Exercise Sheet 12

VU Performance-oriented Computing, Summer Semester 2024

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A) Setup and Basic Execution

Benchmark results with default Lua interpreter compilation (-O2):

test	wall (mean)	wall (stddev)
iter	10.8730	0.036
naive	12.5483	0.322
tail	12.5980	0.046

B) Profiling

Profiling with Gprof

Setup

I first had to recompile the interpreter with the added -pg flag, i had some issues so i had to add -g and -no-pie aswell to get a good output.

Those flags were very easy to set in the Makefile in the /src/ directory as there was a place for custom flags.

Results

gprof output

Figure 1: gprof output

From the flat profile we get by gprof we can see that 77.25% of the time is spent in the function luaV_execute that is therefore a function we should have a look at, here it's also interesting to note that we only have one call to this function, so it might also be something like a main.

Next in the ranking is the function luaD_pretailcall with 10.88% of the execution time, this function has a lot of calls on it (306,946,270), which means that even a small percentage improvement could have a big impact.

The same goes for luaD_precall.

The full profile can be found in gprof_results.txt.

Profiling with Tracy

The modified source code including Makefile can be found in the lua_tracy/ directory. To enable Tracy, compile with make EXTRACFLAGS=-DTRACY_ENABLE. To additionally enable per-opcode zones (refer to the section "Opcode Statistics"), compile with make EXTRACFLAGS=-DTRACY_ENABLE.

The -g flag was added to the default compilation flags in order to make Tracy's source code view would work properly. This did not appear to significantly impact runtime, at least on my machine.

I added the ZoneScoped macro to the following functions:

- luaV_execute
- luaF_findupval
- luaF_closeupval
- luaH_getstr
- luaH_getshortstr
- luaC_step

Additionally, we added a FrameMark to the end of the interpreter loop in an attempt to measure the number of opcodes executed per second, although this is seemingly never hit.

Most of the runtime is spent, unsurprisingly, directly inside the main interpreter loop.

Which Opcodes?

Using the Source view in the Tracy Profiler application, we can narrow down the Lua Virtual Machine opcodes that are actually used by the fib.lua benchmark:

(This is also possible using e.g. perf report, although Tracy's UI is much more navigable.)

- OP MOVE
- OP LOADI
- OP_LOADK
- OP_GETUPVAL
- OP GETTABUP
- OP GETFIELD
- OP_SELF
- OP_ADDI
- OP_ADD
- OP_EQI
- OP_LTI
- OP_GEI
- OP_TESTSET
- OP_CALL
- OP_TAILCALL
- OP_RETURN1
- OP_FORLOOP
- OP_FORPREP
- OP_TFORLOOP
- OP_CLOSURE
- OP VARARG
- OP VARARGPREP
- OP_EXTRAARG

This represents a small subset of the opcodes available in the interpreter, giving a better idea of where optimizations should be made.

Opcode Statistics

To get an idea of which opcodes are executed how often, I added named zones to each opcode in the luaV execute function as follows:

```
vmdispatch (GET_OPCODE(i)) {
    // [...]
    vmcase(OP_LTI) {
        TracyCZoneN(ctx, "OP_LTI", true);
        op_orderI(L, l_lti, Luai_numlt, 0, TM_LT);
        TracyCZoneEnd(ctx);
        vmbreak;
```

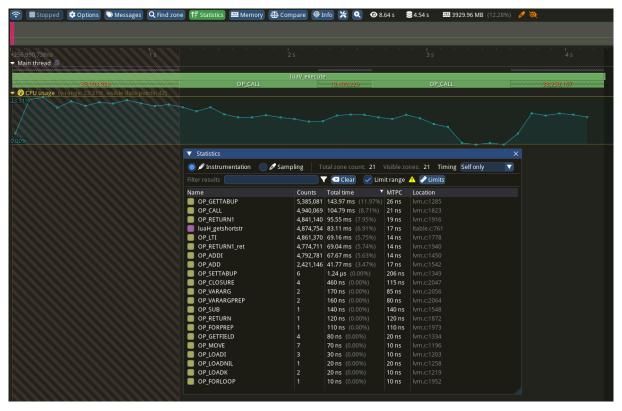
```
}
// [...]
}
```

Special care had to be taken around goto statements in the interpreter, particularly around OP_RETURN1 and the ret label.

As the large number of zones created this way quickly overwhelms the Tracy Profiler host application, I created a modified version of the fib.lua script that calls each fibonacci_ variant only 1/100x as often, and sleeps for one second in between testing the three variants:

Note that this did have a substantial performance impact, increasing the runtime by a factor of 2-4.

