# Genuine equivariant operads

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# April 28, 2017

## Abstract

We build new algebraic structures, which we call genuine equivariant operads, which can be thought of as a hybrid between equivariant operads and coefficient systems. We then prove an Elmendorf type theorem stating that equivariant operads, with their graph model structure, are equivalent to genuine equivariant operads with their projective model structure.

As an application, we build explicit models for the  $N_{\infty}$ -operads of Blumberg and Hill.

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### 1 Introduction

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UNDERLEFTADJ LEM

#### 2 Basic definitions

In this section we recall some definitions that will be used throughout.

Recall that for a diagram category  $\mathcal{D}$  and functor  $\mathcal{I}_{\bullet}$ 

$$\mathcal{D} \xrightarrow{\mathcal{I}_{\bullet}} \mathsf{Cat}$$

$$d \longmapsto \mathcal{I}_{d}$$

$$(2.1)$$

the (covariant) Grothendieck construction  $\mathcal{D} \ltimes \mathcal{I}_{\bullet}$  has objects pairs (d,i) with  $d \in \mathcal{D}, i \in \mathcal{I}_d$ and arrows  $(d, i) \rightarrow (d', i')$  given by pairs

$$(f:d \rightarrow d', g:f_*(i) \rightarrow i'),$$

where  $f_*: \mathcal{I}_d \to \mathcal{I}_{d'}$  is a shorthand for the functor  $\mathcal{I}_{\bullet}(f)$ .

We now discuss a basic property of over and under categories that will be used in §6.4. Given  $\mathcal{J}, \mathcal{C} \in \mathsf{Cat}$  and  $j \in \mathcal{J}$  we will let  $\mathcal{C}^{\downarrow j}$  denote the Grothendieck construction for the functor

$$\mathcal{J} \longrightarrow \mathsf{Cat}$$
 $i \longmapsto \mathcal{C}^{\mathcal{J}(i,j)}$ 

Explicitly, an object of  $\mathcal{C}^{\downarrow j}$  is a pair  $(i, \mathcal{J}(i,j) \xrightarrow{\varphi} \mathcal{C})$  and an arrow  $(i,\varphi) \to (i',\varphi')$  is a pair  $(I: i \to i', \gamma: \varphi \circ I^* \to \varphi').$ 

**Lemma 2.2.** Let  $\mathcal{J} \in \mathsf{Cat}$  be a small category and  $j \in \mathcal{J}$ . One then has adjunctions

$$(-\downarrow j)$$
:  $\mathsf{Cat}_{/\mathcal{J}} \rightleftarrows \mathsf{Cat}: (-)^{\downarrow j}, \qquad (j \downarrow -)$ :  $\mathsf{Cat}_{/\mathcal{J}} \rightleftarrows \mathsf{Cat}: (-)^{j\downarrow}.$ 

*Proof.* Since  $j \downarrow \mathcal{I} = (\mathcal{I}^{op} \downarrow j)^{op}$  by defining  $(\mathcal{C}^{j\downarrow}) = ((\mathcal{C}^{op})^{\downarrow j})^{op}$  one reduces to the leftmost

Given  $\mathcal{I} \xrightarrow{\pi} \mathcal{J}$  and  $\mathcal{C}$  we will show that functors  $\mathcal{I} \downarrow j \xrightarrow{F} \mathcal{C}$  correspond to functors  $\mathcal{I} \xrightarrow{G} \mathcal{C}^{\downarrow j}$ over  $\mathcal{J}$ .

On objects, F associates to each pair  $(i, J: \pi(i) \to j)$  an object  $F(i, J) \in \mathcal{C}$ . One thus sets  $G(i) = (\pi(i), F(i, -))$  and these are clearly inverse processes.

On arrows F associates to  $(i, J' \circ \pi(I)) \xrightarrow{I} (i', J')$  an arrow  $F(i, J' \circ \pi(I)) \xrightarrow{F(I)} F(i', J')$ . One thus defines

$$G(I) = \left(\pi(i) \xrightarrow{\pi(I)} \pi(i'), F(i, (-) \circ \pi(i)) \xrightarrow{F(I)} F(i', -)\right)$$

and again it is clear that these are inverse processes. Finally, the fact that the associativity and unit conditions for F, G coincide is likewise clear.

# 3 Planar and tall maps

## 3.1 Planar structures

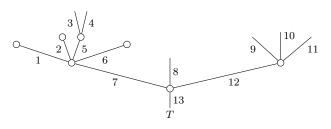
Throughout we will work with trees possessing *planar structures* or, more intuitively, trees embedded into the plane.

Our preferred model for trees will be that of broad posets first introduced by Weiss in [8] and further worked out by the second author in [5]. We now define planar structures in this context.

**Definition 3.1.** Let  $T \in \Omega$  be a tree. A *planar structure* of T is an extension of the descendancy partial order  $\leq_d$  to a total order  $\leq_p$  such that:

• Planar: if  $e \leq_p f$  and  $e \nleq_d f$  then  $g \leq_d f$  implies  $e \leq_p g$ .

**Example 3.2.** An example of a planar structure on a tree T follows, with  $\leq_r$  encoded by the number labels.



PLANAREX EQ

Intuitively, given a planar depiction of a tree T,  $e \leq_d f$  holds when the downward path from e passes through f and  $e \leq_p f$  holds if either  $e \leq_d f$  or if the downward path from e is to the left of the downward path from f (as measured at the node where the paths intersect).

Intuitively, a planar depiction of a tree amounts to choosing a total order for each of the sets of *input edges* of each node (i.e. those edges immediately above that node).

While we will not need to make this last statement precise, we will nonetheless find it convenient to show that Definition 3.1 is equivalent to such choosing total orders for each of the sets of input edges. To do so, we first introduce some notation.

**Notation 3.4.** Let  $T \in \Omega$  be a tree and  $e \in T$  and edge. We will denote

$$I(e) = \{ f \in T : e \leq_d f \}$$

and refer to this poset as the *input path of* e.

We will repeatedly use the following, which is a consequence of [5, Cor. 5.26]

**Lemma 3.5.** If  $e \leq_d f$ ,  $e \leq_d f'$ , then f, f' are  $\leq_d$ -comparable.

**Proposition 3.6.** Let  $T \in \Omega$  be a tree. Then

- (a) for any  $e \in T$  the finite poset I(e) is totally ordered;
- (b) the poset  $(T, \leq_d)$  has all joins, denoted  $\vee$ . In fact,  $\bigvee_i e_i = \min(\bigcap_i I(e_i))$ .

*Proof.* (a) is immediate from Lemma 3.5. To prove (b) we note that  $\min(\bigcap_i I(e_i))$  exists by (a), and that this is clearly the join  $\bigvee e_i$ .

**Notation 3.7.** Let  $T \in \Omega$  be a tree and suppose that  $e <_d b$ . We will denote by  $b_e^{\uparrow} \in T$  the predecessor of b in I(e).

 $\textbf{Proposition 3.8.} \ \ \textit{Suppose} \ e, f \ \ \textit{are} \leq_{\textit{d}} \textit{-incomparable edges of} \ T \ \ \textit{and write} \ b = e \vee f. \ \ \textit{Then}$ 

- (a)  $e <_d b$ ,  $f <_d b$  and  $b_e^{\uparrow} \neq b_f^{\uparrow}$ ;
- (b)  $b_e^{\uparrow}, b_f^{\uparrow} \in b^{\uparrow}$ . In fact  $\{b_e^{\uparrow}\} = I(e) \cap b^{\uparrow}, \{b_f^{\uparrow}\} = I(f) \cap b^{\uparrow}$ ;

PLANARIZE DEF

INPUTPATHS PROP

INCOMPNOTOP

INPUTPATH NOT

ECESSORPROP PROP

(c) if 
$$e' \leq_d e$$
,  $f' \leq_d f$  then  $b = e' \vee f'$  and  $b_{e'}^{\uparrow} = b_e^{\uparrow}$ ,  $b_{f'}^{\uparrow} = b_f^{\uparrow}$ .

*Proof.* (a) is immediate: the condition e = g (resp. f = g) would imply  $f \le_d e$  (resp.  $e \le_d f$ ) while the condition  $b_e^{\dagger} = b_f^{\dagger}$  would provide a predecessor of b in  $I(e) \cap I(f)$ .

For (b), note that any relation  $a <_d h$  factors as  $a \le_d b_a^* <_d b$  for some unique  $b_a^* \in b^{\uparrow}$ , where uniqueness follows from Lemma 3.5. Choosing a = e implies  $I(e) \cap b^{\uparrow} = \{b_e^*\}$  and letting a range over edges such that  $e \le_d a <_d b$  shows that  $b_e^*$  is in fact the predecessor of b.

To prove (c) one reduces to the case e' = e, in which case it suffices to check  $I(e) \cap I(f') = I(e) \cap I(f)$ . But if it were otherwise there would exist an edge a satisfying  $f' \leq_d a <_d f$  and  $e \leq_d a$ , and this would imply  $e \leq_d f$ , contradicting our hypothesis.

**Proposition 3.9.** Let  $c = e_1 \lor e_2 \lor e_3$ . Then  $c = e_i \lor e_j$  iff  $c_{e_i}^{\uparrow} \neq c_{e_j}^{\uparrow}$ . Therefore, all ternary joins in  $(T, \leq_d)$  are binary, i.e.

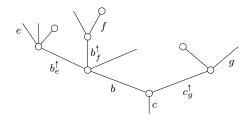
$$c = e_1 \lor e_2 \lor e_3 = e_i \lor e_j \tag{3.10}$$
 TERNJOIN EQ

 $for \ some \ 1 \leq i < j \leq 3, \ and \ (\overline{3.10}) \ \overline{fails} \ \overline{for} \ at \ most \ one \ choice \ of \ 1 \leq i < j \leq 3.$ 

*Proof.* If  $c_{e_i}^{\uparrow} = f_{e_j}^{\uparrow} = f_{e_j}^{\downarrow} = f_$ 

The "therefore" part follows by noting that  $c_{e_1}^{\uparrow}$ ,  $c_{e_2}^{\uparrow}$ ,  $c_{e_3}^{\uparrow}$  can not all coincide, or else c would not be the minimum of  $I(e_1) \cap I(e_2) \cap I(e_3)$ .

**Example 3.11.** In the following example  $b = e \lor f$ ,  $c = e \lor f \lor g$ ,  $c_e^{\uparrow} = c_f^{\uparrow} = b$ .



**Notation 3.12.** Given a set S of size n we write  $Ord(S) \simeq Iso(S, \{1, \dots, n\})$ . We will usually abuse notation by regarding its objects as pairs  $(S, \leq)$  where  $\leq$  is a total order in S.

**Proposition 3.13.** Let  $T \in \Omega$  be a tree. There is a bijection

$$\{planar\ structures\ (T, \leq_p)\} \longrightarrow \prod_{(a^{\uparrow} \leq a) \in V(T)} \mathsf{Ord}(a^{\uparrow})$$

$$\leq_p \longmapsto (\leq_p \mid_{a^{\uparrow}})$$

$$(3.14) \quad \boxed{\mathsf{PLANAR}\ \mathsf{EQ}}$$

Proof. We will keep the setup of Proposition 3.8 throughout: e, f are  $\leq_d$ -incomparable edges and we write  $b = e \vee f$ . BY ANAL PO

and we write  $b = e \lor f$ . PLANAR EQ We first show that (3.14) is injective, i.e. that the restrictions  $\leq_p |_{a}$  determine if  $e <_p f$  holds or not. If  $b_e^{\uparrow} <_p b_f^{\uparrow}$ , the relations  $e \leq_d b_e^{\uparrow} <_p b_f^{\uparrow} \geq_d f$  and Definition 3.1 imply it must be  $e <_p f$ . Dually, if  $b_f^{\uparrow} <_p b_e^{\uparrow}$  then  $f <_p e$ . Thus  $b_e^{\uparrow} <_p b_f^{\downarrow} \Leftrightarrow e <_p f$  and hence (3.14) is indeed injective.

To check that (8.14) is surjective, it suffices (recall that e, f are assumed  $\leq_d$ -incomparable) to check that defining  $e \leq_p f$  to hold iff  $b_e^{\dagger} < b_f^{\dagger}$  holds in  $b^{\dagger}$  yields a planar structure.

Antisymmetry and the total order conditions are immediate, and it thus remains to check the transitivity and planar conditions. Transitivity of  $\leq_p$  in the case  $e'_{\begin{subarray}{c} \begin{subarray}{c} \begin{subarray}{c$ 

TERNARYJOIN PROP

It remains to check transitivity in the hardest case, that of  $e <_p f <_p g$  with remains to check transitivity in the hardest case, that of  $e <_p f <_p g$  with remains incomparable f,g. We write  $c = e \lor f \lor g$ . By the "therefore" part of Proposition 3.9, either (i)  $e \lor f <_d c$ , in which case Proposition 5.9 implies  $c_e^{-} = c_f^{+}$  and transitivity follows; (ii)  $f \lor g <_d c$ , which follows just as (i): (iii)  $e \lor f = f \lor g = c$ , in which case  $c_e^{+} < c_f^{+} < c_g^{+}$  in  $c^{+}$  so that  $c_e^{+} \ne c_g^{+}$  and by Proposition 3.9 it is also  $e \lor g = c$  and transitivity follows. 

PLANARIZE DEF

BENARYJOIN PROP

1.1 The analogue of Proposition 3.1 readily extends to forests  $F \in \Phi$ . The analogue of Proposition

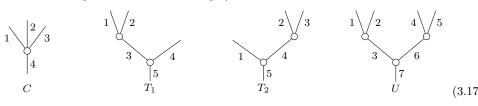
**Remark 3.15.** Definition 3.1 readily extends to forests  $F \in \Phi$ . The analogue of Proposition 3.13 then states that the data of a planar structure is equivalent to total orderings of the nodes of F together with a total ordering of its set of roots. Indeed, this follows by either adapting the proof above or by noting that planar structures on F are clearly in bijection with planar structures on the join tree  $F \star \eta$  (cf. [5, Def. 7.44]), which adds a single edge  $\eta$  to F, serving as the (unique) root of  $F \star \eta$ .

When discussing the substitution procedure in §3.3 we will find it convenient to work with a model for the category  $\Omega$  that possesses exactly one representative of each possible planar structure on each tree or, more precisely, such that the only isomorphisms preserving the planar structures are the identities. On the other hand, using such a model for  $\Omega$  throughout would, among other issues, make the discussion of faces in §3.2 rather awkward. We now outline our conventions to address such issues.

Let  $\Omega^p$ , the category of planarized trees, denote the category with objects pairs  $T_{\leq p} = (T, \leq_p)$  of trees together with a planar structure and morphisms the underlying maps of trees (so that the planar structures are ignored). There is a full subcategory  $\Omega^s \to \Omega^p$ , whose objects we call standard models, of those  $T_{\leq p}$  whose underlying set is one of the sets  $\underline{n} = \{1, 2, \dots, n\}$  and for which  $\leq_p$  coincides with the canonical order.

**Example 3.16.** Some examples of standard models, i.e. objects of  $\Omega^s$ , follow (further, (3.3) can also be interpreted as such an example).

PLANAROMEGAEX1 EQ



Here  $T_1$  and  $T_2$  are isomorphic to each other but not isomorphic to any other standard model in  $\Omega^s$  while both C and U are the unique objects in their isomorphism classes.

Given  $T_{\leq_p} \in \Omega^p$  there is an obvious standard model  $T_{\leq_p}^s \in \Omega^s$  given by replacing each edge by its order following  $\leq_p$ . Indeed, this defines a retraction  $(-)^s : \Omega^p \to \Omega^s$  and a natural transformation  $\sigma : id \Rightarrow (-)^s$  given by isomorphisms preserving the planar structure (in fact, the pair  $((-)^s, \sigma)$  is clearly unique).

**Convention 3.18.** From now on, we will write simply  $\Omega$ ,  $\Omega_G$  to denote the categories  $\Omega^s$ ,  $\Omega_G^s$  of standard models (where planar structures are defined in the underlying forest as in Remark 5.15). Similarly  $O_G$  will denote the model  $O_G^s$  for the orbital category whose objects are the orbital G-sets whose underlying set is one of the sets  $\underline{n} = \{1, 2, \dots, n\}$ .

Therefore, whenever one of our constructions produces an object/diagram in  $\Omega^p$ ,  $\Omega^p_G$ ,  $\mathsf{O}^p_G$  (of trees, G-trees, orbital G-sets with a planarization/total order) we will hence implicitly reinterpret it by using the standardization functor  $(-)^s$ .

Example 3.19. To illustrate our convention, we consider the trees in Example 3.16.

One has subfaces  $F_1 \subset F_2 \subset U$  where  $F_1$  is the subtree with edge set  $\{1, 2, 6, 7\}$  and  $F_2$  is the subtree with edge set  $\{1, 2, 3, 6, 7\}$ , both with inherited tree and planar structures. Applying  $(-)^s$  to the inclusion diagram on the left below then yields a diagram as on the right.



STANDMODEL EX

PLANARCONV CON

Similarly, let  $\leq_{(12)}$  and  $\leq_{(45)}$  denote alternate planar structures for U exchanging the orders of the pairs 1,2 and 4,5, so that one has objects  $U_{\leq_{(12)}}$ ,  $U_{\leq_{(45)}}$  in  $\Omega^p$ . Applying (-)<sup>s</sup> to the diagram of underlying identities on the left yields the permutation diagram on the right.

$$U \xrightarrow{id} U_{\leq_{(45)}} \qquad \qquad U \xrightarrow{(45)} U$$

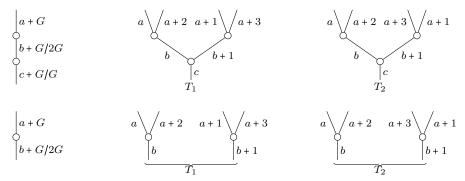
$$U \xrightarrow{id} U_{\leq_{(12)}} U$$

$$U \xrightarrow{(12)} U$$

$$U \xrightarrow{(12)(45)} U$$

**Example 3.20.** An additional reason to leave the use of  $(-)^s$  implicit is that when depicting G-trees it is preferable to choose edge labels that describe the action rather than the planarization (which is already implicit anyway).

For example, when  $G = \mathbb{Z}_{/4}$ , in both diagrams below the orbital representation on the left represents the isomorphism class consisting of the two trees  $T_1, T_2 \in \Omega_G$  on the right.



**Definition 3.21.** A morphism  $S \xrightarrow{\varphi} T$  in  $\Omega$  that is compatible with the planar structures  $\leq_p$  is called a *planar map*.

More generally, a morphism  $F \to G$  in the categories  $\Phi$ ,  $\Phi^G = \Omega^G_{Pe16b}$  of forests, G-forests, G-trees is called a *planar map* if it is an independent map (cf. [5, Def. 5.28]) compatible with the planar structures  $\leq_p$ .

Remark 3.22. The need for the independence condition is justified by [5, Lemma 5.33] and its converse, since non independent maps do not reflect  $\leq_d$  inequalities.

We note that in the  $\Omega_G$  case a map  $\varphi$  is independent iff  $\varphi$  does not factor through a non trivial quotient iff  $\varphi$  is injective on each edge orbit.

**Proposition 3.23.** Let  $F \xrightarrow{\varphi} G$  be an independent map in  $\Phi$  (or  $\Omega$ ,  $\Omega_G$ ,  $\Phi_G$ ). Then there is a unique factorization

$$F \xrightarrow{\simeq} \bar{F} \to G$$

such that  $F \xrightarrow{\simeq} \bar{F}$  is an isomorphism and  $\bar{F} \to G$  is planar.

Proof. We need to show that there is a unique planar structure  $\leq_p^{\overline{F}}$  on the underlying forest of F making the underlying map a planar map. Simplicity of G ensures that for any vertex  $e^{\uparrow} \leq e$  of F the edges in  $\varphi(e^{\uparrow})$  are all distinct while independence of  $\varphi$  likewise ensures that the edges in  $\varphi(e^{\uparrow})$  are all distinct while independence of  $\varphi$  likewise ensures that the edges in  $\varphi(e^{\uparrow})$  are all distinct. The result now follows from (the forest version of) Proposition 3.13: one simply orders each set  $e^{\uparrow}$  and  $\underline{r}_F$  according to its image.

not quite complete... maybe that  $\leq_p$  is the closure of  $\leq_d$  and the vertex relations under transitivity and the planar condition

**Remark 3.24.** Proposition 3.23 says that planar structures can be pulled back along independent many. However, they can not always be pushed forward. As an example, in the notation of (3.17), consider the map  $C \to T_1$  defined by  $1 \mapsto 1$ ,  $2 \mapsto 4$ ,  $3 \mapsto 2$ ,  $4 \mapsto 5$ .

**Remark 3.25.** Given any tree  $T \in \Omega$  there is a unique corolla  $lr(T) \in \Sigma$  and planar tall map  $lr(T) \to T$ . Explicitly, the number of leaves of lr(T) matches that of T, together with the inherited order.

# PULLPLANAR REM

## UNIQCOR REM

OUTTALL SEC

#### 3.2 Outer faces and tall maps

In preparation for our discussion of the substitution operation in §3.3, we now recall some basic notions and results concerning outer subtrees and tree grafting, as in [5, §5].

**Definition 3.26.** Let  $T \in \Omega$  be a tree and  $e_1 \cdots e_n = \underline{e} \le e$  a broad relation in T.

We define the planar outer face  $T_{\underline{e} \leq e}$  to be the subtree with underlying set those edges  $f \in T$  such that

$$f \le_d e, \quad \forall_i e_i \not <_d f, \tag{3.27}$$

OUTERFACE EQ

generating broad relations the relations  $f^{\uparrow} \leq f$  for f satisfying (3.27) and  $\forall i f \neq e_i$ , and planar structure pulled back from T.

**Remark 3.28.** If one forgoes the requirement that  $T_{e \le e}$  be equipped with the pullback planar structure, the inclusion  $T_{\underline{e} \leq e} \to T$  is usually called simply an outer face.

We now recap some basic results.

**Proposition 3.29.** Let  $T \in \Omega$  be a tree.

- (a)  $T_{\underline{e} \leq e}$  is a tree with root e and edge tuple  $\underline{e}$ ;
- (b) there is a bijection

 $\{planar \ outer \ faces \ of \ T\} \leftrightarrow \{broad \ relations \ of \ T\};$ 

- (c) if  $R \to S$  and  $S \to T$  are outer face maps then so is  $R \to T$ ;
- (d) any pair of broad relations  $g \le v$ ,  $fv \le e$  induces a grafting pushout diagram

*Proof.* We first show (a). That  $T_{\underline{e} \leq e}$  is indeed a tree is the content of [5, Prop. 5.20]: more precisely,  $T_{\underline{e} \leq e} = (T^{\leq e})_{\leq \underline{e}}$  in the notation therein. That the root of  $T_{\underline{e} \leq e}$  is e is clear and that the root tuple is  $\underline{e}$  follows from [5, Remark 5.23].

