

Basics of Module Categories

Past 10 weeks

- Lots of examples of monoidal and tensor categories.
- Lots of technology for studying them.

Today

- What are module categories?
- Why should I care?
- How might we classify them?

Definition (Module Category).

,, $(\mathcal{C}, \otimes, \mathbb{1}, a, l, r)$

A **left module category** over a monoidal category \mathcal{C} consists of

- a category M ;
- a bifunctor $\otimes: \mathcal{C} \times M \rightarrow M$; "left action"
- a natural isomorphism $m_{X,Y,M}: (X \otimes Y) \otimes M \xrightarrow{\sim} X \otimes (Y \otimes M)$; "module associativity constraint"
- a natural isomorphism $l_M: \mathbb{1} \otimes M \xrightarrow{\sim} M$; "unit constraint"

for $X, Y \in \text{Ob}(\mathcal{C})$, $M \in \text{Ob}(M)$, such that the following commute for $X, Y, Z \in \text{Ob}(\mathcal{C})$, $M \in \text{Ob}(M)$:

$$\begin{array}{ccccc}
 & ((X \otimes Y) \otimes Z) \otimes M & & & \\
 a_{X,Y,Z} \circ id_M & \swarrow & \searrow m_{X \otimes Y, Z, M} & & \\
 (X \otimes (Y \otimes Z)) \otimes M & & (X \otimes Y) \otimes (Z \otimes M) & & \\
 & \downarrow m_{X, Y \otimes Z, M} & \downarrow m_{X, Y, Z \otimes M} & & \\
 X \otimes ((Y \otimes Z) \otimes M) & \xrightarrow{id_X \otimes m_{Y, Z, M}} & X \otimes (Y \otimes (Z \otimes M)) & & \\
 & & & r_X \circ id_M & \nearrow id_X \otimes l_M \\
 & & & & X \otimes M
 \end{array}$$

If \mathcal{C} is a tensor category, we ask

- that M be locally finite, Abelian;
- that \otimes be bilinear on morphisms, exact in the first variable (why?).

Regular representation: \mathcal{C} is a \mathcal{C} -module category with left action

$\otimes_{\mathcal{C}}$, $m := a$ and unit constraint l .

$\mathcal{C} := \text{Vec}$: every locally finite, Abelian M admits a \mathcal{C} -module structure given by

$$X \otimes M := M^{\oplus n}, \text{ for } X \in [k^n], M \in \mathcal{O}\mathcal{B}(U).$$

$\mathcal{C} := \text{Vec}_G$: module categories

are Abelian categories M

admitting a monoidal functor

$$F : \text{Cat}(G) \rightarrow \text{Aut}_{\otimes}(M).$$

Such an "action of G on M " can always be constructed by taking

$$F(g) : M \rightarrow S_g \otimes M.$$

If $F: \mathcal{C} \rightarrow M$ is a tensor functor,
 $X \otimes Y := F(X) \otimes Y$ defines a \mathcal{C} -module
structure on M .

E.g. $\text{Res}: \text{Rep}(G) \rightarrow \text{Rep}(H)$ induces a
 $\text{Rep}(G)$ -module structure on $\text{Rep}(H)$.

Proposition. There is a bijection
between \mathcal{C} -module structures on M
and monoidal functors $F: \mathcal{C} \rightarrow \text{End}(M)$.
(what about for tensor categories?)

$M := \text{Vec}$: there is a tensor
equivalence $M \cong \text{End}_\mathcal{C}(M)$
induced by the covariant Yoneda
embedding (i.e., $\mathcal{L}: V \mapsto \text{Hom}_M(V, -)$).

$\Rightarrow \text{Vec}$ admits a \mathcal{C} -module
structure if and only if
there is a fibre functor on \mathcal{C} !

Definition. A module subcategory is a full subcategory that is closed under \otimes .

Definition. A \mathcal{C} -module functor consists of

- a functor $F: M \rightarrow N$;
- a natural isomorphism $s_{X,M}: F(X \otimes M) \rightarrow X \otimes F(M)$, $X \in \text{Ob}(\mathcal{C})$, $M \in \text{Ob}(M)$;

satisfying the obvious coherence conditions.

