

COE 308 – Computer Architecture

Assignment 3 Solution

Floating-Point Representation and Arithmetic

1. (4 pts) What is the decimal value of the following single-precision floating-point numbers?

a) 1010 1101 0001 0100 0000 0000 0000 0000 (binary)

b) 0100 0110 1100 1000 0000 0000 0000 0000 (binary)

Solution:

a) Sign is negative

$$\text{Exponent value} = 01011010_2 - 127 = -37$$

$$\text{Significand} = 1.001\ 0100\ 0000\ 0000\ 0000\ 0000_2$$

$$\text{Decimal value} = -1.00101_2 \times 2^{-37} = -1.15625 \times 2^{-37} = -8.412826 \times 10^{-12}$$

b) Sign is positive

$$\text{Exponent value} = 10001101_2 - 127 = 14$$

$$\text{Significand} = 1.100\ 1000\ 0000\ 0000\ 0000\ 0000_2$$

$$\text{Decimal value} = 1.1001_2 \times 2^{14} = 1.5625 \times 2^{14} = 25600$$

2. (3 pts) Show the IEEE 754 binary representation for: -75.4 in ...

a) Single Precision

b) Double precision

Solution:

$$75 = 1001011_2$$

$$0.4 = 0.0110_2 = 0.01100110_2 \dots$$

$$75.4 = 1001011.\overline{0110}_2 = 1.0010110110_2 \times 2^6$$

a) Single-Precision: Biased exponent = $6 + 127 = 133$

$$1\ 10000101\ 0010110110011001100110\textcolor{red}{1}_2 \text{ (rounded to nearest)}$$

b) Double-Precision: Biased exponent = $6 + 1023 = 1029$

$$1\ 10000000101$$

$$0010110110011001100110011001100110011001100110\textcolor{red}{10}_2 \text{ (rounded)}$$

3. (6 pts) $x = 1100\ 0110\ 1101\ 1000\ 0000\ 0000\ 0000\ 0000$ (binary)

and $y = 0011\ 1110\ 1110\ 0000\ 0000\ 0000\ 0000\ 0000$ (binary)

are single-precision floating-point numbers. Perform the following operations showing all work:

a) $x + y$

b) $x * y$

Solution:

$$\text{Value of Exponent}(x) = 10001101_2 - 127 = 14$$

$$x = -1.101\ 1000\ 0000\ 0000\ 0000\ 0000_2 \times 2^{14}$$

$$\text{Value of Exponent}(y) = 01111101_2 - 127 = -2$$

$$y = 1.110\ 0000\ 0000\ 0000\ 0000\ 0000_2 \times 2^{-2}$$

a) $x + y$

$$\begin{array}{r}
 -\ 1.101\ 1000\ 0000\ 0000\ 0000\ 0000_2 \times 2^{14} \\
 +\ 1.110\ 0000\ 0000\ 0000\ 0000\ 0000_2 \times 2^{-2} \\
 \hline
 -\ 1.101\ 1000\ 0000\ 0000\ 0000\ 0000_2 \times 2^{14} \\
 +\ 0.000\ 0000\ 0000\ 0000\ 1110\ 0000_2 \times 2^{14} \text{ (shift right 16)} \\
 \hline
 1\ 0.010\ 1000\ 0000\ 0000\ 0000\ 0000_2 \times 2^{14} \text{ (2's complement)} \\
 0\ 0.000\ 0000\ 0000\ 0000\ 1110\ 0000_2 \times 2^{14} \\
 \hline
 1\ 0.010\ 1000\ 0000\ 0000\ 0000\ 1110\ 0000_2 \times 2^{14} \text{ (add)} \\
 \hline
 -\ 1.101\ 0111\ 1111\ 1111\ 0010\ 0000_2 \times 2^{14} \text{ (2's complement)}
 \end{array}$$

Result is negative and is normalized

All shifted out bits were zeros, so result is also exact

$$x + y = 1\ 10001101\ 101\ 0111\ 1111\ 1111\ 0010\ 0000_2$$

b) $x * y$

$$\text{Biased exponent}(x*y) = 10001101_2 + 01111101_2 - 127$$

$$\text{Biased exponent}(x*y) = 139 = 10001011_2$$

$$\text{Sign}(x*y) = 1 \text{ (negative)}$$

$$\begin{array}{r}
 1.101\ 1000\ 0000\ 0000\ 0000\ 0000_2 \\
 \times 1.110\ 0000\ 0000\ 0000\ 0000\ 0000_2 \\
 \hline
 1\ 1111 \\
 1101\ 1000\ 0000\ 0000\ 0000\ 0000_2 \\
 11011\ 0000\ 0000\ 0000\ 0000\ 000_2 \\
 1.10110\ 0000\ 0000\ 0000\ 0000\ 00_2
 \end{array}$$

$$10.11110\ 1000\ 0000\ 0000\ 0000\ 0000_2$$

Normalize by shifting right 1 bit and increment exponent

$$\text{Significand} = 1.011\ 1101\ 0000\ 0000\ 0000\ 0000_2$$

$$\text{Biased exponent} = 139+1 = 140 = 10001100_2$$

Significand is already rounded

$$x*y = 1\ 10001100\ 011\ 1101\ 0000\ 0000\ 0000\ 0000_2$$

4. (4 pts) $x = 0101\ 1111\ 1011\ 1110\ 0100\ 0000\ 0000\ 0000$ (in binary)
 and $y = 0011\ 1111\ 1111\ 1000\ 0000\ 0000\ 0000\ 0000$ (in binary)
 and $z = 1101\ 1111\ 1011\ 1110\ 0100\ 0000\ 0000\ 0000$ (in binary)
 represent single precision IEEE 754 floating-point numbers. Perform the following operations showing all work:

- $x + y$
- Result of (a) + z
- Why is the result of (b) counterintuitive?

Solution:

- $x = 1.011\ 1110\ 0100\ 0000\ 0000\ 0000_2 \times 2^{64}$
 $y = 1.111\ 1000\ 0000\ 0000\ 0000\ 0000_2 \times 2^0$
 Difference in exponent = 64
 Shift significand of y right by 64 bits and add to x
 The significand bits of y are truncated after rounding
 $x + y = x$ because y is too small with respect to x
 Therefore, $x + y = 1.011\ 1110\ 0100\ 0000\ 0000\ 0000_2 \times 2^{64}$
- Result of (a) is $x = 0\ 10111111\ 011111001000000000000000_2$
 $z = 1\ 10111111\ 011111001000000000000000_2 = -x$
 Therefore, Result of (a) + $z = x - x = 0$
 $0\ 00000000\ 000000000000000000000000_2$
- We are computing $(x+y) + z$ where $z = -x$
 Intuitively $(x+y) + -x = y$ which is not 0
 However, in this example $(x+y) + -x = 0$
 This is because we have limited number of fraction bits

5. (3 pts) IA-32 offers an 80-bit extended precision option with a 1 bit sign, 16-bit exponent, and 63-bit fraction (64-bit significand including the implied 1 before the binary point). Assume that extended precision is similar to single and double precision.

- What is the bias in the exponent?
- What is the range (in absolute value) of normalized numbers that can be represented by the extended precision option?

Solution:

- With a 16-bit exponent, bias = $2^{15} - 1 = 32767$
- largest normalized $\approx 2 \times 2^{32767} = 2^{32768} = 1.415... \times 10^{9864}$
 smallest normalized: $1.0 \times 2^{-32766} = 2.8259... \times 10^{-9864}$