- (b) follows from (a), which shows that  $\underline{e} \leq e$  can be recovered from  $T_{\underline{e} \leq e}$ . (c) follows from the definition of outer face together with [5, Lemma 5.33], which states that the  $\leq_d$  relations on S, T coincide.

Since by (c) both  $T_{\underline{g} \leq v}$  and  $T_{\underline{f}v \leq e}$  are outer faces of  $T_{\underline{f}\underline{g} \leq v}$ , (d) is a restatement of [5,Prop. 5.15].

**Definition 3.31.** A map  $S \xrightarrow{\varphi} T$  in  $\Omega$  is called a *tall map* if

$$\varphi(\underline{l}_S) = \underline{l}_T, \qquad \varphi(r_S) = r_T,$$

where  $l_{(-)}$  denotes the leaf tuple and  $r_{(-)}$  the root. The following is a restatement of [5, Cor. 5.24]

**Proposition 3.32.** Any map  $S \xrightarrow{\varphi} T$  in  $\Omega$  has a factorization, unique up to unique isomorphism.

$$S \xrightarrow{\varphi^t} U \xrightarrow{\varphi^u} T$$

as a tall map followed by an outer face (in fact,  $U = T_{\varphi(l_S) \leq r_S}$ ).

We recall that a face  $F \to T$  is called inner if is obtained by iteratively removing inner edges, i.e. edges other than the root or the leaves. In particular, it follows that a face is inner iff it is tall. The usual face-degeneracy decomposition thus combines with Corollary 3.32 to give the following. **Corollary 3.33.** Any map  $S \xrightarrow{\varphi} T$  in  $\Omega$  has a factorization, unique up to unique isomorphisms,

$$S \xrightarrow{\varphi^{-}} U \xrightarrow{\varphi^{i}} V \xrightarrow{\varphi^{u}} T \tag{3.34}$$
 TRIPLEFACT EQ

as a degeneracy followed by an inner face followed by an outer face.

*Proof.* The factorization (3.34) can be built by first performing the degeneracy-face decomposition and then performing the tall-outer decomposition on the face map.

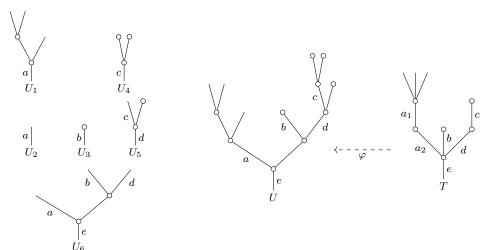
SUBS SEC

#### 3.3 Substitution

One of the key ideas needed to describe operads is that of substitution of tree nodes, a process that we will prefer to repackage in terms of maps of trees. We start by discussing appush EQ example, focusing on the related notion of iterated graftings of trees (as described in (3.30)).

**Example 3.35.** The trees  $U_1, U_2, \dots, U_6$  on the left below can be grafted into the tree U in the middle. More precisely (among other possible grafting orders), one has

$$U = ((((((U_6 \coprod_a U_2)) \coprod_a U_1) \coprod_b U_3) \coprod_d U_5) \coprod_c U_4$$
 (3.36) UFORMULA EQ

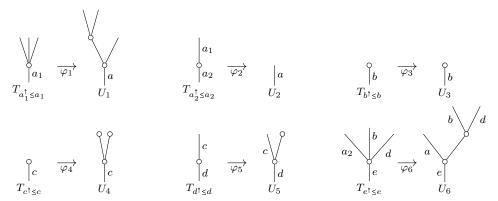


(3.37)

SUBSDATUMTREES EQ

We now consider the tree T, which is built by converting each  $U_i$  into the corollar  $U_i$  (cf. Remark 3.25), and then performing the same grafting operations as in (3.36) at then be regarded as encoding the combinatorics of the iterated grafting in (3.36), with alternative ways to reorder operations in (3.36) in bijection with ways to assemble T out of its nodes. One can now therefore think of the iterated grafting (3.36) as being instead encoded by

the tree T together with the (unique) planar tall maps  $\varphi_i$  below.



(3.38) SUBSDATUMTREES2 EQ

From this perspective, U can now be thought as obtained from T by substituting each of its nodes with the corresponding  $U_i$ . Moreover, the  $\varphi_i$  assemble to a planar tall map  $\varphi: T \to U$  (such that  $a_i \mapsto a, b \mapsto b, \dots, e \mapsto e$ ), which likewise encodes the same information.

Our perspective will then be that data for substitution of tree nodes such as in (3.38) can equivalently be repackaged using planar tall maps.

**Definition 3.39.** Let  $T \in \Omega$  be a tree.

A T-substitution datum is a tuple  $\{U_{e^{\dagger} \leq e}\}_{(e^{\dagger} \leq e) \in V(T)}$  together with tall maps  $T_{e^{\dagger} \leq e} \rightarrow U_{e^{\dagger} \leq e}$ .

Further, a map of planar T-substitution data  $\{U_{e^{\dagger} \leq e}\} \to \{V_{e^{\dagger} \leq e}\}$  is a tuple of tall maps  $\{U_{e^{\dagger} \leq e} \to V_{e^{\dagger} \leq e}\}$  compatible with the chosen maps.

Lastly, a substitution datum is called a *planar T-substitution datum* if the chosen maps are planar (so that  $lr(U_{e^{\uparrow} \le e}) = T_{e^{\uparrow} \le e}$ ) and a morphism of planar data is called a planar morphism if it consists of a tuple of planar maps.

**Definition 3.40.** Let  $T \in \Omega$ .

The Segal core poset Sc(T) is the poset with objects the edge subtrees  $\eta_e$  and vertex substrees  $T_{e^{\uparrow} \leq e}$ . The order relation is given by inclusion.

**Remark 3.41.** Note that the only maps in Sc(T) are inclusions of the form  $\eta_a \subset T_{e^{\dagger} \leq e}$ . In particular, there are no pairs of composable non-identity relations in Sc(T).

Given a T-substitution datum  $\{U_{\{e^{\uparrow} \leq e\}}\}$  we abuse notation by writing

$$U_{(-)}$$
:  $Sc(T) \to \Omega$ 

for the functor  $\eta_a \mapsto \eta$ ,  $T_{e^{\uparrow} \leq e} \mapsto U_{e^{\uparrow} \leq e}$  and sending the inclusions  $\eta_a \subset T_{e^{\uparrow} \leq e}$  to the composites

$$\eta \xrightarrow{a} T_{e^{\uparrow} \leq e} \to U_{e^{\uparrow} \leq e}.$$

**Proposition 3.42.** Let  $T \in \Omega$  be a tree. There is an isomorphism of categories

$$\begin{aligned} \operatorname{Sub}_p(T) & & \longleftarrow & \Omega^{\operatorname{pt}}_{T/} \\ & \{U_{e^{\dagger} \leq e}\} & \longmapsto \left(T \to \operatorname{colim}_{\operatorname{Sc}(T)} U_{(-)}\right) \end{aligned} \tag{3.43}$$

where  $\operatorname{Sub}_p(T)$  denotes the category of planar T-substitution data and  $\Omega_{T/}^{\operatorname{pt}}$  the category of planar tall maps under T.

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*Proof.* We first claim that (i) the  $\operatorname{colim}_{Sc(T)}U_{(-)}$  indeed exists; (ii) for the canonical datum  $\{T_{e^{\dagger} \leq e}\}\$ , it is  $T = \text{colim}_{\mathsf{Sc}(T)} T_{(-)}$ ; (iii) the induced map  $T \to \text{colim}_{\mathsf{Sc}(T)} U_{(-)}$  is planar tall.

The argument is by induction on the number of vertices of T, with the base cases of T with 0 or 1 vertices being immediate, since then T is the terminal object of Sc(T). Otherwise, one can choose a non trivial grafting decomposition so as to write  $T = R \coprod_{e} S$ , resulting in identifications  $Sc(R) \subset Sc(T)$ ,  $Sc(S) \subset Sc(T)$  so that  $Sc(R) \cup Sc(S) = Sc(T)$  and  $Sc(R) \cap Sc(S) = \{\eta_e\}$ . The existence of  $colim_{Sc(T)} U_{(-)}$  is thus equivalent to the existence of the pushout below.

$$\eta \longrightarrow \operatorname{colim}_{\operatorname{Sc}(R)} U_{(-)}$$
 
$$\downarrow \qquad \qquad \downarrow$$
 
$$\operatorname{colim}_{\operatorname{Sc}(S)} U_{(-)} ---- + \operatorname{colim}_{\operatorname{Sc}(T)} U_{(-)}$$
 
$$(3.44) \quad \boxed{ \operatorname{ASSEMBLYGRAFT EQ} }$$

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By induction, the top right and bottom left colimits exist for any  $U_{(-)}$ , equal R and S in the case  $U_{(-)} = T_{(-)}$ , and the maps  $R \to \operatorname{colim}_{\mathsf{Sc}(R)} U_{(-)}$ ,  $S \to \operatorname{colim}_{\mathsf{Sc}(S)} U_{(-)}$  are planar tall. But is now follows that (3.44) is a grafting pushout diagram, so that the pushout indeed exists.

The conditions that  $T = \operatorname{colim}_{\mathsf{Sc}(T)} T_{(-)}$  and  $T \to \operatorname{colim}_{\mathsf{Sc}(T)} U_{(-)}$  is planar tall follow. The fact that the two functors in (3.43) are inverse to each other is clear by the same inductive argument.

Corollary 3.45. Let  $T \in \Omega$  be a tree. There is an isomorphism of categories

$$\mathsf{Sub}(T) \ \Longleftrightarrow \Omega_{T/}^{\mathsf{t}} \tag{3.46}$$

where Sub(T) denotes the category of T-substitution data and  $\Omega_{T/}^{t}$  the category of tall maps under T.

Proof. This is a consequence of Proposition 3.23 together with the previous result with the Proposition 3.23 together with the previous result with the Proposition 3.23 together with the previous result with the previous result with the previous result. functor  $Sub(T) \to \Omega_{T/T}^t$  given by the same formula. Indeed, Proposition 3.13 can be restated as saying that isomorphisms  $T \to T'$  are in bijection with substitution data consisting of isomorphisms, and thus bijectiveness reduces to that in the previous result.

**Remark 3.47.** It follows from the previous proof that, writing  $U = \operatorname{colim}_{Sc(T)} U_{(-)}$ , one has

$$V(U) = \coprod_{(e^{\dagger} \le e) \in V(T)} V(U_{e^{\dagger} \le e}). \tag{3.48}$$

Alternatively, (3.48) can be regarded as a map  $f^*: V(U) \to V(T)$  induced by the planar tall map  $f:T \to U$ . Explicitly,  $f^*(U_{u^{\uparrow} \le u})$  is the unique  $T_{t^{\uparrow} \le t}$  such that  $U_{u^{\uparrow} \le u} \subset U_{t^{\uparrow} \le t}$ . We note that  $f^*$  is indeed contravariant in the tall planar map f.

The following is a converse of sorts to Proposition 3.42.

**Proposition 3.49.** Let  $U \in \Omega$  be a tree. Then:

- (i) given non stick outer subtrees  $U_i$  such that  $V(U) = \coprod_i V(U_i)$  there is a unique tree Tand planar tall map  $T \to U$  such that  $\{U_i\} = \{U_{e^{\uparrow} \leq e}\};$
- (ii) given multiplicities  $m_e \ge 1$  for each edge  $e \in U$ , there is a unique planar degeneracy  $\rho: T \to U \text{ such that } \rho^{-1}(e) \text{ has } m_e \text{ elements};$
- (iii) planar tall maps  $T \to U$  are in bijection with collections  $\{U_i\}$  of outer subtrees such that  $V(U) = \coprod_i V(U_i)$  and  $U_j$  is not an inner edge of any  $U_i$  whenever  $U_j \simeq \eta$  is a stick.

*Proof.* We first show (i) by induction on the number of subtrees  $U_i$ . The base case  $\{U_i\}$  =  $\{U\}$  is immediate, setting  $T = \operatorname{Ir}(U)$ . Otherwise, letting e be edge that is both an inner edge of U and a root of some  $U_i$ , and one can form a pushout diagram

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inducing a nontrivial partition  $\{U_i\} = \{U_i|U_i \hookrightarrow V\} \sqcup \{U_i|U_i \hookrightarrow W\}$ . Existence of  $T \to U$  now follows from the induction hypothesis. For uniqueness, the condition that no  $U_i$  is a stick guarantees that T possesses a single inner edge mapping to e, and thus admits a compatible decomposition as in (3.50), and thus uniqueness too follows by the induction hypothesis.

For (ii), we argue existence by nested induction on the number of vertices |V(U)| and the sum of the multiplicities  $m_e$ . The base case |V(U)| = 0 is  $U_{pECOMPRO} U_{pRO} U_{p$ 

(iii) follows by combining (i) and (ii). Indeed, any planar tall map  $T \to U$  has a unique decomposition  $T \twoheadrightarrow \bar{T} \hookrightarrow U$  as a planar degeneracy followed by a planar inner face, and each of these maps is classified by the data in (b) and (a).

**Lemma 3.51.** Suppose  $T_1, T_2 \hookrightarrow T$  are two outer faces with at least one common edge e. Then there exists an unique outer face  $T_1 \cup T_2$  such that  $V(T_1 \cup T_2) = V(T_1) \cup V(T_2)$ .

*Proof.* If either of  $T_1, T_2$  is the root or a leaf the result is obvious. Otherwise, one can necessarily choose e to be an inner edge of T, in which case all of  $T_1, T_2, T$  admit compatible decompositions (3.50) and the result follows by induction on |V(T)|.

# 4 The genuine equivariant operad monad

We now turn to the task of building the monad encoding genuine equivariant operads.

# 4.1 Wreath product over finite sets

In what follows we will let F denote the usual skeleton of the category of finite sets and all set maps. Explicitly, its objects are the finite sets  $\{1,2,\cdots,n\}$  for  $n\geq 0$ . However, much as in the discussion in Convention 3.18 we will often find it more convenient to regard the elements of F as equivalence classes of finite sets equipped with total orders.

**Definition 4.1.** For a category C, we let  $F \wr C$  denote the opposite of the Grothendieck construction for the functor

$$F^{op} \longrightarrow \mathsf{Cat}$$
 $I \longmapsto \mathcal{C}^I$ 

Explicitly, the objects of  $F \wr C$  are tuples  $(c_i)_{i \in I}$  and a map  $(c_i)_{i \in I} \to (d_j)_{j \in J}$  consists of a pair

$$(\phi: I \to J, (f_i: c_i \to d_{\phi(i)})_{i \in I}),$$

henceforth abbreviated as  $(\phi, (f_i))$ .

The following is immediate.

**Proposition 4.2.** Suppose C has all finite coproducts. One then has a functor as on the left below. Dually, if C has all finite products, one has a functor as on the right below.

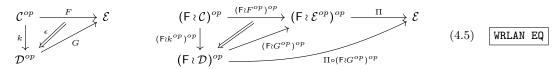
$$\begin{array}{cccc}
\mathsf{F} \wr \mathcal{C} & \xrightarrow{\coprod} & \mathcal{C} & (\mathsf{F} \wr \mathcal{C}^{op})^{op} & \xrightarrow{\Pi} & \mathcal{C} \\
(c_i)_{i \in I} & \longmapsto \coprod_{i \in I} c_i & (c_i)_{i \in I} & \longmapsto \prod_{i \in I} c_i
\end{array}$$

**Lemma 4.3.** Suppose that  $\mathcal{E}$  is a bicomplete category such that coproducts commute with limits in each variable. If the leftmost diagram

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is a right Kan extension diagram then so is the composite of the rightmost diagram. Dually, if in  $\mathcal{E}$  products commute with colimits in each variable, and the leftmost diagram



is a left Kan extension diagram then so is the composite of the rightmost diagram.

*Proof.* Unpacking definitions using the pointwise formula for Kan extensions ([4, X.3.1]), the claim concerning ([4.4]) amounts to showing that for each  $(d_i) \in F \wr \mathcal{D}$  one has natural isomorphisms

$$\lim_{((d_i)\to(kc_j))\in((d_i)\downarrow F\wr C)} \left(\coprod_i F(c_j)\right) \simeq \coprod_i \lim_{(d_i\to kc_i)\in d_i\downarrow C} \left(F(c_i)\right). \tag{4.6}$$

Noting that the canonical factorizations of each  $(\varphi, (f_i)): (d_i)_{i \in I} \to (kc_j)_{j \in J}$  as

$$(d_i)_{i \in I} \to (c_{\phi(i)})_{i \in I} \to (kc_j)_{j \in J}$$

exhibit  $\prod_i (d_i \downarrow \mathcal{C})$  as a coreflexive subcategory of  $(d_i) \downarrow \mathsf{F} \wr \mathcal{C}$ , we see that it is an initial subcategory. Therefore

$$\lim_{((d_i) \to (kc_j)) \in ((d_i) \downarrow \mathsf{FiC})} \left( \coprod_j F(c_j) \right) \simeq \lim_{((d_i) \to (kc_i)) \in \Pi_i(d_i \downarrow \mathcal{D})} \left( \coprod_i F(c_i) \right)$$

and hence (A.6) now follows from the assumption that coproducts commute with limits in each variable.

**Notation 4.7.** Using the coproduct functor  $\mathsf{F}^{\wr 2} = \mathsf{F}^{\wr \{0,1\}} = \mathsf{F} \wr \mathsf{F} \xrightarrow{\sqcup} \mathsf{F}$  (where  $\coprod_{i \in I} J_i$  is ordered lexicographically) and the simpleton  $\{1\} \in \mathsf{F}$  one can regard the collection of categories  $\mathsf{F}^{\wr \{0,\cdots,n\}} \wr \mathcal{C} = \mathsf{F}^{\underline{\imath n}} \wr \mathcal{C}$  as a coaugmented cosimplicial object in Cat. As such, we will denote by

$$\delta^{i} : \mathsf{F}^{2n-1} : \mathcal{C} \to \mathsf{F}^{2n} : \mathcal{C}, \qquad 0 \le i \le n$$

the cofaces obtained by inserting simpletons  $\{1\} \in \mathsf{F}$  and by

$$\sigma^{i} : \mathsf{F}^{i_{n+1}} \wr \mathcal{C} \to \mathsf{F}^{i_{n}} \wr \mathcal{C}, \qquad 0 \le i \le n$$

the codegeneracies obtained by applying the coproduct  $F^{i2} \xrightarrow{\coprod} F$  to adjacent F coordinates.

# 4.2 Equivariant leaf-root and vertex functors

**Definition 4.8.** A morphism  $T \xrightarrow{\varphi} S$  in  $\Omega_G$  is called a *quotient* if the underlying morphism of forests

$$\coprod_{[g] \in G/H} T_{[g]} \to \coprod_{[h] \in G/K} S_{[h]}$$

maps each tree component (or, equivalently, some tree component) isomorphically onto its image component.

We denote the subcategory of G-trees and quotients by  $\Omega_G^q$ .

**Definition 4.9.** The *G*-symmetric category, which we will also call the category of *G*-corollas, is the full subcategory  $\Sigma_G \subset \Omega_G^q$  of those *G*-trees that are corollas, i.e. *G*-trees such that each edge is either a root or a leaf (but not both).

**Definition 4.10.** The *leaf-root functor* is the functor  $\Omega_G^q \xrightarrow{lr} \Sigma_G$  defined by

$$lr(T) = \{leaves of T\} \coprod \{roots of T\}$$

with a broad relation  $l_1 \cdots l_n \leq r$  holding in lr(T) iff its image holds in T and similarly for the planar structure  $\leq_p$ .

**Remark 4.11.** Generalizing Remark 3.25,  $\Gamma(T)$  can alternatively be characterized as being the *unique G*-corolla which admits an also unique (tree-wise) tall planar map  $\Gamma(T) \to T$ . Moreover,  $\Gamma(T)$  can usually be regarded as the "smallest inner face" of T, obtained by removing all the inner edges, although this characterization fails when  $T = G \cdot_H \eta$  is a stick G-tree. Some examples with  $G = \mathbb{Z}_{/4}$  follow.

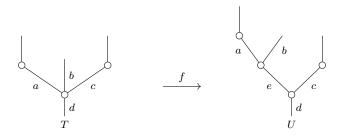
**Remark 4.12.** One consequence of the fact that planarizations can not be pushed forward along tree maps (cf. Remark 3.24) is that  $\operatorname{Ir}:\Omega_G^q \to \Sigma_G$  is not a categorical fibration. maybe add to this.

**Definition 4.13.** Given  $T \in \Omega_G$  we define the set  $V_G(T)$  of *G-vertices* of T to be the orbit set V(T)/G, i.e. the quotient of the vertex set V(T) by its G-action.

Furthermore, we will regard  $V_G(T)$  as an object in F by equipping it with its lexicographic order: i.e. vertex equivalence classes  $[e^{\uparrow} \leq e]$  are ordered according to the planar order  $\leq_p$  of the smallest representative ge,  $g \in G$ .

Remark 4.14. Following Remark 3.47, a planar tall map  $f:T \to U$  of G-trees induces a G-equivariant map  $f^*:V(U) \to V(T)$  and thus also a map of orbits  $f^*:V_G(U) \to V_G(T)$ . We note, however, that  $f^*$  is not in general compatible with the order on  $V_G$ , as is indeed the case even in the non-equivariant case.

A minimal example follows.



In V(T) the vertices are ordered as a < c < d while in V(U) they are ordered as a < e < c < d but the map  $f^*: V(U) \to V(T)$  is given by  $a \mapsto a, c \mapsto c, d \mapsto d, e \mapsto d$ .

