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	${\bf MAT3007-Optimization}$									
Exercise Sheet Nr.:										
Name:	Che	2m /	fh g	Studen	at ID:	181	010	200		
In the creation of this solution sheet, I worked together with:										
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For correction:										
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Grading										

# MAT3007 Assignment 1

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A1.1

(a)

(b)

A1.2

(a)

(b)

(c)

(d)

A1.3

(a)

(b)

A1.4

(a)

(b)

(c)

(d)

## A1.1

(a)

Define n:=# floors above ground, m:=# floors underground, and u:= uniform height of floors. The optimization formulation is as follows:

$$egin{array}{ll} \min & lwd \ l_{w,h,d,n,m,u} \ & subject ext{to} & w \leq l \leq 2w \ & un = h \ & um = d \ & l \leq 40 \ & l \leq h \ & l \leq h \ & l/10 \leq m/(n+m) \leq 1/4 \ & lw(n+m) \geq 10000 \ & lw+2h(l+w) \leq 5000 \ & l,w,h,u \in \mathbb{R}^+ \ & n,m \in \mathbb{N}^+ \end{array}$$

(b)

A feasible point:

```
[l \ w \ h \ d \ n \ m \ u] = [30 \ 30 \ 30 \ 10 \ 30 \ 10 \ 1]
```

(a)

Let x, y be the # product I,II respectively. Then the LP formulation is:

$$\max_{x,y} \qquad \qquad 8x+7y$$
 subject to 
$$x/3+y/4 \leq 100$$
 
$$x/5+y/4 \leq 70$$
 
$$x,y \geq 0$$

(b)

$$egin{array}{ll} \min & -8x - 7y \ \mathrm{subject\ to} & 4x + 3y + s_1 = 1200 \ 4x + 5y + s_2 = 1400 \ x, y, s_1, s_2 \geq 0 \end{array}$$

(c)

Define z:=# overtime assembly hours. LP formulation:

$$\max_{x,y,z} \qquad \qquad 8x+7y-7z$$
 subject to 
$$x/3+y/4 \leq 100+z$$
 
$$x/5+y/4 \leq 70$$
 
$$z \leq 60$$
 
$$x,y,z \geq 0$$

(d)

yielding

$$\mathrm{max\_profit}(x^*,y^*) = 2500$$

with

$$\begin{bmatrix} x^* & y^* \end{bmatrix} = \begin{bmatrix} 255 & 100 \end{bmatrix}$$

(a)

Let  $A = [a_{ij}] \ge 0$  be the actual flow and  $C = [c_{ij}] \ge 0$  the capacity, both from  $v_i$  to  $v_j$ . Then the LP formulation of the problem is:

$$egin{array}{ll} \max & \sum_i a_{in} \ & ext{subject to} & 0 \leq A \leq C \ & \sum_j a_{ji} = \sum_j a_{ij} & orall i, j \in \{1,\dots,n\} \end{array}$$

\*Note that  $c_{ij} := 0$  if  $(i, j) \notin E$ .

(b)

```
% capacity values
C = [0 	 11 	 8]
                   0
    0
        0 10 12
                        0
    0
             0 0 11
             4
                   0 0 15
        0 0 7 0 4
    0
                             0];
cvx_begin
 variable A(6,6)
 maximize (ones(1,6) * A(:,6)) % 6th col sum
 subject to
   for i = 1 : 6*6
     % flow capacity restriction
     0 \ll A(i) \ll C(i)
   end
   for i = 2 : 6-1
     % flow conservation, ith row sum == ith col sum
     A(i,:) * ones(6,1) == ones(1,6) * A(:,i)
   end
cvx_end
```

yielding

$$\max_{\text{flow}}(A^*) = 19$$

with

$$A^* = egin{bmatrix} 0 & 11 & 8 & 0 & 0 & 0 \ 0 & 0 & 1.5947 & 10.2690 & 0 & 0 \ 0 & 0.8637 & 0 & 0 & 9.3653 & 0 \ 0 & 0 & 0.6343 & 0 & 0 & 15 \ 0 & 0 & 0 & 5.3653 & 0 & 4 \ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\min_{x \in \mathbb{R}^n} \ c^ op x \quad ext{s.t.} \quad -\delta \leq Ax - b \leq \delta$$

To show equivalence, it suffices to show equality of two feasible regions (since objective functions are the same).

Say we have  $x \in \mathbb{R}^n$  with  $-\delta \leq Ax - b \leq \delta, \ x \geq 0$ . Then

$$\begin{aligned} -\delta &\leq (Ax-b)_i \leq \delta, \ x \geq 0 \quad \forall i \in \{1,\cdots,m\} \\ \Longrightarrow & |(Ax-b)_i| \leq \delta, x \geq 0 \quad \forall i \in \{1,\cdots,m\} \\ \Longrightarrow & \max_{1 \leq i \leq m} |(Ax-b)_i| \leq \delta, x \geq 0 \\ \Longrightarrow & ||Ax-b||_{\infty} \leq \delta, x \geq 0 \end{aligned}$$

Conversely, suppose  $||Ax - b||_{\infty} \le \delta, x \ge 0$ . We have

$$egin{aligned} ||Ax-b||_{\infty} & \leq \delta, x \geq 0 \ \Longrightarrow & \max_{1 \leq i \leq m} |(Ax-b)_i| \leq \delta, x \geq 0 \ \Longrightarrow & |(Ax-b)_j| \leq \max_{1 \leq i \leq m} |(Ax-b)_i| \leq \delta, x \geq 0 & orall j \in \{1, \cdots, m\} \ \Longrightarrow & -\delta \leq Ax-b \leq \delta, \ x \geq 0 \end{aligned}$$

Therefore two feasible regions are indeed equal, hence the equivalence.

## (b)

Let  $x = \begin{bmatrix} x_1 & x_2 \end{bmatrix}^\top := \begin{bmatrix} \# \text{ salad A} & \# \text{ salad B} \end{bmatrix}^\top$ . The LP can be formulated in the standard form as follows:

(c)

Suppose  $x_0$  is a solution to the constraint. Then

$$25 = [1/4 \quad 1/2] x_0 = 2[1/8 \quad 1/4] x_0 = 2 \cdot 10 = 20,$$

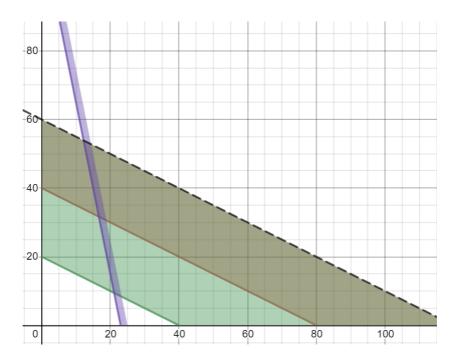
a contradiction. Hence the LP is infeasible.

### (d)

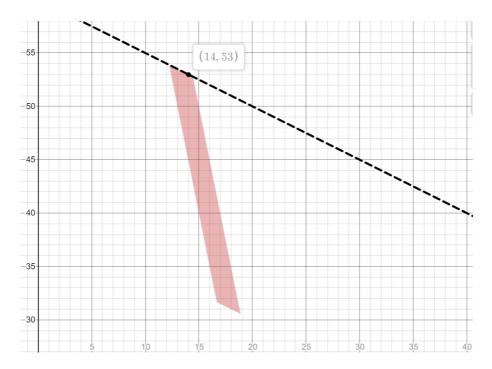
From (a) we have the equivalent formulation of the robust LP:

$$\begin{array}{ll} \max_x & 10x_1 + 20x_2 \\ \text{subject to} & x \geq 0 \\ & (\text{mango}) & -5 \leq x_1/4 + x_2/2 - 25 \leq 5 \\ & (\text{pineapple}) & -5 \leq x_1/8 + x_2/4 - 10 \leq 5 \\ & (\text{strawberry}) & -5 \leq x_1/8 + x_2/4 - 10 \leq 5 \end{array}$$

The following graph shows the three constraints separately given  $x \ge 0$ . (horizontal:  $x_1$ , vertical:  $x_2$ )



The graph below shows the feasible set in red (overlap of three regions above)



where the black dotted line is the contour of the profit function at optimum,  $10x_1 + 20x_2 = 1200$ . We see there's an infinite number of solutions (the entire upper edge of the feasible region) yielding the optimal profit 1200 RMB. The active constraints are:

$$x_1/4 + x_2/2 - 25 \le 5$$
  
 $x_1/8 + x_2/4 - 10 \le 5$ 

Nevertheless,  $(x_1^*, x_2^*) = (14, 53)$  is the only lattice point among the solutions and will be our final production plan. Following this plan, the amounts of fruits used are # mango  $= x_1^*/4 + x_2^*/2 = 30$ ;

# pineapple =  $x_1^*/8 + x_2^*/4 = 15$ ; # strawberry =  $5x_1^* + x_2^* = 123$ .