Definition. A module equivalence is a module functor that gives an equivalence of categories.

Proposition. Given two \mathcal{C} -module categories M_1 and M_2 , we have a natural notion of a direct sum $M := M_1 \oplus M_2$ that is itself a \mathcal{C} -module category.

Definition. Indecomposable means there is no module equivalence with a non-trivial direct sum. Simple means no non-trivial module subcategories.

We have seen some examples of module categories, but there are some issues:

- our module functors do not descend to morphisms of Grothendieck group modules;
- even if we think back to Vec , there is just no hope of classifying all of its module categories!

Definition. Let \mathcal{C} be a tensor category with enough projectives. A \mathcal{C} -module category M is said to be exact if $P \otimes M$ is projective for all $M \in \text{Ob}(M)$ and projective $P \in \text{Ob}(\mathcal{C})$.

Somehow exact categories play an analogous role to projective modules.

$\mathcal{C} := \text{Vec}$: M is exact if and only if it is semisimple.

Any tensor category with enough projectives is exact over itself.

If \mathcal{C} is semisimple, then its exact module categories are semisimple.

$\mathcal{C} := \text{Rep}(G)$: assume semisimple
(i.e., \mathbb{k} does not divide $|G|$).

Every exact, indecomposable \mathcal{C} -module
is equivalent to some $\text{Rep}^*(\tilde{H})$,
where \tilde{H} is the central extension
of some $H \subseteq G$ by $\mathbb{k}^* := \mathbb{k} \setminus \{0\}$
and $\text{Rep}^*(\tilde{H})$ is the category
of \tilde{H} -modules that \mathbb{k}^* acts on
by scalar multiplication.

Let $G := \mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/2\mathbb{Z}$. We
have five subgroups: $G, \{0\}$ and
three isomorphic to $\mathbb{Z}/2\mathbb{Z}$.
This gives us five central extensions,
all of the form $\mathbb{k}^* \times H$. Not every
central extension is of this form, however;
in particular, because $Q_8/Z(Q_8) \cong G$,
 G has a non-trivial central extension
by the quaternion group Q_8 , given by the
sign homomorphism $\varepsilon : \pm 1 \mapsto \pm 1$.

Lemma. Exact module categories have enough projectives.

It follows that all exact module categories with finitely many isomorphism classes of simple objects are finite

Lemma. An object in an exact module category is projective if and only if it is injective.

Moreover, such categories are quasi-Frobenius and hence amenable to homological techniques.

Proposition. An indecomposable, exact \mathcal{C} -module category deategorifies to a simple module over $\text{Gr}(\mathcal{C})$.

If \mathcal{C} is a finite tensor category, there are only finitely-many \mathbb{Z}_+ -modules over $\text{Gr}(\mathcal{C})$ of the form $\text{Gr}(M)$.

Characterization Theorem #1

Let M_1, M_2 be \mathcal{C} -module categories. Then

- If M_1 is exact, then every additive module functor $F: M_1 \rightarrow M_2$ is exact.
- If every module functor $F: M_1 \rightarrow M_2$ is exact, then M_1 is exact.

One great implication of this theorem is that exact module functors descend to morphisms on Grothendieck modules!

We would now like to transition towards algebra objects and how they correspond to certain module categories.

Recall that an *algebra* in \mathcal{C} consists of

- an object $A \in \text{Ob}(\mathcal{C})$;
- a unit morphism $u: \mathbb{1} \rightarrow A$;
- a multiplication morphism $m: A \otimes A \rightarrow A$;

satisfying certain coherence conditions.

$\mathcal{C} := \text{Vec}$: an algebra object is a finite-dimensional, associative, unital \mathbb{k} -algebra.

$\mathcal{C} := \text{Vec}_G$: an algebra object is a G -graded algebra. Moreover, $\mathbb{k}L$ is an algebra for $L \subseteq G$.