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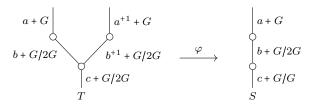
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Note that each element of  $V_G(T)$  corresponds to an unique edge orbit Ge for e not a leaf. As such, we will represent the corresponding G-vertex by  $v_{Ge} = (Ge)^{\dagger} \leq Ge$  (which we interpret as the concatenation of the relations  $f^{\dagger} \leq f$  for  $f \in Ge$ ) and write

$$T_{v_{Ge}} = T_{(Ge)^{\uparrow} \leq Ge} = \coprod_{f \in Ge} T_{f^{\uparrow} \leq f}.$$

We note that  $T_{v_{Ge}}$  is always a G-corolla. Indeed, noting that a quotient map  $\varphi: T \to S$  induces quotient maps  $T_{v_{ge}} \to S_{v_{G\varphi(e)}}$  one obtains a functor

**Remark 4.16.** The need to introduce the  $F \wr C$  categories comes from the fact that general quotient maps do not preserve the number of G-vertices. For a simple example, let  $G = \mathbb{Z}_{/4}$  and consider the quotient map



sending edges labeled a, b, c to the edges with the same name and the edges  $a^{+1}$ ,  $b^{+1}$  to the edges a+1, b+1. We note that T has three G-vertices  $v_{Gc}$ ,  $v_{Gb}$ ,  $v_{Gb+1}$  while S has only two G-vertices  $v_{Gc}$  and  $v_{Gb}$ .  $V(\phi)$  then maps the two corollas  $T_{v_{Gb}}$  and  $T_{v_{Gb+1}}$  isomorphically onto  $T_{S_{Gb}}$  and the corolla  $T_{v_{Gc}}$  non-isomorphically onto  $S_{v_{Gc}}$ .

onto  $T_{S_{Gb}}$  and the corolla  $T_{v_{Gc}}$  non-isomorphically onto  $S_{v_{Gc}}$ .

SubstitutionDatum

Definition 3.39 now immediately generalizes. Here a map is called *rooted* if it induces an ordered isomorphism on the root orbit.

**Definition 4.17.** Let  $T \in \Omega_G$  be a G-tree.

A rooted (resp. planar) T-substitution datum is a tuple  $\{U_{v_{Ge}}\}_{v_{Ge} \in V_G(T)}$  together with rooted (resp. planar) tall maps  $T_{v_{Ge}} \to U_{v_{Ge}} = T_{v_{Ge}}$ .

rooted (resp. planar) tall maps  $T_{v_{Ge}} \to U_{v_{Ge}} = T_{v_{Ge}}$ . Further, a map of rooted (resp. planar) T-substitution data  $\{U_{v_{Ge}}\} \to \{V_{v_{Ge}}\}$  is a tuple of rooted (resp. planar) tall maps  $\{U_{v_{Ge}} \to V_{v_{Ge}}\}$ .

Remark 4.18. To establish the equivariant analogue of Proposition 3.42 we will prefer to repackage equivariant substitution data in terms of non-equivariant terms.

Noting that there are decompositions  $U_{v_{Ge}} = \coprod_{ge \in Ge} U_{ge^{\dagger} \leq ge}$  and letting  $G \ltimes V(T)$  denote the Grothendieck construction for the action of G on the non-equivariant vertices V(T) (often called the action groupoid), it is immediate that an equivariant T-substitution datum is the same as a functor  $G \ltimes V(T) \to \Omega$  whose restriction to  $V(T) \subset G \ltimes V(T)$  is a (non-equivariant) substitution datum.

**Proposition 4.19.** Let  $T \in \Omega_G$  be a G-tree. There are isomorphisms of categories

$$\begin{aligned} \mathsf{Sub}_{\mathsf{p}}(T) & & \longleftarrow & \Omega^{\mathsf{pt}}_{G,T/} \\ \{U_{v_{Ge}}\} & & \longleftarrow & \left(T \to \mathrm{colim}_{\mathsf{Sc}(T)}\,U_{(-)}\right) \end{aligned} \qquad \begin{aligned} & \mathsf{Sub}_{\mathsf{r}}(T) & & \longleftarrow & \Omega^{\mathsf{rt}}_{G,T/} \\ \{U_{v_{Ge}}\} & & \longleftarrow & \left(T \to \mathrm{colim}_{\mathsf{Sc}(T)}\,U_{(-)}\right) \end{aligned}$$

(4.20) SUBDATAUNDERPLANG EQ

Proof. This is a minor adaptation of the non-equivariant analogues Proposition 4.19 and Corollary 3.45. Since Sc(T) inherits a G action, one can form the Grothendieck construction  $G \times Sc(T)$  and by Remark 4.18 equivariant substitution data  $\{U_{v_{Ge}}\}$  therefore induce functors

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 $U_{(-)}: G \ltimes \mathsf{Sc}(T) \to \Omega$ . It is then immediate that  $\mathsf{colim}_{\mathsf{Sc}(T)} U_{(-)}$  inherits a G-action, provided it exists. The key observation is then that, since  $\mathsf{Sc}(T)$  is now a disconnected poset, this colimit is to be interpreted as taken in the category  $\Phi$  of forests rather than in  $\Omega$ .

Additionally, we note that the need to use rooted data comes from the fact that rooted isomorphisms  $T \to T'$  are in bijection with rooted substitution data that are given by isomorphisms, a statement that fails in the absence of the rooted condition.

Remark 4.21. We will need to know that in the planar case each of the maps

$$U_{v_{Ge}} \rightarrow U = \operatorname{colim}_{\mathsf{Sc}(T)} U_{(-)}$$

induced by the previous proof is a planar map of G-trees. This requires two observations: (i) the restrictions to each of the constituent non-equivariant trees  $U_{ge^{\dagger} \leq ge}$  is planar by Proposition 4.19; (ii) the restriction to the roots of  $U_{v_{Ge}}$  is injective and order preserving since it matches the inclusion of the roots of  $T_{v_{Ge}}$ , and the map  $T \to U$  is a planar map of G-trees.

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Remark 4.22. The isomorphisms in Proposition 4.19 are compatible with root pullback of trees. More concretely, any pullback  $\pi: S = \varphi^*T \to T$  induces pullbacks  $\pi_{Ge}: S_{v_{Ge}} \to T_{v_{Ge}}$  for  $v_{Ge} \in V_G(S)$  and one has commutative diagrams

$$\begin{array}{lll} \operatorname{Sub}_{\mathsf{p}}(S) & \longrightarrow \Omega^{\mathsf{pt}}_{G,S/} & \operatorname{Sub}_{\mathsf{r}}(S) & \longrightarrow \Omega^{\mathsf{rt}}_{G,S/} \\ (\pi_{Ge}) & \uparrow & \uparrow_{\pi^*} & (\pi_{Ge}) & \uparrow_{\pi^*} & \\ \operatorname{Sub}_{\mathsf{p}}(T) & \longmapsto \Omega^{\mathsf{pt}}_{G,T/} & \operatorname{Sub}_{\mathsf{r}}(T) & \longmapsto \Omega^{\mathsf{rt}}_{G,T/} \end{array} \tag{4.23}$$

SUBDATAUNDERPLANG2 EQ

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# 4.3 Planar strings

The leaf-root and vertex functors will allow us to reinterpret our results concerning substitution.

**Definition 4.24.** The category  $\Omega_{G,n}$  of substitution n-strings is the category whose objects are strings

$$T_0 \xrightarrow{f_1} T_1 \xrightarrow{f_2} \cdots \xrightarrow{f_n} T_n$$

where  $T_i \in \Omega_G$  and the  $f_i$  are tall planar maps, and arrows are commutative diagrams

where each  $q_i$  is a quotient map.

**Notation 4.26.** Since compositions of planar tall arrows are planar tall and identity arrows are planar tall it follows that  $\Omega_{G,\bullet}$  forms a simplicial object in Cat, with faces given by composing and degeneracies by inserting identities.

Noting that  $\Omega_{G,0} = \Omega_G^q$  and setting  $\Omega_{G,-1} = \Sigma_G$ , the leaf-root functor  $\Omega_G^q \xrightarrow{\operatorname{lr}} \Sigma_G$  makes  $\Omega_{G,\bullet}^q$  into an augmented simplicial object and, furthermore, the maps  $s_{-1} \colon \Omega_{G,n}^q \to \Omega_{G,n+1}^q$  sending  $T_0 \to T_1 \to \cdots \to T_n$  to  $\operatorname{lr}(T_0) \to T_0 \to T_1 \to \cdots \to T_n$  equip it with extra degeneracies.

**Notation 4.27.** We extend the vertex functor to a functor  $V_G: \Omega_{G,n+1} \to \mathsf{F} \wr \Omega_{G,n}$  by

$$V_G(T_0 \to T_1 \to \cdots \to T_n) = (T_{1,v_{Ge}} \to \cdots \to T_{n,v_{Ge}})_{v_{Ge} \in V_G(T_0)}$$
(4.28) VGDEF EQ

where we abuse notation by writing  $T_{i,v_{Ge}}$  for  $T_{i,(f_i \circ \cdots \circ f_1)(v_{Ge})}$ .

Subdatadvaderplang properties a reinterpretation of Proposition 4.19.

SUBSASPULL PROP

Proposition 4.29. The diagram

$$\begin{array}{ccc} \Omega_{G,n+1} & \xrightarrow{V_G} \mathsf{F} \wr \Omega_{G,n} \\ & & \downarrow^{\mathsf{F} \wr d_{0,\cdots,n}} & & \downarrow^{\mathsf{F} \wr d_{0,\cdots,n}} \\ & & \Omega_{G,0} & \xrightarrow{V_G} \mathsf{F} \wr \Sigma_G & & & \end{array} \tag{4.30}$$

is a pullback diagram in Cat.

*Proof.* An object in the pullback ( $\overline{A.30}$ ) object in the pullback ( $\overline{A.30}$ ) object  $\overline{A.70}$   $\Omega_{\overline{PROP}}^{\alpha}$   $\Omega_{\overline{PROP}}^{\alpha}$  is precisely the same as a n-string in Sub(T), and thus by Proposition  $\overline{A.19}$  equivalent to a n+1 planar tall string starting at

The case of arrows is slightly more subtle. A quotient map  $\pi:T\to T'$  induces a Gequivariant poset map  $\pi_*: Sc(T) \to Sc(T')$  (or equivalently, a map of Grothendieck constructions  $G \ltimes Sc(T) \to G \ltimes Sc(T')$  and diagrams as on the left below (where  $v_{Ge}$  ranges over  $V_G(T)$  and  $e' = \varphi(e)$  induce diagrams (of functors  $\mathsf{Sc}(T) \to \Omega$ ) as on the right below.

$$T_{v_{Ge}} \longrightarrow T_{1,v_{Ge}} \longrightarrow \cdots \longrightarrow T_{n,v_{Ge}} \qquad T_{(-)} \longrightarrow T_{1,(-)} \longrightarrow \cdots \longrightarrow T_{n,(-)}$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

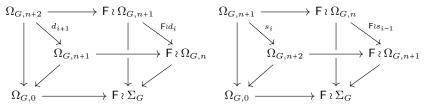
$$T'_{v_{Ge'}} \longrightarrow T'_{1,v_{Ge'}} \longrightarrow \cdots \longrightarrow T'_{n,v_{Ge'}} \qquad T'_{(-)} \circ \pi_* \longrightarrow T'_{1,(-)} \circ \pi_* \longrightarrow \cdots \longrightarrow T'_{n,(-)} \circ \pi_*$$

PTNARROWLOC EQ

Passing to colimits then gives the desired commutative diagram [4.25]. Moreover, diagrams of the form (4.25) clearly induce diagrams as in (4.31) and it is straightforward to check that these are inverse processes.

DSCOM REM

**Remark 4.32.** The diagrams (with back and lower slanted faces instances of (4.30))



commute whenever defined (i.e.  $0 \le i \le n + 1$ ).

INDVNG NOT

Notation 4.33. We will let

$$V_{G,n}:\Omega_{G,n}\to\mathsf{F}\wr\Sigma_G$$

be inductively defined by  $V_{G,n} = \sigma_0 \circ V_{G,n-1} \circ V_G$ .

**Remark 4.34.** When n = 2,  $V_{G,2}$  is thus the composite

$$\Omega_{G,2} \xrightarrow{V_G} \mathsf{F} \wr \Omega_{G,1} \xrightarrow{V_G} \mathsf{F} \wr \mathsf{F} \wr \Omega_{G,0} \xrightarrow{V_G} \mathsf{F} \wr \mathsf{F} \wr \mathsf{F} \wr \Sigma_G \xrightarrow{\sigma^0} \mathsf{F} \wr \mathsf{F} \wr \Sigma_G \xrightarrow{\sigma^0} \mathsf{F} \wr \Sigma_G$$

In light of Remarks 3.47 and 4.14,  $V_{G,n}(T_0 \to \cdots \to T_n)$  is identified with the tuple

$$(T_{n,v_{Ge}})_{v_{Ge} \in V_G(T_n)},$$
 (4.35) VGNISO EQ

though this requires changing the total order in  $V_G(T_n)$ . Rather than using the order induced by  $T_n$ , one instead equips  $V_G(T_n)$  with the order induced lexicographically from the maps  $V_G(T_n) \to V_G(T_{n-1}) \to \cdots \to V_G(T_0)$ , i.e., for  $v, w \in V_G(T_n)$  the condition v < w is determined by the lowest i such that the images of  $v, w \in V_G(T_i)$  are distinct.

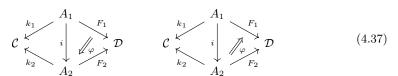
#### 4.4 A monad on spans

WSPAN DEF

**Definition 4.36.** We will write WSpan<sup>l</sup>( $\mathcal{C}, \mathcal{D}$ ) (resp. WSpan<sup>r</sup>( $\mathcal{C}, \mathcal{D}$ )), which we call the category of left weak spans (resp. right weak spans), to denote the category with objects the spans

$$\mathcal{C} \xleftarrow{k} A \xrightarrow{F} \mathcal{D},$$

arrows the diagrams as on the left (resp. right) below



which we write as  $(i, \varphi)$ :  $(k_1, F_1) \to (k_2, F_2)$ , and composition given in the obvious way.

Remark 4.38. There are natural isomorphisms

$$\mathsf{WSpan}^r(\mathcal{C}, \mathcal{D}) \simeq \mathsf{WSpan}^l(\mathcal{C}^{op}, \mathcal{D}^{op}). \tag{4.39}$$

TWISTEDARROWRIGHT EQ

RANLANADJ REM

**Remark 4.40.** The terms left/right<sub>ERSPANISU</sub> tend by the existence of adjunctions (which are seen to be equivalent by using (4.39)

Lan: WSpan<sup>$$l$$</sup> ( $\mathcal{C}, \mathcal{D}$ )  $\rightleftarrows$  Fun( $\mathcal{C}, \mathcal{D}$ ):  $\iota$   
 $\iota$ : Fun( $\mathcal{C}, \mathcal{D}$ )  $\rightleftarrows$  WSpan <sup>$r$</sup>  ( $\mathcal{C}, \mathcal{D}$ )<sup>op</sup>: Ran

where the functors  $\iota$  denote the obvious inclusions (note the need for the  $(-)^{op}$  in the second adjunction) and Lan/Ran denote the left/right Kan extension functors.

We will mainly be interested in the span categories  $\mathsf{WSpan}^l(\Sigma_G^{op}, \mathcal{V}) \simeq \mathsf{WSpan}^r(\Sigma_G, \mathcal{V}^{op})$ .

OMEGAGNA NOT

Notation 4.41. Given a functor  $\pi: A \to \Sigma_G$ , we let  $\Omega_{G,n}^{(A)}$  denote the pullback (in Cat)

$$\Omega_{G,n}^{(A)} \xrightarrow{V_{G,n}^{(A)}} \mathsf{F} \wr A$$

$$\downarrow \qquad \qquad \downarrow$$

$$\Omega_{G,n} \xrightarrow{V_{G,n}} \mathsf{F} \wr \Sigma_{G}$$

Explicitly, the objects of  $\Omega_{G,n}^{(A)}$  are pairs

$$(T_0 \to \cdots \to T_n, (a_e \uparrow \leq_e)_{(e} \uparrow \leq_e) \in V_C(T_n))$$
 (4.42) OMEGAGNA EQ

such that  $\pi(a_{e^{\uparrow} \leq e}) = T_{n,e^{\uparrow} \leq e}$ .

**Remark 4.43.** Our primary interest here will be in the  $\Omega_{G,0}^{(A)}$  construction. Importantly, the composite maps  $\Omega_{G,0}^{(A)} \to \Omega_{G,0} \to \Sigma_G$  allow us to iterate the  $\Omega_{G,0}^{(-)}$  construction. In practice, the role of higher strings  $\Omega_{G,n}^{(A)}$  will then be to provide more convenient models for iterated

Indeed, the content of Proposition 4.29 is then that there are compatible identifications  $\Omega_{G,0}^{(\Omega_{G,n})} \simeq \Omega_{G,n+1}$  which identify  $V_G^{(\Omega_{G,n})}$  with  $V_G$ .

Moreover, since all squares in the diagram

are pullback squares (the top center square is so by induction, the top right square by direct verification, the total top square by definition of  $\Omega_{G,n+1}^{(A)}$  and the bottom left square by

Proposition SUBSASPULL PROP 4.29), we likewise obtain identifications  $\Omega_G^{\left(\Omega_{G,n}^{(A)}\right)} \simeq \Omega_{G.n+1}^{(A)}$ .

**Proposition 4.45.** For any  $A \to \Sigma_G$  there are functors  $d_0^{(A)}: \Omega_{G,1}^{(A)} \to \Omega_G^{(A)}$  and natural isomorphisms

$$\Omega_{G,1}^{(A)} \xrightarrow{V_G} \operatorname{F}{\wr} \Omega_G^{(A)} \xrightarrow{\operatorname{F}{\wr} V_G} \operatorname{F}{\wr} \operatorname{F}{\wr} A$$

$$\downarrow_{\sigma^0} \qquad \qquad \downarrow_{\sigma^0} \qquad \qquad \downarrow_{\sigma^0} \qquad \qquad \downarrow_{\sigma^0} \qquad \qquad \qquad \downarrow_{$$

both natural in  $A \to \Sigma$ . Here naturality of  $\pi^{(-)}$  means that for a functor  $H: A \to B$  with corresponding diagram

one has an equality

$$(\mathsf{F} \wr H)\pi^{(A)} = \pi^{(B)}\Omega^{(H)}_{G,1}$$

(i.e. the two natural isomorphisms between the two distinct functors  $\Omega_{G,1}^{(A)} \Rightarrow \mathsf{F} \wr B$  coincide).

*Proof.* Informally, using the object description in (4.42),  $d_0^{(A)}$  is simply given by the formula

$$d_0^{(A)}\left(T_0 \to T_1, (a_{e^{\dagger} \le e})_{(e^{\dagger} \le e) \in V_G(T_1)}\right) = \left(T_1, (a_{e^{\dagger} \le e})_{(e^{\dagger} \le e) \in V_G(T_1)}\right), \tag{4.48}$$

though one must note that since in (4.42) the order in  $V_G(T_1)$  is induced lexicographically from the string, the two orders for  $V_G(T_1)$  in each side of (4.48) do not coincide. It now follows that the composites  $\sigma^0 \circ (\mathsf{F} \wr V_G^{(A)}) \circ V_G^{(A)}$  and  $V_G^{(A)} \circ d_0^{(A)}$  differ by the natural automorphism  $\pi^{(A)}$  given by the tuple permutations interchanging the two orders in  $V_G(T_1)$  for each  $T_0 \to T_{1\text{\tiny PICUBOIDAB EQ}}$ The commutativity of (4.47) is clear.

**Definition 4.49.** Suppose V has finite products.

We define an endofunctor N of Wspan<sup>r</sup>  $(\Sigma_G, \mathcal{V}^{op})$  by letting  $N(\Sigma_G \leftarrow A \rightarrow \mathcal{V}^{op})$  be the span  $\Sigma_G \leftarrow \Omega_G^{(A)} \rightarrow \mathcal{V}^{op}$  given composition along the diagram

$$\Omega_{G,0}^{(A)} \longrightarrow \mathsf{F} \wr A \longrightarrow \mathsf{F} \wr \mathcal{V}^{op} \xrightarrow{\Pi^{op}} \mathcal{V}^{op} 
\downarrow \qquad \qquad \downarrow 
\Omega_{G,0} \longrightarrow \mathsf{F} \wr \Sigma_{G} 
\downarrow 
\Sigma_{G}$$

and defined on maps of spans in the obvious way.

One has a multiplication  $\mu: N \circ N \Rightarrow N$  given by the natural isomorphisms

4.50) MULTDEFSPAN EQ

where  $\alpha$  is an associativity isomorphism for the product  $\Pi$ . We we note that naturality of  $\mu$  follows from the commutativity of (4.47).

Lastly, there is a unit  $\eta: id \Rightarrow N$  given by the strictly commutative diagrams

MONSPAN PROP

**Proposition 4.52.**  $(N, \mu, \eta)$  form a monad on Wspan<sup>r</sup>  $(\Sigma_G, \mathcal{V}^{op})$ .

*Proof.* The natural transformation component of  $\mu \circ (N\mu)$  is given by the composite diagram

(4.53) ASSOCSPAN1 EQ

whereas the natural transformation component of  $\mu \circ (\mu N)$  is given by

$$\Omega_{G,2}^{(A)} \to \mathsf{F} \wr \Omega_{G,1}^{(A)} \to \mathsf{F}^{\wr 2} \wr \Omega_{G}^{(A)} \to \mathsf{F}^{\wr 3} \wr A \to \mathsf{F}^{\wr 3} \wr \mathcal{V}^{op} \to \mathsf{F}^{\wr 2} \wr \mathcal{V}^{op} \to \mathsf{F}^{\wr 2} \wr \mathcal{V}^{op} \to \mathcal{V}^{op}$$

$$\downarrow^{\sigma^{0}} \qquad \downarrow^{\sigma^{0}} \qquad \downarrow^$$

(4.54) ASSOCSPAN2 EQ

That the rightmost sections of (4.53) and (4.54) coincide follows from compatibility of the associativity isomorphisms for  $\Pi^{op}$ .

For the leftmost sections, note first that, in either diagram, the top right and bottom left paths  $\Omega_{G,2}^{(A)} \to \mathsf{F} \wr A$  differ only by the induced order on  $V_G(T_2)$  for each string  $T_0 \to T_1 \to T_2$ . More explicitly, the top right paths use the order induced lexicographically from the string  $T_0 \to T_1 \to T_2$  while the bottom left paths use the order induced exclusively by  $T_2$ . The two left sections then coincide since are both given by the permutation interchanging these orders, the only difference being that the interpretate stage of (4.53) uses the order induced lexicographically from  $T_0 \to T_2$  while (4.54) uses the order induced lexicographically from  $T_1 \to T_2$ .

