Algorithms and Data Structure 7

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Problem 1

a)

The following is the concrete implementation in Python. Python is by-default generic, since it is not statically typed language.

Note: The python code includes a small unit test. You can simply run

\$ python -m doctest -v stack.py

```
class StackNode:
    A basic linked list implementation
    def __init__(self, data = None, nxt = None):
        self.data = data
        self.next = nxt
class Stack:
    Implementation of Stack According to the code in the assignment
    >>> s = Stack(2)
    >>> s.isEmpty()
    True
    >>> s.push(5)
    >>> s.isEmpty()
    >>> s.current_size
    >>> s.push(4)
    >>> s.current_size
    >>> s.pop()
    >>> s.pop()
    >>> s.isEmpty()
    True
    def __init__(self, size = None):
       if size is None:
            self.size = -1
        elif size <= 0:</pre>
            raise Exception('Invalid size')
        else:
            self.size = size
```

```
self.current_size = -1
    self.top = None
def push(self, data):
    # check the size
    if self.size != -1 and self.current_size == self.size:
        raise Exception('Stack Overflow Exception')
    if self.top is None:
        self.top = StackNode(data) #create a node
        self.current_size = 1 #set the proper size
    else:
        self.top = StackNode(data, self.top) #simply append the linked list
        self.current_size = self.current_size + 1 # increment
def pop(self):
    if self.top is None:
        raise Exception('Stack Underflow Exception')
    val = self.top.data
    self.top = self.top.next # move to next one
    self.current_size = self.current_size - 1
    return val
def isEmpty(self):
    return self.top is None
```

b)

We can implement a queue via two stacks. The first stack store the input, or freshly enqueued elements. The second stack store the output, or the elements that is ready is dequeue. When we try to dequeue, and if the second stack is empty, dump the first stack into the second one(poping from stack 1 and push into stack 2). The implementation in Python is also included here.

Note: The python code includes a small unit test. You can simply run

\$ python -m doctest -v queue.py

```
from stack import Stack

class Queue:
    """

    A simple Queue using two stack

>>> q = Queue()

>>> q.enqueue(2)

>>> q.enqueue(3)

>>> q.enqueue(4)

>>> q.dequeue()

2

>>> q.dequeue()

3

>>> q.dequeue()

4

>>> q.dequeue()
```

```
5
11 11 11
def __init__(self):
    # initalize two Stack
    self.s1 = Stack()
    self.s2 = Stack()
def enqueue(self, data):
    # just put it into stack 2
    self.s1.push(data)
def prepare_dequeue(self):
    # dump stack 1 into stack 2
    if self.s2.isEmpty():
        while not self.s1.isEmpty():
            self.s2.push(self.s1.pop())
def dequeue(self):
    self.prepare_dequeue() #prepare
    if self.s2.isEmpty():
        raise Exception('Queue Underflow Exception')
    return self.s2.pop()
```

Problem 2

a)

To reverse a linked list in-place (or in-situ), we have to abandon our regular functional recursive approach. Instead, we just have to reverse the direction of the pointer. The following Python code is a example of how to reverse a linked list in place. Since we only uses 3 reference, "next", "prev", and "current". The space complexity is obviously constant. Also, since there is only one "while" loop iterate through the linked list, and every operation within the "while" takes constant time, so the time complexity is linear(proportional to the length of the linked list). Thus, the time complexity is O(n).

Note: The python code includes a small unit test. You can simply run

```
$ python -m doctest -v linkedlist.py
```

```
class Node:
    # a basic implementation of linked list node in python
    def __init__(self, data = None, nxt = None):
        self.data = data
        self.next = nxt

def reverse(head):
    """
    test the reverse function
    >>> head = Node(1, Node(2, Node(3)))
    >>> rv = reverse(head)
    >>> rv.data
    3
    >>> rv = rv.next
    >>> rv.data
    2
    >>> rv = rv.next
```

```
>>> rv.data
1
>>> rv.next is None
True
"""
# just call the helper
prev = None
current = head
nxt = None
while current is not None:
    nxt = current.next
    current.next = prev
    prev = current
    current = nxt
return prev
```

b)

Such question is just in-order tranversal of BST. We can write it recursively. Since it will tranverse every element in the tree only once, the total running time is $\Theta(n)$

Note: The python code includes a small unit test. You can simply run

\$ python -m doctest -v binarytree_to_list.py

```
class TreeNode:
    def __init__(self, data, left = None, right = None):
        self.left = left
        self.right = right
        self.data = data
```

```
from linkedlist import Node
from binarytree import TreeNode

def treeToList(root):
    """
    This function will convert a tree into linkedlist
    >>> tree = TreeNode(5, TreeNode(3, TreeNode(2), TreeNode(4)), TreeNode(9))
    >> lst = treeToList(tree)
    >>> lst.data
    2
    >>> lst = lst.next
    >>> lst.data
    3
    >>> lst = lst.next
    >>> lst.data
    4
    >>> lst = lst.next
    >>> lst.data
    5
    >>> lst.data
```

```
# basic reversed in-order traversal of BST
# first right node, then current, then left
rst = None # the result

def tranverse(node):
    nonlocal rst #capture the variable rst
    if node is None:
        return
    tranverse(node.right)
    rst = Node(node.data, rst) #append the result
    tranverse(node.left)

tranverse(root)
return rst
```

c)

We can just build a binary search tree that has no right node. Simply append smaller elements in order on the left, and we will have a binary search tree.

The running time of such operation is $\Theta(n)$ since we will iterate through every element only once.

Note: The python code includes a small unit test. You can simply run

\$ python -m doctest -v list_to_binarytree.py

```
from binarytree import TreeNode
from linkedlist import Node
def listToTree(head):
    # We just simply append everything to the left node
    Convert A sorted linkedlist into binary tree
    \Rightarrow lst = Node(1, Node(2, Node(3, Node(4))))
    >>> tree = listToTree(lst)
    >>> tmp = []
    >>> while tree is not None:
           tmp.append(tree.data)
    . . .
            tree = tree.left
    >>> tmp
    [4, 3, 2, 1]
    rst = None
    while head is not None:
        rst = TreeNode(head.data, rst, None) # append to the left
        head = head.next
    return rst
```