$\mathcal{C} := \text{Rep}(G)$: the algebra of functionals $G \rightarrow \mathbb{k}$ is an algebra, with G acting as right translation.

A right module over an algebra A is an object $M \in \text{Ob}(\mathcal{C})$ together with a morphism $p: A \otimes M \rightarrow M$, subject to certain coherence conditions.

These A -modules form a left \mathcal{C} -module category $\text{Mod}_{\mathcal{C}}(A)$.

If \mathcal{C} has enough projectives (resp. is finite), then $\text{Mod}_{\mathcal{C}}(A)$ has enough projectives (resp. is finite).

$\mathcal{C} := \text{Vec}$, $M := \text{Vec}$: by our last remark, M cannot appear as some $\text{Mod}_{\mathcal{C}}(A)$.

Suppose \mathcal{C} is rigid: then 1 is an algebra, So is $A := X \otimes X^*$ for all $X \in \text{Ob}(\mathcal{C})$, with unit $u := \text{coev}_X$

and multiplication $m := \text{id}_X \otimes \text{ev}_X \otimes \text{id}_{X^*}$.

For all $Y \in \text{Ob}(\mathcal{C})$, $Y \otimes X^*$ is a module over A , and $Y \mapsto Y \otimes X^*$ is a module equivalence $\mathcal{C} \rightarrow \text{Mod}_{\mathcal{C}}(A)$.

Let M be a module category over a finite tensor category \mathcal{C} and fix $M_1, M_2 \in \text{Ob}(M)$. Then the functor $X \mapsto \text{Hom}_M(X \otimes M_1, M_2)$ is represented by $\underline{\text{Hom}}(M_1, M_2) \in \text{Ob}(\mathcal{C})$.

This representative is known as the internal Hom. The Yoneda lemma gives us a \mathcal{C} -module functor

$\underline{\text{Hom}}(-, -): (M_1, M_2) \mapsto \underline{\text{Hom}}(M_1, M_2)$.

Characterization Theorem #2

Let \mathcal{M} be a \mathcal{C} -module category.

(i). If \mathcal{M} is *finite*, then $\mathcal{M} \cong \text{Mod}_{\mathcal{C}}(A)$ for some algebra $A \in \text{Ob}(\mathcal{C})$.

(ii). If \mathcal{M} is *exact* and $M \in \text{Ob}(\mathcal{M})$ is an object such that the isomorphism class $[M]$ generates $\text{Gr}(\mathcal{M})$ as a \mathbb{Z}_+ -module over $\text{Gr}(\mathcal{C})$, then $\mathcal{M} \cong \text{Mod}_{\mathcal{C}}(A)$ for $A := \underline{\text{Hom}}(M, M)$.

Thus we are able to classify all *finite* and all *exact*, "*cyclic*" module categories in terms of *algebras*!

This is the part of the talk where I try to sell you subfactors.
I owe this section to Pinhas and Annael.

One of the most extensively studied "exotic" subfactors is the Haagerup subfactor $N \subseteq M$. Let $M = N \oplus X$ (that is, X is in some sense the "*complement*" of N in M). We find that as an (N, N) -bimodule,

$$X \otimes X \cong N \oplus X \oplus (g \otimes X) \oplus (g^2 \otimes X),$$

where g is a one-dimensional (N, N) -bimodule satisfying

$$g^3 \cong N \quad \text{and} \quad g \otimes X \cong X \otimes g^2.$$

This gives us the Haagerup category. In other words, it is the fusion category with simple objects N , g , g^2 , X , $g \otimes X$ and $g^2 \otimes X$, subject to the fusion rules

$$g^3 \cong N,$$

$$g \otimes X \cong X \otimes g^2,$$

$$X \otimes X \cong N \oplus X \oplus (g \otimes X) \oplus (g^2 \otimes X).$$

It was shown by Pinhas and Noah in 2012 that this category admits three "simple" module categories, given by the algebra objects N , $N \otimes X$ and $N \otimes g \otimes g^2$.

Thank you very much for your time!

