Random Bit Generation and Stream Ciphers

Lecturer: Prof. Dr. Michael Eichberg

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Edition, William Stallings

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Random Numbers

- A number of network security algorithms and protocols based on cryptography make use of random binary numbers:
 - Key distribution and reciprocal **■** wechselseitige authentication schemes
 - Session key generation
 - Generation of keys for the RSA public-key encryption algorithm
 - Generation of a bit stream for symmetric stream encryption

There are two distinct requirements for a sequence of random numbers:

- Randomness
- Unpredictability

Randomness

- The generation of a sequence of allegedly random numbers being random in some well-defined statistical sense has been a concern.
- Two criteria are used to validate that a sequence of numbers is random:

Uniform distribution:

The frequency of occurence of ones and zeros should be

approximately equal.

Independence: No one subsequence in the sequence can be inferred from

the others.

Visualization of a Bad Random Number Generator (RNG)

| Expected distribution of random values in | Distribution of supposedly random values |
|---|--|
| 3D space. | of a bad rng in 3D space. |
| | |

Unpredictability

- The requirement is not just that the sequence of numbers be statistically random, but that the *successive members of the sequence are unpredictable*
- With **true** random sequences each number is statistically independent of other numbers in the sequence and therefore unpredictable.
 - True random numbers have their limitations, such as inefficiency, so it is more common to implement algorithms that generate sequences of numbers that appear to be random.
 - Care must be taken that an opponent not be able to predict future elements of the sequence on the basis of earlier elements.

Pseudorandom Numbers

Cryptographic applications typically make use of algorithmic techniques for random number generation.

- These algorithms are deterministic and therefore produce sequences of numbers that are not statistically random
- If the algorithm is good, the resulting sequences will pass many tests of randomness and are referred to as pseudorandom numbers.

Random and Pseudorandom Number Generators

Note

TRNG: True Random Number Generator **PRNG:** Pseudorandom Number Generator

PRF: Pseudorandom Function

True Random Number Generator (TRNG)

- Takes as input a source that is effectively random.
- The source is referred to as an entropy source and is drawn from the physical environment of the computer:
 - Includes things such as keystroke timing patterns, disk electrical activity, mouse movements, and instantaneous values of the system clock
 - The source, or combination of sources, serve as input to an algorithm that produces random binary output
- The TRNG may simply involve conversion of an analog source to a binary output
- The TRNG may involve additional processing to overcome any bias in the source

Pseudorandom Number Generator (PRNG) and Pseudorandom Function (PRF)

Two different forms of PRNG

Pseudorandom number generator

- An algorithm that is used to produce an open-ended sequence of bits.
- Input to a symmetric stream cipher is a common application for an openended sequence of bits.

Pseudorandom function (PRF)

- Used to produce a pseudorandom string of bits of some fixed length.
- Examples are symmetric encryption keys and nonces.

Pseudorandom Number Generator (PRNG) and Pseudorandom Function (PRF)

- Takes as input a fixed value, called the seed, and produces a sequence of output bits using a deterministic algorithm.
 - Quite often the seed is generated by a TRNG
- The output bit stream is determined solely by the input value or values, so an adversary who knows the algorithm and the seed can reproduce the entire bit stream!
- Other than the number of bits produced there is no difference between a PRNG and a PRF.

PRNG Requirements

- The basic requirement when a PRNG or PRF is used for a cryptographic application is that an adversary who does not know the seed is unable to determine the pseudorandom string.
- The requirement for secrecy of the output of a PRNG or PRF leads to specific requirements in the areas of:
 - Randomness
 - Unpredictability
 - Characteristics of the seed

Randomness

- The generated bit stream needs to appear random even though it is deterministic.
- There is no single test that can determine if a PRNG generates numbers that have the characteristic of randomness
- If the PRNG exhibits randomness on the basis of multiple tests, then it can be assumed to satisfy the randomness requirement

NIST SP 800-22 specifies that the tests should seek to establish three characteristics: (1) Uniformity, (2) Scalability, (3) Consistency

Randomness Tests

SP 800-22 lists 15 separate tests of randomness

Frequency test:

- The most basic test and must be included in any test suite
- Purpose is to determine whether the number of ones and zeros in a sequence is approximately the same as would be expected for a truly random sequence

Runs test:

- Focus of this test is the total number of runs in the sequence, where a run is an uninterrupted sequence of identical bits bounded before and after with a bit of the opposite value
- Purpose is to determine whether the number of runs of ones and zeros of various lengths is as expected for a random sequence

Maurer's universal statistical test:

- Focus is the number of bits between matching patterns.
- Purpose is to detect whether or not the sequence can be significantly compressed without loss of information. A significantly compressible sequence is considered to be non-random.

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Unpredictability

A stream of pseudorandom numbers should exhibit two forms of unpredictability:

Forward unpredictability:

If the seed is unknown, the next output bit in the sequence should be unpredictable in spite of any knowledge of previous bits in the sequence

Backward unpredictability:

- It should not be feasible to determine the seed from knowledge of any generated values.
- No correlation between a seed and any value generated from that seed should be evident.
- Each element of the sequence should appear to be the outcome of an independent random event whose probability is 1/2

The same set of tests for randomness also provides a test of unpredictability: A random sequence will have no correlation with a fixed value (the seed).

