CS202

Title: Sorting and Algorithm Efficiency

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ID: 21802414 Assignment: 1

Description: Comparing Empirical and Theoretical Results of Different Sorting

Algorithms

Question 1

Selection Sort:

Inital Array: 6 1 5 3 7 2 8 4

1th pass:

swap 4,8

6153724 | 8

2th pass:

swap 4, 7

615342 | 78

3th pass:

swap 2,6

21534 | 678

4th pass:

swap 4,5

2143 | 5678

5th pass:

swap 3,4

213 | 45678

6th pass:

swap 3, 3

21 | 345678

7th pass:

swap 1, 2

1 | 2 3 4 5 6 7 8

Key Comparisons: 28

Data Moves: 21

Insertion Sort:

Inital Array: 6 1 5 3 7 2 8 4

1th pass: 1 | 6 5 3 7 2 8 4

2th pass: 15 | 637284

3th pass: 1 3 5 | 6 7 2 8 4

4th pass: 1 3 5 6 | 7 2 8 4

5th pass: 1 2 3 5 6 | 7 8 4

6th pass: 1 2 3 5 6 7 | 8 4

7th pass: 1 2 3 4 5 6 7 | 8

Key Comparisons: 19

Data Moves: 26

Bubble Sort:

Inital Array: 6 1 5 3 7 2 8 4

1th pass:

swap 6, 1

- 16537284
- swap 6,5
- 15637284
- swap 6, 3
- 15367284
- swap 7, 2
- 15362784
- swap 8,4
- 15362748
- 2th pass:
- swap 5, 3
- 13562748
- swap 6, 2
- $1\,3\,5\,2\,6\,7\,4\,8$
- swap 7,4
- 13526478
- 3th pass:
- swap 5, 2
- $1\,3\,2\,5\,6\,4\,7\,8$
- swap 6, 4
- 13254678
- 4th pass:
- swap 3, 2
- 12354678

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swap 5, 4
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 $1\,2\,3\,4\,5\,6\,7\,8$

5th pass:

12345678

Merge Sort:

Inital Array: 6 1 5 3 7 2 8 4

MergeSort called

First portion

6153

Second portion

7284

MergeSort called

First portion

61

Second portion

5 3

MergeSort called

First portion

6

Second portion

1

6

6 1

Merge called

First portion

Second portion

6

16

MergeSort called

First portion

5

Second portion

3

165

1653

Merge called

First portion

3

Second portion

5

1635

Merge called

First portion

13

Second portion

5 6

1356

MergeSort called

First portion

7 2

Second portion

MergeSort called

First portion

7

Second portion

2

13567

135672

Merge called

First portion

2

Second portion

7

135627

MergeSort called

First portion

8

Second portion

4

1356278

13562784

Merge called

First portion

4

Second portion

8

13562748

Merge called

First portion

2 4

Second portion

78

13562478

Merge called

First portion

1234

Second portion

5678

12345678

Key Comparisons: 16

Data Moves: 48

Quick Sort:

Inital Array: 6 1 5 3 7 2 8 4

Partition 1

6 as pivot

Checking if 1 < 6 (pivot)

Checking if 5 < 6 (pivot)

Checking if 3 < 6 (pivot)

Checking if 7 < 6 (pivot)

Checking if 2 < 6 (pivot)

Swap (2,7)

61532784

Checking if 8 < 6 (pivot)

Checking if 4 < 6 (ivot)

Swap (6,4)

41532687

4 as pivot

Checking if 1 < 4

Checking if 5 < 4

Checking if 3 < 4

Swap (3,5)

41352687

Checking if 2 < 4

Swap (2,5)

41325687

Swap (2,4)

21345687

Partition 3

2 as pivot

Checking if 1 < 2

Checking if 3 < 2

Swap (1,2)

12345687

Selecting 8 as pivot

Checking if 7 < 8

Swap (7,8)

12345678

Question 2

Output in Dijkstra

Sorting on Random Arrays

Analysis of Insertion Sort

Array Size	Time Elapsed	countComp	countMove
5000	39.7417	6235702	6240701
10000	160.708	25103015	25113014
15000	357.831	56068725	56083724
20000	638.308	100093616	100113615
25000	991.085	155658723	155683722
30000	1423.08	223610730	223640729
35000	1954.83	307218965	307253964
40000	2543.71	399797892	399837891

Analysis of Merge Sort

Array Size	Time Elapsed	countComp	countMove
5000	1.81195	55259	123616
10000	3.76633	120454	267232
15000	6.01854	189263	417232
20000	8.15281	260876	574464
25000	10.37	334116	734464
30000	12.6657	408650	894464
35000	14.5552	484406	1058928
40000	16.8962	561938	1228928