As for unit conditions,  $\mu \circ (N\eta)$  is represented by

while  $\mu \circ (\eta N)$  is represented by

It is straightforward to check that the composites of the left and right sections of both (4.55) and (4.56) are strictly commutative diagrams, and thus that (4.55) and (4.56) coincide.  $\Box$ 

## The free genuine operad monad

Recalling that  $\mathsf{Wspan}^r(\Sigma_G, \mathcal{V}^{op}) \simeq \mathsf{Wspan}^l(\Sigma_G^{op}, \mathcal{V})$ , Proposition MONSPAN PROP RANLANADJ REM an adjuntion

$$\mathsf{Lan:WSpan}^l(\Sigma_G^{op}, \mathcal{V}) \rightleftarrows \mathsf{Fun}(\Sigma_G^{op}, \mathcal{V}) \colon \iota \tag{4.57}$$

together with a monad N in the leftmost category  $\mathsf{WSpan}^l(\Sigma_G^{op}, \mathcal{V})$ . We now turn to showing that, under reasonable hypothesis on  $\mathcal{V}$ , the composite Lan  $\circ N \circ \iota$  inherits a monad structure from N. The key will be to show that under such conditions the map  $\mathsf{Lan} \circ N \Rightarrow \mathsf{Lan} \circ N \circ \iota \circ \mathsf{Lan}$ is a natural isomorphism.

Recall that following Convention  $\frac{\text{PLANARCONV CON}}{3.18 \text{ our model}}$  for  $O_G$  consists of totally ordered sets. One therefore has root functors

$$\Omega_C^q \xrightarrow{\mathsf{r}} \mathsf{O}_G, \qquad \Sigma_G \xrightarrow{\mathsf{r}} \mathsf{O}_G$$

sending each planar G-tree to its ordered orbital G-set of roots.

Root functors are compatible with the leaf-root functor and the inclusion, i.e. the following commute.

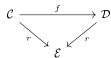
$$\Omega_G^q \xrightarrow{\operatorname{lr}} \Sigma_G \qquad \Sigma_G \hookrightarrow \Omega_G^q$$

$$\downarrow^{\operatorname{r}} \qquad \downarrow^{\operatorname{r}} \qquad (4.58) \quad \boxed{\operatorname{ROOTLEAFTROOTCOM EQ}}$$

Moreover, the diagrams ( $\overline{A.58}$ ) possess some extra structure we will need to make use of. Indeed, both functors are split Grothendieck fibrations: given a map  $\varphi: A \to B$  in  $O_G$  and G-tree T such that r(T) = B we can build a cartesian arrow  $\varphi^*(T) \to T$  by letting  $\varphi^*(T)$ to be the pullback G-tree together with the planar structure on roots given by A and on non-equivariant nodes given by their image via  $\varphi^*(T) \to T$ . It now follows that (4.58) are diagrams of split Grothendieck fibrations.

One advantage of split Grothendieck fibrations (in general) is the following initiality condition on overcategories.

**Definition 4.59.** Suppose we have two split fibrations  $r: \mathcal{C} \to \mathcal{E}$  and  $r: \mathcal{D} \to \mathcal{E}$  and a map of fibrations f as below.



For any  $d \in \mathcal{D}$ , let  $d \downarrow_r \mathcal{C}$  denote the subcategory of  $d \downarrow \mathcal{C}$  of just those maps  $d \to f(c)$  which map to the identity under r.

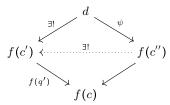
**Lemma 4.60.** In the above scenario,  $d \downarrow_r C$  is inital in  $d \downarrow C$ .

*Proof.* We must show that the appropriate iterated overcategories are non-empty and connected. To that end, suppose we have a map  $\phi: d \to f(c)$  in  $\mathcal{D}$ , and let  $q:=r(\phi):r(d)\to r(c)$ . Since  $\mathcal{C}$  is fibrant over  $\mathcal{E}$ , we have a Cartesian arrow  $q':c'\to c$  lifting  $q:r(d)=r(c')\to r(c)$ . Thus we have diagrams in  $\mathcal{E}$  and  $\mathcal{D}$  of the form

$$r(d) \xrightarrow{id} r(c')$$
 $\downarrow^q$ 
 $r(c)$ 
 $\downarrow^q$ 
 $f(c')$ 
 $\downarrow^q$ 
 $f(c')$ 

As f preserves Cartesian arrows, f(q') is Cartesian, and hence we have a unique lifting  $p(\phi): d \to f(c')$  of r(d) = r(c'). Thus the iterated overcategory is inhabited.

To show it is connected, any other factorization  $d \stackrel{\psi}{\to} f(c'') \to f(c)$  such that  $r(\psi) = id$ , we can again use the fact that q' is Cartesian to produce a lift  $f(c';) \to f(c')$  of r(c'') = r(c'); by uniqueness of the map  $d \to f(c')$ , we have that the diagram below commutes, finishing the proof.



**Remark 4.61.** The above proof can be easily modified to show that  $d \downarrow_r \mathcal{C}$  is in fact a coreflective subcategory of  $d \downarrow \mathcal{C}$ , with  $\phi \mapsto p(\phi)$  the reflection.

Luis: to actually show this would take just as long - functoriality and the fact that the inclusion is a left adjoint both require unique liftings/factorizations

**Definition 4.62.** A split Grothendieck fibration  $A \xrightarrow{\mathsf{r}} \mathsf{O}_G$  is called a *root fibration* and a split Grothendieck fibration diagram



is called a root fibration functor.

ROOTFIBPULL LEM

The relevance of root fibrations is given by the following couple of lemmas.

**Lemma 4.63.** If  $A \to \Sigma_G$  is a root fibration functor then so is  $\Omega_G^{(A)} \to \Omega_G$ , naturally in A.

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*Proof.* We consider the pullback diagram below.

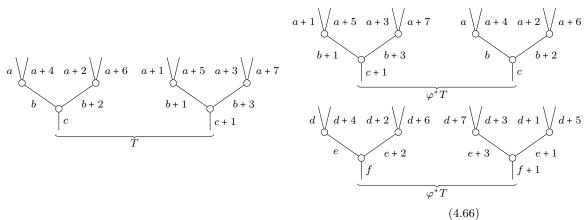
$$\Omega_{G,0}^{(A)} \xrightarrow{V_G^{(A)}} \operatorname{F} \wr A 
\downarrow \qquad \qquad \downarrow 
\Omega_{G,0} \xrightarrow{V_G} \operatorname{F} \wr \Sigma_G$$

$$(4.64) \quad \boxed{\text{ROOTIMPLIESROOT EQ}}$$

The hypothesis that  $A \to \Sigma_G$  is root fibration implies that the rightmost map in (4.64) is a map of split Grothendieck fibrations over  $\mathsf{F} \wr \mathsf{O}_G$ .

Since the map  $V_G$  sends the chosen cartesian arrows in  $\Omega_{G,0}$  (over  $O_G$ ) to chosen cartesian arrows of  $F \wr \Sigma_G$  (over  $F \wr O_G$ ), the result follows.

**Example 4.65.** Let  $G = \mathbb{Z}_{/8}$ . The following exemplifies a pull back along the twist map  $\varphi: G/2G \to G/2G$  (i.e., accounting for order,  $\varphi$  is the permutation (12)), with the topmost representation of  $\varphi^*T$  maintaining the chosen generators for each edge orbit from T and the bottom representation choosing instead the generators to be minimal with regard to the planar structure.



We note that  $(\varphi^*(T))_{v_{Ge}} = \psi^*(T_{v_{Gb}})$  for  $\psi$  the permutation (13)(24) encoded by the composite identifications  $\{1, 2, 3, 4\} \simeq \{e, e + 2, e + 3, e + 1\} \simeq \{b + 1, b + 3, b, b + 2\} \simeq \{3, 4, 1, 2\}$ .

**Lemma 4.67.** Suppose that V is complete and that  $A \to \Sigma_G$  is a root fibration. If the rightmost triangle in

is a right Kan extension diagram then so is the composite diagram.

*Proof.* Unpacking definitions using the pointwise formula for right Kan extensions ([4, X.3.1]), it suffices to check that for each  $T \in \Omega_{G,0}$  the functor

$$T \downarrow \Omega_{G,0}^{(A)} \to V_G(T) \downarrow \mathsf{F} \wr A$$
 (4.69) LANPULLCOMA EQ

is initial. In the course of the proof of Lemma 4.3 it was shown that the subcategory

$$\prod_{v_{Ge} \in V_G(T)} T_{v_{Ge}} \downarrow A$$

is initial in the  $V_G(T) \downarrow \mathsf{F} \wr A$ .

LANPULLCOMA LEM

On the other hand, since  $\Omega_G^{(A)} \to \Omega_G$  is a root fibration functor,  $T \downarrow \Omega_G^{(A)}$  has an initial subcategory  $T \downarrow_{r,\simeq} \Omega_G^{(A)}$  with objects  $(S \in \Omega_G^{(A)}, T \to u(S))$  such that  $T \to u(S)$  is a quotient map that induces an ordered isomorphism on roots. Note that this can be restated as saying that  $T \to u(S)$  is an isomorphism preserving the order of the roots.

The result now follows from the natural isomorphism

$$T\downarrow_{\mathsf{r},\simeq}\Omega_G^{(A)}\simeq\prod_{v_{Ge}\in V_G(T)}T_{v_{Ge}}\downarrow_{\mathsf{r},\simeq}A. \tag{4.70} \qquad \boxed{\mathsf{TDOWNISOA}\ \mathsf{EQ}}$$

To see this, we focus first on the case  $A = \Sigma_G$ . In that case, the left hand side of (4.70)encodes replanarizations of T that preserve the root order. On the other hand, the right hand side encodes replanarizations of all the G-vertices that preserve the order of their roots, or, equivalently, replanarizations of the non-equivariant vertices of T. That these are equivalent is the content of Proposition 3.13.

Note that  $(T \to S) \in (T \downarrow_{r,\simeq} \Omega_G)$  is then encoded by a tuple  $(T_{v_{Ge}} \to \varphi_{v_{Ge}}^* S_{v_{Ge}})_{v_{Ge} \in V_G(T)}$ 

where the pullbacks  $\varphi_{v_{Ge}}^*$  are needed to correct the root order.

The case of general A follows likewise, using the corresponding pullbacks  $\varphi_{v_{Ge}}^*$ .

Note: an addendum is needed to show that (1.70) suffices, since  $T \downarrow_{\mathsf{r},\simeq} \Omega_G^{(A)}$  is not sent

directly to  $\prod_{\substack{V_{Ge} \in V_G(T) \\ ROOTFIBPULL \ LEM}} T_{v_{Ge}} \downarrow_{r,\cong} A$ Lemma 4.63 can be interpreted as saying that, if one defines a category  $\mathsf{Wspan}^l_r(\Sigma_G^{op}, \mathcal{V})$ of rooted spans

$$\Sigma_G^{op} \leftarrow A^{op} \rightarrow \mathcal{V}$$

where  $A \to \Sigma_G$  is a root fibration functor, the monad N built in Proposition 4.52 lifts to a monad  $N_r$  in Wspan<sub>r</sub> ( $\Sigma_G^{op}, \mathcal{V}$ ), and likewise for the adjunction (4.57).

Corollary 4.71. Suppose that finite products in V commute with colimits in each variable. The functors

$$\mathsf{Lan} \circ N_{\mathsf{r}} \Rightarrow \mathsf{Lan} \circ N_{\mathsf{r}} \circ \iota \circ \mathsf{Lan}, \qquad \mathsf{Lan} \circ \iota \Rightarrow i a$$

 $are\ natural\ isomorphisms.$ 

THEMONAD DEF

REPACKAGERES REM

Proof. This follows by combining Lemma 4.67 with Lemma 4.3.

**Definition 4.72.** The genuine equivariant operad monad is the monad  $\mathbb{F}_G$  on  $\mathsf{Fun}(\Sigma_G^{op}, \mathcal{V})$ given by

$$\mathbb{F}_G = \mathsf{Lan} \circ N_{\mathsf{r}} \circ \iota$$

and with multiplication and unit given by the composites

$$\mathsf{Lan} \circ N_{\mathsf{r}} \circ \iota \circ \mathsf{Lan} \circ N_{\mathsf{r}} \circ \iota \stackrel{\simeq}{\leftarrow} \mathsf{Lan} \circ N_{\mathsf{r}} \circ N_{\mathsf{r}} \circ \iota \Rightarrow \mathsf{Lan} \circ N_{\mathsf{r}} \circ \iota$$

$$id \stackrel{\sim}{\Leftarrow} \mathsf{Lan} \circ \iota \Rightarrow \mathsf{Lan} \circ N_{\mathsf{r}} \circ \iota.$$

**Remark 4.73.** The functor  $\mathsf{Lan} \circ N_r \circ \iota$  is isomorphic to  $\mathsf{Lan} \circ N_r \circ \iota$  and this isomorphism is compatible with the multiplication and unit in Definition 4.72, and we will henceforth simply write N rather than  $N_r$ .

From this point of view, the role of root fibrations is to guarantee that Lan  $\circ N \circ \iota$  is indeed a monad, but unnecessary to describe the monad structure itself.

Remark 4.74. Since a map

$$\mathbb{F}_G X = \mathsf{Lan} \circ N_\mathsf{r} \circ \iota X \to X$$

is adjoint to a map

$$N_r \circ \iota X \to \iota X$$

one easily verifies that X is a genuine equivariant operad, i.e. a  $\mathbb{F}_{G}$ -algebra, iff  $\iota X$  is a N-algebra. Moreover, the bar resolution

$$\mathbb{F}_G^{\bullet+1}X$$

is isomorphic to

$$\operatorname{Lan}\left(N^{\bullet+1}\iota X\right).$$

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# 5 Genuine equivariant operads versus equivariant operads

We start by noting that in the case G = \*, genuine operads simply recover the usual notion of symmetric operads, i.e.  $\mathsf{Sym}_*(\mathcal{V}) \simeq \mathsf{Sym}(\mathcal{V})$  and  $\mathsf{Op}_*(\mathcal{V}) \simeq \mathsf{Op}(\mathcal{V})$ , and in what follows we will adopt the notations  $\mathsf{Sym}^G(\mathcal{V})$  and  $\mathsf{Op}^G(\mathcal{V})$  for the corresponding categories of G-objects. Our goal will be to relate these to the categories  $\mathsf{Sym}_G(\mathcal{V})$  and  $\mathsf{Op}_G(\mathcal{V})$  of genuine equivariant sequences and genuine operads.

We will throughout this section fix a total order of G such that the identity e is the first element, though we note that the exact order is unimportant, as any other such choice would lead to unique isomorphisms between the constructions in this section.

We thus have an inclusion functor

$$\iota \colon G \times \Sigma \longleftrightarrow \Sigma_G$$

$$C \longmapsto G \cdot C$$

where  $G \cdot C$  is planarized so that the roots inherit the order of G and each of the individual copies of C inherits the planarization of C. We now consider the associated adjunction.

$$\iota^* : \operatorname{Sym}_G \longrightarrow \operatorname{Sym}^G : \iota_*$$
 (5.1) TWOADJOINTS EQ

Explicitly, we have the formulas

$$\iota^*X(C) = X(G \cdot C), \qquad \iota_*Y(\coprod_i C_i) = \left(\prod_i Y(C_i)\right)^G,$$

where we note that in the formula for  $\iota_*(-)$  the action of G interchanges factors according to the action on the indexing set, i.e. the root orbit of  $\coprod_i C_i$ . As a side note, we note that the formula for  $\iota_*$  is independent of the chosen order of G.

HERE

# 6 Free extensions

Our overall goal in this section will be to produce a description of free genuine operad pushouts, i.e. pushouts of the form

$$\begin{array}{ccc}
\mathbb{F}_G A & \longrightarrow X \\
\downarrow & & \downarrow \\
\mathbb{F}_G B & \longrightarrow Y
\end{array}$$

in the category  $\mathsf{Op}_G$  of genuine equivariant operads.

# 6.1 Extensions over general monads

Any monad T on  $\mathcal C$  one obtains induced monads  $T^{\times l}$  on  $\mathcal C^{\times l}$ , and we will make use of several standard relations between these. In particular, any map  $\alpha: \underline l \to \underline m$  induces a forgetful functor such that for the forgetul functor  $\alpha^*: \mathcal C^{\times l} \to \mathcal C^{\times n}$  one has  $T^{\times \overline l} \alpha^* \simeq \alpha^* T^{\times m}$ .

Indeed, we will need to make use of a slightly more general setup. Letting I denote the identity monad on  $\mathcal{C}$ , and  $K \subset \underline{m}$  be a subset, there is a monad  $T^{\times K} \times I^{\times (\underline{m}-K)}$  on  $\mathcal{C}^{\times m}$ , which we abusively denote simply as  $T^{\times K}$ . Identities then determine maps of monads  $T^J \to T^{\times K}$  whenever  $J \subset K$  and, moreover, there are identifications  $T^{\times \alpha^{-1}(K)}\alpha^* \simeq \alpha^*T^{\times K}$ . One then has the following.

**Proposition 6.1.** The functor

$$T^{\times \alpha^{-1}(K)} \Rightarrow \alpha^* T^{\times K} \alpha_! \tag{6.2}$$

MONADFUNCTORALPHA EQ

adjoint to the identification  $T^{\times \alpha^{-1}(K)} \alpha^* \simeq \alpha^* T^{\times K}$  is a map of monads on  $\mathcal{C}^{\times n}$ .

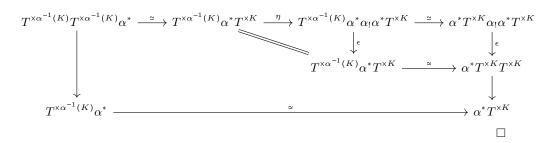
*Proof.* We first note that there are identifications of functors  $(FG)^{\times K} \simeq F^{\times K}G^{\times K}$  which are compatible with the identifications  $F^{\times \alpha^{-1}(K)}\alpha^* \simeq \alpha^*F^{\times K}$  in the sense that the identification  $(FG)^{\times \alpha^{-1}(K)} \circ \alpha^* \simeq \alpha^*(FG)^{\times K}$  matches the composite identification  $F^{\times \alpha^{-1}(K)}G^{\times \alpha^{-1}(K)}\alpha^* \simeq F^{\times K}G^{\times K}$ .

Letting  $\eta, \epsilon$  denote the unit and counit for the  $(\alpha_!, \alpha^*)$  adjunction, (6.2) is then the composite

$$T^{\times \alpha^{-1}(K)} \xrightarrow{\eta} T^{\times \alpha^{-1}(K)} \alpha^* \alpha_! \simeq \alpha^* T^{\times K} \alpha_!.$$

That this is a monad map is the condition that the following multiplication and unit diagrams commute.

We argue only the case of the leftmost multiplication diagram, with commutativity of the unit diagram following by a similar but simpler argument. Since the precomposition  $(-) \circ \alpha^*$  is the left adjoint to the precomposition  $(-) \circ \alpha_!$  this follows from the following diagram.



EXTGENMON SEC

MONADICFUN PROP

TALPHAKMOD REM

COMPPOSTCOMP REM

**Remark 6.3.** Since  $T^{\times K}\alpha_!$  is a right  $\alpha^*T^{\times K}\alpha_!$ -module, Proposition 6.1 implies that it is also a right  $T^{\times \alpha^{-1}(K)}$ -module or, moreover, a right  $T^{\times J}$ -module whenever  $\alpha(J) \subset K$ .

**Remark 6.4.** Combining the precomposition and postcomposition adjunctions, the identification  $T^{\times \alpha^{-1}(K)}\alpha^* \simeq \alpha^* T^{\times K}$  is then adjoint to a functor  $\alpha_! T^{\times \alpha^{-1}(K)} \to T^{\times K}\alpha_!$  which is readily checked to be a map of right  $T^{\times \alpha^{-1}(K)}$ -modules.

More generally, for  $\alpha(J) \subset K$ , the composite  $T^{\times J} \alpha^* \to T^{\times \alpha^{-1}(K)} \alpha^* \simeq \alpha^* T^{\times K}$  is thus adjoint to a map of right  $T^{\times J}$ -modules

$$\alpha_! T^{\times J} \to T^{\times K} \alpha_!.$$
 (6.5)

RIGHTMODULETMAP EQ

RIGHTMODULETMAPAUX EQ

We now unpack the content of (6.5) when  $\alpha: \underline{l} \to *$  is the unique map to the simpleton  $* = \underline{1}$ . In this case we can instead write  $\alpha_! = \coprod_{l} \alpha^* = \Delta$ , and we thus have commutative diagrams

$$\coprod_{J} TT A_{j} \amalg \coprod_{\underline{n}-J} A_{j} \longrightarrow T \left( \coprod_{J} T A_{j} \amalg \coprod_{\underline{n}-J} A_{j} \right) 
\downarrow \qquad \qquad \downarrow \qquad \qquad (6.6)$$

$$\coprod_{J} T A_{j} \amalg \coprod_{\underline{n}-J} A_{j} \longrightarrow T \left( \coprod_{J} A_{j} \amalg \coprod_{\underline{n}-J} A_{j} \right)$$

r the

where the vertical maps come from the right  $T^{\times J}$ -module structure. Writing  $\coprod^a$  for the correduct of T algebras and recalling the canonical identifications  $\coprod_K^a (TA_k) \simeq T(\coprod_K A_k)$ , (6.6) in fact shows that the right  $T^{\times J}$ -module structure on  $T \circ \coprod$  in fact codifies the multiplication maps

$$\coprod_{J}^{a} TTA_{j} \amalg^{a} \coprod_{\underline{l}-J}^{a} TA_{j} \to \coprod_{J}^{a} TA_{j} \amalg^{a} \coprod_{\underline{l}-J}^{a} TA_{j}.$$

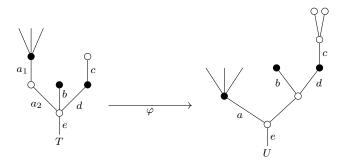
# 6.2 Labeled planar strings

We now translate the results in the previous section to the context of the monad N on  $\mathsf{WSpan}^l(\Sigma^{op}, \mathcal{V})$ . In analogy to the planar string models  $\Omega^{(A)}_{G,n}$  for iterations  $N^{\circ n+1}$  of the monad N, we will find it convenient to build similar string models  $\Omega^{(\underline{A}_J)}_{G,n}$  for  $N \circ \coprod \circ (N^{\times J})^{\circ n}$ .

**Definition 6.7.** A *l-node labeled G-tree* (or just *l-labeled G-tree*) *G*-tree is a pair  $(T, V_G(T) \rightarrow \{1, \cdots, l\})$  with  $T \in \Omega_G$ , which we think of as a *G*-tree together with *G*-vertices labels in  $1, \cdots, l$ . Further, a tall map  $\varphi: T \rightarrow S$  between *l*-labeled trees is called a *label map* if for each *G*-vertex  $v_{Ge}$  of T with label j, the vertices of the subtree  $S_{v_{Ge}}$  are all labeled by j.

Lastly, given a subset  $J \subset \underline{l}$ , a planar label map  $\varphi: T \to S$  is said to be J-inert if for every G-vertex  $v_{Ge}$  of T with label  $j \in J$  it is  $S_{v_{Ge}} = T_{v_{Ge}}$ .

