

# Qiskit 7 - Shores algorithm

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In [109]:

```
import numpy as np
from math import gcd

from qiskit import QuantumRegister, ClassicalRegister, QuantumCircuit, Aer, execute, IBMQ
from qiskit.tools.visualization import circuit_drawer, plot_histogram
```

## 1 - Introduction

### Problem

The problem of factoring integers can be expressed in such way, that based on some integer  $N$ , one want to find the integers  $N_1$  and  $N_2$ , such that:  $N_1 N_2 = N$ , meanwhile:  $1 < N_1, N_2 < N$

Shores algorithm approaches this problem by reducing it into finding the period of a certain function.

### Algorithm

**step 1** Pick an integer  $N$  and use a classical algorithm to determine if it is prime or a power of prime. If so, exit.

**step 2** Randomly choose an integer  $a$  such that  $1 < a < N$ . Perform Euclid's algorithm to determine if the  $GCD(a, N)$  is 1. If not, exit.

**step 3** Use the quantum circuit represented by the unitary operator  $U_{f_{a,N}}$  to find a period  $r$ .

**step 4** If  $a$  is odd, or if  $a^{\frac{r}{2}} \equiv -1 \pmod{N}$ , then return to step 2 and choose another  $a$ .

**step 5** Use Euclid's algorithm to calculate  $GCD((a^{\frac{r}{2}} + 1), N)$  and  $GCD((a^{\frac{r}{2}} - 1), N)$ . Return at least one of the nontrivial solutions.

## 2 - Choose N

So we will work with a non-prime, non power of prime positive integer such that it is manageable on a handful couple of qubits. Allow me to put together the following list of candidates:

- List of possible integers = [4, 6, 8, 10, 12, 14, 15, 18]

In [86]:

```
N = 4
N_bit = format(N, 'b')
print('N =', N)
print('N bit =', N_bit)
```

```
N = 4
N_bit = 100
```

## 3 - Get a

If  $GCD(a, N) = 1$   $a$  is a co-prime of  $N$  and that is what we can use in following steps. Hence we ought to pull random  $a$  and check its common factors against  $N$ , if it is a 1 we keep it.

In [87]:

```
for i in range(N):  
    a = np.random.randint(2, N)  
    GCD = gcd(a, N)  
    if GCD is 1:  
        print('gcd is 1 for a:', a)  
        break
```

gcd is 1 for a: 3

## 4 - Phase estimation

### The modular function

First up is to find the powers of  $a \bmod N$ , that is:

$$a^0 \bmod N, a^1 \bmod N, a^3 \bmod N, \dots$$

In other words we are to find the values  $x$  of the function:

$$f_{a,N}(x) = a^x \bmod N$$

However we are rather interested in the period  $r$  of this function such that:

$$f_{a,N}(r) = a^r \bmod N$$

It is known from a number theory theorem that for any co-prime  $a \leq N$ , the function  $f_{a,N}(r)$  will output a 1 for some  $r \leq N$ . After it hits 1, the sequence of numbers will simply repeat itself. *(isn't this a mechanism we could use for reformulated optimization problems?).*

### State dynamics

Next step is to implement the function  $f_{a,N}$  on a quantum circuit, in order to do that we are to proceed according to the following procedure:

1: Define 2 quantum registers. The first one  $|x\rangle_m$  and the second one  $|y\rangle_n$ . Where  $m$  = the binary length needed to hold  $N$ , and where  $n$  = the binary length needed to hold the periodic base of  $N$ .

|  $q$

2: Place the  $|x\rangle_m$  qubits in an equally weighted superposition:

|  $q$

3: Evaluate the function  $f_{a,N}(x)$  for all the superpositioned possibilities:

|  $q$

4: By measuring the bottom qubits  $|y\rangle_n$  we obtain an estimate  $a^{\bar{x}} \bmod N$  for some  $\bar{x}$ . By the periodicity of  $f_{a,N}$  we also have that:

$$a^{\bar{x}} \equiv a^{\bar{x}+r} \bmod N$$

and,

$$a^{\bar{x}} \equiv a^{\bar{x}+2r} \bmod N$$

such that for any  $s \in \mathbb{Z}$  we have:

$$a^{\bar{x}} \equiv a^{\bar{x}+sr} \bmod N$$

Furthermore we have that out of the  $2^m$  superpositions in  $|x\rangle$  in state  $|\varphi_2\rangle$ , there are  $\frac{2^m}{r}$  of them that has the solution  $\bar{x}$ . This finally gives us:

|  $q$

Where  $t_0$  is the first time the measured the value of  $a^{t_0} \equiv a^{\bar{x}} \bmod N$

To boil this down, we want to estimate  $\bar{x}$  since that gives us a value to plug into the function  $f_{a,N}$  which gives us the period base we are looking for. Hence this stage starts with preparing the states and end with taking the measurements of  $|y\rangle$  such that a satisfying value can be detected. *(Do we use that value and re instantiate the  $x$  qubits? Or is the superpositioned  $x$  state to be 'automatically manipulated' by our measurements of  $y$  such that it simply takes on our preferred value?)*

### Quantum circuit

So how do we construct the black box unitary? Allow us to establish a relationship between the period of the function  $f_{a,N}$  and the phase value of the eigenvalue. this way solving the phase helps us find the period.

