FancyMc Moves Fusion Moves for Multicut Objectives

Anonymous CVPR submission

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Multicuts rule.

1. Introduction

The tale of the multicut

1.1. Related Work

1.1.1 Multicut

- Andres et al. [1]
- Kappes et al. [5]
- Bagon and Galun [2]
- Yarkony et al. [6]
- Beier *et al.* [3]

1.1.2 Fusion Moves

Move making algorithms, in particular fusion moves, have become increasingly popular for energy minimization [?, 4]. For many large scale computer vision applications fusion moves lead to good approximations with state of the art any time performance [4].

2. Name of My Method (Union Fusion Cut)

Global optimal solvers for multicut do not scale beyond ??? [?]. Good approximate solvers for planar graphs exist [3, 6] but have difficulties to find good solutions for non planar graphs [3].

- 2.1. Proposal Generators
- 2.2. Fusion Move Solver
- 3. Experiments
- 4. Conclusion

References

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- [2] S. Bagon and M. Galun. Large scale correlation clustering optimization. CoRR, abs/1112.2903, 2011. 1
- [3] T. Beier, T. Kroeger, J. H. Kappes, U. Koethe, and F. Hamprecht. Cut, Glue & Cut: A Fast, Approximate Solver for Multicut Partitioning. In IEEE Conference on Computer Vision and Pattern Recognition 2014, 2014. 1
- [4] J. H. Kappes, T. Beier, and C. Schnörr. Map-inference on large scale higher-order discrete graphical models by fusion moves. In International Workshop on Graphical Models in Computer Vision, 2014. Oral. 1
- [5] J. H. Kappes, M. Speth, B. Andres, G. Reinelt, and C. Schnörr. Globally optimal image partitioning by multicuts. In EMM-CVPR, pages 31–44. Springer, 2011. 1
- [6] J. Yarkony, A. Ihler, and C. C. Fowlkes. Fast planar correlation clustering for image segmentation. In ECCV. Springer, 2012.

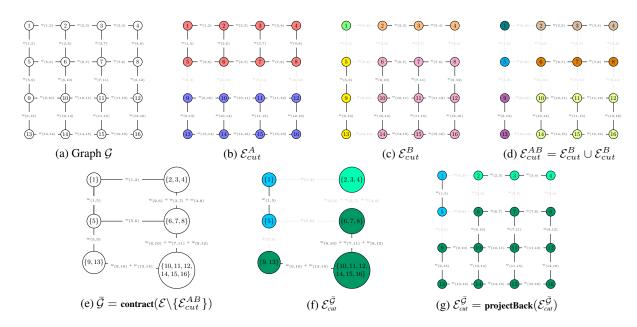


Figure 2: Describe Method here

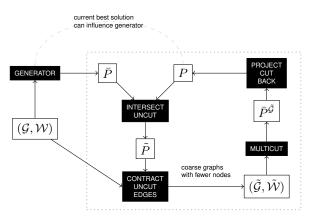


Figure 1: The proposed algorithm works in the following way: Given a graph \mathcal{G} , edge weights \mathcal{W} and a proposal generator, the current best solution P is iteratively improved. A proposal generator generates different versatile proposal partitions \bar{P} . The proposal \bar{P} is intersected with P which results in \tilde{P} . Contracting each each which is not cut in \tilde{P} leads to a coarser graph $\tilde{\mathcal{G}}=(\tilde{\mathcal{V}},\tilde{\mathcal{E}})$ with new edge weights $\tilde{\mathcal{W}}$. If \bar{P} and P have a small fraction of cut edges, $\tilde{\mathcal{G}}$ will be small ($|\tilde{\mathcal{V}}|<<|\mathcal{V}|$ and $|\tilde{\mathcal{E}}|<<|\mathcal{E}|$). The multicut objective on the smaller graph $\tilde{\mathcal{G}}$ can be optimized magnitudes faster than on $\tilde{\mathcal{G}}$. It is guaranteed that the optimal multicut partitioning $\bar{P}^{\tilde{\mathcal{G}}}$ on $\tilde{\mathcal{G}}$ projected back to \mathcal{G} as a lower or equal energy than any of the two input partitions P and \bar{P} , therefore we store the result of fusion as new best state P and repeat the procedure.

