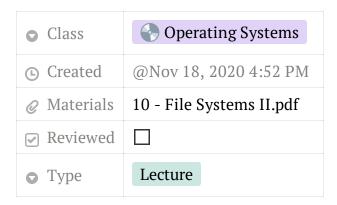


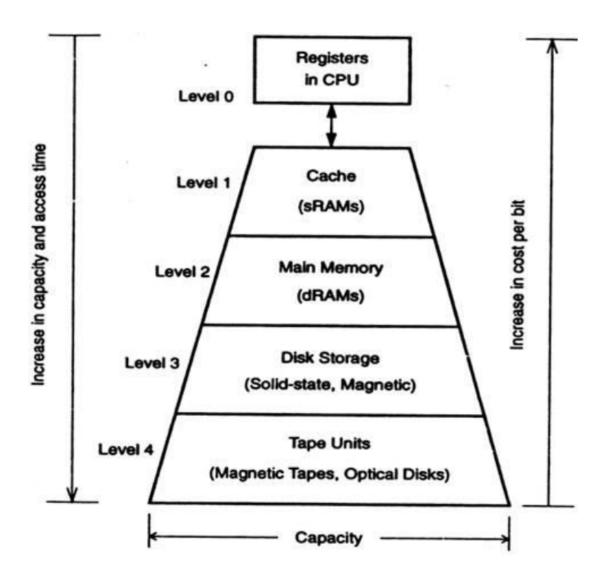
Operating Systems W12L2 - File Systems II



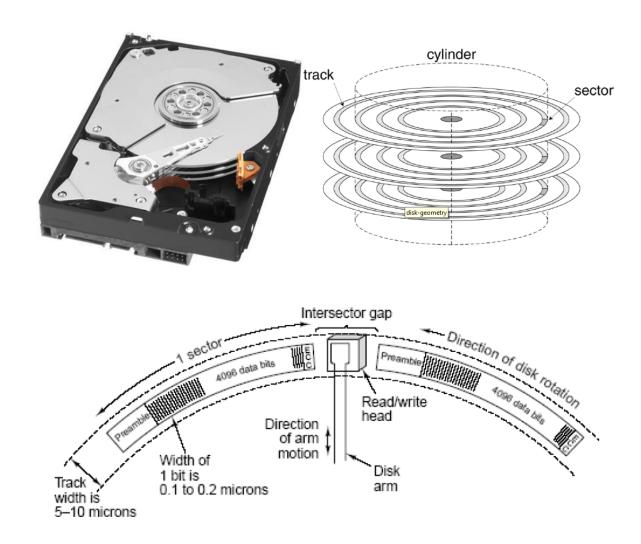
Initial Discussion

▼ Storage Capacity Hierarchy

Disk storage is abstracted by OS as *files*



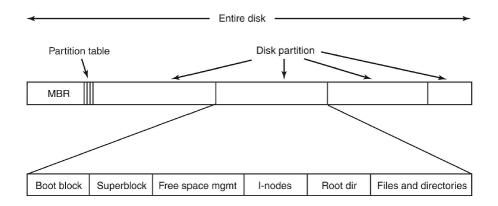
- Going over hard disks today first, SSDs will come later
- Mechanical movements automatically means slower speeds
- Lots of hard drive anatomy
- ▼ Surface is a group of tracks, track is a group of sectors, sector is a group of bytes, and cylinder is several tracks on corresponding surfaces



 Building a reliable program is difficult and part of what defines being "good" at CS (versus, say, programming in C)

MBR and Partition Table

- Starting and ending addresses of each partition
- BIOS executes MBR which finds active partition and executes boot block
- **▼** Diagram of disks and partitions



• A *superblock* records characteristics about file system — Its size is fixed for this since OS needs to depend on it

Implementing File Systems

▼ #1 - Contiguous Allocation

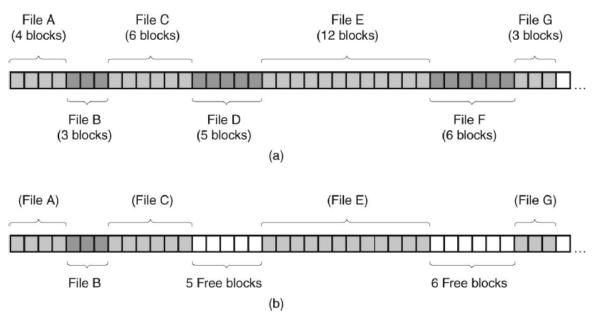
Each file is stored as a contiguous run of disk blocks

Advantages include...

- Simple implementation
- · Great read speed

Disadvantages include...

- Disk could become fragmented
- Need to know size of file when created



After files D and F were deleted

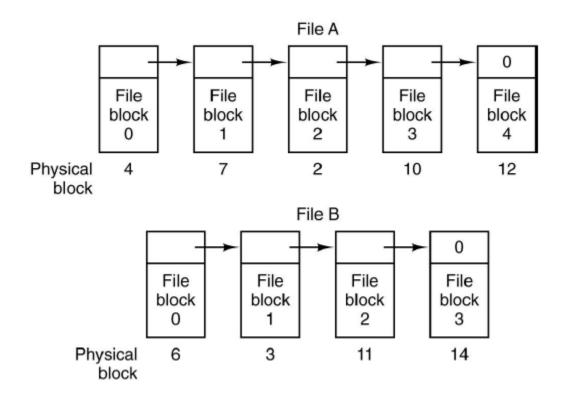
▼ #2 - LinkedList Allocation

The first 'word' of each block is a pointer to the next Advantages include...

- No (external) fragentation because empty blocks can be added to the next
- Directory only needs to store disk address of first block

Disadvantages include...

- Random access is extremely slow
- Data storage is no longer a power of two since pointer takes up a few bytes (Some binary numbers won't map to anything as a result)



These will continue next week