Seed Requirements

- The seed that serves as input to the PRNG must be secure and unpredictable
- The seed itself must be a random or pseudorandom number
- Typically the seed is generated by TRNG

| Generation of seeds | | |
|---------------------|--|--|
| | | |
| | | |

Algorithm Design

Algorithms fall into two categories:

1. Purpose-built algorithms

Algorithms designed specifically and solely for the purpose of generating pseudorandom bit streams.

2. Algorithms based on existing cryptographic algorithms.

Have the effect of randomizing input data.

Three broad categories of cryptographic algorithms are commonly used to create PRNGs:

- Symmetric block ciphers
- Asymmetric ciphers
- Hash functions and message authentication codes

Linear Congruential Generator

An algorithm first proposed by Lehmer that is parameterized with four numbers:

| m | the modulus | m > 0 |
|-------|-----------------------------|-----------------|
| a | the multiplier | 0 < a < m |
| c | the increment | $0 \le c < m$ |
| X_0 | the starting value, or seed | $0 \le X 0 < m$ |

The sequence of random numbers $\{X_n\}$ is obtained via the following iterative equation: $X_{n+1} = (aX_n + c) \mod m$

If m , a , c , and X_0 are integers, then this technique will produce a sequence of integers with each integer in the range $0 \le X_n < m$

The selection of values for a , c , and m is critical in developing a good random number generator.

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Blum Blum Shub (BBS) Generator

- Has perhaps the strongest public proof of its cryptographic strength of any purpose-built algorithm.
- Referred to as a *cryptographically secure pseudorandom bit generator (CSPRBG)*A CSPRBG is defined as one that passes the next-bit-test if there is not a polynomial-time algorithm that, on input of the first k bits of an output sequence, can predict the (k + 1)st bit with probability significantly greater than 1/2
- The security of BBS is based on the difficulty of factoring n.

Blum Blum Shub Block Diagram

Blum Blum Shub Block Diagram

n is the product of two (very large) primes n = pq.

The seed s should be an integer that is co-prime to n (i.e. p and q are not factors of s) and not 1 or 0.

Example - Blum Blum Shub (BBS) Generator

| i | x_i | B __ i |
|----|--------|------------------|
| 0 | 20749 | |
| 1 | 143135 | 1 |
| 2 | 177671 | 1 |
| 3 | 97048 | 0 |
| 4 | 89992 | 0 |
| 5 | 174051 | 1 |
| 6 | 80649 | 1 |
| 7 | 45663 | 1 |
| 8 | 69442 | 0 |
| 9 | 186894 | 0 |
| 10 | 177046 | 0 |

PRNG Using Block Cipher Modes of Operation

Two approaches that use a block cipher to build a PNRG have gained widespread acceptance:

- CTR mode: Recommended in NIST SP 800-90, ANSI standard X.82, and RFC 4086
- OFB mode: Recommended in X9.82 and RFC 4086

Generic Structure of a Typical Stream Cipher

| | Typical Stream Cipher | | | |
|----------------------------------|-----------------------|----------------------------------|--------------------------|--|
| plaintext p_i ciphertext c_i | | key K Initialization Value IV | state σ_i next-st | |
| keystream z_i | | | | |

Stream Cipher Design Considerations

The encryption sequence should have a large period:

A pseudorandom number generator uses a function that produces a deterministic stream of bits that eventually repeats; the longer the period of repeat the more difficult it will be to do cryptanalysis

The keystream should approximate the properties of a true random number stream as close as possible:

There should be an approximately equal number of 1s and 0s

If the keystream is treated as a stream of bytes, then all of the 256 possible byte values should appear approximately equally often

A key length of at least 128 bits is desirable:

The output of the pseudorandom number generator is conditioned on the value of the input key

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The same considerations that apply to block ciphers are valid

With a properly designed pseudorandom number generator a stream cipher can be as secure as a block cipher of comparable key length:

Entropy Sources

- A true random number generator (TRNG) uses a nondeterministic source to produce randomness
- Most operate by measuring unpredictable natural processes such as pulse detectors of ionizing radiation events, gas discharge tubes, and leaky capacitors
- Intel has developed a commercially available chip that samples thermal noise by amplifying the voltage measured across undriven resistors.

Comparison of PRNGs and TRNGs

| | Pseudorandom Number | True Random Number | |
|-------------|---------------------|-----------------------|--|
| | Generators | Generators | |
| Efficiency | Very efficient | Generally inefficient | |
| Determinism | Deterministic | Nondeterministic | |
| Periodicity | Periodic | Aperiodic | |

Conditioning

- A TRNG may produce an output that is biased in some way (such as having more ones than zeros or vice versa)
- Biased: NIST SP 800-90B defines a random process as biased with respect to an assumed discrete set of potential outcomes if some of those outcomes have a greater probability of occurring than do others.
- Entropy rate: NIST 800-90B defines entropy rate as the rate at which a digitized noise source provides entropy
 - Is a measure of the randomness or unpredictability of a bit string.
 - Will be a value between 0 (no entropy) and 1 (full entropy)
- Conditioning algorithms/deskewing algorithms:

Methods of modifying a bit stream to further randomize the bits

Typically conditioning is done by using a cryptographic algorithm to scramble the random bits so as to eliminate bias and increase entropy.

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The two most common approaches are the use of a hash function or a symmetric block cipher.