

CSC3100 Data Structures Tutorial 11: Sorting

Jiawei Lian

The Chinese University of Hong Kong, Shenzhen
School of Data Science

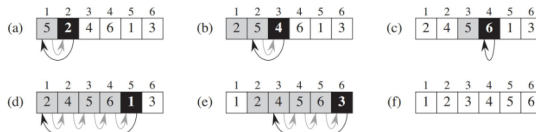
119030043@link.cuhk.edu.cn

December 6, 2021

- 1 Review: Insertion Sort
- 2 Review: Quick Sort
- 3 Problem 1: Insertion Sort List
- 4 Problem 2: Find K Smallest Numbers

- 1 Review: Insertion Sort
- 2 Review: Quick Sort
- 3 Problem 1: Insertion Sort List
- 4 Problem 2: Find K Smallest Numbers

Review: Insertion Sort



INSERTION-SORT(A)

```
1  for  $j = 2$  to  $A.length$ 
2     $key = A[j]$ 
3    // Insert  $A[j]$  into the sorted
      sequence  $A[1 \dots j - 1]$ .
4     $i = j - 1$ 
5    while  $i > 0$  and  $A[i] > key$ 
6       $A[i + 1] = A[i]$ 
7       $i = i - 1$ 
8     $A[i + 1] = key$ 
```

Figure: Insertion sort

Outline

- 1 Review: Insertion Sort
- 2 Review: Quick Sort
- 3 Problem 1: Insertion Sort List
- 4 Problem 2: Find K Smallest Numbers

Main Idea: Divide-and-Conquer

① Divide:

- Take one element $A[q]$ from A as pivot, then partition the rest of A into two sub-arrays $A[p \cdots q - 1]$ and $A[q + 1 \cdots r]$, such that $A[p \cdots q - 1] \leq A[q] < A[q + 1 \cdots r]$

② Conquer:

- Sort $A[p \cdots q - 1]$ and $A[q + 1 \cdots r]$ by recursively calling Quick Sort.

Partition

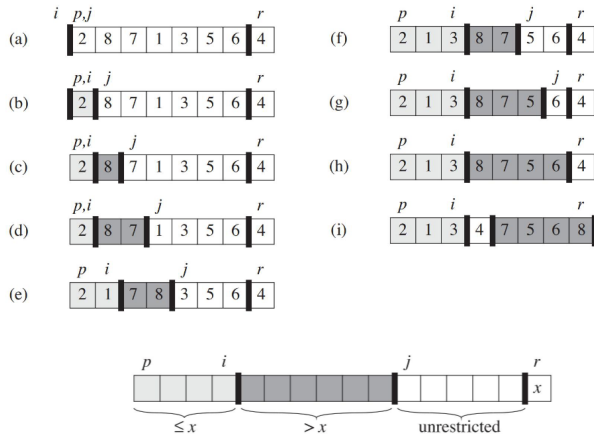


Figure: Partition process

```
PARTITION( $A, p, r$ )  
1   $x = A[r]$   
2   $i = p - 1$   
3  for  $j = p$  to  $r - 1$   
4      if  $A[j] \leq x$   
5           $i = i + 1$   
6          exchange  $A[i]$  with  $A[j]$   
7  exchange  $A[i + 1]$  with  $A[r]$   
8  return  $i + 1$ 
```

Figure: Partition pseudo code

QUICKSORT(A, p, r)

```
1  if  $p < r$   
2       $q = \text{PARTITION}(A, p, r)$   
3      QUICKSORT( $A, p, q - 1$ )  
4      QUICKSORT( $A, q + 1, r$ )
```

Figure: Recursion pseudo code

Outline

- 1 Review: Insertion Sort
- 2 Review: Quick Sort
- 3 Problem 1: Insertion Sort List**
- 4 Problem 2: Find K Smallest Numbers

Problem 1: Insertion Sort List - Solution

```
class Solution {
    public ListNode insertionSortList(ListNode head) {
        ListNode dummy = new ListNode(0), pre;
        dummy.next = head;

        while(head != null && head.next != null) {
            if(head.val <= head.next.val) {
                head = head.next;
                continue;
            }

            pre = dummy;

            while (pre.next.val < head.next.val) pre = pre.next;

            ListNode curr = head.next;
            head.next = curr.next;
            curr.next = pre.next;
            pre.next = curr;
        }
        return dummy.next;
    }
}
```

Figure: Solution of example 1

Outline

- 1 Review: Insertion Sort
- 2 Review: Quick Sort
- 3 Problem 1: Insertion Sort List
- 4 Problem 2: Find K Smallest Numbers

Problem 2: Find K Smallest Numbers

1 Problem statement:

- Given a group of N numbers, determine the first K^{th} smallest numbers, where $k \leq N$.

2 Example:

- Input: $arr = [1, 3, 5, 7, 2, 4, 6, 8]$, $k = 4$
- Output: $[1, 3, 2, 4]$

3 Constraints:

- $0 \leq \text{len}(arr) \leq 100000$
- $0 \leq k \leq \min(100000, \text{len}(arr))$

Find K Smallest Numbers - Solution 1

- 1 Directly sort this group of N numbers in an ascending order
- 2 Return the first k numbers

Find K Smallest Numbers - Solution 2

- ① Read the first K numbers into an array with size of K, sort these numbers in an ascending order
- ② Read the remaining elements one by one
 - Ignore the element if it is larger or equal to the K^{th} element in the array
 - Otherwise, place the element into its correct position
- ③ After checking all the numbers, the remaining array is the answer.

Find K Smallest Numbers - Solution 3

1 Use Quick Sort

```
class Solution {
public:
    int[] smallestK(int[] arr, int k) {
        quickSort(arr, 0, arr.length - 1, k);
        int[] result = new int[k];
        for (int i = 0; i < k; i++) {
            result[i] = arr[i];
        }
        return result;
    }

    private static void quickSort(int[] arr, int startIndex, int endIndex, int k) {
        if (startIndex >= endIndex) {
            return;
        }
        int partitionIndex = quickSortPartition(arr, startIndex, endIndex);
        if (partitionIndex == k) {
            return;
        }
        if (k < partitionIndex) {
            quickSort(arr, startIndex, partitionIndex - 1, k);
        } else {
            quickSort(arr, partitionIndex + 1, endIndex, k);
        }
    }

    private static int quickSortPartition(int[] arr, int startIndex, int endIndex) {
        int pivot = arr[endIndex];
        int i = startIndex;
        for (int j = startIndex; j <= endIndex - 1; j++) {
            if (arr[j] < pivot) {
                int temp = arr[j];
                arr[j] = arr[i];
                arr[i] = temp;
                i++;
            }
        }
        int temp = arr[endIndex];
        arr[endIndex] = arr[i];
        arr[i] = temp;
        return i;
    }
}
```

Figure: Solution of example 2

References I

- <https://leetcode.com/problems/insertion-sort-list/>
- <https://leetcode-cn.com/problems/smallest-k-lcci/>

CSC3100 Data Structures Tutorial 11: Sorting

Jiawei Lian

The Chinese University of Hong Kong, Shenzhen
School of Data Science

119030043@link.cuhk.edu.cn

December 6, 2021