**Example 6.8.** Consider the 2-labeled trees below (for G = \* the trivial group), with black nodes ( $\bullet$ ) denoting labels by the number 1 and white nodes ( $\bullet$ ) labels by the number 2. The planar map  $\varphi$  (sending  $a_i \mapsto a$ ,  $b \mapsto b$ ,  $c \mapsto c$ ,  $d \mapsto d$ ,  $e \mapsto e$ ) is a label map which is  $\{1\}$ -inert.



(6.9) SUBSDATUMTREESLAB EQ

LABELEDTREES EX

**Definition 6.10.** Let  $0 \le s \le n$  and  $J \subset \underline{l}$  be a subset.

We define  $\Omega_{G,n,s}^{J}$  to have as objects n-planar strings

$$T_0 \xrightarrow{f_1} T_1 \xrightarrow{f_2} \cdots \xrightarrow{f_s} T_s \xrightarrow{f_{s+1}} T_{s+1} \xrightarrow{f_{s+2}} \cdots \xrightarrow{f_n} T_n$$
 (6.11) NSTRINGLAB EQ

together with *l*-labelings of  $T_s, T_{s+1}, \dots, T_n$  such that the  $f_r, r > s$  are  $(\underline{l} - J)$ -inert label maps. Arrows in  $\Omega_{G,n,s}^J$  are quotients of strings  $(q_r: T_r \to T_r')$  such that  $q_r, r \le s$  are label maps.

Informally,  $\Omega_{G,n,s}^{\underline{l}}$  consists of *n*-strings such that trees and maps after  $T_s$  are *l*-labeled.

**Remark 6.12.** Our main case of interest will that of s = 0, in which case we abbreviate

 $\Omega^{J}_{G,n} = \Omega^{J}_{G,n,0}. \text{ Indeed, such strings will suffice to build models for } N \circ \coprod_{\text{TALPHAKMOD REM}} \circ (N^{\times J})^{\circ n}$  However, to unpack the right  $N^{\times J}$ -module structure as in Remark 5.3 one further needs to encode composites  $NN \circ \coprod \circ (N^{\times J})^{\circ n-1}$ , a role played by strings  $\Omega^{J}_{G,n,1}.$ 

Notation 6.13. We will further write

$$\Omega_{G,n,-1}^{J} = \coprod_{J} \Omega_{G,n} \coprod \coprod_{\underline{l}-J} \Sigma_{G}, \qquad \Omega_{G,n,n+1}^{J} = \Omega_{G,n}$$

$$\tag{6.14}$$

To justify this convention, we note that a string as in (6.11) can be extended by prepending to it the map  $\operatorname{Ir}(T_0) = T_{-1} \xrightarrow{f_0} T_0$ . If one then attempts to define  $\Omega^J_{G,n,-1}$  by insisting that  $T_{-1}$  also be labeled, it follows that all node labels in each string must coincide, resulting in the coproduct decomposition in (6.14).

There are a number of obvious functors relating the  $\Omega_{G,n,s}^{J}$  categories, which we now make explicit. Given  $s \leq s'$  or  $J \subset J'$  there are forgetful functors

$$\Omega^{J}_{G,n,s} \to \Omega^{J}_{G,n,s'} \qquad \Omega^{J}_{G,n,s} \to \Omega^{J'}_{G,n,s}$$
 (6.15) NKNFGT EQ

OMEGANMINUSONE EQ

 $\Omega^{J}_{G,n,s} \to \Omega^{J}_{G,n,s'} \qquad \Omega^{J}_{G,n,s} \to \Omega^{J'}_{G,n,s} \tag{6.15}$  The simplicial operators in Notation 4.26 generalize to operators (where  $0 \le i \le n, -1 \le j \le n$ )

$$\begin{split} d_i : & \Omega^J_{G,n,s} \to \Omega^J_{G,n-1,s-1} \qquad i < s \qquad \qquad s_j : \Omega^J_{G,n,s} \to \Omega^J_{G,n+1,s+1} \qquad j < s \\ d_i : & \Omega^J_{G,n,s} \to \Omega^J_{G,n-1,s} \qquad s \leq i \qquad \qquad s_j : \Omega^J_{G,n,s} \to \Omega^J_{G,n+1,s} \qquad s \leq j \end{split}$$

which are compatible with the forgetful functors in the obvious way.

**Remark 6.16.** For  $J \subset J'$  the forgetful functor in (6.15) is a fully faithful inclusion. However, and somewhat subtly, this is not the case the for the  $s \le s'$  forgetful functors. Indeed, regarding  $T \to U$  in Examples 6.8 as an object in  $\Omega^2_{*,n,0}$ , changing the label of the  $a_1 \le a_2$ vertex of T from a  $\circ$ -label to a  $\bullet$ -label yields an alternate object  $\bar{T} \to U$  of  $\Omega^2_{*,n,0}$  forgetting to the same object of  $\Omega^{\underline{2}}_{*,n,1}$ , yet  $T \to U$  and  $\overline{T} \to U$  are not isomorphic.

We note that this is a consequence of the fact that substitution data can replace unary nodes by stumps, which have no nodes.

Generalizing Notation 4.33 there is a commutative diagram

$$\begin{array}{ccc} \Omega_{G,n,s}^{J} & \xrightarrow{V_{G,n}} \operatorname{F} \wr \Sigma_{G}^{\amalg l} \\ & \downarrow & \downarrow \\ \Omega_{G,n} & \xrightarrow{V_{G,n}} \operatorname{F} \wr \Sigma_{G} \end{array}$$

where for a labeled string it is  $V_G(T_0 \to \cdots \to T_n) = (T_{n,v_{Ge}})_{V_G(T_n)}$ , where we regard  $T_{n,v_{Ge}} \in$  $\Sigma_G^{ul} \simeq \Omega_{G,-1,-1}^{\underline{l}} \text{ by using the label in } \underline{\text{MECAGNA}} \text{ NOT}$  We now expand Notation 4.41.

**Notation 6.17.** Let  $\underline{A}$  denote a  $\underline{l}$ -tuple  $(\pi_j: A_j \to \Sigma_G)_{\underline{l}}$  of categories over  $\Sigma_G$ . We define  $\Omega_{G,n,s}^{(\underline{A}),J}$  by the pullback diagram

$$\Omega_{G,n,s}^{(\underline{A}),J} \xrightarrow{V_{G,n}^{(\underline{A})}} \mathsf{F} \wr \coprod A_{j} 
\downarrow \qquad \qquad \downarrow 
\Omega_{G,n,s}^{J} \xrightarrow{V_{G,n}} \mathsf{F} \wr \Sigma_{G}^{\mathsf{u}l} 
(6.18) LTUPLEAPULL EQ$$

Explicitly, an object of  $\Omega^{(\underline{A}),J}_{G,n,s}$  consists of a labeled string  $T_0 \to \cdots T_n$  as in (6.11) together with a tuple  $(a_{v_{Ge}})_{V_G(T_n)}$  such that  $a_{v_{Ge}} \in A_j$  if  $v_{Ge}$  has label j and  $\pi_j(a_{v_{Ge}}) = T_{n,v_{Ge}}$ .

The reader may have noticed a certain asymmetry between our definition of the  $V_{G,n}$ functors here versus their analogues in §4.3, where they were defined iteratively in terms of the control of th of simpler functors  $V_G$ . This is because of the possibility that s = -1, in which case (6.14) applies and some caution is needed in that the following result fails.

Proposition 6.19. Suppose  $0 \le s \le n$ . One has a diagram of pullback squares (generalizing (4.44))

$$\Omega_{G,n,s}^{(A),J} \xrightarrow{V_G^{(A)}} \operatorname{F}{}^{\wr}\Omega_{G,n-1,s-1}^{(A),J} \xrightarrow{\operatorname{F}{}^{\wr}V_{G,n}^{(A)}} \operatorname{F}{}^{\wr}\operatorname{F}{}^{\wr}\coprod A_j \xrightarrow{\sigma^0} \operatorname{F}{}^{\wr}\coprod A_j$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

$$\Omega_{G,n,s}^J \xrightarrow{V_G} \operatorname{F}{}^{\wr}\Omega_{G,n-1,s-1}^J \xrightarrow{\operatorname{F}{}^{\wr}V_{G,n}} \operatorname{F}{}^{\wr}\operatorname{F}{}^{\wr}\Sigma_G^{ul} \xrightarrow{\sigma^0} \operatorname{F}{}^{\wr}\Sigma_G^{ul}$$

$$\downarrow \qquad \qquad \downarrow$$

$$\Omega_{G,0} \xrightarrow{V_G} \operatorname{F}{}^{\wr}\Sigma_G$$

$$(6.20)$$
ALLSQUARESJ EQ

such that the composite of the top squares is (6.18). Proof. The  $V_G$  functors are defined just as in (4.28) via the formula

$$V_G(T_0 \to T_1 \to \cdots \to T_n) = (T_{1,v_{Ge}} \to \cdots \to T_{n,v_{Ge}})_{v_{Ge} \in V_G(T_0)}$$

with the strings  $T_{1,v_{Ge}} \to \cdots \to T_{n,v_{Ge}}$  inheriting the extra structure in the obvious way.

Since the top composite square, top center square and top right square are all pullback squares, it remains only to show that the bottom left square is a pullback. This last claim  $\frac{\text{SUBSASPULL FROP}}{\text{FROP}}$ is simply a variation of Proposition 4.29, and follows from the same proof, since both labels and inertness conditions are inherited when assembling substitution data into trees via Proposition 3.42.  $\hfill\Box$ 

#### Bar constructions on spans 6.3

We use the results in the previous sections to obtain a string description of the bar constructions

$$\coprod_{J}^{a} N^{\bullet+1} A_j \coprod^{a} \coprod_{\underline{l}-J}^{a} N A_j.$$

For simplicity, we discuss first the particular case  $\coprod^a N^{\bullet+1}A$ . Writing the span as  $\Sigma_G \leftarrow$  $A \xrightarrow{F} \mathcal{V}$  the identifications  $\Omega_{G,0}^{\left(\Omega_{G,n}^{(A)}\right)} \simeq \Omega_{G,n+1}^{(A)}$  iteratively identify the operator in the bar construction  $N^{\bullet+1}A$  as follows.

The top boundaries  $d_n$  have natural transformation given by

$$\Omega_{G,n}^{(A)} \xrightarrow{V_G^{on}} \mathsf{F}^{\imath n} \wr \Omega_{G,0}^{(A)} \xrightarrow{\mathsf{F}^{\imath n} \wr F_1} \mathsf{F}^{\imath n} \wr \mathcal{V}^{op} \xrightarrow{\Pi^{on}} \mathcal{V}^{op}$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \parallel \qquad \qquad \parallel$$

$$\Omega_{G,n-1}^{(A)} \xrightarrow{V_G^{on}} \mathsf{F}^{\imath n} \wr A \xrightarrow{\mathsf{F}^{\imath n} \wr F} \mathsf{F}^{\imath n} \wr \mathcal{V}^{op} \xrightarrow{\Pi^{on}} \mathcal{V}^{op}$$

$$(6.21)$$

where m is the natural transformation component of the multiplication  $NA \to A$ , and the remaining differentials  $d_i$  for  $0 \le i < n$  are given by

where  $\pi_i^{(A)}$  interchanges lexicographic orders on the *i*-th F coordinate of  $F^{in}$  and  $\alpha_i$  is the natural associativity isomorphism.

## Maybe add degeneracies

Similarly, Proposition 6.19 shows that  $\Omega_{G,n}^{(A)} \simeq \Omega_{G,0}^{\left(\coprod \Omega_{G,n-1}^{(A_j)}\right)}$  so that the top boundaries  $d_n$  in the bar construction  $N \circ \coprod \circ (N^{\times l})^{\circ n} \underline{A}$  are given by

$$\Omega_{G,n}^{(\underline{A})} \xrightarrow{V_G} \mathsf{F} \wr \coprod \Omega_{G,n-1}^{(A_j)} \xrightarrow{V_G^{\circ n-1}} \mathsf{F} \wr \coprod \mathsf{F}^{\imath n-1} \wr \Omega_{G,0}^{(A_j)} \xrightarrow{F_1} \mathsf{F} \wr \coprod \mathsf{F}^{\imath n-1} \wr \mathcal{V}^{op} \xrightarrow{\Pi^{\circ n-1}} \mathsf{F} \wr \mathcal{V}^{op} \xrightarrow{\Pi} \mathcal$$

where  $\underline{m}$  stands for the functor induced by the tuple of multiplication maps  $m_j: NA_j \to A_j$ , and the other boundaries  $d_i$  for  $0 \le i < n$  are given by

where again  $\pi_i^{(\underline{A})}$  interchanges lexicographic orders on the i-th F coordinate and  $\alpha_i$  is again an associativity isomorphism. We note that (6.24) follows directly from (6.22) for I-phaked but that the case i=0, which uses the  $N^{\times l}$  right action on  $N\circ \amalg$  (cf. Remark 6.3), which after unpacked leads to the composite diagram below.

Finally, using the inclusions  $\Omega_{G,n}^{(\underline{A}),J} \hookrightarrow \Omega_{G,n}^{(\underline{A})}$ , one obtains analogous descriptions of the bar constructions  $N \circ \mathfrak{U} \circ (N^{\times J})^{\circ n} \underline{A}$ , depicted below.

TRANSFSIMP SEC

MPSPANREIN LEMMA

#### Transferring simplicial colimits of left Kan extensions 6.4

Given genuine equivariant operads  $X,Y\in\mathsf{Op}_G$  one has an isomorphism

$$X \coprod^{a} Y \simeq \operatorname{colim}_{\Delta^{op}} \left( \mathbb{F}_{G}^{\bullet+1} X \coprod^{a} \mathbb{F}_{G}^{\bullet+1} Y \right)$$

so that combining Remarks 4.74 and Remark 6.4 with the results in the previous section one obtains isomorphisms

$$X \coprod^{a} Y \simeq \operatorname{colim}_{\Delta^{op}} \left( \operatorname{Lan} \left( N^{\bullet + 1} \iota X \coprod^{a} N^{\bullet + 1} \iota Y \right) \right) \tag{6.28}$$

$$\simeq \operatorname{colim}_{\Delta^{op}} \left( \operatorname{Lan} \left( N \circ \coprod \circ \left( N^{\times 2} \right)^{\bullet} (\iota X, \iota Y) \right) \right) \tag{6.29}$$

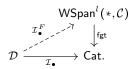
$$\simeq \operatorname{colim}_{\Delta^{op}} \left( \operatorname{Lan}_{\Omega_{G, \bullet}^{2, op} \to \Sigma_{G}^{op}} N_{\bullet}^{(X, Y)} \right) \tag{6.30}$$

where we write  $N_{\bullet}^{(X,Y)}:\Omega_{G,\bullet}^{2,op}\to\mathcal{V}$  for the induced functor.

The purpose of this section will be show that one can repackage formulas such as (6.30) with a single left Kan extension over a category  $\Omega_{G}^{2}=|\Omega_{G,\bullet}^{2}|$  obtained from  $\Omega_{G,\bullet}^{2}$  via realization. tion in Cat.

We note that  $\Omega_{G,\bullet}^2$  together with the corresponding functors to  $\Sigma_G$ ,  $\mathcal{V}^{op}$  can be viewed as a simplicial object  $\Delta^{op} \to \mathsf{WSpan}^l(\Sigma, G^{op}, \mathcal{V})$ , and our first task will be to repackage such functors in terms of Grothendieck constructions.

**Lemma 6.31.** Functors  $F: \mathcal{D} \ltimes \mathcal{I}_{\bullet} \to \mathcal{C}$  are in bijection with lifts



where fgt is the functor forgetting the maps to \* and C.

*Proof.* This is a matter of unpacking notation. The restrictions  $F|_{\mathcal{I}_d}$  to the fibers  $\mathcal{I}_d \subset D \ltimes \mathcal{I}_{\bullet}$ are precisely the functors  $\mathcal{I}_d^F : \mathcal{I}_d \to \mathcal{C}$  describing  $\mathcal{I}_{\bullet}^F(d)$ .

Furthermore, the images  $F((d,i) \to (d',f_*(i)))$  of the pushout arrows over a fixed arrow  $f: d \to d'$  of  $\mathcal{D}$  assemble to a natural transformation

$$\begin{array}{ccc}
\mathcal{I}_{d} & & & \\
I_{d}^{F} & & & \\
I_{d'} & & & \\
\mathcal{I}_{d'} & & & \\
\end{array}$$
(6.32)

which describes  $\mathcal{I}_{\bullet}^{F}(f)$ . It is straightforward to check that the associativity and unitality conditions coincide.

In the cases of interest we will have  $\mathcal{D} = \Delta^{op}$ , so that  $\mathcal{I}_{\bullet}$  can be interpreted as an object  $\mathcal{I}_{\bullet} \in \mathsf{Cat}^{\Delta^{op}}$ . By recalling the standard cosimplicial object  $[\bullet] \in \mathsf{Cat}^{\Delta}$  given by  $[n] = (0 \to 1 \to \cdots \to n)$  one obtains the following definition.

**Definition 6.33.** The left adjoint

$$|-|: \mathsf{Cat}^{\Delta^{op}} \rightleftarrows \mathsf{Cat}: (-)^{[\bullet]}$$

will be called the *realization* functor.

Remark 6.34. More explicitly, one has

$$|\mathcal{I}_{\bullet}| = coeq \left( \prod_{[n] \to [m]} [n] \times \mathcal{I}_m \Rightarrow \prod_{[n]} [n] \times \mathcal{I}_n \right). \tag{6.35}$$

**Example 6.36.** Any  $\mathcal{I} \in \mathsf{Cat}$  induces objects  $\mathcal{I}, \mathcal{I}_{\bullet}, \mathcal{I}^{[\bullet]} \in \mathsf{Cat}^{\Delta^{op}}$  where  $\mathcal{I}$  is the constant simplicial object and  $\mathcal{I}_{\bullet}$  is the nerve  $N\mathcal{I}$  with each level regarded as a discrete category. It is straightforward to check that  $|\mathcal{I}| = |\mathcal{I}_{\bullet}| = |\mathcal{I}^{[\bullet]}| = \mathcal{I}$ .

**Lemma 6.37.** Given  $\mathcal{I}_{\bullet} \in \mathsf{Cat}^{\Delta^{op}}$  one has an identification  $ob(|\mathcal{I}_{\bullet}|) \simeq ob(\mathcal{I}_{0})$ . Furthermore, the arrows of  $|\mathcal{I}_{\bullet}|$  are generated by the image of the arrows in  $\mathcal{I}_{0} \simeq \mathcal{I}_{0} \times [0]$  and the image of the arrows in  $[1] \times ob(\mathcal{I}_{1})$ .

For each  $i_1 \in \mathcal{I}_1$ , we will denote the arrow of  $|\mathcal{I}_{\bullet}|$  induced by the arrow in  $[1] \times \{i_1\}$  by

$$d_1(i_1) \xrightarrow{i_1} d_0(i_1).$$

Proof. We write  $d_{\hat{k}}$ ,  $d_{\hat{k},\hat{l}}$  for the simplicial operators induced by the maps  $[0] \xrightarrow{0 \mapsto k} [n]$ ,  $[1] \xrightarrow{0 \mapsto k, 1 \mapsto l} [n]$  which carbinformally be thought of as the "composite of all faces other than  $d_k$ ,  $d_l$ ". Using (6.35) one has equivalence relations of objects

$$[n] \times \mathcal{I}_n \ni (k, i_n) \sim (0, d_{\hat{k}}(i_n)) \in [0] \times \mathcal{I}_0$$

and since for any generating relation  $(k, i_n) \sim (l, i'_m)$  it is  $d_{\hat{k}}(i_n) = d_{\hat{l}}(i'_m)$  the identification  $ob(|\mathcal{I}_{\bullet}|) \simeq ob(\mathcal{I}_{0})$  follows.

To verify the claim about generating arrows, note that any arrow of  $[n] \times \mathcal{I}_n$  factors as

$$(k, i_n) \rightarrow (l, i_n) \xrightarrow{I_n} (l, i'_n)$$
 (6.38) FACTORIZATIONREAL EQ

for  $I_n: i_n \to i'_n$  an arrow of  $\mathcal{I}_n$ . The  $d_{\hat{l}}$  relation identifies the right arrow in (6.38) with  $(0, d_{\hat{l}}(i_n)) \xrightarrow{d_{\hat{l}}(I_n)} (0, d_{\hat{l}}(i'_n))$  in  $[0] \times \mathcal{I}_0$  while (if k < l) the  $d_{\hat{k},\hat{l}}$  relation identifies the left arrow with  $(0, d_{\hat{k},\hat{l}}(i_n)) \to (1, d_{\hat{k},\hat{l}}(i_n))$  in  $[1] \times \mathcal{I}_1$ . The result follows.

**Remark 6.39.** Given  $\mathcal{I}_{\bullet} \in \mathsf{Cat}^{\Delta^{op}}$ ,  $\mathcal{C} \in \mathsf{Cat}$ , the isomorphisms

$$Hom_{\mathsf{Cat}}(|\mathcal{I}_{ullet}|,\mathcal{C}) \simeq Hom_{\mathsf{Cat}^{\Delta^{op}}}(\mathcal{I}_{ullet},\mathcal{C}^{[ullet]})$$

together with the fact that  $\mathcal{C}^{[\bullet]}$  is always 2-coskeletal show that  $|\mathcal{I}_{\bullet}|$  is determined by the categories  $\mathcal{I}_0, \mathcal{I}_1, \mathcal{I}_2$  and maps between them, i.e. by the truncated version of formula (6.35) with  $n, m \leq 2$ .