- The following screenshot shows
 - 1. the overall execution profile (note the 1-second long OP_CALL zones in between the orange-numbered ones representing the sleep durations).
 - 2. statistics for 100x fibonacci_naive (obtained by right-clicking on the first orange-numbered section, choosing "Limit statistics time range", then opening the Statistics window using the button at the top).



All opcode statistics are shown in the following screenshots.

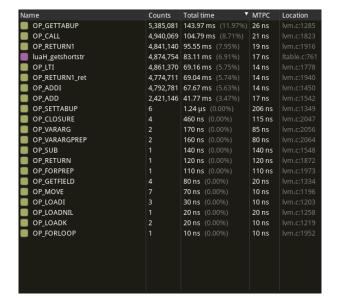


Figure 2: fibonacci_naive

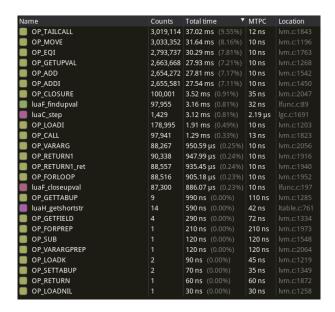


Figure 3: fibonacci_tail



Figure 4: fibonacci_iter

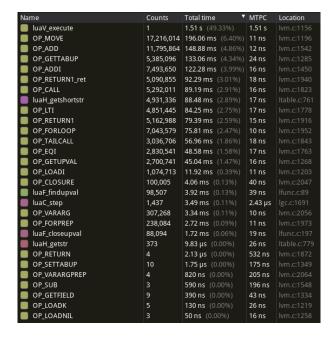


Figure 5: all three

Main Takeaways

- fibonacci_naive makes many, many calls to getshortstr via OP_GETTABUP I assume this has to do with accessing the variable n from within recursive calls.
- OP_GETTABUP is used much more often than OP_SETTABUP and its mean-time-per-call is roughly twice that of other opcodes. Optimising OP_GETTABUP at the expense of OP_SETTABUP is probably a good call for this benchmark.
- fibonacci_tail calls findupval and closeupval roughly once per iteration. This is related to creation of the "inner" closure via OP_CLOSURE.
- The garbage collector (luaC_step) runs most often during fibonacci_tail taking up about as much runtime as findupval.
- OP_MOVE, OP_FORLOOP and OP_ADD represent the bulk of the opcodes used by fibonacci_iter with several million calls each.

Conclusion

We decided to take a look at the functions using the most time and tried to optimize them in exercise D.

C) Code Understanding

Overall process of Lua execution in the interpreter

- Loading: The Lua interpreter reads the script from a file or string.
- Lexical Analysis (Lexing): The script is broken down into tokens.
- Syntax Analysis (Parsing): The tokens are parsed to create an Abstract Syntax Tree (AST).
- Code Generation: The AST is converted into bytecode, which is the intermediate representation of the script.
- Execution: The bytecode is executed by the Lua Virtual Machine (VM).

LUA_USE_JUMPTABLE

The bulk of the main loop of the interpreter is taken up by a switch/case construct, with one case label per opcode supported by the Lua Virtual Machine (e.g. case OP_MOVE:).

When the macro LUA_USE_JUMPTABLE is defined, the switch/case construct is replaced with a manually-defined jump table:

- 1. Instead of a case label, each Virtual Machine opcode is assigned a "regular" label (e.g. L_OP_MOVE:), and the addresses for each of these are stored in the table.
- 2. Immediately after an opcode has been processed, the next opcode is fetched and execution is transferred to its corresponding label using goto as opposed to breaking out of the switch block, then fetching and jumping to the next opcode on the next loop iteration.

While the switch/case construct is also a jump table, the idea of constructing one manually is that, by avoiding a jump back to the top of the interpreter loop before processing the next opcode, performance could be improved.

However as Roberto (creator of Lua) noted in 2018, it may actually cause performance regressions on some systems.

My performance testing seems to reflect this – across ten runs of fib.lua on LCC3, the jump table seems to make performance ever so slightly worse. The difference in runtime is roughly one standard deviation for the iter and tail variants, and one fifth of a standard deviation for the naive variant.

test	wall (mean)	wall (stddev)
iter	10.8730	0.036
iter (jumptable)	10.9105	0.036
naive	12.5483	0.322
naive (jumptable)	12.6001	0.316
tail	12.5980	0.046
tail (jumptable)	12.6484	0.027

(D) Optimization

Varying the compilation flags

We tried different compiler flags – this did not produce any meaningful improvements. With the exception of fibonacci_tail, compiling with -03 or -0fast made performance worse:

test	wall (mean)	wall (stddev)	% vs -02
iter	10.8730	0.036	
iter (-03)	12.2039	0.042	112.24

test	wall (mean)	wall (stddev)	% vs -02
iter (-Ofast)	11.6290	0.031	106.95
naive	12.5483	0.322	
naive (-03)	13.4102	0.296	106.87
naive (-Ofast)	12.7103	0.306	101.29
tail	12.5980	0.046	
tail (-03)	12.5004	0.020	99.23
tail (-Ofast)	12.0485	0.018	95.64

Applying individual flags from the -03 optimization level in addition to -02 also did not have any meaningful impact. The greatest improvement achieved by a single flag is -fsplit-paths for fibonacci_tail, being 0.15 seconds faster than -02 (\sim 1.3%).

