[200], t[200], quo, tj, rj;
    assert(0 <= g && g < m && 1 <= k && k <= m);
    n = gcd_table(m, g, r, q, s, t);
    q[0] = q[n - 1] = 0;
    for (j = 0; j < n && r[j] >= k; ++j);
    if (t[j] > 0) {
        *num = r[j]; *den = t[j];
    }
    else {
        *num = -r[j]; *den = -t[j];
    }
    if (gcd(r[j], t[j]) == 1) return true;
    else {
        quo = (j == n - 1 ? 0 : (k - r[j - 1]) / r[j] + 1);
        rj = r[j - 1] - quo * r[j];
        tj = t[j - 1] - quo * t[j];
        if (gcd(rj, tj) != 1 || (tj > 0 ? tj : -tj) * k > m)
            return false;
        if (tj > 0) {
            *num = rj; *den = tj;
        }
        else {
            *num = -rj; *den = -tj;
        }
        return true;
    }
}

int main() {
    LLI m, g, k, r, t; char c;
    scanf("%lld %lld %lld", &m, &g, &k);
    c = (rat_recon(m, g, k, &r, &t) ? 'y' : 'n');
    printf("%c %lld / %lld\n", c, r, t);
}

```



**/\* Search: Golden section search -----\*/**

*/\* Given an function f(x) with a single local minimum, a lower and upper bound on x, and a tolerance for convergence, this function finds the*

```

    minimizing value of x. f(x) should evaluate globally.
#define GOLD 0.381966
#define move(a,b,c)    x[a]=x[b];x[b]=x[c];fx[a]=fx[b];fx[b]=fx[c]
double f(double x) { return x*x; } // Just an example
double golden(double xlow, double xhigh, double tol) {
    double x[4], fx[4], L;
    int iter = 0, left = 0, mini, i;
    fx[0] = f(x[0]=xlow);
    fx[3] = f(x[3]=xhigh);
    while (1) {
        L = x[3]-x[0];
        if (!iter || left) {
            x[1] = x[0]+GOLD*L;
            fx[1] = f(x[1]);
        }
        if (!iter || !left) {
            x[2] = x[3]-GOLD*L;
            fx[2] = f(x[2]);
        }
        for (mini = 0, i = 1; i < 4; i++)
            if (fx[i] < fx[mini]) mini = i;
        if (L < tol) break;
        if (mini < 2) {
            left = 1;
            move(3,2,1);
        }
        else {
            left = 0;
            move(0,1,2);
        }
        iter++;
    }
    return x[mini];
}

/* Search: KMP String Matching -----*/
/* Given strings T and P, computes the indices of T where P occurs
   as a substring and stores in 'shift'. The return integer is the
   number of indices stored in 'shift' */
#define MAX_LEN 1000
int pi[MAX_LEN];
void compute_prefix(char *P, int m, int *pi) {
    int k = pi[0] = -1;
    for (int q = 1; q < m; ++q) {
        while (k >= 0 && P[k + 1] != P[q]) k = pi[k];

```

```

        if (P[k + 1] == P[q]) ++k;
        pi[q] = k;
    }
}

int kmp_match(char *T, char *P, int *shift) {
    int n, m, q = -1, shifts = 0;
    n = strlen(T); m = strlen(P);
    compute_prefix(P, m, pi);
    for (int i = 0; i < n; ++i) {
        while (q > -1 && P[q + 1] != T[i]) q = pi[q];
        if (P[q + 1] == T[i]) ++q;
        if (q == m - 1) {
            shift[shifts++] = i - m + 1;
            q = pi[q];
        }
    }
    return shifts;
}

int main() {
    char T[MAX_LEN + 1], P[MAX_LEN + 1];
    int shift[MAX_LEN], shifts;
    while (scanf("%s %s", T, P) != -1) {
        shifts = kmp_match(T, P, shift);
        if (shifts) {
            printf("Pattern occurs with shifts:");
            for (int i = 0; i < shifts; ++i) printf(" %d", shift[i]);
            printf("\n\n");
        }
        else printf("No matches.\n\n");
    }
}

.....