Indeed, it can be shown that a sufficient set of generating relations in  $|\mathcal{I}_{\bullet}|$  is given by (i) the relations in  $\mathcal{I}_{0}$  (including relations stating that identities of  $\mathcal{I}_{0}$  are identities of  $|\mathcal{I}_{\bullet}|$ );

(ii) relations stating that for each  $i_0 \in \mathcal{I}_0$  the arrow  $i_0 = d_1(s_0(i_0)) \xrightarrow{s_0(i_0)} d_1(s_0(i_0)) = i_0$  is an identity; (iii) for each arrow  $I_1 : i_1 \to i'_1$  in  $\mathcal{I}_1$  the relation that the square below commutes

$$egin{aligned} d_1(i_1) & \stackrel{i_1}{\longrightarrow} d_0(i_1) \ & d_1(I_1) igg| & \int_{d_0(I_1)} d_0(I_1) \ & d_1(i_1') & \stackrel{i_1'}{\longrightarrow} d_0(i_1') \end{aligned}$$

and (iv) for each object  $i_2 \in \mathcal{I}_2$  the relation that the following triangle commutes.

$$d_{1,2}(i_2) \xrightarrow[d_2(i_2)]{d_1(i_2)} d_{0,1}(i_2)$$

$$d_{0,2}(i_2) \xrightarrow[d_0(i_2)]{d_0(i_2)}$$

**Example 6.40.** For  $\Omega_{G,\bullet}$  the simplicial object of planar strings one has  $|\Omega_{G,\bullet}| = \Omega_G^t$ , the category of G-trees and tall maps. Indeed, arrows of  $\Omega_{G,0}$  and objects of  $\Omega_{G,1}$  are naturally identified with the quotient arrows and planar tall arrows of  $\Omega_G^t$ , which are a generating set of arrows. And likewise, relations in  $\Omega_{G,0}$ , arrows in  $\Omega_{G,1}$  and objects in  $\Omega_{G,2}$  are identified with the relations of  $\Omega_G^t$ .

Analogously, for  $\Omega^{J}_{G,\bullet}$  the simplicial object of planar  $\underline{l}$ -labeled strings that are ( $\{l\} - J$ )-inert, one has  $|\Omega^{J}_{G,\bullet}| = \Omega^{J,t}_{G}$ , the category of  $\underline{l}$ -labeled G-trees and ( $\{l\} - J$ )-inert tall maps.

The following is the key result in this section.

**Proposition 6.41.** Let  $\mathcal{I}_{\bullet} \in \mathsf{Cat}^{\Delta^{op}}$ . Then there is a natural functor

$$\Delta^{op} \ltimes \mathcal{I}_{\bullet} \xrightarrow{s} |\mathcal{I}_{\bullet}|. \tag{6.42}$$

Further, s is final.

PLANARSTRING EX

SOURCEFINAL PROP

**Remark 6.43.** The s in the result above stands for *source*. This is because, for any  $\mathcal{I} \in \mathsf{Cat}$ , the map  $\Delta^{op} \ltimes \mathcal{I}^{[\bullet]} \to |\mathcal{I}^{[\bullet]}| \simeq \mathcal{I}$  is given by  $s(i_0 \to \cdots \to i_n) = i_0$ .

*Proof.* Recall that  $|\mathcal{I}_{\bullet}|$  is the coequalizer (6.35). Given  $(k, g_m) \in [n] \times \mathcal{I}_m$ , we will write  $[k, g_m]$  for the corresponding object in  $|\mathcal{I}_{\bullet}|$ . To simplify notation, we will write objects of  $\mathcal{I}_n$  as  $i_n$  and implicitly assume that  $[k, i_n]$  refers to the class of the object  $(k, i_n) \in [n] \times \mathcal{I}_n$ .

We define s on objects by  $s([n], i_n) = [0, i_n]$  and on an arrow  $(\phi, I_m): (n, i_n) \to (m, i'_m)$ as the composite (note that  $\phi: [m] \to [n]$  and  $I_m: \phi^*(i_n) \to i_m$ )

$$[0, i_n] \to [\phi(0), i_n] = [0, \phi^*(i_n)] \xrightarrow{I_m} [0, i'_m].$$
 (6.44)

TARGETDEFINITON EQ

To check associativity, the cases of of a pair of either two fiber arrows (i.e. arrows where  $\phi$  is the identity) or two pushforward arrows (i.e. arrows where  $I_m$  is the identity) are immediate from (6.44), hence we are left with the case  $([n], i_n) \xrightarrow{I_n} ([n], i'_n) \rightarrow ([m], \phi^*(i'_n))$  of a fiber arrow followed by a pushforward arrow. Noting that in  $\Delta^{op} \ltimes \mathcal{I}_{\bullet}$  this composite can be rewritten as  $([n], i_n) \to ([m], \phi^*(i_n)) \xrightarrow{\phi^*(I_n)} ([m], \phi^*(i'_n))$  this amounts to checking that

$$[0, i_n] \longrightarrow [\phi(0), i_n)] = [0, \phi^*(i_n)]$$

$$I_n \downarrow \qquad \qquad \downarrow^{\phi^*(I_n)}$$

$$[0, i'_n] \longrightarrow [\phi(0), i'_n] = [0, \phi^*(i_n)]$$

$$(6.45)$$

commutes in  $|\mathcal{I}_{\bullet}|$ , which is the case since the left square is encoded by a square in  $[n] \times \mathcal{I}_n$ and the right square is encoded by an arrow in  $[m] \times \mathcal{I}_n$ .

We now turn to showing that s is final.

We now turn to showing that s is inflat. Fix  $j \in \mathcal{I}_0$ . We will show that  $[0,j] \downarrow \Delta^{op} \ltimes \mathcal{I}_{\bullet}$  is indeed connected. By Lemma 6.37 any object in this undercategory has a description (not necessarily unique) as a pair

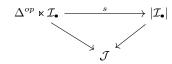
$$\left(\left([n],i_n\right),\left[0,j\right]\xrightarrow{f_1}\cdots\xrightarrow{f_r}s\left([n],i_n\right)\right)$$

where each  $f_i$  is a generating arrow of  $|\mathcal{I}_{\bullet}|$  induced by either an arrow  $I_0$  of  $\mathcal{I}_0$  or object

We will connect this object to the canonical object (([0], h), [0, h] = [0, h]), arguing by induction on r. If  $n \neq 0$ , the map  $d_{\hat{0}}:([n],i_n) \rightarrow ([0],d_{\hat{0}}^*(i_n))$  and the fact that  $s(d_{\hat{0}}^*)$   $id_{[0,d_n^*(i_n)]}$  provides an arrow to an object with n=0 without changing r. If n=0, one can apply the induction hypothesis by lifting  $f_r$  to  $\Delta^{op} \times \mathcal{I}_{\bullet}$  according to one of two cases: (i) if  $f_r$  is induced by an arrow  $I_0$  of  $\mathcal{I}_0$ , the lift of  $f_r$  is simply  $([0], i'_0) \xrightarrow{I_0} ([0], i_0)$ ; (ii) if  $f_r$  is induced by  $i_1 \in \mathcal{I}_1$  the lift is provided by the map  $([1], i_1) \to ([0], d_0(i_1))$ .

In practice, we will need to know that s satisfies the following stronger finality condition with respect to left Kan extensions.

Corollary 6.46. Consider a map  $\mathcal{I}_{\bullet} \to \mathcal{J}$  between  $\mathcal{I}_{\bullet} \in \mathsf{Cat}^{\Delta^{op}}$  and a constant object  $\mathcal{J} =$  $\mathcal{J}_{\bullet} \in \mathsf{Cat}^{\Delta^{\circ p}}$ . Then the source map s

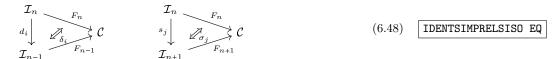


is Lan-final over  $\mathcal{J}$ , i.e. the functors  $s\downarrow j: (\Delta^{op} \ltimes \mathcal{I}_{\bullet}) \downarrow j \to |\mathcal{I}_{\bullet}| \downarrow j$  are final for all  $j\in \mathcal{J}$ .

Is Lan-jinut over J, i.e. the junctions  $3 \downarrow J$  ( $\Delta$  vertex J) while Lemma 2.2 guarantees that since J is clear that  $(\Delta^{op} \ltimes \mathcal{I}_{\bullet}) \downarrow j \simeq \Delta^{op} \ltimes (\mathcal{I}_{\bullet} \downarrow j)$  while Lemma 2.2 guarantees that since J is clear that J is clear t (-)  $\downarrow j$  is a left adjoint,  $|\mathcal{I}_{\bullet}| \downarrow j \simeq |\mathcal{I}_{\bullet}| \downarrow j|$ . One thus reduces to Proposition 6.4.

We end this section with two basic lemmas that will allows us to apply Corollary 6.46 to the tree categories we will be interested in.

**Lemma 6.47.** Let  $\mathcal{I}_{\bullet}^F \in \text{Span}(*,\mathcal{C})^{\Delta^{op}}$  be such that the diagrams



commute up to isomorphism for  $0 < i \le n, 0 \le j \le n$ .

Then the functors  $\tilde{F}_n: \mathcal{I}_n \to \mathcal{C}$  given by the composites

$$\mathcal{I}_n \xrightarrow{d_{1,\cdots,n}} \mathcal{I}_0 \xrightarrow{F_0} \mathcal{C}$$

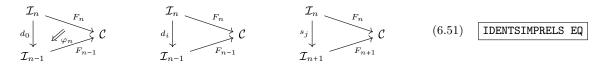
assemble to an object  $\mathcal{I}_{\mathtt{NDENTSIMPRELSISO}}^{\tilde{F}}$  Explicitly such that is isomorphic to  $\mathcal{I}_{\bullet}^{F}$  and such that the corresponding diagrams (6.48) for  $0 < i \le n$ ,  $0 \le j \le n$  are strictly commutative.

*Proof.* This follows by a straightforward verification.

Lemma 6.49. A (necessarily unique) factorization

$$\Delta^{op} \ltimes \mathcal{I}_{\bullet} \xrightarrow{s} |_{\mathcal{I}_{\bullet}}| \mathcal{C}$$
 (6.50) SOURCEFACT EQ

exists iff for the associated object  $\mathcal{I}_{\bullet} \in \mathsf{Span}(*,\mathcal{C})^{\Delta^{op}}$  (cf. Lemma SIMPSPANREIN LEMMA for the associated object  $\mathcal{I}_{\bullet} \in \mathsf{Span}(*,\mathcal{C})^{\Delta^{op}}$  $0 < i \le n$  and degeneracies  $s_j$  for  $0 \le j \le n$  are strictly commutative, i.e. they are given by diagrams



SOURCEFACT LEM

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*Proof.* For the "if" direction, it suffices to note that s sends all pushout arrows of  $\Delta^{op} \times \mathcal{I}_{\bullet}$  for faces  $d_i$ ,  $0 \le j \le n$  and degeneracies  $s_j$ ,  $0 \le j \le n$  to identities and this yields the commutative diagrams (6.51)

For the "only if" direction, this will follow by building a functor  $\mathcal{I}_{\bullet} \xrightarrow{F} \mathcal{C}^{[\bullet]}$  together with the naturality of the source map s (recall that  $|\mathcal{C}^{[\bullet]}| \simeq \mathcal{C}$ ). We define  $\bar{F}_n|_{k\to k+1}$  as the map

$$F_{n-k}d_{0,\dots,k-1} \xrightarrow{\varphi_{n-k}d_{0,\dots,k-1}} F_{n-k-1}d_{0,\dots,k}.$$
 (6.52)

EQUIVALENCEDEF EQ

The claim that  $s \circ (\Delta^{op} \ltimes \bar{F})$  recovers the horizontal map in (6.50) is straightforward, hence the real task is to prove that (6.52) indeed defines a map of simplicial objects.

$$\varphi_{n-1}d_i = \varphi_n, \ 1 < i \qquad \varphi_{n-1}d_1 = (\varphi_{n-1}d_0) \circ \varphi_n, \qquad \varphi_{n+1}s_i = \varphi_n, \ 0 < i, \qquad \varphi_{n+1}s_0 = id_{F_n}$$
(6.53)

Next, note that there is no ambiguity in writing simply  $\varphi_{n-k}d_{0,\cdots,k-1}$  to denote the map (6.52). We now check that  $\bar{F}_{n-1}d_i=d_i\bar{F}_n$ ,  $0 \le i \le n$ , which must be verified after restricting to each  $k \to k+1$ ,  $0 \le k \le n-2$ . There are three cases, depending on i and k:

$$(i < k+1) \varphi_{n-k-1} d_{0,\dots,k-1} d_i = \varphi_{n-k-1} d_{0,\dots,k};$$

$$\begin{array}{lll} (i=k+1) & \varphi_{n-k-1} d_{0,\cdots,k-1} d_i &= \varphi_{n-k-1} d_1 d_{0,\cdots,k-1} &= (\varphi_{n-k-1} d_0 \circ \varphi_{n-k}) d_{0,\cdots,k-1} &= (\varphi_{n-k-1} d_{0,\cdots,k}) \circ (\varphi_{n-k} d_{0,\cdots,k-1}); \end{array}$$

$$(i>k+1) \ \varphi_{n-k-1}d_{0,\cdots,k-1}d_i = \varphi_{n-k-1}d_{i-k}d_{0,\cdots,k-1} = \varphi_{n-k}d_{0,\cdots,k-1}.$$

The case of degeneracies is similar.

DUALRESULTS REM

Remark 6.54. One can twist all results by the opposite functor

$$\Delta \xrightarrow{(-)^{op}} \Delta$$

which sends [n] to itself and  $d_i, s_i$  to  $d_{n-i}, s_{n-i}$ . In doing so, one obtains vertical isomorphisms

$$\begin{array}{ccc} \Delta^{op} \ltimes (\mathcal{J}_{\bullet} \circ (\mathsf{-})^{op}) & \stackrel{s}{\longrightarrow} |\mathcal{J}_{\bullet} \circ (\mathsf{-})^{op}| \\ & \downarrow & \downarrow \simeq \\ & \Delta^{op} \ltimes \mathcal{J}_{\bullet} & \xrightarrow{t} |\mathcal{J}_{\bullet}| \end{array}$$

which reinterpret the "source" functor as what one might call the "target" functor, with  $t([n], i_n) = [n]$  rather than  $s([n], i_n) = [0, i_n]$ .

Corollary 6.46 now says that t is Lan-final and Lemmas 6.47, 6.49 generalize in the

obvious way by replacing s with t and  $d_0$  with  $d_n$ .

#### The category of extension trees 6.5

In this section we combine the previous sections to obtain a compact description of free extension pushouts

$$\begin{array}{ccc}
\mathbb{F}_G A & \longrightarrow & X \\
\downarrow & & \downarrow \\
\mathbb{F}_G B & \longrightarrow & Y
\end{array} (6.55) \quad \boxed{\text{FREEEXT EQ}}$$

as a left Kan extension over a convenient category of trees.

For simplicity, we first explain how to obtain a similar description for the simpler case of a coproduct  $X \coprod^a Y$ . By (6.30), one has a description

$$\begin{split} X & \amalg^a Y \simeq \mathrm{colim}_{\Delta^{op}} \left( \mathsf{Lan}_{\Omega^{2,op}_{G,\bullet} \to \Sigma^{op}_{G}} N^{(X,Y)}_{\bullet} \right) \\ & \simeq \mathsf{Lan}_{\Delta^{op} \ltimes \Omega^{2,op}_{G,\bullet} \to \Sigma^{op}_{G}} N^{(X,Y)}_{\bullet} \end{split}$$

where the second identification follows from formal properties of Grothendieck constructions. Combining the fact that (6.27) consists of natural isomorphisms with (the Remark 6.54  $\frac{\text{IWISTING LEMMA}}{\text{IWISTING LEMMA}} \text{dual of) Lemma } \begin{array}{l} \text{Lemma fo.47}, \text{ yields an isomorphic twisted functor } \tilde{N}_{\bullet}^{(X,Y)} \text{ with strictly commutative } \\ \text{SourceFact Lemma} \\ \text{s}_i \text{ and } d_i \text{ for } i \neq n. \text{ The dual of Lemma } \begin{array}{l} \text{G.49} \\ \text{O.2}, op \\ \text{opp} \end{array} \right] \left( \begin{array}{l} \text{CP}_i \text{ opp} \\ \text{Corollary } \end{array} \right) \left( \begin{array}{l} \text{Writing } \tilde{N}^{(X,Y)} \text{ for the factorization} \right) \text{ and thus the dual of } \\ \text{Corollary } \begin{array}{l} \text{G.46} \text{ finally yields} \end{array} \right)$ 

$$X \coprod^{a} Y \simeq \mathsf{Lan}_{\Omega_{G}^{2,op} \to \Sigma_{G}^{op}} \tilde{N}^{(X,Y)}.$$
 (6.56)

We recall that by Example  $\frac{\text{PLANARSTRING EX}}{6.40, \Omega_G^2 \text{ is simply}}$  the category of 2-labeled trees and tall label

More generally, one has

$$\coprod_{J}^{a} X_{j} \coprod_{l=J}^{a} \mathbb{F}_{G} X_{j} \simeq \operatorname{Lan}_{\Omega_{G}^{J,op} \to \Sigma_{G}^{op}} \tilde{N}^{(\underline{X})}. \tag{6.57}$$

where  $\Omega_G^J$  is the category of  $\underline{l}$ -labeled trees and tall  $(\underline{l} - J)$ -inert label maps.

**Remark 6.58.** We note that the twisting  $\tilde{N}^{(X,Y)}_{\bullet}$  is fairly harmless. For explicitness, we focus on the simplest case of a "unary coprodut", in which case (6.57) is simply recovering the genuine equivariant opgsad  $X_1$  from its bar resolution. In that  $CSSCSPANT COPP \rightarrow V$  is given by the top map in (4.53) or, equivalently, by the top map in (4.54) (we note that, in the notation therein, it is  $A = \Sigma_G$ ). On the other hand, the observable consistency  $\tilde{N}_2^X: \Omega_{G,2}^{op} \rightarrow V$  is given by the left bottom composite in either of (4.53), (4.54). Informally, the role of this twisting is therefore simply that of replacing the order on  $V_G(T_n)$  induced lexicographically by planar strings  $T_0 \to \cdots \to T_n$  with the simpler order induced directly from  $T_n$ .

In what follows we will largely be able to ignore this technicality. Indeed, the role of lexicographic orders in building (6.57) is that of guaranteeing that  $\tilde{N}_{\bullet}$  satisfies the necessary simplicial identities, which are ensured by appealing to the bar construction for the monad

We now turn to the task of building (6.55) as a left Kan extension. One has a colimit

$$\mathbb{F}B \coprod_{\mathbb{F}A} X \simeq \operatorname{colim}_{\Delta^{op}} \Big( \mathbb{F}B \sqcup \mathbb{F}A \sqcup X \not \sqsubseteq \mathbb{F}B \sqcup \mathbb{F}A \sqcup \mathbb{F}A$$

where all differentials are fold maps of  $\mathbb{F}A$  except to the *n*-th differential  $d_n$ , which is induced by the two maps  $\mathbb{F}A \to X$ ,  $\mathbb{F}A \to \mathbb{F}B$ .

the two maps  $\mathbb{F}A \to X$ ,  $\mathbb{F}A \to \mathbb{F}B$ . By the previous discussion each individual object  $X \amalg (\mathbb{F}A)^{\sqcup 2n+1} \amalg \mathbb{F}B$  in (6.59) can be described as a left Kan extension over the tree category  $\Omega_G^{\{X\}}$  where  $\{X\} \subset \{B,A,\cdots,A,X\}$  is a simpleton. The maps in (6.59) can themselves be encoded as span maps between the  $\Omega_G^{\{X\}}$ . To see this, we make (6.59) more precise. Firstly, we write  $\langle n \rangle$  for the poset

$$-\infty \le -n \le -n + 1 \le \dots \le -1 \le 0 \le 1 \le \dots \le n - 1 \le n \le +\infty$$
.

The posets (n) together with antisymmetric (i.e. such that f(-x) = -f(x)) poset maps preserving all three of  $-\infty, 0, +\infty$  then form a simplicial object  $(-): \Delta^{op} \to F$ . (6.55) thus induces a simplicial object  $(B, A, X)_{(n)} \in F \wr Fun(\Sigma_G^{op}, \mathcal{V})$ .

Each level of  $(\iota B, \iota A, \iota X)_{(n)}$  is then a  $N^{\times \{+\infty\}}$ -algebra on  $(\mathsf{WSpan}^l(\Sigma_G^{op}, \mathcal{V}))^{\times (n)}$ , compatibly with the simplicial maps. One thus obtains a bisimplicial object

$$\Sigma_{G}^{op} \leftarrow \Omega_{G,\bullet}^{\{+\infty\}_{\{\bullet\}},op} \xrightarrow{N_{\bullet}^{(B,A,X)_{\{\bullet\}}}} \mathcal{V}$$

<sup>&</sup>lt;sup>1</sup>Indeed, we recall that the opposite simplex category  $\Delta^{op}$  can equivalently described as the category of intervals, i.e. finite ordered posets with distinct top and bottom, along with order maps preserving both top and bottom. (n) can then be regarded as obtained by gluing the interval  $0 \le 1 \le \dots \le n \le +\infty$  with its opposite.

on WSpan  $(\Sigma_G^{op}, \mathcal{V})$  whose realization along the string direction yields the spans

$$\Sigma_G^{op} \leftarrow \Omega_G^{\{+\infty\}_{(\bullet)}, op} \xrightarrow{N^{(B, A, X)_{(\bullet)}}} \mathcal{V}$$

$$(6.60) \quad \boxed{\text{PARTREALSPAN EQ}}$$

discussed above, except now assembled into a simplicial object in WSpan  $(\Sigma_G^{op}, \mathcal{V})$ .

All degeneracies  $s_i$  and differentials  $d_i$  of  $(s_i^{op}, v_i^{op})$  other than the top differential  $d_n$  are induced by maps of described in SURCELANTINAL SCOR by strictly commutative diagrams, so that Lemma 6.47) so as to allow (6.59) to be repackaged as

$$\mathbb{F}B \coprod_{\mathbb{F}A} X \simeq \mathsf{Lan}_{\Omega_G^{e,op} \to \Sigma_G^{op}} N^{(B,A,X)} \tag{6.61}$$

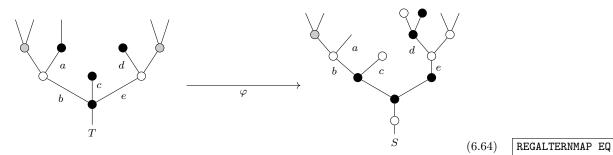
FREEEXTUSEFCOLNEW EQ

where we write  $\Omega_G^e$  for  $|\Omega_G^{\{+\infty\}_{\{\bullet\}}}|$ . We now turn to the task of describing  $\Omega_G^e$ , starting with by defining it directly.

**Definition 6.62.** The extension tree category  $\Omega_c^e$  is the category whose objects are  $\{B, A, X\}$ labeled trees and whose maps  $\varphi:T\to S$  are tall maps of trees such that

- (i) if  $T_{v_{Ge}}$  has an A-label, then  $S_{v_{Ge}}$  =  $T_{v_{Ge}}$  and  $S_{v_{Ge}}$  has an A-label;
- (ii) if  $T_{v_{Ge}}$  has a B-label, then  $S_{v_{Ge}} = T_{v_{Ge}}$  and  $S_{v_{Ge}}$  has either an A-label or a B-label;
- (iii) if  $T_{v_{Ge}}$  has a X-label, then  $S_{v_{Ge}}$  has only A and X-labels.