Analysis of Quick Sort

Array Size Time Elapsed countComp countMove

CS202	2 HW1	SORTING AN	D ALGORITHM EFFICIENCY
5000	1.19592	66414	119349
10000	2.61497	154480	249224
15000	4.00936	236334	390377
20000	5.59431	347663	565425
25000	6.94984	424185	650555
30000	8.88886	536588	944008
35000	10.3557	632630	1027266
40000	11.7163	716116	1157486
Sortin	g on Ascending	Arrays	
Analysis of I	nsertion Sort		
Array Size	Time Elapsed	countComp	countMove
5000	81.0274	12502499	6240701
10000	319.801	50004999	50014998
15000	721.981	112507499	112522498
20000	1275.25	200009999	200029998
25000	1999.04	312512499	312537498
30000	2864.22	450014999	450044998
35000	3903.32	612517499	612552498
40000	5112.17	800019999	800059998
Analysis of N	Merge Sort		
Array Size	Time Elapsed	countComp	countMove

1.1795

2.44087

3.89338

5.30756

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CS202	2 HW1	SORTING AN	ID ALGORITHM EFFICIENCY
25000	6.71511	178756	734464
30000	8.12256	219504	894464
35000	9.3512	260100	1058928
40000	10.8148	298432	1228928
Analysis of 0	Quick Sort		
Array Size	Time Elapsed	countComp	countMove
5000	107.12	12497500	18769996
10000	426.053	49995000	75039996
15000	961.115	112492500	168809996
20000	1703.88	199990000	300079996
25000	2670.63	312487500	468849996
Sortin	g on Descendin	g Arrays	
Analysis of I	nsertion Sort		
Array Size	Time Elapsed	countComp	countMove

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Array Size	Time Elapsed	countComp	countMove
5000	0.049233	4999	6240701
10000	0.096304	9999	19998
15000	0.146081	14999	29998
20000	0.193313	19999	39998
25000	0.239998	24999	49998
30000	0.291311	29999	59998
35000	0.335511	34999	69998
40000	0.388189	39999	79998

Analysis of Merge Sort

Array Size Time Elapsed countComp countMove

CS20	2 HW1	SORTING A	ND ALGORITHM EFFICIENCY
5000	4.40000	22004	122616
5000	1.18923	32004	123616
10000	2.4523	69008	267232
15000	3.90375	106364	417232
20000	5.33883	148016	574464
25000	6.74078	188476	734464
30000	8.16517	227728	894464
35000	9.38636	269364	1058928
40000	10.8665	316032	1228928
Analysis of 0	Quick Sort		
Array Size	Time Elapsed	countComp	countMove
5000	52.9152	12497500	19996
10000	211.594	49995000	39996
15000	475.346	112492500	59996

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ScreenShot from Dijkstra

850.368

1320.85

199990000

312487500

79996

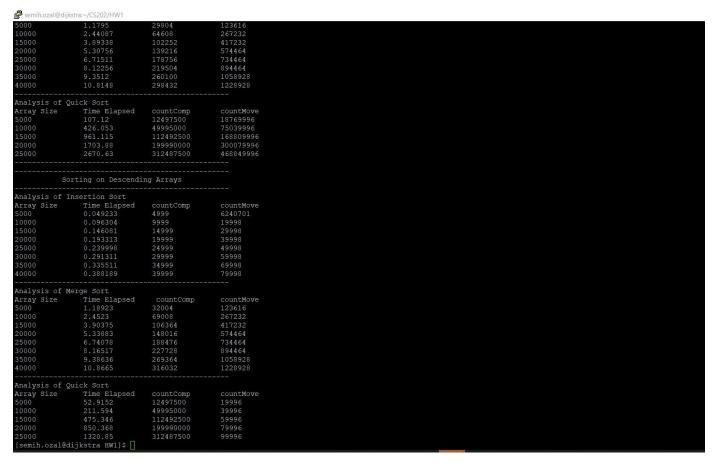
99996

20000

25000

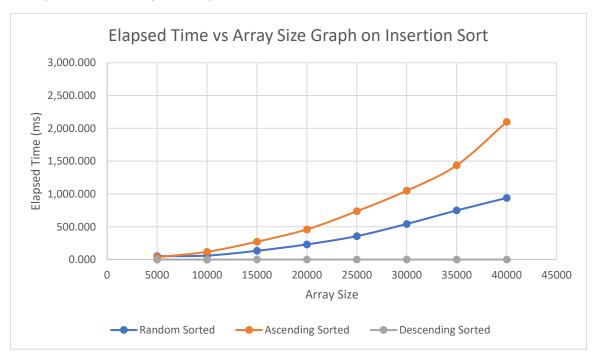
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	ting on Ascendin	g Arrays	
nalysis of in rray Size	sertion Sort Time Elapsed	countComp	countMove
000	81.0274	12502499	6240701
0000	319.801	50004999	50014998
5000	721.981	112507499	112522498
0000	1275.25	200009999	200029998
	1999.04	312512499	312537498
0000	2864.22	450014999	450044998
5000	3903.32	612517499	612552498
0000	5112.17	800019999	800059998
nalysis of Me			
rray Size 000	Time Elapsed 1.1795	countComp 29804	countMove 123616
0000	2.44087	64608	267232
5000	3.89338	102252	417232
0000	5.30756	139216	574464
5000	6.71511	178756	734464
0000	8.12256	219504	894464
5000	9.3512	260100	1058928
0000	10.8148	298432	1228928
nalysis of Qu			
rray Size	Time Elapsed	countComp	countMove
		12497500	18769996
0000	426.053	49995000	75039996
5000	961.115	112492500	168809996
0000	1703.88	199990000	300079996
	2670.63	312487500	468849996
Sor	ting on Descendi	ng Arrays	
nalysis of In	sertion Sort		
rray Size	Time Elapsed	countComp	countMove
000	0.049233	4999	6240701
0000	0.096304	9999	19998
5000	0.146081	14999	29998
	0.193313	19999	39998
	0.239998	24999	49998
	0.291311	29999	59998
5000	0.335511	34999	69998
0000	0.388189	39999	79998