/* Search: Suffix array [O(N log N)] -----*/
/* Notes:
   The build_sarray routine takes in a string S of n
   characters (null-terminated), and constructs two
   arrays sarray and lcp. The properties are:
   - If p = sarray[i], then the suffix of str starting at
     p (i.e. S[p..n-1] is the i-th suffix when all the
     suffixes are sorted in lexicographical order
   - NOTE: the empty suffix is not included in this list,
     so sarray[0] != n.
   - lcp[i] contains the length of the longest common
     prefix of the suffixes pointed to by sarray[i-1]
     and sarray[i]. lcp[0] is defined to be 0.
   - To see whether a pattern P occurs in str, you can

```



```

        look for it as the prefix of a suffix. This can be
        done with a binary search in  $O(|P| \log n)$  time.
    */
// You probably need to include <climits> here.
#define MAXN 100000
int bucket[CHAR_MAX-CHAR_MIN+1];
int prm[MAXN], count[MAXN];
char bh[MAXN+1];
void build_sarray(char *str, int* sarray, int *lcp) {
    int n, a, c, d, e, f, h, i, j, x;
    n = strlen(str);
    memset(bucket, -1, sizeof(bucket));
    for (i = 0; i < n; i++) {
        j = str[i] - CHAR_MIN;
        prm[i] = bucket[j];
        bucket[j] = i;
    }
    for (a = c = 0; a <= CHAR_MAX - CHAR_MIN; a++) {
        for (i = bucket[a]; i != -1; i = j) {
            j = prm[i]; prm[i] = c;
            bh[c++] = (i == bucket[a]);
        }
    }
    bh[n] = 1;
    for (i = 0; i < n; i++)
        sarray[prm[i]] = i;
    x = 0;
    for (h = 1; h < n; h *= 2) {
        for (i = 0; i < n; i++) {
            if (bh[i] & 1) {
                x = i; count[x] = 0;
            }
            prm[sarray[i]] = x;
        }
        d = n - h; e = prm[d];
        prm[d] = e + count[e];
        count[e]++;
        bh[prm[d]] |= 2;
        i = 0;
        while (i < n) {
            for (j = i; (j == i || !(bh[j] & 1)) && j < n; j++) {
                d = sarray[j] - h;
                if (d >= 0) {
                    e = prm[d]; prm[d] = e + count[e];
                    count[e]++; bh[prm[d]] |= 2;
                }
            }
            for (j = i; (j == i || !(bh[j] & 1)) && j < n; j++) {

```

```

                d = sarray[j] - h;
                if (d >= 0 && (bh[prm[d]] & 2)) {
                    for (e = prm[d]+1; bh[e] == 2; e++);
                    for (f = prm[d]+1; f < e; f++)
                        bh[f] &= 1;
                }
            }
            i = j;
        }
        for (i = 0; i < n; i++) {
            sarray[prm[i]] = i;
            if (bh[i] == 2)
                bh[i] = 3;
        }
    }
    h = 0;
    for (i = 0; i < n; i++) {
        e = prm[i];
        if (e > 0) {
            j = sarray[e-1];
            while (str[i+h] == str[j+h])
                h++;
            lcp[e] = h;
            if (h > 0) h--;
        }
    }
    lcp[0] = 0;
}

int main() {
    int sarray[MAXN], lcp[MAXN], i, n, j;
    char S[MAXN], T[MAXN];
    while (scanf("%s", S) == 1) {
        n = strlen(S);
        for (i = 0; i < n; i++)
            S[n+i] = S[i];
        S[n+n] = 0;
        build_sarray(S, sarray, lcp);
        for (i = 0; S[i]; i++)
            if (sarray[i] < n) {
                printf("%3d: %2d [%d]\n", i, lcp[i], n);
                for (j = 0; j < n; j++)
                    printf("%c", S[sarray[i]+j]);
                printf("\n");
            }
    }
}

```