Phase estimation can be described such that if we know  $U$  and  $|\psi\rangle$ , we can estimate  $\phi$ :

$U$  |

Proceeding with the following specifications:

$$\begin{aligned} N &= 4 = 100_{bin} \\ a &= 3 \\ n &= 3 \\ r &= 2 = 01_{bin} \\ freq &= (1, 3) \\ m &= 6 \end{aligned}$$

### Define registers

In [95]:

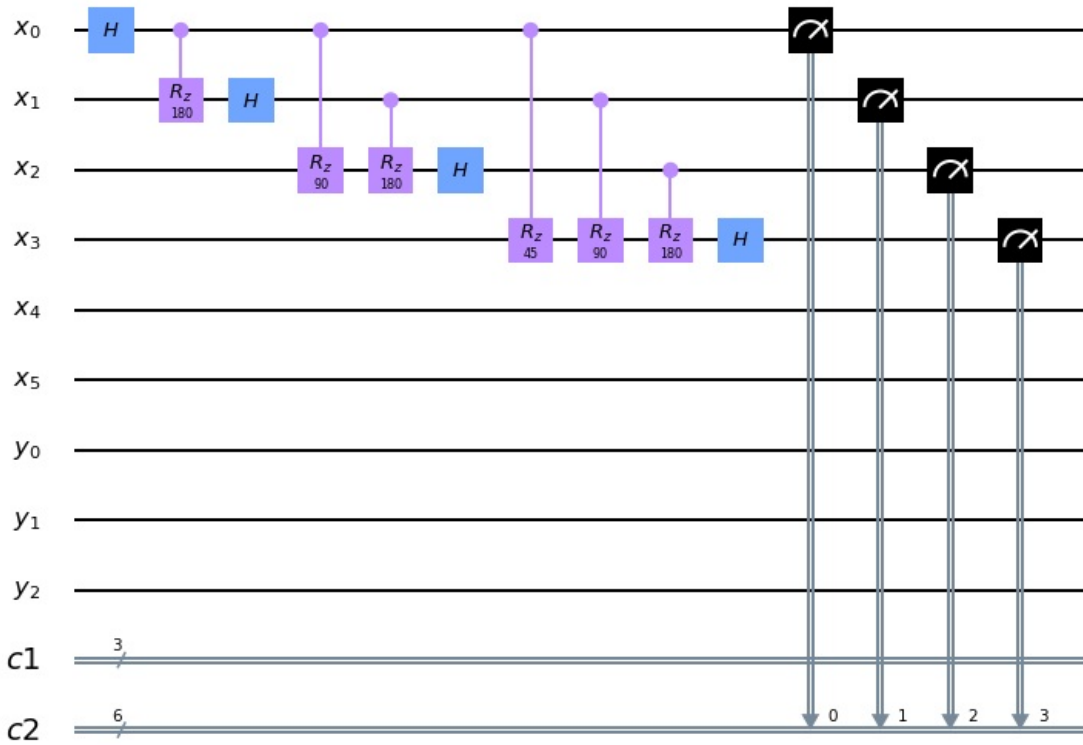
```
x = QuantumRegister(6, 'x')
y = QuantumRegister(3, 'y')
c1 = ClassicalRegister(3, 'c1')
c2 = ClassicalRegister(6, 'c2')
```

## QFT - Quantum Fourier Transform

In [121]:

```
q1 = QuantumCircuit(x, y, c1, c2)
q1.barrier()
q1.h(x[0])
q1.crz(180, x[0], x[1])
q1.h(x[1])
q1.crz(90, x[0], x[2])
q1.crz(180, x[1], x[2])
q1.h(x[2])
q1.barrier()
q1.crz(45, x[0], x[3])
q1.crz(90, x[1], x[3])
q1.crz(180, x[2], x[3])
q1.h(x[3])
q1.barrier()
q1.measure([x[0], x[1], x[2], x[3]], [c2[0], c2[1], c2[2], c2[3]])
circuit_drawer(q1, output='mpl', plot_barriers=False)
```

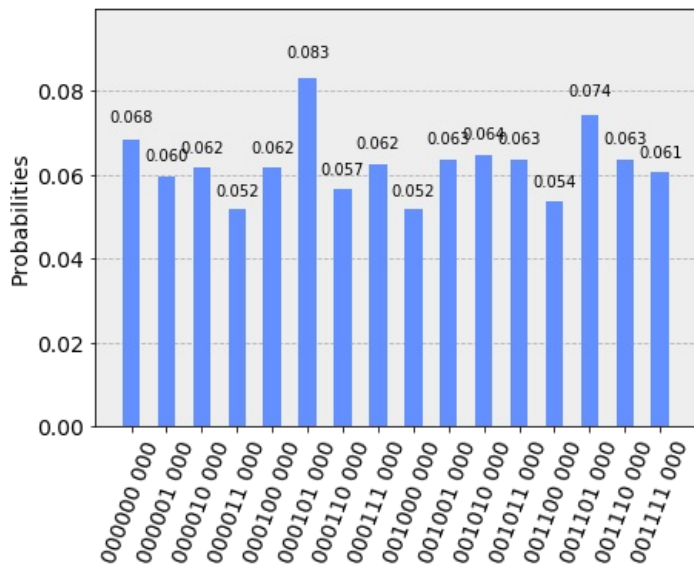
Out[121]:



In [125]:

```
sim = Aer.get_backend('qasm_simulator')
count = execute(q1, sim).result().get_counts()
plot_histogram(count)
```

Out[125]:



To be continued...

In [ ]: