

Assignment 1

Glass Falling

A. Describe the optimal substructure/recurrence that would lead to a recursive solution.

Given n floors and m sheets, we will need to test all the floors as possible initial dropping floors; we'll call the initial floor i ,

If we have 1 floor left, we will have to test the floor.

If we have 0 floors left, we do not need to test anything.

If we have 1 glass sheet left, we need to test each floor from bottom up, or i numbers of times.

Two outcomes can occur from a drop: the glass shatters or it does not shatter.

If the glass shatters, we run the recurrence with one less sheet of glass and with the number of possible floors being $i-1$. We know the critical floor is below i since the glass shattered at i .

If the glass does not shatter, we run the recurrence with $n-i$ floors since we know the critical floor is between i and n .

B. Draw recurrence tree for given (floor=4, sheets=2)

```

DP(4,2) 1 DP(0,1)
          DP(3,2) 1 DP(0,1)
                    DP(2,2) 1 DP(0,1)
                              DP(1,2)
                                  2 DP(1,1)
                                      DP(0,2)
          2 DP(1,1)
            DP(1,2)
          3 DP(2,1)
            DP(0,2)
2 DP(1,1)
  DP(2,2) 1 DP(0,1)
            DP(1,2)
              2 DP(1,1)
                DP(0,2)
3 DP(2,1)
  DP(1,2)
4 DP(3,1)
  DP(0,2)

```

D. How many distinct sub-problems do you end up with given 4 floors and 2 sheets?

There are 8 distinct sub-problems:

(0,1) , (0,2) , (1,1) , (1,2) , (2,1) , (2,2) , (3,1) , (3,2)

E. How many distinct sub-problems do you end up with given n floors and m sheets?

$(n * m)$ distinct sub-problems

F. Describe how you would memorize GlassFallingRecur

A table of size $[\text{floors}+1][\text{sheets}]$ should be created.

`table[0][i] = 0;`

`table[1][i] = 1;`

`table[i][1] = i`

The code would follow the recursive method except before calling `GlassFallingRecur(floors, sheets)`, `table[floor][sheets]` should be checked. If a value exists, we just retrieve the value from the table. If the value does not exist, we will make a recursive call. The table will be completed thus fulfilling the memorization.

Rod Cutting

A. Draw a recursion tree for a rod length of 5.

```
5 0
  1 0
  2 0
    1 0
  3 0
    1 0
    2 0
      1 0
  4 0
    1 0
    2 0
      1 0
    3 0
      1 0
      2 0
        1 0
```

B. On page 370: answer 15.1-2 by coming up with a counterexample, meaning come up with a situation / some input that shows we can only try all the options via dynamic programming instead of using a greedy choice.

| Length | 1 | 2 | 3 | 4 | 5 |
|---------|---|---|----|----|----|
| Price | 1 | 4 | 30 | 36 | 60 |
| Density | 1 | 2 | 10 | 9 | 12 |

Whenever given a rod of length 4, the greedy algorithm would select to cut a portion of size 3 because it has the highest density of available options. This would yield two pieces of size 3 and 1 with corresponding values of 30 and 1; sum value of 31. The greedy algorithm fails to see that a rod of length 4 already has a greater value of 36 compared to the greedy algorithm solution.