**Example 6.63.** The following is an example of a planar map in  $\Omega_G^e$ , where black nodes represent X-labeled nodes, grey nodes represent B-labeled nodes and white nodes represent A-labeled nodes.



**Proposition 6.65.** One has an identification

$$\Omega_G^e \simeq \big|\Omega_G^{\{+\infty\}_{\{\bullet\}}}\big|.$$

*Proof.* We note first that  $\Omega_G^e$  contains all label maps that are  $\{A,B\}$ -inert. In fact, any map of  $\Omega_G^e$  clearly has a unique factorization as such a label map followed by an underlying planar isomorphism of trees that replaces some of the X and B labels with A labels. We will refer to the former as label maps and to the latter as relabel maps.

We recall that  $\Omega_G^{\{+\infty\}_{\{n\}}}$  consists of trees with 2n+3 types of labels: X-labels, B-labels and 2n + 1 distinct types of A-labels. One can equivalently encode such a tree as a string  $T_0 \to \cdots \to T_n$  of relabel maps. Indeed, the A-label nodes of  $T_n$  in such a string are partitioned into 2n+1 types according to that node's labels one the  $T_i$  (which are either all A's, some X's and then A's or some B's and then A's). Moreover, a diagram

$$T_0 \longrightarrow T_1 \longrightarrow \cdots \longrightarrow T_n$$

$$f_0 \downarrow \qquad f_1 \downarrow \qquad \qquad f_n \downarrow$$

$$T'_0 \longrightarrow T'_1 \longrightarrow \cdots \longrightarrow T'_n$$

EXTTREECAT DEF

with  $f_i$  label maps of  $\Omega_G^e$  is then equivalent to a label map  $f_n: T_n \to T'_n$  respecting all 2n+3 labels in  $\Omega_G^{\{+\infty\}_{\{n\}}}$ . Since the string description above is also compatible with the simplicial structure maps in the obvious way, the result is now clear.

Our next task will be that of identifying a convenient Lan-final subcategory  $\bar{\Omega}_G^e \hookrightarrow \Omega_G^e$ . We first introduce the auxiliary notion of alternating trees. We recall the notion of input path (Notation 3.4)  $I(e) = \{f \in T : e \leq_d f\}$  for an edge  $e \in T$ , which naturally extends to T in any of  $\in \Omega$ ,  $\Phi$ ,  $\Omega_G$ ,  $\Phi_G$ .

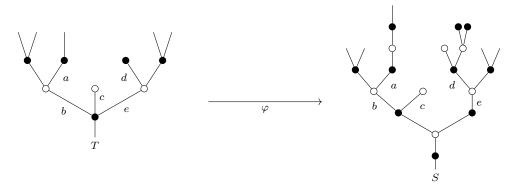
**Definition 6.66.** A G-tree  $T \in \Omega_G$  is called *alternating* if, for all leafs  $l \in T$  one has that the input path I(l) has an even number of elements.

Further, a vertex  $e^{\uparrow} \leq e$  is called *active* if |I(e)| is odd and *inert* otherwise.

Finally, a tall map  $T \xrightarrow{\varphi} T'$  between alternating G-trees is called a tall alternating map if for any inert vertex  $e^{\uparrow} \leq e$  of T one has that  $T'_{e^{\uparrow} \leq e}$  is an inert vertex of T'.

We will denote the category of alternating G-trees and tall alternating maps by  $\Omega_G^a$ .

**Example 6.67.** Two alternating trees (for G = \* the trivial group) and a planar tall alternating map between them follow, with active nodes in black ( $\bullet$ ) and white nodes in white ( $\circ$ ).



(6.68)

REGALTERNMAPLR EQ

The term "alternating" comes from the fact that no adjacent nodes have the same color. We note, however, that there is additional restriction: the "outer" vertices, i.e. those immediately below a leaf or the one immediately above the root, are necessarily black/active (not, however, that this does *not* apply to stumps).

Remark 6.69. One can extend Definition 3.39 to the alternating context by defining a substitution datum to be alternating if it is given by isomorphisms for inert nodes and by alternating maps for active nodes. It is then straightforward to check that Proposition 3.42 and its equivariant analogue Proposition 4.19 extend to give alternating analogues.

**Definition 6.70.**  $\bar{\Omega}_G^e \hookrightarrow \Omega_G^e$  is the full subcategory of (B, A, X)-labeled trees whose underlying trees is alternating, active nodes are labeled by X, and passive nodes are labeled by A or B.

We note that conditions (i) and (ii) in Definition 6.62 imply that maps in  $\bar{\Omega}_G^e$  are underlying alternating maps.

The following establishes the required finality of  $\bar{\Omega}_G^e$  in  $\Omega_G^e$ .

**Proposition 6.71.** For each  $U \in \Omega_G^e$  there exists a unique  $\operatorname{Ir}_X(U) \in \bar{\Omega}_G^e$  together with a unique planar label map of  $\Omega_G^e$ 

$$\operatorname{Ir}_X(U) \to U$$
.

Furthermore,  $\operatorname{Ir}_X : \operatorname{cxtends}$  to a right retraction  $\operatorname{Ir}_X : \Omega_G^e \to \bar{\Omega}_G^e$ .

LXP PROP

Proof. Given U, we form a collection of outer faces  $\{U_i^A\} \sqcup \{U_j^B\} \sqcup \{U_k^X\}$  where the  $U_i^A, U_j^b$  are simply the A, B-labeled nodes and the  $\{U_k^X\}$  are the maximal outer subtrees whose nodes have only X-labels (we note that these may possible be sticks). Lemma 3.51 then guarantees that the  $V_c(U_k^X)$  are disjoint, so that one can apply (the equivariant version of Proposition 3.49) to build

$$T = lr(U) \rightarrow U$$
 (6.72) LRXDEF EQ

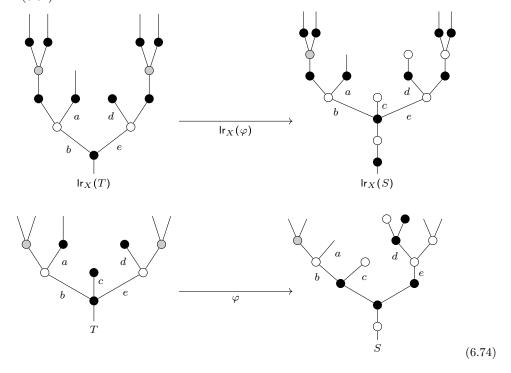
such that  $\{U_{vGe}\}=\{U_i^A\}$  if  $\{U_j^B\}$  if  $\{U_k^X\}$ . T has an obvious (B,A,X)-labeling making (6.72) into a label map, but we must still check  $T\in\bar{\Omega}_G^e$ , i.e. that T is alternating with the X-labeled vertices being precisely the X-labeled vertices. Let us now write any input path of T as  $I(e)=(e=e_n\leq e_{n-1}\leq\cdots\leq e_1\leq e_0)$ . By Lemma 5.51 and maximality of the  $U_k^X$ , no pair of consecutive vertices  $v_{G_{EH}}$  and  $v_{Ge_{i+1}}$  can be both X-labeled. On the other hand, again by Lemma 5.51 any edge of U belongs to some  $U_k^X$  and therefore: (i) at least one of in each pair of consecutive vertices  $v_{Ge_i}$  and  $v_{Ge_{i+1}}$  is X-labeled; (iii) if  $t\in T$  is a root,  $v_{Gr}$  is X-labeled; (iii) if  $t\in T$  is a leaf  $v_{Gl_{n-1}}$  is X-labeled. This suffices to conclude T  $\in \bar{\Omega}_G^e$ , and uniqueness of T is immediate from the uniqueness in Lemma 3.51.

It remains to check that  $lr_X$  in fact defines a functor. We consider the following diagram.

$$|\operatorname{Ir}_X(U) \longrightarrow U$$
 $|\operatorname{Ir}_X(f)| \downarrow \qquad \qquad \downarrow f$ 
 $|\operatorname{Ir}_X(V) \longrightarrow V$ 

When f is a root pullback map, we define  $\operatorname{Ir}_X(f)$  to likewise be a root pullback map. When f is a rooted tall map, writing  $T = \operatorname{Ir}_X(U)$  one has a map of rooted T-substitution data  $\{\operatorname{Ir}_X(V_{v_{G_e}})\} \to \{V_{v_{G_e}}\}$ , which after converted to a tree map yields the desired map  $\operatorname{Ir}_X(f)$ . To check that  $\operatorname{Ir}_X$  respects composition of maps, the only non immediate case is the root pullback followed by a rooted map, in which case this follows from Remark A. A.

**Example 6.73.** The following illustrates the  $lr_X$  construction when applied to the map  $\varphi$  in (6.64).



## 7 Filtration of cellular extensions

Put together, the results in the previous section show that the free extension  $\mathcal{P}[u]$  given by the pushout

$$\mathbb{F}_{G}X \longrightarrow \mathcal{P}$$

$$\downarrow \qquad \qquad \downarrow$$

$$\mathbb{F}_{G}Y \longrightarrow \mathcal{P}[u]$$

$$(7.1) \quad \boxed{\text{CELLEXTPUSH EQ}}$$

is given by a left Kan extension along  $(\bar{\Omega}_G^e)^{op} \xrightarrow{\operatorname{lr}} \Sigma_G^{op}$ . So as to study the homotopical properties of the map  $\mathcal{P} \to \mathcal{P}[u]$  we will identity a suitable filtration of this map, which will in turn be induced by a suitable filtration of the extension tree category  $\bar{\Omega}_G^e$ .

### 7.1 Filtration pieces

We now turn to the task of describing our filtration of  $\bar{\Omega}_{G}^{e}$ .

Firstly, we write  $V_X(T)$  (resp.  $V_Y(T)$ ) to denote the set of (non-equivariant) vertices of T with a X-label (resp. Y-label). We now define the *degree* of  $T \in \bar{\Omega}_G^e$ , denoted |T|, to be the sum  $|T|_X + |T|_Y$ , where  $|T|_X$ ,  $|T|_Y$  are defined by

$$|T|_X = \frac{|V_X(T)|}{|Gr|} = \sum_{Gv \in V_{G,X}(T)} \frac{|Gv|}{|Gr|}, \qquad |T|_Y = \frac{|V_Y(T)|}{|Gr|} = \sum_{Gv \in V_{G,Y}(T)} \frac{|Gv|}{|Gr|}$$

for Gr the root orbit of T.

Intuitively,  $|T|_X$  counts the number of X-labeled vertices in each individual tree component of T.

**Remark 7.2.** One of the key properties of the degrees just defined is that they are invariant under root pullback.

**Definition 7.3.** We define subcategories of  $\bar{\Omega}_G^e$ :

- $\bar{\Omega}_G^e[\leq k]$  (resp.  $\Omega_G^e[k]$ ) is the full subcategory of trees  $T \in \bar{\Omega}_G^e$  with  $|T| \leq k$  (resp. |T| = k);
- $\bar{\Omega}_G^e[\leq k, -]$  (resp.  $\bar{\Omega}_G^e[k, -]$ ) is the full subcategory of  $\bar{\Omega}_G^e[\leq k]$  (resp.  $\bar{\Omega}_G^e[k]$ ) of trees T with  $|T|_Y \neq k$ ;
- $\bar{\Omega}_G^e[k,0]$  is the full subcategory of  $\bar{\Omega}_G^e[k]$  of trees T with  $|T|_X = 0$  (or, equivalently,  $|T|_Y = k$ ).

The above definitions still hold if we replace  $\bar{\Omega}_G^e$  with  $\Omega_G^a$ ; in particular, we have vertical forgetful functors

$$\bar{\Omega}^e_G[k,-] \xrightarrow[\mathrm{fgt}]{} \bar{\Omega}^e_G[k]$$

**Remark 7.4.** The categories  $\bar{\Omega}_G^e[k]$  and  $\bar{\Omega}_G^e[k,-]$  have only rather limited morphisms. In fact, all maps in these categories must be underlying quotients of trees. Indeed, it is clear from Definition 6.62 that maps never lower degree and, moreover, degree is preserved iff  $\mathcal{P}$ -vertices are substituted by  $\mathcal{P}$ -vertices (rather than larger trees in  $\bar{\Omega}_G^e$ , which would necessarily possess X-vertices).

Moreover, we have a clear isomorphism of categories  $\bar{\Omega}_G^e[k,0] \simeq \Omega_G^a[k]$ .

**Lemma 7.5.**  $\bar{\Omega}_G^e[\leq k-1]$  is Ran-initial in  $\bar{\Omega}_G^e[\leq k,-]$  over  $\Sigma_G$ .

In the proof we will make use of the following construction on  $\Omega_{G,e}$ : given  $T \in \Omega_{G,e}$  we will let  $T_{\mathcal{P}}$  denote the result of replacing all X-labeled nodes of T with  $\mathcal{P}$ -labeled nodes.

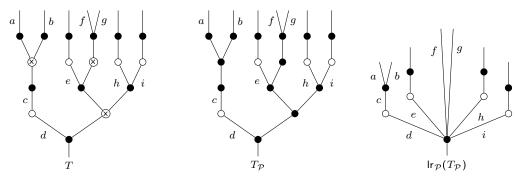
LIMMOR REM

\_LAN\_FINAL\_LEMMA

YINERT REM

**Remark 7.6.** Unlike the  $Ir_{\mathcal{P}}$  construction of Proposition 6.71, which defines a functor  $\operatorname{Ir}_{\mathcal{P}}:\Omega_G^e \to \bar{\Omega}_G^e$ , the construction  $(-)_{\mathcal{P}}$  does not define a full functor  $\Omega_G^e \to \Omega_G^e$ , instead being functorial, and the obvious maps  $T_{\mathcal{P}} \to T$  natural, only with respect to the Y-inert maps of  $\Omega_G^e$ .

**Example 7.7.** Combining the  $(-)_{\mathcal{P}}$  and  $Ir_{\mathcal{P}}$  constructions one obtains a construction sending trees in  $\bar{\Omega}_G^e$  to trees in  $\bar{\Omega}_G^e$ . We illustrate this for the tree  $T \in \bar{\Omega}_G^e$  below, where black nodes are  $\mathcal{P}$ -labeled, white nodes filled with  $\times$  are X-labeled, and empty white nodes are Y-labeled.



Proof of Lemma 7.5. Just as in the proof of Lemma 4.67, for each  $C \in \Sigma_G$ , the under categories  $C \downarrow \bar{\Omega}_G^e[\leq k-1], C \downarrow \bar{\Omega}_G^e[\leq k,-]$  have initial subcategories  $C \downarrow_{r,\simeq} \bar{\Omega}_G^e[\leq k-1].$  $C\downarrow_{\mathsf{r},\simeq}\bar{\Omega}^e_G[\leq k,-]$  of those objects  $(S,q:C\to\mathsf{lr}(S))$  such that q is an ordered isomorphism on roots, and thus an isomorphism in  $\Sigma_G$ . It now suffices to show (cf. ([4, X.3.1])) that for each  $(S, q: C \to \mathsf{Ir}(S))$  in  $C \downarrow_{\mathsf{r},\simeq} \bar{\Omega}^e_G[\leq k, -]$ 

the undercategory

$$(S,q) \downarrow (C \downarrow_{\mathsf{r},\simeq} \bar{\Omega}_G^e [\leq k-1]) \tag{7.8}$$

is non-empty and connected. Moreover, we note that an object in (7.8) is uniquely encoded by a map  $T \to S$  inducing a rooted isomorphism on Ir.

The case  $S \in \Omega_G^e[\leq k-1]$  is immediate. Otherwise, since  $|S|_Y \neq k$  it is  $|S|_Y \neq k$  and the map  $|S|_Y \neq k$  it is immediate. Otherwise, since  $|S|_Y \neq k$  it is immediate. Otherwise, since  $|S|_Y \neq k$  it is immediate.

Otherwise, given any rooted tall map  $T \to S$  with  $T \in \bar{\mathbb{Q}}_{p,p}^e[k_{p,p}] \to \mathbb{Q}_p^e[k_{p,p}]$  (which gives a rooted isomorphism on Ir and thus encodes a unique object of (7.8)). One can then form a diagram

$$S \longleftarrow \operatorname{Ir}_{\mathcal{P}}(S_{\mathcal{P}})$$

$$\uparrow_{Y-\operatorname{inert}} \uparrow \qquad (7.9) \quad \boxed{\mathsf{K-1LANFINAL EQ}}$$

$$T \longrightarrow T' \longleftarrow \operatorname{Ir}_{\mathcal{P}}(T'_{\mathcal{P}})$$

UNDERCATPR EQ

where  $T \to T' \to S$  is the natural factorization such that the second map is Y-inert, i.e., T' is obtained from T by simply relabeling to X those Y-labeled vertices of T that become X-vertices in S. Note that the existence of the right square in (7.9) follows from the map  $T' \to S$  being Y-inert together with Remark 7.6. Since (7.9) becomes a diagram of rooted isomorphism on  $\operatorname{Ir}$  if produces the necessary zigzag connecting the objects  $T \to S$  and  $\operatorname{Ir}_{\mathcal{P}}(S_{\mathcal{P}}) \to S$  in (7.8), finishing the proof.

Similarly to the  $(-)_{\mathcal{P}}$  construction, there is also a construction  $T_Y$  which replaces all X-labels of  $T \in \Omega_G^e$  with Y-labels. Moreover, in this case the construction restricts directly to a construction on  $\bar{\Omega}_G^e$ , which is easily seen to be functorial (and the  $T_Y \to T$  maps natural) with regards to  $\mathcal{P}$ -inert maps. Remark 7.4 thus implies that  $(-)_Y: \bar{\Omega}_G^e[k] \to \bar{\Omega}_G^e[k,0]$  is a left retraction, resulting in the following.

**Lemma 7.10.**  $\bar{\Omega}_G^e[k,0]$  is Ran-initial in  $\bar{\Omega}_G^e[k]$  over  $\Sigma_G$ .

N\_FINALITY\_LEMMA

In what follows we write  $N^e : \Omega_G^{e,op} \to \mathcal{V}$  for the functor in (6.61), and abuse notation by likewise writing  $N^e$  for any of its restrictions to the subcategories in Definition 7.3.

We are now in a position to produce the desired filtration of the map  $\mathcal{P} \to \mathcal{P}[u]$  in (7.1).

**Definition 7.11.** Let  $\mathcal{P}_k$  denote the left Kan extension

Noting that  $\bar{\Omega}_G^e[\leq 0] \simeq \Sigma_G$  (since |T| = 0 only if T is a G-corolla with  $\mathcal{P}$ -labeled vertex) and that  $\bar{\Omega}_G^e$  is the union of (the nerves of) the  $\bar{\Omega}_G^e[\leq k]$ , one has a filtration

$$\mathcal{P} = \mathcal{P}_0 \to \mathcal{P}_1 \to \mathcal{P}_2 \to \cdots \to \operatorname{colim}_k \mathcal{P}_k = \mathcal{P}[u]. \tag{7.12}$$

To analyze (FILT EQ (7.12) homotopically we will further make use of a pushout description of each individual map  $\mathcal{P}_{k-1} \to \mathcal{P}_k$ . To do so, we note that the diagram of inclusions

$$\begin{split} \bar{\Omega}^e_G[k,-] &\longrightarrow \bar{\Omega}^e_G[\leq k,-] \\ & \downarrow & \downarrow \\ \bar{\Omega}^e_G[k] &\longrightarrow \bar{\Omega}^e_G[\leq k] \end{split} \tag{7.13}$$

is a pushout of at the level of nerves. Indeed, this follows since

$$\bar{\Omega}_G^e[k] \cap \bar{\Omega}_G^e[\leq k, -] = \bar{\Omega}_G^e[k, -], \qquad \bar{\Omega}_G^e[k] \cup \bar{\Omega}_G^e[\leq k, -] = \bar{\Omega}_G^e[\leq k],$$

and since a map  $T_{\underline{\mathsf{MINUS}}} = \underline{\mathsf{LAN}}_{G} = [\underbrace{\mathsf{k}}_{\underline{\mathsf{LNMM}}}]_{\underline{\mathsf{LEMMA}}}$  be in one of subcategories in ([7.13] iff T is. Since Lemma [7.5] provides an identification  $\underline{\mathsf{Lan}}_{\bar{\Omega}_{G,e}} = \underline{\mathsf{Lan}}_{\bar{\Omega}_{G,e}} = \underline{\mathsf{Lan}}_$ 

$$\operatorname{Lan}_{\bar{\Omega}_{G,e}[k,-]^{op}} N^{e} \longrightarrow \mathcal{P}_{k-1} 
\downarrow \qquad \qquad \downarrow \qquad \qquad (7.14) \quad \boxed{\text{FILTRATION\_LAN\_SQUARE}} 
\operatorname{Lan}_{\bar{\Omega}_{G,e}[k]^{op}} N^{e} \longrightarrow \mathcal{P}_{k}$$

HERE

PK\_DEFN

#### 7.2 Notation

#### Luis: should this be stated earlier when defining the categorical wreath products?

In order to state our filtration result, we will need to identify another categorical construction. This filtration will be built out of "pushout products over trees of maps of sequences". This subsection is dedicated to making that statement precise.

**Definition 7.15.** Given a map  $y_{\widetilde{D}_G} Y_0 \to Y_1$  of G-symmetric sequences  $\mathcal{V}^{\Sigma_G^{op}}$ , and  $(A, D) \in F \wr \Sigma_G$ , we borrow notation from [1] and define the functor

$$[u]^D: (0 \to 1)^A \to \mathcal{V}$$

as the composite

$$(0 \to 1)^A \to \mathsf{F} \wr \mathcal{V} \xrightarrow{\times} \mathcal{V}$$

where the first map is defined on  $\xi: A \to \{0,1\}$  by

$$(a \mapsto \xi(a)) \mapsto (A, (a \mapsto Y_{\xi(a)}(D(a))))$$

Q\_DEFINITION

LTRATION\_THEOREM

**Definition 7.16.** Given  $u: Y_0 \to Y_1$  and  $(A, D) \in \mathsf{F} \wr \Sigma_G$  as before, let  $\mathsf{pc}(A)$  denote the "punctured cube" subcategory

$$pc(A) := (0 \to 1)^A \setminus \{(1)_a\} \subseteq (0 \to 1)^A,$$

and define

$$Q[u]^D := \operatorname{colim}_{pc(A)}[u]^D$$
,

and

$$[u]^{\square D}: Q[u]^D \to \bigotimes_{a \in A} Y_1(D(a))$$

as the natural map.

#### 7.3 Filtration Result

We can now state our filtration of the cellular extension  $\mathcal{P} \to \mathcal{P}[u]$ .