All mean results are contained in the file named flag_results.csv.

Improving Lua source code

We tried to improve the functions mentioned in *B* - *Profiling with Gprof*. We tried to inline external function calls. However the only function that was no define which is inlined by the compiler anyway was prepCallInfo in luaD_precall in ldo.c. (Note: Only adapted the last case, as it is use for running Lua files and our benchmark is in Lua)

The results were slightly faster (no jump table was also faster than jump table):

```
• naive: 12.3457 s
  • tail: 12.5954 s
  • iter: 10.8881 s
The adapted Code looks like this:
CallInfo *luaD_precall (lua_State *L, StkId func, int nresults) {
 retry:
  switch (ttypetag(s2v(func))) {
   case LUA_VCCL: /* C closure */
     precallC(L, func, nresults, clCvalue(s2v(func))->f);
     return NULL;
   case LUA_VLCF: /* light C function */
     precallC(L, func, nresults, fvalue(s2v(func)));
     return NULL;
   case LUA_VLCL: { /* Lua function */
     CallInfo *ci;
     Proto *p = clLvalue(s2v(func))->p;
     int narg = cast_int(L->top.p - func) - 1; /* number of real arguments */
     int nfixparams = p->numparams;
     int fsize = p->maxstacksize; /* frame size */
     checkstackGCp(L, fsize, func);
     // original code
     //L->ci = ci = prepCallInfo(L, func, nresults, 0, func + 1 + fsize);
     ci = L->ci = next_ci(L); /* new frame */
     ci->func.p = func;
     ci->nresults = nresults;
     ci->callstatus = 0;
     ci->top.p = func + 1 + fsize;
     L->ci = ci;
```

```
// original function
   //l_sinline CallInfo *prepCallInfo (lua_State *L, StkId func, int nret,
                                           int mask, StkId top) {
   //CallInfo *ci = L->ci = next_ci(L); /* new frame */
   //ci->func.p = func;
   //ci->nresults = nret;
   //ci->callstatus = mask;
   //ci->top.p = top;
   //return ci;
   //}
   ci->u.l.savedpc = p->code; /* starting point */
   for (; narg < nfixparams; narg++)</pre>
     setnilvalue(s2v(L->top.p++)); /* complete missing arguments */
   Lua_assert(ci->top.p <= L->stack_last.p);
   return ci;
  default: { /* not a function */
   func = LuaD_tryfuncTM(L, func); /* try to get '__call' metamethod */
   /* return LuaD_precall(L, func, nresults); */
   goto retry; /* try again with metamethod */
}
```

Improving the op calls analyzed with tracy seemed very hard. Probably implementing caching / memoization would rapidly improve the benchmark for a problem like fibonacci calculation. However implementing this in the Lua interpreters source code and not simply adding memoization to the .lua file would exceed the scope of this exercise sheet.

Using a custom memory allocator for the Lua interpreter

As we had good experience with custom memory allocators in terms of performance improvements in previous exercises, we decided to try mimalloc and see if it can improve the performance.

Mimalloc Setup

First we build mimalloc from source, as described in the repository. After that we added the Flag-lmimalloc during the build of the interpreter to use it for all the allocations.

Mimalloc Results

Unfortunatly mimalloc couldn't really speed up the execution times for this scenario.

The baseline results were (avg. from 10 runs):

```
func,wall (mean)
fibonacci_iter(30),10.8779
fibonacci_naive(30),12.7706
fibonacci_tail(30),12.6074
And with mimalloc:
func,wall (mean)
fibonacci_iter(30),10.9223
fibonacci_naive(30),12.5884
fibonacci_tail(30),12.6433
```

We can see that the "improvements" fall within a margin of error. Our possible explanations for that were:

- 1. The program probably is rather compute bound and not so heavy on the memory.
- 2. Maybe Lua already uses a non-default memory allocator / they have their own implementation. A look in the source code, we can see the file lmem.c which contains the implementation of Lua's memory handling. Of course this is completed by the garbage collector in seperate files.