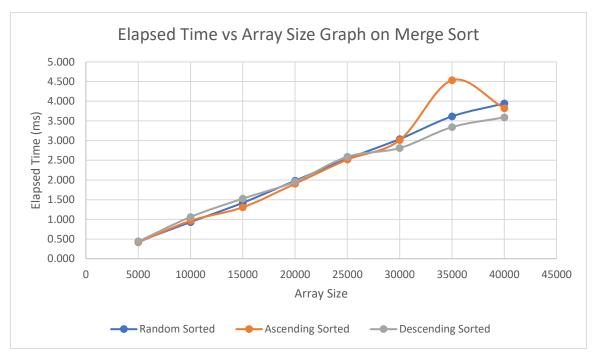


Question3

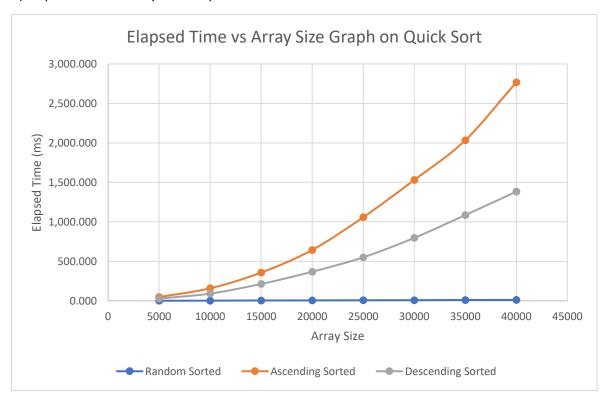
A) Elapsed Time vs Array Size Graph on Insertion Sort



B) Elapsed Time vs Array Size Graph on Merge Sort



C) Elapsed Time vs Array Size Graph on Quick Sort



Analysis:

Insertion Sort Analysis:

As it can be seen from the first graph, the greatest elapsed time belongs to ascending sorted array and if we compare our empirical results to the thereotical one it is quite reasonable. According to the theoretical result, ascending order is the worst-case since our insertion sort is implemented in a descending way and its time complexity is O(n^2). Whereas the best case is when the array is already descending sorted and its time complexity is O(n). Also as we expected the random sorted array gives an averages-case result and its time complexity again O(n^2). Therefore, I can argue that thereotical and empirical results match each other. The reason of why ascending order is the worst case is that is in reverse order, also the number of moves and the number of key comparisons are both in O(n^2) manner.

Merge Sort Analysis:

As it can be seen from the second graph, all elapsed times are quite close to each other and if we compare out empirical results to the theoretical one it is quite reasonable. In Merge Sort algorithm, both worst case and average cases are in O(nlogn) manner so there are very minor differences between ascending, descending, and random sorted arrays. Merge Sort is a quite efficient algorithm if it is compared with insertion, selection or bubble sort. However, by the principle of memory-time tradeoff, this algorithm requires an extra array whose size equals to the size of the original. Therefore, for many large sizes this algorithm could produce a problem and may give an error.

Quick Sort Analysis:

As it can be seen from the third graph, the greatest elapsed times belongs to ascending sorted array and if we compare our empirical results to the theoretical one it is quite reasonable. According to the theoretical result, quicksort is slow when the array is already sorted and if we choose element as the pivot. Similarly, we choose first element as a pivot and descending array is the worst case our implementation works in a descending manner. Therefore, for the sorted arrays the complexity is O(n^2) and both ascending and desceding sorted arrays work in this time complexity. On the other hand random sorted array takes quite minor time compared to sorted arrays and its time complexity is O(nlogn). Furthermore, when I try to compile after the size exceeding 25000, the compiler gives a stack overflow error for both ascending and descending sorted arrays. By my observation, I can argue that quick sort is a quite efficient algorithm if the array is not sorted, because if it is sorted compilers can't calculate excessively recursive and swap operations.

Note: I have added std=c++11 when I try to compile my file because chrono functions works only in this compiler version.