**Theorem 7.17.** Let  $\mathcal{P}$  be a genuine G-operad, and suppose we are given a map of G-symmetric sequences  $u: Y_0 \to Y_1$ . Then we have a filtration in G-sequences of the cellular extension

$$\mathcal{P} = \mathcal{P}_0 \to \mathcal{P}_1 \to \ldots \to \operatorname{colim}(\mathcal{P}_i) = \mathcal{P}[u],$$

where  $\mathcal{P}_{k-1} \to \mathcal{P}_k$  is given by the pushout

$$\operatorname{Lan}_{\bar{\Omega}_{G}^{e}[k,-]^{op}} N^{e} \longrightarrow \mathcal{P}_{k-1}$$

$$\downarrow \qquad \qquad \downarrow$$

$$\operatorname{Lan}_{\bar{\Omega}_{G}^{e}[k]^{op}} N^{e} \longrightarrow \mathcal{P}_{k}$$

where each Lan is over the leaf-root functor associated to the given category.

Levelwise, for each  $C \in \Sigma_G$ , we have a filtration of  $\mathcal{P}(C) \to \mathcal{P}[u](C)$  in  $\mathcal{V}^{\operatorname{Aut}(C)}$ , where  $\mathcal{P}_{k-1}(C) \to \mathcal{P}_k(C)$  is given by the pushout

$$\coprod_{\substack{[T] \in \Omega_G^a[k](C)/\simeq}} \left( \bigotimes_{v \in V_{G,\mathcal{P}}(T)} \mathcal{P}(T_v) \otimes Q[u]^{\mathbb{V}_{G,in}(T)} \right) \otimes_{\operatorname{Aut}(T)} \operatorname{Aut}(C) \longrightarrow \mathcal{P}_{k-1}(C) \\
\downarrow \qquad \qquad \downarrow \\
\coprod_{\substack{[T] \in \Omega_G^a[k](C)/\simeq}} \left( \bigotimes_{v \in V_{G,\mathcal{P}}(T)} \mathcal{P}(T_v) \otimes \bigotimes_{v \in V_{G,in}(T)} Y_1(T_v) \right) \otimes_{\operatorname{Aut}(T)} \operatorname{Aut}(C) \longrightarrow \mathcal{P}_k(C)$$

where the left vertical map is the iterated box product

$$\coprod_{V_{G,\mathcal{P}}(T)} \iota_{\mathcal{P}(T_v)} \, \Box [u]^{\Box \mathbb{V}_{G,in}(T)},$$

 $\iota_{\mathcal{P}(T_v)}$  is the rearranged map recession of the initial object, and  $\Omega_G^a[k](C)$  is as in Definition 7.3.

Proof. The first result is just a restatement of Diagram (7.14), 200

$$l[k]: L_1[k] \to L_2[k].$$

By stacking left Kan extensions, we have

$$L_i[k] \simeq \operatorname{Lan}_{\Omega_G^a[k]}(\operatorname{Lan}_{\bar{\Omega}_G^e[\epsilon_i] \to \Omega_G^a[k]} N^e),$$

where  $[\epsilon_1] = [k, -]$  and  $[\epsilon_2] = [k]$ . Evaluating at  $C \in \Sigma_G$  using pointwise left Kan extensions yields

$$L_{i}[k](C) \simeq \underset{\substack{\Omega_{G}^{a,op}_{r}, C \\ T \stackrel{\rightleftharpoons}{\leftarrow} C}}{\operatorname{colim}} \left( \underset{\substack{\bar{\Omega}_{G}^{e}[\epsilon_{i}]^{op}_{r}, T \\ T^{e} \leftarrow T}}{\operatorname{colim}} N^{e}(T^{e}) \right).$$

Further, we observe that  $\bar{\Omega}_G^e[k]^{op}\downarrow_r T$  is equivalent to the cube category  $(0\to 1)^{V_{G,in}(T)}$  (with only  $\operatorname{Aut}(T)$ -induced isomorphisms), while  $\bar{\Omega}_G^e[k,-]^{op}\downarrow_r T$  is similarly equivalent to the punctured cube category  $\operatorname{pc}(V_{G,in}(T))$ . Unpared Lathers statements in particular that cube categories have terminal objects, Lemma 7.10, and the fact that  $\Omega_G^{a,op}\downarrow_r C$  is a groupoid, combine to imply that l[k] is equivalent to the map

$$\coprod_{[T]\in\Omega_G^{a,op}[k](C)/\simeq} \left( \bigotimes_{V_{G,\mathcal{P}}(T)} \mathcal{P}(T_{Gv}) \otimes Q[u]^{V_{G,in}(T)} \to N^e(T_Y) \right) \otimes_{\operatorname{Aut}(T)} \operatorname{Aut}(C),$$

as desired.

# 8 Model Structures on Genuine Operads

come back: this is all disorganized, internally, externally, everything

replaced \ F with  $\mathcal{F}$  (as opposed to  $\mathbb{F}$ )

In order to encode the homotopical information inspired by  $N_{\infty}$ -operads dicussed in the introduction, we will introduce (semi) model structures on the categories  $\mathcal{V}\mathsf{Op}_G$  and  $\mathcal{V}\mathsf{Op}_G^G$  for a wide range of  $\mathcal{V}$ , which are true model categories in the cases we are most interested in (i.e. for  $\mathcal{V} = \mathsf{sSet}$ ). These model structures will be determined by a choice of weak indexing system, a generalization of the notion defined in  $\boxed{2}$ .

# 8.1 Weak Indexing Systems

We recall certain constructions found in [6] relating graph subgroups, finite H-sets, and systems of categories.

**Definition 8.1.** A G-graph subgroup of  $G \times \Sigma_n$  is a subgroup  $\Lambda \leq G \times \Sigma_n$  such that  $\Lambda \cap \Sigma_n = e$ . Equivalently,  $\Gamma = \Gamma(\phi)$  is the graph of some homomorphism  $G \geq H \xrightarrow{\phi} \Sigma_n$ .

**Definition 8.2.** A *G-vertex family* is a collection

$$\mathcal{F} = \coprod_{n>0} \mathcal{F}_n$$

where each  $\mathcal{F}_n$  is a family of G-graph subgroups of  $G \times \Sigma_n$ , closed under subgroups and conjugation. For a fixed  $\mathcal{F}$ , we call an H-set  $A \in \mathsf{F}^H$   $\mathcal{F}$ -admissible if for some (equivalently, any) choice of bijection  $A \leftrightarrow \{1,\ldots,|A|\}$ , the graph subgroup of  $G \times \Sigma_n$  encoding the induced H-action on  $\{1,\ldots,|A|\}$  is in  $\mathcal{F}_n$ .

**Definition 8.3.** For any G-vertex family  $\mathcal{F}$ , a G-tree  $T \in \Omega_G$  is called  $\mathcal{F}$ -admissible if, for each vertex  $e_1 \dots e_n \leq e$  in V(T), the set  $\{e_1, \dots, e_n\}$  is an  $\mathcal{F}$ -admissible  $\operatorname{Stab}_G(e)$ -set. We let  $\Omega_{\mathcal{F}} \subseteq \Omega_G$  and  $\Sigma_{\mathcal{F}} \subseteq \Sigma_G$  denote the full subcategories spanned by the  $\mathcal{F}$ -admissible trees.

**Definition 8.4.** A G-vertex family is called a weak indexing system if  $\Omega_{\mathcal{F}}$  is a sieve of  $\Omega_{G}$ ; that is, for any map  $f: S \to T$  with  $T \in \Omega_{\mathcal{F}}$ , we have that S (and the map f) are in  $\Omega_{\mathcal{F}}$ .

We note that this always holds for any  $\mathcal F$  is f is an outer face or a quotient (for the latter, this follows from each  $\mathcal F_n$  being closed under subgroups). However, closure under degeneracies implies that all trivial orbits H/H are  $\mathcal F$ -admissible, while closure under inner faces implies that the  $\mathcal F$ -admissible sets are closed under "broad self-induction": if  $A \amalg H/K$  and B are  $\mathcal F$ -admissible H- and K-sets, respectively, then  $A \amalg H \times_K B$  is also an  $\mathcal F$ -admissible H-set. It also implies, in particular, the following:

**Lemma 8.5.** The valence map  $\operatorname{Ir}$  restricts to a map  $\operatorname{Ir}: \Omega_{\mathcal{F}} \to \Sigma_{\mathcal{F}}$ ; that is, if T is  $\mathcal{F}$ -admissible, so is  $\operatorname{Ir}(T)$ .

**Lemma 8.6.** A weak indexing system  $\mathcal{F}$  is an indexing system (in the sense of [2]) if and only if all trivial H-sets are  $\mathcal{F}$ -admissible, for all  $H \leq G$ .

*Proof.* come back

We will define  $\mathcal{F}$ -model structures on  $\mathcal{V}\mathsf{Op}_G$  for any weak indexing system  $\mathcal{F}$ ; the  $\mathcal{F}$ -model structure on  $\mathcal{V}\mathsf{Sym}_G$  will exist for any G-vertex family  $\mathcal{F}$ , but transfering requires the additional closure properties.

NG\_VALENCE\_LEMMA

### 8.2 Semi Model Structures on Genuine Operads

Fix a G-vertex family  $\mathcal{F} = \{\mathcal{F}_n\}$ . For the following section, we fix the following about our base category  $\mathcal{V}$ :

**Definition 8.7.** We say  $\mathcal{V}$  satisfies Assumption 1 (for  $\mathcal{F}$ ) if the following hold:

- 1.  $\mathcal V$  is a cofibrantly-generated Cartesian symmetric monoidal model category, and
- 2. V is has cellular fixed point functors for all finite groups (c.f. [7]).

**Definition 8.8.** The  $\mathcal{F}$ -projective model structure on  $\mathcal{V}\mathsf{Sym} = \mathcal{V}^{\Sigma_G^{op}}$  is the unique model structure induced by the adjunction

$$\mathcal{V}^{\Sigma_G^{op}} \subseteq \mathcal{V}^{\Sigma_F^{op}}$$

Explicitly, a map f is a fibration (resp. weak equivalence) if f(C) is one in  $\mathcal{V}$  for all  $C \in \Sigma_{\mathcal{F}}$ . **Definition 8.9.** A map  $f : \mathcal{O} \to \mathcal{P}$  in  $\mathcal{V}\mathsf{Op}_G$  is called a

- 1.  $\mathcal{F}$ -fibration (resp.  $\mathcal{F}$ -weak equivalence) if  $f(C) : \mathcal{O}(C) \to \mathcal{P}(C)$  is one in  $\mathcal{V}$  for all  $\mathcal{F}$ -admissible G-corollas  $C \in \Sigma_{\mathcal{F}}$ .
- 2.  $\mathcal{F}$ -cofibration if it has the left lifting property against all maps which are both  $\mathcal{F}$ fibrations and  $\mathcal{F}$ -weak equivalences.
- 3.  $\Sigma_{\mathcal{F}}$ -cofibration if f is an  $\mathcal{F}$ -cofibration in  $\mathcal{V}\mathsf{Sym}_G$ .

In particular,  $\mathcal{P} \in \mathcal{V}\mathsf{Op}_G$  will be called  $\mathcal{F}\text{-}cofibrant$  if  $\emptyset \to \mathcal{P}$  is an  $\mathcal{F}$ -cofibration.

**Definition 8.10.** The  $\mathcal{F}$ -model structure on  $\mathcal{V}\mathsf{Op}_G$ , if it exists, is the unique model structure with the above specified weak equivalences and fibrations. Equivalently, it is the transfered model structure along the adjoints

$$\mathcal{V}\mathsf{Op}_G \xrightarrow[\mathsf{fgt}]{\mathbb{F}_G} \mathcal{V}^{\Sigma_G^{op}} \longleftrightarrow \prod_{\mathsf{Ob}(\Sigma_G)} \mathcal{V} \longleftrightarrow \prod_{\mathsf{Ob}(\Sigma_{\mathcal{F}})} \mathcal{V}$$

Using general arguments of [3], this structure would be cofibrantly generated, with generating arrows

$$I_{\mathcal{F}} = \{ \mathbb{F}_{G}(\Sigma_{G}(-,C) \cdot i) \mid C \in \Sigma_{\mathcal{F}}, \ i \in I \}$$

$$J_{\mathcal{F}} = \{ \mathbb{F}_{G}(\Sigma_{G}(-,C) \cdot j) \mid C \in \Sigma_{\mathcal{F}}, \ j \in J \},$$

for I (resp. J) the generating (trivial) cofibrations of  $\mathcal{V}$ .

**Definition 8.11.** If  $\mathcal{F}$  is the complete G-vertex family - so  $\mathcal{F}_n$  is the family of all graph subgroups of  $G \times \Sigma_n$  - we refer to the  $\mathcal{F}$ -model structure as the *genuine* model structure.

**Theorem 8.12.** The genuine semi-model structure on  $\mathcal{V}\mathsf{Op}_G$  exists for all  $\mathcal{V}$  satisfying ASSUMPTION 1.

*Proof.* By general arguments of Kan [3, 11.6.1], it suffices to show that any transfinite composite of cellular extensions  $\mathcal{P} \to \mathcal{P}[u]$ , each built by pushouts

$$\mathbb{F}_G X \longrightarrow \mathcal{P} \\
\downarrow \\
\downarrow \\
\mathbb{F}_G Y \longrightarrow \mathcal{P}[u]$$

where  $u: X \to Y$  is a generating trivial cofibration and  $\mathcal{P}$  is cofibrant in  $\mathcal{V}_{gen}^{\Sigma_{G}^{O}}$ , is itself a weak equivalence in  $V_{gen}^{\Sigma_{gen}^{O}}$ .

In particular, it suffices to show that, for u a genuine level trivial cofibration and  $\mathcal{P}$  genuine level cofibrant, any map  $\mathcal{P} \to \mathcal{P}[u]$  is itself a genuine level trivial cofibration; since these are weaker conditions on u and  $\mathcal{P}$ , and transfinite composition is levelwise, this implies the above result. This case follows from the levelwise filtration from Theorem 7.17, by applying Propositions 8.14 and 7??.

E\_SEMI\_VOP\_G\_THM

**Remark 8.13.** In proving the above, we will often use the fact that, given families  $\mathcal{F} \subseteq \overline{\mathcal{F}}$  of subgroups of a group  $\Pi$ , the identity map

$$\mathcal{V}_{\mathcal{F}}^G \to \mathcal{V}_{\bar{\mathcal{F}}}^G$$

is left Quillen (and in particular preserves cofibrations). In particular, any  $\mathcal{F}$ -cofibration is a genuine cofibration.

**Proposition 8.14.** Fix  $T \in \Omega_G$ , and suppose we are given cofibrations  $f(Gv) \in \mathcal{V}_{gen}^{\operatorname{Aut}(T_{Gv})}$  for all  $Gv \in V_G(T)$ , such that  $f(Gv) = f(\alpha(Gv))$  for all  $\alpha \in \operatorname{Aut}(T)$ . Then the iterated box product

 $f^{\square V_G(T)} = \underset{Gv \in V_G(T)}{\square} f(Gv)$ 

is a cofibration in  $\mathcal{V}_{gen}^{\operatorname{Aut}(T)}$ .

X\_COFIBRANT\_PROP

*Proof.* Using grafting composition  $T \simeq C \circ (T_1^1, \dots, T_1^{k_1}, T_2^1, \dots, T_r^{k_r})$ , we go by induction on  $|V_G(T)|$ . The base cases of  $|V_G(T)| = 0$  or 1 are trivial. Now, we note that we have a decomposition

 $f^{\square V_G(T)} = f(Gv_r) \underset{i=1}{\square} \left( \left( f^{\square V_G(T_i)} \right)^{\square k_i} \right)$ 

where  $Gv_r$  is the root orbit. We will build this map in stages, preserving cofibrancy in each step.

- By induction, for each i,  $f^{\square V_G(T_i)}$  is a cofibration in  $\mathcal{V}_{qen}^{\mathrm{Aut}(T_i)}$ .
- By Proposition ???, each  $(f^{\square V_G(T_i)})^{\square k_i}$  is a cofibration in  $\mathcal{V}_{gen}^{\Sigma_{k_i} \wr \operatorname{Aut}(T_i)}$ .
- $\bullet \ \, \text{By Remark} \overset{\texttt{EXTERINTADJ}, \texttt{EQ}}{???, \ \underset{i=1}{\square}} \left( \left( f^{\square VG}(T_i) \right)^{\square k_i} \right) \text{ is a cofibration in } \mathcal{V}_{gen}^{\Pi\Sigma_{k_i} \wr \texttt{Aut}(T_i)}.$
- By Proposition  $??, f(Gv_r)$  is a cofibration in  $V_{gen}^{\text{Aut}(T)}$ .
- Finally, by the decomposition on Aut(T) and Proposition ??, we have the full  $f^{\Box V_G(T)}$  is a cofibration in  $\mathcal{V}_{qen}^{\mathrm{Aut}(T)}$ , as desired.

If any f(Gv) were trivial, then, during the appropriate step above, the resulting box product would again be trivial, by the same referenced results.

Corollary 8.15. For any G-vertex system  $\mathcal{F}$ , the  $\mathcal{F}$ -semi-model structure on  $\mathcal{VOp}_G$  exists.

Proof. This follows from the proof of Theorem 8.12. In particular, it suffices to show that if  $u:X\to Y$  is a level trivial  $\mathcal F$ -cofibration and  $\mathcal P$  is level  $\mathcal F$ -cofibrant in  $\mathcal V^{\Sigma_G^{op}}$ , then the cellular extension  $\mathcal P\to\mathcal P[u]$  is a level trivial  $\mathcal F$ -cofibration. However, such u are level genuine cofibrations, and such  $\mathcal P$  are level genuine cofibration, and hence  $\mathcal P\to\mathcal P[u]$  is a level trivial genuine cofibration, and hence a level trivial  $\mathcal F$ -cofibration.

#### 8.3 True Model Structures

come back

# 8.4 Preservation of Cofibrant Objects

come back

#### Cofibrant Symmetric Collections 8.5

come back: can do all of this for  $\mathcal{V}^{\Omega_G^{op}}$  and  $\mathcal{V}^{\Sigma_G^{op}}$ , don't need these crazy categories.

We first classify cofibrant objects in  $\mathcal{V}^{\Omega_G^e[k]^{op}}$ , by breaking apart our category  $\Omega_G^e[k]$ .

**Definition 8.16.** Given a tree  $T_0 \in \Omega$ , let  $\Omega_G[T_0]$  denote the full subcategory of  $\Omega_G^q$  spanned by those trees which recieve a map from  $G \cdot T_0$ ; that is, all trees of the form  $G \cdot T_0/N$ .

**Lemma 8.17.**  $\Omega_G[T_0] \simeq O_{\Gamma_{T_0}}$  as categories, where  $\Gamma_{T_0}$  is the family of graph subgroups of  $G \times \Sigma_{T_0}$ .

**Definition 8.18.** Refining the above, given an odd tree  $T_0 \in \Omega^{\text{odd}}$ , let  $T_0(Y) \in \Omega^e$  denote the labeled tree with underlying tree  $T_0$  and all even nodes labeled with Y. Further, let  $\Omega^e[T_0]$  denote the full subcategory of  $\Omega^e$  spanned by those trees  $(S_0, \lambda_{S_0})$  which receive a map from  $T_0(Y)$  in  $\Omega^e$ . Similarly, let  $\Omega_G^e[T_0]$  denote the full subcategory of  $\Omega_G^e$  spanned by those trees  $(S, \lambda_S)$  which recieve a map from  $G \cdot T_0(Y)$ .

We have an inclusion  $\tau: G \times \Omega^e[T_0] \to \Omega^e_G[T_0]$ , refining the inclusion  $G \times \Omega \to \Omega_G$ . This induces an adjunction

$$\mathcal{V}^{G \times \Omega^e[T_0]^{op}} \xrightarrow[\tau_*]{\tau^*} \mathcal{V}^{\Omega^e_G[T_0]^{op}}$$

and we observe that

$$\tau_* X(G \cdot T_0/N, \lambda) \simeq X(G \cdot T_0, q^* \lambda)^N,$$

where we are using the fact that even element in  $\Omega_G^e[T_0]$  has underlying tree of the form  $G \cdot T_0/N$ , and  $q^*\lambda$  is the vertex labeling

$$q^*\lambda: V_G(G \cdot T_0) \to V_G(G \cdot T_0/N) \xrightarrow{\lambda} \{Y, X\}.$$

We further repackage  $\Omega_G^e[T_0]$ . In particular, we observe that if there is a map  $(T', \lambda') \to$  $(T'',\lambda'')$  in  $\Omega_G^e[T_0]$ , it is uniquely determined by the underlying quotient map  $T'\to T''$ . Indeed,  $\Omega_G^e[T_0](T',T'')$  is the subset of  $\Omega_{\Gamma_{T_0}}(G \times \Sigma_{T_0}/N', G \times \Sigma_{T_0}/N'')$ , where  $T' \simeq G \cdot T_0/N'$ and  $T'' \simeq G \cdot T_0/N''$ , of those maps q such that  $v'' \in V_Y(T'')$  implies  $q^{-1}(v'') \subseteq V_Y(T')$ . We can now prove a proposition, modeled on the proof of [7, Theorem 2.10].

**Proposition 8.19.** If  $X \in \mathcal{V}^{\Omega_G^e[T_0]^{op}}$  is cofibrant, then  $\eta_X : X \to \tau_* \tau^* X$  is an isomorphism.

come back, type up these sheets Proof.

First:  $\eta$  is an iso on representables.

Second: carries through across the necessary cellular extensions.

#### G-Operads 8.5.1

**Definition 8.20.** A map  $f: \mathcal{O} \to \mathcal{P}$  in  $\mathcal{V}\mathsf{Op}^G$  is called a

1.  $\mathcal{F}$ -fibration (resp.  $\mathcal{F}$ -weak equivalence) if ...

come back

**Definition 8.21.** The  $\mathcal{F}$ -model structure on  $\mathcal{V}\mathsf{Op}^G$ , ...

**Lemma 8.22.** The F-semi-model structure, if it exists, is the transfered model structure along the adjunction

$$VOp_{\mathcal{T}}^{G} \leftrightharpoons VOp_{G}^{\mathcal{F}}$$

Corollary 8.23. The  $\mathcal{F}(-semi)$ -model structure exists on  $\mathcal{V}\mathsf{Op}^G$  whenever it exists on  $\mathcal{V}\mathsf{Op}_G$ .

**Theorem 8.24.**  $VOp_{\mathcal{F}}^G$  and  $VOp_{\mathcal{G}}^{\mathcal{F}}$  are Quillen equivalent.

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